## TRANSLATION SCHEME

```
Program → D Program | ^
D → Code | Function | Koment
Code → Statement Code | If Code | While Code | ^
Statement → Stmt koment
Stmt → Variable | Input | Output | Return | Chalao
Function → kaam ID { Funct.id = Id.lex} @ FuncT ( PL ) karo Koment Code kaam
khatam Koment
FuncT -> khali | adad { symbolTbl.add(ID.lex,Funct.id, lineNumber}
PL \rightarrow ID  @ adad MPL | ^
MPL \rightarrow |PL|^{\wedge}
-----Rakho-----
Variable \rightarrow rakho ID Type { R.id = Id.lex } R
Type \rightarrow @ adad | ^
\mathbf{R} \rightarrow := Val \{
emit(R.id+"="+ Val.v);
R.v =SymbolTable.add(R.id, INT); }
| ^ { R.v = SymbolTable.add(R.id, INT); }
Val \rightarrow ID \{ Val.v=ID.lex; \}
| Integer { Val.v=Integer.lex; }
| Exp { Val.v=Exp.v; }
| Chalao {Val.v=Chalao.v}
Condition → Cexp RO Cexp { Condition.V = Exp.ex + Ro.lex + Exp.ex }
```

## -----Expression-----

```
// P, S are the recursive vars
E \rightarrow T P
P \rightarrow +T P1 | -T P1 | ^
T \rightarrow F S
S \rightarrow \% F S | / F S | * F S | ^
F \rightarrow ID \mid Digit
```

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## **Actions**

NOTE: New temp function will automatically add that variable in symbol table

```
\begin{split} \textbf{E} &\to & \textbf{T} \quad \{ \, \text{P.i} = \text{T.v} \, \} \quad \textbf{P} \quad \{ \, \text{E.v} = \text{P.s} \, ; \, \} \\ \textbf{P} &\to & + \\ & & \textbf{T} \quad \{ \\ \text{var} = \text{newTemp}(); \\ \text{emit}( \, \text{var} + \text{"="} + \text{P.i} + \text{"+"} + \text{T.val}); \\ \textbf{P}_{1}.i = \text{var}; \\ & \textbf{P}_{1} \quad \{ \, \text{P.s} = \text{P}_{1}.s \, \} \end{split}
\textbf{P} &\to & - \\ & \textbf{T} \quad \{ \\ \text{var} = \text{newTemp}(); \\ \text{emit}(\text{var} + \text{"="} + \text{P.i} + \text{"-"} + \text{T.val}); \\ \textbf{P}_{1}.i = \text{var} \\ \} \\ & \textbf{P}_{1} \quad \{ \text{P.s} = \text{P}_{1}.s \} \end{split}
```

$$P \rightarrow ^{\Lambda} \{P.s = P.i i\}$$

$$T \rightarrow F$$
 {Q.i = F.v} Q {T.v = Q.s}

```
\mathbf{Q} \rightarrow *
         F
var =newTemp();
emit(var + "=" + Q.i + "*" + F.val};
Q_1.i = var
}
         Q_1 = \{Q.s = Q_1.s\}
\textbf{Q} \rightarrow
          F
var =newTemp();
emit(var + "=" + Q.i + "/" + F.val};
Q_1.i = var
}
         Q_1 \qquad \{Q.s = Q_1.s\}
\mathbf{Q} \rightarrow \%
var =newTemp();
emit(var + "=" + Q.i + "%" + F.val};
Q_1.i = var
}
         Q_1 	 \{Q.s = Q_1.s\}
\mathbf{Q} \rightarrow ^{\Lambda} \{Q.s = Q.i\}
\mathbf{F} \rightarrow num \{ F.v = num.lex \}
\mathbf{F} \rightarrow ID { F.v = id.lex }
```

## -----Function Call-----

```
Chalao → chalao ID { PLF.i=0; } ( PLF ) {
var=newTemp();
emit ("call" + ID.lex + PLF.v + "," +var);
Chalao.v = var;
PLF \rightarrow ID \{
emit("param "+ ID.lex);
PLF.i = PLF.i +1; // +1
MPLF.i = PLF.i;
} MPLF { PLF.v = MPLF.v; }
PLF → Integer {
emit ("param"+Integer.lex);
PLF.i=PLF.i+1;
MPLF.i = PLF.i;
} MPLF { PLF.v =MPLF.v; }
PLF \rightarrow ^{\land} \{ PLF.v = PLF.i ; \}
MPLF \rightarrow | { PLF.i = MPLF.i ;} PLF {MPLF.v = PLF.v;}
MPLF → ^ { MPLF.v = MPLF.i ;}
Koment → Comment | ^
IF \rightarrow agar ( Condition ) to phir karo {
InTrue= n;
emit ("if" + Condition.v + goto + ___);
InFalse= n;
Emit ("goto" + )
BackPatch(InTrue)
Koment
Code {
IF end= In;
```

```
emit (goto )
BackPatch(InFalse)
WG
WP
bas karo
BackPatch(IF end)
BackPatch(WG.val)
Koment
WG → warna agar Condition to phir Koment {
InTrue = n;
emit ("if" + Condition.v + goto + );
InFalse = n;
emit( goto ___)
BackPatch(InTrue_)
Code {
WG.v= In; // storing the current line number for Branch Ending
emit (goto ___ )
BackPatch(InFalse)
}
WG \rightarrow ^{\wedge}
WP → warna phir Koment Code
WP → ^
Return-> wapis bhaijo Val { emit ("ret" + Val.v) }
// Todo : Add cascading to it
Input \rightarrow /o InputMsg >> ID { emit("in"+ID.v+"\n") }
InputMsg \rightarrow ^
InputMsg \rightarrow << String { emit ("out" +String .v +"\n") }
```

```
Output \rightarrow dekhao << OutVal { emit ("out" + OutVal.v +"\n" ) } MoreOut MoreOut \rightarrow << OutVal MoreOut { emit ( "out" + OutVal.v +"\n" ) } MoreOut \rightarrow^ OutVal \rightarrow String { String.lex } | Val { Val.v }
```

Note: Backpatch has global access to In, so it patches current line number at the parameter passed to it

```
While → jab tak ( Condition ) karo Koment {

InTrue = n;

emit ("if" + Condition.Value goto _____);

InFalse = n;

Emit (goto ____)

BackPatch(InTrue)}

Code

{ emit( "goto" + InTrue) }

bas karo { BackPatch(InFalse) }

Koment
```