

# A Tour of Type

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Prelude



# Starting from C

```
#include <stdio.h>

int main() {
    // an integer
    int i = 42;
    // a double precision floating point
    double d = 0.1;

    int wth = i + d;
    // ???
    printf("%d\n", wth);
}
```

# So ... WHAT IS TYPE, for real ?

A type is a **classification of data** which tells the compiler or interpreter how the programmer intends to use the data. By Wikipedia

~~Not exactly.~~



# Type as interpretation of bytes

```
#include <stdio.h>
#include <stdint.h>

int main() {
    char a = 'a';
    // ?
    printf("%d\n", a);

    int64_t i = 0x000A214F4C4C4548;
    // ???
    printf("%s", (char*)&i);
}
```

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## Basic Terminologies



# Type Checking: dynamic vs. static

python

```
def main():  
    print("Program start\n")  
  
    im_not_array = 1  
    print(im_not_array[0])
```

C

```
#include <stdio.h>  
  
int main() {  
    printf("program start\n");  
  
    int im_not_array = 0;  
    printf("%d\n", im_not_array[0]);  
}
```

# Type Checking: weak vs. strong

C

```
void Test() {  
    double d = 123.321;  
    int i = d; // ok  
}
```

java

```
public class NotSoWeak {  
    static public void Test() {  
        double d = 123.321;  
        int i2 = d; // error  
        int i = (int)d; // ok  
    }  
}
```

haskell

```
d :: Double  
d = 123.321  
  
i :: Int  
i = floor d -- or i = ceiling d
```



# Type Safety

Type safety is the extent to which a programming language discourages or prevents type errors.

```
void Unsafe() {  
    double d = 2.5;  
    // ???  
    printf("%d\n", d);  
  
    int array[5] = {1, 0, 0, 8, 6};  
    int n = 10086;  
    // out-of-range indexing is not type error  
    printf("%d\n", array[n]);  
}
```

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## Basic Types



# Primitives (maybe..?)

- **Boolean:** true, false
- **Character:** 'a', 'b', '我'
- **Integer:** -1, 2, 3, ...
- **Floating Point:** -1.5, 0.111
- **String:** "hello world", "大吉大利"
- ...

# Array

- represent the sequence of some type
- usually stored in continuous memory
- fast integer indexing
- static or dynamic growing

```
func array() {  
    dynamic_array := make([]int, 0)  
    for i := 0; i < 100; i++ {  
        // the array can grow at runtime  
        append(dynamic_array, i)  
    }  
    for i := 0; i < len(dynamic_array); i++ {  
        fmt.Println(dynamic_array[i])  
    }  
}
```

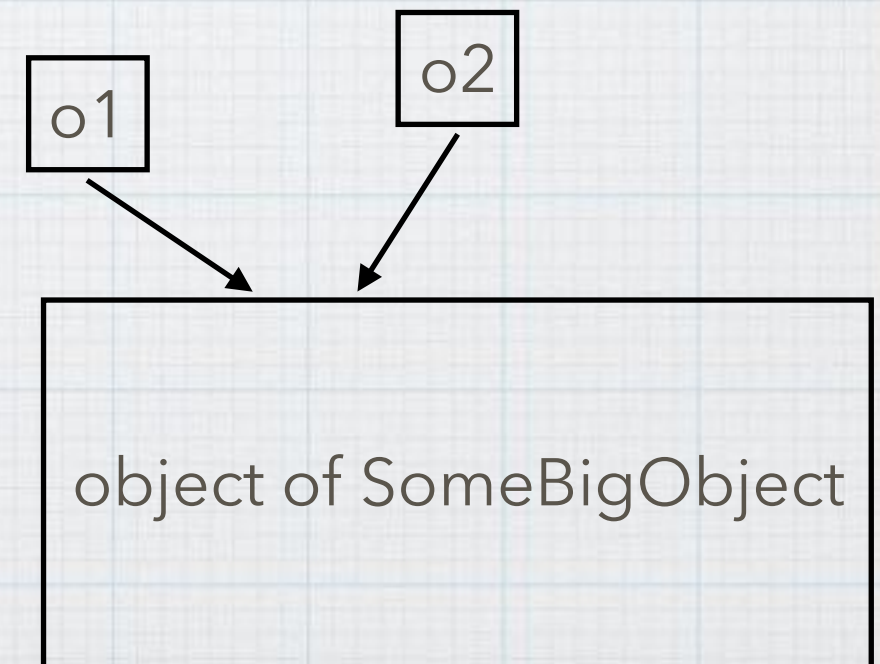


# Pointer & Reference

- point to the actual object of some type
- reduce cost of parameter passing
- allow different variable to share same data
- implicit or explicit dereferencing

```
static void Reference() {  
    SomeBigObject o1 = new SomeBigObject();  
    SomeBigObject o2 = o1;  
    Foo(o1);  
}  
static void Foo(SomeBigObject o) {  
    // do something  
}
```

references to SomeBigObject





# Tuple & Aggregation

- combine multiple types to another type
- promote type abstraction

```
struct Point {  
    double x;  
    double y;  
};
```

```
double Mod(const Point *p) {  
    return sqrt(pow(p->x, 2) + pow(p->y, 2));  
}
```

```
Point Normalize(const Point *p) {  
    double m = Mod(p);  
    Point result = {p->x / m, p->y / m};  
    return result;  
}
```



# Function Type

- the mapping between different types
- specify actions, behaviours or predicates

```
def Sort(a: Array[Int], compare: (Int, Int) => Boolean) = {  
  // Sort the array with special compare function  
}
```

```
def LessThan(i: Int, j: Int) = i < j  
def GreaterThan(i: Int, j: Int) = i > j
```

```
def SortAscent(a: Array[Int]) = Sort(a, LessThan)  
def SortDescent(a: Array[Int]) = Sort(a, GreaterThan)
```



# Generic Type

- allow single interface being used with different type parameters
- promote type abstraction
- parametric polymorphism

```
def Sort[T](a: Array[T], compare: (T, T) => Boolean) = {  
  // ..  
}  
def IntAscent(i: Int, j: Int) = i > j  
def StringAscentByLength(i: String, j: String) = i.length > j.length  
  
def Test(args: Array[String]): Unit = {  
  val a1 = Array(2, 3, 1, 4, 5)  
  val a2 = Array("hello", "generic", "sort")  
  
  Sort(a1, IntAscent) // 5 4 3 2 1  
  Sort(a2, StringAscentByLength) // generic hello sort  
}
```



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Object Oriented

# Type as simulation of real world

```
class Person {  
    public:  
        // methods are behaviours  
        void Kill(Person p) {  
            if (p.weight < weight and p.height < height) {  
                p.alive = false;  
            }  
        }  
        void Suicide() {  
            alive = false;  
        }  
        // .. other "methods"  
    private:  
        // fields are properties  
        double weight;  
        double height;  
        bool alive;  
        string name;  
        // ... other "fields"  
};
```



# Well-known features of OO

- Encapsulation
- Inheritance
- Polymorphism

# Abstract Data Type

Abstract data type is a class of objects whose logical behaviour is defined by a set of values and a set of operations.

```
interface IntArray {
    int get(int index);
    int set(int index, int n);
    int length();
    int[] toArray();
    // ..
}

boolean Has(IntArray arr, int n) {
    for (int i = 0; i < arr.length(); ++i) {
        if (arr.get(i) == n) return true;
    }
    return false;
}
```



# Subtyping

- a form of type polymorphism
- supertypes can be “substituted” by its subtypes where appropriate
- $S$  is a subtype of  $T$  is written as  $S <: T$
- Nominal & Structural subtyping



# Nominal Subtyping

```
class Animal {  
    public String sound() { return "Awww!"; }  
}  
class Dog extends Animal {  
    public String sound() { return "bark!!"; }  
}  
class Cat extends Animal {  
    public String sound() { return "Meow"; }  
}  
class Cow {  
    public String sound() { return "Mooooooooooooo"; }  
}  
boolean Louder(Animal a, Animal b) {  
    return a.sound().length() > b.sound().length();  
}  
void Animals() {  
    Dog dog = new Dog();  
    Cat cat = new Cat();  
    Cow cow = new Cow();  
    Louder(dog, cat); // true  
    // error: Cow is not a subtype of Animal  
    Louder(cow, dog);  
}
```



# Structural Subtyping

```
class Point {  
  x: number;  
  y: number;  
}  
class Line {  
  a: Point;  
  b: Point;  
}  
function Length(l: Line) { return 10086; }
```

```
let p1 = {x: 1, y: 1};  
let p2 = {x: 1, y: 1, z: 1};
```

```
let l1 = {a: p1, b: p1};  
let l2 = {a: p1, b: p1, c: "hello"};  
let l3 = {a: p1, b: p2};
```

```
Length(l1); // exact  
Length(l2); // subtype by width  
Length(l3); // subtype by depth
```



# Bounded type parameter

- able to declare type relationship of type parameters
- more expressive generic types

```
def Concat[A](as1: List[A], as2: List[A]): List[A] = {  
  // ..  
}  
def Concat2[A, B <: A](as1: List[A], as2: List[B]): List[A] = {  
  // ..  
}  
def Join[S <: String](ss: List[S]): String = {  
  // ..  
}
```



# Interesting issues of subtyping

## Subtyping of function types

- $S <: T, U \text{ is a type} \implies U \rightarrow S \text{ ?? } U \rightarrow T$
- $S <: T, U \text{ is a type} \implies S \rightarrow U \text{ ?? } T \rightarrow U$
- $S1 <: T1, S2 <: T2 \implies S1 \rightarrow S2 \text{ ?? } T1 \rightarrow T2$
- $S1 <: T1, S2 <: T2 \implies S1 \rightarrow T2 \text{ ?? } T1 \rightarrow S2$

## Higher Order Subtyping

- $S <: T \implies \text{Array}\langle S \rangle \text{ ?? } \text{Array}\langle T \rangle$
- $S <: T, U \text{ is a generic type} \implies U\langle S \rangle \text{ ?? } U\langle T \rangle$

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## Algebraic Data Type



# Type as set of values

How many values do these types have:

`bool`   `char`

`(bool, char)`

`union {bool b; char c;}`

`()`

`void`

`bool -> char`

# Algebraic Data Type (ADT)

Notations (informal):

- $S$  and  $T$  are ADTs  $\implies S + T$  is ADT (sum/variant)
- $S$  and  $T$  are ADTs  $\implies S * T$  is ADT (product/tuple)
- $S$  and  $T$  are ADTs  $\implies S \rightarrow T$  is ADT (function)

Informally, define  $N(T)$  be the number of values in some ADT  $\tau$ , then

- $N(S + T) = N(S) + N(T)$
- $N(S * T) = N(S) * N(T)$
- $N(S \rightarrow T) = N(T) ^ N(S)$



# Simple Examples

## Function Totality

```
indexOf :: (Array<T>, T) -> int  
type Optional<T> = T * ()  
indexOf :: (Array<T>, T) -> Optional<T>
```

## Currying

```
type Uncurried = (A, B) -> C  
type Curried = A -> (B -> C) // or simply A -> B -> C
```

## List

```
type List<A> = () + A * List<A>
```



# Pattern Matching

- Matching the structure of values of ADTs recursively
- Destructure values into values with “smaller” type

```
data List a = Null | Cons a (List a)
```

```
tail :: List a -> List a
```

```
tail Null = Null
```

```
tail (Cons a t) = t
```

```
tail2 :: List a -> List a
```

```
tail2 Null = Null
```

```
tail2 (Cons a Null) = Null
```

```
tail2 (Cons a (Cons a' t)) = t
```



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Type & Proof

# Type as proposition: the intuition

```
void PrintInteger(int i) {  
    printf("I got an integer: %d!\n", i);  
}
```

```
validateForm :: Form -> Maybe ValidForm
```

```
submitForm :: ValidForm -> Result
```

```
processForm :: Form -> Either Error Result
```

```
processForm form =  
    case validateForm of  
        Just validForm -> Right (submitForm validForm)  
        Nothing        -> Left  invalidFormError
```



# Curry-Howard Isomorphism

- the direct correspondence between mathematical proofs and computer programs
- values as proofs, types as propositions

# Propositional Logic

Proposition	Type
$A$	$A$
$A \wedge B$	$A * B$
$A \vee B$	$A + B$
$A \Rightarrow B$	$A \rightarrow B$
$\neg A$	$A \rightarrow \text{void}$

Currying: the propositional view:

$$A \wedge B \Rightarrow C \Leftrightarrow A \Rightarrow B \Rightarrow C$$

$$(A, B) \rightarrow C \Leftrightarrow A \rightarrow B \rightarrow C$$



# Quantifiers & Dependent type

- a dependent type is a type whose definition depends on a value.
- e.g. `std::array<int, 5>`

Proposition	Type
$(\forall x:A).P(x)$	$(x:A) \rightarrow P(x)$
$(\exists x:A).P(x)$	$(x:A) * P(x)$

`concat` : `Vect m a`  $\rightarrow$  `Vect n a`  $\rightarrow$  `Vect (m + n) a`

$\forall m,n,a. \text{Vect}(m, a) \wedge \text{Vect}(n, a) \Rightarrow \text{Vect}(m + n, a)$



# Induction & Recursion

```
data Natural = Zero
              | Succ Natural
```

```
data List a = Null
            | Cons a (List a)
```

```
-- illustration purpose, need extra refinement
```

```
-- Nil : Vect 0 a
```

```
concat : Vect m a -> Vect n a -> Vect (m + n) a
```

```
-- base case: Vect 0 a -> Vect n a -> Vect (0 + n) a
```

```
concat Nil v = v
```

```
-- induction case:
```

```
-- given Vect k a -> Vect n a -> Vect (k + n) a
```

```
-- prove Vect (1 + k) a -> Vect n a -> Vect (1 + k + n) a
```

```
concat (a :: as) v = a :: concat as v
```



# Predicate & Type: motivation

Expressing equality as result of “function” ?

```
zip : Vect n a -> Vect n b -> Vect n (a, b)
```

```
tryZip : Vect m a -> Vect n b -> Maybe (Vect m (a, b))
```

```
-- did m == n say anything about equality?
```

```
wrongTryZip : Vect m a -> Vect n b -> Maybe (Vect m (a, b))
```

```
wrongTryZip {m} {n} v1 v2 =
```

```
  if m == n
```

```
    then Just (zip v1 v2)
```

```
    else Nothing
```



# Equality as a type

Equality of two things is a proposition!

```
data (==) : A -> B -> Type where  
  Refl : x == x
```

```
zip : Vect n a -> Vect n b -> Vect n (a, b)
```

```
equal : (m: Nat) -> (n: Nat) -> Maybe (m == n)
```

```
tryZip : Vect m a -> Vect n b -> Maybe (Vect m (a, b))
```

```
tryZip {m} {n} v1 v2 =
```

```
  case equal m n of
```

```
    Nothing => Nothing
```

```
    -- the proof here tell compiler m and n are the same
```

```
    (Just Refl) => Just (zip v1 v2)
```



# Summary & Questions