Order statistics: from simple to O(n)

Banin, M.

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When in dobut

How to do it by hand

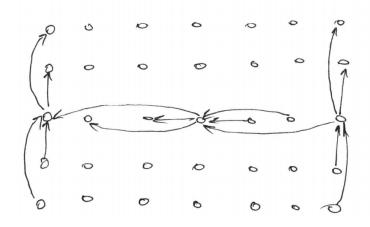
```
template <class Iter, class T>
void my_element(Iter first, Iter nth, Iter last) {
 while (last - first > 3) {
    T pivot = *(first + (last - first) / 2);
    Iter 11 = partition(first, last,
      [pivot](const T& a) -> bool{ return a < pivot; });</pre>
    Iter rr = partition(11, last,
      [pivot](const T&a) -> bool{ return a == pivot; });
   if (nth < 11) last = 11;</pre>
    else if(nth >= 11 && nth < rr) return;
    else if(nth >= rr) first = rr:
  std::sort(first, last);
```

What STL does

```
template<typename _RandomAccessIterator, typename _Size, typename _Compare>
 void
 __introselect(_RandomAccessIterator __first, _RandomAccessIterator __nth,
               _RandomAccessIterator __last, _Size __depth_limit,
               Compare comp)
 while (__last - __first > 3)
       if (__depth_limit == 0)
            std::_heap_select(__first, __nth + 1, __last, __comp);
            // Place the nth largest element in its final position.
            std::iter_swap(__first, __nth);
           return:
        --__depth_limit;
       _RandomAccessIterator __cut =
          std::_unguarded_partition_pivot(__first, __last, __comp);
       if ( cut <= nth)
         __first = __cut;
       else
         __last = __cut;
   std::__insertion_sort(__first, __last, __comp);
                                                4□ → 4周 → 4 = → 4 = → 9 < ○</p>
```

```
/// This is a helper function...
template<typename _RandomAccessIterator, typename _Compare>
  RandomAccessIterator
  __unguarded_partition(_RandomAccessIterator __first,
                        _RandomAccessIterator __last,
                        _RandomAccessIterator __pivot, _Compare __comp)
    while (true)
        while (__comp(__first, __pivot))
          ++ first:
        -- last:
        while (__comp(__pivot, __last))
         -- last:
        if (!(__first < __last))</pre>
          return __first;
        std::iter_swap(__first, __last);
        ++__first;
```

Median-of medians (Blum, Floyd, Pratt, Rivest, Tarjan)



A bit of compexlity theory

$$ightharpoonup \min x = 5^i : x > n = O(n)$$

$$ightharpoonup rac{2}{10} + rac{7}{10} < 1$$

Makeshift strict O(n) implementation

```
template <class Iter, class T>
void lintime_element(Iter first, Iter nth, Iter last) {
    while (last - first > 5) {
        std::vector<T> medians:
        for(Iter it = first; it < last; it += 5){</pre>
            sort(it, std::min(it + 5, last));
            medians.push_back(*std::min(it + 2, last - 1));
        T padding = *max_element(medians.begin(), medians.end());
        while(medians.size() % 5 != 0)
            medians.push_back(padding);
        Iter mid = medians.begin() + (medians.end() - medians.begin()) / 2;
        lintime_element<typename std::vector<T>::iterator, T>(medians.begin(),
                                                               mid.
                                                               medians.end()):
        T pivot = *mid;
```

Benchmark

On 256K-element arrays, we see:

Implementation	Average runtime
sorting in $O(nlogn)$	52 <i>μsec</i>
makeshift strict $O(n)$	34 μsec
makeshift expected $O(n)$	$17~\mu sec$
nth_element call	15 μsec
generating the array	$11~\mu sec$

References

- http://www.cplusplus.com/reference/algorithm/nth_ element/
- ▶ https://en.wikipedia.org/wiki/Quickselect
- ▶ Blum, M.; Floyd, R. W.; Pratt, V. R.; Rivest, R. L.; Tarjan, R. E. (August 1973). "Time bounds for selection"
- Noriyuki Kurosawa. (2016) Quicksort with median of medians considered practic