

Collection of fields → Record or tuple

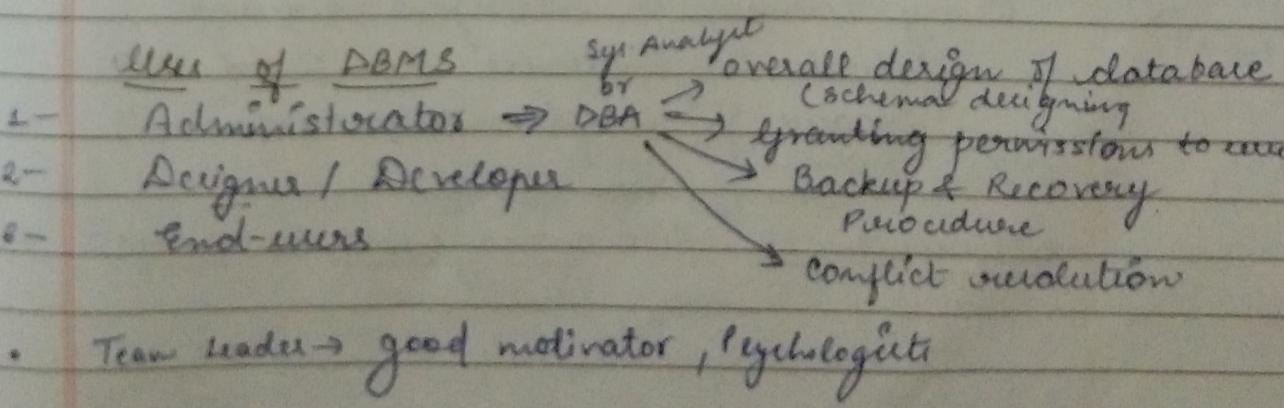
Collection of records → file

Collection of all files → database

- * Database is collection of interrelated data
 - * DBMS → Database + set of program to access those database.
- Retrieval, Insertion,
deletion, updation

Characteristics of DBMS

1. Real World Object
2. Relation based table mean structured data
3. Isolation of data & Application → more file can run without affecting other files.
4. less redundancy
5. Consistency - DBF must be in original state
6. Query language → SQL
7. ACID properties
 ↳ Atomicity, consistency, Isolation, Durability
8. Multiple user & concurrent execution
9. Multiple views
10. Security → backup & recovery mechanism



Advantages of Database BMS

- Control redundancy
- Data Sharing work done in parallel
- Reduced time / Increased throughput
- Multitask - only in read mode can access at one time
- Backup
- Security

Disadvantages

- Costly hardware & software - size
- Higher impact of failure. - complexity

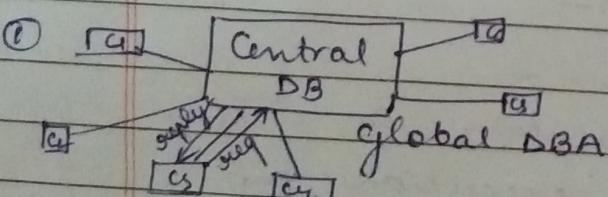
Classification of DBMS / Types

① - Centralized DB

Dumb Terminal

② - Distributed DB

↓
without CPU



It connects user to access stored data from diff. loc. through own app.

disadv

- ① Not easy to update with an extensive → dec. risk of data loss
 Authentication → data consistency maintained as it manages data in central repository

Every terminal is intelligent terminal

Difficult to implement distributed DB because of complex structure

Advantage: - Modular structure possible
 → server fail will not affect entire data set

Difference b/w File management & DBMS

FMS

- ⇒ Small sized program
- Written in C, C++ & Java

DBMS

large sized system
e.g. Oracle, Sybase etc.

⇒ Relatively cheaper

Relatively expensive.

⇒ Structure is simple

⇒ Complex structure.

⇒ Single user

⇒ Multiuser

⇒ Requires less preliminary knowledge of ^{design} design

Requires more preliminary knowledge of design

⇒ less secure

More secure

⇒ Simple backup & recovery process

Strong backup & recovery procedure.

Database System Architecture

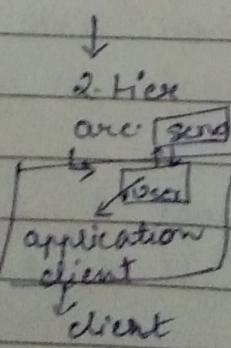
↓
1-tier Architecture

⇒ based on client/server

1-tier



User

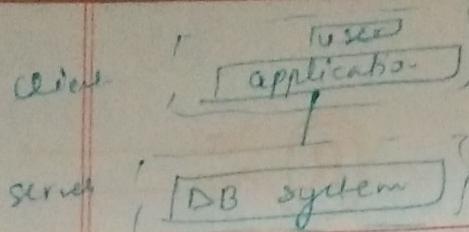


One client/server

client

↓
3-tier
arc.

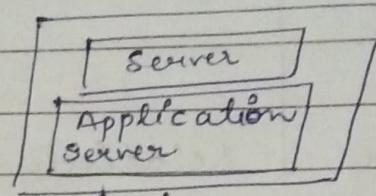
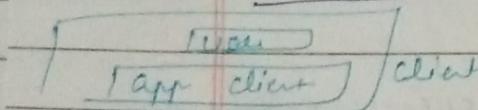
↓
2-tier
arc.



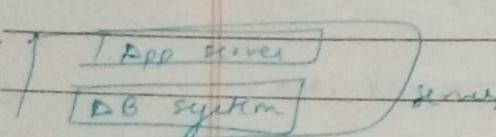
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2-tier → No of clients connected to one server
⇒ Interface → API, ODBC, JDBC

3-tier →



• To reduce load on server, introduce app server layer.



App server comm' with DB system to access data

3 level Architecture ↗ DBMS

user 1

user 2

↓
sub schema

External level

External level

View level

Overall designing

Conceptual schema

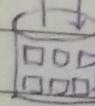
Conceptual level

mapping b/w external & conceptual level

logical level

Physical / Internal schema

Internal / Physical level



Access data from disk

storing data on disk

mapping b/w conceptual & physical level.

to tell how data is to be stored in disk
& what data str you are using.

⇒

Only 1 physical & 1 conceptual but many of external level

Allocating part of conceptual schema to sub schema is advantageous. → multi user can access portion of conceptual as per requirement

Data Independence

Whenever you are modifying certain things which doesn't affect DB architecture.

Physical
Data independence

Logical Data
independence

- Modification at physical level & does not affect higher level (conc. or logical level)
- Modification at conc. or logical level, it doesn't effect phy. or higher level.

Database language

- Data Definition Lang (DDL) → to create DB/ alter the structure or to drop the table
→ specifies DB schema
- Data Manipulation Lang (DML) → Retrieval (Acces)
Updation ↓
Deletion ↓
Insertion
- SQL → supports both DDL & DML
- TCL → Transaction Control Lang → Grant & revoke
- Data Control Lang.
- DDL
- SDL
- VDL → View Definition Lang.
- HL → Host lang.
- 4GL → 4Gen lang.
- TCI

char varchar
fixed size

SQL

Roll	Name	Add	Branch
1	Amit	Lko	IT
2			
3			
4			
5			

⇒ Create a Table

SQL > CREATE Table student

Roll Number(2),

Name varchar(30),

Address varchar(40),

Branch varchar(20),

Mob. Number number(10);

|| Table Created

⇒ To Insert the data

SQL > INSERT INTO

student,

VALUES 1, 'Amit', 'Lko', 'IT', 998...;

|| 1 row updated

⇒ Display all content

SQL > SELECT * FROM student;

SQL > SELECT Roll, Name FROM student;

SQL > SELECT Roll, Name FROM student
WHERE Address = 'Lko';

⇒

Employee (Emp Id, Emp name, Emp designation,
Mob, Salary, address)

Get emp. details who are getting more than
5000 salary

SQL > SELECT * F



- ① SQL > `SELECT * FROM Employee
WHERE Salary > # 50000;`
- ② `SELECT * FROM Employee
WHERE Salary > 50000 AND salary < 100000;
WHERE Salary BETWEEN (50,000, 100000);`
- ③ `SELECT * FROM Employee
WHERE Salary IN (20,000, 50,000, 70,000, 80,000);`
- ④ `SELECT Emp-id, Emp-name FROM employee
WHERE EmpDesignation = "Manager" AND
Address = "LKO".`
- ⑤ // in sorted order of employee name
`SELECT Emp-id, Emp-name FROM employee
* ORDER BY emp-name asc;`
- ⑥ // disp name of emp where name starts with 'C'
`SELECT Emp-name FROM employee
* WHERE emp-name LIKE "C%";
" - uv";`
- ⑦ // Add column age
`ALTER Table Employee
Add Age number(2);`
- ⑧ // by giving 10% hike in salary
`SELECT * FROM employee
WHERE SALARY+= 0.10 * SALARY;`

`UPDATE Table Employee`

`SET Salary = (Salary + Salary * 10) / 100;`

set query to get salary 10% like who are working as manager

⑦ UPDATE Table Employee

where SET Salary = Salary + 0.10 * salary
where Emp designation = "Manager";

⑧ DELETE * FROM Employee

where emp_id = "E10";

⑨ RENAME Employee TO Eng;

Other fn

Functions in SQL

Conversion

Number fn

- Aggregate fn

- character fn

Date func.

Library → Dual

absolute
fn (number fn)

⇒ SQL > SELECT ABS(-200) From Dual;
= 200

⇒ SQL > SELECT CEIL(74.23) From Dual;
= 74

⇒ SQL > SELECT POWER(3,2) From Dual;
= 81

⇒ SQL > SELECT SQRT(81) From Dual;
= 9

⇒ SELECT FLOOR(74.5) From Dual;
= 75

⇒ SQL > $\text{SELECT MOD}(81, 6) \text{ FOR XML Full};$
 $= 3$

⇒ SQL > $\text{SELECT ROUND}(125.5, 0)$ $(125.55, 1)$
 $= 125$ 125.6

⇒ SQL > $\text{SELECT TRUNC}(123.55, 1)$
 Truncate
 $= 123.5$

⇒ SQL > $\text{SELECT SIN}(x), \text{COS}, \text{TAN}$

⇒ SQL > $\text{SELECT SINH}(x)$

⇒ LN(data)

⇒ LOG(b, data)

Employee			
Emp ID	Emp Name	Salary	Address
,		<u>NULL</u>	
20		—	

Aggregate function

⇒ SQL > $\text{SELECT AVG(Salary)} \text{ FROM Employee};$

⇒ SQL > $\text{SELECT COUNT(Salary)} \text{ FROM Employee};$
 $\text{COUNT} * (\text{Salary}) \rightarrow \boxed{\text{All rows}}$.
 $\text{COUNT DISTINCT} (\text{Salary})$

⇒ SQL > $\text{MIN(Salary) / MAX(Salary)}$

~~SELECT * FROM Employee;
WHERE MAX(Salary);~~

SELECT Emp-ID, Emp-Name, Desig., MAX(Salary)

⇒ SQL > SUM (SALARY)

⇒ SQL > STDDEV (SALARY)

⇒ SQL > VARIANCE (SALARY)

Character Function

INITCAP → convert first character to capital, w/
remaining char. lower case.

⇒ SQL > SELECT INITCAP ("RAJIV") FROM Dual
⇒ Rajiv

⇒ SQL > SELECT LENGTH ("RAJIV") FROM Dual
= 5

SUBSTR → to find loc of particular char.

⇒ SELECT SUBSTR ("Delhi_University" 5, 3)
= i-v

⇒ INSTR ("GGSIPU, Delhi", "IPU") FROM Dual
= 4

⇒ GREATEST (85, 95.5, 40.2)
⇒ 95.5

⇒ LEAST (85, 95.5, 40.2)
= 40.2

# Relational Algebra

Set-oriented operation

— Union

— Intersection

— Cartesian Product

— Set-difference

Relation oriented operation

— Selection (σ)— Projection (π) \rightarrow vertical values— Join (\bowtie)

①

R	S	RUS	RTS																								
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$R - S \neq S - R$

 $R \times S \Rightarrow 4$ tuples

Td	Name	S-ID	S-Name

②.

Employee				Employee 2	
E.Id	E.Name	Des.	Salary	E.Id	Qualification

UN
⇒ SQL > SELECT * FROM Employee
WHERE Salary > 40000;

UR
⇒ RA > σ (Employee)
Salary > 40000
:

RA > π (Eid, Employee)
Eid, salary > 40000

Display Emp.Id & Name & Salary where
qualification is 'MBA'

②

①

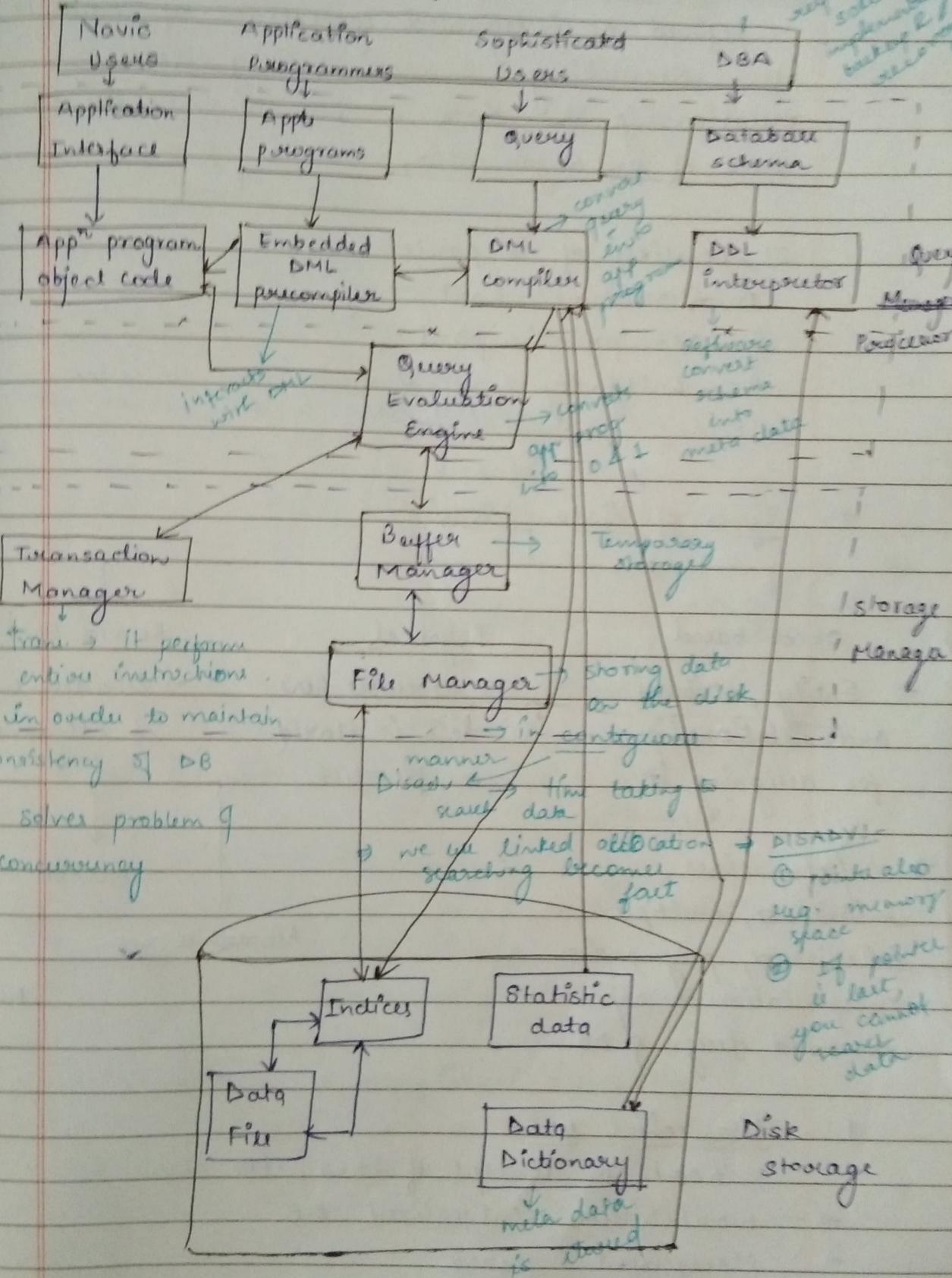
SELECT *
- Emp AS
SELECT E.Id, E.Name, Salary FROM Employee
SELECT E.Id, Qualification
FROM Emp2
WHERE Qualification = "MBA"

RA > σ (Employee, Employee2)
Qualification = "MBA"

SELECT * FROM Employee ∪ Employee2
WHERE Qualification = "MBA"

WHERE E.Id · Employee = E.Id · Employee2

Overall Structure of Database



UN

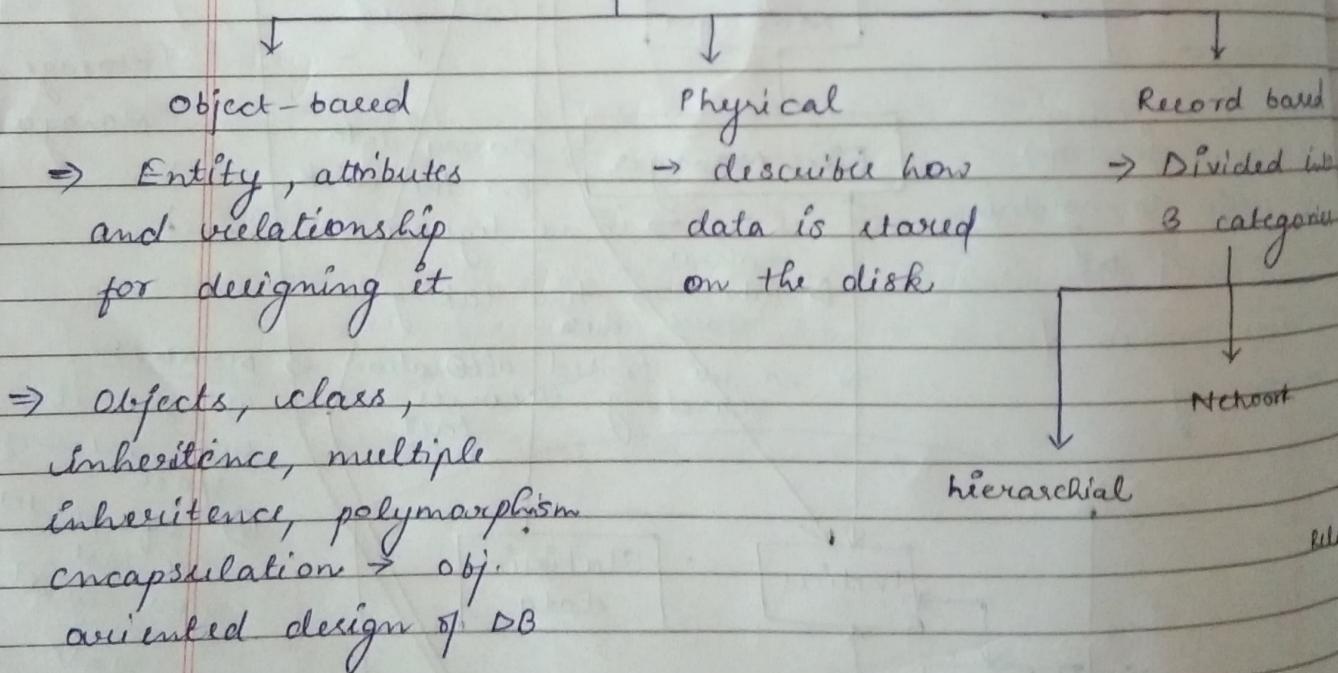
Data Model

Defined as integrated collection of concepts for describing & manipulating data, relationship b/w data & constraints on the data of an organisation.

It consists of 3 parts

- 1 - Structure chart → set of rules that can be used to design db
- 2 - Manipulative part → Retrieval & updation
- 3 - Integrity → Set of rules used to check whether data is accurate or not

Types of Data Model



1. Hierarchical DM

works on concept of tree like structure.

- works like centralized DB
- Management is easy
- If root → overloaded, system will fail.

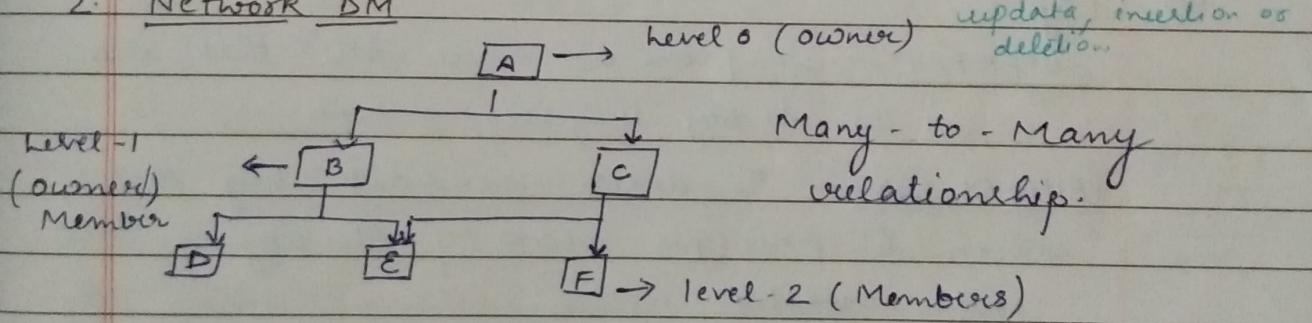
→ for recovery, redundancy is advantageous.

- Advantage
- Simplicity
 - Data sharing
 - Data integrity
 - Data security
 - Efficiency

Disadvantages → It can't support many-to-many relationship.

- Implementation is difficult
- No standard
- Operational anomalies → Inconsistency in data resulting from operation like update, insertion or deletion.

2. Network DM



- Advantage
- Simplicity
 - 1:M, M:1, M:M relationships
 - Data sharing
 - Data independency

Disadvantages → operational anomalies

- Implementation is difficult upto some extent

3. Relational DB

no operational anomalies

RelDBMS → oracle, MySQL

→ Tuple → One complete row

→ Cardinality = No. of rows

→ Degree = No. of attributes/features

Adv :- Can be easily queried.

Disadv → provides inconsistent DB design

Disadv → Hard to cover too high volume
transaction on large amount of data stored,
since it needs to access

→ Key -

Primary key → set of one or more attr. that allow
us to identify uniquely a tuple.

Secondary key → To denote candidate key attr.
chosen as primary means of identifying tuples
within relation.

UNIT - 3. Normalisation

→ process of removing/minimizing the anomalies from a table/relation.

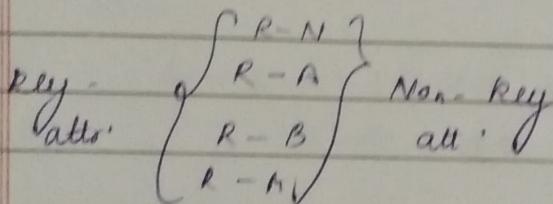
Anomalies :-

- Redundancy - by decomposing relation
- Insertion Anomaly
- Deletion "
- Updation

Functional Dependency :-

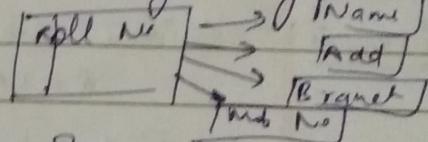
Association of two attributes - determinant & determined.

e.g.: $R \rightarrow N \Rightarrow$ Name is functionally dependent

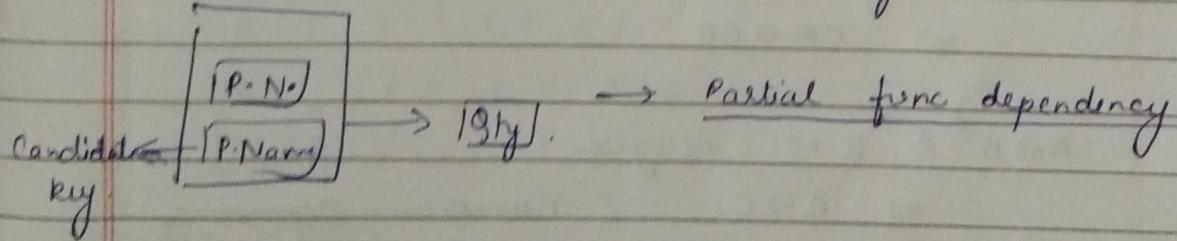


? name

→ When more non-key attr. depend on single key \rightarrow fully func dependency.



→ When non-key attr. depends on more than one key



Rules:

Armstrong axioms

1. Reflexive rule
2. Augmentation rule
3. Transitive dependency :-

1) If Two attr A, B & $B \subseteq A \Rightarrow A \rightarrow B$.

2) If $A \rightarrow B \Rightarrow CA \rightarrow CB$

3) If $A \rightarrow B \wedge B \rightarrow C \Rightarrow A \rightarrow C$

- Union Rule

$$\begin{matrix} A \geq B \\ A \geq C \end{matrix} \quad \left. \begin{matrix} \\ \end{matrix} \right\} \Rightarrow A \rightarrow BC$$

- Decomposition Rule

$$\text{If } A \rightarrow BC \Rightarrow A \rightarrow B \wedge A \rightarrow C$$

- Pseudo-transitivity Rule

$$\begin{matrix} \text{If } A \rightarrow B \wedge BC \rightarrow D \\ \Rightarrow AC \rightarrow D \end{matrix}$$

Example: Given $A \rightarrow B$, $B \rightarrow C$, $BC \rightarrow D$
 $R = (A, B, C, D)$

Find out closure of FDs (F^+)
 $A \rightarrow B$ is given
 $\Rightarrow CA \rightarrow CB$

$A \rightarrow B \wedge B \rightarrow C$ is given
 $\Rightarrow A \rightarrow C$.

(Transitive Dependency)

$A \rightarrow B \wedge BC \rightarrow D$

$\Rightarrow AC \rightarrow D$

(Pseudo Trans. Dep.)

Ques: Let R be relation; $R = (A, B, C, G, H, I)$

and FDs are given $A \rightarrow B$, $A \rightarrow C$, $CG \rightarrow H$,
 $CG \rightarrow I$, $B \rightarrow H$

Find out F^+

$A \rightarrow B$ & $A \rightarrow C$

$\Rightarrow A \rightarrow BC$ (Union)

$CG \rightarrow H$ & $CG \rightarrow I$

$\Rightarrow CG \rightarrow HI$ (Union)

$A \rightarrow B$ & $B \rightarrow H$

$A \rightarrow H$ (Junction).

$A \rightarrow B$ & $CG \rightarrow H$ $AG \rightarrow I$ (Ps)

~~$\Rightarrow A \rightarrow H$~~ $AG \rightarrow I$ (Ps)

Ques $R = (A, B, C, D, E, F)$

FDs: $C \rightarrow F$, $E \rightarrow A$, $EC \rightarrow F$, $A \rightarrow B$, $F \rightarrow D$

Find Key.

E & C are the keys.

$$\{EC\}^+ = \{E, C, F\}$$

$$\{EC\}^+ = \{E, C, F, A\}$$

$$\{EC\}^+ = \{E, C, F, A, B\}$$

$$\{EC\}^+ = \{E, C, F, A, B, D\}$$

Ques $R = (A, B, C, D, E, F, G)$

FDs: $A \rightarrow B$, $BC \rightarrow E$, $BC \rightarrow D$, $AEF \rightarrow G$

$B \rightarrow G$; $B \rightarrow C$

$AEF \rightarrow G$, A, C, F, E, D, G , Find Key (Candidate)

$$A \rightarrow E$$

$$\{BC\}^+ = \{B, C, E, D, G\}$$

$$AC \rightarrow D$$

$$\{AC\}^+ = \{A, C, D, E, G\}$$

$$A \rightarrow G$$

$$\{AEF\}^+ = \{A, E, F, B, G, D, C\}$$