Section : 10

Sparse Matrix and Polynomial Representation

Sparse Matrix Representation -

	1	2	3	4	5	6	7	8	9
1	0	0	0	0	0	0	0	3	0
2	0	0	8	0	0	10	0	0	0
3	0	0	0	0	0	0	0	0	0
4	4	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0	0
6	0	0	2	0	0	0	0	0	0
7	0	0	0	6	0	0	0	0	0
8	0	9	0	0	5	0	0	0	0

8 x 9 (72 element)

72 x 2 = 144 bytes

Sparse Matrix → In numerical ~~and~~ analysis and scientific computing, a sparse matrix or sparse Array is a matrix in which most of the element are zero.

Method to store non-zero elements -

1. Coordinate list (3 column representation)
2. Compared Sparse row

3-Column Representation :-

Row number, column number, value (element)

	row	column	element
No. of rows, no. of column →	8	7	8
	1	8	3
	2	3	8
	2	6	10
	4	1	4
	6	3	2
	7	4	6
	8	2	9
	8	5	5

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Compressed Sparse rows :-

A [3, 8, 10, 4, 2, 6, 9, 5] (Order wise)

IA [0, 1, 3, 3, 7, 7, 9,
0 1 2 3 4 5 6

2+1 3+0 3+4 7+0

~~IA [8, 3, 6, 1, 3, 4, 2, 5]~~

IA [0, 1, 3, 3, 4, 4, 5, 6, 8]
0 1 2 3 4 5 6 7 8

↑
prev + no of elem
sum

IA [8, 3, 6, 1, 3, 4, 2, 5]

Space

⇒ 8 + 9 + 8

⇒ 25 x 2

= 50 bytes

Addition of Sparse Matrices

$A =$

	1	2	3	4	5	6
1	0	0	0	6	0	0
2	0	7	0	0	0	0
3	0	2	0	5	0	0
4	0	0	0	0	0	0
5	4	0	0	0	0	0

$B =$

	1	2	3	4	5	6
1	0	0	0	0	0	0
2	0	3	0	0	5	0
3	0	0	2	0	0	7
4	0	0	0	9	0	0
5	8	0	0	0	0	0

$A + B \rightarrow C =$

	1	2	3	4	5	6
1	0	0	0	6	0	0
2	0	10	0	0	5	0
3	0	2	2	5	0	7
4	0	0	0	9	0	0
5	12	0	0	0	0	0

Non-zero element → 9

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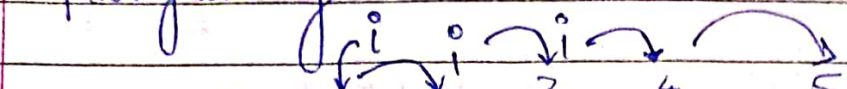
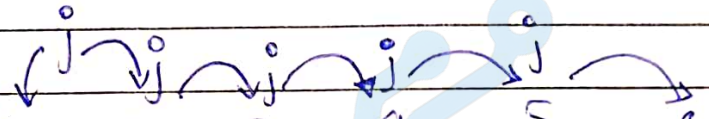
Adding using Coordinate



	0	1	2	3	4	5
A	5	1	2	3	3	5
	6	4	2	~~2~~	4	1
	5	6	(7)	~~2~~	5	4



	0	1	2	3	4	5	6
B	5	2	2	3	3	4	5
	6	2	5	3	6	4	1
	6	(3)	5	2	7	9	8

(rows[i]<row[j]

	0	1	2	3	4	5	6	7	8	9
C	5	1	2	2	3	3	3	~~3~~	~~4~~	5
	6	4	2	5	~~2~~	3	4	~~6~~	4	1
		6	10	5	2	2	5	~~7~~	9	12

rows[i] < rows[j]

copy[i] to c and i++;

column[i] < column[j]

copy[i] to c and i++;

~~rows[i] == rows[j] && column[i] == column[j]~~

Add both element

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rows[i] > row[j]

copy[j] to c and j++;

column[i] > column[j]

copy[j] to c and j++;

rows[i] == rows[j] && column[i] < column[j]

copy[i] to c and i++;

rows[i] == rows[j] && column[i] > column[j]

copy[i] to c and j++;

Array Representation of Sparse Matrix :-

$$A = \begin{array}{c|ccccc} & 1 & 2 & 3 & 4 & 5 \\ \hline 1 & 0 & 0 & 7 & 0 & 0 \\ 2 & 2 & 0 & 0 & 5 & 0 \\ 3 & 9 & 0 & 0 & 0 & 0 \\ 4 & 0 & 0 & 0 & 0 & 4 \end{array}$$

4 x 5

Non-zero element = 5

dimension and no zero element

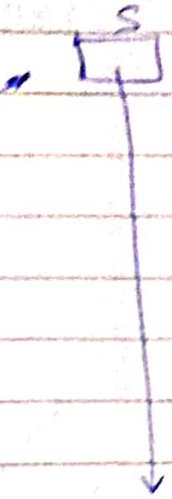
	0	1	2	3	4	5
i	4	1	2	2	3	4
Aj	5	3	1	4	1	5
x	5	7	2	5	9	4

max

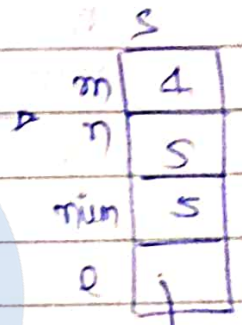
- How to represent
- How to create
- How to add

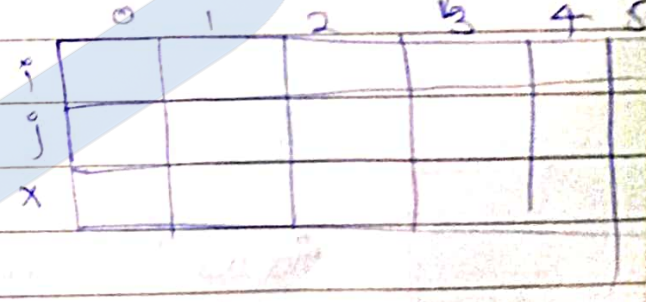
```
Struct Element
{
    int i;
    int j;
    int x;

};
```

```
struct Sparse
{
    int m;
    int n;
    int num;
    struct Element *e;   // Dynamic Array
}                        // Array of element

void create(struct sparse *s)
{
    printf("Enter Dimensions:");
    scanf("%d%d", &s->m, &s->n);
    printf("Enter no. of non-zero");
    scanf("%d", &s->num);

    s->e = new Element[s->num];

    printf("Enter all elements");

    for(int i=0; i<s->num; i++)
    {
        scanf("%d%d%d", &s->e[i].i,
                        &s->e[i].j,
                        &s->e[i].x);
    }
}
```




```
void main()
{
    struct sparse S;
    Create(&S);
}
```





Program for Adding Sparse Matrix →

S_1 =

	1	2	3	4	5
1	0	0	3	0	0
2	4	0	0	0	7
3	0	0	5	0	8
4	0	6	0	0	0

S_2 =

0	0	0	0	2
0	5	0	0	6
4	0	8	0	0
0	0	0	0	9

S_1

m	4
n	5
num	6
e	→

0	1	2	3	4	5
1	2	2	3	3	4
3	1	5	3	5	2
3	4	7	5	8	6

S_2

m	4
n	5
num	6
e	

0	1	2	3	4	5
1	2	2	3	3	4
5	2	5	1	3	5
2	5	6	4	8	9

```
void add(struct sparse *S1, struct sparse *S2)
{
    if (S1->m != S2->m || S1->n != S2->n)
    {
        return 0;
    }
    sparse *sum;

    sum = new sparse;

    sum->m = S1->m; sum->n = S2->n;
```

sum → m, n, num, e

```
sum->e = new Element[s1->num + s2->num];

while( i<s1->num && j<s2->num)
{  if(s1.e[i].i < s2
   if(s1->e[i].i < s2->e[j].i)
   {
        sum->e[k++] = s1->e[i++];

   }
   else if(s1->e[i].i > s2->e[j].i)
   {
        sum->e[k++] = s2->e[j++];

   }
   else
   {
        if(s1.e[i].j < s2.e[j].j)
        {

             sum->e[k++] = s1->e[i++];

        }
        if(s1.e[i].j > s2.e[j].j)
        {
             sum->e[k++] = s2.e[j++];

        }
        else
        {
             sum.e[k++] = s1->e[i++];
             sum->e[k++].xf= s2->e[j++].xs
        }
}
```

Polynomial Representation

coefficient, exponent

$$p(x) = 3x^5 + 2x^4 + 5x^2 + 2x + 7$$

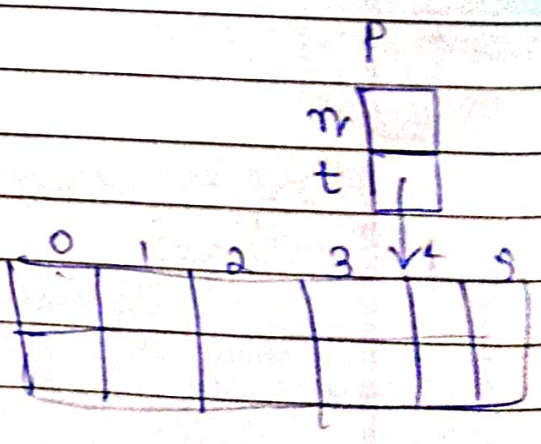
variable

1) Polynomial Representation
2) Evaluation of Polynomial
3) Addition of two polynomials

$n = 5$

coeff	3	2	5	2	7
Exp	5	4	2	1	0

```
struct Term
{
    int coeff;
    int expo;
};

struct Poly
{
    int n;
    struct Term *t;
};
```



```
struct Poly P;
printf("No. of non-zero terms");
scanf("%d", &p.n);

p.t = new term[p.n]

printf("Enter Polynomial terms");
for (i=0; i<p.n; i++)
{
    printf("Term no %d", i+1)
    scanf("%d %d", &p.t[i].coff,
                   &p.t[i].expo);
}
```

Polynomial evaluation →

	p		0	1	2	3	4	
n	5	coff	3	2	5	2	7	
t		exp	5	4	2	1	0	

```
struct Poly P;
int sum=0;
for(i=0; i<p.n; i++)
{
sum+= p.t[i].coff * Pow(x, p.t[i].exp);

}
return sum;
```

Polynomial Addition :-

$P1(x) = 5x^4 + 2x^2 + 5$

P1

		0	1	2
n = 3	coeff	5	2	5
t →	exp	4	2	0

i i

$P2(x) = 6x^4 + 5x^3 + 9x^2 + 2x + 3$

P2

		0	1	2	3	4
n = 5	coeff	6	5	9	2	3
t →	exp	4	3	2	1	0

j j

P3

		0	1	2	3	4
n	coeff	11	5	11	2	8
t →	exp	4	3	2	1	0

K

Algorithm

```
i=j=k=0;

while(i<p1.n && j<p2.n)
{

    if( p1.t[i].Exp > p2.t[j].Exp)
    {

        p3.t[k++] = p1.t[i++] };

    }
    else if( p2.t[j].Exp > p1.t[i].Exp)
    {

        p3.t[k++] = p2.t[j++];

    }
```

```
else
{
    P3.t[k++].Exp = P1.t[i].Exp;
    P3.t[k++].coeff = P1.t[i++].coeff +
                      P2.t[j++].coeff;
}
}
```