Lecture -28 PO Recap'. $E\left[g\left(X,Y\right)\right] = \iint g\left(x,y\right) f\left(x,y\right) ds ds$ Elxty = Elx] + Ely Doesn't assume independence Nois yours of its Nois yours of its (or por collecting problem in a plane Dounbard's walk steplength=1 After n steps, he is Bar D with distance away. E[D²].

x, = (05 0, y,= sina, x = (05 @2 y = sin 02 xn = (050n Ja =sinon ストメント・ナチャンナイフナフェナ・ナタルン

$$D^{2} = (1050, 1050) + (1050)^{2} + (3)$$

$$(5ino) + 5ino) + (1050)^{2}$$

$$= n + \sum_{i \neq j} (050; (050) + (114))$$

$$\sum_{i \neq j} (050; 5ino)$$

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Time (om plexity Analysis (9)
of Quick Sort Algorithm (35, 9, 3, 8, 4, Q) (0) (1), (9), 7) (5, 9, 3, 8, 4, Q) (0) (1), (9), 7) (5,39) 6 (9,8) 10 (11) 14 (17) 3,4,5,6,8,9,10,11,14,17 X- no. of comparisons. = 52 Fix I if is gare compared index of otherwise. T4,9=0 I57=1

 $X = \sum_{i=1}^{n} \sum_{j=i+1}^{n} T_{ij}$ ECX = E [SE Tix] = 3 是正正说 Probability that the its smallest 6 The ish smallest numbers lompared with each other.

In the starting are together). when will is i get compared? $\frac{2}{3-i+1} = EEIIii$ (5, 7.3, 120)17 & 4, 17, 8 I4,8 lesson 4 { 3', 2', 3} (14,17) 2

$$E[X] = \sum_{j=1}^{n-1} \sum_{j=i+1}^{n} \frac{2}{j-i+1}$$

$$= 2 \log (x-i+1)$$

$$= 2 \log (x-i+1) - \log 2$$

$$= 2 \log (x-i+1) - \log 2$$