

2) ACTIVE attacks

It attempts to alter system resources / info

(iv) Denial

It p

facil

dis

## Active attacks

1:- **Masquerade**: someone other

(i) masquerade

→ when one entity pretends to be another entity

eg



2:- **Modification of message**:-

modify, delay

3:- **Reply** :- save the message and send multiple times

4:- **Denial of service**:-

(ii) modification of messages + some portion of the message is altered or the message is delayed or reordered to produce

To much requests on server

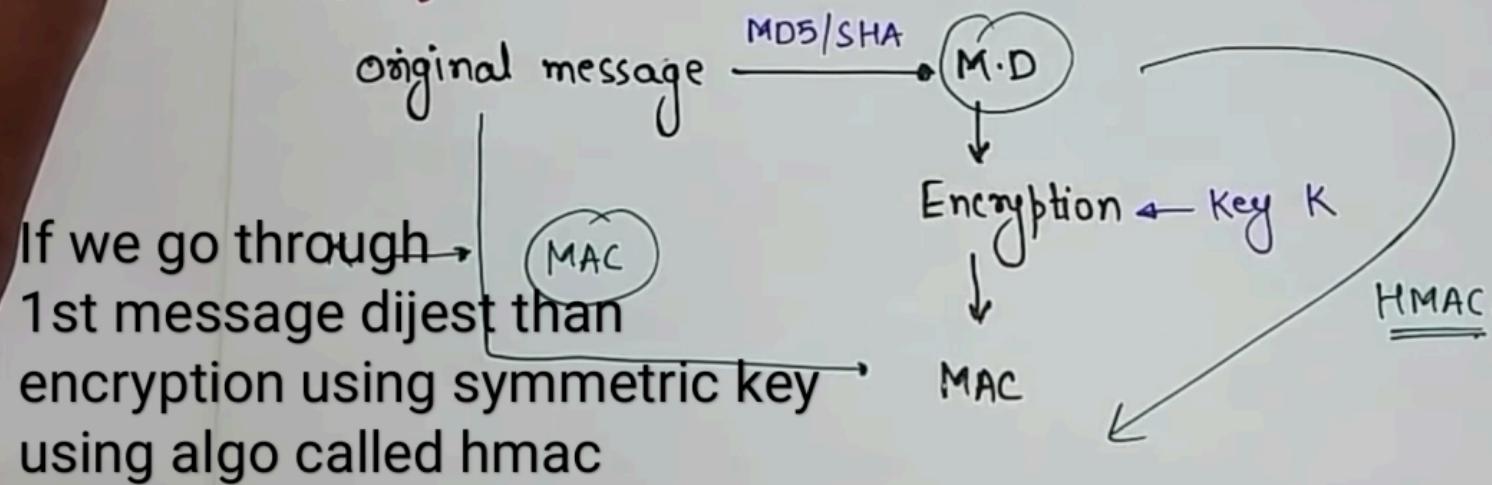




## HMAC [Hash-Based Message Authentication Code]

used for Security implementation in Internet Protocol(IP) and also in SSL(Protocol). → Secure - Socket Layer.

## Concept of HMAC:





MD-5-1

**HASH FUNCTIONS:** [Compression func<sup>n</sup>]

It is a mathematical func<sup>n</sup> - that Converts a numerical i/p Value into another Compressed numerical Value.

↳ o/p is always of fixed Length.

**Features:**

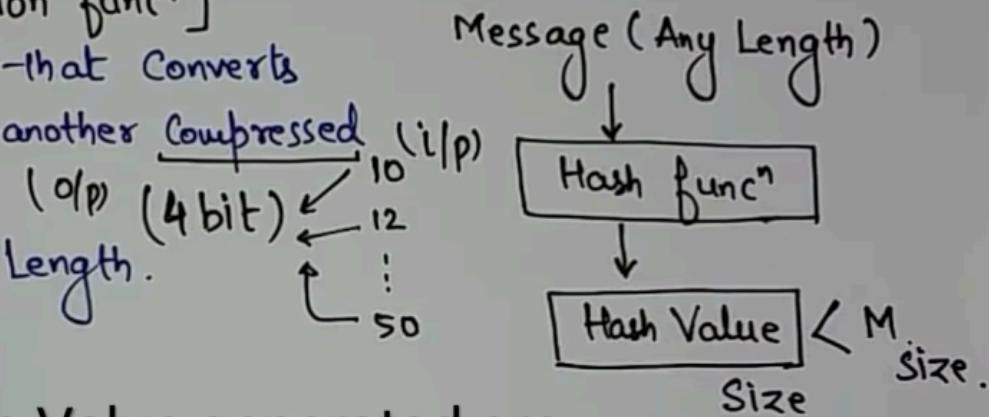
Message Value generated are called message digest

no matter how much length of input if hash function of 4 byte then answer is of 4 byte

**Properties:** (i)

(ii)

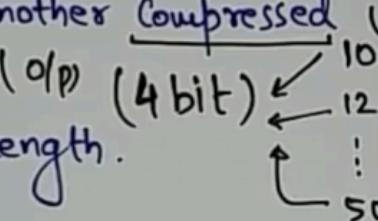
(iii)



**HASH FUNCTIONS:** [Compression func<sup>n</sup>]

It is a mathematical func<sup>n</sup> - that Converts a numerical i/p Value into another Compressed numerical Value.

↳ O/p is always of fixed Length.



Message (Any Length)

Hash func<sup>n</sup>

Hash Value < M

Size

Reverse hash must be impossible

2 ka hash same na ho

**Features:**

(i) fixed Length o/p

(ii) Compression func<sup>n</sup>

(iii) Digest (Smaller rep<sup>n</sup> of larger data) same value ka hash same aya

**Properties:** (i)  $M \rightarrow H$  (Easy)  $H \rightarrow M$  (Very Hard)

(ii)  $M \rightarrow H$  } Same hash Value for Same

$n \dashrightarrow H$  } message everytime.

(iii)  $M_1 \rightarrow H_1$  }  $H_1 = H_2$  <sup>should</sup> not happen.  
 $M_2 \rightarrow H_2$



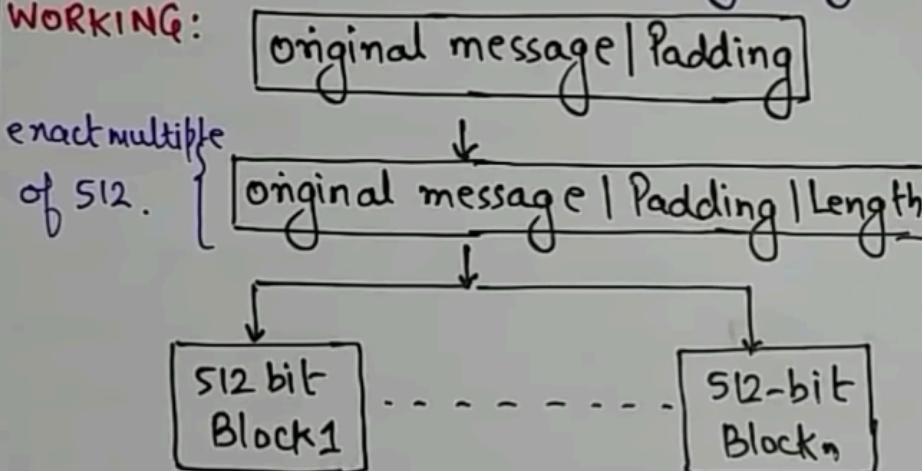
MD5-2

## MD5 [MESSAGE DIGEST]

↳ developed by Ron Rivest.

↳ fast and produces 128-bit message digests.

WORKING:



(iv) Initialize 4-Chaining Variables.  
(32-bit, A, B, C and D)

(i) Padding is done such that total length is 64 bit less than exact multiple of 512.

$$1000 \text{ bits} + 472 = \underline{\underline{1472}}$$

$$512 \times 2 = 1024 \quad \text{64 bit less}$$

$$\begin{array}{r} 512 \times 3 = 1536 \\ \hline 64 \\ \hline 1472 \end{array} \quad \text{than exact multiple of 512.}$$

(ii) Append Original length bef. padding (modulo 64)  
1000 ] length mod  $2^{64}$ .

(iii) Divide it in 512-bit blocks.

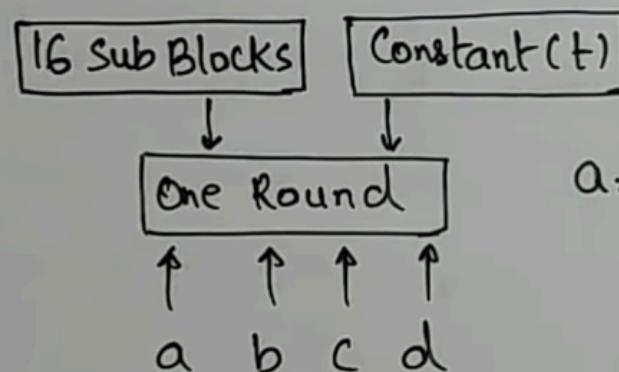
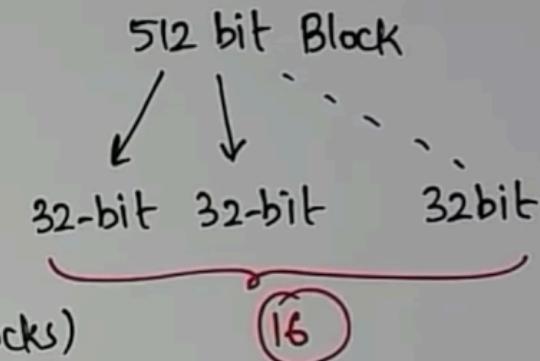


## (V) Process Blocks

↳ Copy four chaining Variables  
into Corresponding Variables.  
 $\{ A=a, B=b, C=c, D=d \}$

↳ divide 512-bit block into 16 (32bit blocks)

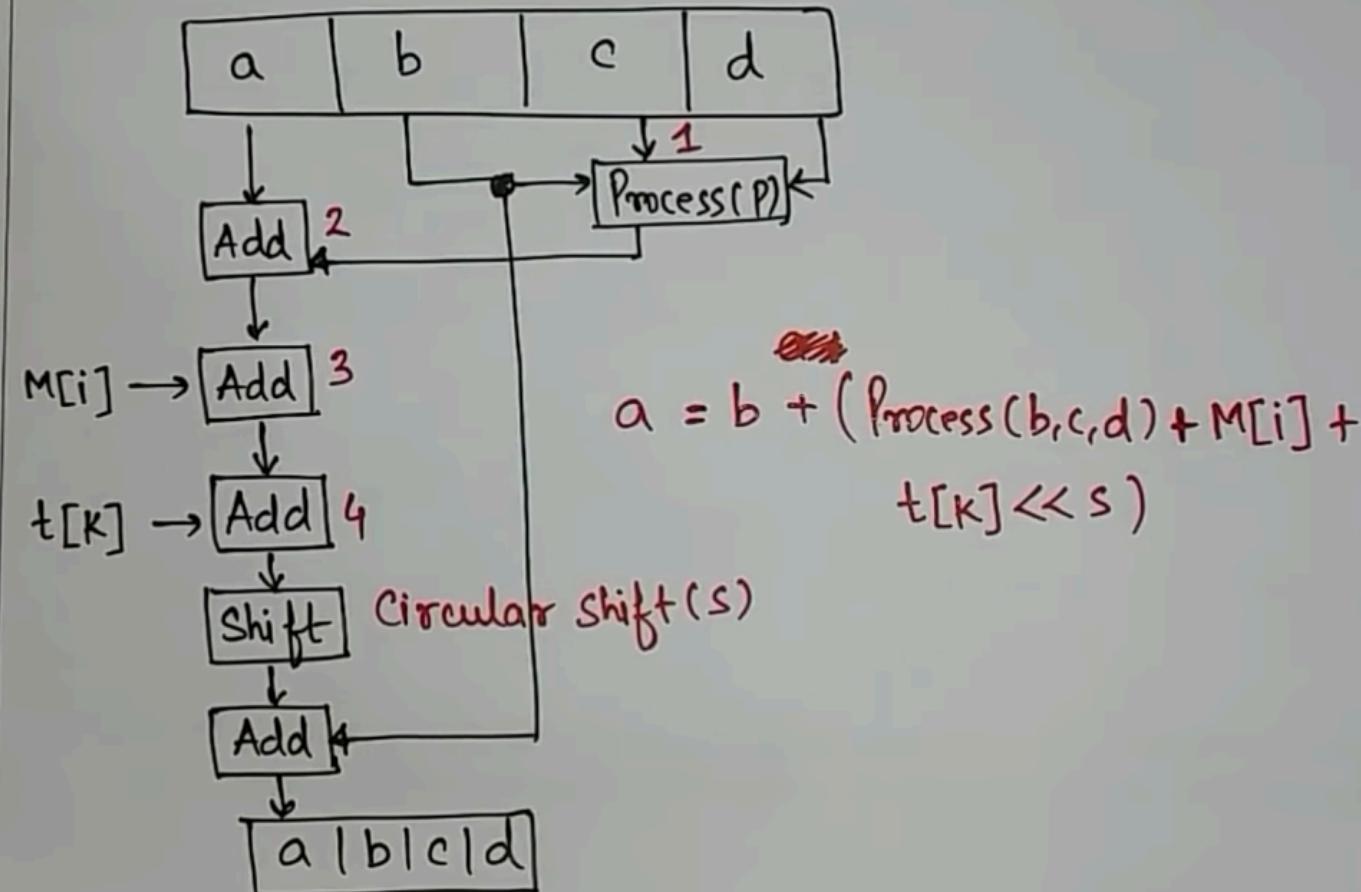
↳ Four Rounds.



$$a = b + ((a + \text{Process}_P(b, c, d) + M[i] + T[k]))$$

<< Shift .

## MD5 Operation:-



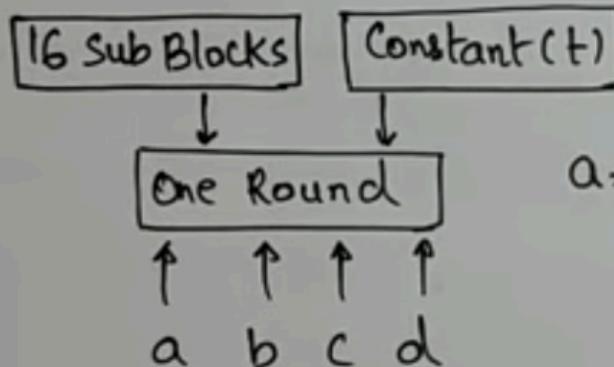
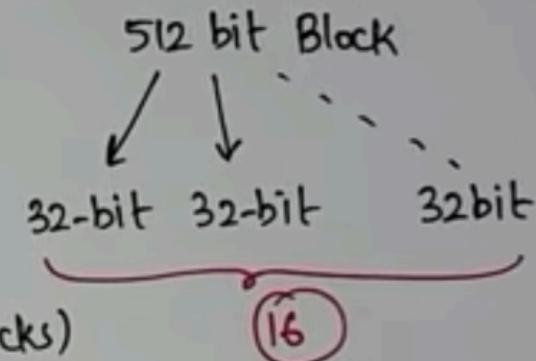


## (v) Process Blocks

↳ Copy four chaining Variables  
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↳ divide 512-bit block into 16 (32bit blocks)

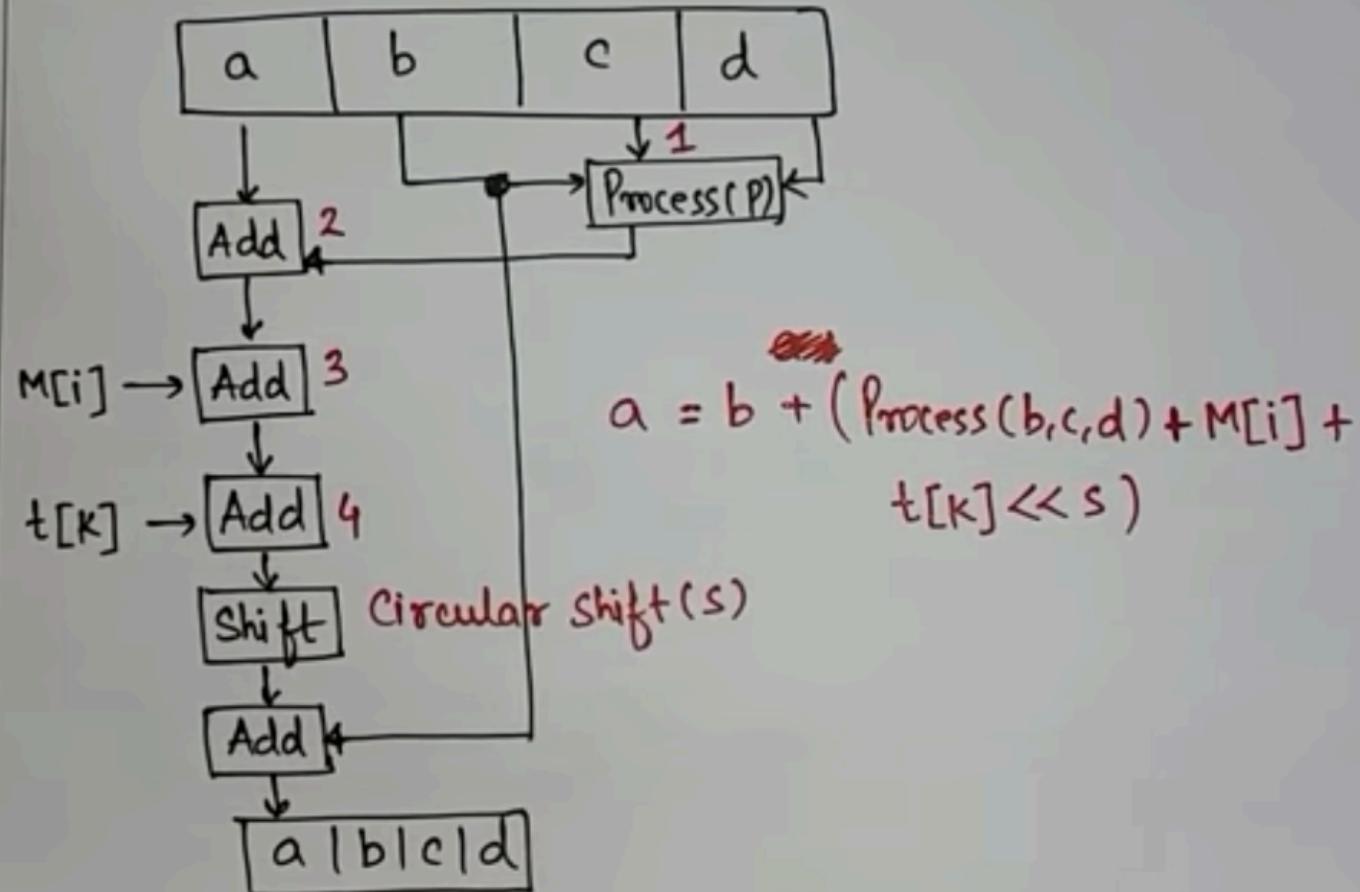
↳ Four Rounds.



$$a = b + ((a + \text{Process}, P(b, c, d) + M[i] + T[k]))$$

$\ll \text{Shift}$ .

## MD5 Operation:-



## → HMAC ALGORITHM

Authentication:

Hashing →  $\times$  keys  
mac → Keys

Hash + mac

Hashing with Keys

1. compute S-bits →

Key

K

2. S || M

Block size of  
plaintext → b bits

3. Hash function

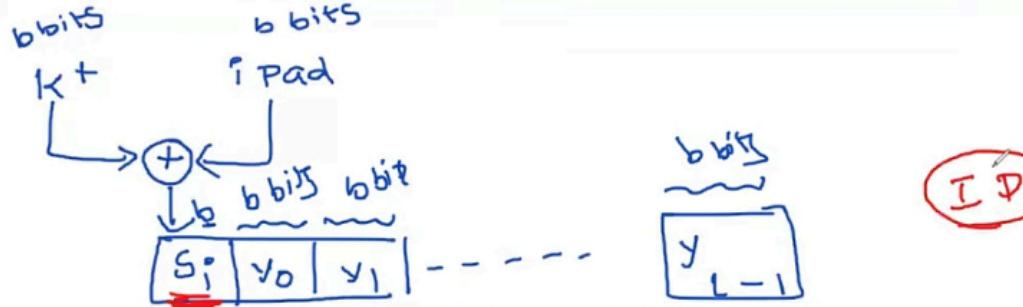
K → K' (b bits)  
100

70 → 30

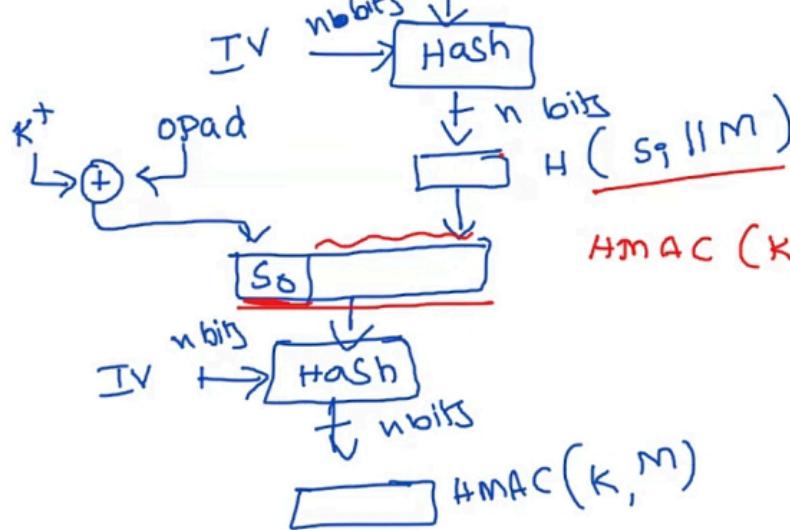
↑  
ipad - 36 hexadecimal

↑  
opad - 5C hexadecimal

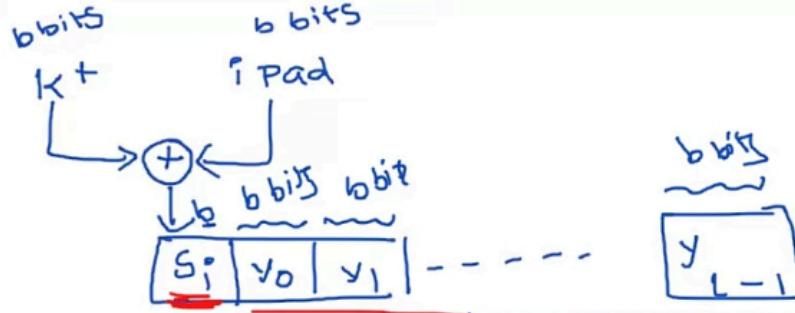
0011 00110 }  
01011100 }



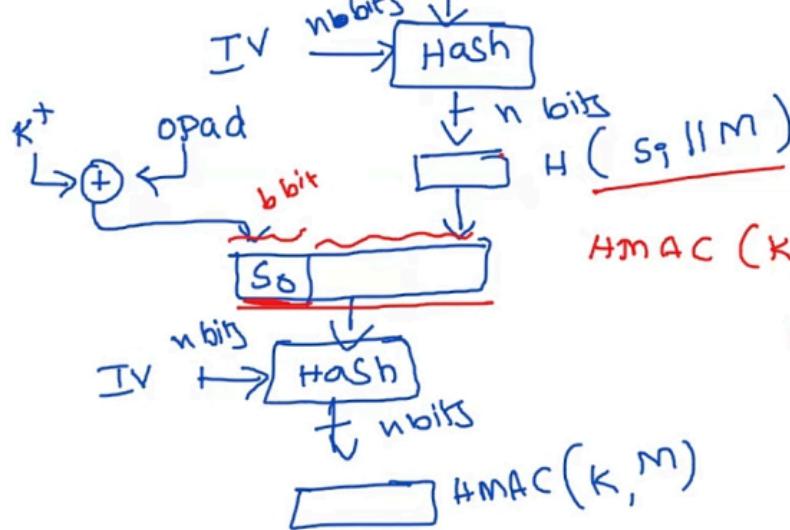
$M$  - PT  
 $y_i$  =  $i$  th block of  $M$   
 $L$  - no. of blocks in  $M$   
 $n$  - length of the hashcode  
 $H$  - Hash function  
 $b$  - no. of bits in a block  
 $K$  - Key



$$\text{HMAC}(K, M) = H \left[ \underline{(K^+ \oplus \text{opad})} \parallel H \left[ \underline{(K^+ \oplus \text{iPad}) \parallel M} \right] \right]$$



IP  
SSL



$$HMAC(K, M) = H[(K^+ \oplus opad) || H((K^+ \oplus i\text{Pad}) || M)]$$

- ① Append 0's to leftend of  $K \rightarrow K^+$
- 2)  $\oplus R$   $K^+$  with iPad  $\rightarrow S^i$
- 3)  $S^i || M$
- 4) Apply  $H$  step 3
- 5)  $\oplus R$   $K^+$  with opad  $\rightarrow S_0$
- 6)  $S_0 ||$

$M = PT$   
 $y_i = i^{\text{th}}$  block of  $M$   
 $L = \text{no. of blocks in } M$   
 $n = \text{length of the hashcode}$   
 $H = \text{Hash function}$   
 $b = \text{no. of bits in a block}$   
 $K = \text{key}$



## ElGamal Cryptography: Asymmetric Key.

### iii) Encryption:

i) Select Random Integer ( $R$ )  $\Rightarrow 4$

ii)  $C_1 = E_1^R \bmod P$ ,  $C_1 = 2^4 \bmod 11 = 5$

iii)  $C_2 = (P \cdot T \times E_2^R) \bmod P = (7 \times 8^4)$

iv)  $C \cdot T = (C_1, C_2) \bmod 11$   
 $\Rightarrow 28672 \bmod 11$

### i) Key Generation:

i) Select Large Prime no. ( $P$ )  $\Rightarrow P = 11$

ii) Select decryption key / Private key ( $D$ )  $= 3$

iii) Select Second part of encryption key or public key ( $E_1$ )  $= 2$

iv) Third part of the encryption key or public key ( $E_2$ ).  $E_2 = \underline{E_1^D \bmod P} = 8$

v) Public Key  $= (E_1, E_2, P)$ , Private key  $= D$   
 $\Rightarrow (2, 8, 11)$ ,  $\hookrightarrow 3$

### iii) Decryption:

$$P \cdot T = [C_2 \times (C_1^D)^{-1}] \bmod P$$

$$P = 11, D = 3, E_1 = 2 \quad \boxed{P \cdot T = 7}$$

$$E_2 = (2)^3 \bmod 11 = 8 \bmod 11 = \\ C_1 = 5$$

# ElGamal Cryptography: Asymmetric Key.

## i) Key Generation:

i) Select Large Prime no. ( $P$ )  $\Rightarrow P = 11$

ii) Select decryption key / Private Key ( $D$ )  $\Rightarrow D = 3$

iii) Select Second part of encryption key or public key ( $E1$ )  $\Rightarrow E1 = 2$

iv) Third part of the encryption key or public key ( $E2$ ).  $E2 = \underline{E1^D \bmod P} \Rightarrow E2 = 2^3 \bmod 11 = 8$

v) Public Key =  $(E1, E2, P)$ , Private Key =  $D$   $E2 = (2)^3 \bmod 11 = 8 \bmod 11 = 8$

$$(125)^{-1} \bmod 11$$

$$(125 \times x) \bmod 11 = 1$$

$$x = 3,$$

$$\Rightarrow (2, 8, 11), \hookrightarrow 3$$

## ii) Encryption:

i) Select Random Integer ( $R$ )  $\Rightarrow 4$

$$ii) C1 = E1^R \bmod P, C1 = 2^4 \bmod 11 = 5$$

$$iii) C2 = (P \cdot T \times E2^R) \bmod P = (7 \times 8^4) \bmod 11$$

$$iv) C.T = (C1, C2) \bmod 11 \Rightarrow 28672 \bmod 11 = 6$$

$$PT = [C2 \times (C1^D)^{-1}] \bmod P$$

$$P = 11, D = 3, E1 = 2 \quad \boxed{P.T = 7}$$

$$E2 = (2)^3 \bmod 11 = 8 \bmod 11 = 8$$

$$C1 = 5, C2 = 6$$

$$\textcircled{7} \longrightarrow \boxed{C.T = (5, 6)} \quad \left\{ \begin{array}{l} P.T = (6 \times (5^3)^{-1}) \bmod 11 \\ \Rightarrow (5^3)^{-1} \bmod 11 \end{array} \right.$$

Encryption.

Decryption

$$\Rightarrow (5^3)^{-1} \bmod 11 \rightarrow 3$$

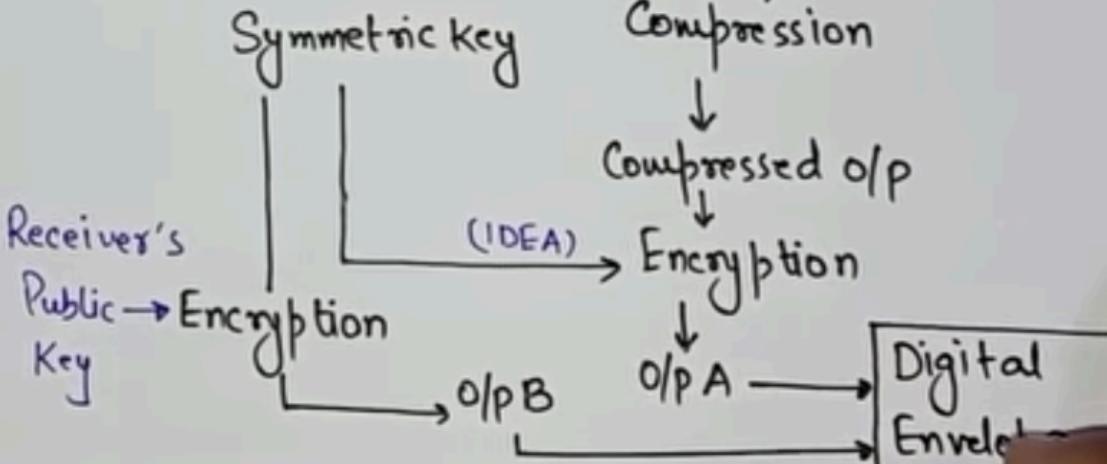
$$(6 \times 3) \bmod 11 \Rightarrow 18 \bmod 11$$

## PRETTY GOOD PRIVACY (PGP) :-

Father of PGP = PHIL Zimmerman.

It is an encryption Program that provides cryptographic privacy and authentication for data communication.  
 ↳ Increases Security of e-mail comm'.

### PGP WORKING:



PGP-1

## PRETTY GOOD PRIVACY(PGP):-

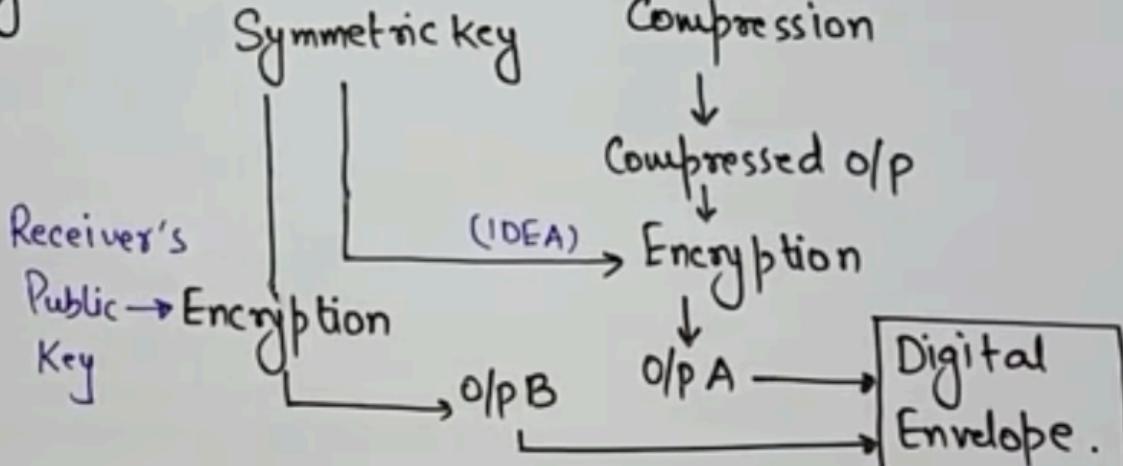
Father of PGP = PHIL Zimmerman.

It is an encryption Program that provides cryptographic privacy, and authentication for data communication.

→ Increases Security of e-mail communication.

**PGP WORKING:** using Public and Private Key

Cryptography .



# IP Security

## Architecture

Authentication header

AH

Authentication Algorithm

Encapsulating Security Payload

ESP

Encryption Algorithm

DOI

Key Management

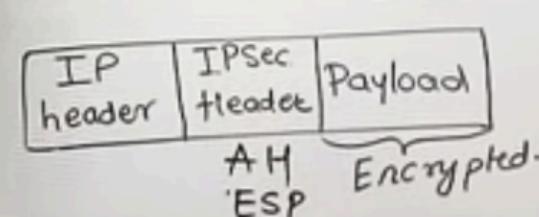
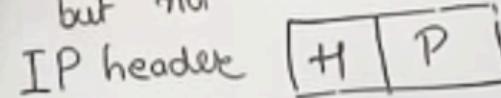
- ipseckey
- Internet Key exchange (IKE) protocol.

## Security Association

- 1) Security Parameter Index
- 2) Security Protocol Identifier (AH/ESP)
- 3) Sequence number Counter (0 to  $2^{32}-1$ )
- 4) AH Information.
- 5) ESP Information.
- 6) Life time of SA
- 7) IPsec Protocol mode
  - Transport Mode
  - Tunnel Mode

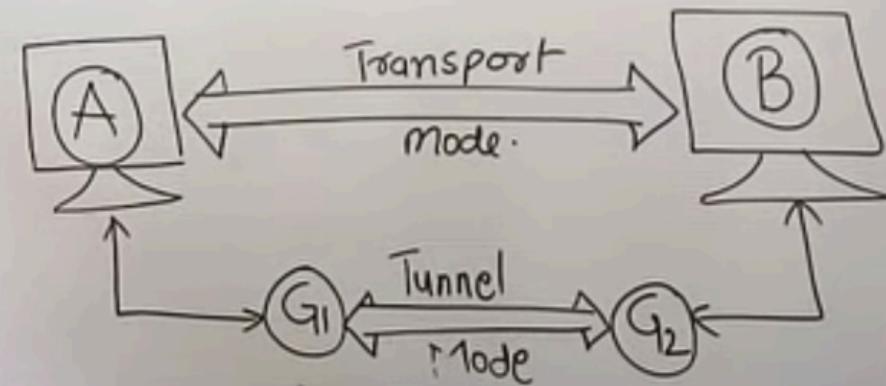
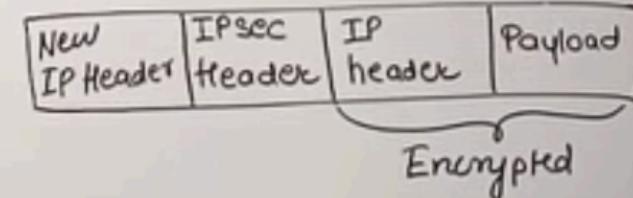
## Transport Mode

- Payload is encrypted but not IP header



## Tunnel Mode

- Payload as well as IP header are encrypted.

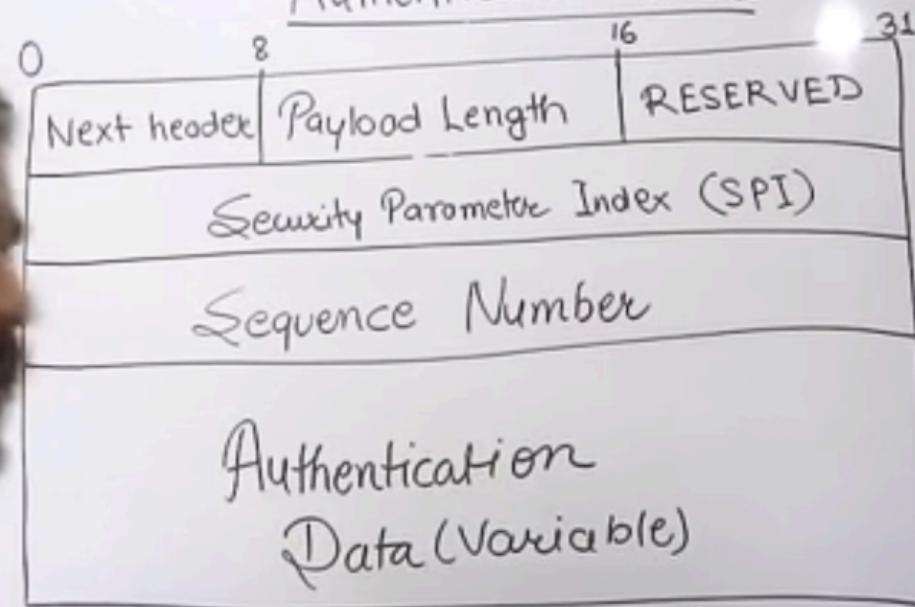


Companies which live from India for these  
enables :-

- ✓ Apera Group
- ✓ Tower research Capital
- ✓ Gravitor Research
- ✓ NK Securities
- ✓ World Quant
- ✓ Quad eye .

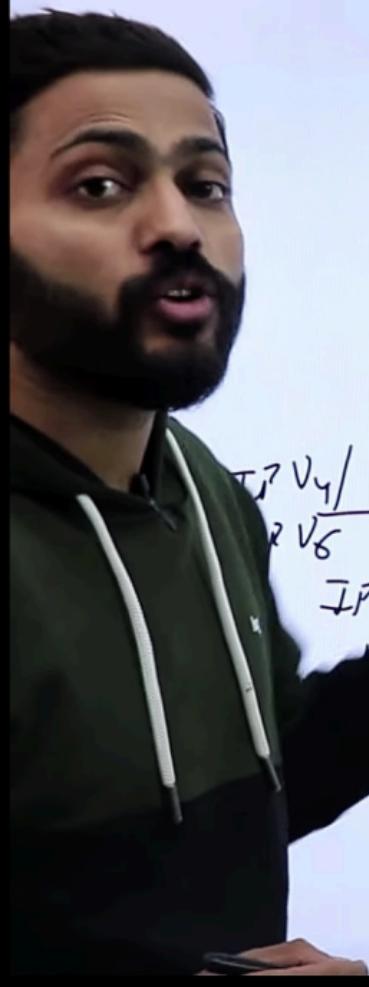


## Authentication header



# IP security (IPSec)

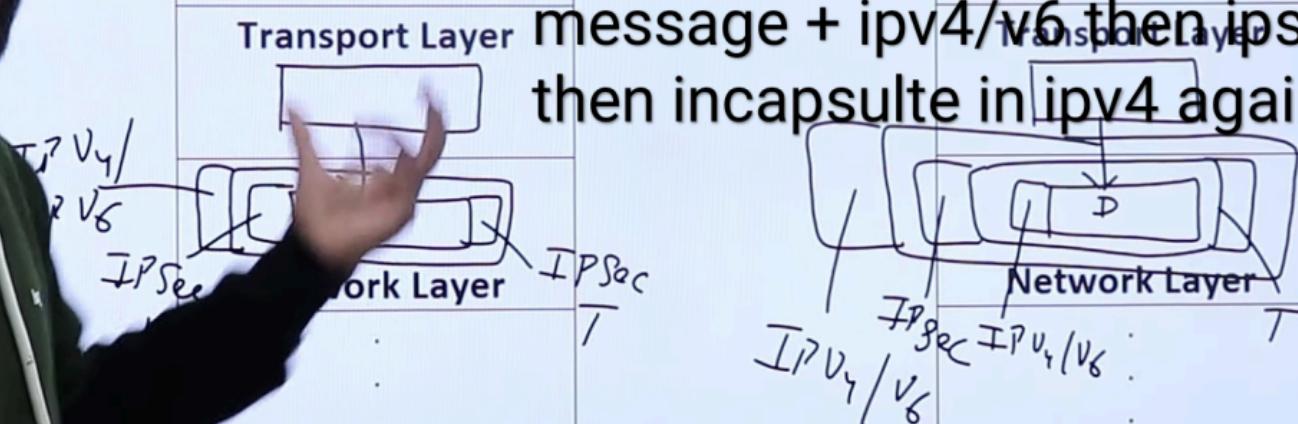
- IETF Standard
- Network layer protocol(Used both by IPv4 & IPv6)
- Uses of IP Security
  - 1. Confidentiality
  - 2. Authentication/Integrity
  - 3. Replay Attack Protection
- Collection of Protocols
  - ✓ 1. Encapsulating Security Payload (ESP)
  - 2. Authentication Header (AH)
  - 3. Internet Key Exchange (IKE)
- Two Modes of Operation(Transport Mode & Tunnel Mode)



## Transport Mode vs Tunnel Mode



Transport mode :- end to end message + ip sec then + ipv4/v6  
Tunnel mode :- Route to router message + ipv4/v6 then ipsec then incapsulate in ipv4 again

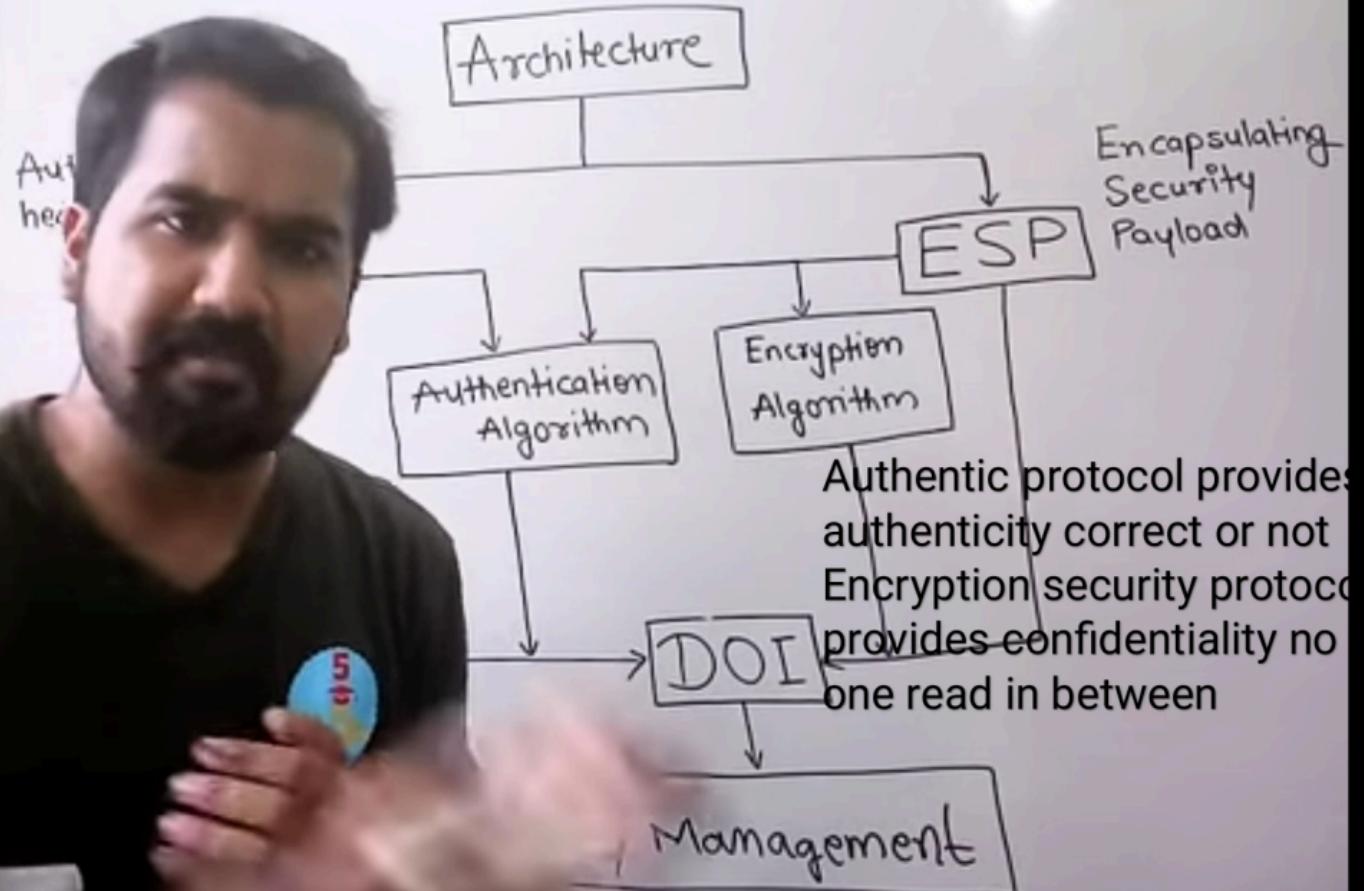


SUBSCRIBED

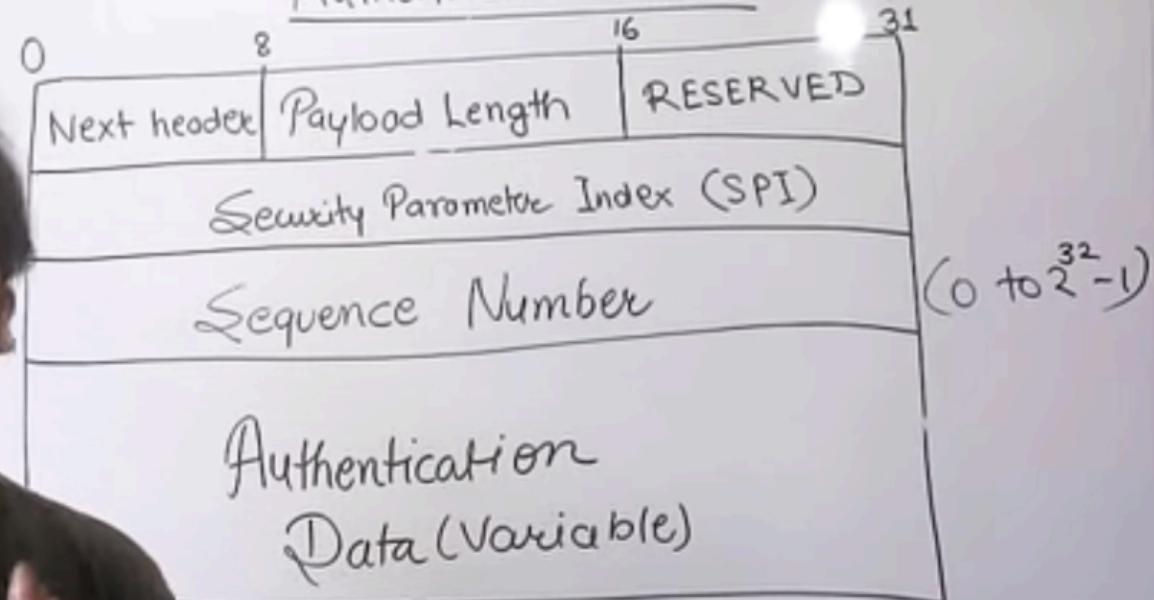


SUBSCRIBE

# IPSecurity



## Authentication header

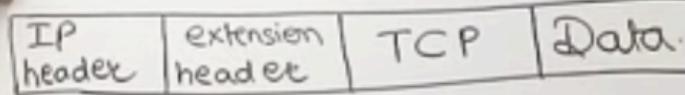


in Authentication header we talk about integrity, authentication in authentication data  
Integrity check value is used

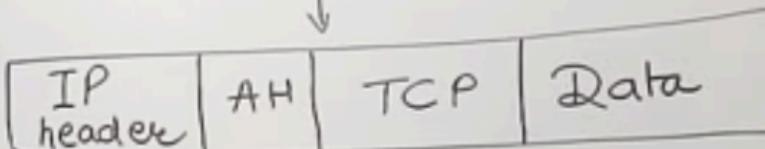


Before Applying  
AH

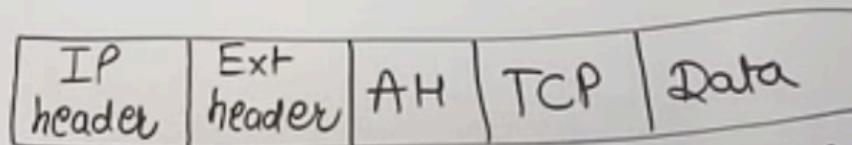
IPV6



IPV4

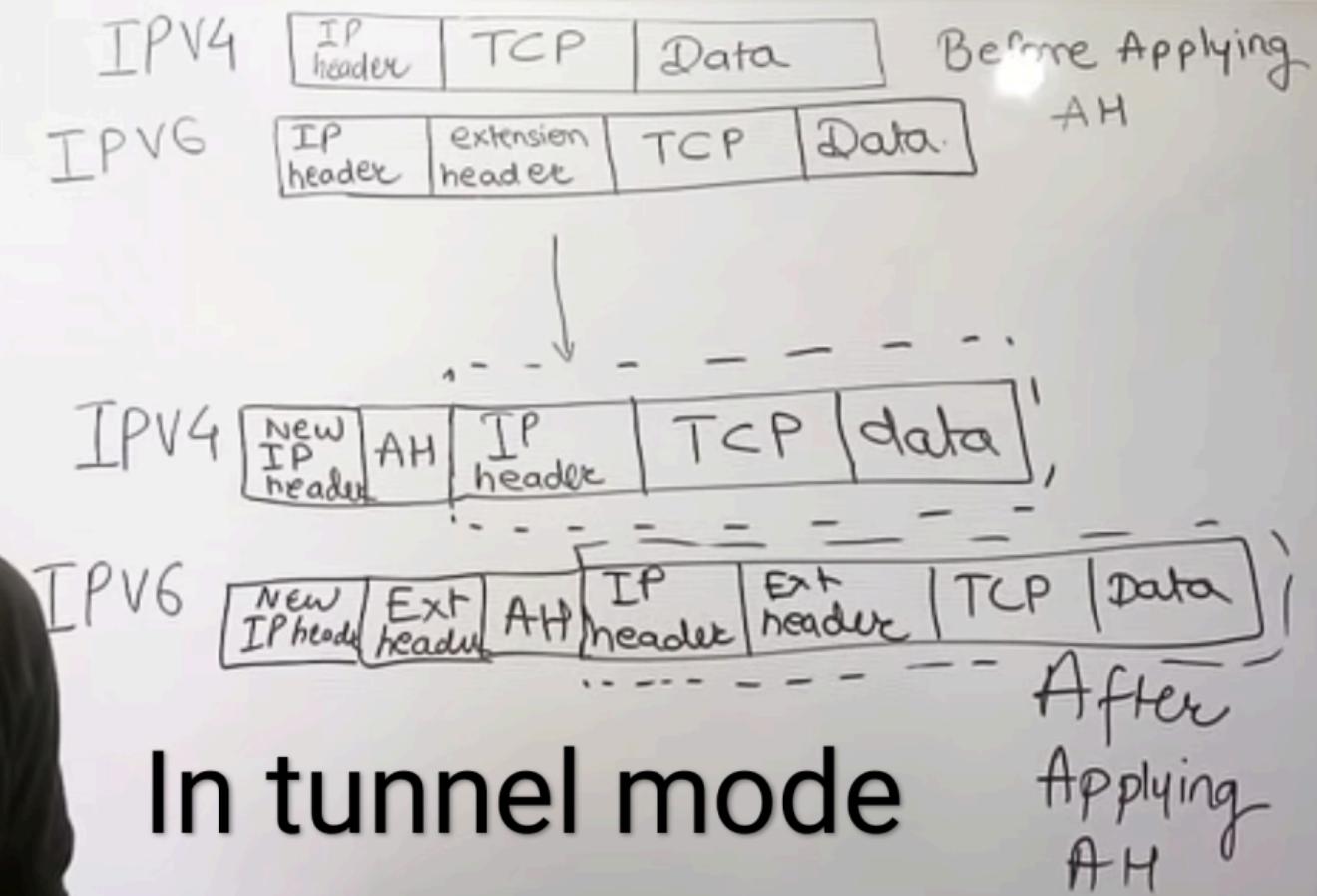


IPV6



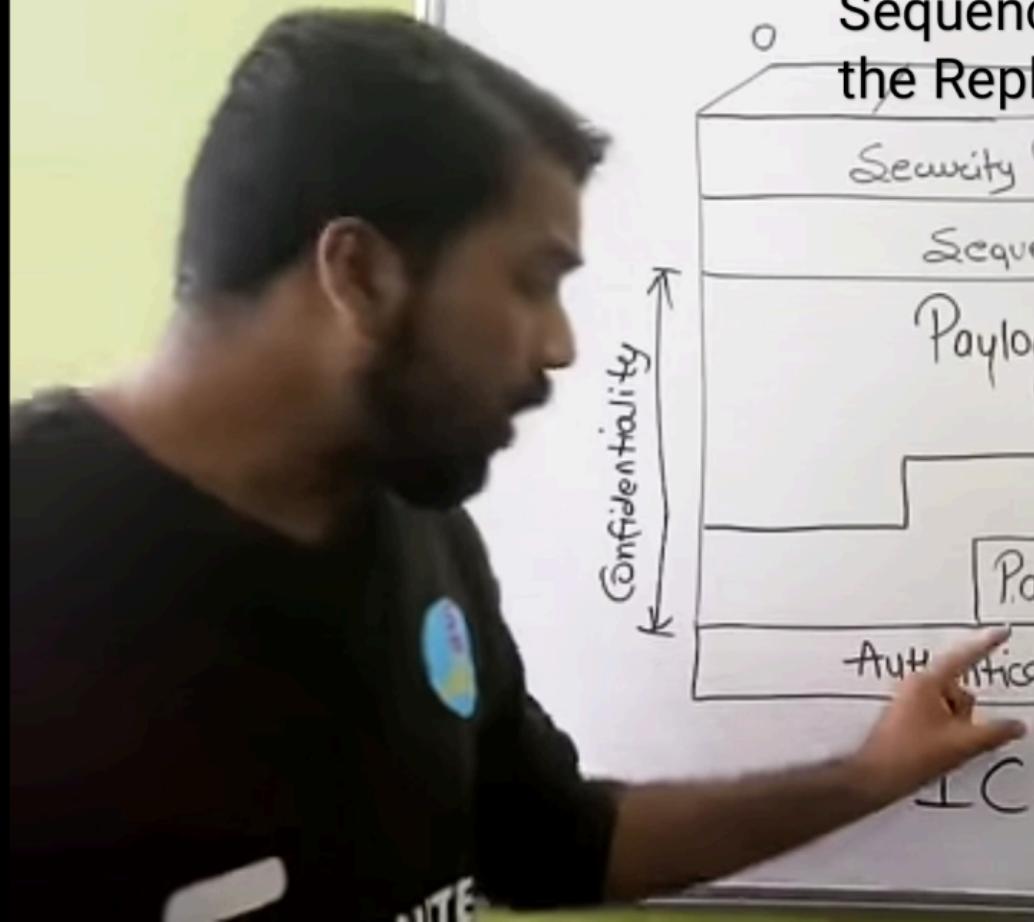
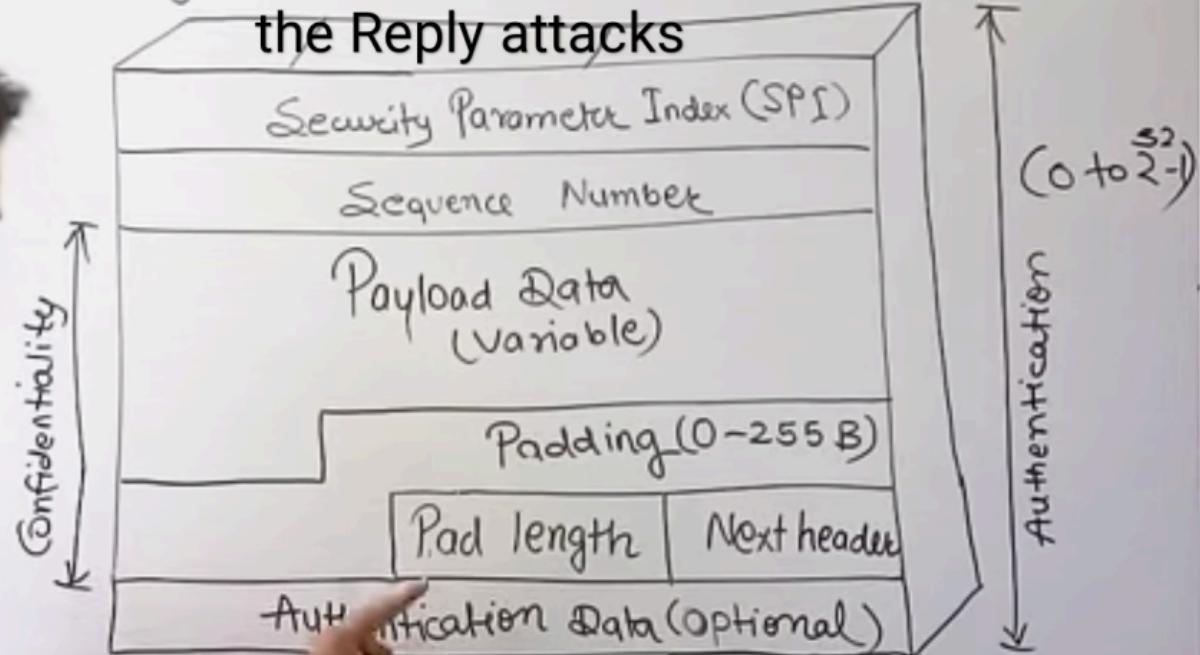
After  
Applying  
AH

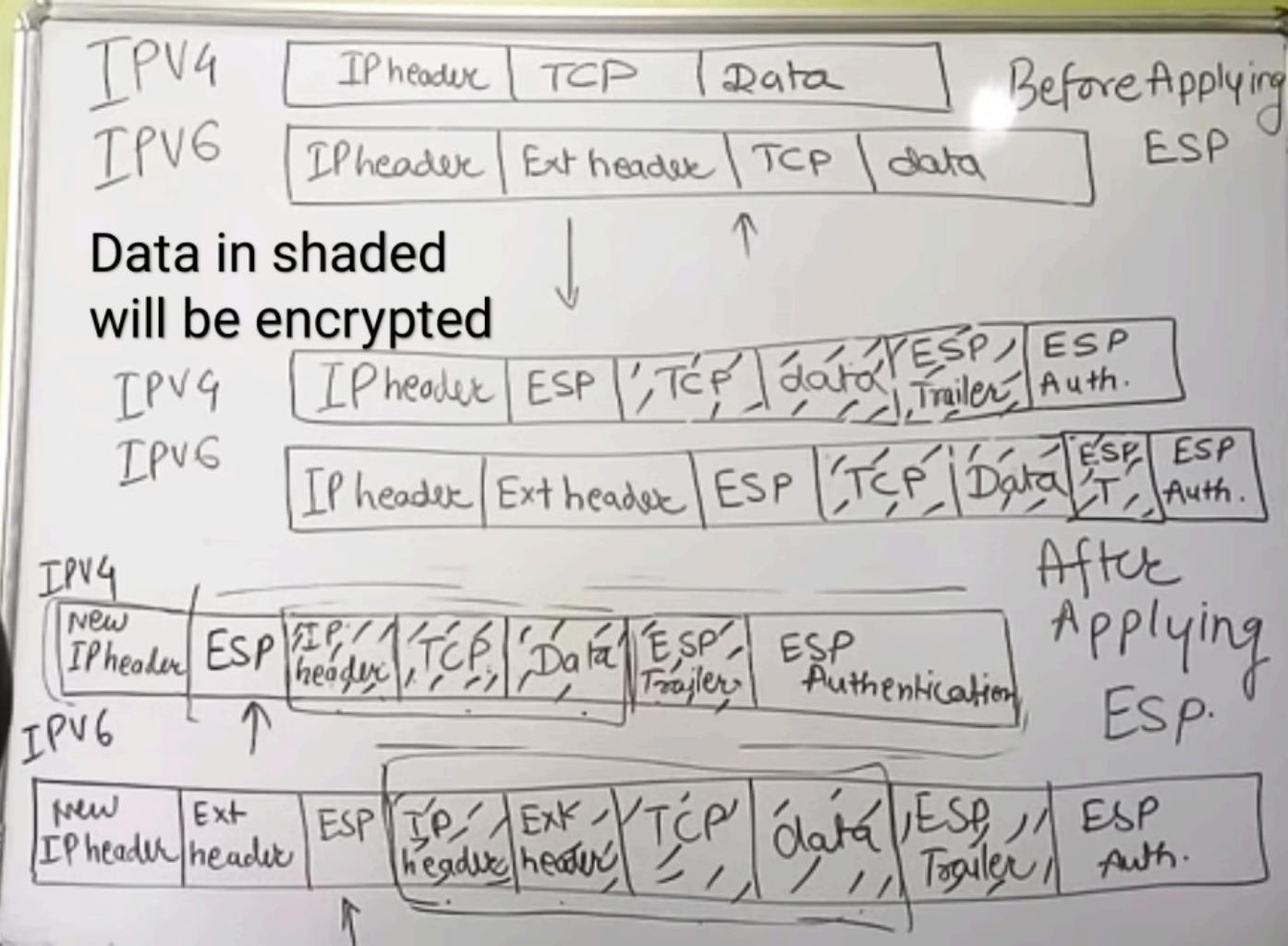
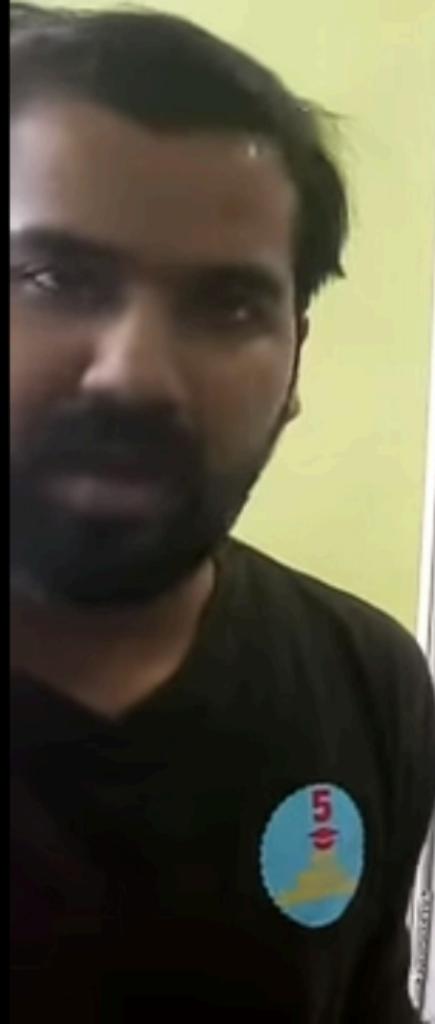
AH in Transport mode



# Encapsulating Security Payload

Sequence number check  
the Reply attacks





## Secure Socket Layer (SSL)

Pending → Current

SSL Handshake Protocol	SSL Change Cipher spec protocol	SSL Alert Protocol	HTTP
------------------------	---------------------------------	--------------------	------

SSL Record Protocol

TCP

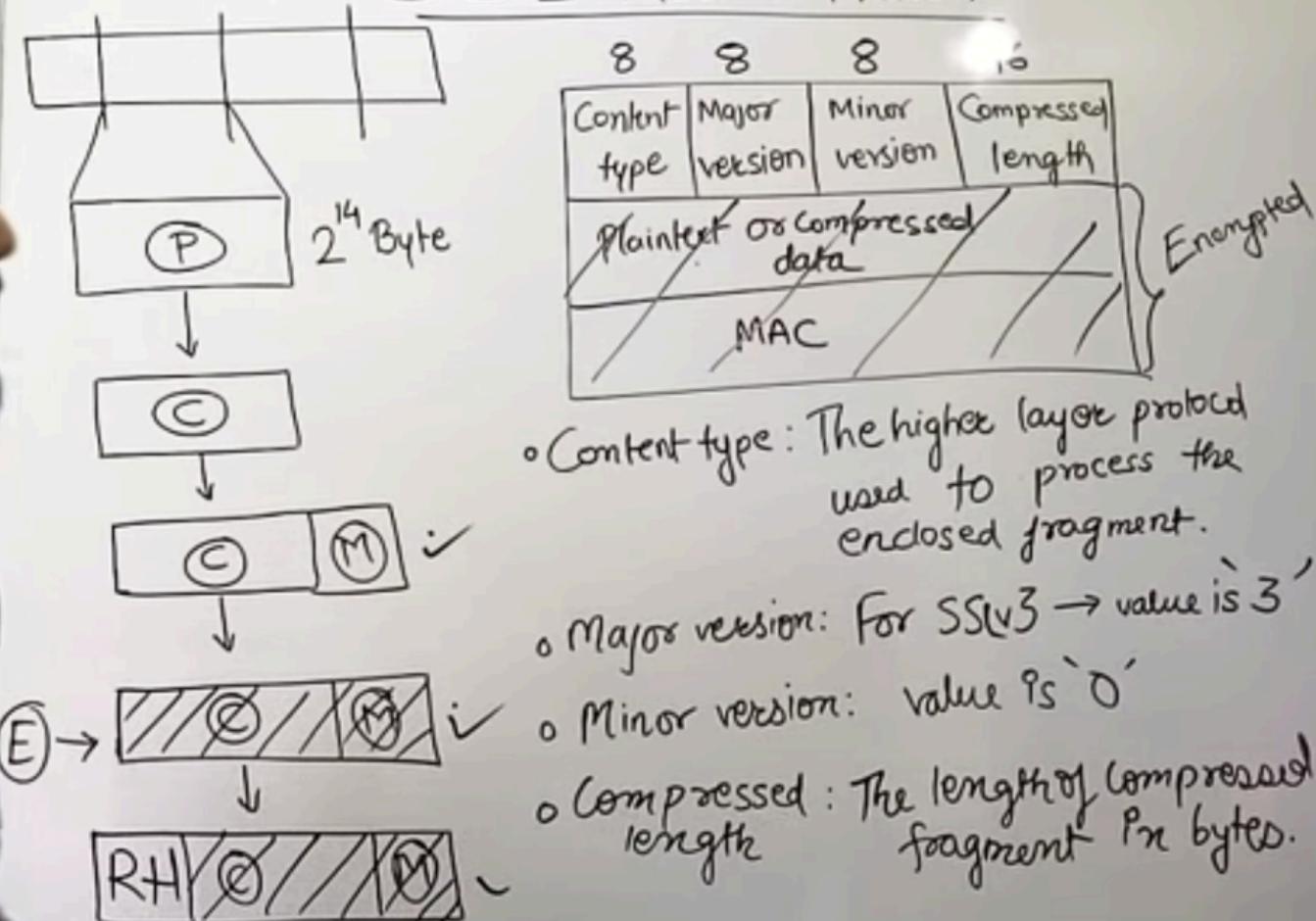
Secure socket Layer is present between Application, transport layer  
Main protocols Handshake deal authentication , record protocol handle integrity and confidentiality

IP

SSL Protocol Stack

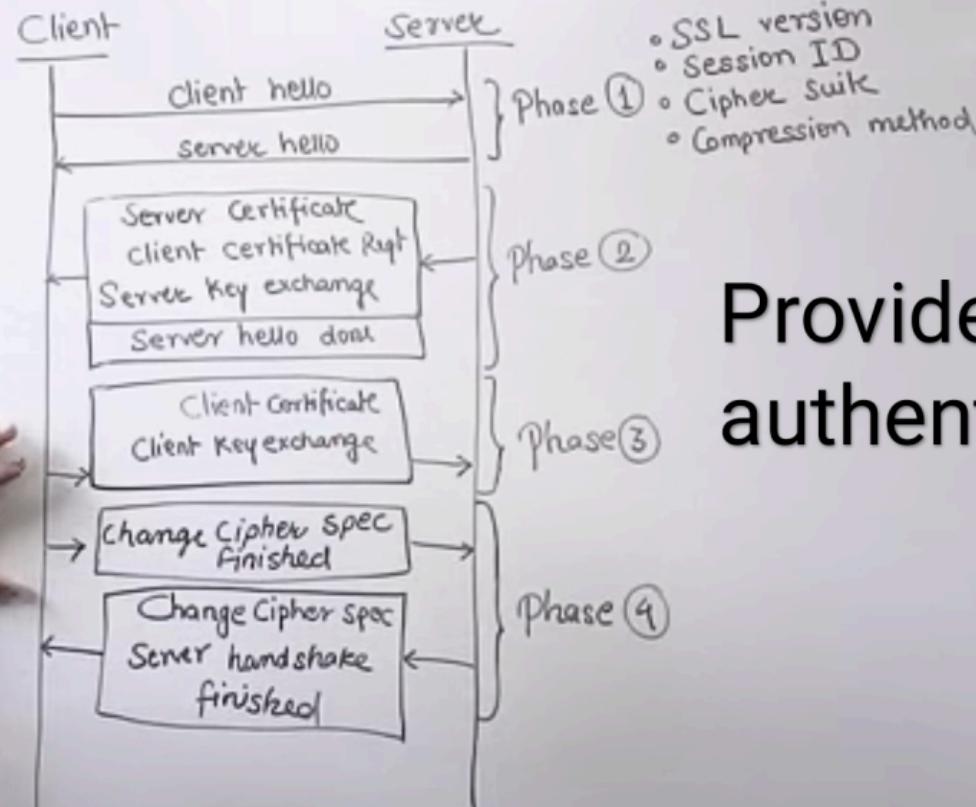
AL  
SSL  
TL

## SSL Record Protocol



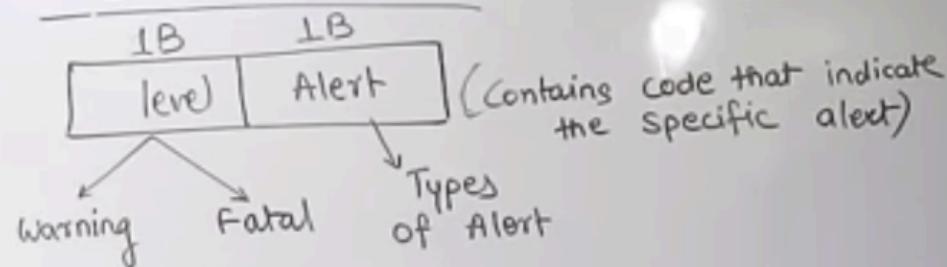
from p to c compression  
must be lossless  
or optional

## SSL Handshake Protocol



Provides  
authentication

## SSL Alert Protocol



Alert msg : ① Description

- Close\_notify → No more message sender
- Unexpected\_message → Incorrect message received.
- bad\_record\_mac → Wrong mac received.
- bad\_certificate → Received a corrupted certificate
- Certificate expired

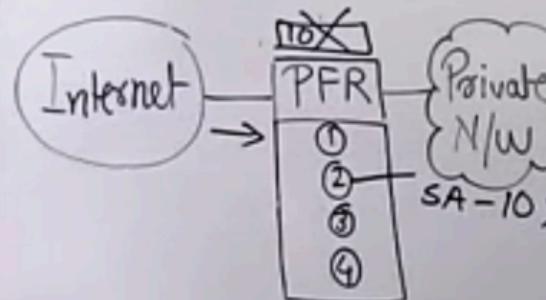
## FIREWALL TYPES

### I) Packet filtering

- Set of Rules
- SA, DA, Port number, Protocols.
- If Rule matched then forward or discard.

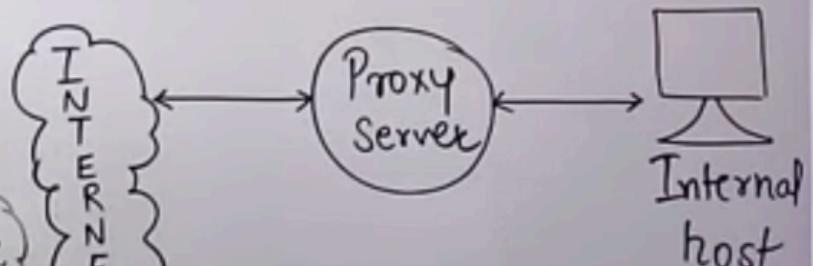
Default action.

- Data / Payload  $\times$



### II) Application - el Gateway (proxy server).

- More Secure
- Processing Overhead.
- Check data / Payload.



## FIREWALL TYPES

### Packet filtering

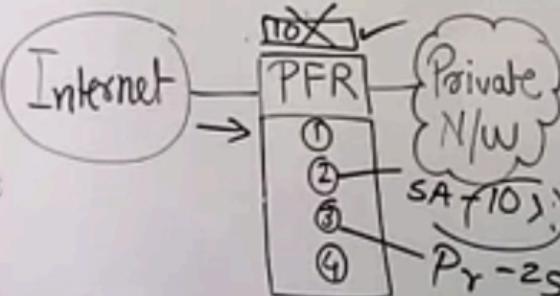
Set of Rules

SA, DA, Port number,  
Protocols.

If Rule matched  
then forward  
or  
discard.



- Default action.
- Data / Payload ~~(X)~~



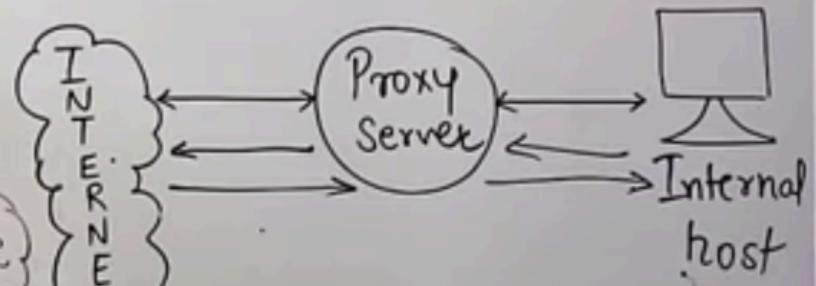
II

Application - el Gateway  
(proxy server).

◦ More Secure

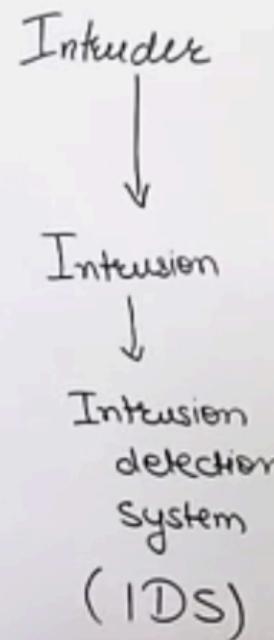
◦ Processing Overhead.

◦ Check data / Payload.



packet filtering is less secure  
it check the source destination  
not data , data may be malicious  
application gateway is secure  
proxy firewall check data also

## Intrusion Detection System



- Outside Intruder (Masquerader)
- Inside Intruder (misfeasor)

intruder are the unauthorised person with criminal mindset to break or steal information  
insider is dangerous

# Intrusion Detection System

## Methods

- Signature Based IDS
  - Pattern
  - Db of attack pattern
  - Detect Known attack
  - Cannot Identify new attack.
- Anomaly Based IDS
  - deviation.

deviation , do task  
other than its normal  
tasks

## IDS TYPES

### I NIDS

- Network Based
- Analysis: Matches traffic to the library of known attack.
- Monitors, Capture & analyze network traffic.
- Detect malicious data present into packets.
- NIDS Analysis very difficult in busy n/w.

### II HIDS

- Host Based
  - installed on individual host or device on network.
  - It monitor data packets from the device only and will alert the admin if suspicious activity is detected
  - Snapshot
- Existing  $\rightleftharpoons$  Previous system
- files deleted or modified