

Representation of Sets

⇒ Roaster or Tabular

⇒ Rule or Set builder.

Set

Finite

Disjoint

Infinite

Null

Singleton

Subset

Universal

Operation on set

⇒ Union ' \cup '

⇒ Intersection ' \cap '

⇒ Complement ' A^c / A' ' = $U - A$

⇒ Relative Complement ($A - B$)

↳ Relative complement of

B w.r.t A

⇒ Symmetric Difference $A \Delta B$ or $A \oplus B = (A - B) \cup (B - A)$

e.g. $A = \{-3, 0, 1, 2\}$ $B = \{1, 2, 3, 4\}$

$A \oplus B = \{-3, 0, 3, 4\}$

Algebra of Sets

⇒ Idempotent law

$$A \cap A = A$$

$$A \cup A = A$$

\Rightarrow Associative law

$$(A \cup B) \cup C = A \cup (B \cup C)$$

$$(A \cap B) \cap C = A \cap (B \cap C)$$

\Rightarrow Commutative law

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A$$

\Rightarrow Distributive law

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

\Rightarrow Identity

$$A \cup \phi = A$$

$$A \cup A = A$$

$$A \cap \phi = \phi$$

$$A \cap A = A$$

\Rightarrow Involution law

$$(A')' = A$$

\Rightarrow Complement law

$$A \cup A' = U$$

$$A \cap A' = \phi$$

\Rightarrow D'Morgan's law

$$(A \cup B)' = A' \cap B'$$

$$(A \cap B)' = A' \cup B'$$

$$(A \cup B)' = A' \cap B'$$

$$\text{L.H.S: } (A \cup B)'$$

$$\text{Let } x \in (A \cup B)'$$

$$x \in U - (A \cup B)$$

$$\Rightarrow x \in U \text{ & } x \notin A \cup B$$

$$\Rightarrow x \in U \text{ & } (x \notin A \text{ & } x \notin B)$$

$$\Rightarrow (x \in U \text{ & } x \notin A) \text{ & } (x \in U \text{ & } x \notin B)$$

$$\Rightarrow x \in U - A \text{ & } x \in U - B$$

$$\Rightarrow x \in A' \text{ & } x \in B'$$

$$\Rightarrow x \in A' \cap B'$$

$$\Rightarrow (A \cup B)' \subseteq A' \cap B' - \textcircled{1} \text{ & } (A \cup B)' \supseteq A' \cap B' - \textcircled{2}$$

From $\textcircled{1} \text{ & } \textcircled{2}$

$$(A \cup B)' = A' \cap B'$$

$$2) (A \cap B)' = A' \cup B'$$

$$\text{L.H.S: } (A \cap B)'$$

$$\text{Let } x \in (A \cap B)'$$

$$\Rightarrow x \in U - (A \cap B)$$

$$\Rightarrow x \in U \text{ & } x \notin A \cap B$$

$$\Rightarrow x \in U \text{ & } (x \notin A \text{ or } x \notin B)$$

$$\Rightarrow (x \in U \text{ & } x \notin A) \text{ or } (x \in U \text{ & } x \notin B)$$

$$\Rightarrow x \in U - A \text{ or } x \in U - B$$

$$\Rightarrow x \in A' \text{ or } x \in B'$$

$$\Rightarrow x \in A' \cup B'$$

$$(A \cap B)' \subseteq A' \cup B' - \textcircled{1} \text{ & } (A \cap B)' \supseteq A' \cup B' - \textcircled{2}$$

From $\textcircled{1} \text{ & } \textcircled{2}$

$$(A \cap B)' = A' \cup B'$$

Cardinal No: or Cardinality ($|A| / n(A)$)

Properties of Cardinal no.

$$* |A \cup B| = |A| + |B| - |A \cap B|$$

$$* |A - B| = |A| - |A \cap B|$$

$$* |B - A| = |B| - |B \cap A|$$

$$* |A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |B \cap C| - |C \cap A| + |A \cap B \cap C|$$

Power set $|P(A)| = 2^{|A|}$

Theorem :- If $A \subseteq B$ then $P(A) \subseteq P(B)$

Countable Set

* Finite & infinite set with one-one correspondence with natural numbers

e.g: $\mathbb{Z} = \{0, \pm 1, \pm 2, \pm 3, \dots\}$
 ↳ countable

\aleph_0 (Aleph not)

$$|N| = |\mathbb{Z}| = |\mathbb{Q}| = \aleph_0$$

$$* \aleph_0 + \aleph_0 = \aleph_0$$

$$* \aleph_0 \times \aleph_0 = \aleph_0$$

$$|A| = \aleph_0 \quad |P(A)| = 2^{\aleph_0} = c \text{ (Cardinal No)} \\ \text{(Uncountable)}$$

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Theorem: Union of countable families of countable sets is countable

Proof: Let $\{A_1, A_2, A_3, \dots, A_n\}$ be countable family of countable sets

$$A_1 = \{a_{11}, a_{12}, a_{13}, \dots, \cancel{a_{1n}}\}$$

$$A_2 = \{a_{21}, a_{22}, a_{23}, \dots, \cancel{a_{2n}}\}$$

$$A_3 = \{a_{31}, a_{32}, a_{33}, \dots, \cancel{a_{3n}}\}$$

$$\vdots \quad \vdots \quad \vdots \quad \vdots \quad \vdots \quad \vdots$$

$$A_n = \{a_{n1}, a_{n2}, a_{n3}, \dots, \cancel{a_{nn}}\}$$

then $\bigcup_{i=1}^n A_i$ can be found in following way

$$A_1 = a_{11}, \cancel{a_{12}}, \cancel{a_{13}}, \dots$$

$$A_2 = \cancel{a_{21}}, a_{22}, \cancel{a_{23}}, \dots$$

$$A_3 = \cancel{a_{31}}, \cancel{a_{32}}, \cancel{a_{33}}, \dots$$

$$A_4 = \cancel{a_{41}}, a_{42}, \cancel{a_{43}}, \dots$$

List the elements a_{11}

$$a_{21} \quad a_{12} \quad \dots$$

$$a_{31} \quad a_{22} \quad a_{13} \quad \dots$$

$$a_{41} \quad a_{32} \quad a_{23} \quad a_{14} \quad \dots$$

$$a_{n1} \quad a_{(n-1)2} \quad a_{(n-2)3} \quad a_{(n-3)4} \quad \dots$$

It is clear that a_{ij} is the j th element of $(i+j-1)$ th row. Thus all the ele. have been counted. Thus $\bigcup_{i=1}^n A_i$ is countable.

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Theorem :- Show $N \times N$ is countable

Proof :- $(1,1), (1,2), (1,3), \dots$
 $(2,1), (2,2), (2,3), \dots$
 $(3,1), (3,2), (3,3), \dots$
 $(4,1), (4,2), (4,3), \dots$

list the elements $(1,1)$

$(2,1), (1,2)$
 $(3,1), (2,2), (1,3)$
 $(4,1), (3,2), (2,3), (1,4)$

Or it is clear that a_{ij} is the j th ele of row. Thus all ele. have been counted.

Thus $\bigcup_{i=1}^n A_i$ is countable.

If A & B are countable then $A \times B$ is countable.

Relation

$R \subseteq A \times B = \{(a_1, b_1), (a_2, b_2), (a_3, b_3), \dots, (a_n, b_n)\}$
 $\emptyset \in R$ (void)
 $A \times B \subseteq R$ (universal)

Total no of relation = 2^{mn}

$$|A| = m \quad |B| = n$$

$$|A \times B| = mn$$

Types of Relation

1. Inverse Relation

R be a relation from A to B .

R^{-1} be a relation from B to A

i.e. $R^{-1} = \{(b, a) : (a, b) \in R\}$

or, $x R y \Rightarrow y R^{-1} x$

e.g. $A = \{2, 3, 5\}$: $B = \{6, 12, 10\}$

$\forall (x, y) \in A \text{ or } B, (x, y) \in R \Leftrightarrow x \text{ divide } y$

$R \& R^{-1} = 2$

$2 R 6, 2 R 8, 2 R 10, 3 R 6, 5 R 10$

$R = \{(2, 6), (2, 8), (2, 10), (3, 6), (5, 10)\}$

$R^{-1} = \{(6, 2), (10, 2), (6, 3), (10, 5)\}$

$\text{Domain}(R) = \text{Range}(R^{-1}) = \{2, 3, 5\}$

$\text{Domain}(R^{-1}) = \text{Range}(R) = \{6, 8, 10\}$

2. Identity Relation

$I_A = \{(x, x) : x \in A\}$

e.g. $A = \{1, 2, 3\}$

$I_A = \{(1, 1), (2, 2), (3, 3)\}$

3. Reflexive Relation

$(a, a) \in R \quad \forall a \in A$

4. Irreflexive.

$$(a,a) \notin R \text{ i.e. } aRa \nvdash aRa.$$

5. Symmetry

$$(a,b) \in R \Rightarrow (b,a) \in R$$

$$\text{i.e. } aRb \Rightarrow bRa \quad \forall a, b \in A$$

6. Asymmetric

$$(a,b) \in R \Rightarrow (b,a) \notin R$$

$$\text{i.e. } aRb \Rightarrow b \not R a \quad \nvdash (b,a) \in R$$

7. Antisymmetric

$$(a,b) \in R \wedge (b,a) \in R$$

$$\Rightarrow a=b$$

$$\text{i.e. } aRb \wedge bRa \Rightarrow a=b \quad \nvdash (a,b) \in A$$

8. Transitive

$$(a,b) \in R \wedge (b,c) \in R \Rightarrow aRc$$

$$\quad \forall a, b, c \in A$$

e.g.

Equivalence Relation

If R is relation in \mathbb{Z} defined by

$$R = \{(x,y) : y-y \leq 2 ; (y-x) \text{ divisible by 6}\}$$

Soln

Reflexive: $xRx \Rightarrow x-x \text{ is divisible by 6}$

$$\Leftrightarrow 0 \text{ is divisible by 6}$$

Symmetric: $xRy \Rightarrow x-y \text{ is divisible by 6}$
 $\Rightarrow -(y-x) \text{ is divisible by 6}$
 $\Rightarrow yRx$

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Tranitive: $x-y$ divisible by 6
 $\Rightarrow y-z$ divisible by 6.
 $\Rightarrow x-z$ divisible by 6

Note:

- ① R is reflexive $\Rightarrow R^{-1}$ is reflexive
- ② R is symmetric $\Leftrightarrow R^{-1}$ is symmetric
- ③ R is antisymmetric $\Leftrightarrow R \cap R^{-1} \subseteq \Gamma$
- ④ $R \in S$ are reflexive $\Rightarrow R \circ S \in R \circ S$ are also reflexive.
- ⑤ $R \in S$ are symmetric $\Rightarrow R \circ S \in R \circ S$ are also symmetric.
- ⑥ $R \in S$ are transitive $\Rightarrow R \circ S$ is transitive.
- ⑦ $R \in S$ are equivalence $\Rightarrow R \circ S$ is equivalence
- ⑧ R is equivalence $\Rightarrow R^{-1}$ is equivalence

Partial Order Rel.

A relation R on set S . is partial order
if reflexive: $aRa ; \forall a \in S$

Antisymmetric: $aRb \wedge bRa \Rightarrow a=b \quad \forall a, b \in S$

Transitive: $aRb \wedge bRc \Rightarrow aRc \quad \forall a, b, c \in S$

A set S together with partial order R
i.e. (S, R) is partial order

Ex ' \geq ' relation is partial order on \mathbb{Z}

Reflexive: $a \geq a$

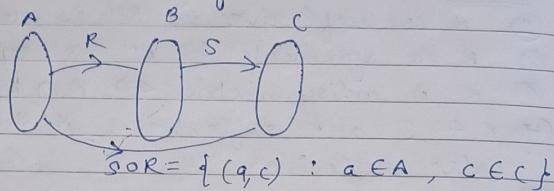
Antisymmetric: $a \geq b \wedge b \geq a \Rightarrow a = b$

Transitive: $a \geq b \wedge b \geq c \Rightarrow a \geq c$

(\mathbb{Z}, \geq)

Composite Relation

If $A, B \& C$ are non empty sets & R be a relation from A to B & S be a relation from B to C then composite relation of $R \& S$ is relation from A to C



$$\text{Ex. } A = \{1, 2, 3\}; B = \{p, q, r\}, C = \{x, y, z\}$$

$$R = \{(1, p), (1, q), (2, q), (3, q)\}$$

$$S = \{(p, x), (q, y), (q, z)\}$$

~~Ans.~~
 $\text{SOR} = \{(1, x), (1, y), (2, x), (3, x)\}$

Function

- Bijective \leftarrow
 - One-one (Injective) $f(x_1) = f(x_2) \Rightarrow x_1 = x_2$
 - Many-one \rightarrow co-domain $x_1 \neq x_2 \Rightarrow f(x_1) = f(x_2)$
 - Onto (Surjective)
 - Into \rightarrow co-domain all diff ele. are not cover.

Composition of Function

Let $f: A \rightarrow B$ & $g: B \rightarrow C$.

Composition of $f \circ g$ $gof: A \rightarrow C$ defines as
 $gof(x) = g[f(x)] \quad \forall x \in A$

Ques $A = \{1, 2, 3\}$; $B = \{a, b\}$; $C = \{r, s\}$
 $f: A \rightarrow B$ as $f(1) = a$; $f(2) = a$; $f(3) = b$.
 $g: B \rightarrow C$ as $g(a) = r$; $g(b) = s$.
 $gof: A \rightarrow C$

$$\begin{aligned} g(f(1)) &= g(f(1)) = g(a) = s. \\ g(f(2)) &= g(a) = s \\ g(f(3)) &= g(b) = s \end{aligned}$$

Property \Rightarrow Associative: $f \circ (g \circ h) = (f \circ g) \circ h$
 $A \rightarrow B \quad B \rightarrow C \quad C \rightarrow D$

$f: A \rightarrow B$ $g: B \rightarrow C$ $h: C \rightarrow D$
 $gof: A \rightarrow C$ $h \circ g: B \rightarrow D$
Hence. $h \circ (gof): A \rightarrow D$
& $(h \circ g) \circ f: A \rightarrow D$.

Dom. of $[h \circ (gof)] = \text{dom. } [(h \circ g) \circ f]$
let $x \in A$; $y \in B$, $z \in C$
i.e. $f(x) = y$ & $g(y) = z$
then

$$\begin{aligned} [h \circ g] \circ f (x) &= h(z) && \text{--- (1)} \\ [h \circ (gof)] (x) &= h(z) && \text{--- (2)} \end{aligned}$$

$$[(h \circ g) \circ f] = [h \circ (gof)]$$

Theorem:

- Let $f: A \rightarrow B$ & $g: B \rightarrow C$
- If f and g are injection then gof is injection
 - If f & g are surjection then gof is surjection

PROOF :-

(a) Let $a_1, a_2 \in A$.
 We have $(gof)a_1 = (gof)a_2$
 $\Rightarrow g[f(a_1)] = g[f(a_2)]$
 $\Rightarrow f(a_1) = f(a_2)$ [g is injection]
 $a_1 = a_2$ [f is injection]
 Thus, gof is injection

(b) Let $c \in C$, then we can find $a \in A$
 s.t. $(gof)(a) = c$
 g is onto $C \Rightarrow b \in B$ s.t. $g(b) = c$
 then since f is onto B , there exist $a \in A$ s.t.
 $f(a) = b$
 Then $(gof)(a) = g[f(a)] = g(b) = c$.

Inverse

Let $f: A \rightarrow B$ then $g: B \rightarrow A$ if
 $gof = I_A$ & $fog = I_B$

Q Show that $f(x) = x^3$ & $g(y) = \sqrt[3]{y}$ $\forall x \in \mathbb{R}$
are inverse to each other

$$fog(x) = f[g(x)] = f[\sqrt[3]{x}] = x = f(x)$$

$$gof(x) = g[f(x)] = g(x^3) = x = f(x)$$

Theorem:

$f: A \rightarrow B$ is one-one & onto then
 $f^{-1}: B \rightarrow A$ is one-one & onto

Proof:

Given :- $f: A \rightarrow B$ is one-one & onto
let $a_1, a_2 \in A$ & $b_1, b_2 \in B$
so $b_1 = f(a_1)$ & $b_2 = f(a_2)$
& $a_1 = f^{-1}(b_1)$; $a_2 = f^{-1}(b_2)$.

f is one-one $f(a_1) = f(a_2) \Leftrightarrow a_1 = a_2$
 $b_1 = f(a_1) \Leftrightarrow f^{-1}(b_1) = f^{-1}(b_2)$.
i.e. $f^{-1}(b_1) = f^{-1}(b_2) \Rightarrow b_1 = b_2$
 $\therefore f^{-1}$ is one-one

f is onto :-

all element at B is associated with
unique ele. at A .

i.e.

For any $a \in A$ is pre image of $b \in B$
where $b \in f(a) \Rightarrow a \in f^{-1}(b)$
i.e. for $b \in B$ f^{-1}
So f^{-1} is onto

Principle of Mathematical Induction

- 1) $s(n)$ is true for $n=1, 2, 3$ i.e. $s(1)$ is true.
- 2) $s(k)$ is true $\Rightarrow s(k+1)$ is true.

Qn $1+2+3+\dots+n = \frac{n(n+1)}{2} = EN$

For $n=1$

$$s(1) = \frac{1(1+1)}{2} = 1$$

For $n=k$

$$s(k) = \frac{k(k+1)}{2}$$

For $n=k+1$, we have to prove

$$s(k+1) = \frac{(k+1)(k+2)}{2}$$

$$\begin{aligned} \Rightarrow 1+2+3+\dots+k+(k+1) &= \frac{k(k+1)+k+1}{2} \\ &= (k+1)\left[\frac{k}{2}+1\right] \\ &= \frac{(k+1)(k+2)}{2} \end{aligned}$$

$$\text{Q} \quad 1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{1}{6}(n+1)(2n+1)$$

For $n=1$

$$S(1) = \frac{1}{6}(2)(3)$$

$$= 1$$

For $n=k$

$$S(k) = \frac{k}{6}(k+1)(2k+1)$$

For $n=k+1$

$$1^2 + 2^2 + 3^2 + \dots + k^2 + (k+1)^2 = \frac{k}{6}(k+1)(2k+1) + (k+1)^2$$

$$= (k+1) \left[\cancel{\frac{k(2k+1)}{6}} + \cancel{k+1} \right]$$

$$= (k+1) \left[\cancel{\frac{2k(2k+1)}{6}} + \cancel{6k+6} \right]$$

$$= \frac{1}{6}(k+1) [2k^2 + k + 6]$$

=

$$= (k+1) \left[\frac{k(2k+1)}{6} + k+1 \right]$$

$$= (k+1) \left[\frac{2k^2 + k + 6k + 6}{6} \right]$$

$$= \frac{(k+1)}{6} [2k^2 + 7k + 6]$$

$$= \frac{(k+1)}{6} [2k^2 + 4k + 3k + 6]$$

$$= \frac{(k+1)}{6} [2k(k+2) + 3(k+2)] = \frac{(k+1)}{6} [(k+2)(2k+3)]$$

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Q. Prove by mathematical induction
 $6^{n+2} + 7^{2n+1}$ is divisible by 43 $\forall n \in \mathbb{Z}$

Proof:

For $n=1$
 $6^{1+2} + 7^{2+1} = 6^3 + 7^3$

$= 559$
which is divisible by 43.

For $n=k$

$$\Rightarrow 6^{k+2} + 7^{2k+1} \text{ is divisible by 43}$$
$$6^{k+2} + 7^{2k+1} = 43m, \forall m \in \mathbb{Z}$$

For $n=k+1$

$$\begin{aligned} & 6^{k+3} + 7^{2k+3} \\ &= 6^{k+2+1} + 7^{2k+1+2} \\ &= 6^{k+2} \cdot 6 + 7^{2k+1} \cdot 7^2 \\ &= 6 \cdot [6^{k+2} + 7^{2k+1}] (43+6) \\ &= 6 \left[6^{k+2} + 7^{2k+1} \right] + 43 \cdot 7^{2k+1} \\ &= 6 \cdot 43m + 43 \cdot 7^{2k+1} \\ &= 43 \left[6m + 7^{2k+1} \right] \end{aligned}$$