Analysis of recursive algorithms results in recursive equations. If T(n) is the time taken to find the max element in an array of size n, then.

- ▶ $T(n) = c_1$, if n = 1
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$$T(n) = c_2 \cdot (n-1) + c_1$$

- ▶ At times recursive equations can be tricky.
- ▶ Master Theorem is used to solve recurrence equations in asymptotic terms.

Recurrence relations can also be used to define a sequence

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$$x_{n+1} = c.x_n$$
, given $(n \ge 0; x_0 = 1)$

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$$x_n = c^n$$

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Given:

- ▶ $n \ge 0$
- ▶ the value of x₀
- ▶ the set $\{b_1, b_2, \cdots\}$

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$$x_n = b_n.b_{n-1}.b_{n-2}\cdots b_1.x_0$$

Let us raise the ante further: $x_{n+1} = b_{n+1}.x_n + c_{n+1}$

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Hint: Reduce it to the form $y_{n+1} = y_n + d_{n+1}$

Example: $x_{n+1} = 3.x_n + n$, where $(n \ge 0; x_0 = 0)$

If $x_n = 3^n y_n$, then we get

$$y_{n+1} = y_n + n/3^{n+1}$$
, where $n \ge 0$; $y_0 = 0$

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Finally we get:
$$x_n = 3^n \sum_{i=1}^{n-1} j/3^{j+1}$$

Till now we have only considered first-order recursive equations.

Note:

- ▶ In first-order recursive equations, the current value depends only on the previous value.
- ▶ In second-order recursive equations, the current value depends on the previous **two** values.

Let the equation be: $x_{n+1} = a.x_n + b.x_{n-1}$

Given:

- ▶ $n \ge 1$
- ▶ the value of x_0, x_1 and a, b

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Hint: Call for a trial solution (as you solve second-order differential equations)

Let the equation be:
$$x_{n+1} = a.x_n + b.x_{n-1}$$

Follow the following steps:

- Let the trial solution be $x_n = \alpha^n$, and substitute it in the given equation
- We obtain the quadratic equation $\alpha^2 = a \cdot \alpha + b$
- ▶ If α_+ and α_- are the distinct roots, the general solution is $x_n = c_1.\alpha_+^n + c_2.\alpha_-^n$
- ► The constants c_1 and c_2 will be determined so that x_0, x_1 have the assigned values.

Example

$$L_0 = 100000, L_1 = 200000, \text{ and } L_n = (L_{n-1} + L_{n-2})/2$$

- ▶ The characteristic polynomial is · · ·
- ▶ The roots are · · ·
- ▶ The general solution is · · ·
- ▶ The values of c_1 and c_2 are · · ·

Example

$$L_0 = 100000, L_1 = 200000, \text{ and } L_n = (L_{n-1} + L_{n-2})/2$$

- ▶ The characteristic polynomial is $x^2 x/2 1/2$
- ▶ The roots are 1 and -1/2
- ▶ The general solution is $L_n = c_1 + c_2(-1/2)^n$
- $c_1 = 500000/3$ and $c_2 = -200000/3$.