

## UNIT - 1

### Introduction

#### Basic Terminology:-

Data - values or set of values

Information - Meaningful data

Data structure - Data may be organized in many different ways. The logical or mathematical model of a particular orgn of data is called DS.

It is of 2 type -

1. Linear DS - Array, linked list, stack, Queue

2. Non-linear DS - Tree, graph

Linear DS - Processing of data items is possible in linear fashion i.e. data can be processed one by one sequentially

Non-linear DS - A DS in which insertion & deletion is not possible in a linear fashion

#### Elementary Data organization

field → Record → File

(Field - single elementary unit of infn)

Record - collectn of fields

File - collectn of records

It efficiently process certain collection of data.

so, data are also organized into more complex type of structures (array, linked list, graph, tree etc.)

steps for study of such DS -

→ logical / Mathematical description of structure

→ Implementation of the structure on a computer

→ Quantitative analysis of the structure

    (Memory needed to store the structure)  
    Time

Data may be organized in many different ways as studied above. The logical or mathematical model of a particular org'n of data is called array.

### Array:-

- It is list of a finite no. of similar <sup>or homogeneous</sup> data elements  
→ Elements are stored in successive memory locations  
→ Elements are referenced by an index set  
for ex.

Elements of array A =  $a_1, a_2, a_3, \dots, a_n$  : subscript notation  
=  $A(1), A(2), \dots, A(N)$  : Parenthesis notation (FIRER AND, FORTRAN, BASIC)  
=  $A[1], A[2], \dots, A[N]$  : Bracket notation (PASCAL, C, C++)

Here ~~A~~  $\{a\}$ , A - subscripted variable  
    1 - subscript

Linear Array = 1D array

Single dimensional array  
= Each element is referenced by one subscript

Ex. student[5] = {A, B, C, D, E}

Multi Dimensional Array = 2D array

= Each element is referenced by 2 or more subscripts

Ex. → Matrices in Mathematics

→ Tables in Business app'n

→  $A[2][2] = \begin{bmatrix} \cdot & \cdot \\ \cdot & \cdot \end{bmatrix}_{2 \times 2}$

Student

1	A
2	B
3	C
4	D
5	E

student[i][j] = A

Sales

Dept	1	2	3
store	1	2	3
1	25	50	100
2	200	300	400
3	500	600	700

circle [1][1] = 25

## Arrays

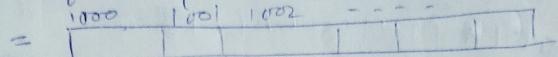
No bound checking concept in C. It's upto programmers to check the UB.

Definition: An array is a consecutive set of memory locations (not always).

### 1. Data Array / Linear Array

Representation of linear array in memory:-

Memory = sequence of addressed locations



### Address calculation

$\text{Loc}(A[K])$  = address of element  $A[K]$  of the array A

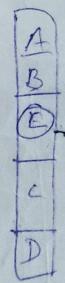
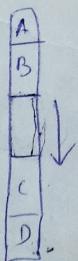
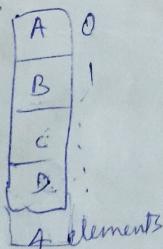
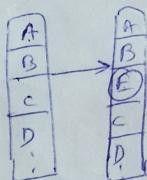
Base (A) = Base address of A  
(<sup>Array name is actually a pointer to the 1<sup>st</sup> location of memory block</sup>)  
(allocated to the name of the array)

$$\text{Loc}(A[K]) = \text{Base}(A) + w(K-LB)$$

w = no of words per memory cell of array

### Analysis regarding Insertion & deletion -

#### Insertion



(Move C,D downward by 1 position)

Create space  
to insert E

5 elements  
(Insert E & increment the array length by 1)

## INSERT (LA, N, K, ITEM)

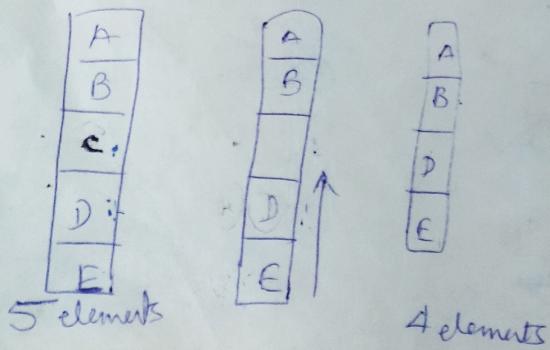
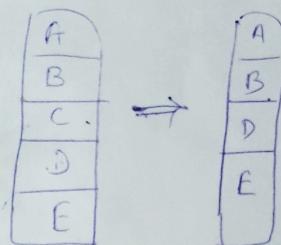
LA - Linear array with N elements

K - Positive integer such that  $K \leq N$

ITEM is inserted into K<sup>th</sup> position in LA

- ① set  $J := N - 1$  // Initialize counter
- ② Repeat ③ & ④ while  $J \geq K$
- ③     set  $LA[J+1] := LA[J]$  // move J<sup>th</sup> element downward
- ④     set  $J := J - 1$  // decrease counter / end of loop
- ⑤ set  $LA[K] := ITEM$  // Insert element
- ⑥ set  $N := N + 1$  // Reset N
- ⑦ Exit

## Deletion



DELETE (LA, N, K, ITEM)

- ① set ITEM := LA[K]
- ② Repeat for J = K to N-1  
and
- ③ set LA[J] := LA[J+1] // (Move J+1<sup>st</sup> element)  
↑  
end of loop
- ④ N := N-1 // Reset the no. of elements in LA
- ⑤ EXIT

### Traversal

TRAVERSE (LA, LB, UBR)

LA - Linear array

LB - Lower bound

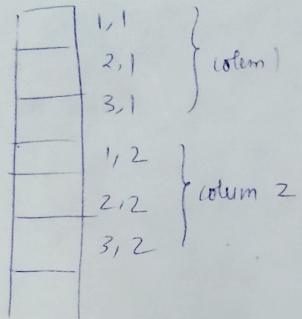
UB - upper bound

- ① Repeat for K = LB to UB  
Apply PROCESS to LA[K]  
(End of loop)
- ② Exit

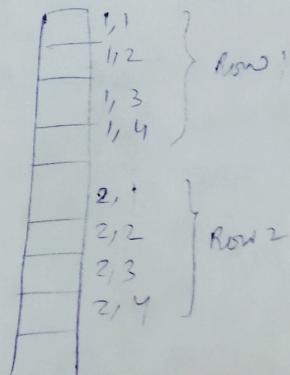
## Multidimensional Array

2D array:

$$\begin{bmatrix} & 1 & 2 & 3 & 4 \\ 1 & - & - & - & - \\ 2 & - & - & - & - \\ 3 & - & - & - & - \end{bmatrix} = (3 \times 4)$$



Column-major order



Row-major order

## Matrix Programs

Traversal

```
for (i=1 to n)
  {
    for (j=1 to m)
      printf("%dA[i][j]);
```

## Matrix Addition

```
for (i=1 to n)
  {
    for (j=1 to m)
      c[i][j] = A[i][j] + B[i][j]
```

$$\begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix} + \begin{bmatrix} 7 & 8 & 9 \\ 10 & 11 & 12 \end{bmatrix} = \begin{bmatrix} 8 & 10 & 12 \\ 14 & 16 & 18 \end{bmatrix}$$

### Matrix Multiplication

$$\begin{matrix} \text{Matrix} & \text{Multiplication} \\ \left[ \begin{matrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{matrix} \right]_{2 \times 3} & \left[ \begin{matrix} 7 & 8 & 9 \\ 10 & 11 & 12 \\ 13 & 14 & 15 \end{matrix} \right]_{3 \times 3} = \left[ \begin{matrix} 7+22+42 \\ 24+20+39 \\ 3+24+45 \end{matrix} \right]_{3 \times 1} \end{matrix}$$

Mat(A, B, C, M, P, N)

A - MXP matrix array

B - P X N matrix array

C - Product of A & B is stored in this M X N matrix

→① Repeat ② to ④ for I=1 to M

→②      Repeat ③ to ④ for J=1 to N

→③      set C[I,J] := 0 // initialization C[1,1]

→④      Repeat for K=1 to P

$$C[I,J] := C[I,J] + A[I,K] * B[K,J]$$

[end of inner loop]

[end of step ② loop]

[end of step ① loop]

→⑤ exit

### INITIALIZATION OF ARRAYS

#### Initializing an 1D array

int A[5] = {1, 2, 3, 4, 5};

char A[5] = {'h', 'e', 'l', 'l', 'o'};

char A[8] = "Hello";

#### Initializing a 2D array

int A[2][2] = {{1, 2}, {3, 4}};

int A[3][3] = {{1, 2, 3}, {4, 5, 6}, {7, 8, 9}};

#### Initializing a 3D array

int A[3][3][3] = {{{1, 2, 3}, {4, 5, 6}, {7, 8, 9}}, {{1, 2, 3}, {4, 5, 6}, {7, 8, 9}}, {{1, 2, 3}, {4, 5, 6}, {7, 8, 9}}};

A[1][1][0] = 16

A[0][0][0] = 0

A[0][0][2] = 0

### Application of arrays -

→ To list some elements of a record.

→ Matrices in Mathematics

→ Tables in Business app'

### SPARSE Matrices & Vectors -

Matrices with a relatively high proportion of zero entries are called SPARSE Matrices.

There are 2 types of n square sparse matrices

#### D Triangular Matrix -

Non-zero entries can only occur (on) or (below)/above the main diagonal.

$$\begin{bmatrix} \times & \times & \\ \times & \times & \times \\ \times & \text{non-zero} & \times \\ \times & \text{zero} & \times \\ \times & \times & \times \end{bmatrix}$$

lower triangular

$$\begin{bmatrix} \times & \times & \times & \times \\ \times & \times & \times & \times \end{bmatrix}$$

non-zero

upper triangular

Ex.

$$\left[ \begin{array}{cccc} 1 & & & \\ 2 & 5 & & \\ 3 & 6 & 8 & \\ 4 & 7 & 9 & 10 \end{array} \right] = \left[ \begin{array}{cccc} 1 & 0 & 0 & 0 \\ 2 & 5 & 0 & 0 \\ 3 & 6 & 8 & 0 \\ 4 & 7 & 9 & 12 \end{array} \right]$$

### ③ Tri-diagonal Matrix

non-zero entries can only occur on the diagonal or on elements immediately above or below the diagonal.

$$\begin{bmatrix} x & & \\ x & x & x \\ & x & x & x \\ & & x & x & x \\ & & & x & x \\ & & & & x & x \end{bmatrix}$$

Ex.  $\begin{bmatrix} 1 & 2 & & & & \\ 3 & 4 & 5 & & & \\ & 6 & 7 & 8 & & \\ & & 9 & 10 & & \end{bmatrix} \equiv \begin{bmatrix} 1 & 2 & 0 & 0 & & \\ 3 & 4 & 5 & 0 & & \\ 0 & 6 & 7 & 8 & & \\ 0 & 0 & 9 & 10 & & \end{bmatrix}$

NOTE :- Natural method of representing matrices in memory as 2D array may not be suitable for sparse matrices i.e. one may save space by storing only non-zero entries. (to avoid wastage of memory space which will caused by storing zeros in memory)

### Array Representation of polynomials

1	1	2	3	4	5
0	1	2	3	4	5

$$1 + x + 2x^2 + 3x^3 + 4x^4 + 5x^5$$

$$\text{Polynomial Addition: } (1 + x + x^3) + (2x + x^2 + 2x^3) = (1 + 2x + x^2 + 3x^3)$$

$$\begin{array}{r} \boxed{1 \ 1 \ 0 \ 1} \\ + \boxed{0 \ 1 \ 1 \ 2} \\ \hline \boxed{1 \ 2 \ 1 \ 3} \end{array}$$

ORDERED LIST - Linear list

It is one of the simplest & most commonly found data object.

Ex -

① days of the week  
(Mon, Tue, Wed, Thu, Fri, Sat, Sun)

② values in card deck

(2, 3, 4, 5, 6, 7, 8, 9, 10, Jack, Queen, King, Ace)

③ floors of a building

(basement, lobby, mezzanine, first, second, third)

④ Years the United States fought in World War II

(1941, 1942, 1943, 1944, 1945)

It may be empty or can be written as  $(a_1, a_2, \dots, a_n)$   
where  $a_i$  are atoms of some set S.

### Difference b/w Array & List:

Array	List
It is a set consisting of a <u>fixed</u> no. of data items	It is a ordered set consisting of a <u>variable</u> no. of data items
Insertion & deletion operations can be performed using pointers	

## Character string in C - (%s) (%oc)

"Length of a string = no of characters of the string + 0"

A string of char's is stored in successive elements of a char array & terminated by the Null char ('0').  
to store the Null char

char msg[5] = "Hello" means

msg[0] = 'H'  
msg[1] = 'e'  
msg[2] = 'l'  
msg[3] = 'l'  
msg[4] = 'o'  
msg[5] = '\0' = NULL = 0

ASCII value of Null is 0

## Library functions for String Manipulation -

strcpy - To copy a string, including the null char

~~String.h~~  
char A[10], B[10];  
gets(A);  
strcpy(B, A);  
printf("A: %s and B: %s\n", A, B);

Strncpy - similar to strcpy, but it allows the no of char's to be copied.

- If the source is shorter than the dest, then dest is padded with null char upto the length specified.

char \*strncpy(char \*dst, const char \*src, size\_t len);

strcat - void to append a str to the end of dst "str"

→ char \*strcat (char \*dst, const char \*src);  
→ strcat (A, B); // where char A[10] = {.....}; char B[10];

strncat -

char \*strncat (char \*dst, const char \*src,  
size\_t N);

strcmp - To compare 2 strings, it returns no > 0

if (strcmp (first, second) == 0)  
puts ("The 2 strings are equal")

strncpy - It compares first N char of each string

int strncpy (const char \*first, const char \*second,  
size\_t N);

strlen - It returns the length of the string

(not counting the null char)

→ size\_t strlen (const char \*str);

→ int length;  
gets (name);  
length = strlen (name);

WAP

strrev - To reverse the string

- strrev (str);

strlwr - convert a string into lower case

strlwr (str);

- It can be used by ASCII values

Difference = 32

(lower case = ASCII value - 32)

## Arrays of strings :-

It is just a 2D array of chars.

char A[3][6] = { "Ryan",  
"Tom",  
"Jackie",  
};

~~no~~  
no of strings

max no of char per string.  
(if no is not mentioned, compiler  
will take the longest length of the  
name as the max length)

WAP to check a string as a palindrome - Page 4/25  
of Balaji

### Array as Parameter:-

display (int \*, int);  
main()

\* Pts  
2 parameters

```
{
    int num [] = {1, 2, 3, 4, 5, 6};
    display (& num[0], 6);
}

display (int *x, int n)
{
    int i;
    for (i=0; i <= (n-1); i++)
    {
        printf ("\n element=%d", *x);
        x++;
    }
}
```

void display (int +)
main()
{
 int n [] = {1, 2, 3};
 display (&n[0]);
}

single parameter

void display (int \*a)
{
 for (int i = 0; i <= 2; i++)
 {
 printf ("%d", \*a);
 a++;
 }
}

void length (char +);
main()

length of string

char a[10] = "computer";
length (&a[0]);

void length (char \*a)
{
 for (int c=0; \*a != '\0'; c++)
 a++;
}