

→ more important note

→ Heading

→ Subheading or important point

→ Data Structure

→ Complexities

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DSA-SUPREME

CODES

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If you like these notes,
Give me job opportunities

And give party to bhaiya after placement

IMPORTANT POINTS

- Focus on work on skill and build networks those work in companies and can talk to HR for you
- Internship- do if you want to learn or convert it as PPO
- Rough sols. and dry run are very important
- Do documentation to promote and mail everything you discuss with HR or team
- Think twice, Code once
- To clarify any code you are confused, use cout statements every where to know what is going on in the code
- Code all approaches you can think of and can find & understand from google
- Revise all incorrect & skipped questions in quizes regularly
- Watch sol. only after attempting the question
- Interviewer will ask Time & Space complexity after every sol. you give

→ 2 websites

- cplusplus.com
- cppreference

→ Think on paper → with 5 testcases atleast

→ Write readable codes

→ In interviews, tell approaches of questions using example

→ Signs of beginner → no logic build



can't solve new ques

forget the approach

memorizing the ans

Improve → with more no. of questions

with time

with dry run on mind logic, not on code

with your new approach

(even brute force)

with a lot of practice on syntax

with consistency

→ Focus on placement, not on feeling

You are not studing to get fun

You are studing to get placed

First placement, then fun and interest.

Also focus on health of you and your parents

Bhaad me gya interest yaar, placement ke baad karna

Jo duniya sunna chahati hai, use sunado

Growth is important

Week 4 [Connect] Class 1:45 - 2:01

→ 20-30 interview experience

↳ to break Google pattern

→ Web Dev and DSA → both are important

CP → do if you enjoy it.

→ At least learn one more approach if you are doing questions with map

Many times if you tell map approach, interviewer asks another approach

- Tag all khatarnak questions
 - Otherwise you will forget them after some time
- Revise and code all, every after 2 weeks
- For 6th sem student (like me),
 - At least do potd on leetcode everyday along with course
- Make notes
- In dev, you should have at least 2 major projects
 - 1 → in dot batch
 - 2 → group project if you want
- Focus on Networking,
After 5 months, you should have atleast 2 friends in every company
- All questions doing in batch, should be on your tips

- Focus on accuracy more than speed
- Resume is made company specific
- Make sure others don't take your credits
- Higher position $\propto \frac{\text{Development}}{\text{DSA Complexity}}$
- Interview experience before interview
- Flex se kuch nhi hota
- Comparision se kuch nhi hota
- Don't forcefully try to increase the speed of DSA learning,
Try to increase the number of hours spent everyday in DSA learning
- if ques seems tough in another datatype but easy in another datatype, then first do that in later datatype and after solving the sol. change datatypes

IMPORTANT C++ NOTES

→ -1 in binary 1.....1

- 2^{31} in binary 10...0

→ In left shift, vacant bit will be filled by 0

In right shift, vacant bit will be filled by MSB

1 → in -ve no.

0 → in +ve no.

$(-1) \gg n \longrightarrow$ (-1) no change
any no. of times

→ int a = 5;

cout << (++a) * (++a);

a * (++a), a=6

a * a, a=7

7 * 7 = 49

cout << (a++) * (a++);

5 * (a++), a=6

5 * a, a=7

5 * 7 = 35

cout << (++a) * (a++);

cout << (a++) * (a++);

a * (a++), a=6

5 * (a++), a=6

a * 6, a=7

5 * 6, a=

7 * 6 = 42

30

→ Continue cannot be used in switch case
can only used in loops

→ int main () {

`return 0;` —————> returns to OS

} 0 is used to represent successful execution

→ A cpp file can't have more than 1 main function

→ main() can't have return type other than int
in offline compiler

→ % is heavy operator, so try to use it less

→ array name returns address of first index

Integer Array -

`cout << arr ;` → `A1` → address of first index of array
address of integer

→ To find max., initialize ans with INT-MIN

To find min., initialize ans with INT-MAX

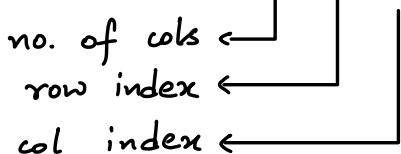
→ xor → cancels out same element

$$\rightarrow 0 \wedge \text{ans} = \text{ans}$$
$$0 \wedge 1 = 1$$
$$0 \wedge 0 = 0$$

→ Mapping of 2D array with memory 1D array

$$\text{Linear-index} = c * i + j$$

$i = \text{linear-index} / c$
 $j = \text{linear-index \% } c$



→ Arrays and 2D arrays of any datatype

↳ pass by reference

→ Search space

↳ range of ans (start and end) in binary search questions

→ store mid in ans if needed in binary search ques

→ In binary search questions

```
while (s <= e){  
    ans = mid;  
    s = mid + 1;  
}  
  
left  
search ← e = mid - 1;  
}  
  
right  
search
```

convert
ans = mid
s = mid + 1

```
while (s < e){  
    s = mid; → right search  
    left  
    search ← e = mid - 1;  
}  
  
OR  
s = mid + 1 → right  
search  
left  
search ← e = mid;  
}
```

convert
ans = mid, e = mid - 1

→ cout << (int)(-3.74);

→ - [int(3.74)]

→ - (3) → -3

cout << (-22)/7; → -3

cout << 22/(-7); → -3

→ In ASCII → '0' → 48
 'A' → 65
 'a' → 97

→ In sorted array, try to apply binary search or
2 pointer / 3pointer approach

→ To maximize or minimize
try to use binary search using concept of search space

- find search space of answer
- find mid
- if $\text{isPossibleAns}(\text{mid}) \rightarrow$ go to left / right acc.
to the ques.

You can use above approach also in other
ques like find the duplicate number in an array
having elements range 1 to n

→ Week 6 Lecture 2

→ Magical line for recursion

- ↳ Ek case solve krdo baaki recursion smblial lega
- ↳ Just believe on it, dont doubt on recursion

→ Week 7 Lecture 2

→ Week 7 Lecture 3

→ If the recursion ques seems tough,
use void function and pass ans by reference
in function

- If in ques. array elements are in range $[1, n]$ or $[0, n]$, try to treat elements as indices of an array
(Week 3 Assignment)
- Try to make cnts [max element present in array] instead of making cnts [max limit given in question]
Array more than 10^6 ints cant be formed

→ `int digit = 7;`
`string s;`
`s.push_back (digit + '0');`
`cout << s; —→ 7`

`s.push_back(digit);`
`cout << s; —→ contains`
`char with`
`ascii value = digit`

'7' to 7 —→ '7' - '0'

7 to '7' —→ 7 + '0'

ERROR / BAD PRACTICE

If let's go page

→ 13

31, 33, 43, 70, 71, 72, 73, 75, 76, 79, 86, 87, 89, 90, 93, 95

→ 14

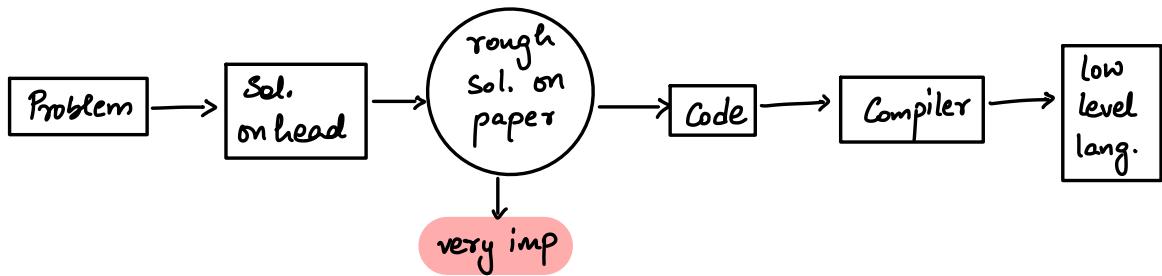
104

LET'S GO

Thought process to solve a problem-

W1-L1

- Understand a problem
- input values
- find approach



Algorithm - Sequence of steps

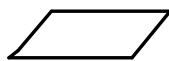
Flowchart - Graphical representation of algo

Components -



terminator

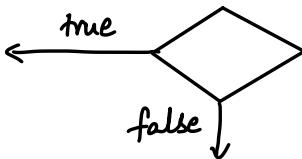
for start / end



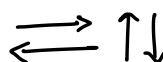
for input /output read /write



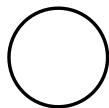
computation / process / declaration



decision making block
takes condition



flow



Connector
takes function

Pseudo Code - Generic way of writing algo

Dry Run → Very Important to understand any topic

W1-L2

IDE - Replit, VS-Code

```
# include <iostream>           → preheader file contains implementation of keyword
using namespace std;
int main () {
    cout << "Namaste Bharat";
}
```

region where scope of keyword is defined

used to point on console/standard display

→ using standard implementation of cout choosing from multiple types of namespace

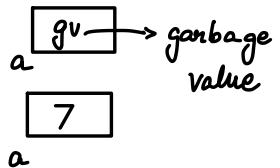
→ to end any statement

→ string

cout << endl; → for next line

cout << '\n';

int a; \longrightarrow a is an integer
 cin >> a; \longrightarrow input a from user
 ex - 7



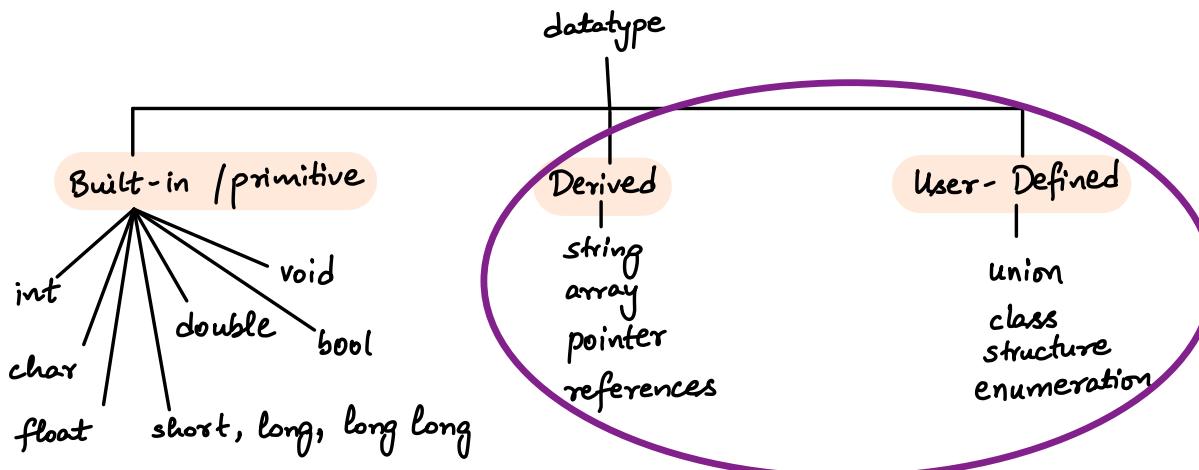
Variables

named memory location

int \downarrow
 datatype
 a = 5;
 ↓
 variable
 name

Datatypes

type of data



int - 4 byte - 32 bits in memory

\longrightarrow -2^{31} to $2^{31}-1$ in signed int
 \longrightarrow 0 to $2^{32}-1$ in unsigned int

char - 1 byte - 8 bits in memory

\longrightarrow 2^8 different chars.

ASCII

↳ char maps with numerical ASCII value

char \leftrightarrow ASCII value \rightarrow store in memory

bool \rightarrow 1 byte \rightarrow 8 bits

true - 1

false - 0

↳ because minimum addressable memory is
1 byte

We cannot address 1 bit in memory

float \rightarrow 4 byte \rightarrow 32 bits

double \rightarrow 8 byte \rightarrow 64 bits

long long \rightarrow 8 byte \rightarrow 64 bits

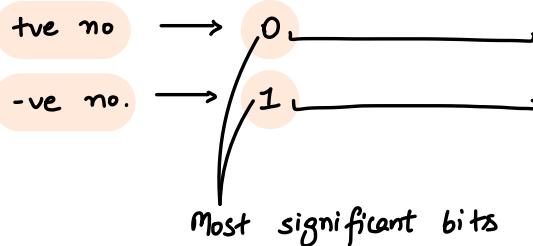
short \rightarrow 2 byte \rightarrow 16 bits

long \rightarrow 4 byte \rightarrow 32 bits

How data is stored

int a=5;

↳ 32 bits 0...00101
 29 bits



How -ve number is stored in memory

In 2's complement form

→ 1's complement + 1

→ reverse all bits

int a = -7;

7 → 0.....00111 } 32 bits

ignore -ve sign
find binary equivalent

1's (7) → 1.....11000

find 2's complement

2's (7) → 1.....11001

→ this is how -7 will be stored in memory

How to read -ve no. present in memory

→ take 2's complement

1....11001

→ 1's complement → 0...00110

2's complement → 0...00111

→ + 7

-7

Interesting problem



how computer know these are 4 chars or a single integer

↳ Using datatype

↳ tell 2 things

- ↳ type of data used
- ↳ space used in memory

Signed vs Unsigned

↓
↳ 0, +ve
+ve, -ve, 0

↳ by default

int - 4 byte - 32 bits in memory

↳ total no. of combinations - 2^{32}

signed int

-2^{31} to $2^{31}-1$

unsigned int

0 to $2^{32}-1$

} range

(1) 0...0

011...1

0.....0

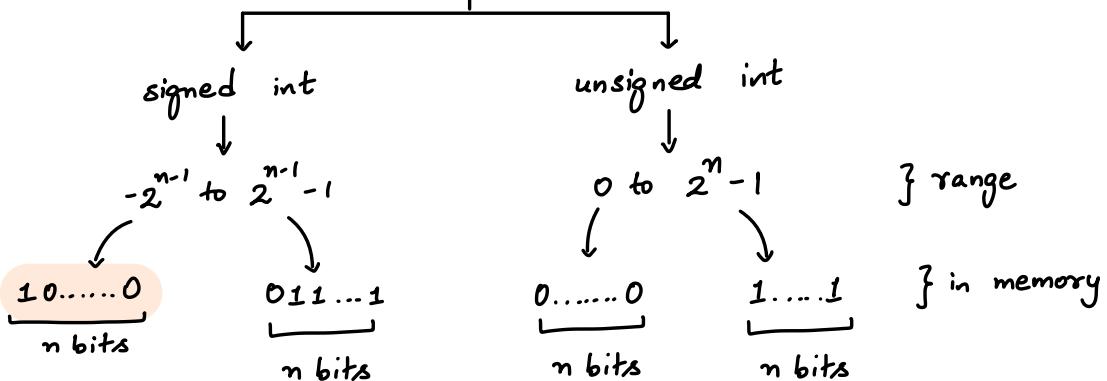
1....1

$2^{15} \rightarrow \underline{10...0} \rightarrow -2^{31}$

} in memory

General Formula

n bits in memory
↳ total no. of combinations - 2^n



$-1 \rightarrow \underbrace{1\dots1}_{n\text{ bits}}$

Typecasting

↳ convert one type of data to another

implicit typecasting

ex- char ch = 97;
cout << ch; → (a)

explicit typecasting

ex- char ch = (char) 97;
cout << ch; → (a)

overflow ex- char ch = 9999;
cout << ch;

9999 → 10011100001111
binary conversion stores only last 8 bits

so ch stores 00001111 in memory
 ↓
 ↓
 acc. to ASCII table

Operators -

Arithmetic Operator

→ +, -, *, /, %

int op int → int

float op int } float
 int op float } float

float op float }

bool op bool → int

bool act as 0 or 1

double op int }
 int op double } double
 double op double }
 float op double }
 double op float }

3 → int → by default → cout << sizeof(3.0);
 3.0 → float / double not int ↓
 ⑧

Relational Operator

(a op b)

>, <, >=, <=, !=, ==

Output - 0 or 1

false ← true

these are different things

Assignment Operators

=

Logical Operators

↳ when you have multiple conditions

$(a \& \& b)$ → and → true if both are true

$(a || b)$ → or → true if any one is true

$(!a)$ → not → negate the result

Output - 0 or 1
false ↕ true

$(\text{cond}1 \&\& \text{cond}2 \&\& \text{cond}3)$

if cond1 is false

compiler will not check further
as ans will already false

$(\text{cond}1 || \text{cond}2 || \text{cond}3)$

if cond1 is true

compiler will not check further
as ans will already true

Conditions

if (cond.){
execute
}

if

if (cond){
execute 1
}

else {
execute 2
}

if - else

W1-L3

if (cond1)
execute 1
else if (cond2)
execute 2

if - else if

```

if (cond 1)
    execute 1
else if ( cond 2)
    execute 2
else if (...)
else
    execute n

```

if - else if - else

```

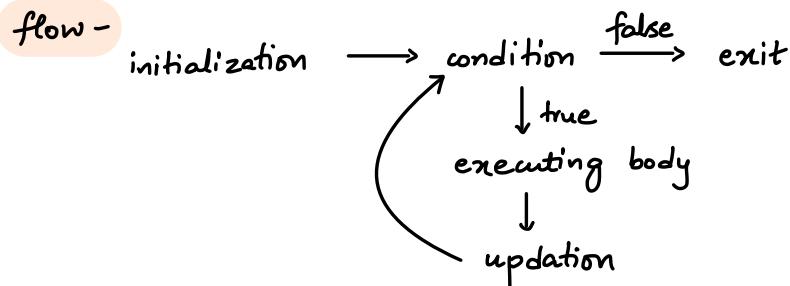
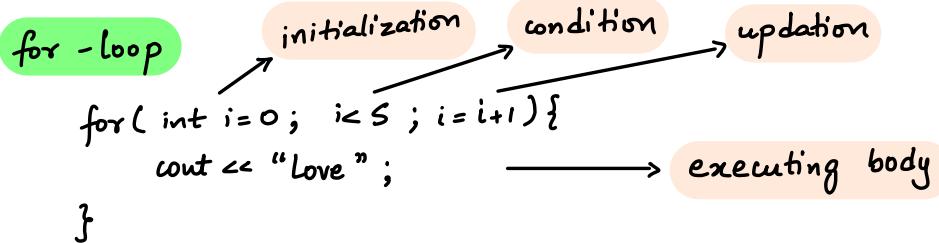
if (cond 1)
    execute 1
else {
    if () { }
    else () { }
}

```

nested if - else

Loops

↳ to do something repeatedly



initialization
condition
updation } none is mandatory
 one or multiple i,c,u can be added
multiple c → i>5, i<10; → i>5 & & i<10

patterns -

Generally 2 loops → outer loop() {
 inner loop() {
 }
 }
 cout << endl ;
}
for rows
for cols

→ a op = b → a = a op b

op → +, -, *, /, %, <<, >>

cin in if()

```
int num;  
if (cin >> num) {  
    cout << "hello";  
}  
else {  
    cout << "hi";  
}
```

it will not give error

output -
hello

for all input values of num
↓

0, true, -ve

cout in if()

```
int num = 0;  
if (cout << num << endl) {  
    cout << "hello";  
}  
else {  
    cout << "hi";  
}
```

it will not give error

output -
0
hello

for all values of num
↓

0, true, -ve

HLL - High level language

↳ human readable and user friendly

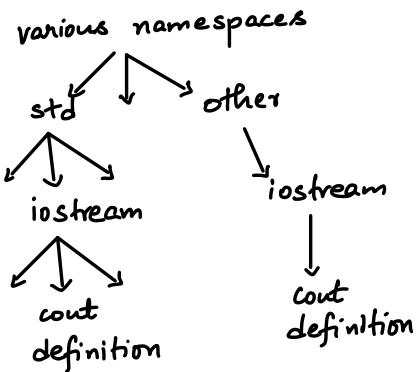
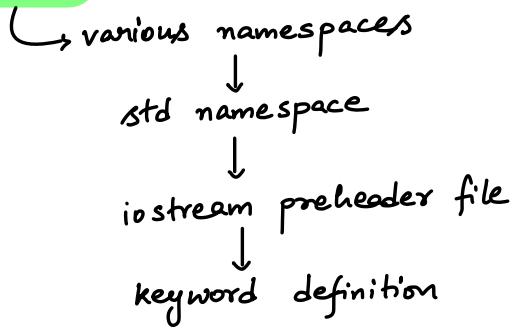
W1-L4

C++, C - Middle Level language

namespace → to avoid collision

↓
multiple definitions of a single keyword

hierarchy



float f = 2.0 + 100;

cout << f ; → output -

102 or 102.0
compiler dependent

float f = 2.7;

int n = 157;

int diff = n-f;

cout << diff ;

output -

154

explanation -

$$n-f = 157 - 2.7 = 154.3$$

int diff = n-f

diff = 154

ternary operator -

W1-HW

↳ syntax

variable = (condition) ? expression2 : expression3

(condition)? variable = expression2 : variable = expression3

2 and 3 cant be statements, they must be exp.

ex- return (a > b) ? a : b ; → CORRECT

(a > b) ? return a : return b ; → WRONG
ERROR

these are statements, not expression

by default -

cout << sizeof(2.3); → 8

↳ float

cout << sizeof(a); → 4 → int

↳ -(2^{31} -1) to $2^{31}-1$

cout << sizeof(-2 31) → 8

↳ long long

patterns

W2-L1

↳ how to think

→ finding formula for rows and cols

$n=5$

row	stars
0	0
1	0
2	1
3	2
4	3

→ formula -

0 to $< n-1$

$n-1$

-1

0

1

2

3

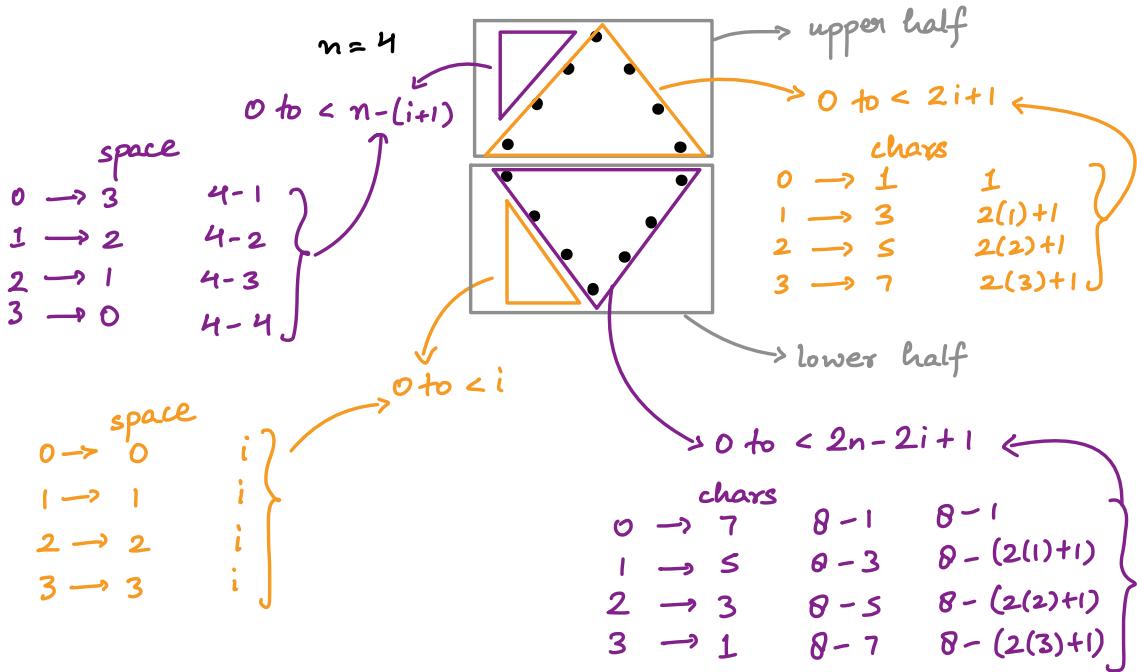
no. of times loop runs

as condition fails ($0 < -1$)

→ to do anything n times

↳ `for(i=0 ; i<n ; i++) {}`

→ break the complex patterns



Bitwise Operators

W2-L2

→ use on bit level

And $(a \& b)$ 1 if both bits are 1

Or $(a | b)$ 1 if any or both bits are 1

not $(\sim a)$ negate the result

nor $(a ^ b)$ same values $\rightarrow 0$
diff. values $\rightarrow 1$

~ 5

$\sim 5 \rightarrow 1 \dots 1010$

$\sim 5 \rightarrow 1 \dots 1010$

↳ how compiler read this

↳ 2's complement

$0 \dots 0101 \rightarrow 1$'s complement

$0 \dots 0110 \rightarrow 2$'s complement

-6

So $\sim 5 = -6$

Left and right shift operators

<<

shift all bits to left

* by 2 (not in every case)

↳ if MSB is 1 and

2nd MSB is 0

>>

shift all bits to right

/ by 2 (not in every case)

↳ in -1

$a = a \ll b$ a left shifts, b times \rightarrow result $\rightarrow a \times 2^b$

$a = a \gg b$ a right shifts, b times \rightarrow result $\rightarrow \frac{a}{2^b}$
b cant be -ve

↳ in case of -ve

$a = 5;$

↳ gives 8^v

$a = a \ll 1;$ $a = 10$

$a = 5;$

in left shift \rightarrow filled with 0

$a = a \ll 2;$ $a = 20$

in right shift \rightarrow filled with

in +ve no. 0 or 1
in -ve no.

right shift in -ve number

-ve no. in memory $\rightarrow 1 \dots$

↓ right shift

1 1 ...

↳ signed bit is used to fill
the vacant bit

ex-

$s \rightarrow 0 \dots 0 101$

$-s \rightarrow 1 \dots 1 011$

$-s \gg 1 \rightarrow 1 \dots 1 01 \rightarrow -3$

$-1 \gg 1 \rightarrow -1$

left shift in number where MSB is 1

and 2nd MSB is 0

no. $\rightarrow 1 0 \dots \rightarrow$ -ve no.

left shift $\hookrightarrow 0 \dots$

\hookrightarrow +ve no.

Pre- Post → Increment / Decrement Operator

pre- increment

↳ $++a$

↳ first increment by 1, then use

post - increment

↳ $a++$

↳ first use then increment by 1

pre- decrement

↳ $--a$

↳ first decrement by 1, then use

post - decrement

↳ $a--$

↳ first use then decrement by 1

int a = 5;

cout << (++a) * (++a);

output -

49

↳ due to operator precedence

→ links.txt in repo

break and continue

break

↳ exit from that loop

continue

↳ skip that iteration

Variable Scoping -

```
int g= 25;           -----> global variable
int main(){
    int a;          -----> declaration
    int b= 5;        -----> initialization
    b = 10;         -----> updation
    //int b= 15;     -----> redefinition is not allowed
    int c= 7;
    g= 30;
    cout << g;      -----> 30
    if (true){
        int b= 15;
        cout << b;      -----> 15
        cout << c;      -----> 7
        g= 50;
        cout << g;      -----> 50
    }
    cout << a;      -----> gv
    cout << b;      -----> 10
    cout << c;      -----> 7
    cout << g;      -----> 50
}
```

Making global variable is very **BAD PRACTICE**

Operator Precedence

- order of priority of operator
- no need to remember
- use brackets properly

Switch Case

```
switch (expression) {
```

```
    case value1 :
```

executing body 1

```
    break ;
```

```
    case value2 :
```

executing body 2

```
    break ;
```

:

```
    case value n :
```

executing body n

```
    break ;
```

```
    default :
```

executing body

```
}
```

without break

→ all below executing body will also execute

→ continue cannot be used in switch case

→ can only use in loops

can also have
nested switch
case

not
mandatory

Function -

- program linked with well defined task
- why
 - reusable
 - readable
- without
 - bulky
 - lengthy
 - buggy if mistake in any place

syntax -

```
return type function name ( input parameters ) {
    executing body
}
```

void → empty / no value void x; → ERROR

cout << sizeof(int) → 4

cout << sizeof(void) → ERROR

```
int main(){
    return 0;
}
```

→ returns 0 to Operating System

→ 0 is used as means of successful execution

- a cpp file cant have more than 1 main functions
- main cant have return type other than int in offline compiler

Function Call Stack

function call \leftrightarrow function invoke

Stack

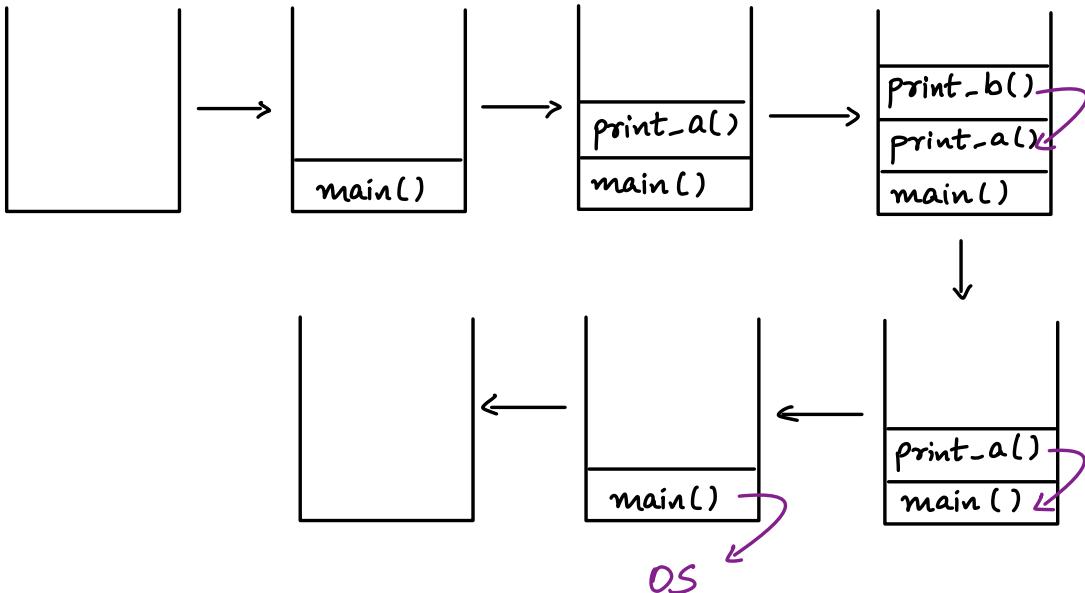
\hookrightarrow Last In First Out

- \hookrightarrow tells what functions are active
- \hookrightarrow which function calls which
- \hookrightarrow local variables of function
- \hookrightarrow return type of function

ex -

```
int main() {
    int a=5;
    print_a(a)
    return 0;
}
} return to OS
```

```
void print_a(int a){
    cout << a;
    int b=3;
    print_b(b);
}
} Output - 5 3
```



Pass by value → value is passed to another variable

↳ a copy will be created of variables

```
int main() {
```

```
    int a=5;
```

```
    printNumber(a);
```

```
    cout << a;
```

```
}
```

A₁

5

a

argument

diff. memory
locations

```
void printNumber(int a){
```

```
    cout << a; → 5
```

```
    a++;
```

```
    cout << a; → 6
```

```
}
```

A₂

5

a

A₂

6

a

new copy

Address Of Operator &

```
int n=5;
```

```
cout << &n;
```

→ output -
address of n

```
int main() {  
    int a=3;  
    int b=4;  
    int sum = add(a,b);
```

```
    cout << sum;  
    return 0;
```

```
}
```

sum

7

A₁

```
int add(int a, int b){  
    int result = a+b;  
    return result;
```

returning a value

7

A₂

copying the
returning value in sum

Function Order

Order 1

```
int add (int a, int b) {  
    return a+b;  
}  
  
int main () {  
    int a= 3;  
    int b = 5;  
    int sum= add (a,b);  
    cout << sum;  
    return 0;  
}
```

function
declaration
and
definition

Order 2

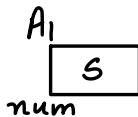
```
function declaration  
{  
int add (int a, int b);  
  
int main () {  
    int a= 3;  
    int b = 5;  
    int sum= add (a,b);  
    cout << sum;  
    return 0;  
}  
  
int add (int a, int b) {  
    return a+b;  
}
```

% operator → heavy operator

↳ so try to use it less

```
int num=5;
```

Symbol table



`int num <--> A1`

variable address mapping is stored

datatype variable name memory location of variable

`cout << num;`

go to symbol table → A₁ (address) → print 5 (value)

`cout << sizeof(num);`

go to symbol table → int datatype → print 4

while loop

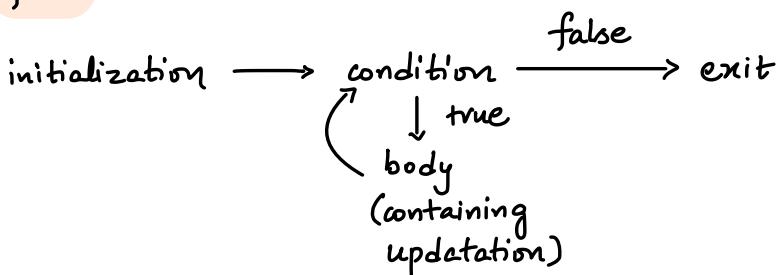
`int i=0;` → initialization

`while (i < 5) {` → condition

`cout << i << endl;` `i++;` → body

`}` → updation

flow



left and right shift operators

int a = 2;

$a \ll 1;$ → no change

$\text{cout} \ll a;$ → 2

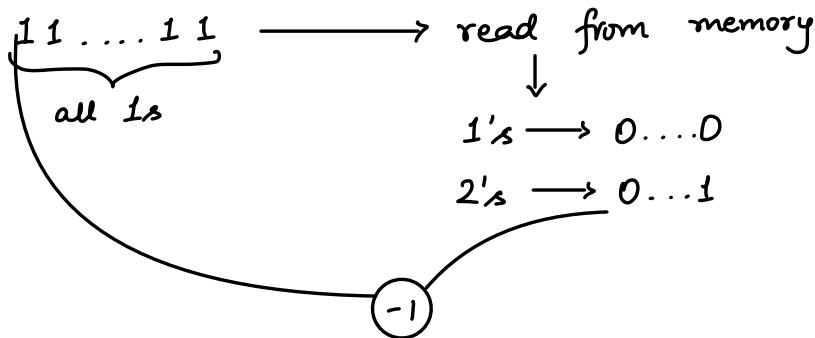
$a = a \ll 1;$ → change → left shift by 1

$\text{cout} \ll a;$ → 4

right shift in -ve no.

↳ link in links.txt in repo

How -1 is stored in memory -



$$\sim a = -(a+1) \quad \text{and} \quad \sim(\sim a) = a$$

ex - $a = 5; \rightarrow 0\dots0101$

$a = \sim a; \rightarrow 1\dots1010 \rightarrow -6 \rightarrow -(5+1)$

$a = -6;$

↳ $1\dots1010$

read

-6

$a = \sim a; \rightarrow 0\dots0101$

↳ 5 → $-(-6+1)$

Number System

↳ method to represent numeric values using digits

Decimal Number System

↳ base 10

↳ digits → 0 to 9

Binary Number System

→ base 2

↳ digits → 0, 1

→ used in CPU, memory, computer

→ 0 → power off

→ 1 → power on

→ number, images, all files & folder are in binary

Decimal to Binary

→ divide no. by 2

→ store remainder

→ repeat above steps until no. is 0 or 1

→ reverse the bits so obtained

Binary to Decimal

- multiply each bit with its place value
 - ↳ base i
- add all products
 - ↳ 2^i

Time & Space Complexity

W3-R

Time Complexity

- amount of time taken by an algo as a function of length of input
- not actual time
- it defines CPU operations
- use case -
 - to make efficient programs
 - ask by interviewer after every sol. you give

Space Complexity

- amount of space taken by an algo as a function of length of input

Units to represent Complexity

Big O \longrightarrow upper bound \longrightarrow worst case

Theta $\Theta \longrightarrow$ average case

Omega $\Omega \longrightarrow$ lower bound \longrightarrow best case

Big O Complexities

$O(1)$ \longrightarrow Constant time

$O(n)$ \longrightarrow Linear time

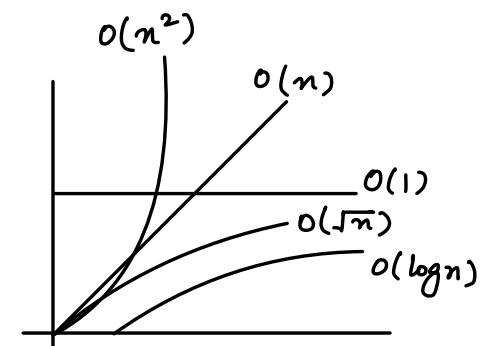
$O(\log_2 n)$ \longrightarrow Logarithmic time

$O(n^2)$ \longrightarrow Quadratic time

$O(n^3)$ \longrightarrow Cubic time

nesting \longrightarrow multiply \longrightarrow { } { }

in sequence \longrightarrow add \longrightarrow { }



$$f(n) = 4n^4 + \frac{n^3}{5} + \log n + n \log n \longrightarrow O(n^4)$$

Complexity Order

$$O(1) < O(\log_2 n) < O(\sqrt{n}) < O(n) < O(n \log_2 n) < O(n^2)$$

$$< O(n^3) < O(2^n) < O(n!) < O(n^n)$$

ARRAY

W3-L1

- Data Structure to store similar items
 - ↳ same datatype
- Continuous memory location
- use case
 - ↳ for multiple same kind of data
`int a[30000];` → 30000 variables are ready

continuous memory location

↳ memory wastage
if needable memory is present but not in continuous way

`int a = 5;`

A



a

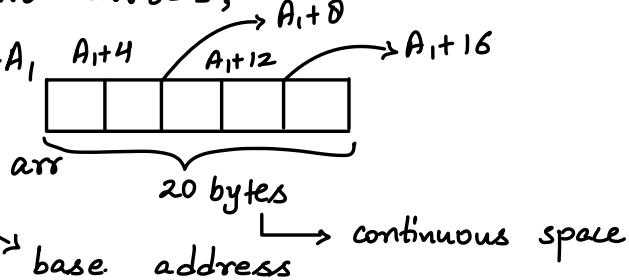
symbol table

`int a ↔ A`

`int arr ↔ A1`

Declaration

`int arr[5];`



`cout << arr ;` —————— A_1 —————— address of first index of array
address of integer

`cout << &arr ;` —————— A_1 —————— base address of array arr
address of array

Initialization

diff. things

`int arr [] = { 2, 4, 6, 8, 10 };`

`int arr2 [5] = { 2, 4, 6, 8, 10 };`

`int arr3 [10] = { 2, 4, 6, 8, 10 };` —————— remaining 5 will be 0

`//int arr4 [4] = { 2, 4, 6, 8, 10 };` —————— ERROR

`int arr5 [10] = { 0 };` —————— initializing all values with 0

Making array at runtime

`int n;`

`cin >> n;`

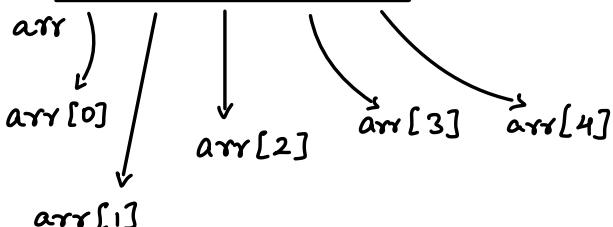
`int arr[n];` —————— BAD PRACTICE

Index and Access in memory

`int arr[5] = { 10, 20, 30 };` —————— 0^{th} based indexing

↳ 0 to $n-1$

A_1	0	1	2	3	4
	10	20	30	0	0



$\text{arr}[i] \longrightarrow \text{value at address } [A_1 + (i * 4)]$

BA
index
due to int (datatype size)

that's why 0 based indexing

taking input in array

$\hookrightarrow \text{cin} >> \text{arr}[i];$

due to internal working

Arrays and Function

$\hookrightarrow \text{func (int arr[], int size) } \{$

}

arr here is not
an array, it is
pointer pointing to
first index of array
(will learn ahead)

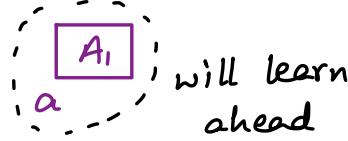
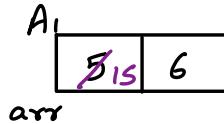
array is passed by reference
because pointer is passed by
value

updation in actual array

always pass size alongwith arr

```
int main() {
    int arr[] = {5, 6};
    int size = 2;
    func (arr, size);
    return 0;
}
```

```
void func ( int a[], int size) {
    a[0] = a[0] + 10;
}
```



; will learn
ahead

`sizeof(int);` → 4 → in bytes

`int arr [5];`

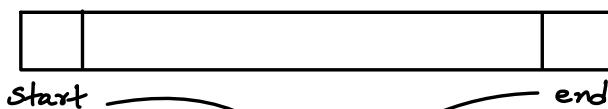
`sizeof(arr);` → 20 → in bytes
→ not 8 (size of pointer)

linear search in array

INT_MIN and INT_MAX

- to find max. , start ans with INT_MIN
- to find min. , start ans with INT_MAX

2 pointer approach



use of 2 variables as extreme points

To find size of array

`int arr [] = { 1, 2, 3, 4 };`

`int size = sizeof(arr) / sizeof(int);`

→ datatype

Vector

W3-L2

- Data structure
- Same as array but dynamic
 - size not fixed
- default size → 0
- if gets full and new items are inserted
 - size gets doubled
- pass by value in functions

Initialization

vector <int> arr {10, 20, 30}; →

10	20	30
----	----	----

vector <int> arr (5); →

0	0	0	0	0
---	---	---	---	---

vector <int> arr (5, -2); →

-2	-2	-2	-2	-2
----	----	----	----	----

int n;
size →

size	value
------	-------

cin >> n; → let n = 5

vector <int> arr(n); →

0	0	0	0	0
---	---	---	---	---

vector <int> arr (n, 10); →

10	10	10	10	10
----	----	----	----	----

Insertion -

arr.push_back(5);

Remove

arr.pop_back();

Size -

arr.size();

→ no. of elements it stores

declaration

vector <int> arr;

→ arr.size() → 0

→ arr.capacity() → 0

Empty or Not

arr.empty(); \longrightarrow true if empty

Capacity -

arr.capacity(); \longrightarrow * by 2 if arr gets fully filled
and a new element is inserted

→ no. of elements it can store

→ in initialization, capacity = size in all methods of initialization

sizeof(arr); \longrightarrow compiler dependent
initially

cout << arr; \longrightarrow give ERROR

→ Xor \longrightarrow cancels out same element

$$0 \wedge \text{ans} = \text{ans} \quad \begin{cases} 0^1 = 1 \\ 0^0 = 0 \end{cases}$$

for each loop

```
for (auto val: arr){  
    cout << val << ' ';  
}
```

2D Arrays

W3-L3

→ array of arrays
→ use case

→ to work on multiple rows and columns

Declaration -

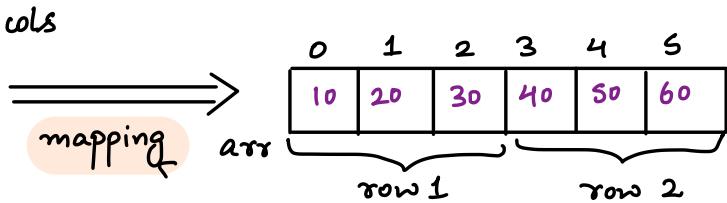
int arr [m][n]; → m*n elements
 | → cols → 0 to n-1
 | → rows → 0 to m-1

int arr [2][3];

visualize -

	cols		
0	10	20	30
1	40	50	60
rows	0	1	2

in memory -



Access -

arr [i][j];
 | → col index $0 \leq j < n$
 | → row index $0 \leq i < m$

Mapping -

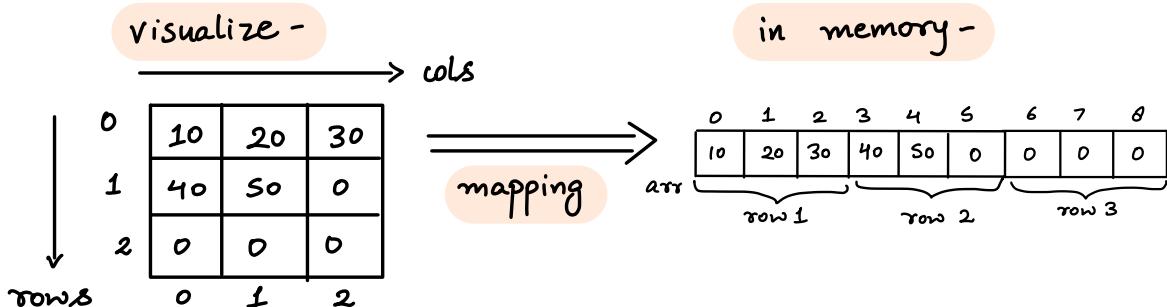
$$\text{linear_index} = c * i + j$$

no. of cols col index
row index

$$i = \text{linear_index} / c$$
$$j = \text{Linear_index \% c}$$

Initialization -

int arr[3][3] = {{10, 20, 30}, {40, 50}}; same
int arr[3][3] = {10, 20, 30, 40, 50};



2D Arrays and function -

→ pass by reference

func (arr[] [500], int rows, int cols)

actually arr
here is
int (*)[500]
(array of pointers),
not a 2D array

→ cannot leave blank

why → if dont know, put large value

500 tells size of array of pointers

→ this value and no. of cols
in array passed in function

call should be same

2D Array -

→ not an array of pointers

→ array of arrays

```

func( int a [ ] [ 3 ], int rows , int cols ) {
    cout << & a ; ----- add. of a -----> A4
    cout << a ;
    cout << a [ 0 ];
    cout << & a [ 0 ];
    cout << & a [ 0 ] [ 0 ];
    cout << a [ 0 ] [ 0 ]; -----> 10
    cout << sizeof( a ); -----> 4
    cout << sizeof( a [ 0 ]); -----> 12
}

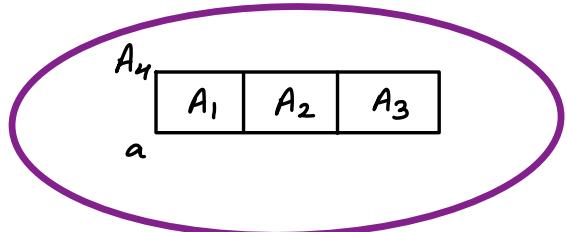
```

```

int main () {
    int arr [ 2 ] [ 3 ] = { 10, 20, 30, 40, 50 };
    func( arr, 2, 3 );
    return 0;
}

```

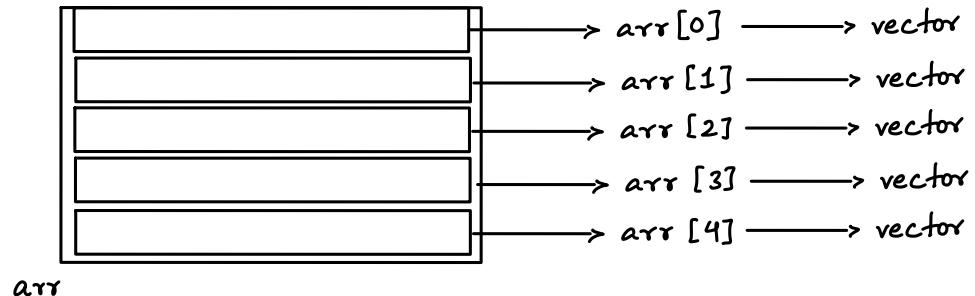
	A ₁	A ₂	A ₃
A ₁	10	20	30
A ₂	40	50	0
A ₃	0	0	0



2D Vector

Declaration -

```
vector<vector<int>> arr;
```



Declaration -

```
vector<vector<int>> arr;  
vector<vector<int>> arr(m);
```

Number of rows

arr.size()

Number of cols

arr[i].size()

→ in ith row
→ size of ith row

Initialization

```
vector<vector<int>> arr (rows, (vector<int>(cols, value)));
```

rows → no. of rows in arr

cols → no. of cols in arr

→ size of 1D vector

size initial value in 2D
 ↑ vector

initialization of 1D
vectors in arr

value → initial value in all elements

of all 1D vectors

```
vector<vector<int>> arr(2, vector<int>(4, 101));
```

101	101	101	101
101	101	101	101

arr

auto keyword -

→ automatically replace with required data type

```
int arr [4]
```

```
for( auto i = 0; i < 4; i++)
    cout << arr [i];
```

: Operator

→ belongs to operator

→ generally used for sequential access

→ map, set, array, vector

```
ex - vector<int> arr {1, 2, 3};
```

```
for( int i : arr)
    cout << i;           → 1 2 3
```

```
for( auto i : mapping)
    cout << i;
```

```
for( auto i : sett)
    cout << i;
```

Searching and Sorting

W4-L1

Searching

Linear Search

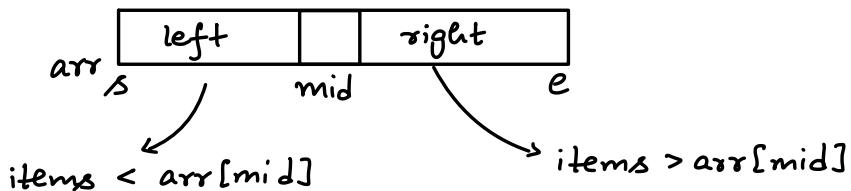
```
int linearSearch ( vector <int> arr, int target ) {  
    int n = arr.size();  
    for( int i=0; i<n ; i++ ) {  
        if ( target == arr[i] )  
            return i;  
    }  
    return -1;  
}
```

T.C - $O(n)$

S.C. - $O(1)$

Binary Search -

- condition → sorted order → monotonic function
- binary → 2 → left and right search



```

int binarySearch ( int arr[], int n , int target ) {

    int s= 0, e=n-1;

    int mid = s+(e-s)/2;

    while( s<=e) {

        int element = arr[mid];

        if (target == element)
            return mid;

        else if ( target < element)
            e=mid - 1; —→ search in left subarray

        else
            s= mid + 1; —→ search in right
                           subarray

        mid = s + (e-s)/2;
    }

    return -1;
}

```

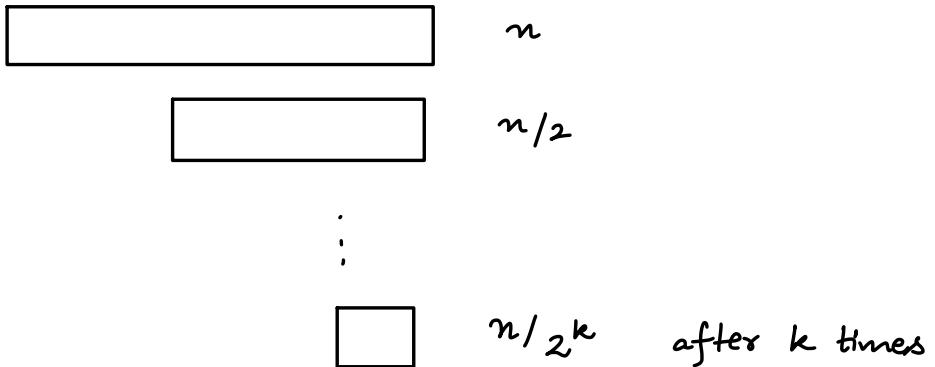
issue in $mid = \frac{s+e}{2}$; —→ int overflow
 if $s+e > INT_MAX$

└—————> so use $mid = s + \frac{e-s}{2}$;

T.C - $O(\log_2 n)$

S.C . - $O(1)$

T.C. of binary Search



at last $\frac{n}{2^k} = 1$

$$k = \log_2 n$$

So loop runs $\log_2 n$ times

$$\text{T.C.} = O(\log_2 n)$$

W4-L2

→ Search Space

 └ find range of search space (start & end)
 in ques of. Binary Search

→ store mid in ans if needed

→ In binary search questions

```
while (s <= e) {  
    ans = mid;  
    s = mid + 1;  
}  
left  
search ← e = mid - 1;  
}  
  
right  
search
```

convert

```
ans = mid  
s = mid + 1  
left  
search ← e = mid - 1;
```

OR

```
ans = mid, e = mid - 1  
left  
search ← e = mid;  
s = mid + 1 → right  
search
```

cout << (int)(-3.74);

W4-L3

→ - [int(3.74)]

→ - (3) → - 3

cout << (-22)/7; → - 3

cout << 22/(-7); → - 3

Types of ques in binary search

→ 1st type → classic questions

→ lower bound

upper bound

peak in mountain array

can also find array is sorted or not

pivot in sorted rotated array

search in sorted rotated array

↓ pivot index = n-1

→ 2nd type → find in search space (range)

- predicate function
- logic to decide either left or right
- sqrt of a no.
- divide 2 numbers

Advance Binary Search Problems

→ Book allocation

Painters Partition

Aggressive Cows

Roti / Paratha Spoj

Eko Spoj

→ 3rd type → observation in index

- missing element in sorted array
- add appearing element in array

Sorting

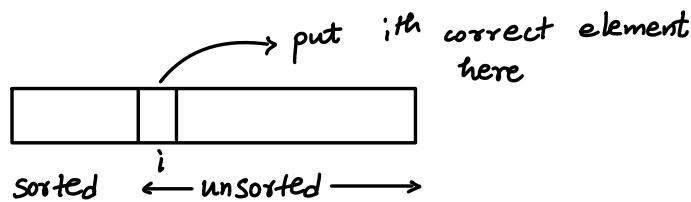
W4-L5

- putting all elements either in increasing order or decreasing order

Selection Sort -

- select minimum element and put it in its right position
- select correct element for i^{th} index
 - ↳ i^{th} minimum element

```
void selectionSort ( int arr[], int n) {  
    for( int i=0; i<n-1; i++) {  
        mini = i;  
        for( int j=i+1 ; j<n ; j++) {  
            if ( arr[mini] > arr[j] )  
                mini = j;  
        }  
        swap (arr[i], arr[mini]);  
    }  
}
```



T.C. $O(n^2)$

S.C. $O(1)$

Use Case - for small size array

- ↳ put i^{th} lowest element to its correct position

Bubble Sort -

→ in i^{th} round put i^{th} largest element to its correct position using adjacent comparisions

```
void bubbleSort ( int arr[], int n){  
    for( int i=0 ; i<n-1 ; i++) {  
        bool swapping = false;  
        for( int j= 0 ; j<n-i-1 ; j++) {  
            if (arr[j] > arr[j+1]) {  
                swap ( arr [j], arr [j+1]);  
                swapping = true;  
            }  
        }  
        if (swapping == false)  
            break;  
    }  
}
```

T.C. - $O(n^2)$ → reverse sorted
Worst and average case

$O(n)$ → best case → already sorted

S.C. - $O(1)$

Use Case - To put i^{th} largest element to its correct position

Insertion Sort -

→ take an element and insert it on its correct position by shifting

```
void insertionSort( int arr[], int n){
```

```
    for( int i = 1; i < n; i++ ) {
```

```
        int curr = arr[i];
```

```
        int j = i - 1;
```

```
        for( ; j >= 0; j-- ) {
```

```
            if( arr[j] > curr )
```

```
                arr[j+1] = arr[j]; // shifting
```

```
            else
```

```
                break;
```

```
}
```

```
        arr[j+1] = curr; // inserting
```

```
}
```

```
}
```

T.C. - $O(n^2)$ → worst & average case

$O(n)$ → best case

S.C. - $O(1)$

Use Case - When array is small or when array is partially sorted

Inbuilt sort function

- sort (arr.begin(), arr.end());
- algo used is Intro sort
 - hybrid of quick sort, heap sort, insertion sort
- min. time than any of other sort

T. C. - $O(n \log n)$

S. C. - Not Defined

Stable and Unstable Algorithm

Stable → order preserve after sorting
 ↳ of same values

Stable Sorting Algo

2 1 **2** 2 3
 ↓
 after sorting

if $A[i] = A[j]$, $i < j$
then $A[i]$ comes first before
 $A[j]$ after sorting too

1 2 **2** 2 3

Stable - Bubble Sort, Insertion Sort, Merge Sort,
Count Sort

Other (unstable) sorting algorithms can be made stable
by some changes

T.C. and S.C Sorting Algorithm Table

Algo	S. C.	Worst T.C.	Avg. T.C.	Best T.C.
Selection	$O(1)$	$O(n^2)$	$O(n^2)$	$O(n^2)$
Bubble	$O(1)$	$O(n^2)$	$O(n^2)$	$O(n)$
Insertion	$O(1)$	$O(n^2)$	$O(n^2)$	$O(n)$
sort func. (intro sort)	not defined	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
merge	$O(n)$	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
quick	$O(n)$	$O(n \log n)$	$O(n \log n)$	$O(n^2)$

CHAR ARRAYS & STRING

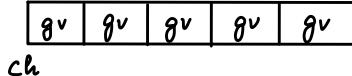
W5-L1

Char Arrays -

→ Data structure → used to store data

Not Datatype → tells type of data

char ch[5];



char ch[] = {'U', 'j', 'j', 'w', 'a', 'l'}
ch

char ch[] = "Ujjwal";

ch

Taking input in char array

char ch[7];

cin >> ch[i];

cin >> ch; → by default NULL char will

insert at end

↙ ↘

'\0'

NULL char refers string termination

→ cin reads until it gets any white space

↳ space ''
tab '\t'
endl '\n'

cout << ch; → cout print until it gets delimiter

```
char ch [10];
```

cin >> ch ; → Ujjwal

store in buffer memory

'U'	'j'	'j'	'w'	'a'	'l'	'\0'
-----	-----	-----	-----	-----	-----	------

buffer

copy in memory of ch

'U'	'j'	'j'	'w'	'a'	'l'	'\0'	gv	gv	gv
-----	-----	-----	-----	-----	-----	------	----	----	----

ch 0 1 2 3 4 5 6 7 8 9

cout << ch ; → Ujjwal → stops after 'l'

because of '\0' char

Overflow -

```
char ch [4];
```

cin >> ch ; → Ujjwal

cout << ch ; → Compiler Dependent

sizeof (char array)

```
char ch [] = "Ujjwal";
```

cout << sizeof(ch) ; → 7

6 + 1

NULL char

get line

`cin.getline (char array, max char to write, delimiter);`

char where taking input stops
↓
by default '\n' → enter

ex - `cin.getline (ch, 50);`

`cin.getline (ch, 50, '');`

Char arrays and function

→ pass by reference

`func (char ch[]){`

}

Size of char array → $\frac{\text{sizeof}(ch)}{\text{sizeof}(char)}$

→ 1

Some inbuilt functions of char array

`strlen(ch);`

`strcmp(ch1, ch2);`

`strcpy(ch1, ch2);`

String 8

- Datatype
- Not Data Structure
- Dynamic char array
- NULL char at last of string

string str; → empty string created

cin >> str; → Ujjwal

['U' | 'j' | 'j' | 'w' | 'a' | 'l' | '\0']

cout << str; → Ujjwal str

cout << str.length(); → 6, not 7

getline -

string str;

getline (cin, str);

char array

char ch [] = "B_abba-\r";

cout << ch; → B_abba-\r

cout << arr.size () → 9

ch[1] = '\0';

ch[6] = '\0';

cout ch; → B

stops just as it gets NULL char

string

string str = "B_abba-\r";

cout << str; → B_abba-\r

cout << str.length () → 8

cout << str.size () → 8

str[1] = '\0';

str[6] = '\0';

cout << str; → Babbar

Runs till the length of string

Sort function in strings -

W5-L3

```
string str = "babbar";
```

```
sort(str.begin(), str.end()); -----> aabbbr
```

```
sort (str.begin(), str.end(), greater<char>());
```

-----> rbbbbaa

Custom Comparator -

```
bool cmp (char a, char b) {
```

return a < b; -----> can be any function
according to need

```
}
```

```
bool cmp2 (char a, char b) {
```

return a > b;

```
}
```

```
int main () {
```

string str = "babbar";

```
sort (str.begin(), str.end(), cmp);
```

cout << str; -----> aabbbr

string str2 = "babbar";

```
sort (str.begin(), str.end(), cmp2);
```

cout << str2; -----> rbbbbaa

return 0;

```
}
```

Hash Map → WILL LEARN LATER

→ Data structure

→ data stored in key-value pair

Initialization

`map<key datatype, value datatype> map-name;`

`map<int, char> m;` → ordered map

`m[0] = 'a';`

`m[1] = 'b';`

`m[25] = 'z';`

`cout << m[0];` → 'a'

`cout << m[25];` → 'z'

`cout << m[20];` → " " → NULL char

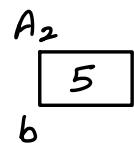
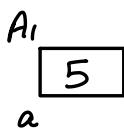
`cout << (int) m[20];` → 0

SYMBOL TABLE

W6-L1

- Data Structure
- Stores mapping of variable name, datatype & memory location address → Done by OS
- entries in symbol table can't be changed

int a = 5;
int b = 5;



Symbol Table
name → address
a A₁
b A₂

→ entry of a in symbol table is made

cout << a; → 5

& Address Of Operator

cout << &a; → A₁ → address of

cout << &b; → A₂

hexadecimal
↑
variable a in
memory

Not 100%
Sure

{ A₂ = A₁ ± 4 → because of consecutive
variables in memory
↓
int datatype
of a and b

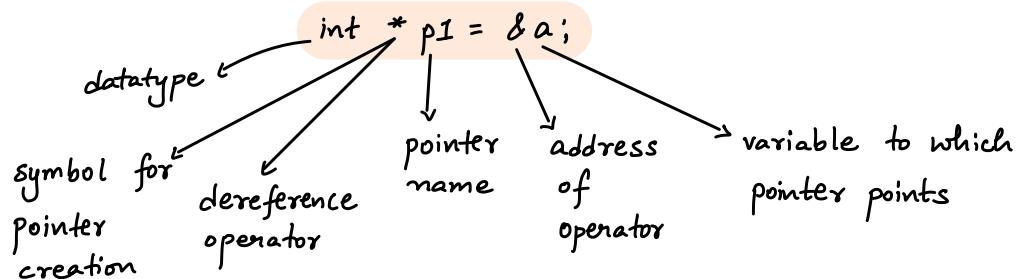
POINTERS

→ stores address

→ NOT a datatype, just a variable storing address of another variable

```
int a = 5;
```

`int * p1 = &a;` p1 is a pointer to integer datatype



```
string s = "Ujjwal 2327";
```

`string * ptr = &s` → ptr is a pointer to string datatype

```
int a = 5;
```

```
int * ptr = &a;
```

```
cout << a; → 5
```

```
cout << * a; → ERROR
```

```
cout << &a; → A1
```

```
cout << ptr; → A1
```

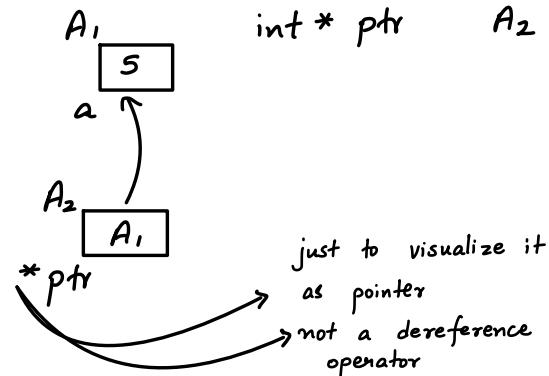
```
cout << * ptr; → 5
```

```
cout << &ptr; → A2
```

symbol table

int a A₁

int * ptr A₂



$* \text{ptr}$ → value at (address stored in ptr)
→ dereference operator

$\& \text{ptr}$ → address of ptr

size of pointer

$\text{sizeof}(\text{ptr})$; → 8
64 bit architecture

depends on
architecture
compiler implementation
memory organization

Always
it stores address,
datatype does not matter

use case of pointer

- dynamic memory allocation
- memory management
- pointer arithmetic
 - ↳ go from one location to other
- pass by reference in array
- to create pointer to function
 - ↳ passing a function inside another function as an argument

int * ptr;
cout << ptr;

VERY BAD PRACTICE

Segmentation fault

A
 $* \text{ptr}$
gv

ptr points to a memory location that may not be of its program

Segmentation fault -

→ using other's memory

NULL Pointer

```
int * ptr = 0;
```

```
int * ptr2 = NULL;
```

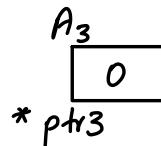
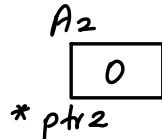
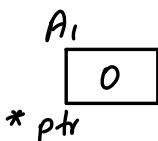
```
int * ptr3 = nullptr;
```

```
cout << ptr ; → 0
```

```
cout << *ptr ; → ERROR
```

```
cout << &ptr ; → A1
```

Same



Segmentation fault

Arithmetic Operations In Pointers

```
int a=5;
```

```
a++;
```

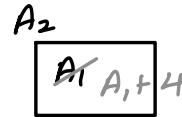
```
int * ptr = &a;
```

```
*ptr++;
```

```
ptr++;
```



a



*ptr

A₁ to A₁+3 has already be taken by integer a, so next address will be A₁+4

int a = 10;

int *ptr = &a;

a 10

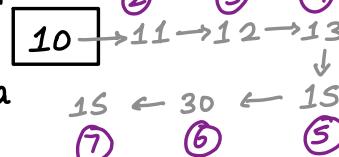
&a A₁

ptr A₁

*ptr 10

&ptr A₂

A₁



① *ptr * 2

② (*ptr) ++

③ ++ (*ptr)

④ a = a + 1

⑤ *ptr = *ptr + 2 15

⑥ *ptr = *ptr * 2 30

⑦ *ptr = *ptr
2 15

⑧ *(ptr++) 15

⑨ *(++ptr) 9v

int a = 5;

int *ptr = a; → ERROR

Copying a pointer

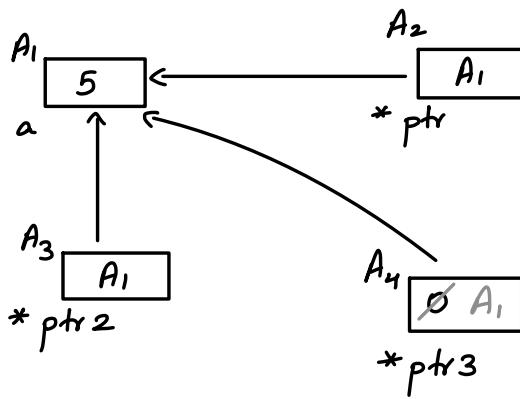
int a = 5;

int *ptr = &a;

{ int *ptr2 = ptr;

{ int *ptr3 = 0;

{ ptr3 = ptr;



ARRAYS & POINTERS

W6-L2

`int arr[] = {10, 20, 30, 40, 50};`

→ CONSTANT pointer to the first element of array

cannot change

→ also does not have separate memory

`arr[i]` → element at index i

$\& \text{arr}[i]$ → address of $\text{arr}[i]$

`cout << arr[0];` \longrightarrow 10

`cout << &arr[0];` → A₁

{ cout << arr ; } → A1

```
cout << &arr;
```

array name returns add. of
1st index of array

1st index of array

→ from symbol table (add. of array)

same address unlike pointers

int *ptr = arr ;

`cout << * ptr;` → 10

`cout << ptr ;` \longrightarrow A

`cout << &ptr;` \longrightarrow `A_2`

`cout << *arr ;` \longrightarrow 10

`cout << *(&arr);` ——————> 10

`cout << *(& arr[0]);` —> 10

Symbol table

type	name	address
int (*)[5]	arr	A ₁
int *	ptr	A ₂

1000

`cout << *arr + 1;` \longrightarrow 11

`cout << *(arr + 1);` \longrightarrow 20

`cout << *(arr + 2);` \longrightarrow 30

`cout << arr[2];` \longrightarrow 30

`cout << 2[arr];` \longrightarrow 30

`arr[i] \longleftrightarrow *(arr + i) \longleftrightarrow i[arr]`

`arr ++;` \longrightarrow ERROR \longrightarrow Entry in symbol table

`ptr ++;` \longrightarrow A₁ + 4 \longrightarrow cant be change

→ You can access subpart of an array using pointers

`int *ptr = arr;`

`cout << *(ptr + 2);` \longrightarrow 40

`cout << *(ptr + 100);` \longrightarrow gv / segmentation fault /
out of bound error

Arrays / Array pointers

`int arr [] = {10, 20, 30};`

① `cout << arr;` \longrightarrow A₁ } same
`cout << &arr;` \longrightarrow A₁

Pointers / Normal pointers

`int *ptr = arr;`

① `cout << ptr;` \longrightarrow A₁ } diff.
`cout << &ptr` \longrightarrow A₄

② `arr ++;` \longrightarrow ERROR

② `ptr ++;` \longrightarrow VALID

③ `sizeof(arr);` \longrightarrow 3 * 4 = 12

③ `sizeof(ptr)` \longrightarrow 8
↓
size of address

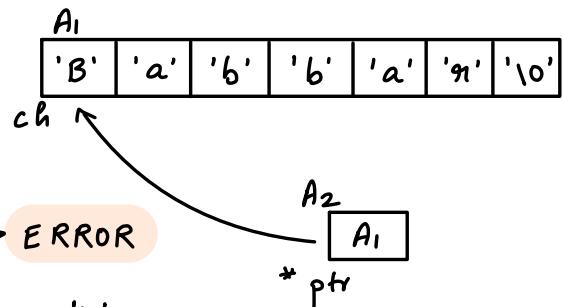
```
int * p1 = arr;
```

```
int * p2 = &arr; -----> ERROR
```

```
int * p3 = &arr[0];
```

CHAR ARRAYS & POINTERS

```
char ch [] = "Babbar";
```



```
char * ptr = ch;
```

```
char * p = &ch; -----> ERROR
```

```
char * p2 = &ch[0]; -----> valid
```

```
cout << ch; -----> Babbar
```

not an address

```
cout << ptr; -----> Babbar
```

not an address

whole string from that
location until NULL char

cout implementation
is diff. in char
pointers and char arrays

```
cout << &ch; -----> A1
```

```
cout << ch[0]; -----> B
```

```
cout << &ch[0]; -----> Babbar
```

```
cout << *ptr; -----> B
```

```
cout << &ptr; -----> A2
```

$$ch[i] \longleftrightarrow * (ch + i) \longleftrightarrow i[ch]$$

→ ch, &ch[0], and ptr values are addresses
but due to diff. cout implementation in char pointers
and char arrays,
cout << ch , cout << &ch[0] and cout << ptr
will give Babbar

char ch[] = "Sherbano";

char *ptr = ch;

cout << ch ; → Sherbano

cout << *ch ; → S

cout << &ch ; → A₁

cout << *(ch + 3); → n

cout << ptr ; → Sherbano

cout << &ptr ; → A₂

cout << *(ptr + 3); → n

cout << ptr + 2; → erbano

cout << *ptr ; → S

cout << ptr + 8; → " → NULL char

cout << ptr + 9; → gv

`cout << ch[0];` → S

`cout << &ch[0];` → Sherban

`cout << &(*ch);` → Sherban

CHAR AND POINTER

`char ch = 'k';`

`char * ptr = &ch;`

`cout << ch;` → k

`cout << &ch;` → k.....

→ gv
print until it gets '\0'

`cout << ptr;` → k.....

`cout << &ptr;` → A₂ → gv
print until it gets '\0'

`cout << *ptr;` → k

Behind The Scenes

`char ch[10] = "Babber";`

→ 2 step process

① |'B'| 'a'| 'b'| 'b'| 'a'| 'g'| '\0'|

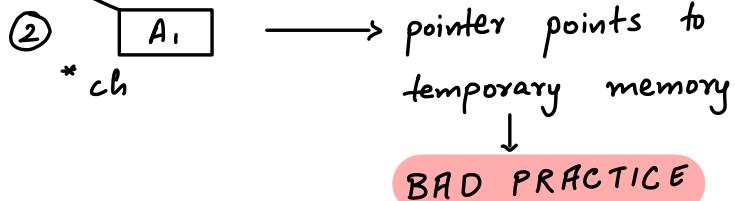
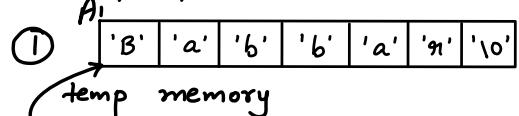
temp memory } copy

② |'B'| 'a'| 'b'| 'b'| 'a'| 'g'| '\0'| gv | gv | gv |

ch

```
char * ch = "Babbar"
```

→ 2 step process



POINTERS WITH FUNCTIONS

→ pass by value

a copy of pointer is made

→ but simulates pass by reference

→ as in case of arrays & functions

→ pointer is passed

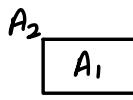
```
int func ( int arr[]){
```

cout << arr; → A1

cout << *arr; → 10

cout << &arr; → A2

cout << sizeof(arr); → 8



copy of pointer
is created

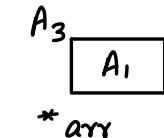
int func2 (int * arr) {

cout << arr; → A₁

cout << * arr; → 10

cout << &arr; → A₃

cout << sizeof(arr); → 8

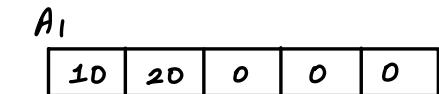


copy of pointer
is created

}

int main() {

int arr [5] = { 10, 20 };



cout << arr; → A₁

cout << * arr; → 10

cout << &arr; → A₁

cout << sizeof(arr); → 5 * 4 = 20

func (arr);

func (arr 2);

return 0;

}



→ whole array will not pass

→ only array pointer / base address will pass

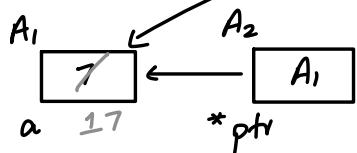
→ pass by reference for array

→ copy of pointer pointing to base address will be created

```

int main() {
    int a = 7;
    int * ptr = &a;
    update (ptr);
    return 0;
}

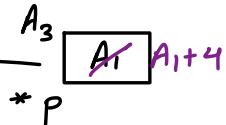
```



```

void update( int * p) {
    ① *p = *p + 10;
    ② p++;
}

```



$$\begin{aligned}
 *(&A_1) &= *(A_1) + 10 \\
 *(&A_1) &= 7 + 10 = 17 \\
 a &= 17
 \end{aligned}$$

Function to Pointers in HW

Basic Mathematics for DSA

W6-L3 R

Sieve of Eratosthenes Theorem

→ to find no. of prime numbers between 1 & n

Steps-

- make an array of size n and mark them all as primes from 2 to n
- Start from 2 till end, mark all no. > 2 comes in the table of 2 as non prime
- Do above step for numbers 2 to $<n$ if they are marked prime
- Count all remaining marked prime numbers

```

int countPrimes ( int n) {
    if (n <= 1)
        return 0;

    vector <int> isPrime (n, true);
    isPrime [0] = isPrime [1] = 0; // making both 0

```

```

    int ans = 0;
    for( int i=2 ; i<n ; i++){
        if ( isPrime [i]){
            ans++;
            for( int j= 2*i ; j<n ; j+= i)
                isPrime [j] = false;
        }
    }
    return ans;
}

```

i ≠ i < n
↳ to optimize T.C.
as 2i to (i-1)*i are
already marked when
i = 2 to (i-1)

$$TC - O(n * \log(\log n)) \quad SC - O(n)$$

$$n \left[\frac{n}{2} + \frac{n}{3} + \frac{n}{5} + \frac{n}{7} + \frac{n}{11} + \dots \right]$$

$$n^2 \left[\frac{1}{2} + \frac{1}{3} + \frac{1}{5} + \frac{1}{7} + \frac{1}{11} + \dots \right]$$

$$n \log(\log n)$$

Same with **segmented sieve** → having high & low

GCD/HCF and LCM

→ Greatest Common Divisor

$$\rightarrow \gcd(a, b) = \gcd(a - b, b) \quad a > b$$

$$\gcd(a, b) = \gcd(b - a, a) \quad b > a$$

$$\gcd(a, b) = \gcd(a \% b, b) \quad a > b$$

$$\gcd(a, b) = \gcd(b \% a, a) \quad b > a$$

takes more time bcz
of \% op.

Use first method as \% is heavy operator
and computer takes more time

→ Euclid Algorithm

Apply above formula till one of the parameter becomes 0

And other one will be GCD

```
int gcd( int a, int b ) {
```

```
    if (a == 0)
```

```
        return b;
```

```
    else if (b == 0)
```

```
        return a;
```

```
    while (a > 0 && b > 0) {
```

```
        if (a > b)
```

```
            a = a - b;
```

```
        else
```

```
            b = b - a;
```

```
}
```

}

```
return (a == 0) ? b : a;
```

$$TC = O(\min(a, b))$$

LCM

$$\text{lcm} * \text{gcd} = a * b$$

$$\text{lcm} = \frac{a * b}{\text{gcd}}$$

Modulo Arithmetic

$$\rightarrow a \% n \longrightarrow [0, n)$$

$$\rightarrow (a + b) \% n = (a \% n) + (b \% n)$$

$$(a - b) \% n = (a \% n) - (b \% n)$$

$$(a * b) \% n = (a \% n) * (b \% n)$$

$$(\dots ((a \% n) \% n) \dots \% n) = a \% n$$

Fast Exponentiation

$$\rightarrow a^b = a^{b/2} * a^{b/2}, \quad b \text{ is even}$$

$$a^b = [a^{b/2} * a^{b/2}] * a, \quad b \text{ is odd}$$

```
int fastExponentiation ( int a, int b){  
    int ans = 1;  
    while (b){  
        if (b & 1)  
            ans = ans * a;  
        a = a * a ;  
        b >>= 1; // b = b >> 1 or b = b / 2  
    }  
    return ans  
}
```

dry run code on 2^5 , if confusion

T.C. - $O(\log n)$

Learn wild, void and dangling pointers from dashboard
after learning dynamic allocation

MULTI LEVEL POINTER

W6-L4

int a = 5;

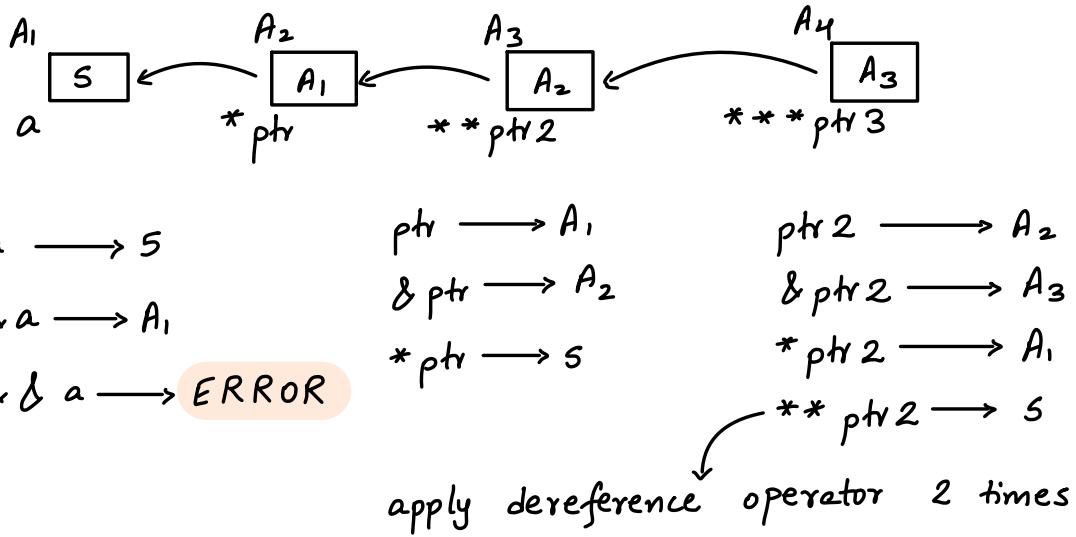
int *ptr = &a;

int **ptr2 = &ptr → double pointer

→ ptr2 is a pointer to int * data

int ***ptr3 = &ptr2;

→ ptr3 is a pointer to int ** data



$ptr3 \rightarrow A_3$

$***ptr3 \rightarrow A_1$

$\&ptr3 \rightarrow A_4$

$***ptr3 \rightarrow 5$

$*ptr3 \rightarrow A_2$

```

int main(){
    int a = 5;
    int * ptr = &a;
    int ** ptr2 = &ptr;
    func( ptr2 );
    return 0;
}

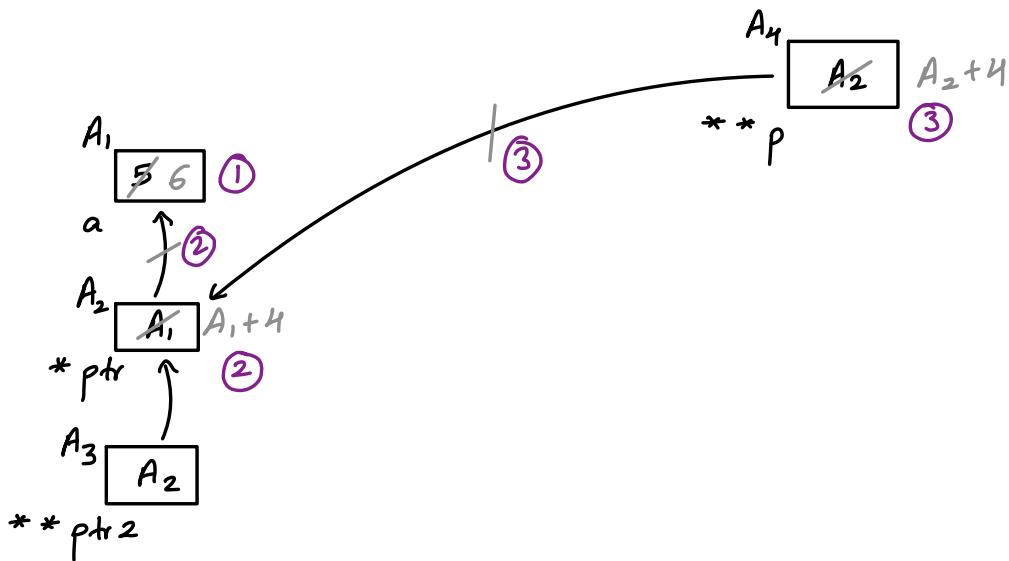
```

```

void func( int *** p ){
    ① (** p) ++ ;
    ② (* p) ++ ;
    ③ p ++ ;
}

both are same
func( &ptr );

```



```

int a = 5;
int * p = &a;
int ** q = p; -----> ERROR

```

REFERENCE VARIABLE

- alternate of pointers
 - $a, *ptr, **ptr \longrightarrow a$
diff. names of same memory location
 - very confusing
- diff. names of same variable
 - same memory location
- only new entry in symbol table,
no other memory will allocate
- can access a variable by diff. names

int a = 5;

A₁

Symbol Table

a	A ₁
b	A ₁

int &b = a;

5
a, b

Use Case

- reference variable can't be set to NULL
pointers can be set to NULL
So more safety in reference variable
Always points to valid object / variable
- pointers are difficult to understand,
more readability in reference variable
- generally used to implement **PASS BY REFERENCE** concept

PASS BY REFERENCE

- reference variable passes in function
- does not create copy

```
int main () {  
    int a = 5;  
    update (a);  
    update2 (&a);  
    int *ptr = &a;  
    //update3 (&a); → ERROR  
    update3 (ptr);  
    return 0;  
}
```

```
void update( int &x){
```

① $x++;$

}

PASS BY
REFERENCE

```
void update2( int * p){
```

② $(*p)++;$

}

PASS BY
VALUE

A₂

A₁

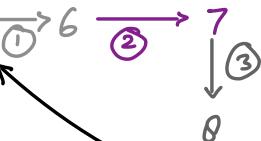
*p

A₁

5

a, x

↑



A₃

A₁

*ptr

A₄

A₁

*pt

```
void update3( int * &pt){
```

③ $(*pt)++;$

}

```

vector<int> func( vector<int> &arr) {
    for (int i=0; i < arr.size(); i++)
        arr[i]++;
    return arr;
}

```

int solve (int & arr) { this is reference variable
of type int

```

} cout << arr;

```

```

int main () {

```

```

vector<int> arr { 1, 2, 3, 4 };

```

```

vector<int> ans = func(arr);

```

```

for ( auto i: ans)

```

```

cout << i ; → 2 3 4 5

```

solve(arr); → ERROR vector != int

```

int arr2 [ ] = { 1, 2, 3 };

```

solve(arr2); → ERROR int[3] != int

```

return 0;
}

```

RETURN BY REFERENCE

→ return a variable , not a value

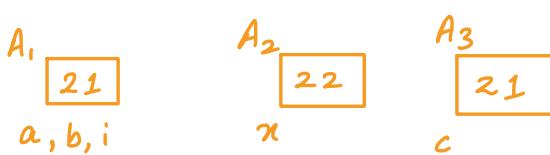
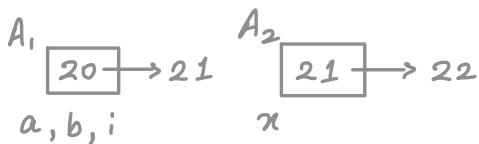
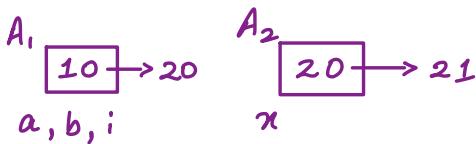
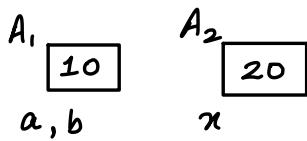
→ To implement atleast 1 of 2 conditions
must be true.

1. passing a reference variable
2. passing a global variable

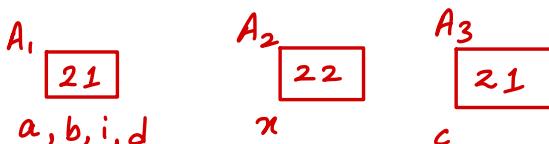
passing a reference variable

```
int main (){  
    int a = 10;  
    int & b = a;  
    int x = 20;  
    ref(a) = x;  
    x++;  
    ref(b) = x;  
    x++;  
    int c = ref(a);  
    int &d = ref(a);  
  
    return 0;  
}
```

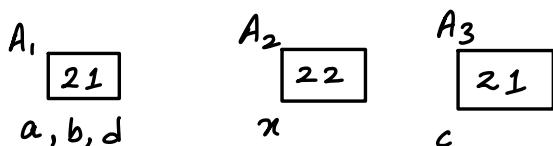
```
int & ref( int &i){  
    return i;  
}
```



`int c = ref(a) = i;`
`int c = 21;`



`int &d = ref(a) = i;`
`int &d = i;`



passing a global variable

```
int x = 5;
```

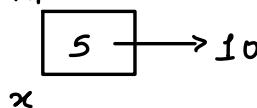
```
int main () {
```

```
    ref () = 10;
```

```
    x = 10;  
    ref2 () = 20; → ERROR  
    return 0;
```

```
}
```

A,



x

```
int & ref () {
```

```
    return x;
```

```
}
```

```
int ref2 () {
```

```
    return x;
```

```
}
```

returning a value, not variable

10 = 20;

↳ that's why ERROR

pass by value but return by reference

```
int main () {
```

```
    int a = 10;
```

```
    int & b = a;
```

```
    int x = 20;
```

```
    ref(a) = x;
```

```
    x++;
```

```
    ref(b) = x;
```

```
    x++;
```

BAD PRACTICE

ERROR

```
int & ref (int temp) {
```

```
    return temp;
```

```
}
```

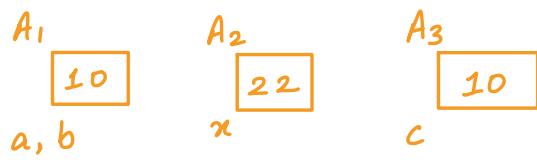
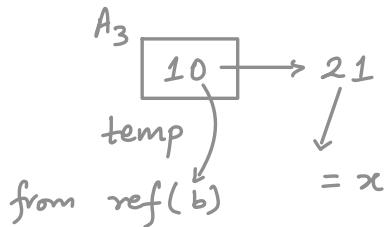
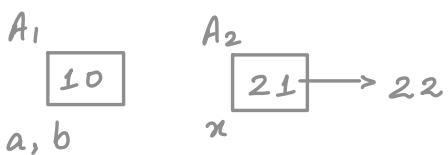
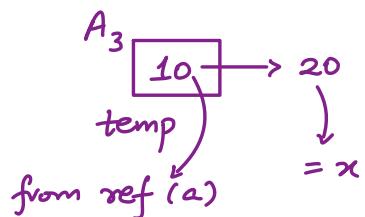
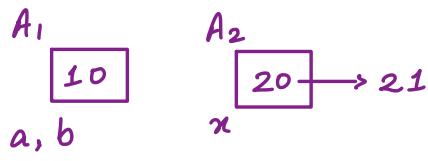
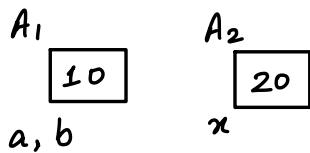
returning a variable stored
in temporary memory

```

int c = ref(a);
int &d = ref(a);
return 0;
}

```

lets assume it will run , then this would happen



```

int c = ref(a) = temp;
int c = temp;
int c = 10

```



d (variable of main function refers to
a temporary memory

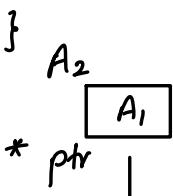
→ BAD PRACTICE

DON'T DO THIS
AT HOME

Have similarity

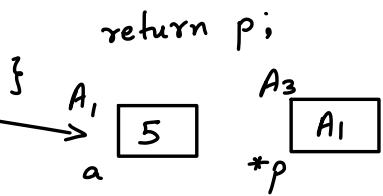
IMP QUESTION

```
int main () {
    int * ptr = solve ();
    return 0;
}
```



pointing to a temp
memory, 'a' variable
will finished outside
solve function so value
of 'a' can be change with time

```
int * solve () {
    int a=5;
    int * p = &a;
    return p;
}
```



→ BAD PRACTICE

Some Important Questions

W6-A

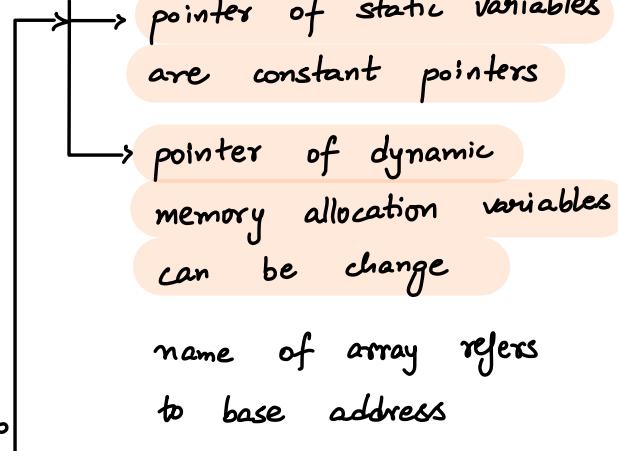
→ `int *ptr = 0;`
`int a = 10;`
`*ptr = a; -----> runtime ERROR`
`cout << *ptr;`

→ `char *ch = 'a';`
`char *ptr = &a;`
`ch++;`
`cout << *ptr; -----> 'b'`

→ `int a[] = {1, 2, 3, 4, 5};`

`int *p = a++; -----> ERROR`

`cout << *p;`



→ `char ch = "hello";`
`cout << ch; -----> hello`
`ch++; -----> ERROR`
`cout << ch;`

name of array refers
to base address

→ double arr [] = { 2.5, 7.9, 50.25 } ;

double * ptr = arr ;

ptr = ptr + 0.25 ; → ERROR

cout << ptr ;

pointer arithmetic op-
can only be done
using integers

→ int arr [] = { 1, 2, 3, 4, 5 } ;

int * ptr1 = arr ;

int * ptr2 = arr + 3 ;

cout << ptr2 - ptr1 ; → 3, NOT 12

→ int a = 5 ;

int * ptr = &a ;

int ** ptr2 = &ptr ;

ptr2 = &a ; → ERROR → int ** cant convert

cout << *ptr2 ; to int *

→ int a = 5 ;

int * p = &a ;

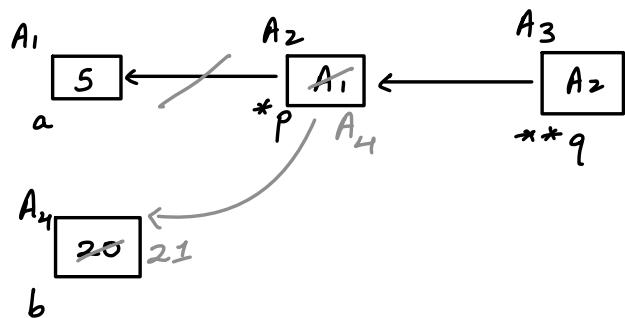
int ** q = &p ;

int b = 20 ;

* q = &b ;

(*p)++ ;

cout << a << ' ' << b ; → 5 21



→ int a = 5;

int * ptr = &a;

int ** ptr2 = &ptr;

int ** ptr3 = &(&a); → ERROR

cout << *** ptr3;

address of (address in
memory location)

(doesn't make any sense)

W6-L5

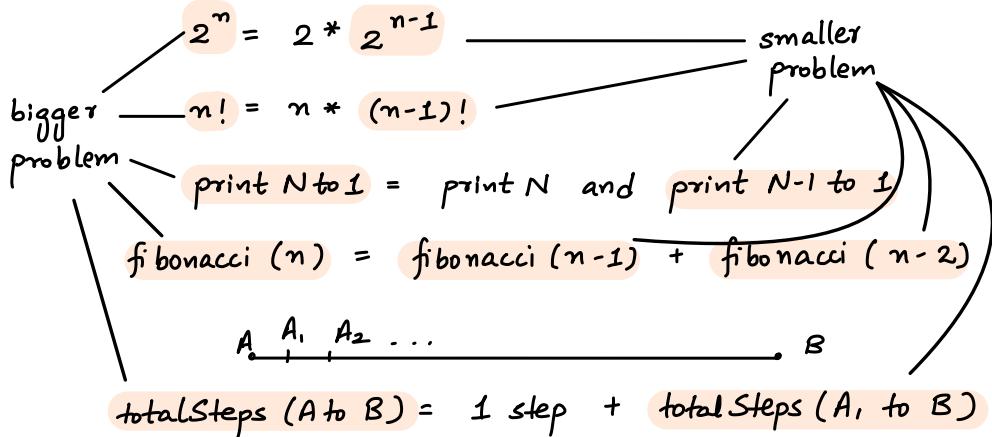
POITER IMP DOUBT

(IN DASH BOARD)

RECURSION

W7-L1

- When a function calls itself
- When sol. of bigger problem depends on similar smaller problem (s)



Components of Recursion

- Base Case / Stopping cond. → return statement
- Recurrence Relation (recursive call)
- Processing → optional

Head Recursion

- recurrence relation before processing

Tail Recursion

- recurrence relation after processing

```

int main () {
    int n = 4;
    print (n);
    return 0;
}

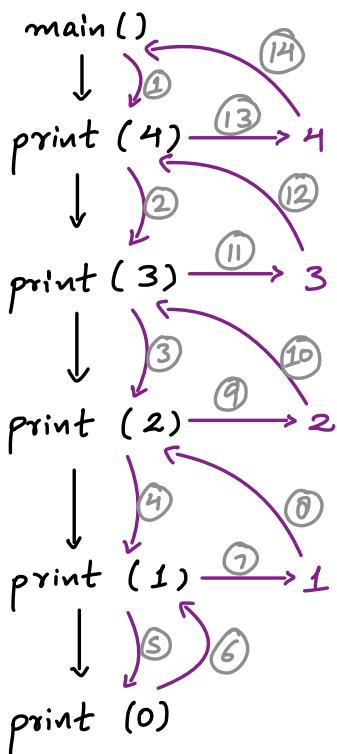
```

```

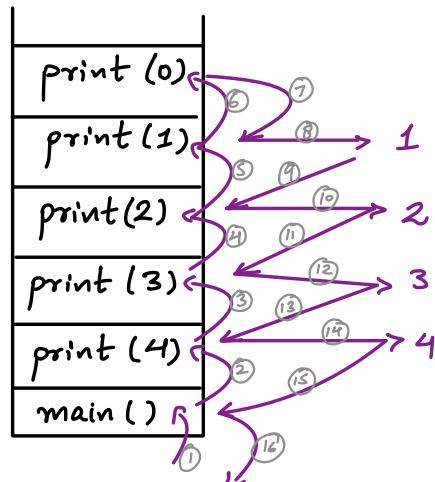
void print ( int n) {
    // base condition
    if ( n==0)
        return;
    // recurrence relation
    print (n-1);
    // processing
    cout << n;
}

```

Recursion Tree



Call Stack



Output - 1 2 3 4

```

int main(){
    int n = 4,
        ans = fib(n);
    cout << ans;
    return 0;
}

```

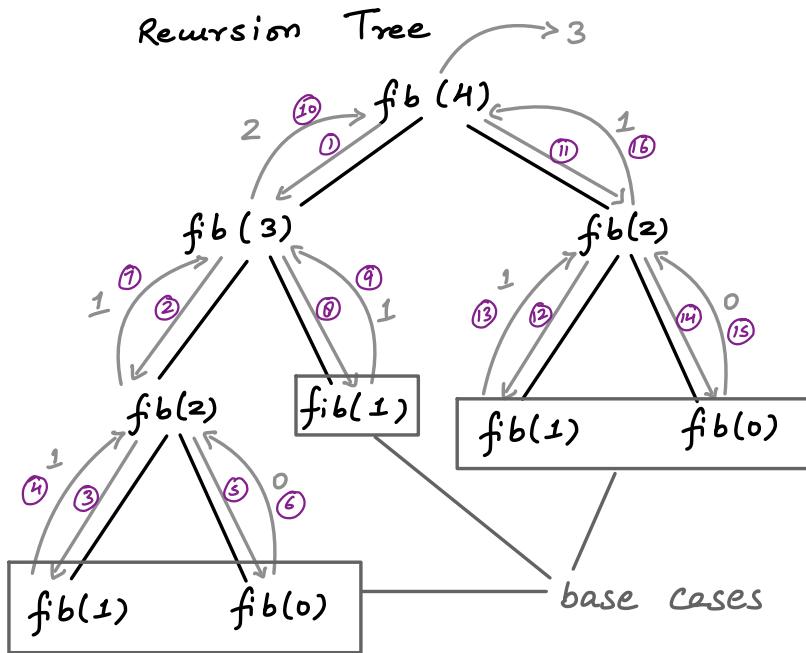
```

int fib( int n){
    if( n == 0 || n == 1)
        return n;
    return fib(n-1) + fib(n-2);
}

```

$n \rightarrow 0 \ 1 \ 2 \ 3 \ 4$

$\text{fib}(n) \rightarrow 0 \ 1 \ 1 \ 2 \ 3$



Magical line for Recursion

Ek case solve krdo, baaki recursion sambhal lega

Just believe on it, dont doubt on recursion

W7-L2

→ func(vector<int> arr, int i, int &ans)

pass by reference

if we want to retain value of ans after function call

→ Try to pass everything by reference IF POSSIBLE

└ SC Yes

TC Yes (no copying of variables again & again)

→ to reverse answer, you can use recursion

Integer literal with a leading 0

n = 0100 ; → interpreted as an octal number

cout << n; → 64 → convert 0100 octal to decimal

cin >> n; → (0100)

cout << n; → 100 → no conversion as
cin reads 0100 as 100

Entry in function call stack

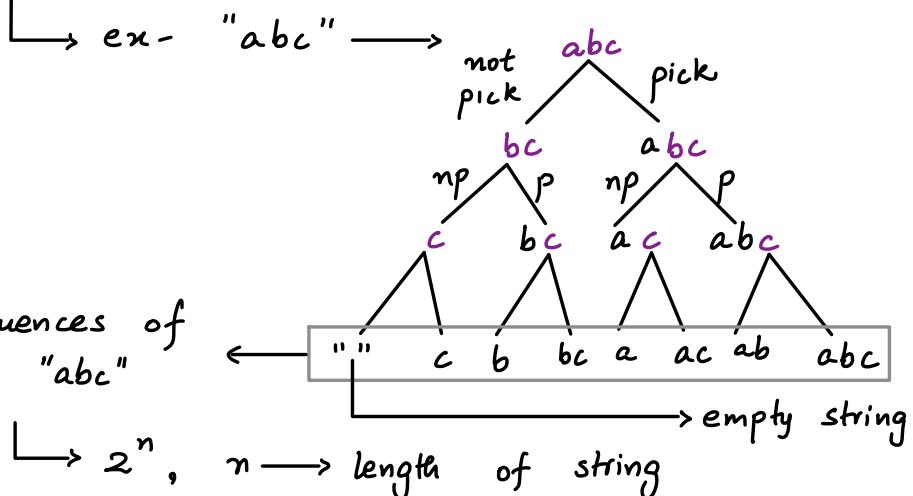
W7-L3

→ Computer has track of variables of that function about their memory location and all in symbol table

→ Expression can't be pass by reference

Subsequences

→ Made by either including or excluding every element, but order remains same



Arithmetic Operations on Strings

s1 = "Ujjwal";

s2 = "2327";

cout << "Ujjwal" + "2327"; → ERROR

s1 = s1 + s2;

cout << s1; → "Ujjwal2327" → Concatenate

- , * , / , % → NOT VALID in strings

Coin change problem

W7-L4

→ find min. coins equivalent to given amount

```
int coinChange( vector<int> arr, int amount) {  
    if (amount == 0)  
        return 0;  
    if (amount < 0)  
        return INT_MAX; // means INVALID or  
                           can't be changed  
    int mini = INT_MAX;  
    for (int i=0; i<arr.size(); i++) {  
        mini = min (mini, coinChange (arr, target - arr[i]));  
    }  
    // now min. coins req. has been found  
    if (mini == INT_MAX)  
        return INT_MAX; // means INVALID or  
                           can't be changed  
    else  
        return mini + 1;  
}
```

Type 1 - inclusion exclusion pattern → for using
either this or that subsequence questions

Type 2 - for loop in recursion

→ If the recursion ques seems tough,
use void function and pass ans by reference
in function

TC & SC of Recursive Solutions

W7-R

T.C. → Can be found by recursive relation
or recursive tree

S.C. → found using max depth of function call stack
└→ max space used in any particular time
instant in function call stack

Linear Search

$$f(n) = k + f(n-1) \quad \text{Base case} \rightarrow \text{for } n=0$$

Recursive Relation Method -

$$T(n) = k + T(n-1)$$

$$T(n-1) = k + T(n-2)$$

.

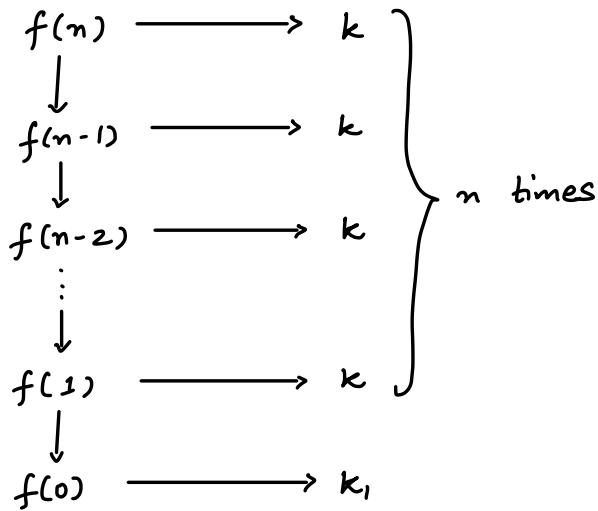
$$T(0) = k,$$

$$T(n) = nk + k_1$$

$$T(n) = O(n)$$

T.C - $O(n)$

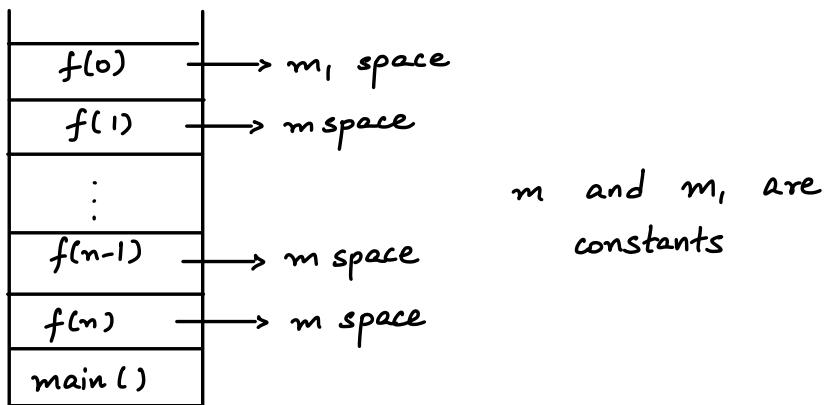
Recursive Tree Method - Best method



$$T(n) = nk + k_1$$

$$T(n) = O(n)$$

Space Complexity -



$$SC = mn + m_1$$

$$SC = O(n)$$

→ max depth

Binary Search

$$f(n) = k + f(n/2); \quad \text{base case} \rightarrow \text{for } n=0$$

Recursive Relation Method

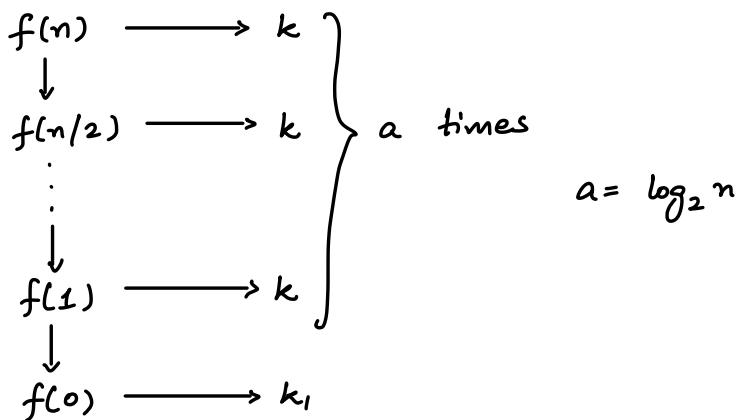
$$\begin{aligned} T(n) &= k + T(n/2) \\ T(n/2) &= k + T(n/4) \\ &\vdots \\ T(1) &= k + T(0) \\ T(0) &= k_1 \end{aligned} \quad \left. \begin{array}{l} \\ \\ \\ \end{array} \right\} a \text{ times} \quad a = \log_2 n$$

$$T(n) = ak + k_1$$

$$T(n) = k \log_2 n + k_1$$

$$T(n) = O(\log_2 n)$$

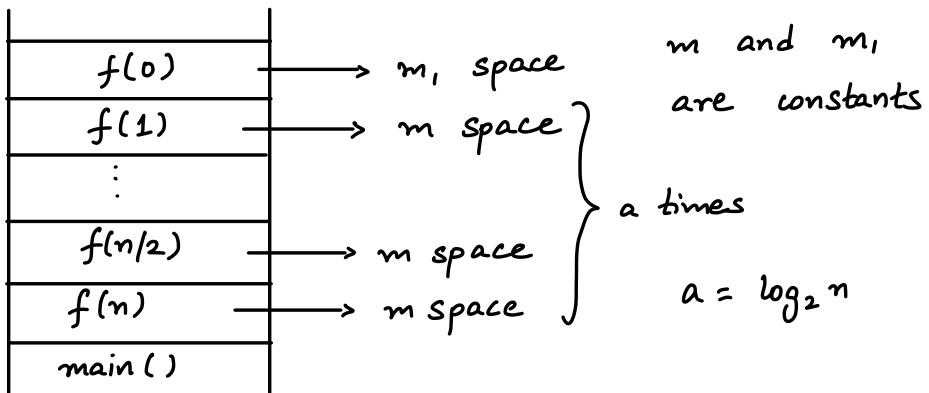
Recursive Tree Method



$$TC = k \log_2 n + k_1$$

$$TC = O(\log_2 n)$$

Space Complexity -



$$SC = m \log_2 n + m_1$$

$$SC = O(\log_2 n)$$

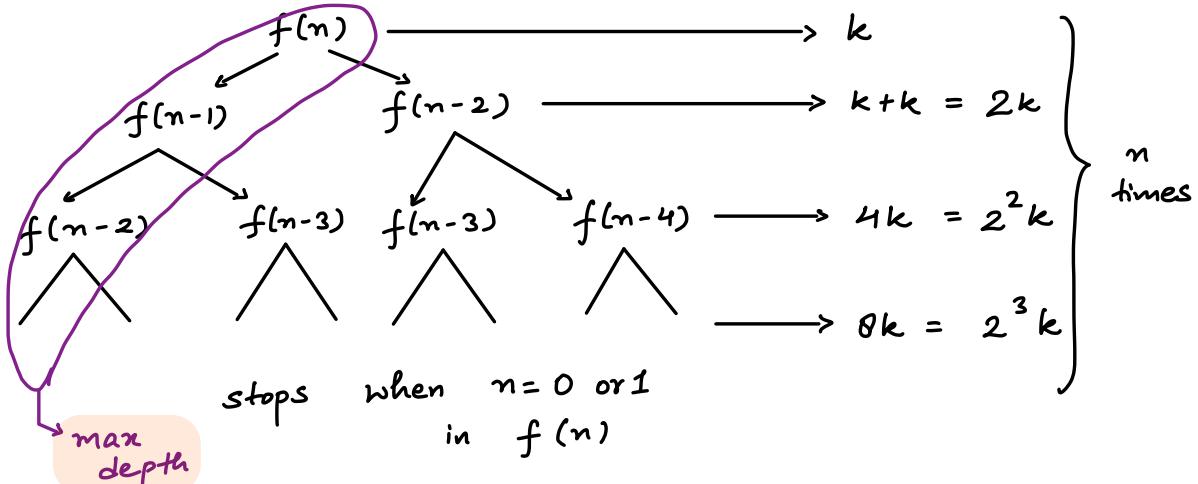
max depth

Fibonacci Series

$$f(n) = k + f(n-1) + f(n-2)$$

base case \rightarrow for $n=0$ or 1

Recursive Tree Method



$$T.C \leq [k + 2k + 2^2k + \dots + 2^{n-1}k]$$

$$T.C. \leq k * [1 + 2 + 2^2 + \dots + 2^{n-1}]$$

$$T.C \leq k * (2^n - 1)$$

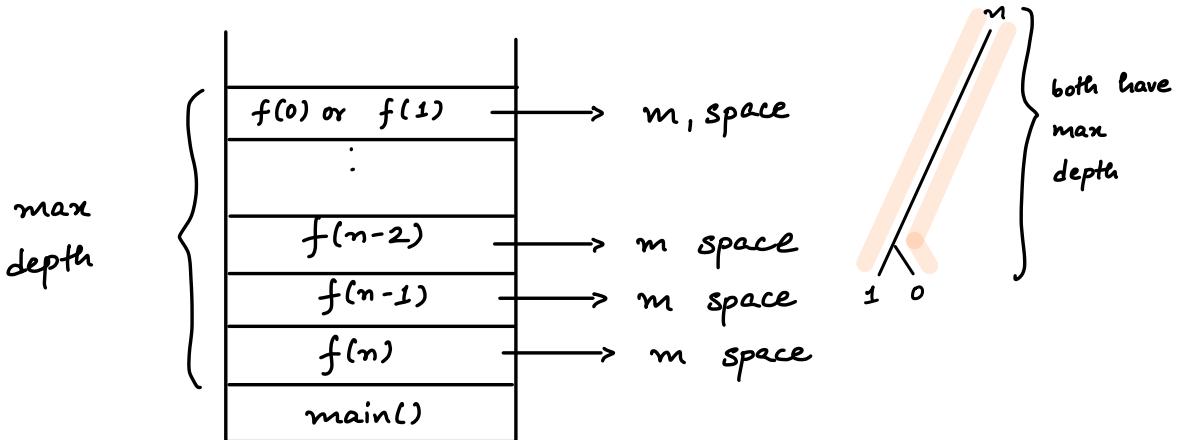
$$T.C = O(2^n)$$

assume
for simplicity
and for large
numbers

→ exponential TC

→ very bad

Space Complexity



$$SC = m(n-1) + m,$$

$$SC = O(n)$$

→ max depth

if this is such a case

fun(0)	→ n^3 space / time
fun(1)	→ n^2 space / time
:	
fun(n-1)	→ n space / time
fun(n)	→ k space / time
main()	

$$S.C. - n * n^3$$

↳ max SC in any call

$$T.C. - n * n^3$$

↳ max T.C. in any call

Divide & Conquer

→ using recursion

W8-L1

Merge Sort

- Divide the array in 2 halves
- Sort them
- Merge those 2 sorted subarrays

Standard Merge Sort

```
void merge( vector <int> &arr, int s, int e) {  
    int mid = e - (s + e)/2;  
    int len1 = s - mid + 1;  
    int len2 = e - mid;  
    vector <int> left (len1);  
    vector <int> right (len2);  
    int i=0, j=0, k=s;  
  
    while (i < len1)  
        left [i++] = arr [k++];  
  
    while (j < len2)  
        right [j++] = arr [k++];  
  
    i=0, j=0, k=s;
```

```

while ( i < len1 && j < len2 ) {
    if ( left [ i ] <= right [ j ] )
        arr [ k ++ ] = left [ i ++ ];
    else
        arr [ k ++ ] = right [ j ++ ];
}
while ( i < len1 )
    arr [ k ++ ] = left [ i ++ ];
while ( j < len2 )
    arr [ k ++ ] = right [ j ++ ];
}

void mergeSort ( vector < int > & arr, int s, int e ) {
    if ( s >= e )
        return ;
    int mid = s + ( e - s ) / 2;
    mergeSort ( arr, s, mid );
    mergeSort ( arr, mid + 1, e );
    merge ( arr, s, e );
}

```

T.C. - $O(n \log n)$

S.C. - $O(n)$

IMPLEMENT MERGE SORT - GFG

Quick Sort

→ Partition

Choose one element named pivot and put it on its correct position

Put elements less than pivot in left subarray and otherwise in right subarray

→ Sort left & right subarrays recursively

```
int partition( vector<int> &arr, int s, int e){
    int pivotIndex = s; —————> This is one way
    int pivot = arr[s];
Other ways - choose last  
- choose random
↓  
randomised  
quick sort
    int cnt = 0;
    for( int i=s+1; i<=e ; i++ ){
        if( arr[i] > pivot )
            cnt++;
    }
    int correctPosition = s + cnt;
    swap( arr[pivotIndex], arr[correctPosition] );
    pivotIndex = correctPosition
}
```

```

int i=s, j=e;
while( i < correctPosition && j > correctPosition) {
    while( arr[i] < pivot)
        i++;
    while( arr[j] >= pivot)
        j--;
    if( i < correctIndex && j > correctIndex)
        swap( arr[i++], arr[j--]);
}
return correctPosition;
}

void quickSort( vector<int> &arr, int s, int e) {
if( s >= e)
    return;
int p = partition( arr, s, e);
quickSort( arr, s, p-1);
quickSort( arr, p+1, e);
}

```

T.C. - $O(n \log n)$ \longrightarrow log_n number of function calls

$O(n^2)$ \longrightarrow for both ascending & descending sort

S.C. - $O(n)$ \longrightarrow n number of function calls

Backtracking

- specific form of recursion
 - explore all possible solutions and
 - don't check the solution you have already discarded
 - brute force
 - used for reference variables passed in functions

Permutation in a string

$s = "abc"$ permutations - abc, acb, bac, bca,

permutations - abc, acb, bac, bca,
 | | | abc

cab , cba

> total - n!

