CSC 520, Spring 2020

Principles of Programming Languages

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Plan

- **Yesterday**
  - Closures to create "private" variables
  - High-order function curry
  - Reasoning about functions
  - Useful higher-order functions: exists? all? filter map foldr
- **Today**
  - Continuations

Continuations

- **Code that represents "the rest of the computation."**
- **Not a normal function call because continuations never return (think "goto with arguments")**
- **Different coding styles**
  - **Direct style: Last action of a function is to return a value.**
  - **Continuation-passing style (CPS): Last action of a function is to "throw" a value to a continuation.**

Uses of continuations

- **A style of coding that can mimic exceptions**
- **Callbacks in GUI frameworks**
- **Some languages**
  - provide a construct for capturing the current continuation and giving it a name k.
  - Control can be resumed at captured continuation by throwing to k.
- **Compiler representation**
  - Compilers for functional languages often convert direct-style user code to CPS…
  - Because CPS matches control flow of assembly

Implementation

- **First-class continuations require compiler support**
- **We are going to simulation continuations with function calls in tail position**
- **Tail position is defined inductively:**
  - **The body of a function is in tail position.**
  - **When (if e1 e2 e3) is in tail position, so are e2 and e3.**
  - **When (let (...) e) is in taile position, so is e, similar for letrec and let*.**
  - **When (begin e1 ... en) is in tail position, so is en.**
- **Idea: The last thing that is executed**

How functions finish

Direct: `return answer;`

True CPS: `throw k answer;`

uScheme: `(k answer)`

Motivating Ex: From existence to witness

Design Problem: Missing Value

Provide a witness to existence:

```
(witness p? xs) == x, where (member x xs),
                            provided (exists? p? xs)
```

Problem: What if there exists no such `x`?

Ideas?

Solution: A New Interface

Success and failure continuations!

Contract written using properties (not algorithmic):

```
(witness-cps p? xs succ fail) = (succ x)
      ; where x is in xs and (p? x)

(witness-cps p? xs succ fail) = (fail)
      ; where (not (exists? p? xs))
```

Your turn: Refine the laws


```
(witness-cps p? xs succ fail) = (succ x)
    ; where x is in xs and (p? x)

(witness-cps p? xs succ fail) = (fail)
    ; where (not (exists? p? xs))

(witness-cps p? '() succ fail) = ?

(witness-cps p? (cons z zs) succ fail) = ?
    ; when (p? z)

(witness-cps p? (cons z zs) succ fail) = ?
    ; when (not (p? z))
```

Coding with continuations


```
(define witness-cps (p? xs succ fail)
    (if (null? xs)
        (fail)
        (let ((x (car xs)))
            (if (p? x)
                (succ x)
                (witness-cps p? (cdr xs) succ fail)))))
```

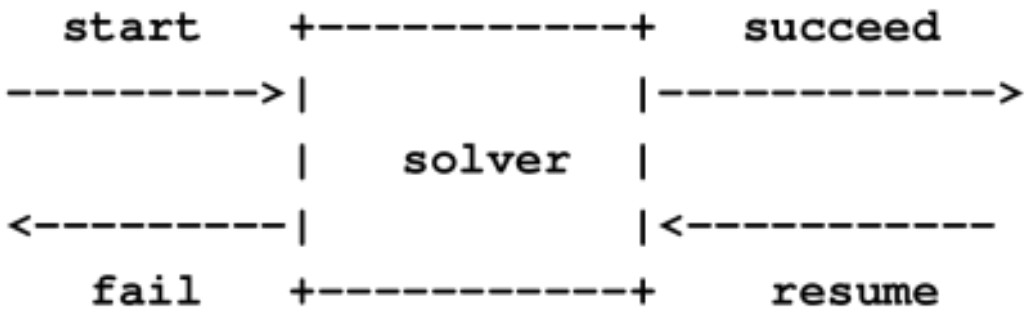
Are all tail positions continuations or recursive calls?

➔ *Do activity*

Extended example: A SAT Solver

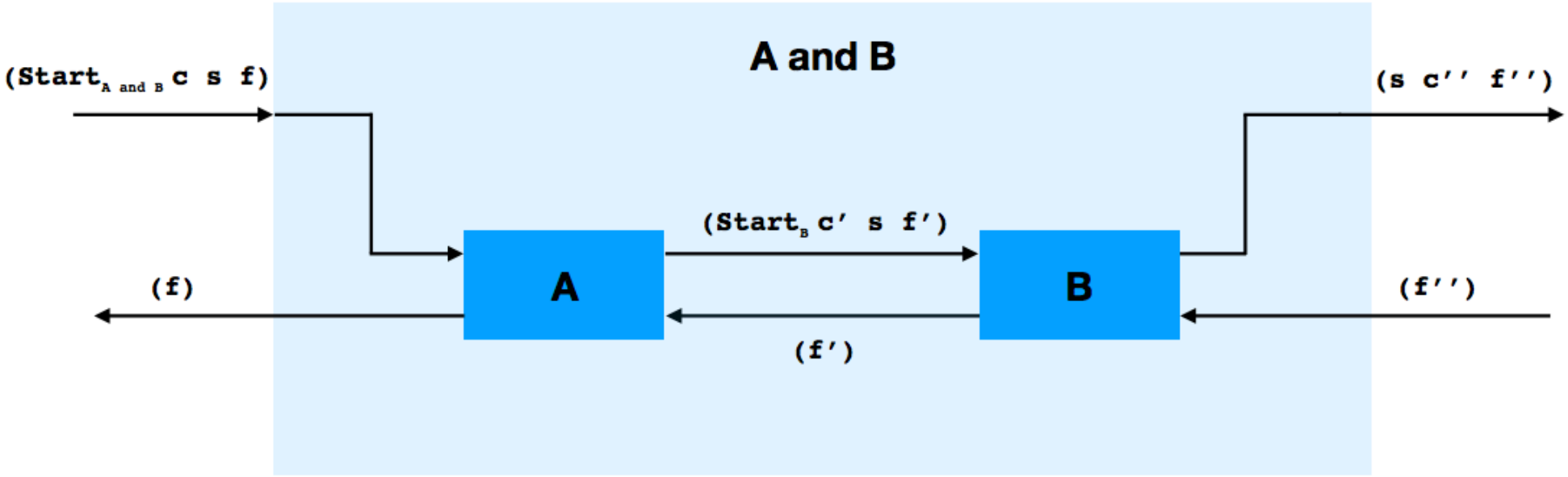

```
;; Find a satisfying assignment if one exists

(val f1 '(and x y z w p q (not x)))

(val f2 '(not (or x y)))

(val f3 '(not (and x y z)))

(val f4 '(and (or x y z) (or (not x) (not y) (not z))))
```

Continuations for Search



start	Gets partial solution, `fail`, `succeed` (On homework, "solution" is assignment)
`fail`	Partial solution won't work (no params)
`succeed`	Gets improved solution + `resume`
`resume`	If improved solution won't work, try another (no params)

Given boxes for "A" and "B", we can build a box for "A and B"



Given boxes for "A" and "B", we can build a box for "A or B