

- Partiti gets created based on the requirement of the incoming process

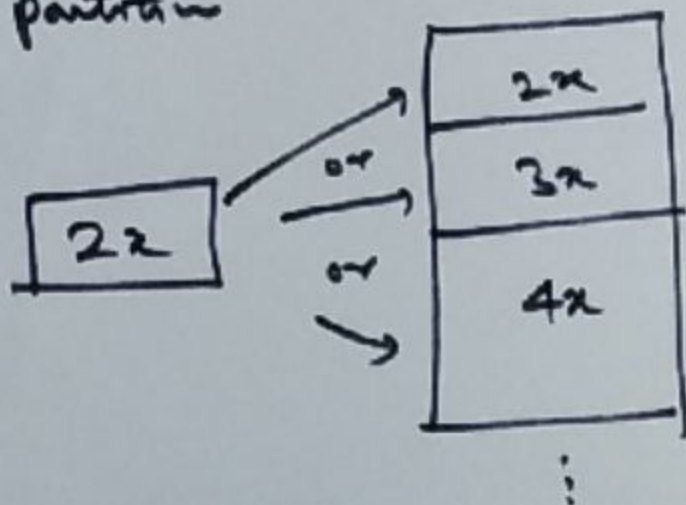
## Issues on Partitioning

— Placement Algos —

— Memory Wasted, —  
(Fragmentation)

Problem

— Program can go into multiple partitions



# Placement Algorithm

- ① First-fit — from beginning
- ② Best-fit — Minimum wastage  
↳ Best possible fit of memory
- ③ Worst-fit — Maximum wastage of memory
- ④ Next-fit → starts from point of last allocation

## First fit

X  
incoming  
program

NE ↓  
Search

A
B
C
D
E

start  
↓, searching  
First  
fit (X)

Y  
incoming  
program

Z  
incoming  
program

A - Z

B - Z

C - Z

D - Z

G - Z

Allocated

Z  
BE wait

Best  
available



Best possible  
Best Available  
(sub-optimal)

different  
than Best fit  
(optimal)

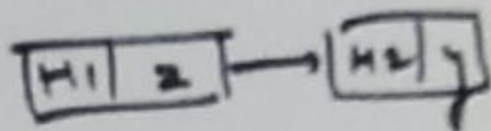
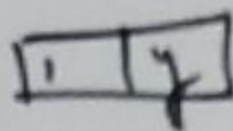
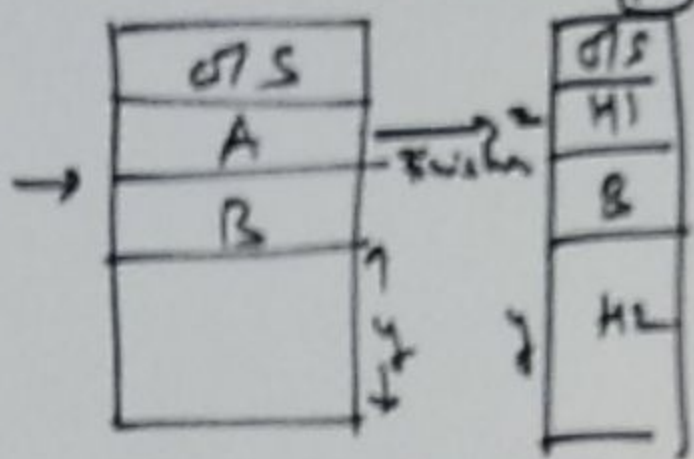
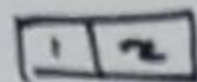
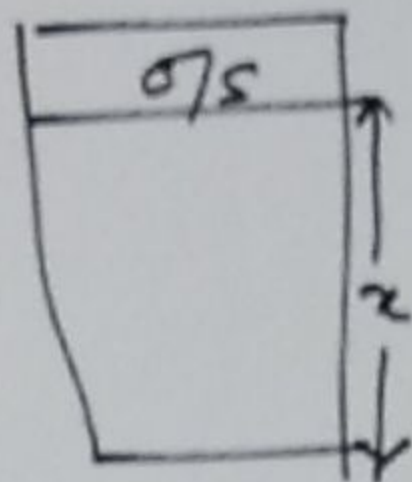
Worst fit  $\rightarrow$  largest space

## Variable Partitioning

↳ Hole → Every memory allocation creates a partition & a hole

$$H \xrightarrow[A]{\text{Allocate}} \left. \begin{matrix} A \\ H-A \end{matrix} \right\}$$

- Hole list — scattered across memory



Placement  
Algo

Fixed Partitioning  
(Partition)

Variable Partitioning  
(Hde)

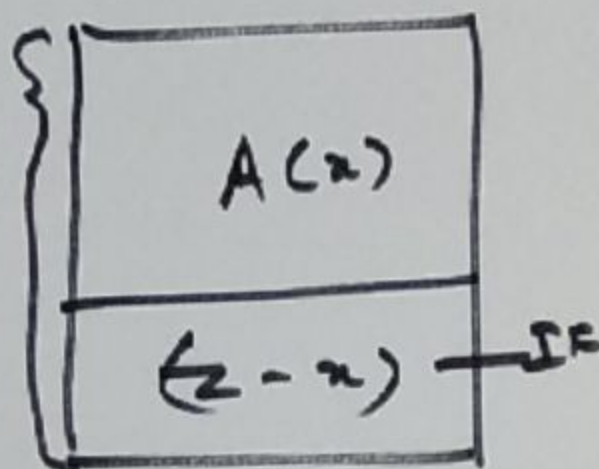


# Memory Fragmentation

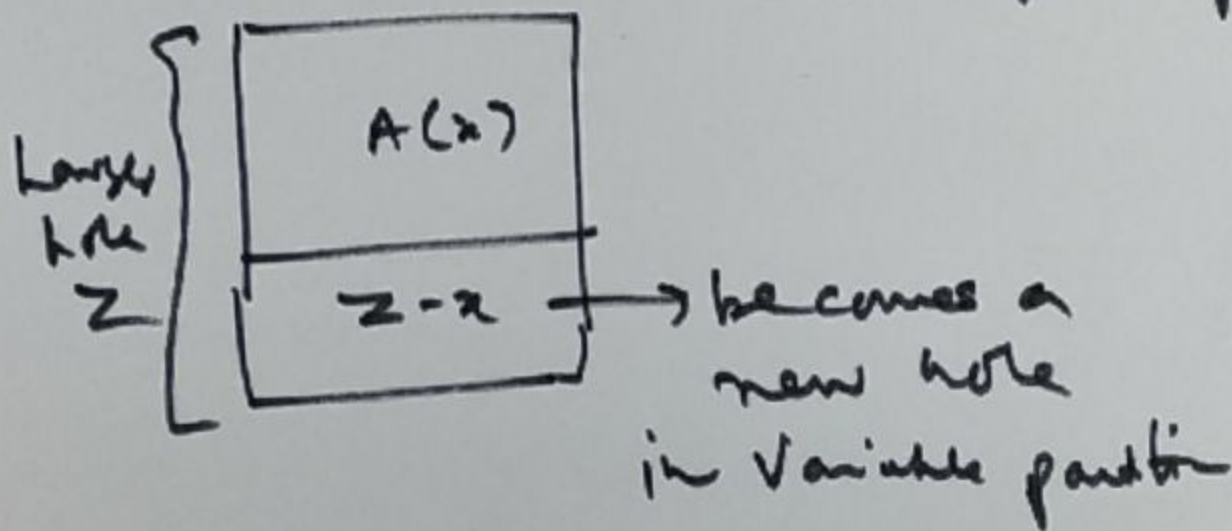
⑪

- ① Internal Fragmentation (IF)  
(inside a partition/page...  
memory block)

A  
(size =  $x$ )  
( $Z > x$ )



Partition size =  $Z$   
(Fixed partition)



(12)

— All memory partitions are created to the size of powers of 2 (closest)

A  
(20K) 3 bytes → 4 bytes  
5 bytes → 8 bytes

Theory → No IF  
Implementation → Yes IF  
(less → high)

② External Fragmentation  
(outside the partition)



D  
increasing  $\xrightarrow{\text{load}}$

575
H1
A
H2
B
H3
C

$$D > H1$$

$$D > H2$$

$$D > H3$$

H1, H2, H3  
are scattered

contiguous  
(compactness)

$$D < H1 + H2$$

$$D < H2 + H3$$

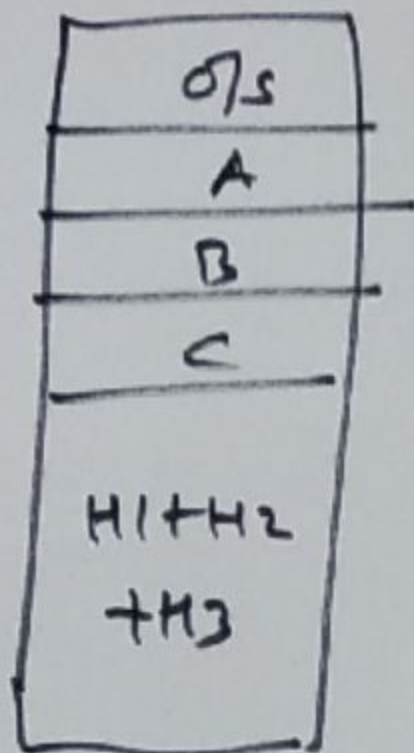
$$D < H1 + H3$$

$$D < \boxed{H1 + H2 + H3}$$

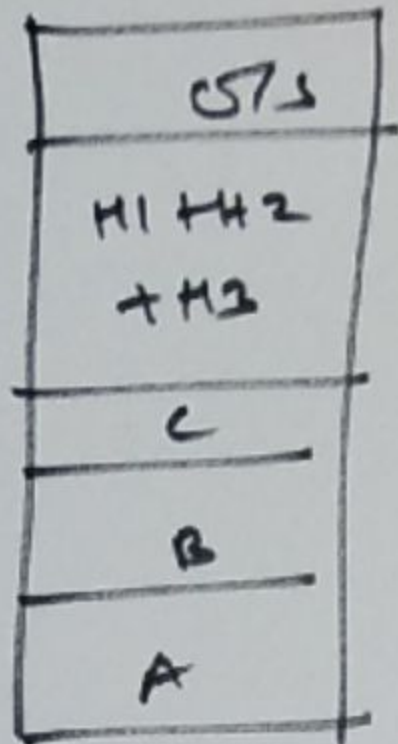
EF

No  
EF

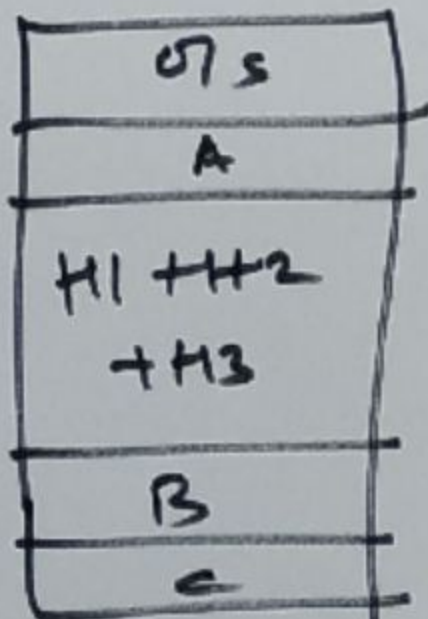
$$\boxed{D > H1 + H2 + H3}$$



or



or



Compact<sup>n</sup>  
 ↓  
 solution  
 to  
 FP  
 (in many  
 cases)

—Memory Manager attempt to  
 optimize the relocation effort

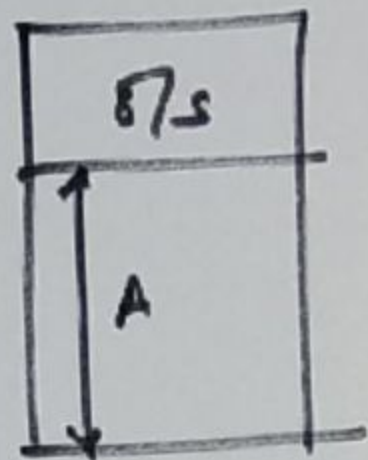


### ③ Table Fragmentation (TF) ⑮

↳ Data structures related  
to memory management  
Used

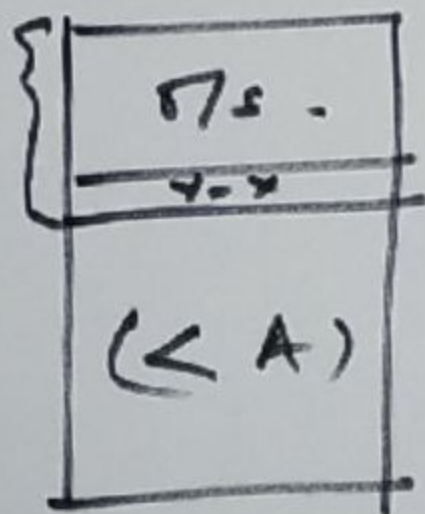
Memory  
Management  
Scheme 1  
(TF = X)

—————  
Data  
structure  
X



Memory  
Management  
Scheme 2  
(TF = Y)

—————  
Data  
structure  
Y



$(Y > X)$

$(Y - X)$  extra.

Total frag = IF + EF + TF