

Bitcoin Mechanics

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Properties of Bitcoin

- Uses Elliptic Curve Public Keys (secp256k1) and ECDSA signature algorithm
- UTXO Blockchain (Unspent Transaction Output)
- Block size limit: 1 MB (~2000 transactions/block)
- Block time : ~10 minutes
- Consensus: Proof of Work (coming later)

Propagation Time

According to the paper "*Information propagation in the bitcoin network*" by Decker and Wattenhofer (2013):

The **median time** until a node receives a block is **6.5 seconds** whereas **the mean** is at **12.6 seconds**.

The long tail of the distribution means that even **after 40 seconds there still are 5% of nodes that have not yet received the block**

- It is hard to maintain data consistency and avoid double spending attack (rf lecture 1)

The Bitcoin solution : Mining

Confirming transaction into blocks

- Miners validate every transaction broadcasted on the network and add them to a mempool of unconfirmed transactions
- Approximately every 10 minutes, *one node is selected* (see consensus later) to create a block containing all unconfirmed transactions and broadcast that block to the network to be added to the blockchain
- All blocks validate the new node before adding it to their own copy of the blockchain

How is this solving the problems

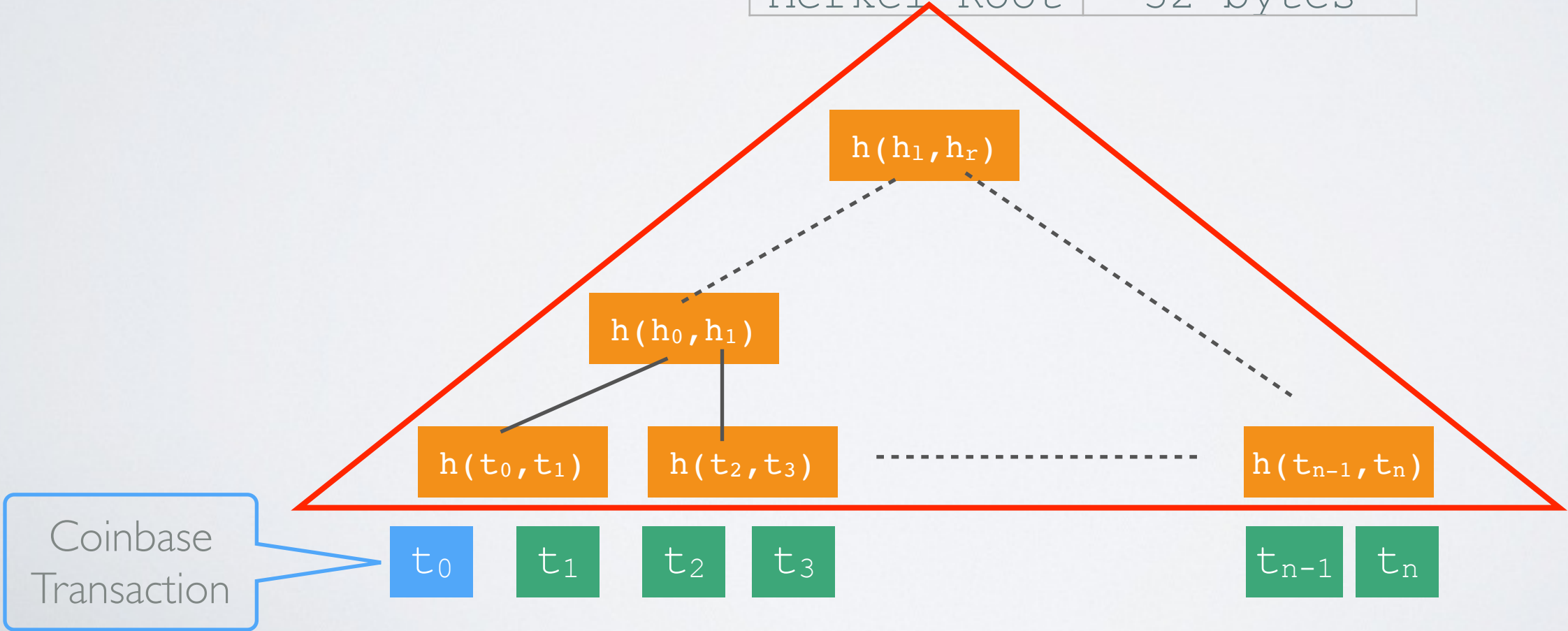
- ➔ New transactions cannot use input UTXO of transactions that have not been confirmed yet
- ✓ Data is always consistent and not double spending attack unless two different blocks are mined at the same time (see consensus problem)

Mining awards

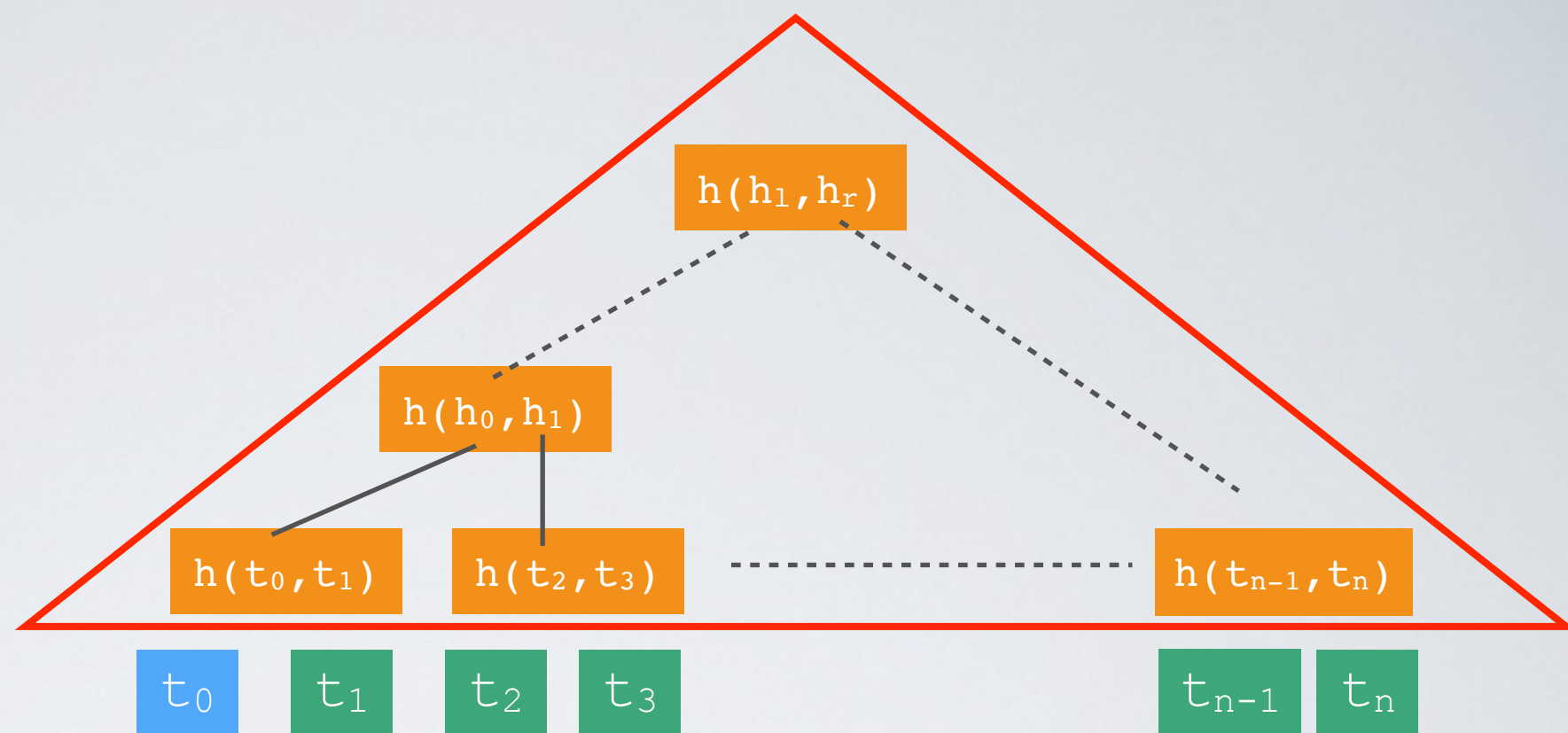
- ➡ Miners verify/broadcast blocks transactions and broadcast and are rewarded for that work
 - **Coinbase transaction** (first transaction in the block)
Currently 6.25 BTC - Block reward halves every four years
The only way BTC is created (max 21M BTC in total)
 - and/or **Transaction Fees** (chosen by the issuer)

Anatomy of the Bitcoin blockchain

Block Hash	32 bytes	→
Block Id	4 bytes	
→ Old Hash	32 bytes	
Time	4 bytes	
Size	4 bytes	
Nonce	4 bytes	
Merkel Root	32 bytes	



Properties of Merkle Trees



Why using a Merkle-Tree to record all transactions in the block rather than a "flat list" such as $H(T_0 + T_1 + \dots + T_n)$?

- ✓ Because it is easier to prove than a given transaction belongs to a block
- Using flat list, need all n transactions to build the proof
- Using "Merkle Tree, just $\log_2(n)$ intermediate hashes to build the proof

Let's look at some blocks