

Tethys
A Toy Functional Programming Language
with a System $F\omega$ -based Core Calculus
<https://github.com/ThePuzzlemaker/tethys>

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January 15, 2022

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1 Introduction

This “paper” (which is really just a well-typeset, but somewhat informal write-up) introduces Tethys, a toy functional programming language based on a System $F\omega$ -based core calculus. Hence the title.

There are two parts of Tethys: the surface language, and the core calculus. The core calculus is the intermediate representation of Tethys which is used for type checking and inference, and for interpretation. The surface language is the higher-level interface that is eventually desugared by the compiler/interpreter to the core calculus.

The reference implementation in Rust will not use any particular “tricks” in terms of interpretation, instead just using a simple tree-walk interpreter or similar.

This language was created in order to conduct informal research (i.e., not actually discovering anything interesting, probably) on type systems, especially bidirectional typechecking and polymorphism. Tethys is named as such as it is the name of the co-orbital moon to Calypso; as my work on this language is “co-orbital”, so to speak, to my work on [Calypso](#).

Please note that I am not an expert in pretty much any subject this writeup covers. If you notice something you don’t understand, or that you think is a mistake, don’t fret to point it out.

1.1 Background

Note: This section is taken from [some slides I prepared for a lightning talk](#) in the [r/PL Discord server](#). If it seems a bit un-prose-y, that’s why.

Originally, I thought I needed to use System $F\omega$, an extension of the polymorphic typed lambda calculus (System F), with the extension of type lambdas. However, it has some problems:

- System F is already undecidable enough, with respect to inference. Adding *more* types of types does **not** help.
- Theoretically, bidirectional typechecking (my chosen strategy) doesn’t *need* unification, but it sure makes it nice. And having type lambdas means you need higher order unification, which is scary because anything including “higher-order” in its name is automatically scary. Also it’s undecidable, so you have to restrict yourself to higher-order pattern unification, which is still quite scary.
- I just can’t really find much good introductory literature on System $F\omega$ (at least, literature that I can understand).

As it turns out, System $F\omega$ is more expressive than F, but **way** much more of a pain. Not even Haskell, the paragon of kitchen sink type system features has it. (By default, because of course GHC’s gonna have it as an extension).

Here are some of the key theoretical differences:

- Within the type system, type constructors are opaque. Defining a type just creates opaque term-level and type-level constants, which don’t normalize to the type or term

they “really” mean. Of course, these are stripped out after typechecking, once they’re no longer needed.

- Type constructors are injective, i.e. for some type constructor C and types X and Y , if $C\ X$ and $C\ Y$ are the same type, then X and Y are the same type.
- Differently named type constructors are never equal, even if they “mean” the same thing, i.e. for some type constructors $C1$ and $C2$, and type T , $C1\ T$ never equals $C2\ T$, no matter how they are defined.

Note that this is pretty much nominal typing! (If my understanding of nominal typing is correct.)

2 The Surface Language

This section has not been started yet.

3 The Declarative Core Calculus

This section is, unsurprisingly, work-in-progress.

4 References

4.1 Acknowledgements

Thanks to Brendan Zabarauskas and András Kovács, who have helped with many type system details. Many thanks to Mark Barbone (aka MBones/mb64), who provided a code sample of their bidirectionally typed (with unification) algorithm for typechecking higher-rank System F (Barbone, 2021), and who also noted some issues they had with “Complete and Easy Bidirectional Typechecking for Higher-Rank Polymorphism” (Dunfield and Krishnaswami, 2020).

4.2 References

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