

CHAPTER 2

Markov chains

2.1. Examples

EXAMPLE 2.1.1 (Markov chain with two states). Consider a phone which can be in two states: “free” = 0 and “busy” = 1. The set of the states of the phone is

$$E = \{0, 1\}.$$

We assume that the phone can randomly change its state in time (which is assumed to be discrete) according to the following rules.

1. If at some time n the phone is free, then at time $n + 1$ it becomes busy with probability p or it stays free with probability $1 - p$.
2. If at some time n the phone is busy, then at time $n + 1$ it becomes free with probability q or it stays busy with probability $1 - q$.

Denote by X_n the state of the phone at time $n = 0, 1, \dots$. Thus, $X_n : \Omega \rightarrow \{0, 1\}$ is a random variable and our assumptions can be written as follows:

$$\begin{aligned} p_{00} &:= \mathbb{P}[X_{n+1} = 0 | X_n = 0] = 1 - p, & p_{01} &:= \mathbb{P}[X_{n+1} = 1 | X_n = 0] = p, \\ p_{10} &:= \mathbb{P}[X_{n+1} = 0 | X_n = 1] = q, & p_{11} &:= \mathbb{P}[X_{n+1} = 1 | X_n = 1] = 1 - q. \end{aligned}$$

We can write these probabilities in form of a *transition matrix*

$$P = \begin{pmatrix} 1 - p & p \\ q & 1 - q \end{pmatrix}.$$

Additionally, we will make the following assumption which is called the *Markov property*: Given that at some time n the phone is in state $i \in \{0, 1\}$, the behavior of the phone after time n does not depend on the way the phone reached state i in the past.

PROBLEM 2.1.2. Suppose that at time 0 the phone was free. What is the probability that the phone will be free at times 1, 2 and then becomes busy at time 3?

SOLUTION. This probability can be computed as follows:

$$\mathbb{P}[X_1 = X_2 = 0, X_3 = 1] = p_{00} \cdot p_{00} \cdot p_{01} = (1 - p)^2 p.$$

PROBLEM 2.1.3. Suppose that the phone was free at time 0. What is the probability that it will be busy at time 3?

SOLUTION. We have to compute $\mathbb{P}[X_3 = 1]$. We know the values $X_0 = 0$ and $X_3 = 1$, but the values of X_1 and X_2 may be arbitrary. We have the following possibilities:

- (1) $X_0 = 0, X_1 = 0, X_2 = 0, X_3 = 1$. Probability: $(1 - p) \cdot (1 - p) \cdot p$.
- (2) $X_0 = 0, X_1 = 0, X_2 = 1, X_3 = 1$. Probability: $(1 - p) \cdot p \cdot (1 - q)$.

- (3) $X_0 = 0, X_1 = 1, X_2 = 0, X_3 = 1$. Probability: $p \cdot q \cdot p$.
(4) $X_0 = 0, X_1 = 1, X_2 = 1, X_3 = 1$. Probability: $p \cdot (1 - q) \cdot (1 - q)$.

The probability we look for is the sum of these 4 probabilities:

$$\mathbb{P}[X_3 = 1] = (1 - p)^2 p + (1 - p)(1 - q)p + p^2 q + p(1 - q)^2.$$

EXAMPLE 2.1.4 (Gambler's ruin). At each unit of time a gambler plays a game in which he can either win 1€ (which happens with probability p) or he can lose 1€ (which happens with probability $1 - p$). Let X_n be the capital of the gambler at time n . Let us agree that if at some time n the gambler has no money (meaning that $X_n = 0$), then he stops to play (meaning that $X_n = X_{n+1} = \dots = 0$). We can view this process as a Markov chain on the state space $E = \{0, 1, 2, \dots\}$ with transition matrix

$$P = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & \dots \\ 1 - p & 0 & p & 0 & 0 & \dots \\ 0 & 1 - p & 0 & p & 0 & \dots \\ 0 & 0 & 1 - p & 0 & p & \dots \\ \dots & \dots & \dots & \dots & \dots & \dots \end{pmatrix}.$$

2.2. Definition of Markov chains

Let us consider some system. Assume that the system can be in some states and that the system can change its state in time. The set of all states of the system will be denoted by E and called the *state space* of the Markov chain. We always assume that the state space E is a finite or countable set. Usually, we will denote the states so that $E = \{1, \dots, N\}$, $E = \mathbb{N}$, or $E = \mathbb{Z}$.

Assume that if at some time the system is in state $i \in E$, then in the next moment of time it can switch to state $j \in E$ with probability p_{ij} . We will call p_{ij} the *transition probability* from state i to state j . Clearly, the transition probabilities should be such that

- (1) $p_{ij} \geq 0$ for all $i, j \in E$.
(2) $\sum_{j \in E} p_{ij} = 1$ for all $i \in E$.

We will write the transition probabilities in form of a *transition matrix*

$$P = (p_{ij})_{i, j \in E}.$$

The rows and the columns of this matrix are indexed by the set E . The element in the i -th row and j -th column is the transition probability p_{ij} . The elements of the matrix P are non-negative and the sum of elements in any row is equal to 1. Such matrices are called *stochastic*.

DEFINITION 2.2.1. A *Markov chain* with state space E and transition matrix P is a stochastic process $\{X_n : n \in \mathbb{N}_0\}$ taking values in E such that for every $n \in \mathbb{N}_0$ and every states $i_0, i_1, \dots, i_{n-1}, i, j$ we have

$$(2.2.1) \quad \mathbb{P}[X_{n+1} = j | X_n = i] = \mathbb{P}[X_{n+1} = j | X_0 = i_0, X_1 = i_1, \dots, X_{n-1} = i_{n-1}, X_n = i] \\ = p_{ij},$$

provided that $\mathbb{P}[X_0 = i_0, \dots, X_{n-1} = i_{n-1}, X_n = i] \neq 0$ (which ensures that the conditional probabilities are well-defined).

Condition (2.2.1) is called the *Markov property*.

In the above definition it is not specified at which state the Markov chain starts at time 0. In fact, the initial state can be in general arbitrary and we call the probabilities

$$(2.2.2) \quad \alpha_i := \mathbb{P}[X_0 = i], \quad i \in E,$$

the *initial probabilities*. We will write the initial probabilities in form of a row vector $\alpha = (\alpha_i)_{i \in E}$. This vector should be such that $\alpha_i \geq 0$ for all $i \in E$ and $\sum_{i \in E} \alpha_i = 1$.

THEOREM 2.2.2. *For all $n \in \mathbb{N}_0$ and for all $i_0, \dots, i_n \in E$ it holds that*

$$(2.2.3) \quad \mathbb{P}[X_0 = i_0, X_1 = i_1, \dots, X_n = i_n] = \alpha_{i_0} p_{i_0 i_1} p_{i_1 i_2} \dots p_{i_{n-1} i_n}.$$

PROOF. We use the induction over n . The induction basis is the case $n = 0$. We have $\mathbb{P}[X_0 = i_0] = \alpha_{i_0}$ by the definition of initial probabilities, see (2.2.2). Hence, Equation (2.2.3) holds for $n = 0$.

Induction assumption: Assume that (2.2.3) holds for some n . We prove that (2.2.3) holds with n replaced by $n + 1$. Consider the event $A = \{X_0 = i_0, X_1 = i_1, \dots, X_n = i_n\}$. By the induction assumption,

$$\mathbb{P}[A] = \alpha_{i_0} p_{i_0 i_1} p_{i_1 i_2} \dots p_{i_{n-1} i_n}.$$

By the Markov property,

$$\mathbb{P}[X_{n+1} = i_{n+1} | A] = p_{i_n i_{n+1}}.$$

It follows that

$$\begin{aligned} \mathbb{P}[X_0 = i_0, X_1 = i_1, \dots, X_n = i_n, X_{n+1} = i_{n+1}] &= \mathbb{P}[X_{n+1} = i_{n+1} | A] \cdot \mathbb{P}[A] \\ &= p_{i_n i_{n+1}} \cdot \alpha_{i_0} p_{i_0 i_1} p_{i_1 i_2} \dots p_{i_{n-1} i_n} \\ &= \alpha_{i_0} p_{i_0 i_1} p_{i_1 i_2} \dots p_{i_{n-1} i_n} p_{i_n i_{n+1}}. \end{aligned}$$

This completes the induction. □

REMARK 2.2.3. If $\mathbb{P}[A] = 0$, then in the above proof we cannot use the Markov property. However, in the case $\mathbb{P}[A] = 0$ both sides of (2.2.3) are equal to 0 and (2.2.3) is trivially satisfied.

THEOREM 2.2.4. *For every $n \in \mathbb{N}$ and every state $i_n \in E$ we have*

$$\mathbb{P}[X_n = i_n] = \sum_{i_0, \dots, i_{n-1} \in E} \alpha_{i_0} p_{i_0 i_1} \dots p_{i_{n-1} i_n}.$$

PROOF. We have

$$\begin{aligned} \mathbb{P}[X_n = i_n] &= \sum_{i_0, \dots, i_{n-1} \in E} \mathbb{P}[X_0 = i_0, X_1 = i_1, \dots, X_n = i_n] \\ &= \sum_{i_0, \dots, i_{n-1} \in E} \alpha_{i_0} p_{i_0 i_1} \dots p_{i_{n-1} i_n}, \end{aligned}$$

where the last step is by Theorem 2.2.2. □

2.3. n -step transition probabilities

NOTATION 2.3.1. If we want to indicate that the Markov chain starts at state $i \in E$ at time 0, we will write \mathbb{P}_i instead of \mathbb{P} .

DEFINITION 2.3.2. The n -step transition probabilities of a Markov chain are defined as

$$p_{ij}^{(n)} := \mathbb{P}_i[X_n = j].$$

We will write these probabilities in form of the n -step transition matrix $P^{(n)} = (p_{ij}^{(n)})_{i,j \in E}$.

By Theorem 2.2.4 we have the formula

$$p_{ij}^{(n)} = \sum_{i_1, \dots, i_{n-1} \in E} p_{ii_1} p_{i_1 i_2} \cdots p_{i_{n-1} j}.$$

The next theorem is crucial. It states that the n -step transition matrix $P^{(n)}$ can be computed as the n -th power of the transition matrix P .

THEOREM 2.3.3. We have $P^{(n)} = P^n = P \cdot \dots \cdot P$.

PROOF. We use induction over n . For $n = 1$ we have $p_{ij}^{(1)} = p_{ij}$ and hence, $P^{(1)} = P$. Thus, the statement of the theorem is true for $n = 1$.

Let us now assume that we already proved that $P^{(n)} = P^n$ for some $n \in \mathbb{N}$. We compute $P^{(n+1)}$. By the formula of total probability, we have

$$p_{ij}^{(n+1)} = \mathbb{P}_i[X_{n+1} = j] = \sum_{k \in E} \mathbb{P}_i[X_n = k] \mathbb{P}[X_{n+1} = j | X_n = k] = \sum_{k \in E} p_{ik}^{(n)} p_{kj}.$$

On the right hand-side we have the scalar product of the i -th row of the matrix $P^{(n)}$ and the j -th column of the matrix P . By definition of the matrix multiplication, this scalar product is exactly the entry of the matrix product $P^{(n)}P$ which is located in the i -th row and j -th column. We thus have the equality of matrices

$$P^{(n+1)} = P^{(n)}P.$$

But now we can apply the induction assumption $P^{(n)} = P^n$ to obtain

$$P^{(n+1)} = P^{(n)}P = P^n \cdot P = P^{n+1}.$$

This completes the induction. □

In the next theorem we consider a Markov chain with initial distribution $\alpha = (\alpha_i)_{i \in E}$ and transition matrix P . Let $\alpha^{(n)} = (\alpha_j^{(n)})_{j \in E}$ be the distribution of the position of this chain at time n , that is

$$\alpha_j^{(n)} = \mathbb{P}[X_n = j].$$

We write both $\alpha^{(n)}$ and α as *row* vectors. The next theorem states that we can compute $\alpha^{(n)}$ by taking α and multiplying it by the n -step transition matrix $P^{(n)} = P^n$ from the right.

THEOREM 2.3.4. We have

$$\alpha^{(n)} = \alpha P^n.$$

PROOF. By the formula of the total probability

$$\alpha_j^{(n)} = \mathbb{P}[X_n = j] = \sum_{i \in E} \alpha_i \mathbb{P}_i[X_n = j] = \sum_{i \in E} \alpha_i p_{ij}^{(n)}.$$

On the right-hand side we have the scalar product of the row α with the j -th column of $P^{(n)} = P^n$. By definition of matrix multiplication, this means that $\alpha^{(n)} = \alpha P^n$. \square

2.4. Invariant measures

Consider a Markov chain on state space E with transition matrix P . Let $\lambda : E \rightarrow \mathbb{R}$ be a function. To every state $i \in E$ the function assigns some value which will be denoted by $\lambda_i := \lambda(i)$. Also, it will be convenient to write the function λ as a row vector $\lambda = (\lambda_i)_{i \in E}$.

DEFINITION 2.4.1. A function $\lambda : E \rightarrow \mathbb{R}$ is called a *measure* on E if $\lambda_i \geq 0$ for all $i \in E$.

DEFINITION 2.4.2. A function $\lambda : E \rightarrow \mathbb{R}$ is called a *probability measure* on E if $\lambda_i \geq 0$ for all $i \in E$ and

$$\sum_{i \in E} \lambda_i = 1.$$

DEFINITION 2.4.3. A measure λ is called *invariant* if $\lambda P = \lambda$. That is, for every state $j \in E$ it should hold that

$$\lambda_j = \sum_{i \in E} \lambda_i p_{ij}.$$

REMARK 2.4.4. If the initial distribution α of a Markov chain is invariant, that is $\alpha P = \alpha$, then for every $n \in \mathbb{N}$ we have $\alpha P^n = \alpha$ which means that at every time n the position of the Markov chain has the same distribution as at time 0:

$$X_0 \stackrel{d}{=} X_1 \stackrel{d}{=} X_2 \stackrel{d}{=} \dots$$

EXAMPLE 2.4.5. Let us compute the invariant distribution for the Markov chain from Example 2.1.1. The transition matrix is

$$P = \begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix}.$$

The equation $\lambda P = \lambda$ for the invariant probability measure takes the following form:

$$(\lambda_0, \lambda_1) \begin{pmatrix} 1-p & p \\ q & 1-q \end{pmatrix} = (\lambda_0, \lambda_1).$$

Multiplying the matrices we obtain the following two equations:

$$\begin{aligned} \lambda_0(1-p) + \lambda_1 q &= \lambda_0, \\ \lambda_0 p + \lambda_1(1-q) &= \lambda_1. \end{aligned}$$

From the first equation we obtain that $\lambda_1 q = \lambda_0 p$. Solving the second equation we obtain the same relation which means that the second equation does not contain any information not contained in the first equation. However, since we are looking for invariant *probability* measures, we have an additional equation

$$\lambda_0 + \lambda_1 = 1.$$

Solving this equation together with $\lambda_1 q = \lambda_0 p$ we obtain the following result:

$$\lambda_0 = \frac{q}{p+q}, \quad \lambda_1 = \frac{p}{p+q}.$$

PROBLEM 2.4.6. Consider the phone from Example 2.1.1. Let the phone be free at time 0. What is (approximately) the probability that it is free at time $n = 1000$?

SOLUTION. The number $n = 1000$ is large. For this reason it seems plausible that the probability that the phone is free (busy) at time $n = 1000$ should be approximately the same as the probability that it is free (busy) at time $n + 1 = 1001$. Denoting the initial distribution by $\alpha = (1, 0)$ and the distribution of the position of the chain at time n by $\alpha^{(n)} = \alpha P^n$ we thus must have

$$\alpha^{(n)} \approx \alpha^{(n+1)} = \alpha P^{n+1} = \alpha P^n \cdot P = \alpha^{(n)} P.$$

Recall that the equation for the invariant probability measure has the same form $\lambda = \lambda P$. It follows that $\alpha^{(n)}$ must be approximately the invariant probability measure:

$$\alpha^{(n)} \approx \lambda.$$

For the probability that the phone is free (busy) at time $n = 1000$ we therefore obtain the approximations

$$p_{00}^{(n)} \approx \lambda_0 = \frac{q}{p+q}, \quad p_{01}^{(n)} \approx \lambda_1 = \frac{p}{p+q}.$$

Similar considerations apply to the case when the phone is busy at time 0 leading to the approximations

$$p_{10}^{(n)} \approx \lambda_0 = \frac{q}{p+q}, \quad p_{11}^{(n)} \approx \lambda_1 = \frac{p}{p+q}.$$

Note that $p_{00}^{(n)} \approx p_{10}^{(n)}$ and $p_{01}^{(n)} \approx p_{11}^{(n)}$ which can be interpreted by saying that the Markov chain almost forgets its initial state after many steps. For the n -step transition matrix we therefore may conjecture that

$$\lim_{n \rightarrow \infty} P^n = \lim_{n \rightarrow \infty} \begin{pmatrix} p_{00}^{(n)} & p_{01}^{(n)} \\ p_{10}^{(n)} & p_{11}^{(n)} \end{pmatrix} = \begin{pmatrix} \lambda_0 & \lambda_1 \\ \lambda_0 & \lambda_1 \end{pmatrix}.$$

The above considerations are not rigorous. We will show below that if a general Markov chain satisfies appropriate conditions, then

- (1) The invariant probability measure λ exists and is unique.
- (2) For every states $i, j \in E$ we have $\lim_{n \rightarrow \infty} p_{ij}^{(n)} = \lambda_j$.

EXAMPLE 2.4.7 (Ehrenfest model). We consider a box which is divided into 2 parts. Consider N balls (molecules) which are located in this box and can move from one part to the other according to the following rules. Assume that at any moment of time one of the N balls is chosen at random (all balls having the same probability $1/N$ to be chosen). This ball moves to the other part. Then, the procedure is repeated. Let X_n be the number of balls at time n in Part 1. Then, X_n takes values in $E = \{0, 1, \dots, N\}$ which is our state space. The transition probabilities are given by

$$p_{0,1} = 1, \quad p_{N,N-1} = 1, \quad p_{i,i+1} = \frac{N-i}{N}, \quad p_{i,i-1} = \frac{i}{N}, \quad i = 1, \dots, N-1.$$

For the invariant probability measure we obtain the following system of equations

$$\lambda_0 = \frac{\lambda_1}{N}, \quad \lambda_N = \frac{\lambda_{N-1}}{N}, \quad \lambda_j = \frac{N-j+1}{N} \lambda_{j-1} + \frac{j+1}{N} \lambda_{j+1}, \quad j = 1, \dots, N-1.$$

Additionally, we have the equation $\lambda_0 + \dots + \lambda_N = 1$. This system of equations can be solved directly, but one can also guess the solution without doing computations. Namely, it seems plausible that after a large number of steps every ball will be with probability $1/2$ in Part 1 and with probability $1/2$ in Part 2. Hence, one can guess that the invariant probability measure is the binomial distribution with parameter $1/2$:

$$\lambda_j = \frac{1}{2^N} \binom{N}{j}.$$

One can check that this is indeed the unique invariant probability measure for this Markov chain.

EXAMPLE 2.4.8. Let X_0, X_1, \dots be independent and identically distributed random variables with values $1, \dots, N$ and corresponding probabilities

$$\mathbb{P}[X_n = i] = p_i, \quad p_1, \dots, p_N \geq 0, \quad \sum_{i=1}^N p_i = 1.$$

Then, X_0, X_1, \dots is a Markov chain and the transition matrix is

$$P = \begin{pmatrix} p_1 & \dots & p_N \\ \dots & \dots & \dots \\ p_1 & \dots & p_N \end{pmatrix}.$$

The invariant probability measure is given by $\lambda_1 = p_1, \dots, \lambda_N = p_N$.

2.5. Class structure and irreducibility

Consider a Markov chain on a state space E with transition matrix P .

DEFINITION 2.5.1. We say that state $i \in E$ *leads* to state $j \in E$ if there exists $n \in \mathbb{N}_0$ such that $p_{ij}^{(n)} \neq 0$. We use the notation $i \rightsquigarrow j$.

REMARK 2.5.2. By convention, $p_{ii}^{(0)} = 1$ and hence, every state leads to itself: $i \rightsquigarrow i$.

THEOREM 2.5.3. *For two states $i, j \in E$ with $i \neq j$, the following statements are equivalent:*

- (1) $i \rightsquigarrow j$.
- (2) $\mathbb{P}_i[\exists n \in \mathbb{N} : X_n = j] \neq 0$.
- (3) *There exist $n \in \mathbb{N}$ and states $i_1, \dots, i_{n-1} \in E$ such that $p_{ii_1} \dots p_{i_{n-1}j} > 0$.*

PROOF. We prove that Statements 1 and 2 are equivalent. We have the inequality

$$(2.5.1) \quad p_{ij}^{(n)} \leq \mathbb{P}_i[\exists n \in \mathbb{N} : X_n = j] \leq \sum_{n=1}^{\infty} \mathbb{P}_i[X_n = j] = \sum_{n=1}^{\infty} p_{ij}^{(n)}.$$

If Statement 1 holds, then for some $n \in \mathbb{N}$ we have $p_{ij}^{(n)} > 0$. Hence, by (2.5.1), we have $\mathbb{P}_i[\exists n \in \mathbb{N} : X_n = j] > 0$ and Statement 2 holds. If, conversely, Statement 2 holds, then

$\mathbb{P}_i[\exists n \in \mathbb{N} : X_n = j] > 0$. Hence, by (2.5.1), $\sum_{n=1}^{\infty} p_{ij}^{(n)} > 0$, which implies that at least one summand $p_{ij}^{(n)}$ must be strictly positive. This proves Statement 1.

We prove the equivalence of Statements 1 and 3. We have the formula

$$(2.5.2) \quad p_{ij}^{(n)} = \sum_{i_1, \dots, i_{n-1} \in E} p_{ii_1} \cdots p_{i_{n-1}j}.$$

If Statement 1 holds, then for some $n \in \mathbb{N}$ we have $p_{ij}^{(n)} > 0$ which implies that at least one summand on the right-hand side of (2.5.2) must be strictly positive. This implies Statement 3. If, conversely, Statement 3 holds, then the sum on the right-hand side of (2.5.2) is positive which implies that $p_{ij}^{(n)} > 0$. Hence, Statement 1 holds. \square

DEFINITION 2.5.4. States $i, j \in E$ *communicate* if $i \rightsquigarrow j$ and $j \rightsquigarrow i$. Notation: $i \longleftrightarrow j$.

THEOREM 2.5.5. $i \longleftrightarrow j$ is an equivalence relation, namely

- (1) $i \longleftrightarrow i$.
- (2) $i \longleftrightarrow j \iff j \longleftrightarrow i$.
- (3) $i \longleftrightarrow j, j \longleftrightarrow k \Rightarrow i \longleftrightarrow k$.

PROOF. Statements 1 and 2 follow from the definition. We prove Statement 3. If $i \longleftrightarrow j$ and $j \longleftrightarrow k$, then, in particular, $i \rightsquigarrow j$ and $j \rightsquigarrow k$. By Theorem 2.5.3, Statement 3, we can find $r \in \mathbb{N}$, $s \in \mathbb{N}$ and states $u_1, \dots, u_{r-1} \in E$ and $v_1, \dots, v_{s-1} \in E$ such that $p_{iu_1} p_{u_1 u_2} \cdots p_{u_{r-1} j} > 0$ and $p_{jv_1} p_{v_1 v_2} \cdots p_{v_{s-1} k} > 0$. Multiplying both inequalities, we get

$$p_{iu_1} p_{u_1 u_2} \cdots p_{u_{r-1} j} p_{jv_1} p_{v_1 v_2} \cdots p_{v_{s-1} k} > 0.$$

By Theorem 2.5.3, Statement 3, we have $i \rightsquigarrow k$. In a similar way one shows that $k \rightsquigarrow i$. \square

DEFINITION 2.5.6. The *communication class* of state $i \in E$ is the set $\{j \in E : i \longleftrightarrow j\}$. This set consists of all states j which communicate to i .

Since communication of states is an equivalence relation, the state space E can be decomposed into a disjoint union of communication classes. Any two communication classes either coincide completely or are disjoint sets.

DEFINITION 2.5.7. A Markov chain is *irreducible* if every two states communicate. Hence, an irreducible Markov chain consists of just one communication class.

DEFINITION 2.5.8. A communication class C is *open* if there exist a state $i \in C$ and a state $k \notin C$ such that $i \rightsquigarrow k$. Otherwise, a communication class is called *closed*.

If a Markov chain once arrived in a closed communication class, it will stay in this class forever.

EXERCISE 2.5.9. Show that a communication class C is open if and only if there exist a state $i \in C$ and a state $k \notin C$ such that $p_{ik} > 0$.

THEOREM 2.5.10. If the state space E is a finite set, then there exists at least one closed communication class.

PROOF. We use a proof by contradiction. Assume that there is no closed communication class. Hence, all communication classes are open. Take some state and let C_1 be the communication class of this state. Since C_1 is open, there is a path from C_1 to some other communication class $C_2 \neq C_1$. Since C_2 is open, we can go from C_2 to some other communication class $C_3 \neq C_2$, and so on. Note that in the sequence C_1, C_2, C_3, \dots all classes are different. Indeed, if for some $l < m$ we would have $C_l = C_m$ (a “cycle”), this would mean that there is a path starting from C_l , going to C_{l+1} and then to $C_m = C_l$. But this is a contradiction since then C_l and C_{l+1} should be a single communication class, and not two different classes, as in the construction. So, the classes C_1, C_2, \dots are different (in fact, disjoint) and each class contains at least one element. But this is a contradiction since E is a finite set. \square

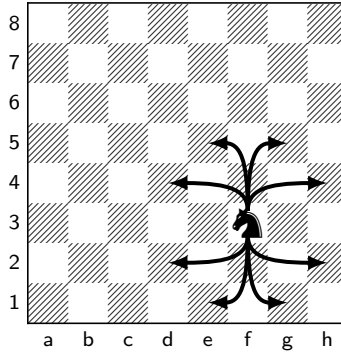
2.6. Aperiodicity

DEFINITION 2.6.1. The *period* of a state $i \in E$ is defined as

$$\gcd\{n \in \mathbb{N} : p_{ii}^{(n)} > 0\}.$$

Here, gcd states for the greatest common divisor. A state $i \in E$ is called *aperiodic* if its period is equal to 1. Otherwise, the state i is called *periodic*.

EXAMPLE 2.6.2. Consider a knight on a chessboard moving according to the usual chess rules in a random way. For concreteness, assume that at each moment of time all moves of the knight allowed by the chess rules are counted and then one of these moves is chosen, all moves being equiprobable.



This is a Markov chain on a state space consisting of 64 squares. Assume that at time 0 the knight is in square i . Since the knight changes the color of its square after every move, it cannot return to the original square in an odd number of steps. On the other hand, it can return to i in an even number of steps with non-zero probability (for example by going to some other square and then back, many times). So,

$$p_{ii}^{(2n+1)} = 0, \quad p_{ii}^{(2n)} > 0.$$

Hence, the period of any state in this Markov chain is 2.

EXAMPLE 2.6.3. Consider a Markov chain on a state space of two elements with transition matrix

$$P = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix},$$

We have

$$p_{ii}^{(2n+1)} = 0, \quad p_{ii}^{(2n)} = 1.$$

Hence, the period of any state in this Markov chain is 2.

EXERCISE 2.6.4. Show that in the Ehrenfest Markov chain (Example 2.4.7) every state is periodic with period 2.

LEMMA 2.6.5. *Let $i \in E$ be any state. The following conditions are equivalent:*

- (1) *State i is aperiodic.*
- (2) *There is $N \in \mathbb{N}$ such that for every natural number $n > N$ we have $p_{ii}^{(n)} > 0$.*

PROOF. If Statement 2 holds, then for some sufficiently large n we have $p_{ii}^{(n)} > 0$ and $p_{ii}^{(n+1)} > 0$. Since $\gcd(n, n+1) = 1$, the state i has period 1. Hence, Statement 1 holds.

Suppose, conversely, that Statement 1 holds. Then, we can find $n_1, \dots, n_r \in \mathbb{N}$ such that $\gcd\{n_1, \dots, n_r\} = 1$ and $p_{ii}^{(n_1)} > 0, \dots, p_{ii}^{(n_r)} > 0$. By a result from number theory, the condition $\gcd\{n_1, \dots, n_r\} = 1$ implies that there is $N \in \mathbb{N}$ such that we can represent any natural number $n > N$ in the form $n = l_1 n_1 + \dots + l_r n_r$ for suitable $l_1, \dots, l_r \in \mathbb{N}$. We obtain that

$$p_{ii}^{(l_1 n_1 + \dots + l_r n_r)} \geq (p_{ii}^{(n_1)})^{l_1} \cdot \dots \cdot (p_{ii}^{(n_r)})^{l_r} > 0.$$

This proves Statement 2.

LEMMA 2.6.6. *If state $i \in E$ is aperiodic and $i \rightsquigarrow j$, then j is also aperiodic.*

REMARK 2.6.7. We can express this by saying that aperiodicity is a *class property*: If some state in a communication class is aperiodic, then all states in this communication class are aperiodic. Similarly, if some state in a communication class is periodic, then all states in this communication class must be periodic. We can thus divide all communication classes into two categories: the aperiodic communication classes (consisting of only aperiodic states) and the periodic communication classes (consisting only of periodic states).

DEFINITION 2.6.8. An irreducible Markov chain is called aperiodic if some (and hence, all) states in this chain are aperiodic.

PROOF OF LEMMA 2.6.6. From $i \rightsquigarrow j$ it follows that $i \rightsquigarrow j$ and $j \rightsquigarrow i$. Hence, we can find $r, s \in \mathbb{N}_0$ such that $p_{ji}^{(r)} > 0$ and $p_{ij}^{(s)} > 0$. Since the state i is aperiodic, by Lemma 2.6.5 we can find $N \in \mathbb{N}$ such that for all $n > N$, we have $p_{ii}^{(n)} > 0$ and hence,

$$p_{jj}^{(n+r+s)} \geq p_{ji}^{(r)} \cdot p_{ii}^{(n)} \cdot p_{ij}^{(s)} > 0.$$

It follows that $p_{jj}^{(k)} > 0$ for all $k := n + r + s > N + r + s$. By Lemma 2.6.5, this implies that j is aperiodic. \square

2.7. Recurrence and transience

Consider a Markov chain $\{X_n : n \in \mathbb{N}_0\}$ on state space E with transition matrix P .

DEFINITION 2.7.1. A state $i \in E$ is called *recurrent* if

$$\mathbb{P}_i[X_n = i \text{ for infinitely many } n] = 1.$$

DEFINITION 2.7.2. A state $i \in E$ is called *transient* if

$$\mathbb{P}_i[X_n = i \text{ for infinitely many } n] = 0.$$

A recurrent state has the property that a Markov chain starting at this state returns to this state infinitely often, with probability 1. A transient state has the property that a Markov chain starting at this state returns to this state only finitely often, with probability 1.

The next theorem is a characterization of recurrent/transient states.

THEOREM 2.7.3. Let $i \in E$ be a state. Denote by f_i the probability that a Markov chain which starts at i returns to i at least once, that is

$$f_i = \mathbb{P}_i[\exists n \in \mathbb{N} : X_n = i].$$

Then,

- (1) The state i is recurrent if and only if $f_i = 1$.
- (2) The state i is transient if and only if $f_i < 1$.

COROLLARY 2.7.4. Every state is either recurrent or transient.

PROOF. For $k \in \mathbb{N}$ consider the random event

$$B_k = \{X_n = i \text{ for at least } k \text{ different values of } n \in \mathbb{N}\}.$$

Then, $\mathbb{P}_i[B_k] = f_i^k$. Also, $B_1 \supset B_2 \supset \dots$. It follows that

$$\mathbb{P}_i[X_n = i \text{ for infinitely many } n] = \mathbb{P}_i[\cap_{k=1}^{\infty} B_k] = \lim_{k \rightarrow \infty} \mathbb{P}_i[B_k] = \lim_{k \rightarrow \infty} f_i^k = \begin{cases} 1, & \text{if } f_i = 1, \\ 0, & \text{if } f_i < 1. \end{cases}$$

It follows that state i is recurrent if $f_i = 1$ and transient if $f_i < 1$. □

Here is one more characterization of recurrence and transience.

THEOREM 2.7.5. Let $i \in E$ be a state. Recall that $p_{ii}^{(n)} = \mathbb{P}_i[X_n = i]$ denotes the probability that a Markov chain which started at state i visits state i at time n . Then,

- (1) The state i is recurrent if and only if $\sum_{n=1}^{\infty} p_{ii}^{(n)} = \infty$.
- (2) The state i is transient if and only if $\sum_{n=1}^{\infty} p_{ii}^{(n)} < \infty$.

PROOF. Let the Markov chain start at state i . Consider the random variable

$$V_i := \sum_{n=1}^{\infty} \mathbb{1}_{\{X_n = i\}}$$

which counts the number of returns of the Markov chain to state i . Note that the random variable V_i can take the value $+\infty$. Then,

$$\mathbb{P}_i[V_i \geq k] = \mathbb{P}[B_k] = f_i^k, \quad k \in \mathbb{N}.$$

Thus, the expectation of V_i can be computed as follows:

$$(2.7.1) \quad \mathbb{E}_i[V_i] = \sum_{k=1}^{\infty} \mathbb{P}_i[V_i \geq k] = \sum_{k=1}^{\infty} f_i^k.$$

On the other hand,

$$(2.7.2) \quad \mathbb{E}_i[V_i] = \mathbb{E}_i \sum_{n=1}^{\infty} \mathbb{1}_{\{X_n=i\}} = \sum_{n=1}^{\infty} \mathbb{E}_i \mathbb{1}_{\{X_n=i\}} = \sum_{n=1}^{\infty} p_{ii}^{(n)}.$$

CASE 1. Assume that state i is recurrent. Then, $f_i = 1$ by Theorem 2.7.3. It follows that $\mathbb{E}_i[V_i] = \infty$ by (2.7.1). (In fact, $\mathbb{P}_i[V_i = +\infty] = 1$ since $\mathbb{P}[V_i \geq k] = 1$ for every $k \in \mathbb{N}$). Hence, $\sum_{n=1}^{\infty} p_{ii}^{(n)} = \infty$ by (2.7.2)

CASE 2. Assume that state i is transient. Then, $f_i < 1$ by Theorem 2.7.3. Thus, $\mathbb{E}_i V_i < \infty$ by (2.7.1) and hence, $\sum_{n=1}^{\infty} p_{ii}^{(n)} < \infty$ by (2.7.2). \square

The next theorem shows that recurrence and transience are class properties: If some state in a communicating class is recurrent (resp. transient), then all states in this class are recurrent (resp. transient).

THEOREM 2.7.6.

1. If $i \in E$ be a recurrent state and $j \rightsquigarrow i$, then j is also recurrent.
2. If $i \in E$ be a transient state and $j \rightsquigarrow i$, then j is also transient.

PROOF. It suffices to prove Part 2. Let i be a transient state and let $j \rightsquigarrow i$. It follows that there exist $s, r \in \mathbb{N}_0$ with $p_{ij}^{(s)} > 0$ and $p_{ji}^{(r)} > 0$. For all $n \in \mathbb{N}$ it holds that

$$p_{ii}^{(n+r+s)} \geq p_{ij}^{(s)} p_{jj}^{(n)} p_{ji}^{(r)}.$$

Therefore,

$$\sum_{n=1}^{\infty} p_{jj}^{(n)} \leq \frac{1}{p_{ij}^{(s)} p_{ji}^{(r)}} \sum_{n=1}^{\infty} p_{ii}^{(n+r+s)} \leq \frac{1}{p_{ij}^{(s)} p_{ji}^{(r)}} \sum_{n=1}^{\infty} p_{ii}^{(n)} < \infty,$$

where the last step holds because i is transient. It follows that state j is also transient. \square

Theorem 2.7.6 allows us to introduce the following definitions.

DEFINITION 2.7.7. A communicating class is called recurrent if at least one (equivalently, every) state in this class is recurrent. A communicating class is transient if at least one (equivalently, every) state in this class is transient.

DEFINITION 2.7.8. An irreducible Markov chain is called recurrent if at least one (equivalently, every) state in this chain is recurrent. An irreducible Markov chain is called transient if at least one (equivalently, every) state in this chain is transient.

The next theorem states that it is impossible to leave a recurrent class.

THEOREM 2.7.9. *Every recurrent communicating class is closed.*

PROOF. Let C be a non-closed class. We need to show that it is not recurrent. Since C is not closed, there exist states i, j so that $i \in C$, $j \notin C$ and $i \rightsquigarrow j$. This means that there exists $m \in \mathbb{N}$ so that $p_{ij}^{(m)} = \mathbb{P}_i[X_m = j] > 0$. If the event $\{X_m = j\}$ occurs, then after

time m the chain cannot return to state i because otherwise i and j would be in the same communicating class. It follows that

$$\mathbb{P}_i[\{X_m = j\} \cap \{X_n = i \text{ for infinitely many } n\}] = 0.$$

This implies that

$$\mathbb{P}_i[X_n = i \text{ for infinitely many } n] < 1.$$

Therefore, state i is not recurrent. \square

If some communicating class contains only finitely states and the chain cannot leave this class, then it looks very plausible that the chain which started in some state of this class will return to this state infinitely often (and, in fact, will visit any state of this class infinitely often), with probability 1. This is stated in the next theorem.

THEOREM 2.7.10. *Every finite closed communicating class is recurrent.*

PROOF. Let C be a closed communicating class with finitely many elements. Take some state $i \in C$. A chain starting in i stays in C forever and since C is finite, there must be at least one state $j \in C$ which is visited infinitely often with positive probability:

$$\mathbb{P}_i[X_n = j \text{ for infinitely many } n \in \mathbb{N}] > 0.$$

At the moment it is not clear whether we can take $i = j$. But since i and j are in the same communicating class, there exists $m \in \mathbb{N}_0$ so that $p_{ji}^{(m)} > 0$. From the inequality

$$\mathbb{P}_j[X_n = j \text{ for infinitely many } n] > p_{ji}^{(m)} \cdot \mathbb{P}_i[X_n = j \text{ for infinitely many } n] > 0$$

it follows that state j is recurrent. The class C is then recurrent because it contains at least one recurrent state, namely j . \square

So, in a Markov chain with finitely many states we have the following equivalencies

- (1) A communicating class is recurrent if and only if it is closed.
- (2) A communicating class is transient if and only if it is not closed.

LEMMA 2.7.11. *Consider an irreducible, recurrent Markov chain with an arbitrary initial distribution α . Then, for every state $j \in E$ the number of visits of the chain to j is infinite with probability 1.*

PROOF. Exercise. \square

2.8. Recurrence and transience of random walks

EXAMPLE 2.8.1. A *simple random walk on \mathbb{Z}* is a Markov chain with state space $E = \mathbb{Z}$ and transition probabilities

$$p_{i,i+1} = p, \quad p_{i,i-1} = 1 - p, \quad i \in \mathbb{Z}.$$

So, from every state the random walk goes one step to the right with probability p , or one step to the left with probability $1 - p$; see Figure 1. Here, $p \in [0, 1]$ is a parameter.

THEOREM 2.8.2. *If $p = \frac{1}{2}$, then any state of the simple random walk is recurrent. If $p \neq \frac{1}{2}$, then any state is transient.*

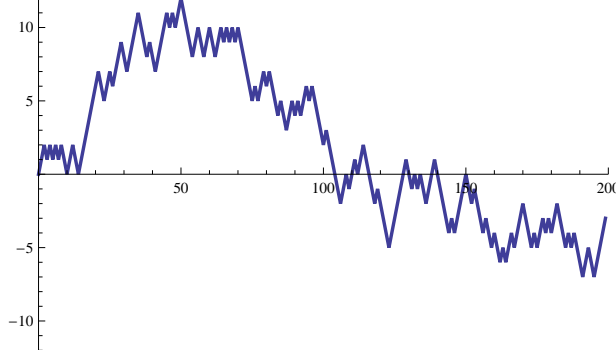


FIGURE 1. Sample path of a simple random walk on \mathbb{Z} with $p = \frac{1}{2}$. The figure shows 200 steps of the walk.

PROOF. By translation invariance, we can restrict our attention to state 0. We can represent our Markov chain as $X_n = \xi_1 + \dots + \xi_n$, where ξ_1, ξ_2, \dots are independent and identically distributed random variables with Bernoulli distribution:

$$\mathbb{P}[\xi_k = 1] = p, \quad \mathbb{P}[\xi_k = -1] = 1 - p.$$

CASE 1. Let $p \neq \frac{1}{2}$. Then, $\mathbb{E}\xi_k = p - (1 - p) = 2p - 1 \neq 0$. By the strong law of large numbers,

$$\lim_{n \rightarrow \infty} \frac{1}{n} X_n = \lim_{n \rightarrow \infty} \frac{\xi_1 + \dots + \xi_n}{n} = \mathbb{E}\xi_1 \neq 0 \quad \text{a.s.}$$

In the case $p > \frac{1}{2}$ we have $\mathbb{E}\xi_1 > 0$ and hence, $\lim_{n \rightarrow \infty} X_n = +\infty$ a.s. In the case $p < \frac{1}{2}$ we have $\mathbb{E}\xi_1 < 0$ and hence, $\lim_{n \rightarrow \infty} X_n = -\infty$ a.s. In both cases it follows that

$$\mathbb{P}[X_n = 0 \text{ for infinitely many } n] = 0.$$

Hence, state 0 is transient.

CASE 2. Let $p = \frac{1}{2}$. In this case, $\mathbb{E}\xi_k = 0$ and the argument of Case 1 does not work. We will use Theorem 2.7.5. The n -step transition probability from 0 to 0 is given by

$$p_{00}^{(n)} = \begin{cases} 0, & \text{if } n = 2k + 1 \text{ odd,} \\ \frac{1}{2^{2k}} \binom{2k}{k}, & \text{if } n = 2k \text{ even.} \end{cases}$$

The Stirling formula $n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$, as $n \rightarrow \infty$, yields that

$$p_{00}^{(2k)} \sim \frac{1}{\sqrt{\pi k}}, \quad \text{as } k \rightarrow \infty.$$

Since the series $\sum_{k=1}^{\infty} \frac{1}{\sqrt{k}}$ diverges, it follows that $\sum_{n=1}^{\infty} p_{00}^{(n)} = \sum_{k=1}^{\infty} p_{00}^{(2k)} = \infty$. By Theorem 2.7.5, this implies that 0 is a recurrent state. \square

EXAMPLE 2.8.3. The *simple, symmetric random walk on \mathbb{Z}^d* is a Markov chain defined as follows. The state space is the d -dimensional lattice

$$\mathbb{Z}^d = \{(n_1, \dots, n_d) : n_1, \dots, n_d \in \mathbb{Z}\}.$$

Let e_1, \dots, e_d be the standard basis of \mathbb{R}^d , that is

$$e_1 = (1, 0, 0, \dots, 0), \quad e_2 = (0, 1, 0, \dots, 0), \quad e_3 = (0, 0, 1, \dots, 0), \quad \dots, \quad e_d = (0, 0, 0, \dots, 1).$$

Let ξ_1, ξ_2, \dots be independent and identically distributed d -dimensional random vectors such that

$$\mathbb{P}[\xi_i = e_k] = \mathbb{P}[\xi_i = -e_k] = \frac{1}{2d}, \quad k = 1, \dots, d, \quad i \in \mathbb{N}.$$

Define $S_n = \xi_1 + \dots + \xi_n$, $n \in \mathbb{N}$, and $S_0 = 0$. The sequence S_0, S_1, S_2, \dots is called the simple symmetric random walk on \mathbb{Z}^d . It is a Markov chain with transition probabilities

$$p_{i, i+e_1} = p_{i, i-e_1} = \dots = p_{i, i+e_d} = p_{i, i-e_d} = \frac{1}{2d}, \quad i \in \mathbb{Z}^d.$$

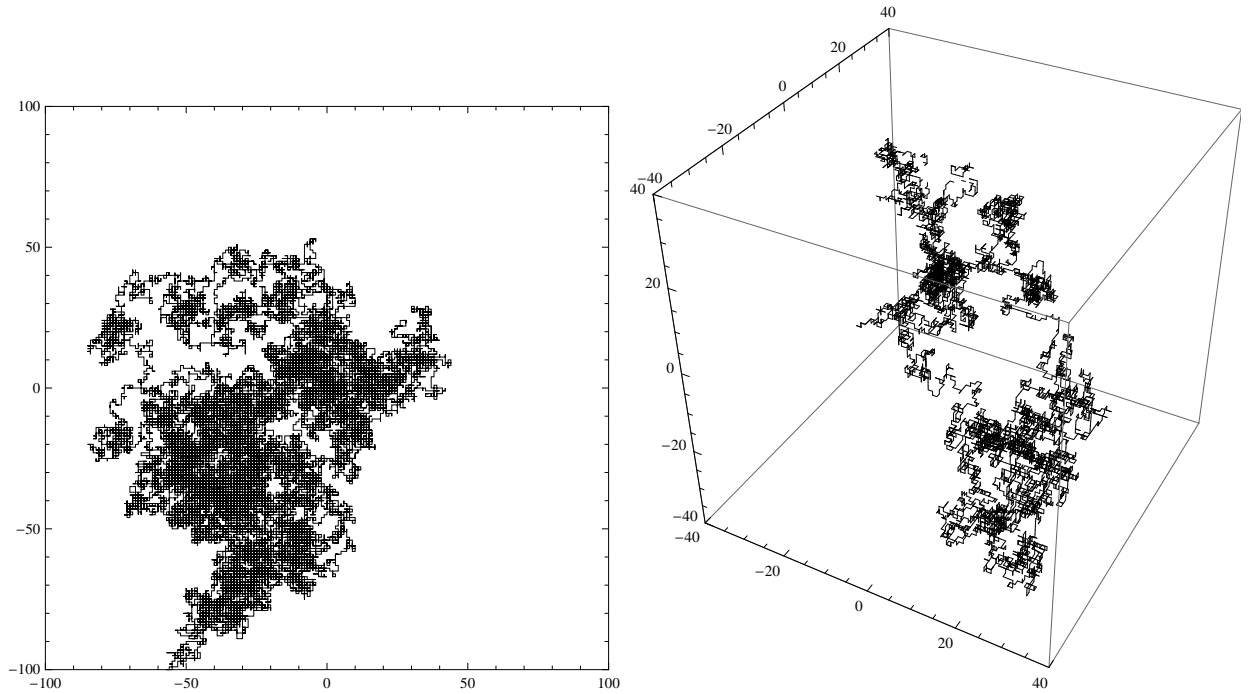


FIGURE 2. *Left:* Sample path of a simple symmetric random walk on \mathbb{Z}^2 . *Right:* Sample path of a simple symmetric random walk on \mathbb{Z}^3 . In both cases the random walk makes 50000 steps.

THEOREM 2.8.4 (Pólya, 1921). *The simple symmetric random walk on \mathbb{Z}^d is recurrent if and only if $d = 1, 2$ and transient if and only if $d \geq 3$.*

PROOF. For $d = 1$ we already proved the statement in Theorem 2.8.2.

Consider the case $d = 2$. We compute the n -step transition probability $p_{00}^{(n)}$. For an odd n this probability is 0. For an even $n = 2k$ we have

$$p_{00}^{(2k)} = \frac{1}{4^{2k}} \sum_{i=0}^k \binom{2k}{i, i, k-i, k-i} = \frac{1}{4^{2k}} \binom{2k}{k} \sum_{i=0}^k \binom{k}{i} \binom{k}{k-i} = \left(\frac{1}{2^{2k}} \binom{2k}{k} \right)^2 \sim \frac{1}{\pi k},$$

as $k \rightarrow \infty$, where the last step is by the Stirling formula. The harmonic series $\sum_{k=1}^{\infty} \frac{1}{k}$ diverges. Therefore, $\sum_{n=1}^{\infty} p_{00}^{(n)} = \infty$ and the random walk is recurrent in $d = 2$ dimensions. Generalizing the cases $d = 1, 2$ one can show that for an arbitrary dimension $d \in \mathbb{N}$ we have, as $k \rightarrow \infty$,

$$p_{00}^{(2k)} \sim \frac{1}{(\pi k)^{d/2}}.$$

Since the series $\sum_{k=1}^{\infty} k^{-d/2}$ is convergent for $d \geq 3$ it holds that $\sum_{n=1}^{\infty} p_{00}^{(n)} < \infty$ and the random walk is transient in $d = 3$ dimensions. \square

2.9. Existence and uniqueness of the invariant measure

The next two theorems state that any irreducible and recurrent Markov chain has a *unique invariant measure* λ , up to a multiplication by a constant. This measure may be finite (that is, $\sum_{i \in E} \lambda_i < +\infty$) or infinite (that is, $\sum_{i \in E} \lambda_i = +\infty$).

First we provide an explicit construction of an invariant measure for an irreducible and recurrent Markov chain. Consider a Markov chain starting at state $k \in E$. Denote the time of the first return to k by

$$T_k = \min\{n \in \mathbb{N} : X_n = k\} \in \mathbb{N} \cup \{+\infty\}.$$

The minimum of an empty set is by convention $+\infty$. For a state $i \in E$ denote the expected number of visits to i before the first return to k by

$$\gamma_i = \gamma_i^{(k)} = \mathbb{E}_k \sum_{n=0}^{T_k-1} \mathbb{1}_{\{X_n=i\}} \in [0, +\infty].$$

THEOREM 2.9.1. *For an irreducible and recurrent Markov chain starting at state $k \in E$ we have*

- (1) $\gamma_k = 1$.
- (2) *For all $i \in E$ it holds that $0 < \gamma_i < \infty$.*
- (3) $\gamma = (\gamma_i)_{i \in E}$ *is an invariant measure.*

PROOF.

STEP 1. We show that $\gamma_k = 1$. By definition of T_k , we have $\sum_{n=0}^{T_k-1} \mathbb{1}_{\{X_n=k\}} = 1$, if the chain starts at k . It follows that $\gamma_k = \mathbb{E}_k 1 = 1$.

STEP 2. We show that for every state $j \in E$,

$$(2.9.1) \quad \gamma_j = \sum_{i \in E} p_{ij} \gamma_i.$$

(At this moment, both sides of (2.9.1) are allowed to be infinite, but in Step 3 we will show that both sides are actually finite). The Markov chain is recurrent, thus $T_k < \infty$ almost surely. By definition, $X_{T_k} = k = X_0$. We have

$$\gamma_j = \mathbb{E}_k \sum_{n=1}^{T_k} \mathbb{1}_{\{X_n=j\}} = \mathbb{E}_k \sum_{n=1}^{\infty} \mathbb{1}_{\{X_n=j, n \leq T_k\}} = \sum_{n=1}^{\infty} \mathbb{P}_k[X_n = j, T_k \geq n].$$

Before visiting state j at time n the chain must have been in some state i at time $n - 1$, where $i \in E$ can be, in general, arbitrary. We obtain that

$$\gamma_j = \sum_{i \in E} \sum_{n=1}^{\infty} \mathbb{P}_k[X_n = j, X_{n-1} = i, T_k \geq n] = \sum_{i \in E} \sum_{n=1}^{\infty} p_{ij} \mathbb{P}_k[X_{n-1} = i, T_k \geq n].$$

Introducing the new summation variable $m = n - 1$, we obtain that

$$\gamma_j = \sum_{i \in E} p_{ij} \sum_{m=0}^{\infty} \mathbb{E}_k \mathbb{1}_{\{X_m=i, T_k \geq m+1\}} = \sum_{i \in E} p_{ij} \mathbb{E}_k \sum_{m=0}^{T_k-1} \mathbb{1}_{\{X_m=i\}} = \sum_{i \in E} p_{ij} \gamma_i.$$

This proves that (2.9.1) holds.

STEP 3. Let $i \in E$ be an arbitrary state. We show that $0 < \gamma_i < \infty$. Since the chain is irreducible, there exist $n, m \in \mathbb{N}_0$ such that $p_{ik}^{(m)} > 0$ and $p_{ki}^{(n)} > 0$. From (2.9.1) it follows that

$$\gamma_i = \sum_{l \in E} p_{li}^{(n)} \gamma_l \geq p_{ki}^{(n)} \gamma_k = p_{ki}^{(n)} > 0.$$

On the other hand, again using (2.9.1), we obtain that

$$1 = \gamma_k = \sum_{l \in E} p_{lk}^{(m)} \gamma_l \geq p_{ik}^{(m)} \gamma_i.$$

This implies that $\gamma_i \leq 1/p_{ik}^{(m)} < \infty$. □

The next theorem states the uniqueness of the invariant measure, up to multiplication by a constant.

THEOREM 2.9.2. *Consider an irreducible and recurrent Markov chain and fix some state $k \in E$. Then, every invariant measure λ can be represented in the form*

$$\lambda_j = c \gamma_j^{(k)} \quad \text{for all } j \in E,$$

where c is a constant (not depending on j). In fact, $c = \lambda_k$.

REMARK 2.9.3. Hence, the invariant measure is unique up to a multiplication by a constant. In particular, the invariant measures $(\gamma_i^{(k_1)})_{i \in E}$ and $(\gamma_i^{(k_2)})_{i \in E}$, for different states $k_1, k_2 \in E$, differ by a multiplicative constant.

PROOF. Let λ be an invariant measure.

STEP 1. We show that $\lambda_j \geq \lambda_k \gamma_j^{(k)}$ for all $j \in E$. We will *not* use the irreducibility and the recurrence of the chain in this step. The invariance of the measure λ implies that

$$\lambda_j = \sum_{i_0 \in E} \lambda_{i_0} p_{i_0 j} = \sum_{i_0 \neq k} \lambda_{i_0} p_{i_0 j} + \lambda_k p_{kj}.$$

Applying the same procedure to λ_{i_0} , we obtain

$$\begin{aligned}\lambda_j &= \sum_{i_0 \neq k} \left(\sum_{i_1 \neq k} \lambda_{i_1} p_{i_1 i_0} + \lambda_k p_{k i_0} \right) p_{i_0 j} + \lambda_k p_{k j} \\ &= \sum_{i_0 \neq k} \sum_{i_1 \neq k} \lambda_{i_1} p_{i_1 i_0} p_{i_0 j} + \left(\lambda_k p_{k j} + \lambda_k \sum_{i_0 \neq k} p_{k i_0} p_{i_0 j} \right).\end{aligned}$$

Applying the procedure to λ_{i_1} and repeating it over and over again we obtain that for every $n \in \mathbb{N}$,

$$\lambda_j = \sum_{i_0, i_1, \dots, i_n \neq k} \lambda_{i_n} p_{i_n i_{n-1}} \cdots p_{i_1 i_0} p_{i_0 j} + \lambda_k \left(p_{k j} + \sum_{i_0 \neq k} p_{k i_0} p_{i_0 j} + \cdots + \sum_{i_0, \dots, i_{n-1} \neq k} p_{k i_0} p_{i_0 i_1} \cdots p_{i_{n-1} j} \right).$$

Noting that the first term is non-negative, we obtain that

$$\lambda_j \geq 0 + \lambda_k \mathbb{P}_k[X_1 = j, T_k \geq 1] + \lambda_k \mathbb{P}_k[X_2 = j, T_k \geq 2] + \cdots + \lambda_k \mathbb{P}_k[X_n = j, T_k \geq n].$$

Since this holds for every $n \in \mathbb{N}$, we can pass to the limit as $n \rightarrow \infty$:

$$\lambda_j \geq \lambda_k \sum_{n=1}^{\infty} \mathbb{P}_k[X_n = j, T_k \geq n] = \lambda_k \gamma_j^{(k)}.$$

It follows that $\lambda_j \geq \lambda_k \gamma_j^{(k)}$.

STEP 2. We prove the converse inequality. Consider $\mu_j := \lambda_j - \lambda_k \gamma_j^{(k)}$, $j \in E$. By the above, $\mu_j \geq 0$ for all $j \geq 0$ so that $\mu = (\mu_j)_{j \in E}$ is a measure. Moreover, this measure is invariant because it is a linear combination of two invariant measures. Finally, note that by definition, $\mu_k = 0$. We will prove that this implies that $\mu_j = 0$ for all $j \in E$. By the irreducibility of our Markov chain, for every $j \in E$ we can find $n \in \mathbb{N}_0$ such that $p_{jk}^{(n)} > 0$. By the invariance property of μ ,

$$0 = \mu_k = \sum_{i \in E} \mu_i p_{ik}^{(n)} \geq \mu_j p_{jk}^{(n)}.$$

It follows that $\mu_j p_{jk}^{(n)} = 0$ but since $p_{jk}^{(n)} > 0$, we must have $\mu_j = 0$. By the definition of μ_j this implies that $\lambda_j = \lambda_k \gamma_j^{(k)}$. \square

We can now summarize Theorems 2.9.1 and 2.9.2 as follows:

THEOREM 2.9.4. *A recurrent, irreducible Markov chain has unique (up to a constant multiple) invariant measure.*

This invariant measure may be finite or infinite. However, if the Markov chain has only finitely many states, then the measure must be finite and we can even normalize it to be a *probability* measure.

COROLLARY 2.9.5. *A finite and irreducible Markov chain has a unique invariant probability measure.*

PROOF. A finite and irreducible Markov chain is recurrent by Theorem 2.7.10. By Theorem 2.9.1, there exists an invariant measure $\lambda = (\lambda_i)_{i \in E}$. Since the number of states in E is finite by assumption and $\lambda_i < \infty$ by Theorem 2.9.1, we have $M := \sum_{i \in E} \lambda_i < \infty$ and hence, the measure λ is finite. To obtain an invariant *probability* measure, consider the measure $\lambda'_i = \lambda_i/M$.

To show that the invariant probability measure is unique, assume that we have two invariant probability measures $\nu' = (\nu'_i)_{i \in E}$ and $\nu'' = (\nu''_i)_{i \in E}$. Take an arbitrary state $k \in E$. By Theorem 2.9.2, there are constants c' and c'' such that $\nu'_i = c' \gamma_i^{(k)}$ and $\nu''_i = c'' \gamma_i^{(k)}$, for all $i \in E$. But since both ν' and ν'' are *probability* measures, we have

$$1 = \sum_{i \in E} \nu'_i = c' \sum_{i \in E} \gamma_i^{(k)}, \quad 1 = \sum_{i \in E} \nu''_i = c'' \sum_{i \in E} \gamma_i^{(k)}.$$

This implies that $c' = c''$ and hence, the measures ν' and ν'' are equal. \square

Above, we considered only irreducible, recurrent chains. What happens if the chain is irreducible and transient? It turns out that in this case everything is possible:

- (1) It is possible that there is no invariant measure at all (except the zero measure).
- (2) It is possible that there is a unique (up to multiplication by a constant) invariant measure.
- (3) It is possible that there are at least two invariant measures which are not constant multiples of each other.

EXERCISE 2.9.6. Consider a Markov chain on \mathbb{N} with transition probabilities $p_{i,i+1} = 1$, for all $i \in \mathbb{N}$. Show that the only invariant measure is $\lambda_i = 0$, $i \in \mathbb{N}$.

EXERCISE 2.9.7. Consider a Markov chain on \mathbb{Z} with transition probabilities $p_{i,i+1} = 1$, for all $i \in \mathbb{Z}$. Show that the invariant measures have the form $\lambda_i = c$, $i \in \mathbb{Z}$, where $c \geq 0$ is constant.

EXERCISE 2.9.8. Consider a simple random walk on \mathbb{Z} with $p \neq \frac{1}{2}$. Show that any invariant measure has the form

$$\lambda_i = c_1 + c_2 \left(\frac{p}{1-p} \right)^i, \quad i \in \mathbb{Z},$$

for some constants $c_1 \geq 0$, $c_2 \geq 0$.

2.10. Positive recurrence and null recurrence

The set of recurrent states of a Markov chain can be further subdivided into the set of positive recurrent states and the set of negative recurrent states. Let us define the notions of positive recurrence and null recurrence.

Consider a Markov chain on state space E . Take some state $i \in E$, assume that the Markov chain starts at state i and denote by T_i the time of the first return of the chain to state i :

$$T_i = \min\{n \in \mathbb{N} : X_n = i\} \in \mathbb{N} \cup \{+\infty\}.$$

Denote by m_i the expected return time of the chain to state i , that is

$$m_i = \mathbb{E}_i T_i \in (0, \infty]$$

Note that for a transient state i we always have $m_i = +\infty$ because the random variable T_i takes the value $+\infty$ with strictly positive probability $1 - f_i > 0$, see Theorem 2.7.3. However, for a recurrent state i the value of m_i may be both finite and infinite, as we shall see later.

DEFINITION 2.10.1. A state $i \in E$ is called *positive recurrent* if $m_i < \infty$.

DEFINITION 2.10.2. A state $i \in E$ is called *null recurrent* if it is recurrent and $m_i = +\infty$.

REMARK 2.10.3. Both null recurrent states and positive recurrent states are recurrent. For null recurrent states this is required by definition. For a positive recurrent state we have $m_i < \infty$ which means that T_i cannot attain the value $+\infty$ with strictly positive probability and hence, state i is recurrent.

THEOREM 2.10.4. *Consider an irreducible Markov chain. Then the following statements are equivalent:*

- (1) *Some state is positive recurrent.*
- (2) *All states are positive recurrent.*
- (3) *The chain has invariant **probability** measure $\lambda = (\lambda_i)_{i \in E}$.*

Also, if these statements hold, then $m_i = \frac{1}{\lambda_i}$ for all $i \in E$.

PROOF. The implication $2 \Rightarrow 1$ is evident.

PROOF OF $1 \Rightarrow 3$. Let $k \in E$ be a positive recurrent state. Then, k is recurrent and all states of the chain are recurrent by irreducibility. By Theorem 2.9.1, $(\gamma_i^{(k)})_{i \in E}$ is an invariant measure. However, we need an invariant *probability* measure. To construct it, note that

$$\sum_{j \in E} \gamma_j^{(k)} = m_k < \infty$$

(since k is positive recurrent). We can therefore define $\lambda_i = \gamma_i^{(k)} / m_k$, $i \in E$. Then, $\sum_{i \in E} \lambda_i = 1$, and $(\lambda_i)_{i \in E}$ is an invariant probability measure.

PROOF OF $3 \Rightarrow 2$. Let $(\lambda_i)_{i \in E}$ be an invariant probability measure. First we show that $\lambda_k > 0$ for every state $k \in E$. Since λ is a probability measure, we have $\lambda_l > 0$ for at least one $l \in E$. By irreducibility, we have $p_{lk}^{(n)} > 0$ for some $n \in \mathbb{N}_0$ and by invariance of λ , we have

$$\lambda_k = \sum_{i \in E} p_{ik}^{(n)} \lambda_i \geq p_{lk}^{(n)} \lambda_l > 0.$$

This proves that $\lambda_k > 0$ for every $k \in E$.

By Step 1 from the proof of Theorem 2.9.2 (note that this step does not use recurrence), we have for all $j \in E$,

$$\lambda_i \geq \lambda_k \gamma_i^{(k)}.$$

Hence,

$$m_k = \sum_{i \in E} \gamma_i^{(k)} \leq \sum_{i \in E} \frac{\lambda_i}{\lambda_k} = \frac{1}{\lambda_k} < \infty.$$

It follows that k is positive recurrent, thus establishing statement 2.

PROOF THAT $m_k = \frac{1}{\lambda_k}$. Assume that statements 1,2,3 hold. In particular, the chain is recurrent and by Theorem 2.9.2, we must have $\lambda_i = \lambda_k \gamma_i^{(k)}$ for all $i \in E$. It follows that

$$m_k = \sum_{i \in E} \gamma_i^{(k)} = \sum_{i \in E} \frac{\lambda_i}{\lambda_k} = \frac{1}{\lambda_k},$$

thus proving the required formula. \square

EXAMPLE 2.10.5. Any state in a *finite* irreducible Markov chain is positive recurrent. Indeed, such a chain has an invariant probability measure by Corollary 2.9.5.

EXAMPLE 2.10.6. Consider a simple symmetric random walk on \mathbb{Z} or on \mathbb{Z}^2 . This chain is irreducible. Any state is recurrent by Pólya's Theorem 2.8.4. We show that in fact, any state is *null* recurrent. To see this, note that the measure assigning the value 1 to every state $i \in E$ is invariant by the definition of the chain. By Theorem 2.9.2, any other invariant measure must be of the form $\lambda_i = c$, $i \in E$, for some constant $c \geq 0$. However, no measure of this form is a probability measure. So, there is no invariant probability measure and by Theorem 2.10.4, all states must be null recurrent.

2.11. Convergence to the invariant probability measure

We are going to state and prove a “strong law of large numbers” for Markov chains. First recall that the usual strong law of large numbers states that if ξ_1, ξ_2, \dots are i.i.d. random variables with $\mathbb{E}|\xi_1| < \infty$, then

$$(2.11.1) \quad \frac{\xi_1 + \dots + \xi_n}{n} \xrightarrow[n \rightarrow \infty]{a.s.} \mathbb{E}\xi_1.$$

The statement is not applicable if $\mathbb{E}|\xi_1| = \infty$. However, it is an exercise to show that if ξ_1, ξ_2, \dots are i.i.d. random variables which are a.s. *nonnegative* with $\mathbb{E}\xi_1 = +\infty$, then

$$(2.11.2) \quad \frac{\xi_1 + \dots + \xi_n}{n} \xrightarrow[n \rightarrow \infty]{a.s.} +\infty.$$

Consider a Markov chain $\{X_n : n \in \mathbb{N}_0\}$ with initial distribution $\alpha = (\alpha_i)_{i \in E}$. Given a state $i \in E$, denote the number of visits to state i in the first n steps by

$$V_i(n) = \sum_{k=0}^{n-1} \mathbb{1}_{\{X_k=i\}}.$$

THEOREM 2.11.1. Consider an irreducible Markov chain $\{X_n : n \in \mathbb{N}_0\}$ with an arbitrary initial distribution $\alpha = (\alpha_i)_{i \in E}$.

1. If the Markov chain is transient or null recurrent, then for all $i \in E$ it holds that

$$(2.11.3) \quad \frac{V_i(n)}{n} \xrightarrow[n \rightarrow \infty]{} 0 \quad a.s.$$

2. If the Markov chain is positive recurrent with invariant probability measure λ , then for all $i \in E$ it holds that

$$(2.11.4) \quad \frac{V_i(n)}{n} \xrightarrow[n \rightarrow \infty]{} \lambda_i \quad a.s.$$

PROOF. If the chain is transient, then $V_i(n)$ stays bounded as a function of n , with probability 1. This implies (2.11.3). In the sequel, let the chain be recurrent.

For simplicity, we will assume in this proof that the chain starts in state i . Denote the time of the k -th visit of the chain to i by S_k , that is

$$\begin{aligned} S_1 &= \min \{n \in \mathbb{N} : X_n = i\}, \\ S_2 &= \min \{n > S_1 : X_n = i\}, \\ S_3 &= \min \{n > S_2 : X_n = i\}, \end{aligned}$$

and so on. Note that S_1, S_2, S_3, \dots are a.s. finite by the recurrence of the chain. Let also $\xi_1, \xi_2, \xi_3, \dots$ be the excursion times between the returns to i , that is

$$\xi_1 = S_1, \quad \xi_2 = S_2 - S_1, \quad \xi_3 = S_3 - S_2, \quad \dots$$

Then, $\xi_1, \xi_2, \xi_3, \dots$ are i.i.d. random variables by the Markov property.

By definition of $V_i(n)$ we have

$$\xi_1 + \xi_2 + \dots + \xi_{V_i(n)-1} \leq n \leq \xi_1 + \xi_2 + \dots + \xi_{V_i(n)}.$$

Dividing this by $V_i(n)$ we get

$$(2.11.5) \quad \frac{\xi_1 + \xi_2 + \dots + \xi_{V_i(n)-1}}{V_i(n)} \leq \frac{n}{V_i(n)} \leq \frac{\xi_1 + \xi_2 + \dots + \xi_{V_i(n)}}{V_i(n)}.$$

Note that by recurrence, $V_i(n) \xrightarrow[n \rightarrow \infty]{} \infty$ a.s.

CASE 1. Let the chain be *null* recurrent. It follows that $\mathbb{E}\xi_1 = \infty$. By using (2.11.2) and (2.11.5), we obtain that

$$\frac{n}{V_i(n)} \xrightarrow[n \rightarrow \infty]{a.s.} \infty.$$

This proves (2.11.3). □

CASE 2. Let the chain be *positive* recurrent. Then, by Theorem 2.10.4, $\mathbb{E}\xi_1 = m_i = \frac{1}{\lambda_i} < \infty$. Using (2.11.1) and (2.11.5) we obtain that

$$\frac{n}{V_i(n)} \xrightarrow[n \rightarrow \infty]{a.s.} \frac{1}{\lambda_i}.$$

This proves (2.11.4). □

In the next theorem we prove that the n -step transition probabilities converge, as $n \rightarrow \infty$, to the invariant probability measure.

THEOREM 2.11.2. *Consider an irreducible, aperiodic, positive recurrent Markov chain $\{X_n : n \in \mathbb{N}_0\}$ with transition matrix P and invariant probability measure $\lambda = (\lambda_i)_{i \in E}$. The initial distribution $\alpha = (\alpha_i)_{i \in E}$ may be arbitrary. Then, for all $j \in E$ it holds that*

$$\lim_{n \rightarrow \infty} \mathbb{P}[X_n = j] = \lambda_j.$$

In particular, $\lim_{n \rightarrow \infty} p_{ij}^{(n)} = \lambda_j$ for all $i, j \in E$.

REMARK 2.11.3. In particular, the theorem applies to any irreducible and aperiodic Markov chain with finite state space.

For the proof we need the following lemma.

LEMMA 2.11.4. *Consider an irreducible and aperiodic Markov chain. Then, for every states $i, j \in E$ we can find $N = N(i, j) \in \mathbb{N}$ such that for all $n > N$ we have $p_{ij}^{(n)} > 0$.*

PROOF. The chain is irreducible, hence we can find $r \in \mathbb{N}_0$ such that $p_{ij}^{(r)} > 0$. Also, the chain is aperiodic, hence we can find $N_0 \in \mathbb{N}$ such that for all $k > N_0$ we have $p_{ii}^{(k)} > 0$. It follows that for all $k > N_0$,

$$p_{ij}^{(k+r)} > p_{ii}^{(k)} p_{ij}^{(r)} > 0.$$

It follows that for every $n := k + r$ such that $n > N_0 + r$, we have $p_{ij}^{(n)} > 0$. \square

PROOF OF THEOREM 2.11.2. We use the “coupling method”.

STEP 1. Consider *two* Markov chains called $\{X_n : n \in \mathbb{N}_0\}$ and $\{Y_n : n \in \mathbb{N}_0\}$ such that

- (1) X_n is a Markov chain with initial distribution α and transition matrix P .
- (2) Y_n is a Markov chain with initial distribution λ (the invariant probability measure) and the same transition matrix P .
- (3) The process $\{X_n : n \in \mathbb{N}_0\}$ is independent of the process $\{Y_n : n \in \mathbb{N}_0\}$.

Note that both Markov chains have the same transition matrix but different initial distributions. Fix an arbitrary state $b \in E$. Denote by T be the time at which the chains meet at state b :

$$T = \min\{n \in \mathbb{N} : X_n = Y_n = b\} \in \mathbb{N} \cup \{+\infty\}.$$

If the chains do not meet at b , we set $T = +\infty$.

STEP 2. We show that $\mathbb{P}[T < \infty] = 1$. Consider the stochastic process $W_n = (X_n, Y_n)$ taking values in $E \times E$. It is a Markov chain on $E \times E$ with transition probabilities given by

$$\tilde{p}_{(i,k),(j,l)} = p_{ij} p_{kl}, \quad (i, k) \in E \times E, \quad (j, l) \in E \times E.$$

The initial distribution of W_0 is given by

$$\mu_{(i,k)} = \alpha_i \lambda_k, \quad (i, k) \in E \times E.$$

Since the chains X_n and Y_n are aperiodic and irreducible by assumption of the theorem, we can apply Lemma 2.11.4 to obtain for every $i, j, k, l \in E$ a number $N = N(i, j, k, l) \in \mathbb{N}$ such that for all $n > N$ we have

$$\tilde{p}_{(i,k),(j,e)}^{(n)} = p_{ij}^{(n)} p_{ke}^{(n)} > 0.$$

Thus, the chain W_n is irreducible. Also, it is an exercise to check that the probability measure $\tilde{\lambda}_{(i,k)} := \lambda_i \lambda_k$ is invariant for W_n . Thus, by Theorem 2.10.4, the Markov chain W_n is positive recurrent and thereby recurrent. Therefore, $T < \infty$ a.s. by Lemma 2.7.11.

STEP 3. Define the stochastic process $\{Z_n : n \in \mathbb{N}_0\}$ by

$$Z_n = \begin{cases} X_n, & \text{if } n \leq T, \\ Y_n, & \text{if } n \geq T. \end{cases}$$

Then, Z_n is a Markov chain with initial distribution α and the same transition matrix P as X_n and Y_n . (The Markov chain Z_n is called the coupling of X_n and Y_n). The chain Y_n starts with the invariant probability measure λ and hence, at every time n , Y_n is distributed according to λ . Also, the chain Z_n has the same initial distribution α and the same transition

matrix P as the chain X_n , so that in particular, the random elements X_n and Z_n have the same distribution at every time n . Using these facts, we obtain that

$$|\mathbb{P}[X_n = j] - \lambda_j| = |\mathbb{P}[X_n = j] - \mathbb{P}[Y_n = j]| = |\mathbb{P}[Z_n = j] - \mathbb{P}[Y_n = j]|.$$

By definition of Z_n , we can rewrite this as

$$\begin{aligned} |\mathbb{P}[X_n = j] - \lambda_j| &= |\mathbb{P}[X_n = j, n < T] + \mathbb{P}[Y_n = j, n \geq T] - \mathbb{P}[Y_n = j]| \\ &= |\mathbb{P}[X_n = j, n < T] - \mathbb{P}[Y_n = j, n < T]| \\ &\leq \mathbb{P}[T > n]. \end{aligned}$$

But we have shown in Step 2 that $\mathbb{P}[T = \infty] = 0$, hence $\lim_{n \rightarrow \infty} \mathbb{P}[T > n] = 0$. It follows that

$$\lim_{n \rightarrow \infty} \mathbb{P}[X_n = j] = \lambda_j,$$

thus establishing the theorem. □