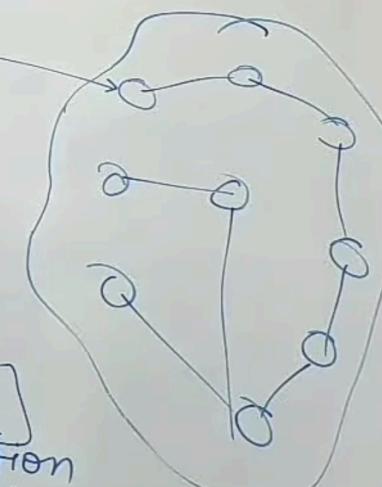


## Wireless Sensor Networks

Sensor  
node

→ SU  
→ CU  
→ PU  
→ SU  
→ ADC

Options  
Location  
Power





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## Wireless Sensor N/W (WSN):

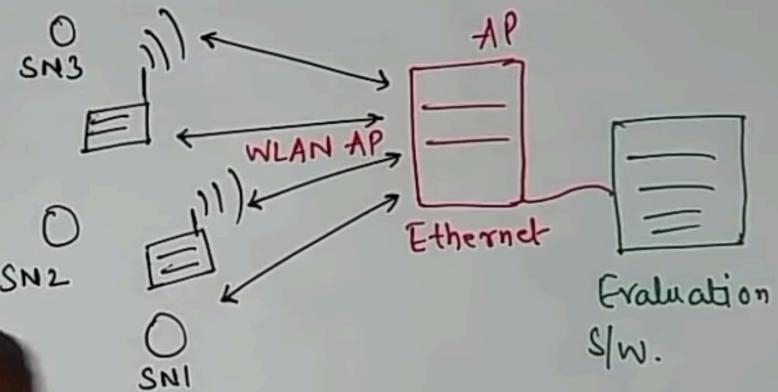
### Components:

(i) Sensors: They capture the measured variable in a data acquisition N/W. For further processing, sensor signal is converted into

### (ii) wireless Sensors

They receive signals from Sensors.

(iii) WLAN Access Points: which are wirelessly connected to



(iv) Evaluation S/w: for any data analysis, WLAN access point is connected to an evaluation unit.



## Wireless Sensor N/W (WSN):

### Components:

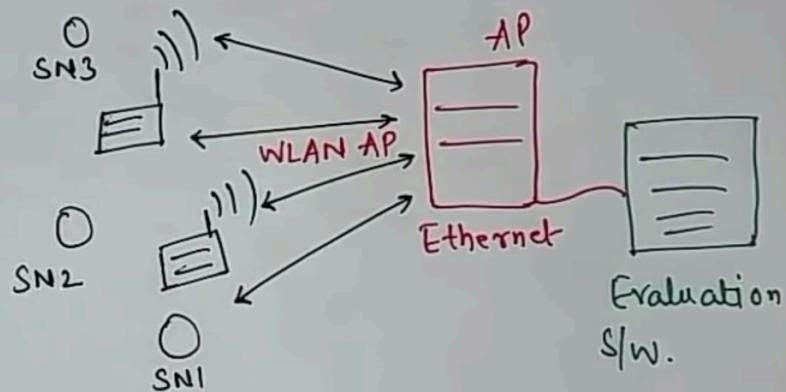
(i) Sensors: They capture -the measured variable in a data acquisition N/W.

for further processing sensor signal is converted into electrical Signal.

(ii) wireless sensor nodes or Radio nodes:

They receive -the Sensor data from Sensors.

(iii) WLAN Access Point: Receives sensor data which are transferred by sensor nodes wirelessly.



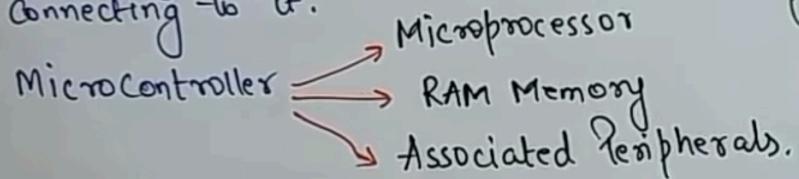
(iv) Evaluation S/w: for any data analysis, WLAN access point is connected to an evaluation unit.



## Components of a wireless Sensor Node:

(i) Microcontroller: This is a Computer-on-a chip which is very small in size.

↳ Capable of doing powerful tasks including controlling the func<sup>n</sup> of other devices connecting to it.



(ii) Transceiver: This is transmitter-receiver that is used for comm<sup>n</sup> purposes -to send and receive data, and commands.

(iv) External Memory: (WSN) nodes usually use flash memories.

Small in Size

Reasonable storage capacity.

(v) Power Source: Power is stored in the form of batteries.



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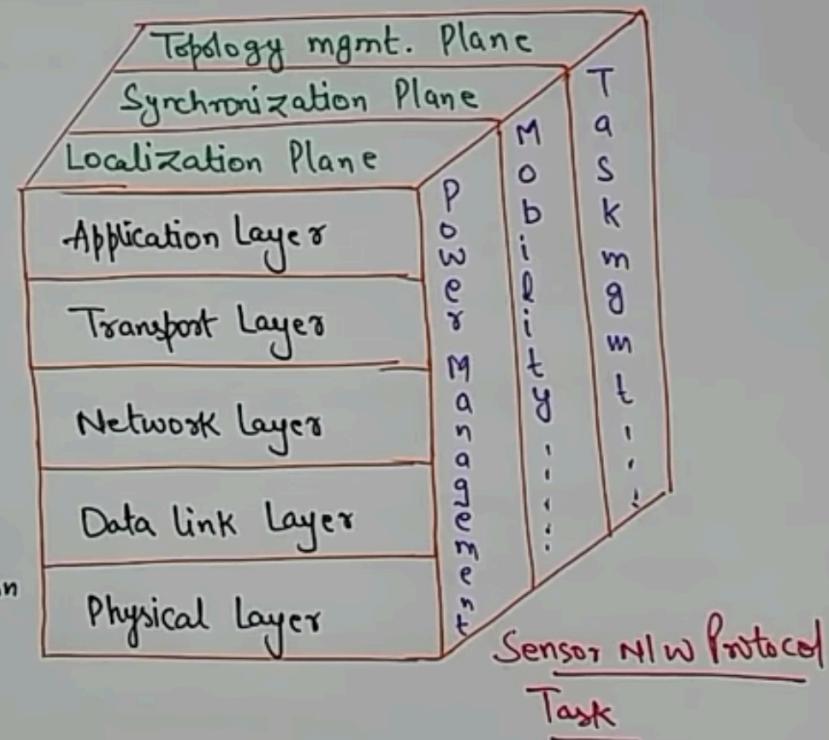
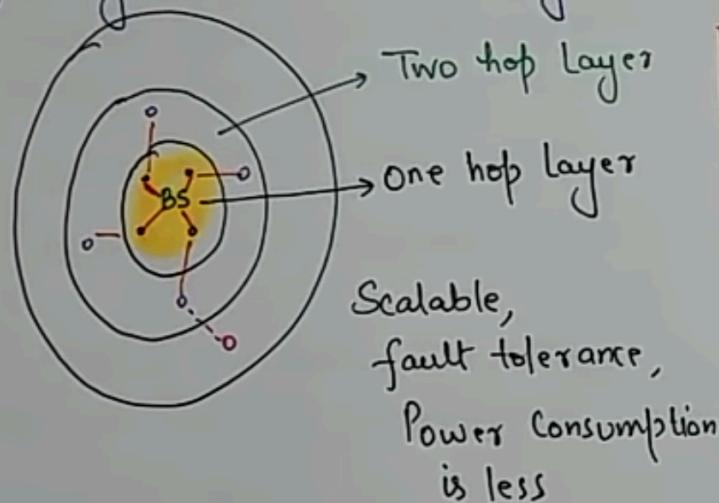
### Adhoc Wireless NW Vs. Sensor NW:      Issues and Challenges in designing Sensor NW:

- i) No. of nodes in Sensor NW can be more as compared to adhoc NW.
- ii) Sensor Nodes are more prone to failure and energy drain.
- iii) Sensor Nodes → Data Centric  
Adhoc NW → Address Centric
- iv) Most of the adhoc nw routing protocols can't be implemented to Sensor NW.
- (i) No Regular topology.
- (ii) Infrastructure less.
- (iii) Resource Constraints
  - Energy
  - bandwidth.
- (iv) H/W Should be energy efficient.  
↳ microcontroller.
- (v) Security Issues.
  - Enemy territory.
  - No central tower.
- (vi) Adaptable to changing connectivity  
Synchronization.



## Sensor NW Architecture:

(i) Layered NW Architecture: It includes five Layers and three cross Layers.



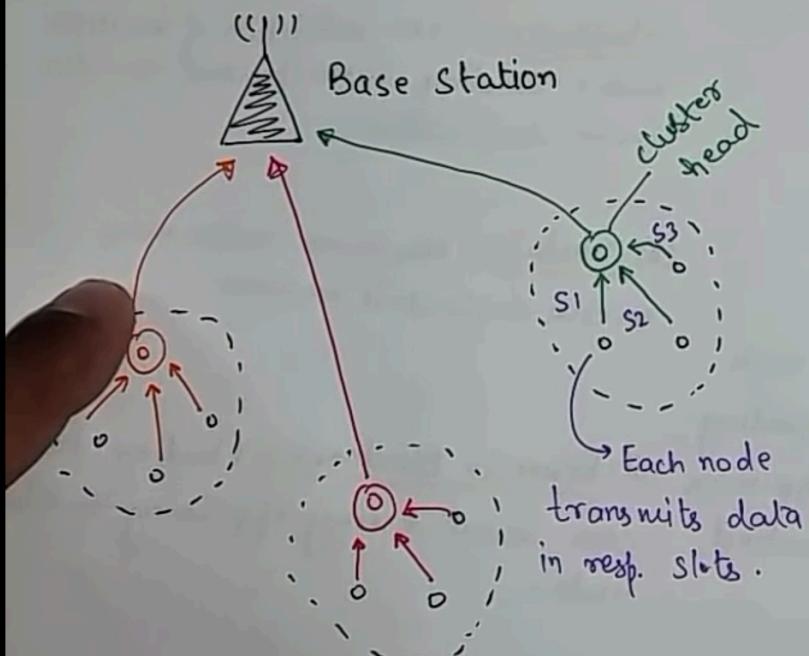


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### (ii) Clustered NW Architecture:



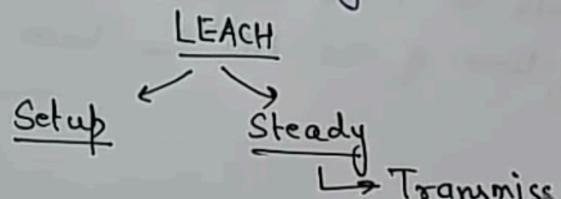
LEACH Protocol: Low Energy Adaptive clustering hierarchy.

↳ 2-tier hierarchy clustering architecture

↳ distributed algo to organize the Sensor nodes into clusters.

↳ cluster head nodes create TDMA Schedules.

↳ Energy Efficiency = Data fusion.





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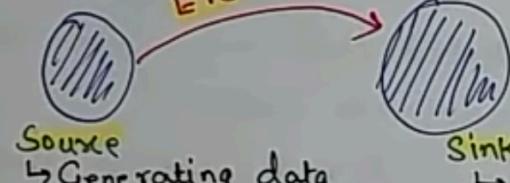
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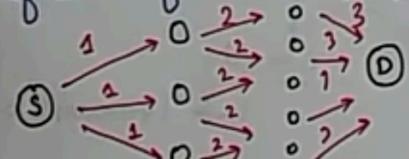
Data Dissemination: It is -the process by which queries or data are routed in -the Sensor NW.

"Info" to be reported / Sent .

Event



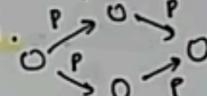
FLOODING: Simplest Design. In this each node receiving data repeats it by broadcasting -the data -to every neighbour unless -the max hop lifetime of -the data has been reached.



- ③ ↗ man
- hop life time.

ShortComings of flooding:

(i) Implosion:- No restriction of multiple nodes sending same packet -to -the Same destination.

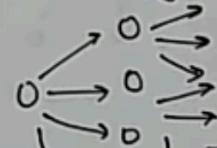


(ii) Overlap :- Neighbour nodes may receive duplicated message.

↳ overlapping regions.



(iii) Resource Blindness :- Flooding doesn't care about energy efficiency of the nodes.





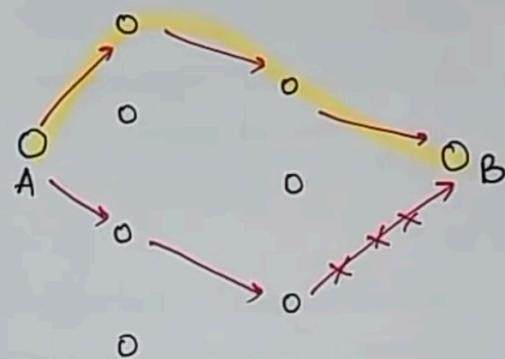
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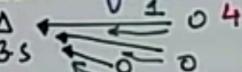
Gossiping: It is -the enhancement of flooding. In this, when a node receives data, it randomly chooses a neighbour and sends -the data -to it.

- ↳ It avoids -the problem of Implosion,
- ↳ Contributes -to -the latency of NW.





Data Gathering in Sensor NW: Objective is to transmit the sensed data from each Sensor node to a BS.

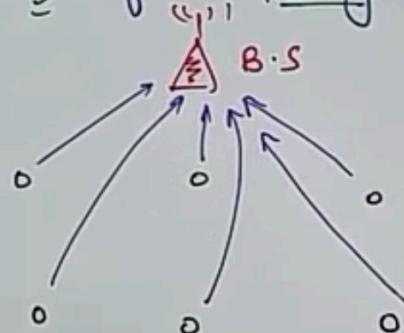


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Algorithm Goal = Maximize the no. of rounds before node dies and NW becomes inoperable.

↳ Minimize Energy and reduce delay.

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↳ expensive in terms of energy consumed. (Energy ↑)  
↳ Media Delay is also large  
↳ Performs Poorly.



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Enhancement over (LEACH) protocol.

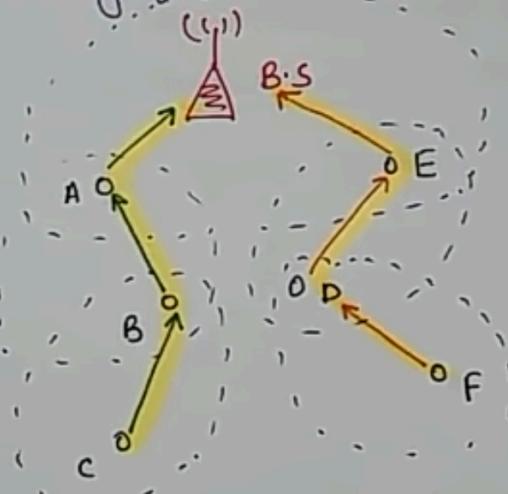
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- ↳ Equal distribution of energy consumption across all nodes.

#### Data Gathering with Regasis:

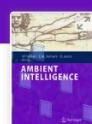
- ↳ Greedy algo is used.
- Chain of Sensor nodes is constructed starting from -the farthest node from BS.



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## TinyOS: An Operating System for Sensor Networks

[P. Levis, S. Madden, J. Polastre, R. Szewczyk, K. Whitehouse, A. Woo, D. Gay, J. Hill, M. Welsh, E. Brewer & D. Culler](#)

Chapter

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### Abstract

We present TinyOS, a flexible, application-specific operating system for sensor networks, which form a core component of ambient intelligence systems. Sensor networks consist of (potentially) thousands of tiny, low-power nodes, each of which execute concurrent, reactive programs that must operate with severe memory and power constraints. The sensor network challenges of limited resources, event-centric concurrent applications, and low-power operation drive the design of TinyOS. Our solution combines flexible, fine-grain components with an execution model that supports complex yet safe concurrent operations. TinyOS meets these challenges well and has become the platform of choice for sensor network research; it is in use by over a hundred groups worldwide, and supports a broad range of applications and research topics. We provide a qualitative and quantitative evaluation of the system, showing that it supports complex, concurrent programs with very low memory requirements (many applications fit within 16KB of memory, and the core OS is 400 bytes) and efficient, low-power operation. We present our experiences with TinyOS as a platform for sensor network innovation and applications.

### Keywords

[Sensor Network](#) [Wireless Sensor Network](#) [Power Management](#) [Execution Model](#) [Data Race](#)

*These keywords were added by machine and not by the authors. This process is experimental and the keywords may be updated as the learning algorithm improves.*

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## Network Layer

This layer is used to send datagrams from the source network to the destination network. we use IPv4 and IPv6 protocols as host identification that transfers data in packets.

### IPv4

This is a protocol address that is a unique and numerical label assigned to each device connected to the network. an IP address performs two main functions host and location addressing. IPv4 is an IP address that is 32-bit long.

### IPv6

It is a successor of IPv4 that uses 128 bits for an IP address. it is developed by the IETF task force to deal with long-anticipated problems.

## Link Layer

Link-layer protocols are used to send data over the network's physical layer. it also determines how the packets are coded and signaled by the devices.

### Ethernet

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## Ethernet

It is a set of technologies and protocols that are used primarily in LANs. it defines the physical layer and the medium access control for wired ethernet networks.

## WiFi

It is a set of LAN protocols and specifies the set of media access control and physical layer protocols for implementing wireless local area networks.

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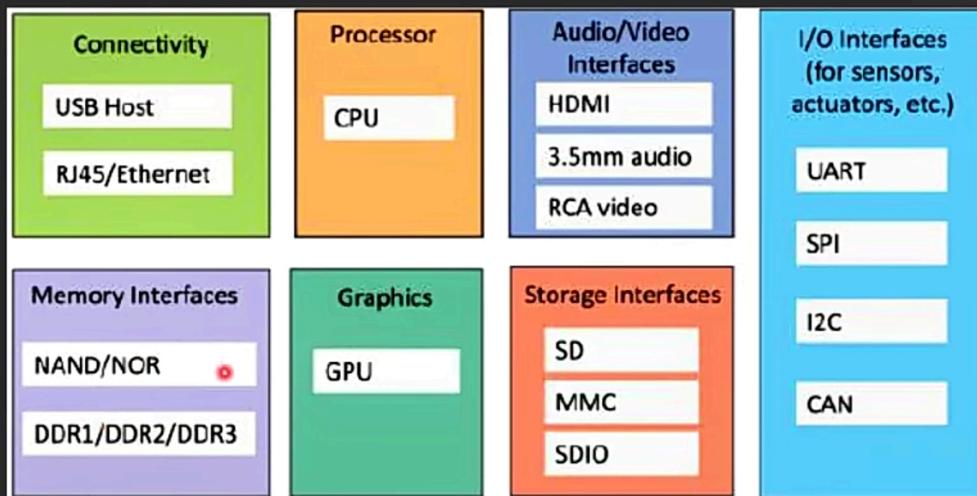
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# Physical Design of IoT - Generic Block Diagram



## Logical Design of IOT

IOT Functional Block

IOT Communicational Block

IOT Communicational API's

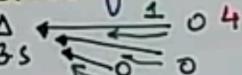


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Data Gathering in Sensor NW: Objective is to transmit the sensed data from each Sensor node to a BS.

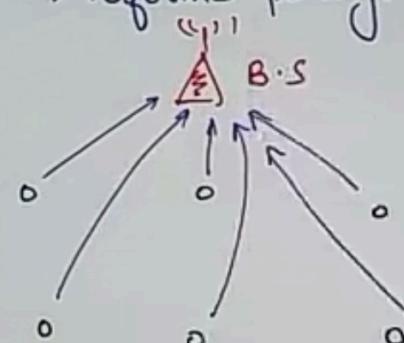


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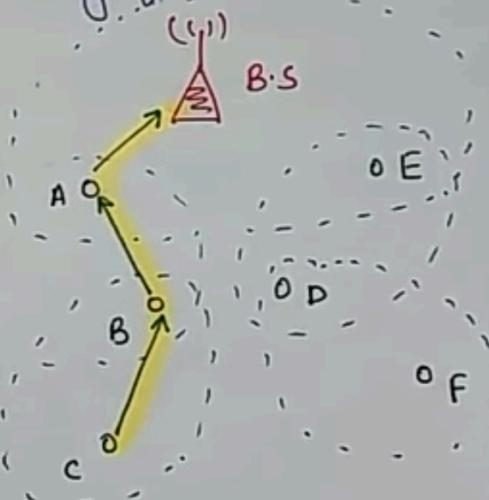
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iii) Binary Scheme: Chain-based Scheme, iv) Chain-Based Three Level Scheme:  
which nodes are classified into different levels.

not applicable for non-CDMA Sensor nodes. At each level no. of nodes are reduced from pr. level.

S<sub>1</sub> S<sub>2</sub> S<sub>3</sub> S<sub>4</sub> S<sub>5</sub> S<sub>6</sub> S<sub>7</sub> B.S

S<sub>1</sub> → S<sub>2</sub> S<sub>3</sub> → S<sub>4</sub> S<sub>5</sub> → S<sub>6</sub> S<sub>7</sub> B.S

(L3) S<sub>2</sub> → S<sub>4</sub> S<sub>6</sub> → S<sub>7</sub>

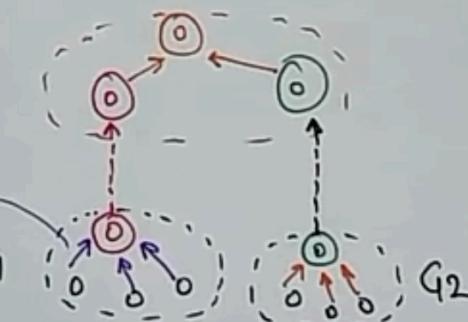
(L4) S<sub>4</sub> → S<sub>7</sub>

S<sub>7</sub> → B.S

Each node transmits alone.

It is applicable for non- CDMA Sensor nodes.

↪ CHAIN is constructed, but chain is divided into no. of groups to space out simultaneous transmission.



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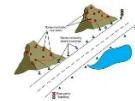
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## About: Virtual sensor network

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A virtual sensor network (VSN) in computing and telecommunications is an emerging form of collaborative wireless sensor networks. In contrast to early wireless sensor networks that were dedicated to a specific application (e.g., target tracking), VSNs enable multi-purpose, collaborative, and resource efficient WSNs. The key idea difference of VSNs is the collaboration and resource sharing. By doing so nodes achieve application objectives in a more resource efficient way. These networks may further involve dynamically varying subset of sensor nodes (e.g., when the phenomenon migrates sensors that detect the phenomenon changes with time) and/or users (users that are accessing the network changes with time). A VSN can be formed by providing logical connectivity among collaborative sensors. Nod



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telecommunications is an emerging form of collaborative wireless sensor networks.<sup>[1]</sup> In contrast to early wireless sensor networks that were dedicated to a specific application (e.g., target tracking), VSNs enable multi-purpose, collaborative, and resource efficient WSNs. The key idea difference of VSNs is the collaboration and **resource sharing**. By doing so nodes achieve application objectives in a more resource efficient way. These networks may further involve dynamically varying subset of **sensor nodes** (e.g., when the **phenomenon** migrates sensors that detect the phenomenon changes with time) and/or users (users that are accessing the network changes with time).

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# TinyOS

[Article](#) [Talk](#)

This article needs additional citations for verification.  
(December 2006) [Learn more](#)

**TinyOS** is an embedded, component-based [operating system](#) and platform for low-power wireless devices, such as those used in [wireless sensor networks](#) (WSNs), [smartdust](#), [ubiquitous computing](#), [personal area networks](#), [building automation](#), and [smart meters](#). It is written in the programming language nesC, as a set of cooperating tasks and processes. It began as a collaboration between the [University of California, Berkeley](#), [Intel Research](#), and [Crossbow Technology](#), was released as [free](#) and [open-source software](#) under a [BSD license](#), and has since grown into an international consortium, the [TinyOS Alliance](#).



TinyOS applications are written in the programming language [nesC](#), a dialect of the [C language](#) optimized for the memory limits of sensor networks. Its supplementary tools are mainly in the form of [Java](#) and [shell script](#) front-ends. Associated libraries and tools, such as the nesC compiler and [Atmel AVR](#) binutils toolchains, are mostly written in C.

TinyOS programs are built of [software components](#), some of which present hardware abstractions. Components are connected to each other using [interfaces](#). TinyOS provides interfaces and components for common abstractions such as packet communication, routing, sensing, actuation and storage.

TinyOS is fully [non-blocking](#): it has one [call stack](#). Thus, all [input/output](#) (I/O) operations that last longer than a few hundred [microseconds](#) are asynchronous and have a [callback](#). To enable the native [compiler](#) to better optimize across call boundaries, TinyOS uses nesC's features to link these callbacks, called events, statically. While being non-blocking enables TinyOS to maintain high concurrency with one stack, it forces programmers to write complex logic by stitching together many small event handlers. To support larger computations, TinyOS provides tasks, which are similar to a [Deferred Procedure Call](#) and



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TOSThreads are unmaintained and have been deprecated.<sup>[1]</sup>

TinyOS code is statically linked with program code and is compiled into a small binary, using a custom [GNU toolchain](#). Associated utilities are provided to complete a development platform for working with TinyOS.

