

Indirect Search Methods

Gradient of a function

The gradient of a function in an n -dimensional component vector is given by

$$\nabla f_{nx_1} = \left\{ \begin{array}{l} \frac{\partial f}{\partial x_1}, \\ \frac{\partial f}{\partial x_2}, \\ \vdots \\ \frac{\partial f}{\partial x_n} \end{array} \right\} \quad \dots (6.56)$$

✓ Steepest Descent (Cauchy) method \rightarrow the use of the negative of the gradient vector as a direction for minimization was first made by Cauchy in 1847.

In this method, we start from an initial trial point x_i and iteratively move along the steepest descent directions until the optimum point is found. This method is summarized as

1. Start with an arbitrary initial point x_i .
2. Find the search directions s_i as

$$s_i = -\nabla f_i = -\nabla f(x_i)$$

3. Determine the optimal step length d_i^* in the direction ~~s_i~~ and get

$$x_{i+1} = x_i + d_i^* s_i = x_i - d_i^* \nabla f_i \quad (6.70)$$

4. Test the new point, x_{i+1} for optimality - If x_{i+1} is optimum, stop the process, otherwise, go to Step 5.

⑤. Set new iteration number $i = i + 1$ and go to Step 2.

Q. Minimize $f(n_1, n_2) = n_1 - n_2 + 2n_1^2 + 2n_1n_2 + n_2^2$
starting from the point $x_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$

Iteration 1.

The gradient of f is given by

$$\nabla f = \begin{Bmatrix} \frac{\partial f}{\partial n_1} \\ \frac{\partial f}{\partial n_2} \end{Bmatrix} = \begin{Bmatrix} 1 + 4n_1 + 2n_2 \\ -1 + 2n_1 + 2n_2 \end{Bmatrix}$$

$$\Rightarrow \nabla f_1 = \begin{Bmatrix} 1 + 4 \cdot 0 + 2 \cdot 0 \\ -1 + 2 \cdot 0 + 2 \cdot 0 \end{Bmatrix} = \begin{Bmatrix} 1 \\ -1 \end{Bmatrix}$$

Therefore

$$s_1 = -\nabla f = -\begin{Bmatrix} 1 \\ -1 \end{Bmatrix} = \begin{Bmatrix} -1 \\ 1 \end{Bmatrix}$$

To find λ_1^* , we need to find the optimum step length λ_1^* . For this, we minimize

$$\text{we minimize } f(x_1 + \lambda_1 s_1) = f\left(\begin{Bmatrix} -\lambda_1 \\ \lambda_1 \end{Bmatrix}\right)$$

$$= -\lambda_1 - \lambda_1 + 2x_1^2 - 2x_1\lambda_1 + \lambda_1^2$$

$$= \lambda_1^2 - 2\lambda_1$$

$$\text{Now } \frac{\partial f}{\partial \lambda_1} = 2\lambda_1 - 2 = 0 \Rightarrow \lambda_1^* = 1$$

we obtain

$$x_2 = x_1 + d_1^* s_1 = \begin{Bmatrix} 0 \\ 0 \end{Bmatrix} + 1 \begin{Bmatrix} -1 \\ 1 \end{Bmatrix} = \begin{Bmatrix} -1 \\ 1 \end{Bmatrix}$$

$$\text{as } \nabla f_2 = \nabla f(x_2) = \begin{Bmatrix} 1-4+2 \\ -1+2+2 \end{Bmatrix} = \begin{Bmatrix} -1 \\ 1 \end{Bmatrix} \neq \begin{Bmatrix} 0 \\ 0 \end{Bmatrix}, x_2 \text{ is not optimum.}$$

Iteration 2.

$$s_2 = -\nabla f_2 = \begin{Bmatrix} 1 \\ 1 \end{Bmatrix}$$

$$\begin{aligned} f(x_1 + d_2 s_2) &= f\left(\begin{Bmatrix} -1 \\ 1 \end{Bmatrix} + d_2 \begin{Bmatrix} 1 \\ 1 \end{Bmatrix}\right) = f(-1 + d_2, 1 + d_2) \\ &= 5d_2^2 - 2d_2 - 1 \end{aligned}$$

$$\frac{df}{dd_2} = 10d_2 - 2 = 0 \Rightarrow d_2^* = 0.2$$

$$\begin{aligned} x_3 &= x_2 + d_2^* s_2 = \begin{Bmatrix} -1 \\ 1 \end{Bmatrix} + \frac{1}{5} \begin{Bmatrix} 1 \\ 1 \end{Bmatrix} = \begin{Bmatrix} -1+0.2 \\ 1+0.2 \end{Bmatrix} = \begin{Bmatrix} -0.80 \\ 1.20 \end{Bmatrix} \quad \text{circled} \\ &= \begin{Bmatrix} -0.80 \\ 1.20 \end{Bmatrix} \quad \text{circled} \end{aligned}$$

$$= \begin{Bmatrix} -1+0.2 \\ 1+0.2 \end{Bmatrix} = \begin{Bmatrix} -0.80 \\ 1.20 \end{Bmatrix} \Rightarrow x_3 = \begin{Bmatrix} -0.80 \\ 1.20 \end{Bmatrix}$$

$$\text{now, } \nabla f_3 = \nabla f(x_3) = \begin{Bmatrix} 1-3.20+2.4 \\ -1-1.6+2.40 \end{Bmatrix} = \begin{Bmatrix} 0.2 \\ -0.2 \end{Bmatrix}$$

similarly we calculate x_4, x_5 and so on.

Newton's Method.

Consider the quadratic approximation of the function $f(x)$ at $x = x_i$ using the Taylor's series expansion

$$f(x) = f(x_i) + \nabla f_i^T (x - x_i) + \frac{1}{2} (x - x_i)^T [J_i] (x - x_i)$$

here $[J_i] = [\nabla]_{x_i}$ is the matrix of second partial derivatives (6.95)
at point x_i (Hessian Matrix).

By set the partial derivatives of Equation (6.95) equal to zero for minimum of $f(x)$, we obtain

$$\frac{\partial f(x)}{\partial x_j} = 0, \quad j = 1, 2, \dots, n \quad \text{--- (6.96)}$$

Equation (6.96) and (6.95) give

$$\nabla f = \nabla f_i + [J_i] (x - x_i) = 0 \quad \text{--- (6.97)}$$

If $[J_i]$ is non-singular, equation (6.97) can be solved to obtain an improved approximation ($x = x_{i+1}$) as

$$x_{i+1} = x_i - [J_i]^{-1} \nabla f_i \quad \text{--- (6.98)}$$

The sequence of points x_1, x_2, \dots, x_{i+1} can be shown to converge to the actual solution x^* from any initial point x_1 sufficiently close to the solution x^* , provided that $[J_i]$ is non-singular.

Q. Minimize $f(x_1, x_2) = x_1 x_2 + 2x_1^2 + 2x_1 x_2 + x_2^2$ by taking the starting point as $\{0\} = x_1$.

To find x_2 , we require $[J_1]^{-1}$

$$[J_1] = \begin{bmatrix} \frac{\partial f}{\partial x_1} & \frac{\partial f}{\partial x_2} \\ \frac{\partial f}{\partial x_2} & \frac{\partial f}{\partial x_1} \end{bmatrix} = \begin{bmatrix} 4 & 2 \\ 2 & 2 \end{bmatrix}$$

$$[J_1]^{-1} = \frac{1}{4} \begin{bmatrix} 2 & -2 \\ -2 & 4 \end{bmatrix} = \begin{bmatrix} \frac{1}{2} & -\frac{1}{2} \\ -\frac{1}{2} & 1 \end{bmatrix}$$

$$\begin{aligned} g_1 = \nabla f_i &= \left\{ \frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2} \right\}_{x_1} = \left\{ 1+4x_1+2x_2, -1+2x_1+2x_2 \right\}_{(0,0)} \\ &= \left\{ 1, -1 \right\} \end{aligned}$$

Now

$$\begin{aligned} x_2 &= x_1 - [J_1]^{-1} \nabla f_i = \{0\} - \begin{bmatrix} \frac{1}{2} & -\frac{1}{2} \\ -\frac{1}{2} & 1 \end{bmatrix} \begin{bmatrix} 1 \\ -1 \end{bmatrix} \\ &= \{0\} - \begin{bmatrix} \frac{1}{2} + \frac{1}{2} \\ -\frac{1}{2} - 1 \end{bmatrix} = \{-1, +3\} \end{aligned}$$

To see x_2 is the optimum point, we calculate

$$\begin{aligned} g_2 = \nabla f_i &= \left\{ \frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial x_2} \right\}_{(-1, 3)} = \left\{ 1+4x_1+2x_2, -1+2x_1+2x_2 \right\}_{(-1, 3)} \\ &= \{-1-4+3, -1-2+3\} = \{0\} \end{aligned}$$

as $\frac{\partial f}{\partial x_1} = 0$, x_1 at the optimum point.

Q. Minimize $f(x_1, x_2) = 100(x_1^2 - x_2)^2 + (1 - x_1)^2$ taking

$$\nabla f = \begin{Bmatrix} \frac{\partial f}{\partial x_1} \\ \frac{\partial f}{\partial x_2} \end{Bmatrix} = \begin{Bmatrix} 200(x_1^2 - x_2) \cdot 2x_1 + 2(1 - x_1) \cdot (-1) \\ 200(x_1^2 - x_2)(-1) \end{Bmatrix}$$
$$= \begin{Bmatrix} 400(4+2) \cdot (-2) + 2(1+2)(-1) \\ 200(4+2)(-1) \end{Bmatrix} = \begin{Bmatrix} -4806 \\ -1200 \end{Bmatrix}$$

Q₂ - $f = x_1^2 + x_2^2 - 2x_1 - 4x_2 + 5$, $x_1 = \begin{Bmatrix} 1 \\ 2 \end{Bmatrix}$ by Newton's
Method.

Univariate Method

Univariate method

can be summarized as follows

- ① choose an arbitrary starting point x_0 and set $i = 1$.
- ② Find the search direction s_i as

$$s_i^T = \begin{cases} (1, 0, 0, \dots, 0) & \text{for } i = 1, n+1, 2n+1, \dots \\ (0, 1, 0, \dots, 0) & \text{for } i = 2, n+2, 2n+2, \dots \\ (0, 0, 1, \dots, 0) & \text{for } i = 3, n+3, 2n+3, \dots \\ \vdots & \\ (0, 0, 0, \dots, 1) & \text{for } i = n, 2n, 3n, \dots \end{cases}$$

- ③ Determine whether d_i should be +ve or negative. For the current direction s_i , this means find whether the function value decreases in the +ve or -ve direction. For this we take a small probe length (ϵ) and evaluate $f_i = f(x_i)$, $f^+ = f(x_i + \epsilon s_i)$ and $f^- = f(x_i - \epsilon s_i)$. If $f^+ < f_i$, s_i will be the correct direction for decreasing the value of f and if $f^- < f_i$, $-s_i$ will be the correct one. If both f^+ and f^- are greater than f_i we take x_i as the minimum along the direction s_i .

- ④ Find the optimal step length d_i^* such that

$$f(x_i \pm d_i s_i) = \min_{d_i} (x_i \pm d_i s_i)$$

where + or - sign has to be used depending upon whether s_i or $-s_i$ is the direction for decreasing the function value.

- (5) Set $x_{i+1} = x_i \pm d_i s_i$ depending on the direction for decreasing the function value, and

$$f_{i+1} = f(x_{i+1}).$$

- (6) Set the new value of $i = i+1$ and go to Step 2.

Continue this procedure until no significant change is achieved in the value of the objective function.

Q. Minimize $f(\gamma_1, \gamma_2) = \gamma_1 - \gamma_2 + 2\gamma_1^2 + 2\gamma_1\gamma_2 + \gamma_2^2$

with the starting point $(0, 0)$

Let Probe length (ε) as 0.01 .

Iteration $i=1$

Step 1. Choose the search directions s_i as $s_i = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$

Step 2. To find whether the value of f decreases along s_i or $-s_i$, we use the probe length ε . Since

$$f_1 = f(x_1) = f(0, 0) = 0, \quad x_1 = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

$$f^+ = f(x_1 + \varepsilon s_i) = f\left(\begin{pmatrix} 0 \\ 0 \end{pmatrix} + \varepsilon \begin{pmatrix} 1 \\ 0 \end{pmatrix}\right) = f(\varepsilon)$$

$$= 0.01 - 0 + 2(0.0001) + 0 + 0 = 0.0102 > f_1$$

$$f^- = f(x_1 - \varepsilon s_1) = f(-\varepsilon, 0) = -0.01 - 0 + 2(0.0001)$$

$$-s_1 \text{ is the correct direction} \quad +0+0 = -0.9998 < f_1$$

To find the optimum step length.

$$f(x_1 + d_1 s_1) = f\left(\begin{pmatrix} 0 \\ 0 \end{pmatrix} + d_1 \begin{pmatrix} 1 \\ 0 \end{pmatrix}\right) = \cancel{\{f(0, -d_1)\}}$$

$$= f(-d_1, 0)$$

$$\text{as } \frac{\partial f}{\partial d_1} = -1 - 1 + 4d_1 = 0 \Rightarrow d_1 = \frac{1}{4} \text{ we have } d_1^* = \frac{1}{4}$$

$$\text{Set } x_2 = x_1 - d_1^* s_1 = \begin{pmatrix} 0 \\ 0 \end{pmatrix} - \frac{1}{4} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} -1/4 \\ 0 \end{pmatrix}$$

$$f_2 = f(x_2) = f(-1/4, 0) = -1/8$$

iteration i=2, choose the search direction s_2 as $s_2 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$

$$d_2 = f(x_2) = -\frac{1}{8} = -1/25$$

$$f^+ = f(x_2 + \varepsilon s_2) = f\left(\begin{pmatrix} -1/4 \\ 0 \end{pmatrix} + 0.01 \begin{pmatrix} 0 \\ 1 \end{pmatrix}\right)$$

$$= f(-0.25, 0.01) = -0.1399 < f_2$$

$$f^- = f(x_2 - \varepsilon s_2) = f(-0.125, -0.01) = -0.1099 > f_2$$

$\Rightarrow s_2$ is the correct direction for decreasing the value of f from x_2 .

We minimize $f(x_L + d_2 s_2)$ to find d_2^*
here

$$f(x_L + d_2 s_2) = f\left(\begin{pmatrix} -0.25 \\ 0 \end{pmatrix} + d_2 \begin{pmatrix} 0 \\ 1 \end{pmatrix}\right)$$

$$= f(-0.25, d_2) = d_2^2 - 1 \cdot s_{d_2} - 0.125$$

$$\frac{\partial f}{\partial d_2} = 2d_2 - 1 \cdot s = 0 \Rightarrow d_2^* = \frac{1.5}{2} = 0.75$$

Set

$$x_3 = x_L + d_2^* s_2 = \begin{pmatrix} -0.25 \\ 0 \end{pmatrix} + 0.75 \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$= \begin{pmatrix} -0.25 \\ 0.75 \end{pmatrix}$$

$$f(x_3) = -0.6875$$

Q. Minimize $= x_1 - x_2 + x_3 + 2x_1^L + 2x_2^2 - x_3^2 + 2x_1 x_3 + 4x_2 x_3 - 6x_1 x_2$

$$x_1 = \begin{pmatrix} 1 \\ 2 \end{pmatrix}, s_1 = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \epsilon = 0.01$$

$$d_1 = f(x_1) = 7$$

$$d_1^+ = 6.9502 < f_1$$

$$d_1^- = 7.0502$$

$\Rightarrow s_1$ is the correct direction

$-s_1$ is the correct direction

$$\Rightarrow x_3 = x_2 - d_2 s_2$$

$$= \begin{pmatrix} 9/5 \\ 2-d_2 \end{pmatrix}$$

~~Note~~ $f(x_2 + d_1 s_1) = f\left(\begin{pmatrix} 1+d_1 \\ 2 \end{pmatrix}\right) = 7 - 5d_1 + 2d_1^2$

$$\frac{\partial f(x_2 + d_1 s_1)}{\partial d_1} = -5 + 4d_1 \Rightarrow d_1 = 5/4$$

$$\Rightarrow x_2 = \begin{pmatrix} 9/4 \\ 2 \end{pmatrix}$$

$$f(x_2) = d_2 = \frac{107}{25} = 4.28$$

$$d_2^+ = 4.2822$$

$$d_2^- = 4.2782 < f(x_2) = f_2$$

Geometric Programming.

Polyomial - the objective function $f(x)$ is given by the sum of several component costs $v_i(x)$ as

$$f(x) = v_1 + v_2 + \dots + v_n$$

In many cases, the component cost v_i can be expressed as power functions of the type

$$v_i = c_i x_1^{a_{i1}} x_2^{a_{i2}} \dots x_n^{a_{in}}$$

the coefficients c_i are +ve constants, the exponents a_{ij} are real constants (+ve, zero, -ve) and the variables x_1, x_2, \dots, x_n are taken to be +ve. Functions of because of the coefficients and variables and real exponents are called polynomials. For example

$$f(x_1, x_2, x_3) = 6 + 3x_1 - 8x_2 + 7x_3 + 2x_1x_3 - 3x_1x_3 + \frac{4}{3}x_2x_3 + \frac{8}{7}x_1^2 - 9x_2^2 + x_3^2 \text{ is a second-degree polynomial in variables } x_1, x_2, x_3$$

while

$$g(x_1, x_2, x_3) = x_1x_2x_3 + x_1^2x_2 + 4x_3 + \frac{2}{x_1} + 5x_3^{-1/2}$$

is a polynomial

Unconstrained minimization Problem

unconstrained minimizing problem

Find $x = \{x_1, x_2, \dots, x_n\}$ that minimizes the objective function

$$f(x) = \sum_{j=1}^N v_j(x) = \sum_{j=1}^N (y_j x_1^{a_{1j}} x_2^{a_{2j}} \dots x_n^{a_{nj}}) \quad \text{--- (P.3)}$$

where $y_j > 0, a_{ij} > 0$ and a_{ij} is real const.

Solution of an unconstraint geometric problem
using Differential calculus,

as minimizing objective function is

$$f(x) = \sum_{j=1}^N v_j(x) = \sum_{j=1}^N (c_j x_1^{a_{1j}} x_2^{a_{2j}} \dots x_n^{a_{nj}})$$

For maxima or minima

$$\frac{\partial f(x)}{\partial x_k} = \sum_{j=1}^N \frac{\partial v_j}{\partial x_k} = 0 \quad k = 1, \dots, n, \text{ means } n \text{ variables.}$$

If we multiply

$$x_k \frac{\partial f}{\partial x_k} = \sum_{j=1}^N a_{kj} v_j(x) = 0, \quad k = 1, \dots, n \quad \text{--- (P.5)}$$

To find the minimizing vector $x^* = \{x_1^*, x_2^*, \dots, x_n^*\}$

we have

$$\sum_{j=1}^N a_{kj} v_j(x^*) = 0 \quad \Rightarrow \quad k = 1, 2, \dots, n \quad \text{--- (P.6)}$$

now, divide the (P.6) by f^* (minim value of f)

$$\sum_{j=1}^N a_{kj} \frac{v_j(x^*)}{f^*} = \sum_{j=1}^N a_{kj} \Delta_j^* \quad \text{--- (P.7)}$$

$$\text{where } \Delta_j^* = \frac{v_j(x^*)}{f^*}$$

this relation (P.7) is orthogonality condition

as

$$f(x) = \sum_{j=1}^N v_j(x)$$

$$\frac{\partial f}{\partial x_k} = \sum_{j=1}^N \frac{\partial v_j}{\partial x_k}$$

$$\Rightarrow \sum_{j=1}^N \frac{v_j(x)}{f^*} = 1 \Rightarrow \sum_{j=1}^N \Delta_j^* = 1$$

$$\Rightarrow \Delta_1 + \Delta_2 + \dots + \Delta_N = 1$$

this condition (P.8) is called the (P.8)
normality condition. now,

$$f^* = \left(\frac{v_1^*}{\Delta_1^*} \right)^{\Delta_1^*} \left(\frac{v_2^*}{\Delta_2^*} \right)^{\Delta_2^*} \dots \left(\frac{v_N^*}{\Delta_N^*} \right)^{\Delta_N^*} \quad \text{--- (P.12)}$$

or relation (8.12) can be written as

$$f^* = \left(\frac{c_1}{\Delta_1} \right)^{\alpha_1} \left(\frac{c_2}{\Delta_2} \right)^{\alpha_2} \left(\frac{c_3}{\Delta_3} \right)^{\alpha_3} \cdots \left(\frac{c_N}{\Delta_N} \right)^{\alpha_N}$$

$\therefore n = \text{no. of variables}, N = \text{no. of terms in the objective function}$

If $N = n+1$, there will be as many linear simultaneous equations as there are unknowns and we can find a unique solution.

If $N - n - 1 = 0$, the problem is said to have a zero degree of difficulty. If $N > n+1$, we have more no. of variables than the equations, then sometimes this method is not applicable.

Unknown α_j can be determined uniquely from the orthogonality and normality conditions.

Solve the problem

$$f(x) = 80x_1x_2 + 40x_2x_3 + 20x_1x_3 + \frac{80}{x_1x_2x_3} \quad \text{solve}$$

The problem, it is a general polynomial.

$$c_1 = 80, c_2 = 40, c_3 = 20, c_4 = 80$$

Compare with

$$\begin{aligned} f(x) &= \sum_{j=1}^N (c_j x_1^{q_{1j}} x_2^{q_{2j}} x_3^{q_{3j}} \cdots x_n^{q_{nj}}) \\ &= c_1 x_1^{q_{11}} x_2^{q_{21}} x_3^{q_{31}} + c_2 x_1^{q_{12}} x_2^{q_{22}} x_3^{q_{32}} \\ &\quad + c_3 x_1^{q_{13}} x_2^{q_{23}} x_3^{q_{33}} + c_4 x_1^{q_{14}} x_2^{q_{24}} x_3^{q_{34}} \end{aligned}$$

$$\begin{bmatrix} q_{11} & q_{12} & q_{13} & q_{14} \\ q_{21} & q_{22} & q_{23} & q_{24} \\ q_{31} & q_{32} & q_{33} & q_{34} \end{bmatrix} = \begin{bmatrix} 1 & 0 & 1 & -1 \\ 1 & 1 & 0 & -1 \\ 0 & 1 & 1 & -1 \end{bmatrix}$$

The orthogonality condition

The orthogonality and normality conditions are given by

$$\begin{bmatrix} 1 & 1 & 0 \end{bmatrix}$$

$$\sum_{j=1}^N \Delta_j q_{kj} = 0 \quad k=1 \dots n \quad (1.7)$$

$$\Rightarrow \Delta_1 q_{k1} + \Delta_2 q_{k2} + \Delta_3 q_{k3} + \Delta_4 q_{k4} = 0$$

$$\left\{ \begin{array}{l} \Delta_1 q_{11} + \Delta_2 q_{12} + \Delta_3 q_{13} + \Delta_4 q_{14} = 0 \\ \Delta_1 q_{21} + \Delta_2 q_{22} + \Delta_3 q_{23} + \Delta_4 q_{24} = 0 \\ \Delta_1 q_{31} + \Delta_2 q_{32} + \Delta_3 q_{33} + \Delta_4 q_{34} = 0 \end{array} \right.$$

$$\left. \begin{array}{l} \Delta_1 q_{41} + \Delta_2 q_{42} + \Delta_3 q_{43} + \Delta_4 q_{44} = 0 \\ \text{Orthogonal conditions} \end{array} \right\}$$

$$\Rightarrow \sum_{j=1}^N \Delta_j = 1 \quad \text{in orthogonal condition}$$

$$\Rightarrow \Delta_1 + \Delta_2 + \Delta_3 + \Delta_4 = 1. \quad (1)$$

$$\Delta_1 + 0 + \Delta_3 + -\Delta_4 = 0$$

$$\Delta_1 + \Delta_2 + 0 - \Delta_4 = 0$$

$$0 + \Delta_2 + \Delta_3 - \Delta_4 = 0$$

Solving these eqns

$$\Delta_4 = \Delta_1 + \Delta_3 = \Delta_1 + \Delta_2 \Rightarrow \Delta_2 = \Delta_3$$

$$\text{and } \Delta_4 = \Delta_1 + \Delta_3 = \Delta_2 + \Delta_3 \Rightarrow \Delta_1 = \Delta_2$$

$$\Rightarrow \Delta_1 = \Delta_2 = \Delta_3 \Rightarrow \Delta_4 = 2\Delta_1$$

$$\text{from eqn (1), } \cancel{\Delta_1 + \Delta_2 + \Delta_3} \quad \Delta_1 + \Delta_1 + \Delta_1 + 2\Delta_1 = 1$$

$$\Rightarrow \Delta_1 = \frac{1}{5} \Rightarrow \Delta_2 = \Delta_3 = \Delta_4 = \frac{1}{5}, \Delta_1 = \frac{2}{5}$$

So, optimal value of the objective function is

$$\begin{aligned} f^* &= \sum_{j=1}^N \left(\frac{y_j}{\Delta_j} \right) \Delta_j = \left(\frac{80}{1/5} \right)^{1/5} \left(\frac{40}{1/5} \right)^{1/5} \left(\frac{20}{1/5} \right)^{1/5} \left(\frac{80}{2/5} \right)^{1/5} \\ &= (400)^{1/5} \times (200)^{1/5} \times (100)^{1/5} \times (800)^{1/5} \\ &= (400 \times 200 \times 100 \times 800)^{1/5} \\ &= (32 \times 10^{10})^{1/5} = 200 \end{aligned}$$

Now to find y_1, y_2, y_3, \dots

$$v_j^* = \Delta_j f^*$$

$$\Rightarrow v_1^* = 80 \Delta_1 f^* = \Delta_1 f^* = \frac{1}{5} \times 200 = 40$$

$$v_2^* = 40 \Delta_2 f^* = \Delta_2 f^* = \frac{1}{5} \times 200 = 40$$

$$v_3^* = 20 \Delta_3 f^* = \Delta_3 f^* = \frac{1}{5} \times 200 = 40$$

$$v_4^* = \frac{80}{\Delta_4} f^* = \Delta_4 f^* = \frac{2}{5} \times 200 = 80$$

Solving these eqn we get

$$\gamma_2^* = \frac{1}{2} \gamma_1^* \Rightarrow \frac{1}{\gamma_2^*} = 2 \gamma_1^*$$

$$\gamma_3^* = \frac{1}{\gamma_2^*} = \frac{1}{2 \gamma_1^*}$$

$$\gamma_3^* = \frac{2}{\gamma_1^*}$$

$$\Rightarrow \frac{2}{\gamma_1^*} = \frac{1}{\gamma_2^*} = 2 \gamma_1^* \Rightarrow \gamma_1^* = 1 \Rightarrow \gamma_1^* = 1.$$

$$\gamma_2^* = \frac{1}{2}, \quad \gamma_3^* = 2$$

Solution of an unconstrained Geometric Problem using
Arithmetic-Geometric inequality

Geometrical Primal Problem (Unconstrained)

$$\text{Find } X = \left\{ \begin{array}{c} x_1 \\ x_2 \\ \vdots \\ x_n \end{array} \right\} \text{ so that } M_{\text{Primal}}(X) = \sum_{j=1}^N c_j x_1^{a_{1j}} x_2^{a_{2j}} \cdots x_n^{a_{nj}}$$

$$x_1 > 0, x_2 > 0, \dots, x_n > 0$$

then Geometric dual of Primal Problem is

$$\text{Find } \Delta = \left\{ \begin{array}{c} \Delta_1 \\ \vdots \\ \Delta_N \end{array} \right\} \text{ so that}$$

$$\text{Max } V(\Delta) = \prod_{j=1}^N \left(\frac{c_j}{\Delta_j} \right)^{\Delta_j} \quad \text{or}$$

$$\log \{ \text{Max } V(\Delta) \} = \log \left[\prod_{j=1}^N \left(\frac{c_j}{\Delta_j} \right)^{\Delta_j} \right]$$

subject to the constraint

$$\sum_{j=1}^N \Delta_j = 1$$

$$\sum_{j=1}^N a_{ij} \Delta_j = 0, \quad i = 1, 2, \dots, n$$

dual and Primal of Bimetric Problem

dual problem/bimetric dual.

$$\text{Find } x = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$$

so that

$$g_0(x) = f(x) \rightarrow \text{minimum}$$

subject to constraints

$$g_1(x) \leq 1$$

$$g_2(x) \leq 1$$

:

$$g_m(x) \leq 1$$

with

$$g_0(x) = \sum_{j=1}^{N_0} c_{0j} x_1^{a_{01j}} x_2^{a_{02j}} \dots x_n^{a_{0nj}}$$

$$g_1(x) = \sum_{j=1}^{N_1} c_{1j} x_1^{a_{11j}} x_2^{a_{12j}} \dots x_n^{a_{1nj}}$$

$$g_2(x) = \sum_{j=1}^{N_2} c_{2j} x_1^{a_{21j}} x_2^{a_{22j}} \dots x_n^{a_{2nj}}$$

$$g_m(x) = \sum_{j=1}^{N_m} c_{mj} x_1^{a_{m1j}} x_2^{a_{m2j}} \dots x_n^{a_{mnj}}$$

$$\text{Find } d = \begin{pmatrix} d_{01} \\ d_{02} \\ \vdots \\ d_{0N_0} \end{pmatrix}$$

$$d = \begin{pmatrix} d_{11} \\ d_{12} \\ \vdots \\ d_{1N_1} \\ d_! \\ \vdots \\ d_{mN_m} \end{pmatrix}$$

so that

$$d_{kj} \quad k=0, j=1 \dots N_k$$

$$v(d) = \prod_{k=0}^m \prod_{j=1}^{N_k} \left(\frac{c_{kj}}{d_{kj}} \right) \Rightarrow \text{maximum}$$

subject to the constraints

$$d_{01} \geq 0, d_{02} \geq 0 \dots d_{0N_0} \geq 0, d_{11} \geq 0, d_{12} \geq 0$$

$$d_{1N_1} \geq 0 \dots d_{mN_m} \geq 0 \quad \therefore d_{mN_m} \geq 0$$

$$\sum_{j=1}^{N_0} d_{0j} = 1$$

$$\sum_{k=0}^m \sum_{j=1}^{N_k} a_{kij} d_{kj} = 0, i = 1, 2, \dots, n$$

(c_{kj} are +ve and a_{kij} are real numbers)

here $K = 1 \text{ to } m$
 $n = \text{total no. of variables}$

$g_0 = f$ = Primal function

$m = \text{no. of Primal constraints}$

$N = N_0 + N_1 + \dots + N_m = \text{total number of terms in the}$

number of terms in the

Primal function

$N - n - 1 = \text{degree of difficulty}$

of problem

$v = \text{dual function}$
 $d_{01}, d_{02}, \dots, d_{mN_m} = \text{dual var}$

in the normality constraint

and the orthogonality

constraints

$$d_{kj} \geq 0, j = 1, 2, \dots, N_k$$

$k = 0 \text{ to } m$

$$N = N_0 + N_1 + \dots + N_m$$

number of dual variables

$n+1$ number of dual constraints

Solve the Problem

$$J(D, \alpha) = 100D^4\alpha^4 + 50D^3\alpha^6 + 20D^2\alpha^8 + 250D^0\alpha^2$$

Here, $C_1 = 100, C_2 = 50, C_3 = 20, C_4 = 250$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \end{bmatrix} = \begin{bmatrix} 1 & 2 & 0 & -5 \\ 0 & 0 & -1 & 2 \end{bmatrix}$$

The orthogonality and normality conditions.

$$\begin{bmatrix} 1 & 2 & 0 & -5 \\ 0 & 0 & -1 & 2 \\ 1 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} D_1 \\ D_2 \\ D_3 \\ D_4 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \quad \text{---(1)}$$

here $N = 4, n = 2, N > n+1 = 2+1$

\Rightarrow these equations (1) do not yield the required D_j ($j=1 \text{ to } 4$). So, solving

D_1, D_2, D_3 in term of D_4 from (1)

$$D_1 + 2D_2 - 5D_4 = 0$$

$$-D_3 + 2D_4 = 0$$

$$D_1 + D_2 + D_3 + D_4 = 1$$

from (1) $D_3 = 2D_4$

$$D_1 = -2D_2 + 5D_4$$

$$-2D_2 + 5D_4 + D_2 + 2D_4 + D_4 = 1$$

$$-D_2 + 8D_4 = 1$$

$$\Rightarrow D_2 = 8D_4 - 1$$

so

$$D_1 = \cancel{-2D_2} - \cancel{2(1-8D_4)} + \cancel{5D_4}$$

$$= \cancel{-2} + 16D_4 + 5D_4 =$$

$$D_1 = -2(8D_4 - 1) + 5D_4 = -16D_4 + 2 + 5D_4$$

$$\Rightarrow \boxed{\begin{aligned} D_1 &= 2 - 11D_4 \\ D_2 &= 8D_4 - 1 \\ D_3 &= 2D_4 \end{aligned}}$$

the dual problem can be written as

Maximize $V(D_1, D_2, D_3, D_4)$

$$= \left(\frac{C_1}{D_1}\right)^{D_1} \left(\frac{C_2}{D_2}\right)^{D_2} \left(\frac{C_3}{D_3}\right)^{D_3} \left(\frac{C_4}{D_4}\right)^{D_4}$$

$$= \left(\frac{100}{2-11D_4}\right)^{2-11D_4} \left(\frac{50}{8D_4-1}\right)^{8D_4-1} \left(\frac{20}{2D_4}\right)^{2D_4} \left(\frac{300}{D_4}\right)^{D_4}$$

taking log on both sides, get

$$\log V = (g - 11D_4) [\ln 100 - \ln(2 - 11D_4)] + (8D_4 - 1) [\ln 50 - \ln(8D_4 - 1)] + g D_4 [\log 2.0 - \log 2D_4] + D_4 [\ln 300 - \ln D_4]$$

necessary condition for maximization

$$\frac{\partial}{\partial D_4} (\ln V) = -11 [\ln 100 - \ln(2 - 11D_4)] + \frac{11}{2 - 11D_4} + 8 [\ln 50 - \ln(8D_4 - 1)] + \frac{8}{8D_4 - 1} + g [\ln 2.0 - \ln 2D_4] + 2D_4 \left[\frac{-1}{2D_4} \right] + 1 [\ln 300 - \ln D_4] + D_4 \left[-\frac{1}{D_4} \right] = 0$$

$$\Rightarrow -11 [2 - \ln(2 - 11D_4)] + 11 + 8 \left[\ln \frac{50}{8D_4 - 1} \right] - 8 + 2 \log \left(\frac{g_0}{g D_4} \right) - 2 + \log \frac{300}{D_4} - 1 = 0$$

$$\Rightarrow \cancel{\ln}$$

$$\Rightarrow -\ln \left[\frac{(100)^8}{(50)^8 (20)^2 (300)} \right] + \ln \left[\frac{(2 - 11D_4)^8}{(8D_4 - 1)^8 (2D_4)^2 D_4} \right] = 0$$
 ~~$\cancel{2 - 11D_4 = 0}$~~

$$\Rightarrow \frac{(2 - 11D_4)^8}{(8D_4 - 1)^8 (2D_4)^2 D_4} = \frac{(100)^8}{(50)^8 (20)^2 (300)}$$

$$= \frac{(1 \times 10^2)^8}{5^8 \times 10^8 2^2 10^2 \times 3 \times 10^2} = \frac{1 \times 10^{22}}{5^8 \times 10^{12} \times 12}$$

$$= \frac{10^{10}}{5^8 \times 12} = 2130$$

here get value of D_4 by trial method.

$$D_4^* \approx 0.147, D_1^* = 0.385, D_2^* = 0.175$$

$$D_3 = 0.294$$

$$\Rightarrow V^* = f^* = \left(\frac{100}{0.385} \right)^{0.385} \left(\frac{50}{0.175} \right)^{0.175} \left(\frac{20}{0.294} \right)^{0.294} \times \left(\frac{300}{0.147} \right)^{0.147} = 242$$

$$V_1^* = D_1^* f^* = 0.385 \times 242 = 92.2$$

$$U_2^* = 42.4$$

$$U_3^* = 71.1$$

$$U_4^* = 35.6$$

$$\text{Given } V_i = 100 \text{ D}^2 = 92.2 \\ \Rightarrow D^2 = 0.922. D^2 = 0.281 \text{ m}^3/\text{kg. N}$$

Example 8.3 - zero degree of difficulty Problem
the optimization problem can be stated as

Find $X = \begin{Bmatrix} x_1 \\ x_2 \\ x_3 \end{Bmatrix}$ so as to minimize

$$f(X) = 20x_1x_3 + 40x_2x_3 + 80x_1x_2 \text{ subject to} \\ \frac{80}{x_1x_2x_3} \leq 10 \text{ or } \frac{8}{x_1x_2x_3} < 1$$

Ans here, $n = \text{no. of variables} = 3$
 $N_0 = \text{no. of terms in the objective function} = 3$

$$N_1 = 1$$

$m = \text{total no. of the constraint} = 1$

$N_k = \text{number of terms in } k\text{th constraint}$

mean

$N_1 = 1 = \text{no. of terms in 1st constraint.}$

$N = N_0 + N_1 + \dots + N_m = \text{Total number of terms in}$
 $\text{the polynomials, mean there}$

$$N = 3 + 1 = 4$$

and $N - n - 1 = 4 - 3 - 1 = 0 \Rightarrow \text{zero-degree}$

of difficulty problem.

so, dual problem can be stated as

Find $d = \begin{Bmatrix} d_{01} \\ d_{02} \\ d_{03} \\ d_{11} \end{Bmatrix}$ to maximize

$$\begin{aligned}
 v(d) &= \prod_{k=0}^1 \prod_{j=1}^{N_k} \left(\frac{c_{kj}}{d_{kj}} \sum_{l=1}^{N_k} d_{kl} \right)^{d_{kj}} \\
 &= \prod_{j=1}^{N_0=3} \left(\frac{c_{0j}}{\sum_{l=1}^{N_0} d_{0l}} \right)^{d_{0j}} \prod_{j=1}^{N_1=1} \left(\frac{c_{1j}}{\sum_{l=1}^{N_1} d_{1l}} \right)^{d_{1j}} \\
 &= \prod_{j=1}^3 \left(\frac{c_{0j}}{\sum_{l=1}^{N_0} d_{0l}} (d_{01} + d_{02} + d_{03}) \right)^{d_{0j}} \prod_{j=1}^{N_1=1} \left(\frac{c_{1j}}{\sum_{l=1}^{N_1} d_{1l}} (d_{11}) \right)^{d_{1j}} \\
 &= \frac{c_{01}}{d_{01}} (d_{01} + d_{02} + d_{03})^{d_{01}} \times \\
 &\quad \frac{c_{02}}{d_{02}} (d_{01} + d_{02} + d_{03})^{d_{02}} \times \\
 &\quad \frac{c_{03}}{d_{03}} (d_{01} + d_{02} + d_{03})^{d_{03}} \left(\frac{c_{11}}{d_{11}} d_{11} \right)^{d_{11}}
 \end{aligned}$$

(E)

subject to the constants are.

$$\sum_{j=1}^{N_0} d_{0j} = 1 \rightarrow \text{Normality constraint. constraint} \\
 \Rightarrow d_{01} + d_{02} + d_{03} = 1$$

and orthogonal constraints are

$$\sum_{k=0}^m \sum_{j=1}^{N_k} a_{kij} d_{kj} = 0, \quad i = 1, 2, \dots, n$$

$$\Rightarrow \sum_{k=0}^1 \sum_{j=1}^{N_k} a_{kij} d_{kj} = 0$$

$$\sum_{k=0}^1 \sum_{j=1}^{N_0} a_{0ij} d_{0j} = 0$$

$$\sum_{j=1}^{N_0} a_{0ij} d_{0j} + \sum_{j=1}^{N_1} a_{1ij} d_{1j} = 0$$

$$\sum_{j=1}^3 a_{0ij} d_{0j} + \sum_{j=1}^1 a_{1ij} d_{1j} = 0$$

$$a_{011} d_{01} + a_{012} d_{02} + a_{013} d_{03} + a_{111} d_{11} = 0$$

$$a_{021} d_{01} + a_{022} d_{02} + a_{023} d_{03} + a_{121} d_{11} = 0$$

$$q_{031}d_{01} + q_{032}d_{02} + q_{033}d_{03} + q_{131}d_{11} = 0$$

$$d_{ij} \geq 0, j=1,2,3, d_{ii} \geq 0$$

In this Problem

$$q_{011} = 1, q_{021} = 0, q_{031} = 0,$$

$$q_{012} = 0, q_{022} = 1, q_{032} = 1, q_{013} = 1, q_{023} = 1,$$

$$q_{033} = 0, q_{111} = -1, q_{121} = -1, q_{131} = -1.$$

So, Problem can be written as

$$V(d) = \left[\frac{20}{d_{01}} (d_{01} + d_{02} + d_{03}) \right]^{d_{01}}$$

subject to

$$d_{01} + d_{02} + d_{03} = 1$$

$$d_{01} + d_{03} - d_{11} = 0$$

$$d_{02} + d_{03} - d_{11} = 0$$

$$d_{01} + d_{02} - d_{11} = 0$$

$$\Rightarrow d_{01} = d_{02} = d_{03} = \frac{1}{3}, d_{11} = \frac{2}{3}.$$

Then the maximum value of V or minimum

value of x_0 is given by

$$V = x_0^{\frac{1}{3}} = (60)^{\frac{1}{3}} (120)^{\frac{1}{3}} (240)^{\frac{1}{3}} (8)^{\frac{1}{3}} = 400 \text{ P.M.}$$

Complementary Geometric Programming

Geometric Programming to include any rational function of polynomial terms and called the method of complementary geometric programming.

Let the complementary geometric programming problem be stated as follows

minimize $R_0(x)$ subject

$$R_K(x) \leq 1 \quad K = 1, 2, \dots, m, \text{ where}$$

$$R_K(x) = \frac{A_K(x) - B_K(x)}{C_K(x) - D_K(x)}, \quad K = 0, 1, 2, \dots, m$$

where $A_K(x), B_K(x), C_K(x)$ and $D_K(x)$ are polynomial and possibly some of them may be absent. To solve the problem stated, we introduce a new variable $x_0 > 0$, constrained to satisfy the relation $x_0 \geq R_0(x)$

i.e. $\frac{R_0(x)}{x_0} \leq 1$ so the problem may be

stated as

minimize π_0

$$\text{subject to } \frac{A_k(x) - B_k(x)}{C_k(x) - D_k(x)} \leq 1, \quad k = 0, 1, 2, \dots, m$$

where $A_0(x) = R_0(x)$, $C_0(x) = \pi_0$, $B_0(x) = 0$
and $D_0(x) = 0$.

thus any complementary geometric programming problem (ChP) can be stated in the standard form

minimize π_0 subject to

$$\frac{P_k(x)}{Q_k(x)} \leq 1, \quad k = 1, 2, \dots, m \quad \dots (8.71)$$

$$x = \begin{Bmatrix} x_0 \\ x_1 \\ \vdots \\ x_n \end{Bmatrix} \geq 0, \text{ where} \quad \dots (8.72)$$

where $P_k(x)$ and $Q_k(x)$ are polynomials of

the form

$$P_k(x) = \sum_{j=0}^n c_{kj} \prod_{i=0}^n (x_i)^{a_{kij}} = \sum_j P_{kj}(x) \quad \dots (8.73)$$

$$Q_k(x) = \sum_j d_{kj} \prod_{i=0}^n (x_i)^{b_{kij}} = \sum_j Q_{kj}(x) \quad \dots (8.74)$$

$$P_k(x) = \sum_j (c_{kj}) (x_1)^{a_{k1j}} x_2^{a_{k2j}} x_3^{a_{k3j}}$$

$$= \sum_j P_{kj}(x)$$

$$= P_{k1}(x) + P_{k2}(x) + P_{k3}(x) + \dots$$

$$= c_{k1} (x_1)^{a_{k11}} (x_2)^{a_{k21}} (x_3)^{a_{k31}}$$

$$+ c_{k2} (x_1)^{a_{k12}} (x_2)^{a_{k22}} (x_3)^{a_{k32}}$$

$$Q_k(x) = Q_{k1}(x) + Q_{k2}(x) + Q_{k3}(x)$$

$$Q_1(x) = q_{11}(x) + q_{12}(x) + q_{13}(x)$$

① Minimize x_1 subject to

$$-4x_1^2 + 4x_1 \leq 1$$

$$x_1 + x_2 \geq 1, x_1 \geq 0, x_2 \geq 0$$

or L.P. can be stated as

minimize x_1 subject to ①

$$+x_2 \leq 1 + 4x_1^2 \quad ②$$

$$\frac{4x_1}{1+4x_1^2} \leq 1 \quad ③$$

$$\frac{1}{x_1} + \frac{4}{x_1} \leq 1 \quad x_1 \geq 1$$

$$\Rightarrow \frac{x_1 + x_2 - 1}{x_1 x_2} \leq \frac{1}{4}$$

$$\Rightarrow \frac{1/x_1}{1 + x_1/x_1} \leq \frac{x_1^{-1}}{1 + x_1^{-1}/4} \leq 1 \quad ④$$

* now we start the process that in eq ② and ④, denominators convert into single term

initial starting

Point $x^{(1)} = \{1\}$
(assume float)

2(a) Minimized $f(x_1, x_2) = 2x_1^2 + x_2^2 \quad \left\{ \begin{array}{l} 1 \\ 2 \end{array} \right\}$ Univariate method

Iteration → $S_1 = \{1\}$, $X_1 = \{1\}$

$$f_1 = f(X_1) = f(1, 2) = 2 \cdot 1^2 + 2^2 = 6, \epsilon = 0.01$$

$$f^+ = f(X_1 + \epsilon S_1) = f(1 + \epsilon, 2) = f(1.01, 2)$$

$$= 2(1.01)^2 + 4 = 2.0402 + 4 = 6.0402$$

here $f^+ > f_1$

$$f^- = f(X_1 - \epsilon S_1) = f(1, 2 - \epsilon(1, 0)) = f(1 - \epsilon, 2)$$

$$= f(0.99, 2) = 5.9602 < f_1 \quad f^+ > f_1$$

⇒ $-S_1$ will be the correct direction

For optimum length, we minimize

~~$f(X_1 - d_1 S_1) = f(1 - d_1, 2)$~~

$$= 2(1 - d_1)^2 + 4 = 2(1 + d_1^2 - 2d_1) + 4$$

$$\frac{\partial f}{\partial d_1} = 2(2d_1 - 2) = 0 \Rightarrow d_1 = 1$$

$$\therefore X_2 = X_1 - d_1 S_1 = (1, 2) - 1(1, 0) = (0, 2)$$

$$\Rightarrow f(X_2) = 0 + 2^2 = 4$$

2nd iteration, $S_2 = \{0\}$

$$f_2^+ = f(X_2 + \epsilon S_2) = f(0, 2) + 0.01(0, 1) = f(0, 2.01)$$

$$= 0.0 + 4.0401 = f_2^+ > f_2$$

$$f_2^- = f(X_2 - \epsilon S_2) = f(0, 2) - 0.01(0, 1) = f(0, 1.99)$$

$$= 3.9601 \Rightarrow f_2^- < 4 = f_2$$