# Taraxa PBFT

#### Version 1.0

### February 01, 2025

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When user u starts round (p, r), they reset their timer u to 0 and other constants as follows:

- If r = 1:  $-\lambda_r \in (500, 1500)$   $-\Lambda_r = 4000$ • Otherwise if r >= 2:
  - $-\lambda_r = 2000$  $-\Lambda_r = 17000$

## 1.2 Round (p, r) Voting Instructions

Round (p, r) Constants

The voting instructions are as follows:

Step 1: **Proposal** - When  $timer_u = 0$ :

- If r=1 or r>1 and u has received a next-quorum for  $\bot$  from round (p,r-1), then u assembles a new block proposal  $B_u$  and propagates  $B_u$  and  $H(B_u)$ .
- Otherwise, if r > 1 and u has received a next-quorum for  $H(B') \neq \bot$  from (p, r 1), then u propagates H(B').

Step 2: **Filtering** - When  $timer_u = 2\lambda_r$ :

- If r=1 or if r>1 and u has received a next-quorum for  $\bot$ , then u selects the proposal with the minimum credential and soft-votes for it
- Otherwise, if r > 1 and u has received a next-quorum for  $H(B') \neq \bot$ , then u soft-votes for H(B').

Step 3: Certifying - While  $timer_u \in (2\lambda_r, \max(4\lambda_r, \Lambda_r))$ :

- If u receives a soft-quorum for H(B) and a valid block B with H(B) = H(B), then u cert-votes for H(B).

Step s=2n, where  $n\in(2,\infty)$ : First Finishing Step - When  $timer_u=\max(4\lambda_r,\Lambda_r)+(s-4)\lambda_r$ :

- If i has certified some value v for round r, he next-votes v.
- Else if  $(r \ge 2$  and i has seen 2t+1 next-votes for  $\bot$  for round r-1), he next-votes  $\bot$ .
- Else he next-votes his starting value  $st_i^r$ .

Step s=2n+1, where  $n\in(2,\infty)$ : **Second Finishing Step** - When  $\limsup_{n\to\infty} 4\lambda_r, \Lambda_r) + (s-5)\lambda_r + 100$ ms:

- If i sees 2t+1 soft-votes for some value  $v \neq \bot$  for round r, then i next-votes v.
- If  $(r \ge 2 \text{ and } i \text{ sees } 2t+1 \text{ next-votes for } \bot \text{ for round } r-1 \text{ and } i \text{ has not certified in round } r)$ , then  $i \text{ next-votes } \bot$ .