



## Unit-3-bas-26

### BAS-26 OPTIMIZATION TECHNIQUES

**Course category** : Basic Sciences & Maths (BSM)  
**Pre-requisite Subject** : NIL  
**Contact hours/week** : Lecture : 3, Tutorial : 1, Practical: 0  
**Number of Credits** : 4  
**Course Assessment methods** : Continuous assessment through tutorials, attendance, home assignments, quizzes and Three Minor tests and One Major Theory Examination  
**Course Outcomes** : The students are expected to be able to demonstrate the following knowledge, skills and attitudes after completing this course

1. To find the root of a curve using iterative methods.
2. To interpolate a curve using Gauss, Newton's interpolation formula.
3. Use the theory of optimization methods and algorithms developed for various types of optimization problems.
4. To apply the mathematical results and numerical techniques of optimization theory to Engineering problems.

#### Topics Covered

**UNIT-I** 9  
**Classical Optimization Techniques:** Single variable optimization, Multi-variable with no constraints. Non-linear programming: One Dimensional Minimization methods. Elimination methods: Fibonacci method, Golden Section method.  
**UNIT-II** 9  
**Linear Programming: Constrained Optimization Techniques:** Simplex method, Solution of System of Linear Simultaneous equations, Revised Simplex method, Transportation problems, Karmarkar's method, Duality Theorems, Dual Simplex method, Decomposition principle.  
**UNIT-III** 9  
**Non-Linear Programming: Unconstrained Optimization Techniques:** Direct search methods: Random jumping method, Univariate method, Rosenbrock's method. Indirect search methods: Steepest Descent method, Cauchy-Newton Methods, Newton's method.  
**UNIT-IV** 9  
**Geometric Programming:** Polynomial, Unconstrained minimization problem, Degree of difficulty. Solution of an unconstrained **Geometric** Programming problem. Constrained minimization complementary Geometric Programming, Application of Geometric Programming.

#### Books & References

1. Engineering Optimization- S.S. Rao, New Age International
2. Applied Optimal Design-E.J. Haug and J.S. Arora; Wiley New York
3. Optimization for Engineering Design-Kalyanmoy Deb; Prentice Hall of India

You tube Channel:- [https://www.youtube.com/watch?v=l\\_fhwyndmIU](https://www.youtube.com/watch?v=l_fhwyndmIU)

## Indirect Search Methods

### Gradient of a function

the gradient of a function in an  $n$ -dimensional component vector is given by

$$\nabla f_{n \times 1} = \begin{Bmatrix} \partial f / \partial x_1 \\ \partial f / \partial x_2 \\ \vdots \\ \partial f / \partial x_n \end{Bmatrix} \quad \dots (6.56)$$

✓ Steepest Descent (Cauchy) Method → the use of the negative of the gradient vector as a direction for minimization was first made by Cauchy in 1847. ✓  
In this method, we start from an initial trial point  $x_1$  and iteratively move along the steepest descent directions until the optimum point is found. this method is summarized as ✓

1. Start with an arbitrary initial point  $x_1$ .
2. Find the search direction  $S_i$  as

$$S_i = -\nabla f_i = -\nabla f(x_i)$$

3. Determine the optimal step length  $d_i^*$  in the direction  $S_i$  and set

$$✓ \underline{x_{i+1} = x_i + d_i^* S_i = x_i - d_i^* \nabla f_i} \quad (6.70)$$

$f(x)$  Analytic

$$\nabla f(x_i) \quad x_i = \{x_1, x_2\}$$
$$\nabla f(x_i) = \begin{Bmatrix} \partial f / \partial x_1 \\ \partial f / \partial x_2 \end{Bmatrix}$$

4. Test the new point,  $x_{i+1}$  for optimality. If  $x_{i+1}$  is optimum, stop the process, otherwise, go to step 5.

5. Set new iteration number  $i = i+1$  and go to step 2.

Q. Minimize  $f(x_1, x_2) = x_1 - x_2 + 2x_1^2 + 2x_1x_2 + x_2^2$  starting from the point  $x_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$

Iteration 1.

the gradient of  $f$  is given by

$$\nabla f = \begin{Bmatrix} \partial f / \partial x_1 \\ \partial f / \partial x_2 \end{Bmatrix} = \begin{Bmatrix} 1 + 4x_1 + 2x_2 \\ -1 + 2x_1 + 2x_2 \end{Bmatrix}$$

$$\Rightarrow \nabla f_1 = \begin{Bmatrix} 1 + 4 \cdot 0 + 2 \cdot 0 \\ -1 + 2 \cdot 0 + 2 \cdot 0 \end{Bmatrix} = \begin{Bmatrix} 1 \\ -1 \end{Bmatrix}$$

therefore

$$s_1 = -\nabla f_1 = -\begin{Bmatrix} 1 \\ -1 \end{Bmatrix} = \begin{Bmatrix} -1 \\ 1 \end{Bmatrix}$$

To find  $x_2$ , we need to find the optimum step

length  $d_1^*$ . for this, we minimize

$$\text{we minimize } f(x_1 + d_1 s_1) = f\left(\begin{bmatrix} -d_1 \\ d_1 \end{bmatrix}\right)$$

$$= -d_1 - d_1 + 2d_1^2 - 2d_1^2 + d_1^2$$

$$= d_1^2 - 2d_1$$

$$\text{now } \frac{\partial f}{\partial d_1} = 2d_1 - 2 = 0 \Rightarrow d_1^* = 1$$

$$\begin{matrix} (0,0) & (1,-1) \\ \downarrow & \downarrow \\ \begin{bmatrix} 1-4+2 \\ -1-2+2 \end{bmatrix} & \Rightarrow \begin{bmatrix} -1 \\ -1 \end{bmatrix} \end{matrix}$$

$$x_1 + d_1 s_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix} + d_1 \begin{bmatrix} -1 \\ 1 \end{bmatrix} = \begin{bmatrix} -d_1 \\ d_1 \end{bmatrix}$$

$$x_1 = -1, x_2 = 1$$

1.0  
xj  
tce

we obtain

$$x_2 = x_1 + d_1^* S_1 = \begin{Bmatrix} 0 \\ 0 \end{Bmatrix} + 1 \begin{Bmatrix} -1 \\ 1 \end{Bmatrix} = \begin{Bmatrix} -1 \\ 1 \end{Bmatrix} \checkmark$$

$$\text{as } \nabla f_2 = \nabla f(x_2) = \begin{Bmatrix} 1-4+2 \\ -1-2+2 \end{Bmatrix} = \begin{Bmatrix} -1 \\ -1 \end{Bmatrix} \checkmark$$

$\neq \begin{Bmatrix} 0 \\ 0 \end{Bmatrix}$ ,  $x_2$  is not optimum.

Iteration 2.

$$S_2 = -\nabla f_2 = \begin{Bmatrix} 1 \\ 1 \end{Bmatrix}$$

$$f(x_2 + d S_2) = f\left(\begin{Bmatrix} -1 \\ 1 \end{Bmatrix} + d_1 \begin{Bmatrix} 1 \\ 1 \end{Bmatrix}\right) = f(-1+d_2, 1+d_2) \\ = 5d_2^2 - 2d_2 - 1$$

$$\frac{df}{dd_2} = 10d_2 - 2 = 0 \Rightarrow d_2^* = \frac{1}{5}$$

$$x_3 = x_2 + d_2^* S_2 = \begin{Bmatrix} -1 \\ 1 \end{Bmatrix} + \frac{1}{5} \begin{Bmatrix} 1 \\ 1 \end{Bmatrix} = \begin{Bmatrix} -0.8 \\ 1.2 \end{Bmatrix}$$

$$= \begin{Bmatrix} -1+0.2 \\ 1+0.2 \end{Bmatrix} = \begin{Bmatrix} -0.80 \\ 1.20 \end{Bmatrix} \Rightarrow x_3 = \begin{Bmatrix} -0.80 \\ 1.20 \end{Bmatrix} \checkmark$$

$$\text{now, } \nabla f_3 = \nabla f(x_3) = \begin{Bmatrix} 1-3.20+2.4 \\ -1-0.8+2.40 \end{Bmatrix} = \begin{Bmatrix} 0.2 \\ -0.2 \end{Bmatrix} \neq 0$$

similarly we calculate  $x_4$ , and so on.

$$\nabla f_i \neq 0$$

$$x_1 = -1 + d_2 \\ x_2 = 1 + d_2$$

$\Rightarrow$  this sol<sup>n</sup> is not optimum

$\Rightarrow$

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### Newton's Method.

consider the quadratic approximation of the function  $f(x)$  at  $x = x_i$  using the Taylor's series expansion

$$f(x) = f(x_i) + \nabla f_i^T (x - x_i) + \frac{1}{2} (x - x_i)^T [J_i] (x - x_i) \quad (6.95)$$

where  $[J_i] = [J]_{x_i}$  is the matrix of second partial derivatives at point  $x_i$  (Hessian Matrix)

By set the partial derivatives of Equation (6.95) equal to zero for minimum of  $f(x)$ , we obtain

$$\frac{\partial f(x)}{\partial x_j} = 0, \quad j = 1, 2, \dots, n \quad (6.96)$$

Equation (6.96) and (6.95) give

$$\nabla f = \nabla f_i + [J_i] (x - x_i) = 0 \quad (6.97)$$

If  $[J_i]$  is non-singular, equation (6.97) can be solved to obtain an improved approximation  $(x = x_{i+1})$  as

$$x_{i+1} = x_i - [J_i]^{-1} \nabla f_i \quad (6.98)$$

The sequence of points  $x_1, x_2, \dots, x_{i+1}$  can be shown to converge to the actual solution  $x^*$  from any initial point  $x_1$  sufficiently close to the solution  $x^*$ , provided that  $[J_i]$  is non-singular

$$x_{i+1} = x_i - [J_i]^{-1} \nabla f_i$$

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$$x_1 = 0.5, \quad x_2 = 1, \quad x_3 = 1.1$$

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$$2.005$$

$$x^* \rightarrow [1 - u]$$

$$x_0 = 0.5 \quad (3.5)$$

a. Minimize  $f(x_1, x_2) = x_1 x_2 + 2x_1^2 + 2x_1 x_2 + x_2^2$  by taking

the starting point as  $\begin{bmatrix} 0 \\ 0 \end{bmatrix} = x_0$

b. To find  $x_1$ , we require  $[J_1]^{-1}$

$$[J_1] = \begin{bmatrix} \frac{\partial^2 f}{\partial x_1^2} & \frac{\partial^2 f}{\partial x_1 \partial x_2} \\ \frac{\partial^2 f}{\partial x_1 \partial x_2} & \frac{\partial^2 f}{\partial x_2^2} \end{bmatrix} = \begin{bmatrix} 4 & 2 \\ 2 & 2 \end{bmatrix}$$

$$[J_1]^{-1} = \frac{1}{4} \begin{bmatrix} 2 & -2 \\ -2 & 4 \end{bmatrix} = \begin{bmatrix} \frac{1}{2} & -\frac{1}{2} \\ -\frac{1}{2} & 1 \end{bmatrix}$$

$$g_1 = \nabla f_i = \begin{bmatrix} \frac{\partial f}{\partial x_1} \\ \frac{\partial f}{\partial x_2} \end{bmatrix}_{x_1} = \begin{bmatrix} 1 + 4x_1 + 2x_2 \\ -1 + 2x_1 + 2x_2 \end{bmatrix}_{[0,0]} = \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

Now

$$f(x) = f(x_0) + (x - x_0) f'(x_0)$$

$$+ \frac{(x - x_0)^2}{2!} f''(x_0) + \dots$$

$$(x_1, x_2, \dots, x_n)$$

$$[J_i] = [J]_{x_i}$$

$$f(x)$$

$$\frac{\partial f}{\partial x} = 0$$

$$\rightarrow 0$$

$$\nabla f = \nabla f_i + [J_i] (x - x_i) = 0$$

$$0 = [J_i]^{-1} \nabla f_i$$

$$+ (x - x_i)$$

$$x - x_i = -[J_i]^{-1} \nabla f_i$$

$$x = x_i - [J_i]^{-1} \nabla f_i$$

$$x_0 = 1, \quad x_1 = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$$

$$x_{i+1} = x_i - [J_i]^{-1} \nabla f_i$$

$$x_2 = x_1 - [J_1]^{-1} \nabla f_1$$

$$x_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix}, [J] = \begin{bmatrix} 4 & 2 \\ 2 & 2 \end{bmatrix}$$

$$\frac{\partial f}{\partial x_1} = 1 + 4x_1 + 2x_2$$

$$\frac{\partial f}{\partial x_2} = -1 + 2x_1 + 2x_2$$

$$1 + 4x_1 + 2x_2$$

$$= \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

$$x_2 = x_1 - [1, 1]^{-1} \nabla f_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix} - \begin{bmatrix} \frac{1}{2} & -\frac{1}{2} \\ -\frac{1}{2} & 1 \end{bmatrix} \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

$$= \begin{bmatrix} 0 \\ 0 \end{bmatrix} - \begin{bmatrix} \frac{1}{2} + \frac{1}{2} \\ -\frac{1}{2} - 1 \end{bmatrix} = \begin{bmatrix} -1 \\ +3/2 \end{bmatrix} \Rightarrow x_2 = \begin{bmatrix} -1 \\ 3/2 \end{bmatrix}$$

To see  $x_2$  is the optimum point, we check

$$g_2 = \nabla f_2 = \begin{bmatrix} \frac{\partial f}{\partial x_1} \\ \frac{\partial f}{\partial x_2} \end{bmatrix} = \begin{bmatrix} 1 + 4x_1 + 2x_2 \\ -1 + 2x_1 + 2x_2 \end{bmatrix}$$

$$= \begin{bmatrix} 1 - 4 + 3 \\ -1 - 2 + 3 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

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Try

$$\nabla f_1 = (\nabla f)_{x=x_1} = \begin{bmatrix} 1 + 4x_1 + 2x_2 \\ 1 + 2x_1 + 2x_2 \end{bmatrix}_{(0,0)}$$

$$= \begin{bmatrix} 1 \\ -1 \end{bmatrix} \text{ if } \nabla f_1 = 0$$

$$\nabla f_1 \neq 0$$

Proceed further

$$g_2 = \nabla f_2 = \begin{bmatrix} 1 + 4x_1 + 2x_2 \\ -1 + 2x_1 + 2x_2 \end{bmatrix}_{(-1, 3/2)}$$

$$= \begin{bmatrix} 1 - 4 + 3 \\ -1 - 2 + 3 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

as here  $\nabla f_2 = 0$

$$\Rightarrow x_2 = \begin{bmatrix} -1 \\ 3/2 \end{bmatrix} \text{ is the}$$

Optimum soln.

as  $g_2 = 0$ ,  $x_2$  is the optimum point.

Q. Minimize  $f(x_1, x_2) = 10(x_1^2 - x_2)^2 + (1 - x_1)^2$  taking

$$\nabla f = \begin{bmatrix} \frac{\partial f}{\partial x_1} \\ \frac{\partial f}{\partial x_2} \end{bmatrix} = \begin{bmatrix} 20(x_1^2 - x_2) \cdot 2x_1 + 2(1 - x_1) \cdot (-1) \\ 20(x_1^2 - x_2) \cdot (-1) \end{bmatrix}_{(-2, -2)}$$

$$= \begin{bmatrix} 40(4 + 2) \cdot (-2) + 2(1 + 2) \cdot (-1) \\ 20(4 + 2) \cdot (-1) \end{bmatrix} = \begin{bmatrix} -480 \\ -120 \end{bmatrix}$$

Q2.  $f = x_1^2 + x_2^2 - 2x_1 - 4x_2 + 5$ ,  $x_1 = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$  by Newton's Method.

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## Univariate Method

Univariate method can be summarized as follows

1) Choose an arbitrary starting point  $x_1$  and set  $i=1$ .

2) Find the search direction  $S_i$  as

$$S_i^T = \begin{cases} (1, 0, 0, \dots, 0) & \text{for } i=1, n+1, 2n+1, \dots \\ (0, 1, \dots, 0) & \text{for } i=2, n+2, 2n+2, \dots \\ (0, 0, 1, \dots, 0) & \text{for } i=3, n+3, 2n+3, \dots \\ \vdots & \vdots \\ (0, 0, 0, \dots, 1) & \text{for } i=n, 2n, 3n, \dots \end{cases}$$

3) Determine whether  $d_i$  should be +ve or negative. For the current direction  $S_i$ , this means find whether the function value decreases in the +ve or -ve direction. For this, we take a small probe length ( $\epsilon$ ) and evaluate  $f_i = f(x_i)$ ,  $f^+ = f(x_i + \epsilon S_i)$  and  $f^- = f(x_i - \epsilon S_i)$ . If  $f^+ < f_i$ ,  $S_i$  will be the correct direction for decreasing the value of  $f$  and if  $f^- < f_i$ ,  $-S_i$  will be the correct one. If both  $f^+$  and  $f^-$  are greater than  $f_i$ , we take  $x_i$  as the minimum along the direction  $S_i$ .

4) Find the optimal step length  $d_i^*$  such that

$$f(x_i \pm d_i^* S_i) = \min_{d_i} f(x_i \pm d_i S_i)$$

where + or - sign has to be used depending upon whether  $S_i$  or  $-S_i$  is the direction for decreasing the function value.

5) Set  $x_{i+1} = x_i \pm d_i^* S_i$  depending on the direction for decreasing the function value, and

$$f_{i+1} = f(x_{i+1}).$$

6) Set the new value of  $i = i+1$  and go to step 2. Continue this procedure until no significant change is achieved in the value of the objective function.

Q. Minimize  $f(x_1, x_2) = x_1 - x_2 + 2x_1^2 + 2x_1x_2 + x_2^2$

$$\begin{array}{l|l} x_1 & (x_1, x_2, x_3) \\ \hline (2) & f_1(1, 0) \\ S_1 = (1, 0) & S_2 = (0, 1, 0) \\ S_2 = (0, 1) & S_3 = (0, 0, 1) \\ S_3 = (1, 0) & S_4 = S_1 \\ S_4 = (0, 1) & S_5 = S_2 \end{array}$$

$\epsilon \rightarrow$  small probe length

$$f_i = f(x_i)$$

$$f_i^+ = f(x_i + \epsilon S_i)$$

$$f_i^- = f(x_i - \epsilon S_i)$$

if  $f^+ < f_i$ ,  $S_i$  will be the correct direction

$f_i^- < f_i \Rightarrow -S_i$  will be correct direction.

$$\epsilon = 0.01, 0.001$$

$$x_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$



Q. Minimize  $f(x_1, x_2) = x_1 - x_2 + 2x_1^2 + 2x_1x_2 + x_2^2$   
 with the starting point  $(0,0)$   
 Let probe length  $(\epsilon)$  as  $0.01$ ,  
 iteration  $i=1$

Step 1. Choose the search direction  $S_1$  as  $S_1 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$

Step 3. To find whether the value of  $f$  decreases along  $S_1$  or  $-S_1$ , we use the probe length  $\epsilon$ . Since

$$f_1 = f(x_1) = f(0,0) = 0, \quad x_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$f^+ = f(x_1 + \epsilon S_1) = f\left(\begin{pmatrix} 0 \\ 0 \end{pmatrix} + \epsilon \begin{pmatrix} 1 \\ 0 \end{pmatrix}\right) = f(\epsilon)$$

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$$x_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix},$$

$$S_1 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$f_1 = f(x_1) = f(0,0) = 0$$

$$f_1^+ = f(x_1 + \epsilon S_1) = f\left(\begin{bmatrix} 0 \\ 0 \end{bmatrix} + 0.01 \begin{bmatrix} 1 \\ 0 \end{bmatrix}\right)$$

$$= f(0.01, 0) = f(\epsilon, 0)$$

$$= 0.0102 > f_1$$

$$\Rightarrow f_1^+ > f_1 \quad (f_1^+ \neq f_1)$$

$$f_1^- = f(x_1 - \epsilon S_1)$$

$$= f\left(\begin{bmatrix} -\epsilon \\ 0 \end{bmatrix}\right) = f\left(\begin{bmatrix} -0.01 \\ 0 \end{bmatrix}\right)$$

$$= -0.9998 < f_1 = 0$$

$$\Rightarrow f_1^- < f_1$$

$\Rightarrow -S_1$  is the correct direction.

$$f(x_1 - d_1 S_1) = f(-d_1, 0)$$

$$= -d_1 + 2d_1^2$$

$$\frac{\partial f}{\partial d_1} = -1 + 4d_1 = 0$$

$$\Rightarrow d_1 = 1/4$$

$$\Rightarrow d_1^* = 1/4$$

$$x_2 = x_1 - d_1 S_1$$

$$= \begin{bmatrix} 0 \\ 0 \end{bmatrix} - \frac{1}{4} \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$= \begin{bmatrix} -1/4 \\ 0 \end{bmatrix}$$

Step 2.  $x_2, \epsilon, S_2 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$

$$f_2 = f(x_2)$$

$$f^- = f(x_1 - \epsilon S_1) = f(-\epsilon, 0) = -0.01 - 0 + 2(0.0001) + 0 + 0 = 0.0102 > f_1 \checkmark$$

$-S_1$  is the correct direction for minimizing  $f$  from  $x_1$ .

To find the optimum step length  $d_1^*$ , we minimize

$$f(x_1 - d_1 S_1) = f\left(\begin{pmatrix} 0 \\ 0 \end{pmatrix} - d_1 \begin{pmatrix} 1 \\ 0 \end{pmatrix}\right) = f(-d_1, 0)$$

$$= f(-d_1, 0)$$

$$= (-d_1) - 0 + 2(-d_1)^2 + 0 + 0 = -d_1 + 2d_1^2$$

$$\frac{\partial f}{\partial d_1} = -1 + 4d_1 = 0 \Rightarrow d_1 = 1/4 \text{ we have } d_1^* = 1/4$$

$$\text{Set } x_2 = x_1 - d_1^* S_1 = \begin{bmatrix} 0 \\ 0 \end{bmatrix} - \frac{1}{4} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} -1/4 \\ 0 \end{bmatrix}$$

$$f_2 = f(x_2) = f(-1/4, 0) = -1/8 \checkmark$$

iteration  $i=2$ , choose the search direction  $S_2$  as  $S_2 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$

$$f_2 = f(x_2) = -1/8 = -0.125$$

$$f^+ = f(x_2 + \epsilon S_2) = f\left(\begin{bmatrix} -1/4 \\ 0 \end{bmatrix} + 0.01 \begin{bmatrix} 0 \\ 1 \end{bmatrix}\right)$$

$$= f(-0.25, 0.01) = -0.1399 < f_2 = -0.125$$

$$f^- = f(x_2 - \epsilon S_2) = f(-0.25, -0.01) = -0.1099 > f_2$$

$\Rightarrow S_2$  is the correct direction for decreasing the value of  $f$  from  $x_2$ .

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$$f < f_2^+ \Rightarrow S_2$$

$$f < f_2^+ \Rightarrow S_2$$

$$f_2 < f_2^- \Rightarrow -S_2$$

we minimize  $f(x_L + d_L S_L)$  to find  $d_L^*$  ✓  
 here  
 $f(x_L + d_L S_L) = f\left(\begin{pmatrix} -0.25 \\ 0 \end{pmatrix} + d_L \begin{pmatrix} 0 \\ 1 \end{pmatrix}\right)$   
 $= f(-0.25, d_L) = d_L^2 - 1.5d_L - 0.125$  ✓  
 $\frac{\partial f}{\partial d_L} = 2d_L - 1.5 = 0 \Rightarrow d_L^* = \frac{1.5}{2} = 0.75$  ✓

Set  $x_3 = x_L + d_L^* S_L = \begin{pmatrix} -0.25 \\ 0 \end{pmatrix} + 0.75 \begin{pmatrix} 0 \\ 1 \end{pmatrix}$   
 $= \begin{pmatrix} -0.25 \\ 0.75 \end{pmatrix}$  ✓ ✓

$$f(x_3) = -0.6875$$

Q. Minimize  $= x_1 - x_2 + x_3 + 2x_1^2 + 2x_2^2 - x_3^2 + 2x_1x_3 + 4x_2x_3 - 6x_1x_2$  ✓

$x_1 = \begin{pmatrix} 1 \\ 2 \\ 1 \end{pmatrix}, S_1 = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}, \epsilon = 0.01$   $-S_2$  is the correct direction

$f_1 = f(x_1) = 7$  ✓

$f_1^+ = 6.9502 < f_1$  ✓

$f_1^- = 7.0502$  ✓

$\Rightarrow S_1$  is the correct direction

$f(x_1 + d_1 S_1) = f\left(\begin{pmatrix} 1+d_1 \\ 2 \\ 1 \end{pmatrix}\right) = 7 - 5d_1 + 2d_1^2$

$\frac{\partial f(x_1 + d_1 S_1)}{\partial d_1} = -5 + 4d_1 \Rightarrow d_1 = 5/4$  ✓

$\Rightarrow x_2 = \begin{pmatrix} 9/4 \\ 2 \\ 1 \end{pmatrix}$  ✓

$f(x_2) = f_2 = \frac{107}{25} = 4.28$

$f_2^+ = 4.2822$

$f_2^- = 4.2782 < f(x_2) = f_2$  ✓

$$f_2 = f(x_2)$$

$$f_2^+ = f(x_2 + \epsilon S_2)$$

$$f_2^- = f(x_2 - \epsilon S_2)$$

$$\frac{\partial f_3}{\partial x_2} = 0$$

$$1 + \frac{S}{4} = 2/4$$

## Geometric Programming.

Posynomial - the objective function  $f(x)$  is given by the sum of several component costs  $U_i(x)$  as

$$f(x) = U_1 + U_2 + \dots + U_n$$

In many cases, the component cost  $U_i$  can be expressed as power functions of the type

$$U_i = c_i x_1^{a_{i1}} x_2^{a_{i2}} \dots x_n^{a_{in}}, \text{ where}$$

the coefficients  $c_i$  are +ve constants, the exponent  $a_{ij}$  are real constants (+ve, zero, -ve) and the variables  $x_1, x_2, \dots, x_n$  are taken to be +ve. Functions  $f$  because of +ve coefficients and variables and real exponents are called posynomials. For example

$$f(x_1, x_2, x_3) = 6 + 3x_1 - 8x_2 + 7x_3 + 2x_1x_3 - 3x_1x_3 + \frac{4}{3}x_2x_3 + \frac{8}{7}x_1^2 - 9x_2^2 + x_3^2 \text{ is a second-degree polynomial in variables } x_1, x_2, x_3$$

while

$$g(x_1, x_2, x_3) = x_1x_2x_3 + x_1x_2 + 4x_3 + \frac{2}{x_1x_2} + 5x_3^{-1/2}$$

is a posynomial

# Unconstrained minimization Problem

Unconstrained minimizing problem

Find  $x = \begin{Bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{Bmatrix}$  that minimizes the objective function

$$J(x) = \sum_{j=1}^N U_j(x) = \sum_{j=1}^N (c_j x_1^{a_{j1}} x_2^{a_{j2}} \dots x_n^{a_{jn}}) \quad \text{--- (P.5)}$$

where  $c_j > 0, x_j > 0$  and  $a_{ij}$  is real const.

Solution of an unconstrained geometric Problem using Differential calculus,

As minimizing objective function is

$$J(x) = \sum_{j=1}^N U_j(x) = \sum_{j=1}^N (c_j x_1^{a_{j1}} x_2^{a_{j2}} \dots x_n^{a_{jn}})$$

For maxima or minima

$$\frac{\partial J(x)}{\partial x_k} = \sum_{j=1}^N \frac{\partial U_j}{\partial x_k} = 0$$

$k = 1, \dots, n$ , means  $n$  variables.

Use multipliers

$$x_k \frac{\partial J}{\partial x_k} = \sum_{j=1}^N a_{kj} U_j(x) = 0, \quad k = 1 \text{ to } n \quad \text{--- (P.6)}$$

To find the minimizing vector  $x^* = \begin{Bmatrix} x_1^* \\ x_2^* \\ \vdots \\ x_n^* \end{Bmatrix}$  we have

$$\sum_{j=1}^N a_{kj} U_j(x^*) = 0 \quad \rightarrow k = 1, 2, \dots, n \quad \text{--- (P.6)}$$

now, divide (P.6) by  $J^*$  (min-value of  $J$ )

$$\sum_{j=1}^N a_{kj} \frac{U_j(x^*)}{J^*} = \sum_{j=1}^N a_{kj} \Delta_j^* \quad \text{--- (P.7)}$$

where  $\Delta_j^* = \frac{U_j(x^*)}{J^*}$

this relation (P.7) is orthogonality condition

$$\Delta_j^* = \frac{U_j(x^*)}{J^*} \quad \text{--- (P.7)}$$

$$\Rightarrow \sum_{j=1}^N \frac{U_j(x^*)}{J^*} = 1 \Rightarrow \sum_{j=1}^N \Delta_j^* = 1$$

this condition (P.8) is called as normality condition. now,  $\Delta_j^*$

$$J^* = \begin{pmatrix} \frac{U_1(x^*)}{J^*} \\ \frac{U_2(x^*)}{J^*} \\ \vdots \\ \frac{U_N(x^*)}{J^*} \end{pmatrix} \begin{pmatrix} \Delta_1^* \\ \Delta_2^* \\ \vdots \\ \Delta_N^* \end{pmatrix} \quad \text{--- (P.12)}$$

$$f^{\circ} = \left( \frac{c_1}{\Delta_1^{\circ}} \right)^{\Delta_1^{\circ}} \left( \frac{c_2}{\Delta_2^{\circ}} \right)^{\Delta_2^{\circ}} \left( \frac{c_3}{\Delta_3^{\circ}} \right)^{\Delta_3^{\circ}} \dots \left( \frac{c_N}{\Delta_N^{\circ}} \right)^{\Delta_N^{\circ}}$$

$$f^{\phi} = \left( \frac{c_1}{\Delta_1^{\phi}} \right)^{\Delta_1^{\phi}} \left( \frac{c_2}{\Delta_2^{\phi}} \right)^{\Delta_2^{\phi}} \left( \frac{c_3}{\Delta_3^{\phi}} \right)^{\Delta_3^{\phi}} \dots \left( \frac{c_N}{\Delta_N^{\phi}} \right)^{\Delta_N^{\phi}}$$

if  $N = n+1$ , there will be as many linear <sup>objective function</sup> simultaneous eq. as the

if  $N - n - 1 = 0$ , the problem is said to have a zero degree of difficulty. If  $N > n + 1 \Rightarrow$  we have

more no. of variables than the equations, then sometimes this method is not applicable.

Unknown  $\vec{o}_j$  can be determined uniquely from the orthogonality and normality conditions

$$J(x) = 80x_1x_2 + 40x_1x_3 + 20x_1x_3 + \frac{80}{x_1x_2} \text{ solve}$$

$c_1 = 80, c_2 = 40, c_3 = 20, c_4 = 80$

$$J(x) = \sum_{j=1}^N (j^{a_1} j^{a_2} \dots j^{a_n})$$

$$= c_1 a_{11} x_1 a_{21} x_2 a_{31} x_3 + c_2 a_{12} x_1 a_{22} x_2 a_{32} x_3 + c_3 a_{13} x_1 a_{23} x_2 a_{33} x_3 + c_4 a_{14} x_1 a_{24} x_2 a_{34} x_3$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \end{bmatrix} = \begin{bmatrix} 1 & 0 & 1 & -1 \\ 1 & 1 & 0 & -1 \\ 0 & 1 & 1 & -1 \end{bmatrix}$$

the orthogonality and normality conditions are given by

$$\begin{bmatrix} 1 & 10 \end{bmatrix}$$

$$\sum_{j=1}^N \Delta_j^* a_{kj} = 0 \quad k=1 \dots n \quad (2.7)$$

$$j=1$$

$$\Delta_1 a_{k1} + \Delta_2 a_{k2} + \Delta_3 a_{k3} + \Delta_4 a_{k4} = 0$$

$$\Delta_1 a_{11} + \Delta_2 a_{12} + \Delta_3 a_{13} + \Delta_4 a_{14} = 0$$

$$\Delta_1 a_{21} + \Delta_2 a_{22} + \Delta_3 a_{23} + \Delta_4 a_{24} = 0$$

$$\Delta_1 a_{31} + \Delta_2 a_{32} + \Delta_3 a_{33} + \Delta_4 a_{34} = 0$$

or orthogonal conditions

$$\sum_{j=1}^N \Delta_j = 1 \quad \text{in orthogonal condition}$$

$$\Rightarrow \Delta_1 + \Delta_2 + \Delta_3 + \Delta_4 = 1 \quad (3)$$

$$\Delta_1 + 0 + \Delta_3 + \Delta_4 = 0$$

$$\Delta_1 + \Delta_2 + 0 + \Delta_4 = 0$$

$$0 + \Delta_2 + \Delta_3 + \Delta_4 = 0$$

Solving these 4 eqs

$$\Delta_4 = \Delta_1 + \Delta_3 = \Delta_1 + \Delta_2 \Rightarrow \Delta_2 = \Delta_3$$

$$\text{and } \Delta_4 = \Delta_1 + \Delta_3 = \Delta_2 + \Delta_3 \Rightarrow \Delta_1 = \Delta_2$$

$$\Rightarrow \Delta_1 = \Delta_2 = \Delta_3 \Rightarrow \Delta_4 = 2\Delta_1$$

$$\text{from eq (3), } \Delta_1 + \Delta_1 + \Delta_1 + \Delta_1 + \Delta_1 + \Delta_1 + \Delta_1 + \Delta_1 = 1$$

$$\Rightarrow \Delta_1 = \frac{1}{5} \Rightarrow \Delta_1^* = \Delta_2^* = \Delta_3^* = \frac{1}{5}, \Delta_4^* = \frac{2}{5}$$

So, optimal value of the objective function is

$$J^* = \sum_{j=1}^n \left( \frac{c_j}{\Delta_j^*} \right) \Delta_j^* = \left( \frac{80}{1/5} \right)^{1/5} \left( \frac{40}{1/5} \right)^{1/5} \left( \frac{20}{1/5} \right)^{1/5} \left( \frac{80}{2/5} \right)^{1/5}$$

$$= (400)^{1/5} \times (200)^{1/5} \times (100)^{1/5} \times (80)^{1/5}$$

$$= (400 \times 200 \times 100 \times 80)^{1/5}$$

$$= (32 \times 10^6)^{1/5} = 200$$

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Solving these eqn is

$$\gamma_2^* = \frac{1}{2} \gamma_1^* \Rightarrow \frac{1}{\gamma_2^*} = 2 \gamma_1^*$$

$$\gamma_3^* = \frac{1}{\gamma_2^*} = \left( \frac{1}{2} \gamma_1^* \right)^{-1}$$

$$\gamma_3^* = \frac{2}{\gamma_1^*}$$

$$\Rightarrow \frac{2}{\gamma_1^*} = \frac{1}{\gamma_2^*} = 2 \gamma_1^* \Rightarrow \gamma_2^* = 1 \Rightarrow \gamma_1^* = 1.$$

$$\gamma_2^* = \frac{1}{2}, \gamma_3^* = 2$$

Solution of an unconstrained geometric Problem using  
Arithmetic-Geometric inequality

Geometric and Primal Problem (unconstrained)

$$\text{Find } x = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \text{ so that } \min_{x \geq 0} \sum_{j=1}^n c_j x_j \text{ a.t. } x_1 \dots x_n$$

$$x_1 > 0, x_2 > 0, \dots, x_n > 0$$

then geometric dual of Primal Problem is

$$\text{Find } \Delta = \begin{pmatrix} \Delta_1 \\ \vdots \\ \Delta_n \end{pmatrix} \text{ so that}$$

$$\text{Max } V(\Delta) = \prod_{j=1}^n \left( \frac{c_j}{\Delta_j} \right)^{\Delta_j} \text{ or}$$

$$\log \{ \text{Max } V(\Delta) \} = \log \left[ \prod_{j=1}^n \left( \frac{c_j}{\Delta_j} \right)^{\Delta_j} \right]$$

subject to the constraints

$$\sum_{j=1}^n \Delta_j = 1$$

$$\sum_{j=1}^n a_{ij} \Delta_j = 0, \quad i = 1, 2, \dots, n$$

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Dual and Primal of Linear Problem

Primal Problem

Find  $x = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$  such that

$z_0(x) = f(x) \rightarrow \text{Minimum}$

subject to constraints

$g_1(x) \leq 1$

$g_2(x) \leq 1$

$\vdots$

$g_m(x) \leq 1$

with

$z_0(x) = \sum_{j=1}^n c_{0j} x_j$

$g_1(x) = \sum_{j=1}^n c_{1j} x_j$

$g_2(x) = \sum_{j=1}^n c_{2j} x_j$

$\vdots$

$g_m(x) = \sum_{j=1}^n c_{mj} x_j$

where  $c_{ij}$  are real numbers

Dual Problem

Find  $d = \begin{pmatrix} d_1 \\ d_2 \\ \vdots \\ d_m \end{pmatrix}$  such that

$z_1(d) = \sum_{i=1}^m d_i$

subject to constraints

$d_1 \geq 0, d_2 \geq 0, \dots, d_m \geq 0$

$d_{01} \geq 0, d_{02} \geq 0, \dots, d_{0n} \geq 0$

$d_{11} \geq 0, d_{12} \geq 0, \dots, d_{1n} \geq 0$

$d_{21} \geq 0, d_{22} \geq 0, \dots, d_{2n} \geq 0$

$\vdots$

$d_{m1} \geq 0, d_{m2} \geq 0, \dots, d_{mn} \geq 0$

where  $d_{ij}$  are real numbers

Relationship between Primal and Dual

Let  $K = 1$  to  $m$  (no. of constraints)

$n = \text{total no. of variables}$

$z_0 = f = \text{Primal function}$

$m = \text{no. of Primal constraints}$

$N = n_0 + n_1 + \dots + n_m = \text{total number of terms in the Primal function}$

$N - n - 1 = \text{degree of difficulty of Problem}$

Dual function

$z_1 = d_1 + d_2 + \dots + d_m$

$d_{01}, d_{02}, \dots, d_{0n}$  are the normality constraints

$d_{11}, d_{12}, \dots, d_{1n}$  are the orthogonality constraints

$d_{k1} \geq 0, d_{k2} \geq 0, \dots, d_{kn} \geq 0$

$k = 0$  to  $m$

$N = n_0 + n_1 + \dots + n_m$

number of dual variables

$n + 1$  number of dual constraints

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Solve the problem

$$f(b, \theta) = 100b^2\theta^2 + 50b^2\theta^3 + 20b^2\theta^4 - \frac{1}{2}30b^2\theta^2$$

$$\text{RHS: } c_1 = 100, c_2 = 50, c_3 = 20, c_4 = 30$$

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \end{bmatrix} = \begin{bmatrix} 1 & 2 & 0 & -5 \\ 0 & 0 & -1 & 2 \end{bmatrix}$$

the orthogonality and normality condition is

$$\begin{bmatrix} 1 & 2 & 0 & -5 \\ 0 & 0 & -1 & 2 \\ 1 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \quad (1)$$

$$\text{here } m = 4, n = 2, N = m + 1 = 2 + 1$$

$\Rightarrow$  these equations (1) do not yield the

desired  $b_i$  ( $i = 1, 2, 3, 4$ ). So, solving

$b_1, b_2, b_3$  in terms of  $b_4$ , from (1)

$$\begin{cases} b_1 + 2b_2 - 5b_4 = 0 \\ -b_3 + 2b_4 = 0 \\ b_1 + b_2 + b_3 + b_4 = 1 \end{cases} \quad (11)$$

$$\text{from (11) } b_3 = 2b_4$$

$$b_1 = -2b_2 + 5b_4$$

$$-2b_2 + 5b_4 + b_2 + 2b_4 + b_4 = 1$$

$$-b_2 + 8b_4 = 1$$

$$\Rightarrow b_2 = 8b_4 - 1$$

so

$$b_1 = -2(8b_4 - 1) + 5b_4$$

$$= -16b_4 + 2 + 5b_4$$

$$b_1 = -11b_4 + 2$$

$$\begin{cases} b_1 = 2 - 11b_4 \\ b_2 = 8b_4 - 1 \\ b_3 = 2b_4 \end{cases}$$

the dual problem can be written as

Maximize  $V(b_1, b_2, b_3, b_4)$

$$= \left(\frac{c_1}{b_1}\right)^{b_1} \left(\frac{c_2}{b_2}\right)^{b_2} \left(\frac{c_3}{b_3}\right)^{b_3} \left(\frac{c_4}{b_4}\right)^{b_4}$$

$$= \left(\frac{100}{2-11b_4}\right)^{2-11b_4} \left(\frac{50}{8b_4-1}\right)^{8b_4-1} \left(\frac{20}{2b_4}\right)^{2b_4} \left(\frac{30}{b_4}\right)^{b_4}$$

taking log on both sides, get

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$$\begin{aligned}
 \log V &= (2-11\Delta_4) [\log 100 - \ln(2-11\Delta_4)] \Big|_{\ln=10} \\
 &+ (8\Delta_4-1) [\ln 50 - \ln(8\Delta_4-1)] \\
 &+ 2\Delta_4 [\log 20 - \log 2\Delta_4] + \Delta_4 [\ln 300 - \ln \Delta_4] \\
 \text{necessary condition for minimization} \\
 \frac{\partial}{\partial \Delta_4} (\log V) &= -11 [\ln 100 - \ln(2-11\Delta_4)] + \\
 &+ (8\Delta_4-1) \frac{1}{2-11\Delta_4} + 8 [\ln 50 - \ln(8\Delta_4-1)] \\
 &+ (8\Delta_4-1) \left( \frac{1}{8\Delta_4-1} \right) + 2 [\ln 20 - \ln 2\Delta_4] \\
 &+ 2\Delta_4 \left( \frac{1}{2\Delta_4} \right) + 1 [\ln 300 - \ln \Delta_4] \\
 &+ \Delta_4 \left[ -\frac{1}{\Delta_4} \right] = 0 \\
 \Rightarrow -11 [2 - \ln(2-11\Delta_4)] + 11 + 8 \left[ \ln \frac{50}{8\Delta_4-1} \right] - 8 \\
 + 2 \log \left( \frac{20}{2\Delta_4} \right) - 2 + 2 \log \frac{300}{\Delta_4} - 1 &= 0 \\
 \Rightarrow \frac{2.7}{\Delta_4}
 \end{aligned}$$

$$\begin{aligned}
 \Rightarrow \Delta_4 &= \ln \left[ \frac{(100)^{11}}{(50)^8 (20)^2 (300)} \right] + \ln \left[ \frac{(2-11\Delta_4)^{11}}{(8\Delta_4-1)^8 (2\Delta_4)^2 \Delta_4} \right] = 0 \\
 \Rightarrow \frac{(2-11\Delta_4)^{11}}{(8\Delta_4-1)^8 (2\Delta_4)^2 \Delta_4} &= \frac{(100)^{11}}{(50)^8 (20)^2 (300)} \\
 &= \frac{(1 \times 10^2)^{11}}{5^8 \times 10^8 \times 2^2 \times 10^2 \times 3 \times 10^2} = \frac{1 \times 10^{22}}{5^8 \times 10^{12} \times 3 \times 10^2} \\
 &= \frac{10^{10}}{5^8 \times 10^{12}} = 2.130 \\
 \text{Now set value of } \Delta_4 &\text{ by trial method.} \\
 \Delta_4^* &\approx 0.147, \Delta_1^* = 0.385, \Delta_2^* = 0.175 \\
 \Delta_3 &= 0.294 \\
 \Rightarrow V^* = f^* &= \left( \frac{100}{0.385} \right)^{0.385} \left( \frac{50}{0.175} \right)^{0.175} \left( \frac{20}{0.294} \right)^{0.294} \\
 &\times \left( \frac{300}{0.147} \right)^{0.147} = 2.42 \\
 V_1^* = \Delta_1^* &= 0.385 \times 2.42 = 92.2 \\
 V_2^* &= 42.4 \\
 V_3^* &= 71.1 \\
 V_4^* &= 35.6
 \end{aligned}$$

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$$V^* = 100 D^* = 31.2$$

$$\Rightarrow D^* = 0.312, \delta^* = 0.281 \text{ m}^3/\text{s. m}$$

Example 8-3 - zero degree of difficulty problem  
 the optimization problem can be stated as  
 Find  $X = \begin{Bmatrix} x_1 \\ x_2 \\ x_3 \end{Bmatrix}$  so as to minimize  
 $f(X) = 20x_1x_2 + 40x_2x_3 + 80x_1x_3$  subject to  
 $\frac{80}{x_1x_2x_3} \leq 10$  or  $\frac{8}{x_1x_2x_3} < 1$

Ans. First,  $n = \text{no. of variables} = 3$   
 $N_0 = \text{no. of terms in the objective function} = 3$   
 $N_1 = 1$   
 $m = \text{total no. of the constraint} = 1$   
 $N_k = \text{number of terms in } k^{\text{th}} \text{ constraint}$   
 mean  $N_1 = 1 = \text{no. of terms in 1st constraint.}$   
 $N = N_0 + N_1 + \dots + N_m = \text{Total number of terms in the polynomial, mean, here}$   
 $N = 3 + 1 = 4$   
 and  $n - n - 1 = 4 - 3 - 1 = 0 \Rightarrow \text{zero-degree of difficulty problem}$   
 So, dual problem can be stated as

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Find  $\lambda = \begin{Bmatrix} \lambda_{01} \\ \lambda_{02} \\ \lambda_{03} \\ \lambda_{11} \end{Bmatrix}$  to maximize

$$\begin{aligned} v(\lambda) &= \frac{1}{\prod_{k=0}^M \prod_{j=1}^{N_k}} \left( \frac{c_{kj}}{\lambda_{kj}} \sum_{\ell=1}^{N_k} \lambda_{k\ell} \right)^{\lambda_{kj}} \\ &= \frac{1}{\prod_{j=0}^3 \prod_{\ell=1}^{N_0}} \left( \frac{c_{0j}}{\lambda_{0j}} \sum_{\ell=1}^{N_0} \lambda_{0\ell} \right)^{\lambda_{0j}} \frac{1}{\prod_{j=1}^{N_1} \prod_{\ell=1}^{N_1}} \left( \frac{c_{1j}}{\lambda_{1j}} \sum_{\ell=1}^{N_1} \lambda_{1\ell} \right)^{\lambda_{1j}} \\ &= \frac{1}{\prod_{j=0}^3} \left( \frac{c_{0j}}{\lambda_{0j}} (\lambda_{01} + \lambda_{02} + \lambda_{03}) \right)^{\lambda_{0j}} \frac{1}{\prod_{j=1}^{N_1}} \left( \frac{c_{1j}}{\lambda_{1j}} (\lambda_{11}) \right)^{\lambda_{1j}} \\ &= \frac{c_{00}}{\lambda_{00}} (\lambda_{01} + \lambda_{02} + \lambda_{03})^{\lambda_{01}} \times \\ &= \frac{c_{01}}{\lambda_{01}} (\lambda_{01} + \lambda_{02} + \lambda_{03})^{\lambda_{02}} \times \\ &= \frac{c_{02}}{\lambda_{02}} (\lambda_{01} + \lambda_{02} + \lambda_{03})^{\lambda_{03}} \times \\ &= \frac{c_{03}}{\lambda_{03}} (\lambda_{01} + \lambda_{02} + \lambda_{03})^{\lambda_{03}} \left( \frac{c_{11}}{\lambda_{11}} \lambda_{11} \right)^{\lambda_{11}} \end{aligned}$$

subject to the constraints are.

$$\sum_{j=1}^{N_0} \lambda_{0j} = 1 \rightarrow \text{Normality constraint. Constant}$$

$$\Rightarrow \lambda_{01} + \lambda_{02} + \lambda_{03} = 1$$

and orthogonal constraints are

$$\sum_{k=0}^M \sum_{j=1}^{N_k} a_{kij} \lambda_{kj} = 0, \quad i = 1, 2, \dots, n$$

$$\Rightarrow \sum_{k=0}^M \sum_{j=1}^{N_k} a_{kij} \lambda_{kj} = 0$$

$$\sum_{j=1}^{N_0} a_{0ij} \lambda_{0j} = 0$$

$$\sum_{j=1}^{N_0} a_{0ij} \lambda_{0j} + \sum_{j=1}^{N_1} a_{1ij} \lambda_{1j} = 0$$

$$\sum_{j=1}^3 a_{0ij} \lambda_{0j} + \sum_{j=1}^1 a_{1ij} \lambda_{1j} = 0$$

$$a_{011} \lambda_{01} + a_{012} \lambda_{02} + a_{013} \lambda_{03} + a_{111} \lambda_{11} = 0$$

$$\Rightarrow a_{011} \lambda_{01} + a_{012} \lambda_{02} + a_{013} \lambda_{03} + a_{111} \lambda_{11} = 0$$

$$a_{021} \lambda_{01} + a_{022} \lambda_{02} + a_{023} \lambda_{03} + a_{121} \lambda_{11} = 0$$

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$$a_{031}d_{01} + a_{032}d_{02} + a_{033}d_{03} + a_{131}d_{11} = 0$$

$$d_{0j} \geq 0, j=1,2,3, d_{11} \geq 0$$

in this problem

$$a_{011} = 1, a_{021} = 0, a_{031} = 0$$

$$a_{012} = 0, a_{022} = 1, a_{032} = \frac{1}{2}, a_{013} = 1, a_{023} = 1,$$

$$a_{033} = 0, a_{111} = -1, a_{121} = -1, a_{131} = -1.$$

So, problem can be written as

$$V(d) = \int \frac{20}{d_1} (d_1 + d_{02} + d_{03}) d d_1$$

subject to

$$d_{01} + d_{02} + d_{03} = 1$$

$$d_{01} + d_{03} - d_{11} = 0$$

$$d_{02} + d_{03} - d_{11} = 0$$

$$d_{01} + d_{02} - d_{11} = 0$$

$$\Rightarrow d_{01}^* = \frac{1}{3}, d_{02}^* = \frac{1}{3}, d_{03}^* = \frac{1}{3}, d_{11}^* = \frac{2}{3}.$$

thus the maximum value of  $V$  or minimum value of  $x_0$  is given by

$$V^* = x_0^* = (60)^{1/3} (100)^{1/3} (240)^{1/3} (8)^{1/3} = 480 \text{ Bm}$$

Complementary Geometric Programming

Geometric programming to include any rational function of polynomial terms and called the method of complementary geometric programming.

At the complementary geometric programming problem be stated as follows

Minimize  $f_0(x)$  subject

$$R_k(x) \leq 1 \quad k=1, 2, \dots, m, \text{ where}$$

$$R_k(x) = \frac{A_k(x) \cdot B_k(x)}{C_k(x) \cdot D_k(x)}, \quad k=0, 1, 2, \dots, m$$

where  $A_k(x)$ ,  $B_k(x)$ ,  $C_k(x)$  and  $D_k(x)$  are polynomial and possibly some of them may be absent. To solve the problem stated, we introduce a new variable  $x_0 > 0$ , constrained to satisfy the relation  $x_0 \geq R_0(x)$

$$\text{i.e. } \frac{R_0(x)}{x_0} \leq 1 \text{ so the problem may be}$$

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restated as  
 Minimise  $x_0$   
 subject to  $\frac{A_k(x) - B_k(x)}{C_k(x) - D_k(x)} \leq 1, k = 1, 2, \dots, m$   
 where  $A_0(x) = B_0(x), C_0(x) = x_0 + B_0(x) = 0$   
 and  $D_0(x) = 0$ .

thus any complementary geometric programming  
 problem (C.G.P.) can be stated in the  
 standard form

minimise  $x_0$  subject to  
 $\frac{P_k(x)}{Q_k(x)} \leq 1, k = 1, 2, \dots, m \quad (P.71)$

$x = \begin{Bmatrix} x_0 \\ x_1 \\ \vdots \\ x_n \end{Bmatrix} > 0$ , where  $(P.72)$

here  $P_k(x)$  and  $Q_k(x)$  are polynomials of

$$P_k(x) = \sum_j c_{kj} \prod_{i=0}^n (x_i)^{a_{kij}} = \sum_j p_{kj}(x) \quad (P.73)$$

$$Q_k(x) = \sum_j d_{kj} \prod_{i=0}^n (x_i)^{b_{kij}} = \sum_j q_{kj}(x) \quad (P.74)$$

$$\begin{aligned} p_k(x) &= \sum_j (c_{kj}) (\gamma_j)^{a_{k1j}} \gamma_2^{a_{k2j}} \gamma_3^{a_{k3j}} \dots \\ &= \sum_j p_{kj}(x) \end{aligned}$$

$$= p_{k1}(x) + p_{k2}(x) + p_{k3}(x) + \dots$$

$$= c_{k1}(\gamma_1)^{a_{k11}} (\gamma_2)^{a_{k21}} (\gamma_3)^{a_{k31}} \dots \\ + c_{k2}(\gamma_1)^{a_{k12}} (\gamma_2)^{a_{k22}} (\gamma_3)^{a_{k32}} \dots$$

$$Q_k(x) = q_{k1}(x) + q_{k2}(x) + q_{k3}(x) \dots$$

$$Q_1(x) = q_{11}(x) + q_{12}(x) + q_{13}(x) \dots$$

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7.6 Minimize  $x_1$  subject to

$$-4x_1^2 + 4x_2 \leq 1$$

$$x_1 + x_2 \geq 1, \quad x_1 > 0, x_2 > 0$$

Ans. (L.P.F.) can be stated as

Minimize  $x_1$  subject to

$$x_2 \leq 1 + 4x_1^2$$

$$\frac{x_2}{1 + 4x_1^2} \leq 1$$

$$\frac{1}{x_1 + x_2} \leq 1$$

$$\frac{x_1 + x_2}{x_1 x_2} \leq 1$$

$$\frac{1}{x_1} + \frac{1}{x_2} \leq 1$$

$$\frac{1}{x_1} \leq \frac{1}{x_2}$$

$$\frac{1}{1 + x_1/x_2} \leq 1$$

$$1 + x_1/x_2 \geq 1$$

$$x_1/x_2 \geq 0$$

instead of this  
 Point  $x^* = (1, 1)$   
 (Active Point)

(12) Piecewise  $f(x_1, x_2) = 2x_1^2 + 3x_2^2$  (Convexity check)  
 Iterative:  $\downarrow$ ,  $S_1 = \left\{ \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right\}$ ,  $x_1 = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$   
 $f_1 = f(x_1) = f(1, 1) = 2 \cdot 1^2 + 3 \cdot 1^2 = 5$ ,  $\ell = 0.01$   
 $f^* = f(x_1 + \ell S_1) = f(1+1, 1) = f(2, 1) = 2$   
 $= 2(1+1)^2 + 1 = 6.0002$   
 Next  $x_2 > x_1$   
 $f^* = f\{x_1 - \ell S_1\} = f\{1 - 1, 1\} = f(0, 1) = 3$   
 $= f(-1, 1) = 5.9998 < 5$   
 $\Rightarrow S_1$  will be the correct direction  $\Rightarrow$   $f^* = 3$   
 For optimum, let  $\nabla f(x_1)$  be minimum  
 $f(x_1) = f(x_1 - \ell_1 S_1) = \frac{1}{2} \ell_1^2 \cdot d(1 - \ell_1, 2)$   
 $= \frac{1}{2} (1 - \ell_1)^2 + 2 = 2(1 + \ell_1^2 - 2\ell_1) + 4$   
 $\frac{d}{d\ell_1} = 2(2\ell_1 - 2) = 0 \Rightarrow \ell_1 = 1$   
 $\Rightarrow x_2 = x_1 - \ell_1 S_1 = (1, 1) - 1 \cdot \begin{pmatrix} 1 \\ 0 \end{pmatrix} = (0, 1)$   
 $\Rightarrow d(x_1) = 0 + 3 = 4$   
 and iteration,  $S_2 = \left\{ \begin{pmatrix} 1 \\ 1 \end{pmatrix} \right\}$   
 $f_2^* = f(x_2 + \ell S_2) = f\left\{ \begin{pmatrix} 0 \\ 1 \end{pmatrix} + 0.01 \begin{pmatrix} 1 \\ 1 \end{pmatrix} \right\} = f(0, 2.01)$   
 $= 0 + 4.0401 \Rightarrow f_2^* > f_1^*$   
 $f_2^* = d(x_2 - 0.01 S_2) = f(0, 0.99) = 2.9601$   
 $\Rightarrow \frac{d}{d\ell_2} = 4 < 4 = d_2$