

Lecture 8

RECURSION

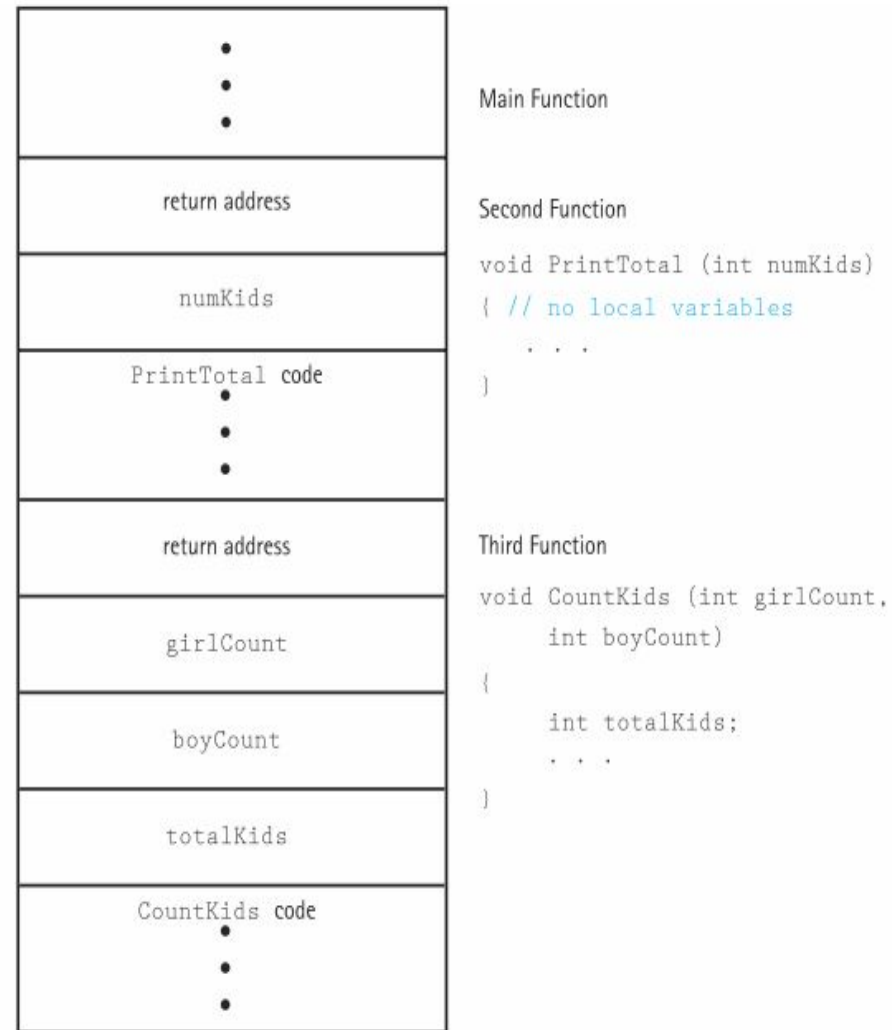
September 23, 2021
Thursday

FUNCTION CALLS

- If the function has formal parameters, they have to be initialized to the values passed as actual parameters.
- The system has to know where to resume the execution after the function is completed.
- Any other function can call this function, call from **main ()** ?
- Hence, a return address must be stored in main memory.
- But do we only require return address..?

STATIC ALLOCATION

- The function formal parameters and local variables are bound to actual addresses in memory at compile time.
- Think of having a 60 seats and sending invites to only 60 people. You can't entertain the 61st person.



DYNAMIC ALLOCATION

- A runtime stack is a much better solution to preserve the required information about a function.
- The distinction among local variables of different functions.
- What about the distinction between local variables of the same function which has been called more than once.

ACTIVATION RECORD

- The system characterizes each function
 - By the contents of all local variables
 - By the value of function's parameters.
 - By return address indicating

The data area containing all this information is called
Activation Stack Or **Stack Frame**

ACTIVATION RECORD

- Activation record is allocated on run-time stack.
- Exists as long as function owning it is executing.
- Consider it as the private repository which provides all the necessary information for proper execution and where to return after this function completes execution.
- Activation records usually have short lifespan because they are allocated on function entry and deallocated upon function exit.

Which Activation Record lives the most?

ACTIVATION RECORD | COMPONENTS

- Values for all parameters to the function.
 - Base address in case of array.
 - Address of variables passed by reference,
 - copies of all data items in case of pass by value.
- Local Variables which can be stored elsewhere in memory.
 - Their pointers will be saved in this case.
- Return Address to resume the control by the caller.
 - The address of the caller's instruction immediately following the call.

ACTIVATION RECORD | COMPONENTS

- Dynamic Link
 - A pointer to the caller's activation record.
- Returned Value
 - If function is not declared as void.
 - The size of the activation record may vary from one call to the other.
 - The returned value is placed right above the activation record of the caller.

main () calls f1 ()
f1 () calls f2 ()
f2 calls f3 ()

*Activation
record
of f3 ()*

Parameters and
local variables

Dynamic link

Return address

Return value

*Activation
record
of f2 ()*

Parameters and
local variables

Dynamic link

Return address

Return value

*Activation
record
of f1 ()*

Parameters and
local variables

Dynamic link

Return address

Return value

*Activation
record
of main ()*

RECURSIVE CALL | AN ANATOMY

$$x^n = \begin{cases} 1 & \text{if } n = 0 \\ x \cdot x^{n-1} & \text{if } n > 0 \end{cases}$$

```
/* 102 */ double power (double x, unsigned int n) {  
/* 103 */     if (n == 0)  
/* 104 */         return 1.0;  
        // else  
/* 105 */     return x * power(x,n-1);  
}
```

RECURSIVE CALL | AN ANATOMY

call 1	power(x, 4)				
call 2		power(x, 3)			
call 3			power(x, 2)		
call 4				power(x, 1)	
call 5					power(x, 0)
call 5					1
call 4				x	
call 3			x · x		
call 2		x · x · x			
call 1	x · x · x · x				

RECURSIVE CALL | AN ANATOMY

```
int main ( ) {  
    ...  
/* 136 */    y = power(5.6, 2);  
    ...  
}
```

call 1	power(5.6, 2)
call 2	power(5.6, 1)
call 3	power(5.6, 0)
call 3	1
call 2	5.6
call 1	31.36

<i>Third call to power ()</i>			$\begin{Bmatrix} 0 \leftarrow \text{SP} \\ 5.6 \\ (105) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 0 \leftarrow \text{SP} \\ 5.6 \\ (105) \\ 1.0 \end{Bmatrix}$	$\begin{Bmatrix} 0 \\ 5.6 \\ (105) \\ 1.0 \end{Bmatrix}$			
<i>Second call to power ()</i>		$\begin{Bmatrix} 1 \leftarrow \text{SP} \\ 5.6 \\ (105) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 1 \\ 5.6 \\ (105) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 1 \\ 5.6 \\ (105) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 1 \leftarrow \text{SP} \\ 5.6 \\ (105) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 1 \leftarrow \text{SP} \\ 5.6 \\ (105) \\ 5.6 \end{Bmatrix}$	$\begin{Bmatrix} 1 \\ 5.6 \\ (105) \\ 5.6 \end{Bmatrix}$	
<i>First call to power ()</i>	$\begin{Bmatrix} 2 \leftarrow \text{SP} \\ 5.6 \\ (136) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 2 \\ 5.6 \\ (136) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 2 \\ 5.6 \\ (136) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 2 \\ 5.6 \\ (136) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 2 \\ 5.6 \\ (136) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 2 \\ 5.6 \\ (136) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 2 \leftarrow \text{SP} \\ 5.6 \\ (136) \\ ? \end{Bmatrix}$	$\begin{Bmatrix} 2 \leftarrow \text{SP} \\ 5.6 \\ (136) \\ 31.36 \end{Bmatrix}$
<i>AR for main ()</i>	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$	$\begin{Bmatrix} : \\ y \\ : \end{Bmatrix}$
	(a)	(b)	(c)	(d)	(e)	(f)	(g)	(h)

Key: SP Stack pointer
AR Activation record
? Location reserved
for returned value

TAIL RECURSION

- Use of only one Recursive Call.
- Recursive call is the last statement of the function.
- There are no earlier recursive calls.
 - Neither Directo nor indirect

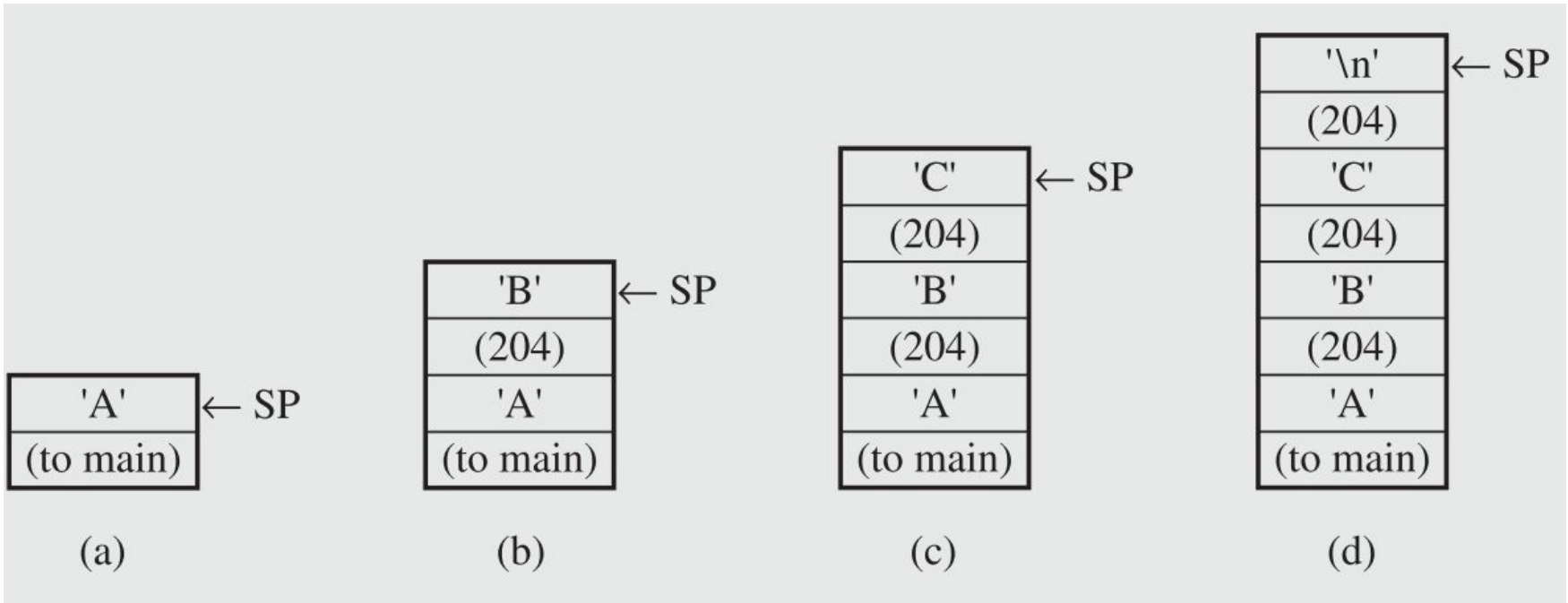
Tail recursion is used in languages like Proglo, which does not have a loop construct.

```
void tail ( int i ) {  
    if ( i > 0 ){  
        cout<<i<<' '  
        tail ( i - 1 );  
    }  
}
```

NON-TAIL RECURSION | INVERT A LINE

```
void reverse() {                /* 200    */
    char ch;                    /*      */
    cin.getch(ch);              /* 201    */
    if ( ch != '\n'){           /* 202    */
        reverse();              /* 203    */
        cout.put(ch);           /* 204    */
    }
}
```

NO-TAIL RECURSION



NESTED RECURSION

A function is not only defined in terms of itself, but also is used as one of the parameters.

$$h(n) = \begin{cases} 0 & \text{if } n = 0 \\ n & \text{if } n > 4 \\ h(2 + h(n)) & \text{if } n \leq 4 \end{cases}$$

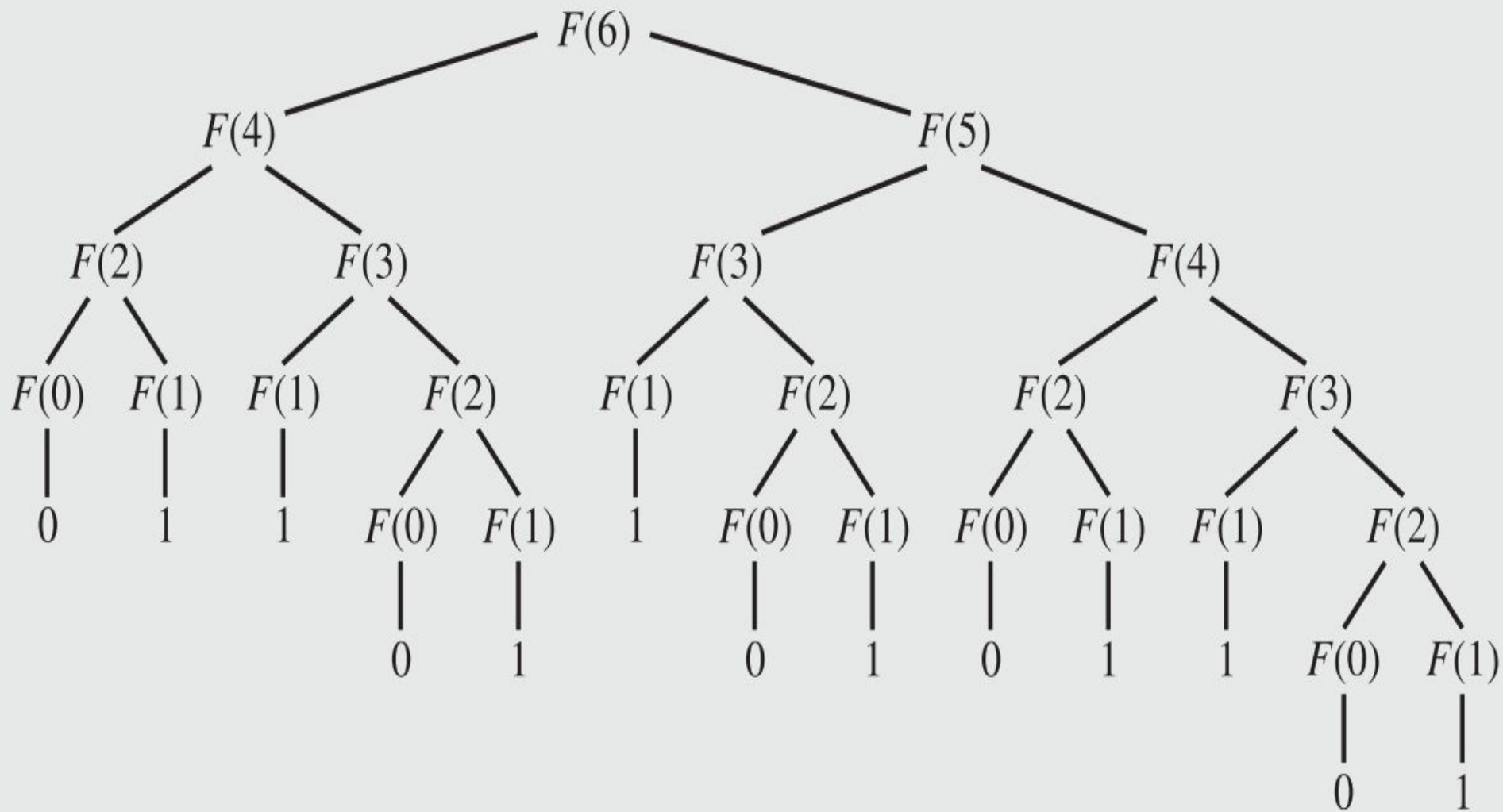
EXCESSIVE RECURSION

- If the recursion is too deep ($5.6^{100,000}$), we will run out of space on stack.
 - Program crashes.
 - In case of power range of variable might not allow such a big number as well.
- But what if some recursive functions repeat themselves for a very long time even for very simple cases.
 - Consider the case of Fibonacci numbers.

$$\text{Fib}(n) = \begin{cases} n & \text{if } n < 2 \\ \text{Fib}(n-2) + \text{Fib}(n-1) & \text{otherwise} \end{cases}$$

EXCESSIVE RECURSION

```
unsigned long Fib ( unsigned long n ) {  
    if (n < 2)  
        return n;  
    else  
        return Fib (n-2) + Fib (n-1);  
}
```



NUMBER OF ADDITIONS Vs NUMBER OF CALLS

n	Fib (n+1)	Number of Additions	Number of Calls
6	13	12	25
10	89	88	177
15	987	986	1,973
20	10,946	10,945	21,891
25	121,393	121,392	242,785
30	1,346,269	1,346,268	2,692,537

NUMBER OF ADDITIONS Vs NUMBER OF CALLS

```
unsigned long FibIterative (unsigned long n) {  
    if (n < 2 )  
        return n;  
    else {  
        long i = 2, temp, current = 1, last =0;  
        for (; i <= n; ++i){  
            temp = current;  
            current += last;  
            last = temp;  
        }  
        return current;  
    }  
}
```

ACTIVITY

Check Recursively If given

Word is Palindrome?

Sentence is Palindrome ignoring the space?

Given a C++ Program check if all the opening brackets

({ [< have their corresponding closing >] })

Find an element in 2D DynamicArray