Speciality: O(1) time complexity, for all the dictionary operations. HASH TABLE Hagh table data structure stores elements in Key-value pairs where pairs where Key - Unique integer that is used for indesing the values. Value-Data that are associated with Keys. Key Value Dictionosy operations: INSERT, SEARCH AND DELETE Different Names: Hosh Mab Object Dictionary

Java

Ja In HashMaps, we can't have duplicate keys

Date Type

Date Type Map (Key, Value) map = new HashMap (>U); i. map. put () 11 To add a Key, value paix ii map removel) 11 To remove ici. map. get (Key) 11 Ovet the value at a particular key iii map. contains key() (/Return a boolean if the key is in the Hash table. iv. map. contains Value () 1/ It has to go over all the elements, to check whether the value

O(n) exists or not.

v. map. Keyset () 1 Keturns au the Keys 11 Keturns all the Key Value pairs. vi. map. entryset () MI FIR. HEY STR Maps: Key - Value Maxilla of - Laylating in Philippin De Charlet Sets: Key Adrailed Health - A set only contains Keys. And it doesn't allow duplicates. g: Get (Integer) set = new HoshSet (1); int [] numbers = {1,2,3,3,2,1,43; Jor (int num: numbers)
Set.add (num); System.out.println(set) //output: {1,2,3,4}. Methods: set. contains AU() set. add()

set. semove()

set. size() set. dear () Set. remove AII.() set contains) or inchis at it is property of the potential of

Hausing (Hash Function) In a hosh table, a new index is processed using the keys. And the element corresponding to at Key ! that key is stored in the index. This process is Known () as Hashing. Hash Collision - When the hash function generates the some index for to be stored in that indered. This is known as Hash Collisson + Bimply: When two Keys generates some indext to store their values by going through Some Kind of hosh bunction. Kesolving Hosh Collisions: [hosh I = Key / table - size]

Gome Kind of hosh function.

Resolving Hosh Collisions: [hosh I = Key / table_size]

In Integer form

Closed / Separate Chaining: Elements are stored

in the same index by using a

doubly kinked dist.

Open Addressing: By Probing (searching for an

Open Addressing: By Probing (searching for an

Linear Probing: (hash 1 + i) // table_size

(squared) Quadratic Probing: (hash 1 + i² hash 2)// table_size

Double Hash: (hash 1 + i² hash 2)// table-size

hash 2 (Key): prime - (Key 1. prime)

```
Some basic algorithms:
To find the brequency of all the Keys in Hash Map.
        Str = " a mango man"
        Map (String, Integer) map = new HashMap ();
       Jos (Chas Key: Str. to Chas Array())
 int count = map.containskey(key) ?map.get(key):0;
Count the 3
        System.out.println(map);
      // Output: 2a=3, -=2, m=2, n=2, g=1, 0=1
```