

Section-11

of
Linked Lists

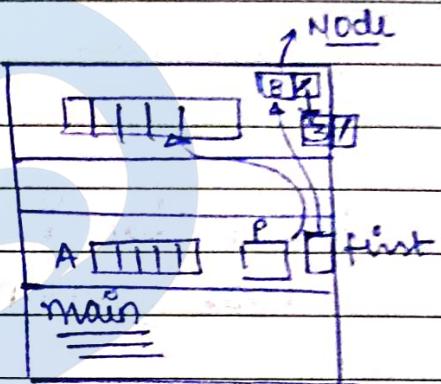
Keep practice of linked lists.

why we need Dynamic Data Structure Linked List

i) Problem with Arrays

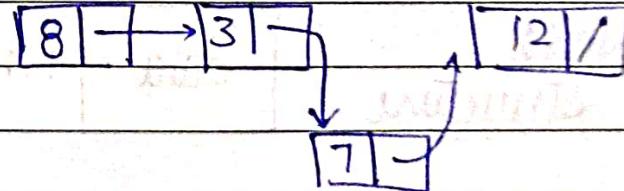
ii) Difference between Array & Linked List

`int A[10];` → Not available to store
more 10 element in this
array.



`int A[5];`
`int *p=new int[5];`

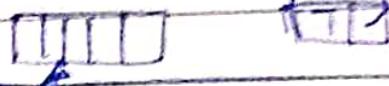
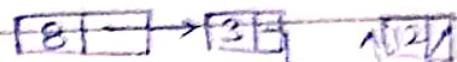
Together, element space and pointer space together, we call
it as a Node.



(Insert) between
3 and 12 node

Linked List

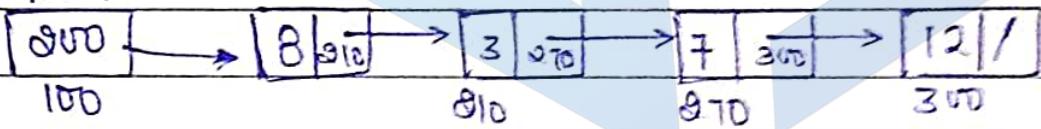
- what is linked list?
- what is Node?
- Node Structure
- Create a Node
- Access Node



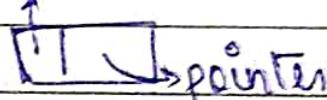
main

Linked list is a collection of nodes where each node contains data and points to next node.

first

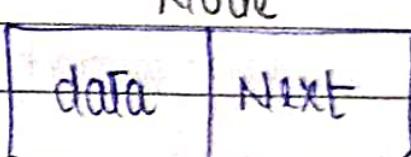


data



use of linked list :-

Node
Structure



- stacks and queue
- image browser
- music player
- Browser
- Hashmap and Hashset

class → by default → private
struct → by default → public

C Program / C++ Program

struct Node

}

2 — int data;

2 — struct Node *next;

}

↓

4 bytes

↳ size of nodes

→ This type of structure is
called as self-referential
structure

struct Node *p;

p = (struct Node*) malloc (size of (struct Node));

p = new Node;

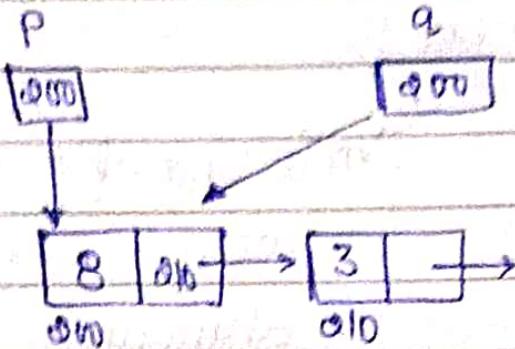
p → data = 10;

p → next = 0;

P
500

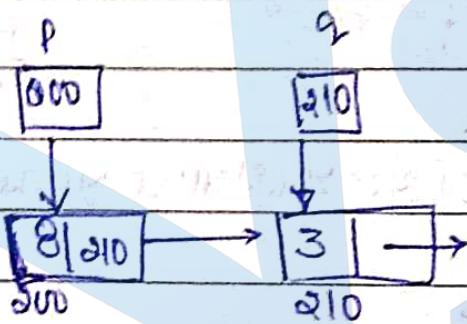
10 | /
500

More about linked lists +

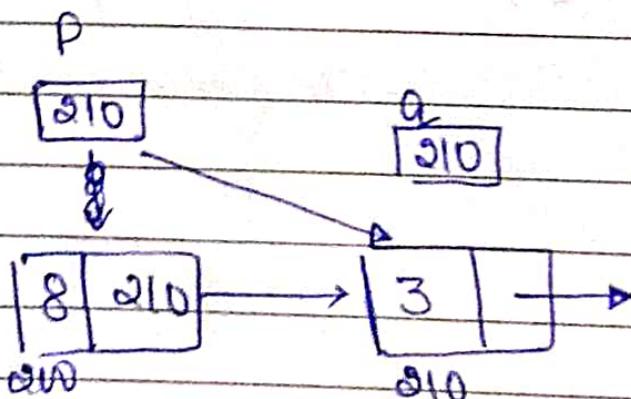


$q = p \rightarrow$

$q = p \rightarrow \text{next}$



$p = p \rightarrow \text{next}$



1.

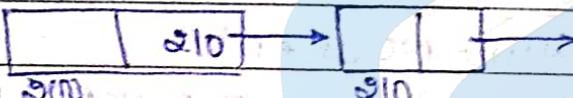
P



struct Node *p=NULL;

2.

P



if (p == NULL) { condition 1 }

if (p == 0) { condition 1 }

if (!p) { condition 1 }

if (p != NULL) { condition 2 }

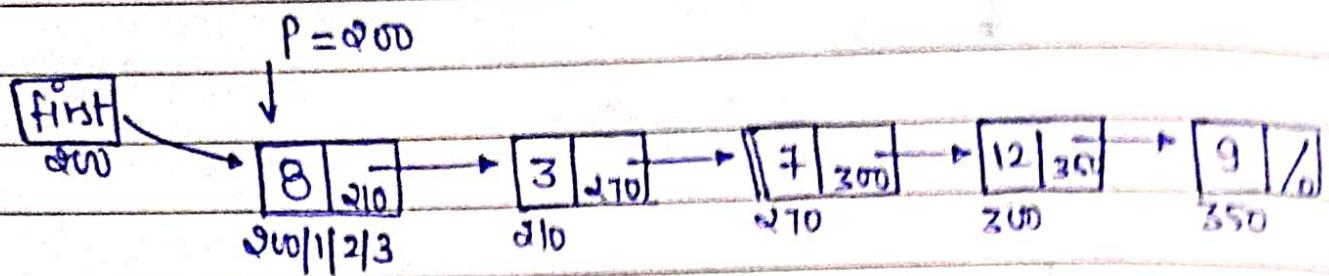
if (p != 0) { condition 2 }

if (p) { condition 2 }

if (p -> next == NULL) { Means last node of
linked list }

if (p -> next != NULL) { means some node left in
linked list }

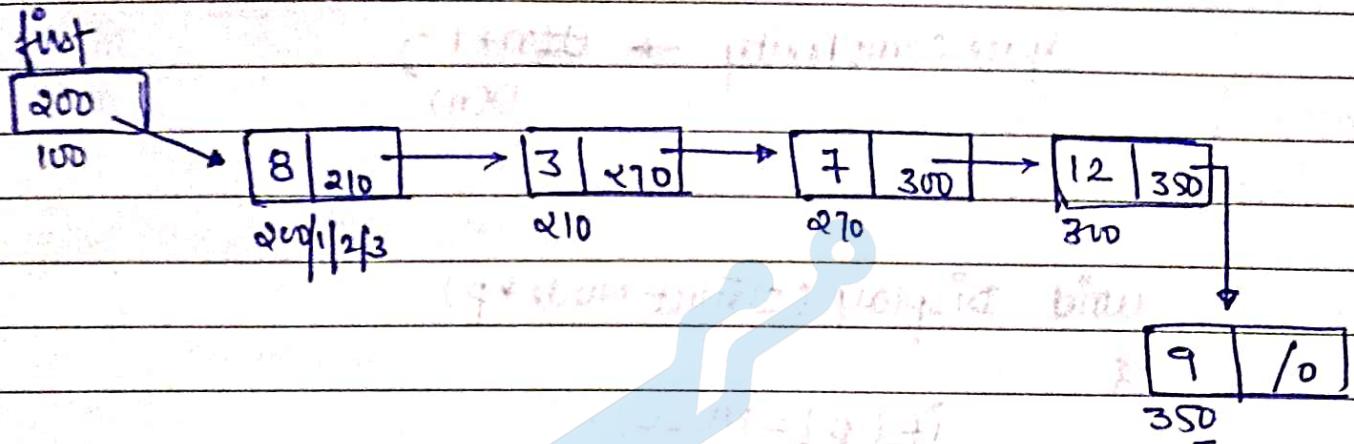
Display Linked List :-



```
// struct Node *p = first;  
display(struct Node *p){  
    while(p != 0) ← तक प नुमा रहे हो जाए  
        {  
            cout << p->data;  
            p = p->next;  
        }  
}
```

display(first);

Recursive Display of linked list :-

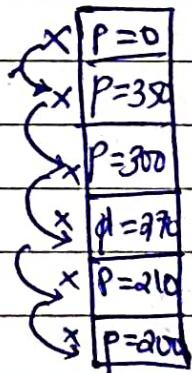


```
void Display(struct Node *p)
```

```
{
```

```
    if (p != null)
    {
        cout << p->data;
        display(p->next);
    }
}
```

Stack



$d(200)$

8

$d(210)$

3

$d(210)$

7

$d(300)$

12

$d(350)$

9

$d(0)$

X

Ques

$0/p \rightarrow 8\ 3\ 7\ 12\ 9$

Time complexity $\rightarrow O(n)$

Space complexity $\rightarrow \Theta(n+1) \downarrow O(n)$

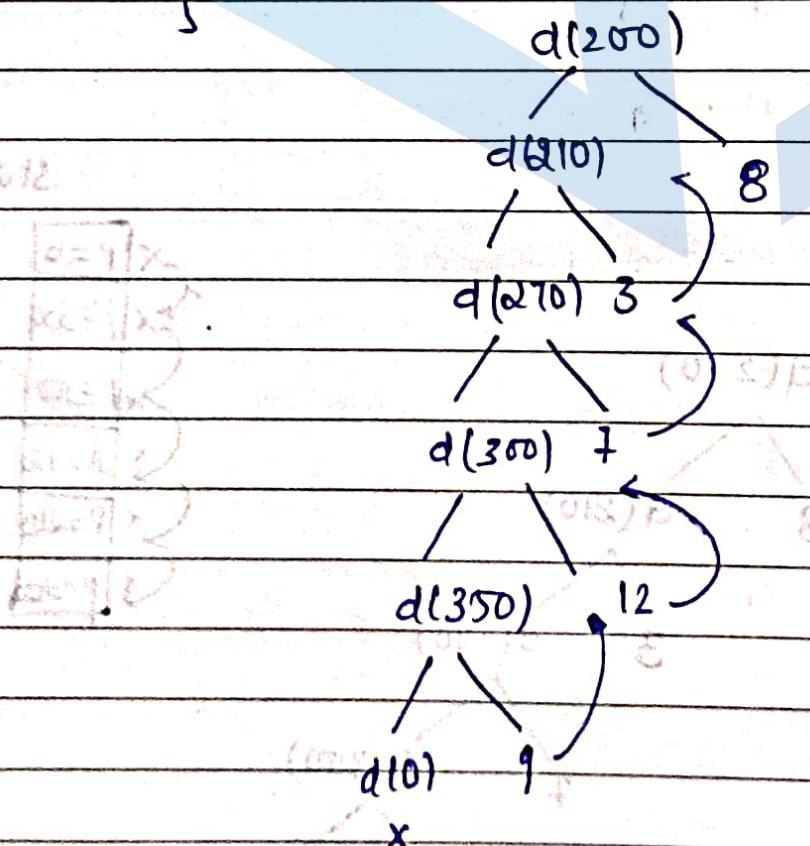
void Display (struct Node *p)

if (p == NULL)

{
 Display (p → next);

 cout << p → data; }

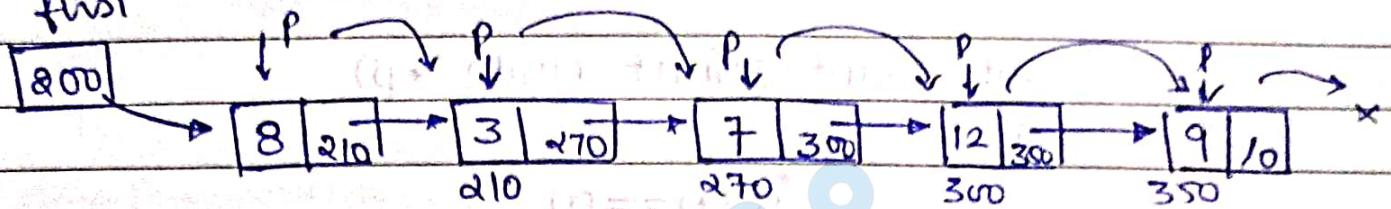
}



Output: 9 12 7 3 8

Counting Nodes in a linked list →

first



count = 0;
for (p ≠ NULL)

int count (struct Node *p)

{

 int c = 0;

 while (p != 0)

 {

 c++;

 p = p->next;

 }

 return c;

}

Time complexity → O(n)

Space complexity → O(1)

Recursive function for count

```
int count(struct Node *p)
```

```
if (p==0)
```

```
return 0;
```

```
else
```

```
return count(p->next)+1;
```

```
}
```

$$\text{count}(200) = ?$$

$$c(210) + 1 \Rightarrow 5$$

$$c(270) + 4 \Rightarrow 4$$

$$c(300) + 1 \Rightarrow 3$$

$$c(350) + 1 \Rightarrow 2$$

$$c(0) + 1 = 1$$

X

↓

0

Time $\rightarrow O(n)$

Space $\rightarrow O(n)$

```
int count (struct Node *p)
```

```
{
```

```
    if (p == NULL)
```

```
{
```

```
        return 1 + count (p->next);
```

```
}
```

```
else
```

```
{
```

```
    return 0;
```

```
}
```

```
}
```

```
int count (struct Node *p)
```

```
{
```

```
    int x = 0;
```

```
    if (p)
```

```
{
```

```
        x = count (p->next);
```

```
    return x + 1;
```

```
}
```

```
else
```

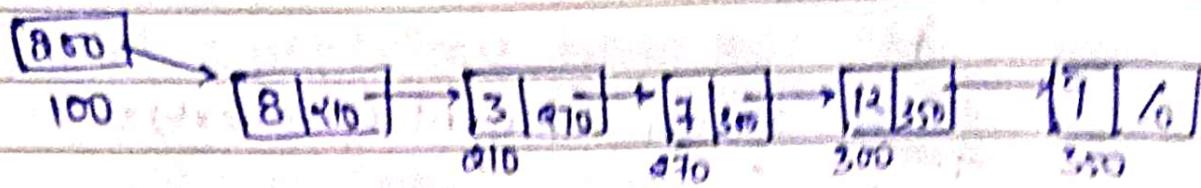
```
{
```

```
    return 0;
```

```
}
```

```
}
```

Sum of All Element in a Linked List



```
int Add (struct Node *p)
```

```
{
```

```
    int sum=0;  
    while(p)
```

```
{
```

```
        sum=sum+p->data;
```

```
        p=p->next;
```

```
}
```

```
return sum;
```

```
3
```

Time complexity $\rightarrow O(n)$

Space $\rightarrow O(1)$

recursive function

int Add(struct Node *p)

{

if (p == 0)

{

return 0;

{
else

}

return Add(p → next) + p → data;

}

Time $\rightarrow O(n)$

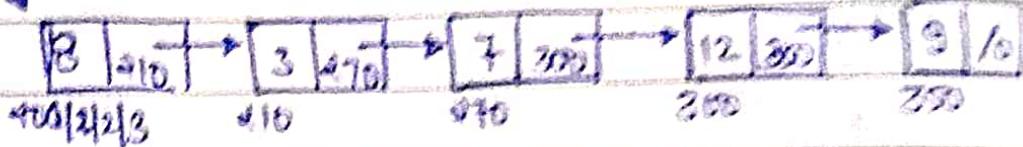
Space $\rightarrow O(1)$

Maximum Element in a Linked List

first

300

100



$$\text{max} = \text{INT_MIN};$$

$$\text{max} = -32768$$

int max(struct Node *p)

{

int max = INT_MIN;

while(p)

{

if (p->data >= max)

{

max = p->data;

}

p = p->next

}

return max;

{

$T \rightarrow O(n)$

$S \rightarrow O(1)$

Recursive function

```
int max(Node *p){
```

```
    if(p==0)
```

```
        return INT-MIN;
```

```
    int x=0;
```

```
    else
```

```
}
```

```
    x=max(p->next);
```

```
    if(x>p->data)
```

```
{
```

```
    return x;
```

```
}
```

```
else
```

```
{
```

```
    return p->data;
```

```
}
```

```
}
```

```
(p->data>x)?x:p->data;
```

```
?
```

```
(x>p->data)?x:p->data;
```

```
int max1(struct Node *p)
```

```
{
```

```
    int x=0;
```

```
    if(p==0)
```

```
        return MIN-INT;
```

```
x=max(p->next);
```

```
return x>p->data? x : p->data;
```

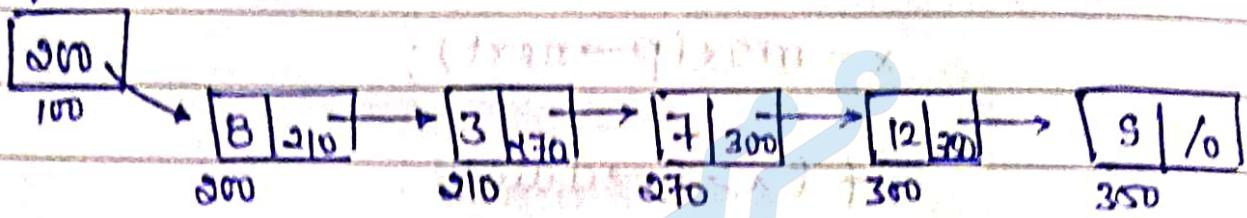
```
}
```

Searching in a linked lists :-

i - Linear search

ii - Binary search (Not suitable
for Linked lists)

first



key = 12

Node *

int search(struct Node *p, int key)

{
 while(p != NULL)

{

if (p->data == key)

{

return (p);

}

p = p->next;

}
return 0;

Recursive function \rightarrow

node * search (struct node * p, int key)

{

if ($p == \text{NULL}$)

{

return NULL;

}

if ($\text{key} == p \rightarrow \text{data}$)

{

return p;

}

return search($p \rightarrow \text{next}$, key);

$T \rightarrow O(1)$
 $S \rightarrow O(n)$

($p \neq \text{NULL}$, $q \neq \text{NULL}$) where q is

$\text{NULL} = 0 \times \text{address}$

$\text{NULL} = \{\}$ direct access

($q \rightarrow \text{data} == p \rightarrow \text{data}$) \Rightarrow

$\text{data} == \text{data}$

$\text{data} == \text{data}$

Improving Searching in Linked List :-

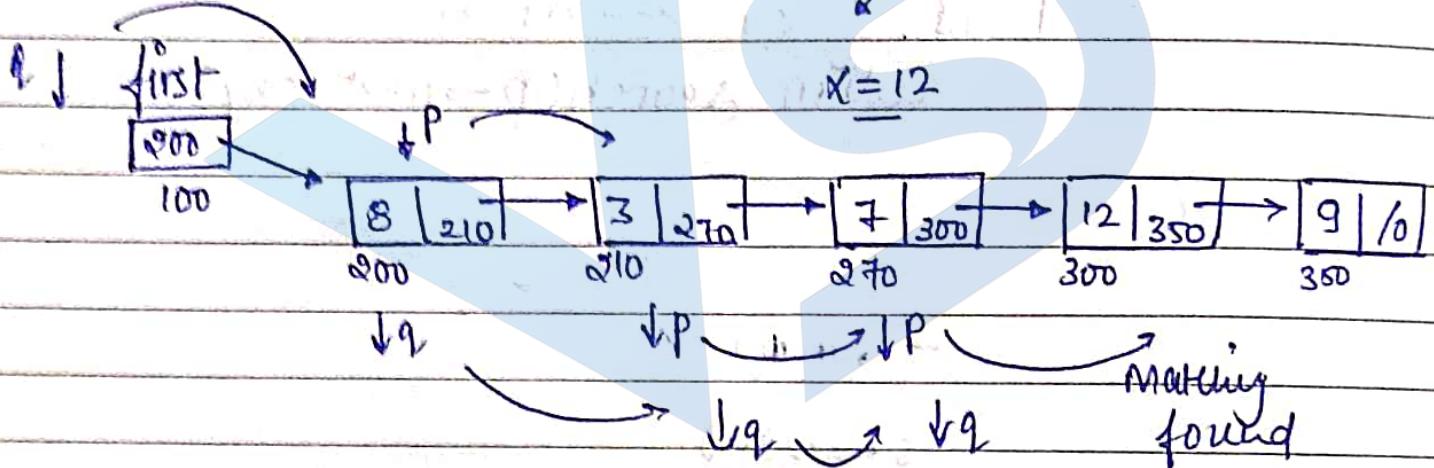
Improving the linear search

Two methods

Transportation

Transposition

move to head



Node* search (Node* p, int key)

{

 Node *q = NULL;

 while ($p \neq \text{NULL}$)

{

 if ($\text{key} == p \rightarrow \text{data}$)

{

$q \rightarrow \text{next} = p \rightarrow \text{next};$

$p \rightarrow \text{next} = \text{first};$

$\text{first} = p;$

}

$q = p;$

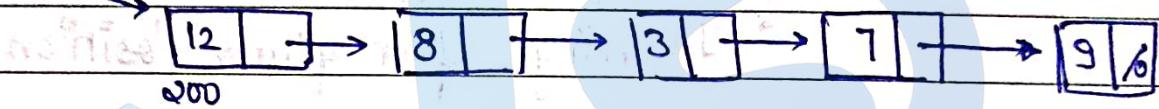
$p = p \rightarrow \text{next};$

}

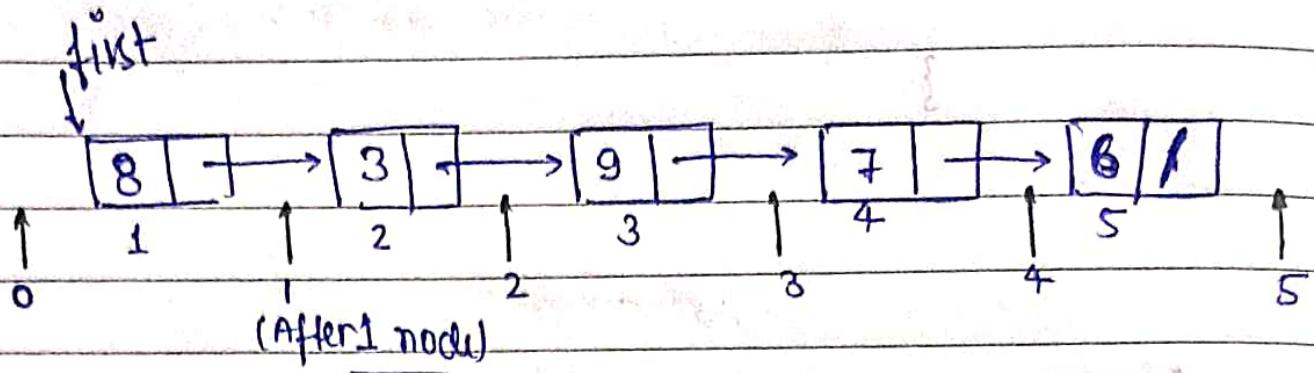
After search

first

200



Inserting in a linked list :-



1. Insert before first

2. Inserting after given Position

- Insert before first →



O(1)

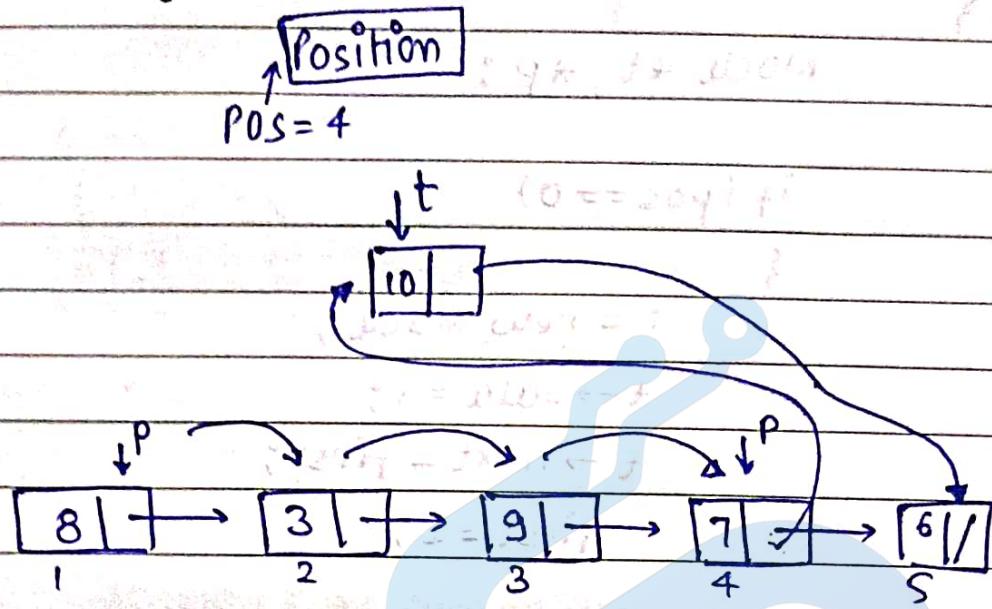
Node *t = new Node;

t->data = x;

t->next = first;

first = t;

• Inserting after given Position



```
Node *t = new Node;
```

```
t->data = x;
```

O(n)

```
p = first;
```

```
for(int i=0 ; i<pos-1; i++)
```

min $\rightarrow O(1)$

max $\rightarrow O(n)$

```
{ p = p->next;
```

```
}
```

```
t->next = p->next;
```

```
x = p->data;
```

$t = t \times 10 + x$ \Rightarrow $t = t \times 10 + p->data$

$t = t \times 10 + q \Rightarrow t = t \times 10 + t \times 10 + p->data$

```
void insert (int pos, int x)
```

```
{
```

```
Node *t, *p;
```

```
if (pos == 0)
```

```
{
```

```
t = new Node;
```

```
t->data = x;
```

```
t->next = first;
```

```
first = t;
```

```
}
```

```
else if (pos > 0)
```

```
{
```

```
p = first;
```

```
for (i=0; i<pos-1 && p; i++)
```

```
{
```

```
p = p->next;
```

```
}
```

```
if (p)
```

```
{
```

```
t = new Node;
```

```
t->data = x;
```

```
t->next = p->next;
```

```
p->next = t;
```

```
}
```

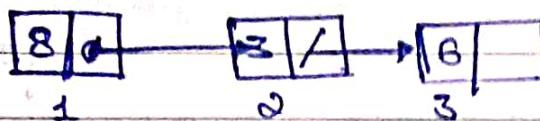
```
}
```

```
}
```

Creating a linked list using insert →

first

'



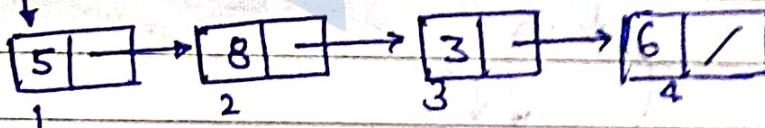
Insert(0, 8)

Insert(1, 3)

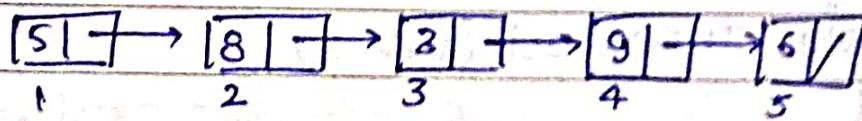
Insert(2, 6)

Insert(0, 5)

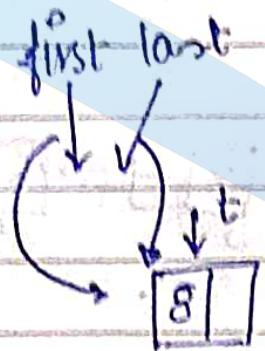
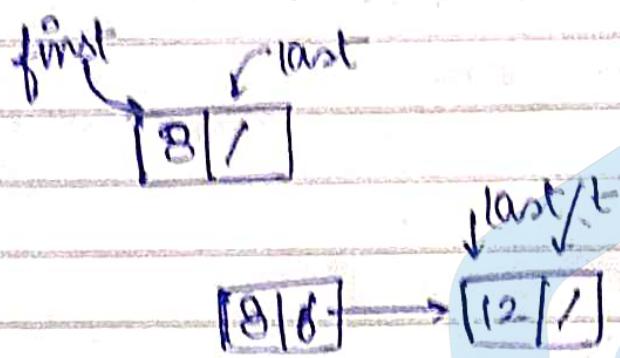
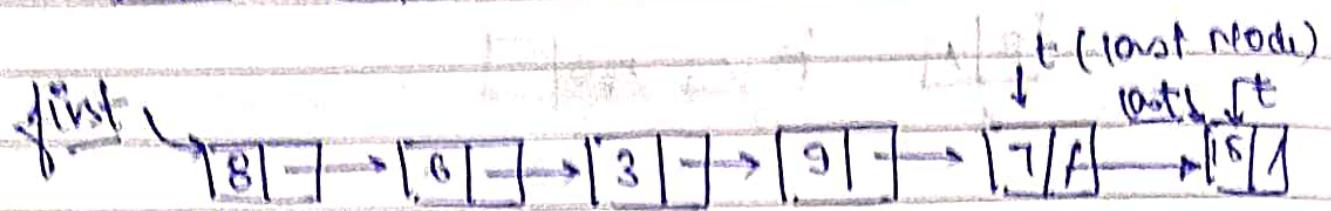
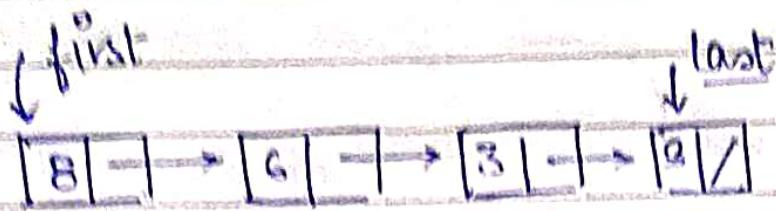
first



Insert(3, 9)



Creating a linked list by inserting at last → ~~initial~~



Void InsertLast(int x)

{

Node *t = new Node;

t → data = x;

t → next = NULL;

if (first == NULL)

{

first = last = t;

}

else

{

last → next = t;

last = t;

}

insertLast(8)

first last

[8] []

(9 words) 8125
(9 words) 8124
(9 words) 8123
(4 words) 8120

insertLast(3)

first last

[8] → [3] []

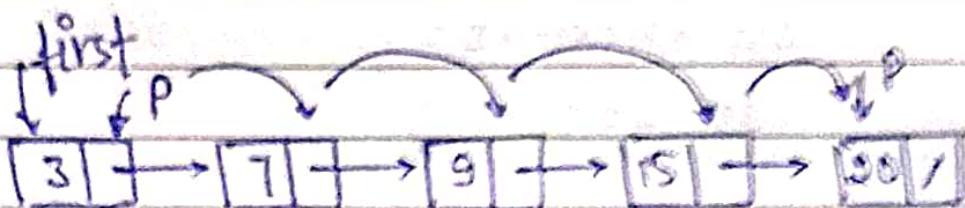
insertLast(9)

first

last

[8] → [3] → [9] []

Inserting in a Sorted linked lists :-



$x = 18 \rightarrow (\text{insert})$

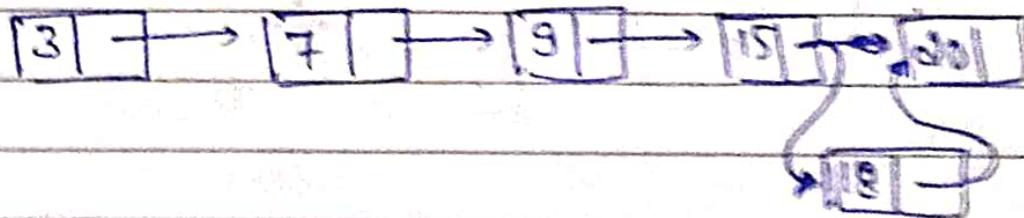
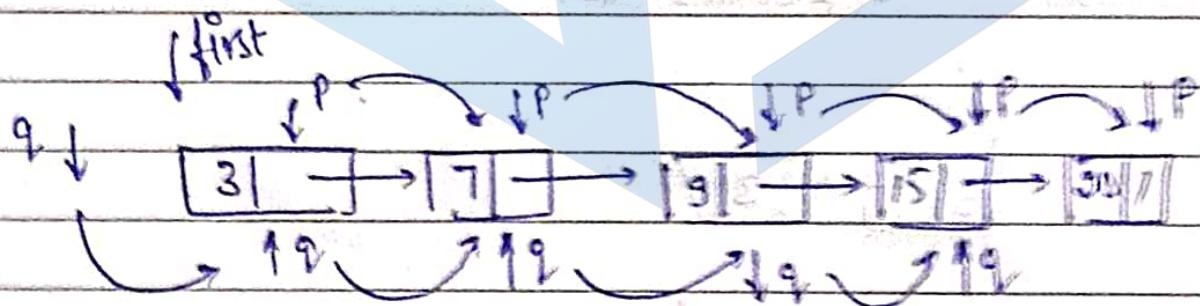
$3 < 18$ (move P)

$7 < 18$ (move P)

$9 < 18$ (move P)

$15 < 18$ (move P)

$20 > 18$ (stop P)



$p = \text{first};$
 $q = \text{NULL};$

while ($p \neq \text{NULL}$ & $p \rightarrow \text{data} < x$)

{

$q = p;$

$p = p \rightarrow \text{next};$

$\min \rightarrow O(1)$

$\max \rightarrow O(n)$

}

$t = \text{new Node};$

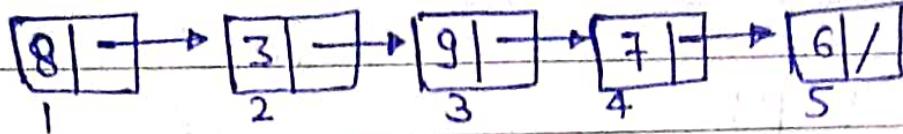
$t \rightarrow \text{data} = x;$

$t \rightarrow \text{next} = q \rightarrow \text{next};$

$q \rightarrow \text{next} = t;$

Deleting from linked lists →

↓
first



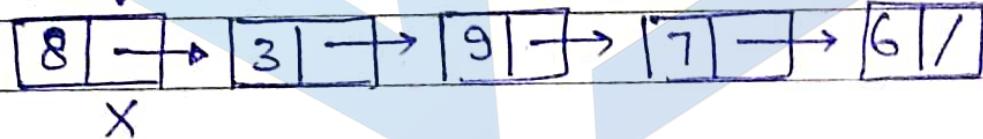
(Node → next)

1. Delete first Node

2. Delete a Node at given position

• Delete first Node →

↓ first ↓ p ↓ first



Node *p = first;

first = first → next;

x = p->data;

delete p;

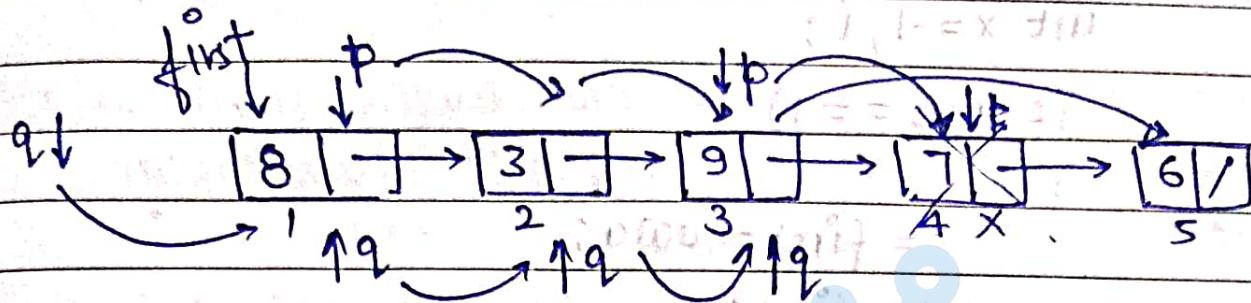
Time → O(1)

- Delete a node at given position

(e.g. 4th node from first)

1. P = 4th node

[Pos = 4]



Node *p = first;

Node *q = NULL;

Max $\rightarrow O(n)$

Min $\rightarrow O(1)$

for (i=0; i<pos-1; i++)

{

 q = p;

 ++i; q->p = p->next; i

}

 q->next = p->next;

 x = p->data;

 delete p;

int delti(int pos)

}

Node *p, *q;

int x=-1, i;

if (pos == 1)

{

x = first->data;

p = first;

first = first->next;

delti p;

}

else

{

p = first;

q = NULL;

for (i=0; i<pos-1 && p; i++)

{

q = p;

p = p->next;

}

if (p)

{

q->next = p->next;

x = p->next;

delti p;

}

}

return x;

}

Important!!!

Why we need linked list?

→ We need linked list when we want to store and manipulate data that has an unknown or variable size. Linked list is a dynamic data structure that can grow and shrink at runtime by allocating and deallocating memory. This means that we don't need to specify the initial size of linked list nor waste memory for unused elements.

Difference between Array and linked list

Array

- Arrays are stored in contiguous location.

- fixed in size.

- Memory is allocated at compile time.

- Use less memory than linked lists.

linked lists

- linked list are not stored in contiguous location.

- Dynamic size.

- Memory is allocated at run time.

- Use more memory because it stores both data and the address of next node.

linked list

- Elements can be accessed easily.

- Element accessing requires the traversal of whole linked list.

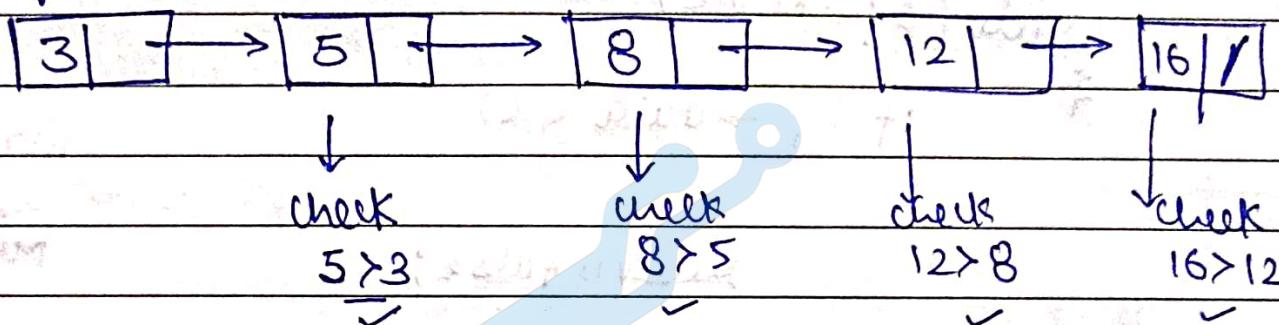
- Insertion and deletion operation takes time.
- Insertion and deletion operation is faster.

why we are not using vector array because when a vector is resized, it may be necessary to copy all of the elements to a new memory location. This can be time consuming, especially for large vectors.

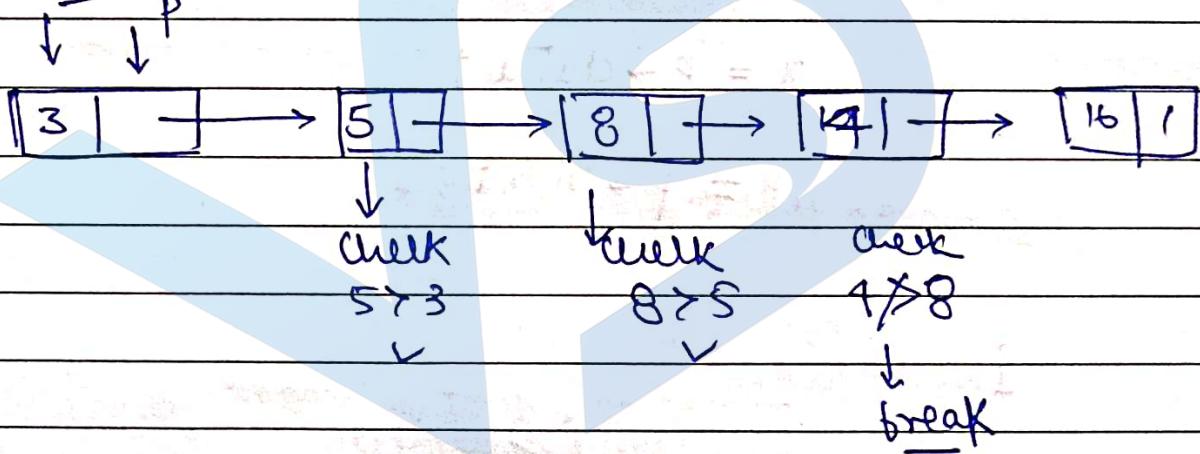
check if

check if a linked list is sorted

first



first



$$x = -32768; 388 \not\leq 16$$

$$3 > \underline{-32768}$$

$$5 > 3$$

$$8 > 5$$

$$12 > 8$$

$$\underline{16 > 12}$$

int x = INT_MIN;

node *p = first;

while ($p \neq \text{NULL}$)

{ if ($p \rightarrow \text{data} < x$)

return false;

max $\rightarrow O(n)$

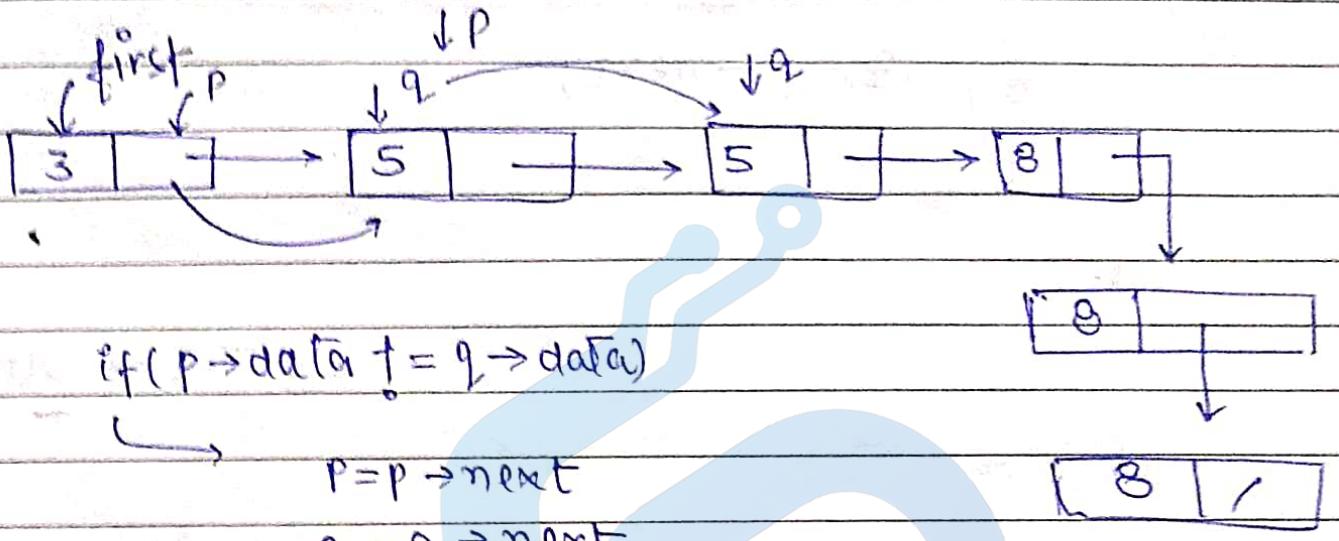
min $\rightarrow O(1)$

$x = p \rightarrow \text{data};$

$p = p \rightarrow \text{next};$

}

Remove Duplicates from Sorted Linked List



```
node *p = first;
node *q = first → next;
while (q != NULL)
{ if (p → data != q → data)
    {
        p = q;
        q = q → next;
    }
    else
    {
    }
}
```

$p \rightarrow \text{next} = q \rightarrow \text{next};$

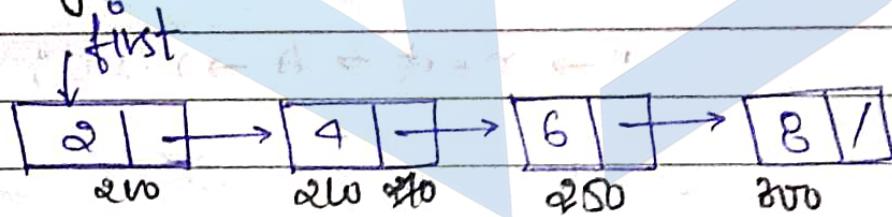
delete $q;$

$q = p \rightarrow \text{next};$

}

Time T.C $\rightarrow O(n)$

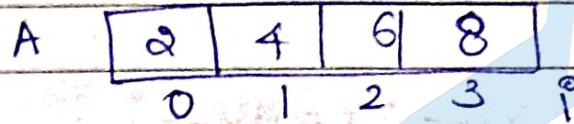
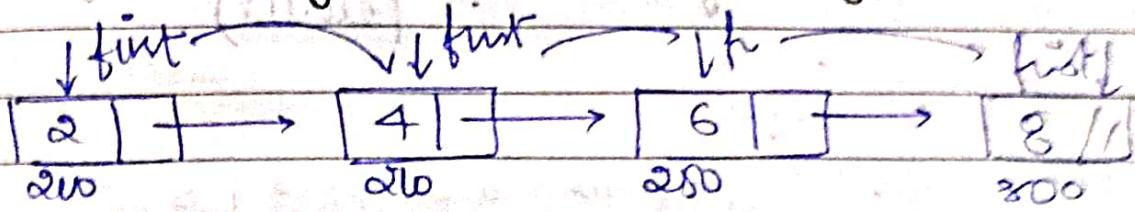
Reversing a linked list :-



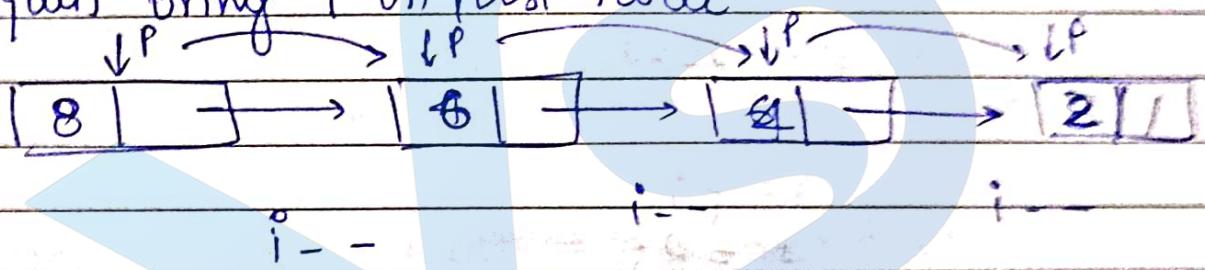
Methods

- 1. Reversing element (swapping)
- 2. Reversing links
- 3. Recursive Reverse

Method 1 :- Reversing element (swapping)



Again bring 'P' on first node



Node * p = first;

i = 0;

while (p != NULL)

T.C → O(n)

S.C → O(n)

```

    {
        A[i] = p->data;
        p = p->next;
        i++;
    }

```

p = first;
i--;

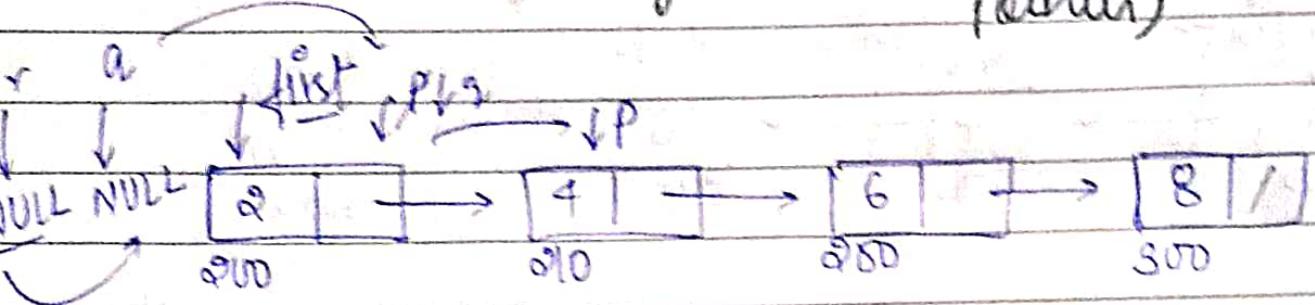
while (p != NULL)

```

    {
        p->data = A[i];
        i--;
        p = p->next;
    }

```

Method 2 :- Running links (Running using Sliding pointer)



Node $p = \text{first}$, $q = \text{NULL}$, $r = \text{NULL}$;

while ($p \neq \text{NULL}$)

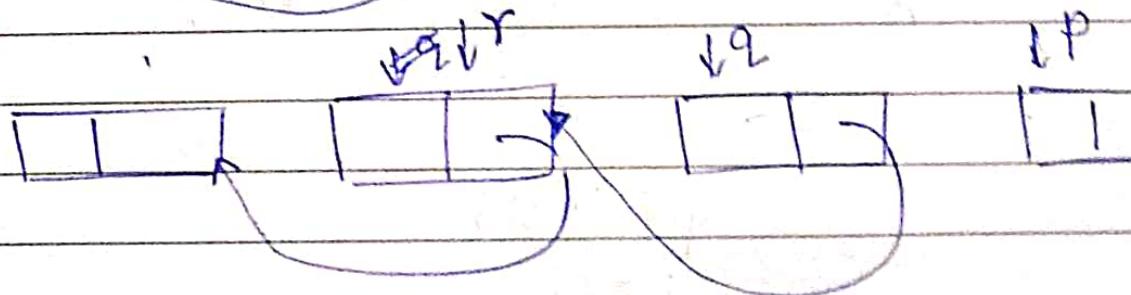
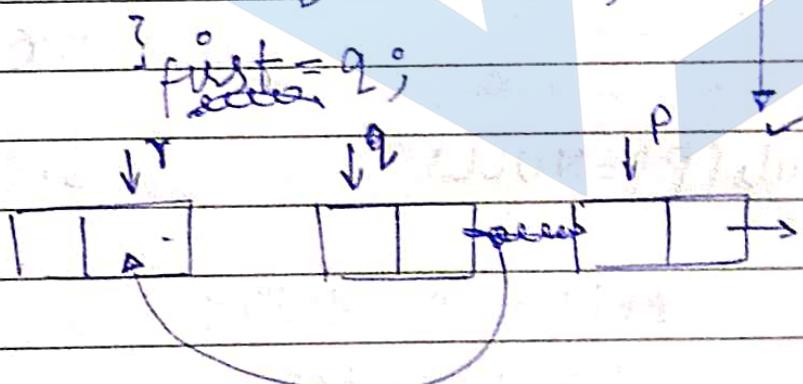
{

$r = p$;

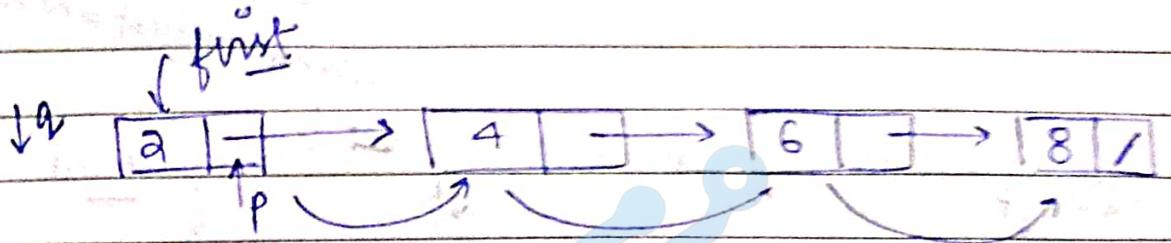
$q = p$;

$p = p \rightarrow \text{next}$;

$q \rightarrow \text{next} = r$;



Method 3:- Recursive Reverse for linked list



```
void reverse(Node *q, Node *p)  
{
```

```
    if (p != NULL)
```

```
{
```

```
    reverse(p, p->next);
```

~~↑ Now in reversing~~

```
    p->next = q;
```

```
}
```

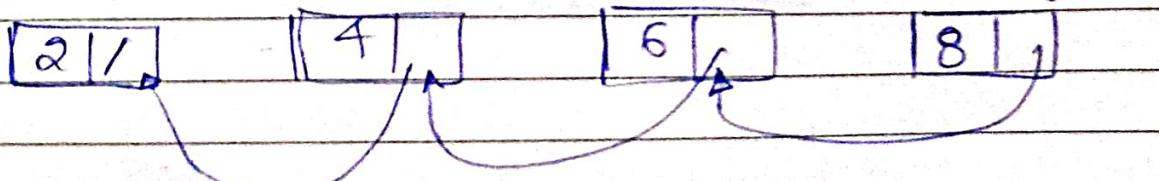
```
use q = first ->
```

```
{
```

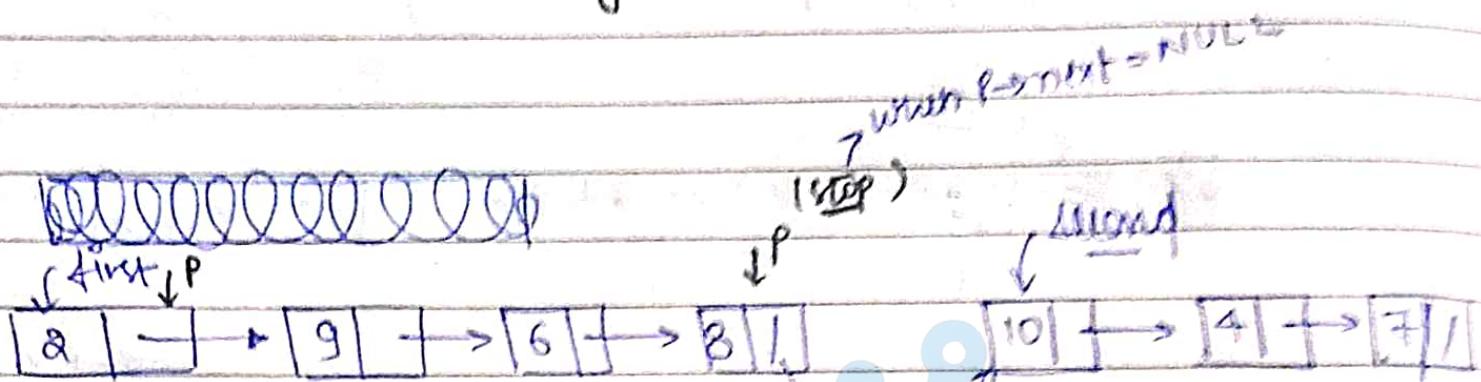
```
    first = q;
```

```
}
```

```
{
```



Concatenating of linked lists :-



Node $\&p = \text{first}$

while ($p \rightarrow \text{next} \neq \text{NULL}$)

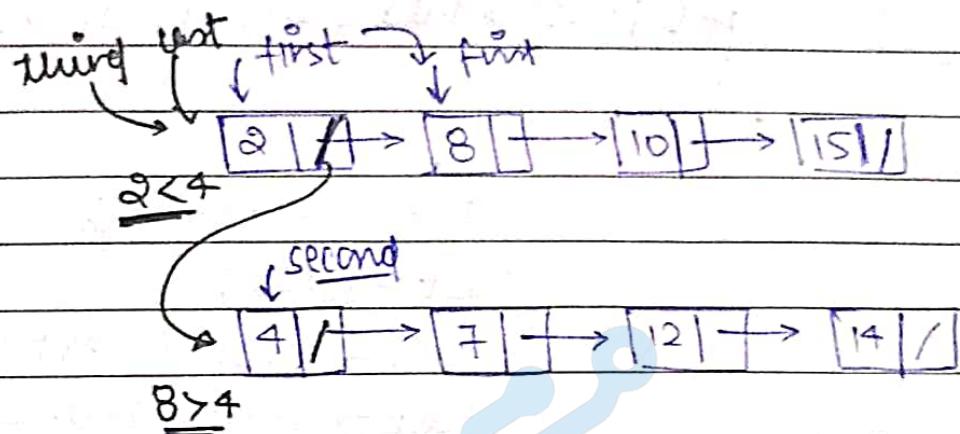
$O(n)$

$p = p \rightarrow \text{next};$

$p \rightarrow \text{next} = \text{second};$

$\text{second} = \text{NULL};$

Merging 2 linked lists :-



Merging → combining two sorted lists into

third (Point to first) last (last node of merge mixed list) single sorted linked lists

node:-

node * third, * last;

if (first → data < second → data)

{

 third = last = first;

 first = first → next;

 last → next = NULL;

}

else

{

 third = last = second;

 second = second → next;

 last → next = NULL;

{

S.C \rightarrow O(1)

T.C \rightarrow O(m+n)

while (first != NULL && second != NULL)

{

 if (first->data < second->data)

{

 last->next = first;

 last = first;

 first = first->next;

 last->next = NULL;

 else

{

 last->next = second;

 last = second;

 first = second = second->next;

 last->next = NULL;

}

}

 if (first != NULL)

{

 last->next = first;

{

 else

{

 last->next = second;

{

Check for loop in linked list :-

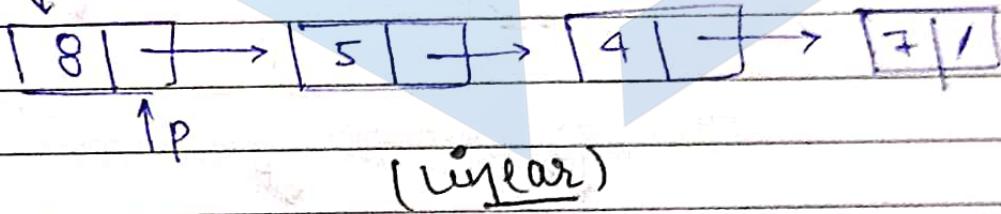
first



last node of linked list pointing some node
of the linked list.

→ loop in linked list

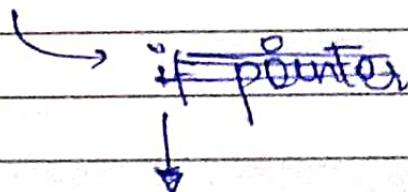
first



Checking linked list is linear →

↳ if pointer become NULL

Checking linked list is loop →



first / last

8 → 5 → 7 → 7 → 3 → 9

P → move one step

Q → move two steps

first P Q
8 → 5 → 7 → 7 → 3 → 9

Again move

P Q
8 → 5 → 7 → 7 → 3 → 9

P Q
8 → 5 → 7 → 7 → 3 → 9

P Q
8 → 5 → 7 → 7 → 3 → 9

Meeting again

$O(n) \rightarrow T.C$

if isLoop(NULL *f)

}

Node *p, *q;

p = q = f;

do

{

p = p->next;

q = q->next;

q = q != NULL ? q->next : NULL;

} while (p != q);

if (p == q)

{

return true;

}

else

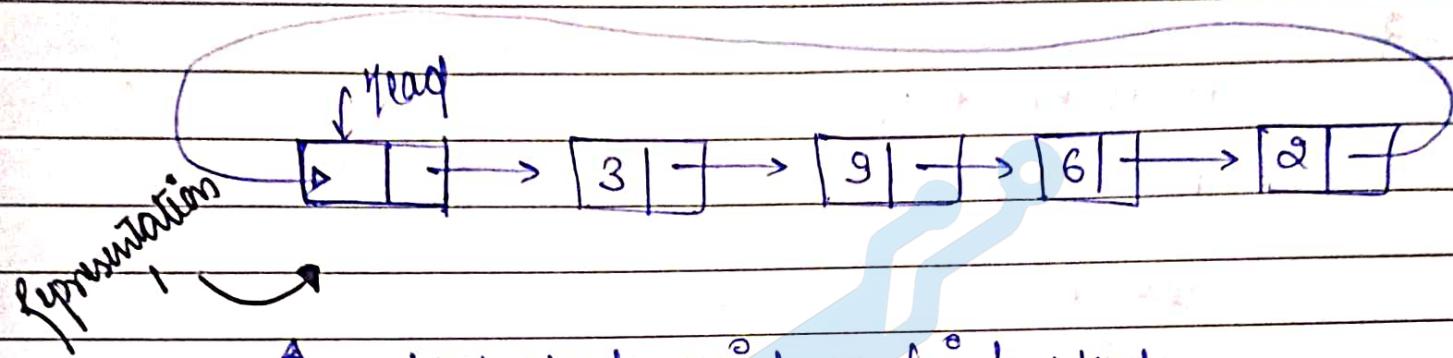
{

return false;

}

{

Circular linked list :-



Commonly used

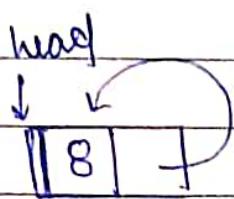
Last Node point on first Node

circular linked list

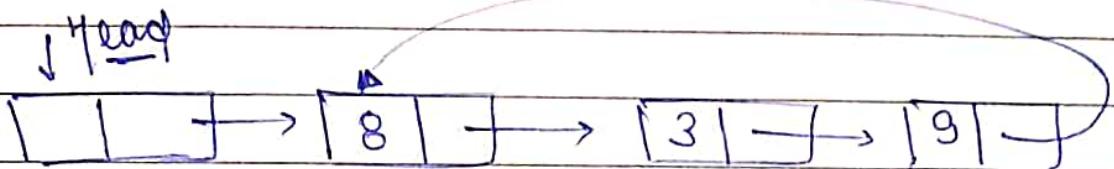
Collection of nodes which are circularly connected

head is one first and last node.

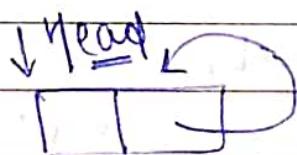
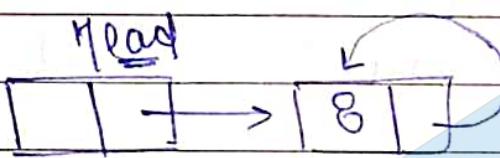
Benefits :-



head

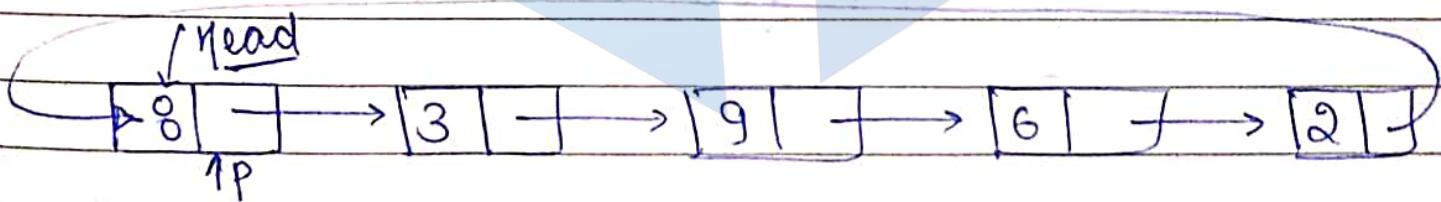


Representation 2



~~head~~
Outcome of
empty linked
list is NULL

Display a Circular linked list :-



void Display (Node *p)

{

do
while (p!=head)

{

cout << p->data;

p = p->next;

} while (p!=head);

}

Display (head)

Recursion Display function

```
void Display( Node *p )
```

```
{
```

```
    static int flag = 0;
```

```
    if ( p == head || flag == 0 )
```

```
{
```

```
        flag = 1;
```

```
        cout << p->data;
```

```
        Display( p->next );
```

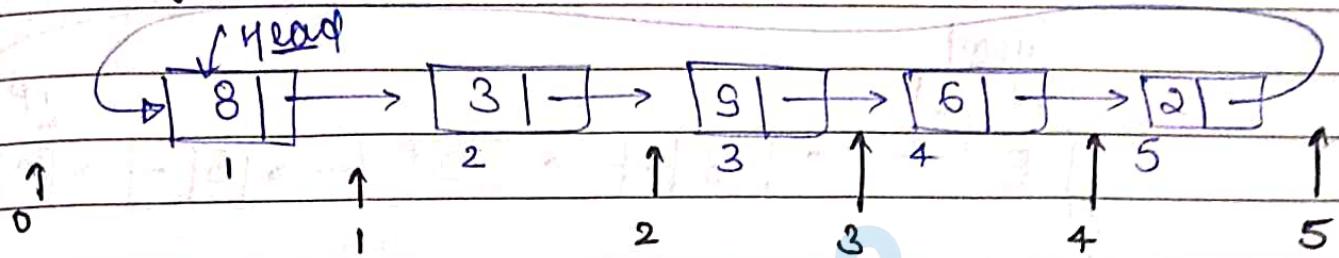
```
- }
```

```
flag = 0;
```

```
}
```

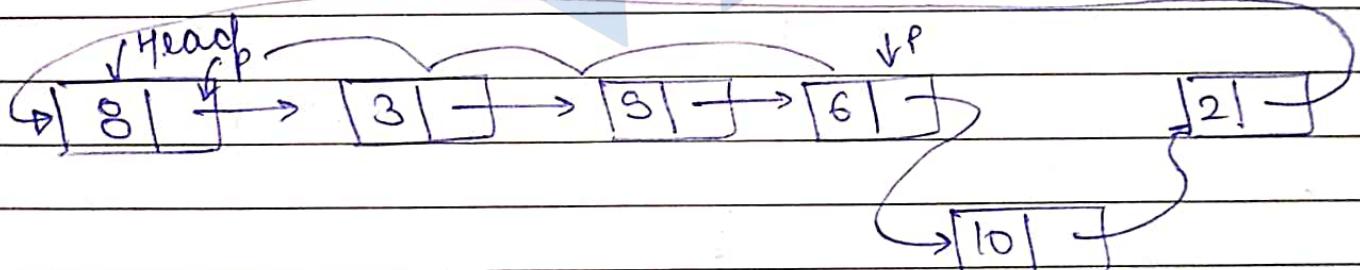
```
Display( head );
```

Inserting in a Circular Linked List :-



- Insert before head
- Insert at any other position

~~• Insert before head~~ • Insert at any other position



Node *t;

Node *p = Head;

for(int i=0; i<pos-1; i++)

p = p->next;

t = new Node;

t->data = x;

t->next = p->next;

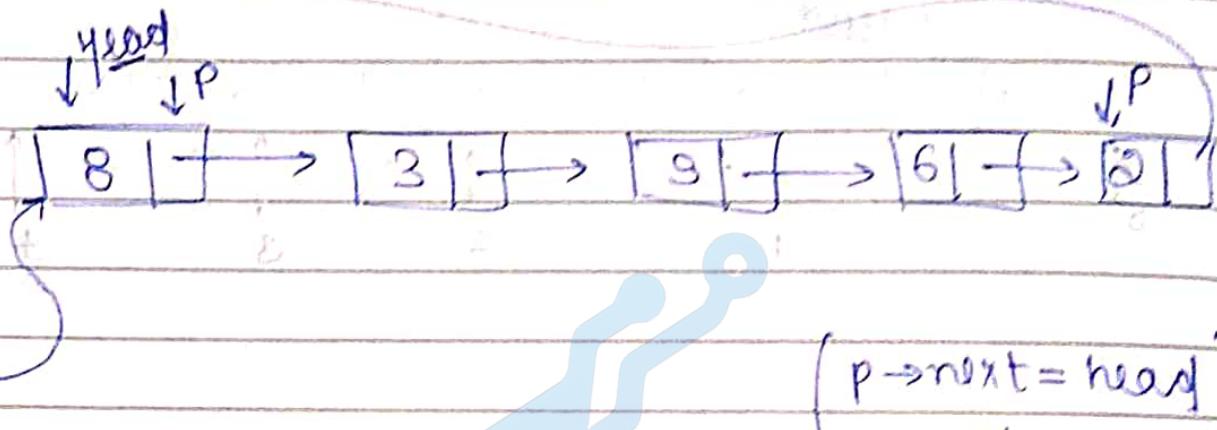
p->next = t;

p->next = t;

t->next = p->next;

p->next = t;

- Insert before head node:



$\text{node } *p = \text{head};$

$\text{node } *t = \text{newNode};$

$t \rightarrow \text{data} = x;$

$t \rightarrow \text{next} = \text{head};$

while($p \rightarrow \text{next} \neq \text{head}$)

}

$p = p \rightarrow \text{next};$

}

$p \rightarrow \text{next} = t;$

$\text{head} = t;$

combine

```
void Insert (int pos, int x)
```

```
{
```

```
Node *t, *p;
```

```
int i;
```

```
if (pos == 0)
```

```
t = new Node;
```

```
t->data = x;
```

```
if (head == NULL)
```

```
{
```

```
head = t;
```

```
head->next = head;
```

```
}
```

```
else
```

```
{
```

```
p = head;
```

```
while (p->next != head)
```

```
{
```

```
p = p->next;
```

```
{
```

```
p->next = t;
```

```
t->next = head;
```

```
head = t;
```

```
{
```

```
{
```

use

{

$p = \text{head};$
 $\text{for } (p \neq \text{null}, i=0; i < \text{pos}-1; i++)$

{

$p = p \rightarrow \text{next};$

{

$t = \text{new Node};$

$t \rightarrow \text{data} = x;$

$t \rightarrow \text{next} = p \rightarrow \text{next};$

$p \rightarrow \text{next} = t;$

{

}

copy

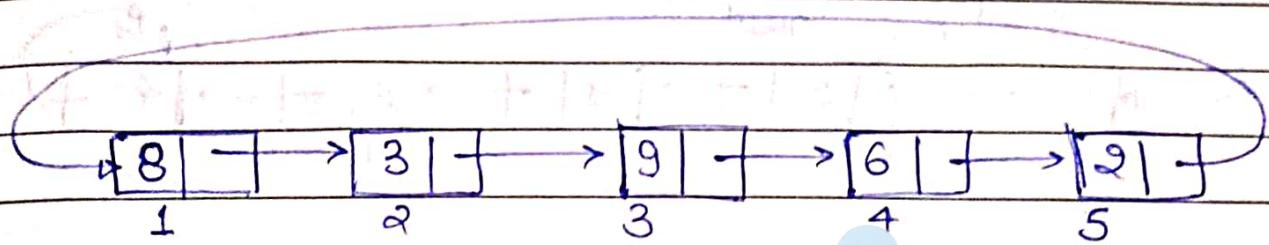
copy function

copy function

function = $\text{copy}(p)$

$t = \text{new Node}$

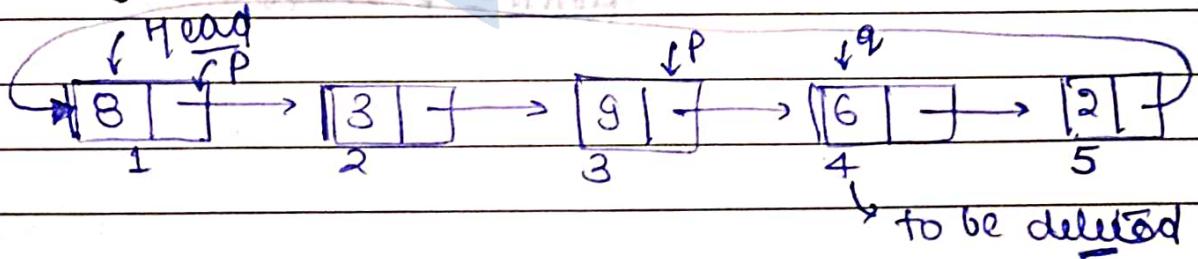
Deleting from Circular linked list :-



1. Deleting Head Node

2. Deleting a Node from given position.

• Deleting a Node from given position :-



P = Head;

for (i=0 ; i< pos-2 ; i++)

{

P = P->next;

}

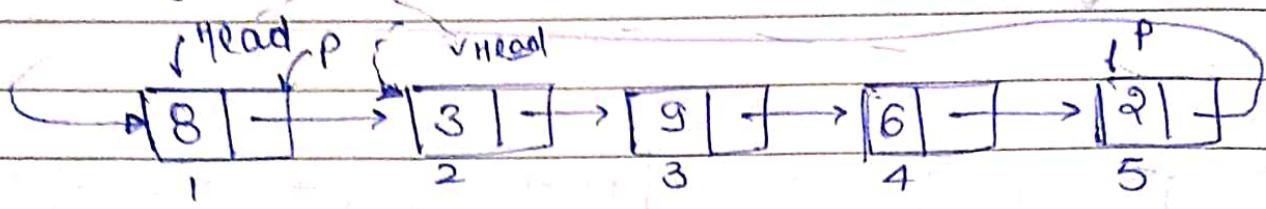
B) q = P->next;

P->next = q->next;

x = q->data;

delete q;

• Deleting Head Node :-



$p = \text{Head};$

$\text{while } (p \rightarrow \text{next} \neq \text{head})$

{

$p = p \rightarrow \text{next};$

3

$p \rightarrow \text{next} = \text{head} \rightarrow \text{next};$

$x = \text{head} \rightarrow \text{data};$

~~$\text{head} = p \rightarrow \text{next};$~~

~~$\text{delete head};$~~

$\text{Head} = p \rightarrow \text{next};$

int Delete(int pos)

{

Node *p, *q;

if (pos == 1)

{

p = Head;

while (p->next != Head)

p = p->next;

x = Head->data;

if (p == Head)

{

delete Head;

Head = NULL;

{

else

{

p->next = Head->next;

delete Head;

Head = p->next;

}

else

{

p = Head;

for (i=0; i<pos-2; i++)

p = p->next;

q = p->next;

x = q->next;

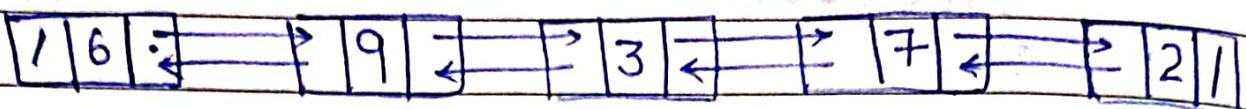
delete q;

} return x;

}

Doubly linked list :-

first



Node



Struct Node

{

struct Node *prev;

int data;

struct Node *next;

}

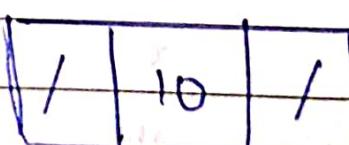
struct Node *t;

t = new Node;

t->prev=NULL;

t->data=10;

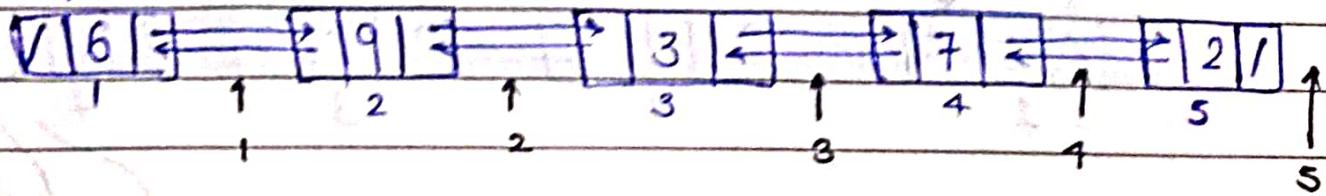
t->next=NULL;



t

Insert in a Doubly Linked List :-

(first)

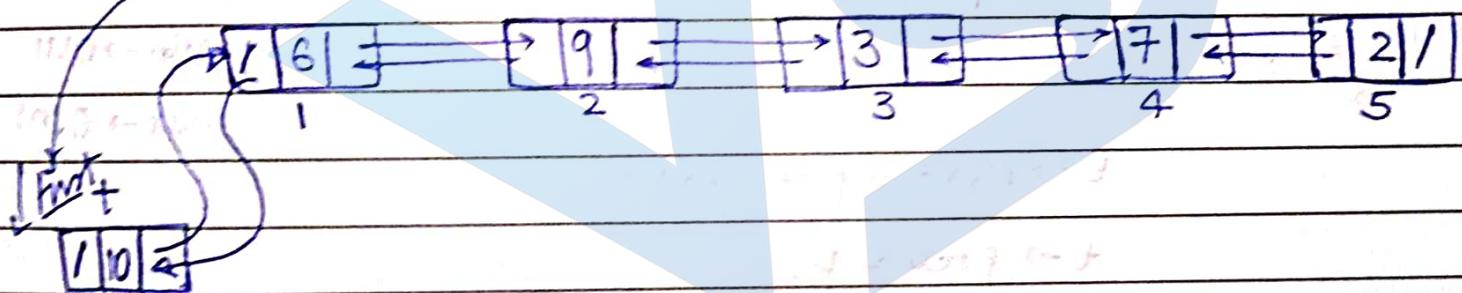


- Before First Node

- At any given Index

- Before first Node :-

ffirst



Node *t = new Node

t → data = x;

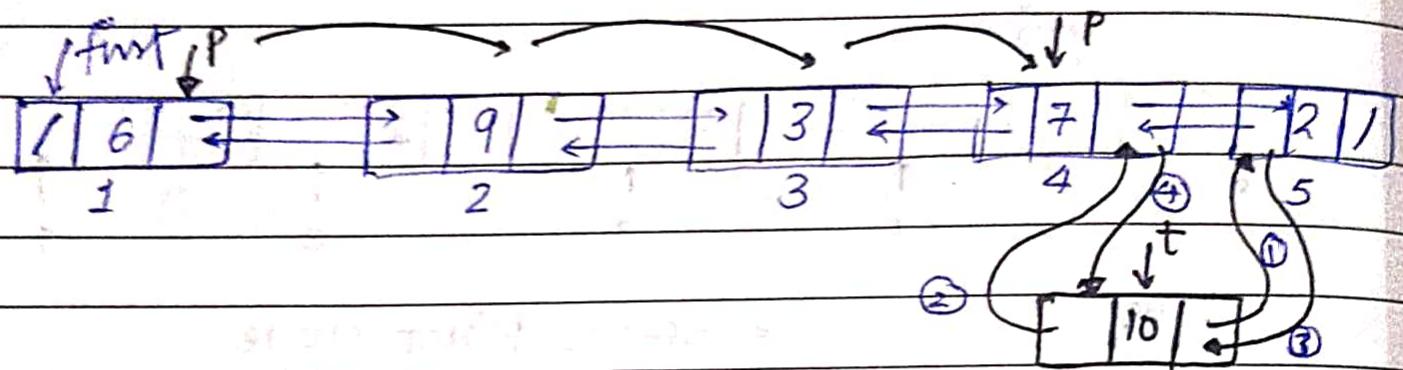
t → prev = NULL;

t → next = first;

first → prev = t;

first = t;

- At any given index :- (pos = 4)



Node *t = new Node;

t->data = x;

for (int i=0; i<pos-1; i++)

{

p = p->next;

t->next = p->next;

t->prev = p;

if (p->next == NULL)

{

p->next->prev = t;

}

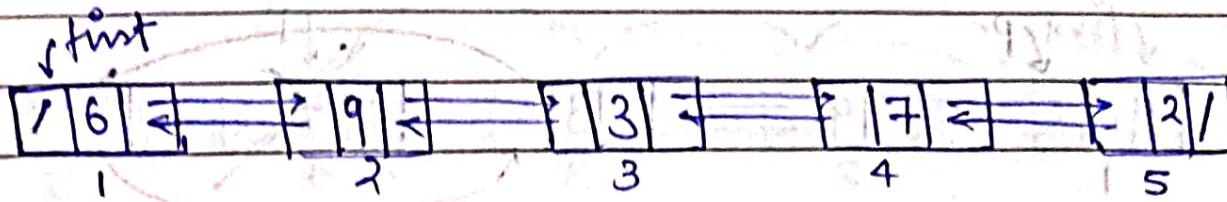
p->next = t;

$O(n)$

$O(1)$

$O(n)$

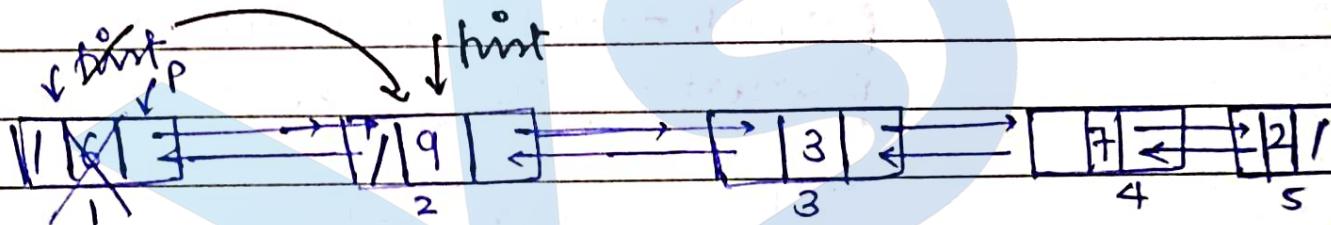
Deleting from Doubly linked list: more about



1. Delete first Node

2. Delete from given index.

- Delete first Node :-



$p = \text{first};$

$\text{first} = \text{first} \rightarrow \text{next};$

$x = p \rightarrow \text{data};$

$\text{delete } p;$

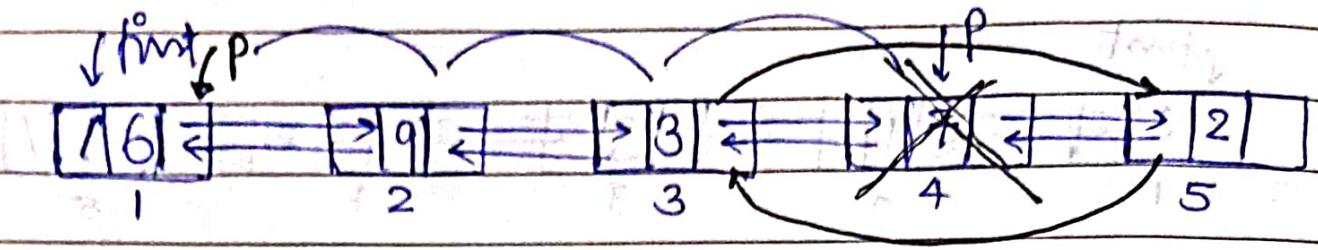
O(1)

$\{$
if ($\text{first} == \text{NULL}$)

$\text{first} \rightarrow \text{prev} = \text{NULL};$

$\}$

- Delete from Given index :- [pos = 4] ~~pos = 4~~



$p = \text{first};$
 $\{$
 $\text{for } (i=0; i < \text{pos}-1; i++)$
 $\}$

$p = p \rightarrow \text{next};$

$p \rightarrow \text{prev} \rightarrow \text{next} = p \rightarrow \text{next};$

$p \rightarrow \text{next} \rightarrow \text{prev} \neq p \rightarrow \text{prev};$

$\text{if}(p \rightarrow \text{next} != \text{NULL})$

$\{$

$p \rightarrow \text{next} \rightarrow \text{prev} = p \rightarrow \text{prev};$

$\}$

~~delete p;~~

$x = p \rightarrow \text{data};$

~~delete p;~~

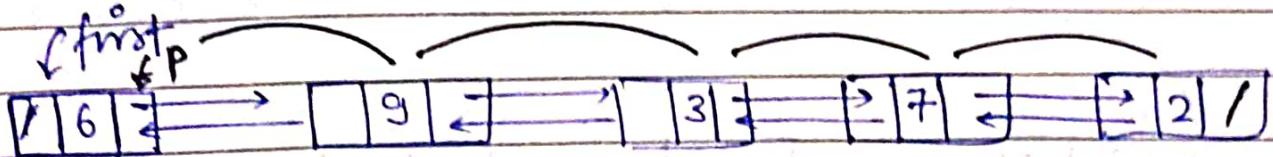
~~free~~

T.C

$\hookrightarrow \text{Min} \rightarrow O(1)$

$\text{Max} \rightarrow \underline{\underline{O(n)}}$

Reverse a Doubly linked list



1. Display.

2. Reverse.

• Display

p = first;

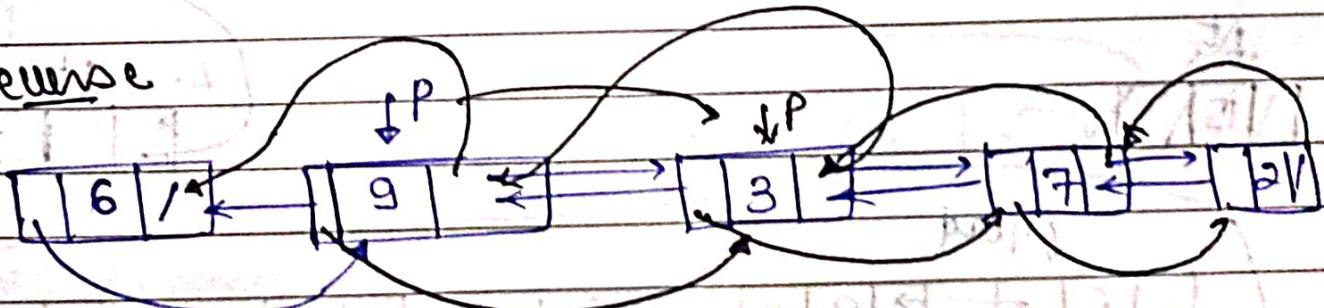
while (p != NULL)

{

cout << p->data;

p = p->next;

• Reverse



p = first;
while (p != NULL)

{

temp = p->next;

p->next = p->prev;

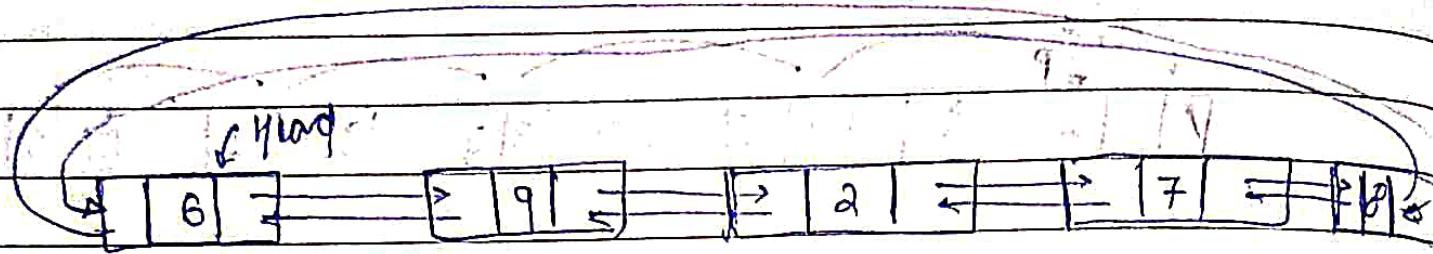
p->prev = temp;

p = p->prev;

}
if (p->next == NULL)
{
first = p;

}

Linear Doubly linked list :-



P = Head;

while do

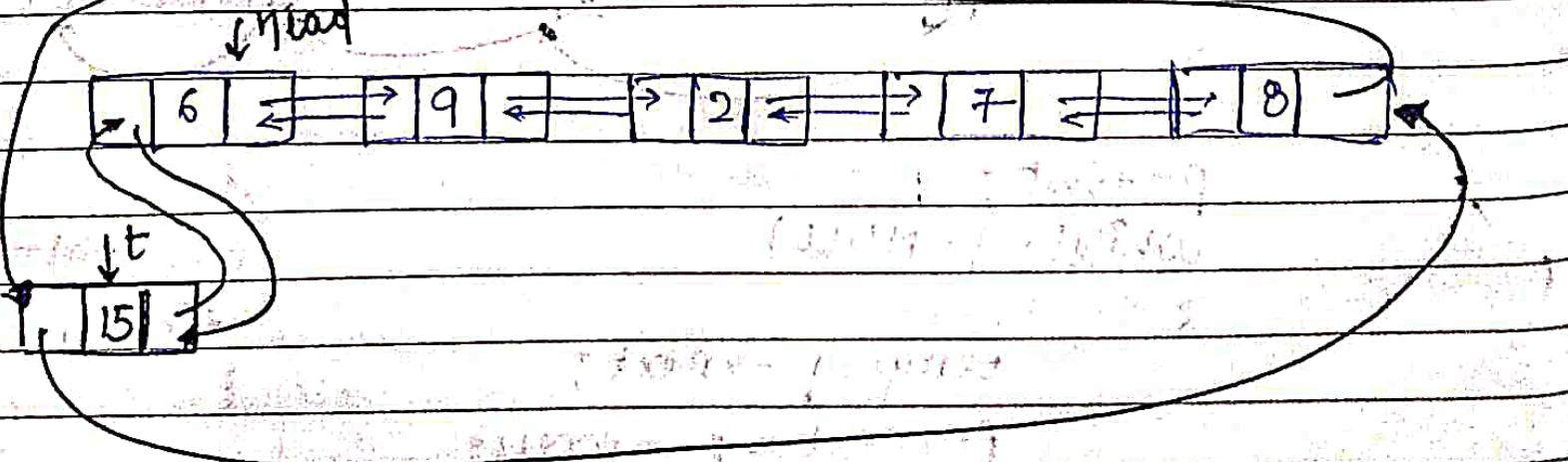
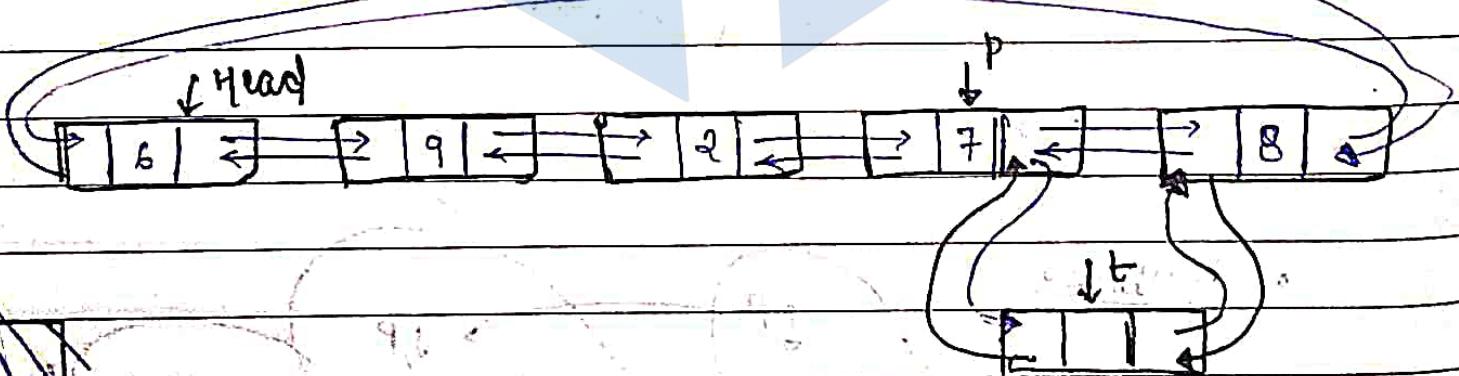
{

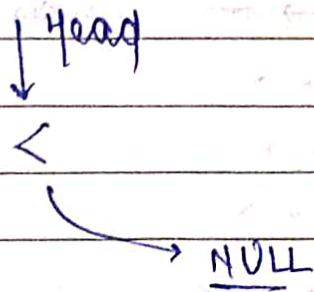
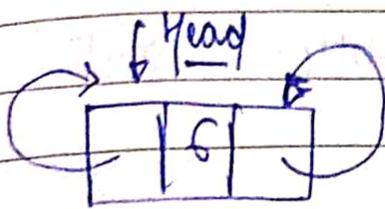
cout << p->data;

p = p->next;

} while (p != Head)

• Insert :-

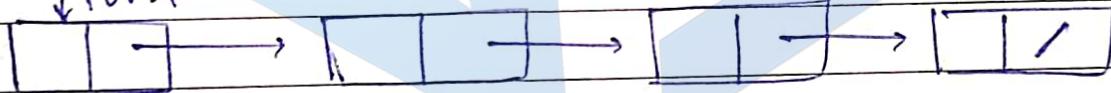




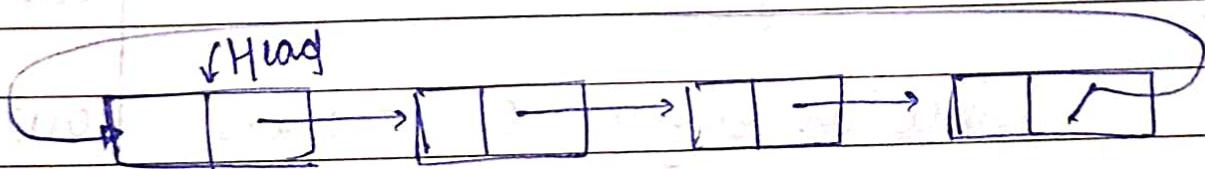
- Insert
- Delete

• Comparison of linked list

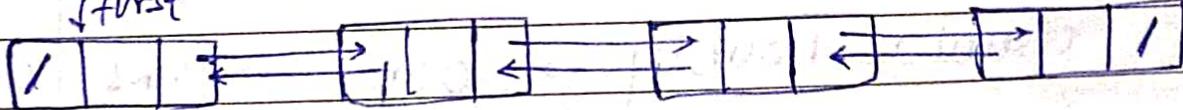
Linear Singly



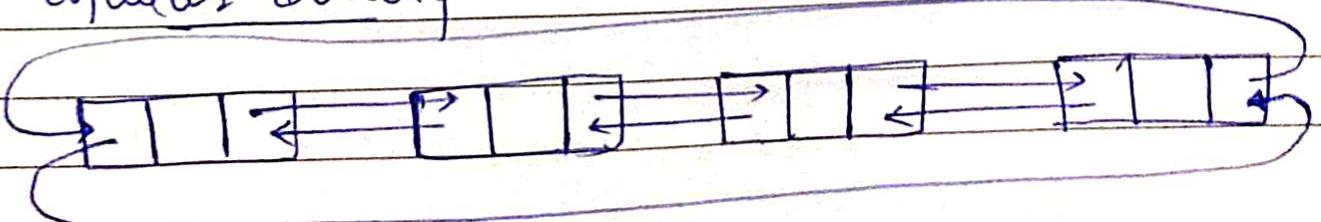
circular singly



linear Doubly



circular doubly



Linear and Circular singly \rightarrow 2 pointers
space \rightarrow Linear and Circular Doubly \rightarrow 4 pointers

insert \rightarrow

delete \rightarrow

Traverse \rightarrow

↳ Singly \rightarrow linear direction

↳ Circular singly \rightarrow linear and circular

↳ Linear Doubly \rightarrow

• forward and backward

→ Circular Doubly \rightarrow

forward, Backward and
Circular

Insert \downarrow

Before head

(Modified)

Linear singly \rightarrow $O(1) \rightarrow 1$ link

Linear
At position

Min Max
 $O(1)$ $O(n)$ 2 links

Circular singly $\rightarrow O(n) \rightarrow 2$ links

$O(1)$ $O(n)$ 4 links

Linear Doubly $\rightarrow O(1) \rightarrow 3$ links

$O(1)$ $O(n)$ 3 links

Circular Doubly $\rightarrow O(n) \rightarrow 4$ links

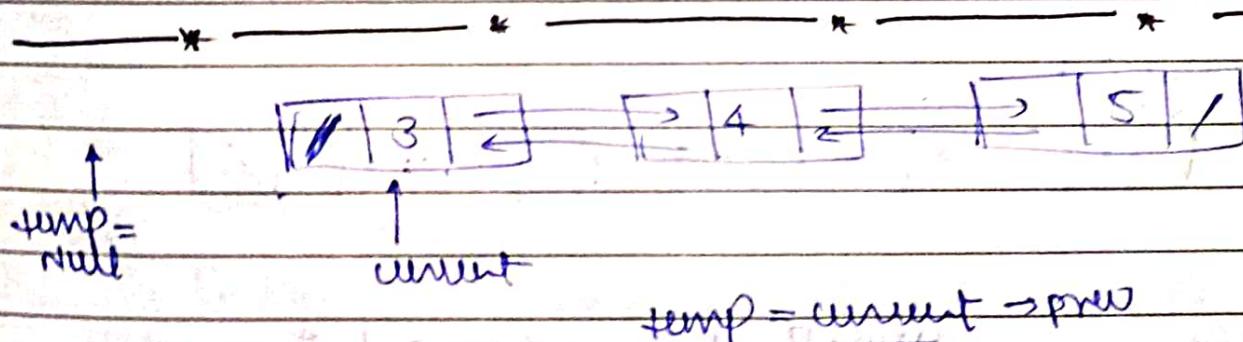
$O(1)$ $O(n)$ 4 links

Delete

	Time complexity	Space complexity	Operations	At given position
linear singly	$O(1)$	Head Modified No links	$O(1)$	$O(n)$ no of links
circular singly	$O(n)$	Link	$O(n)$	$O(n)$ 1 link
doubly singly	$O(1)$	Link	$O(1)$ $O(n)$	2 links
circular doubly	$O(n)$	& Link	$O(1)$ $O(n)$	2 links

Linear singly \rightarrow Queue

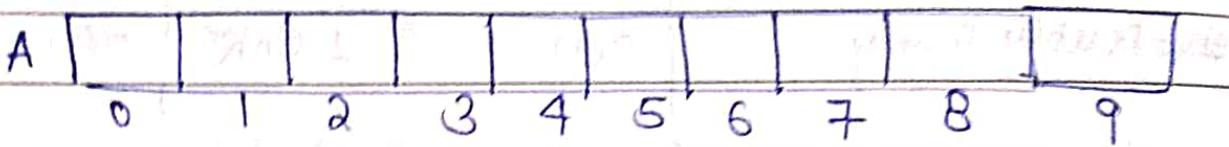
Doubly Circular \rightarrow Best



$temp = \underline{current} \rightarrow prev$

Comparison of Array with linked list:-

Array vs linked list

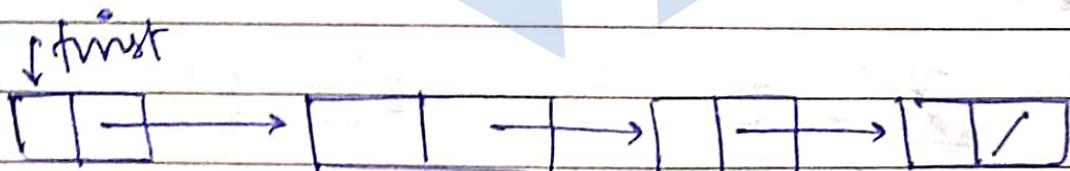


Created in memory

Array → Heap
→ Stack

size → fixed

fixed size, so the chance of getting perfectly utilizes is less
↳ space left or not sufficient to store no. of element



↓ first

linked list → created in memory

↓

heap

size → variable size

When not sure about How many element is used, then used linked list.

Take extra space as compared to array

Array linked list

Access Randomly

- Not Accessed sequentially.

faster Access

slower Access

Insert → O(1) (At last)

At first
O(n)

Insert → O(n) (At last)

At first
O(1)

It is costly
(to move data)

It is not costly

Delete → O(1) (last)

At first
O(n)

Delete → O(n) (last)

O(1) → At first

Linear / binary search
O(n) O(log n)

Linear search
O(n)

Binary search (log n)
~~(not perform)~~

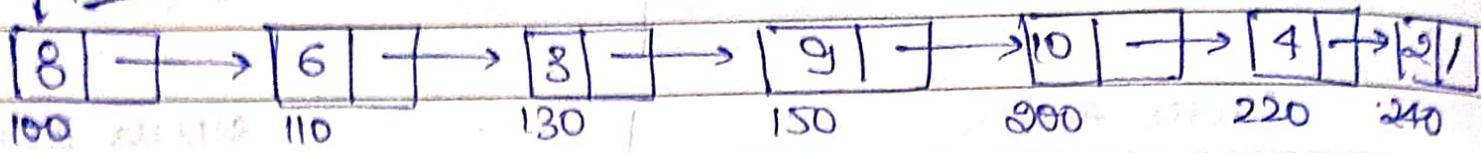
insertion

O(n-1)

insertion and
merge sort

Student Challenge:- Finding Middle Element of a linked list

first

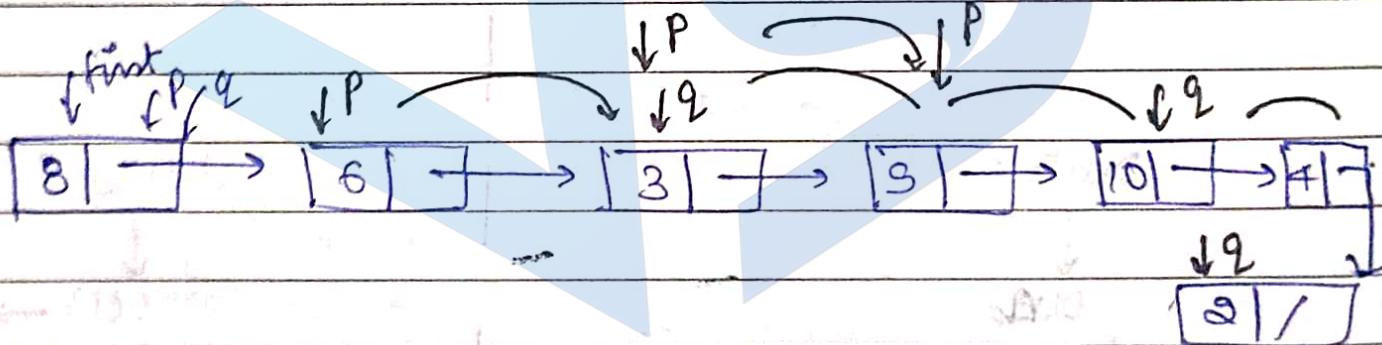


O(n)

1. Find length \rightarrow 7

O(n) \rightarrow 2. Headed to Middle $\rightarrow \frac{1}{2} \rightarrow 3.5 - (\frac{1}{2} \cdot 3.5) = \frac{1}{2}$

O(2n)



$p \rightarrow$ one times

\rightarrow move

$q \rightarrow$ two times

q.next

while ($q_1 = \text{NULL}$)

$q = q \rightarrow \text{next};$

if ($q_1 = \text{NULL}$)

$q = q \rightarrow \text{next};$

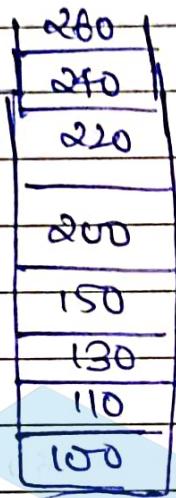
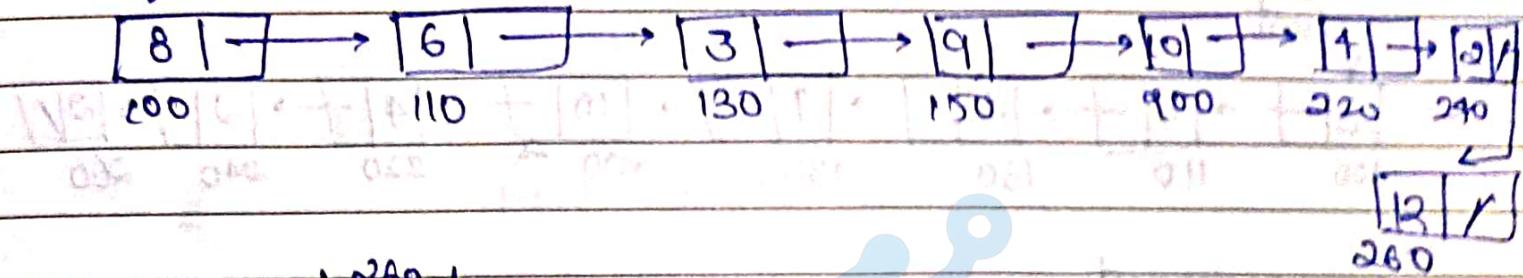
if ($q_1 = \text{NULL}$)

$p = p \rightarrow \text{next};$

cout << p->data;

Using stacks :- for stack overflow & underflow detection

first



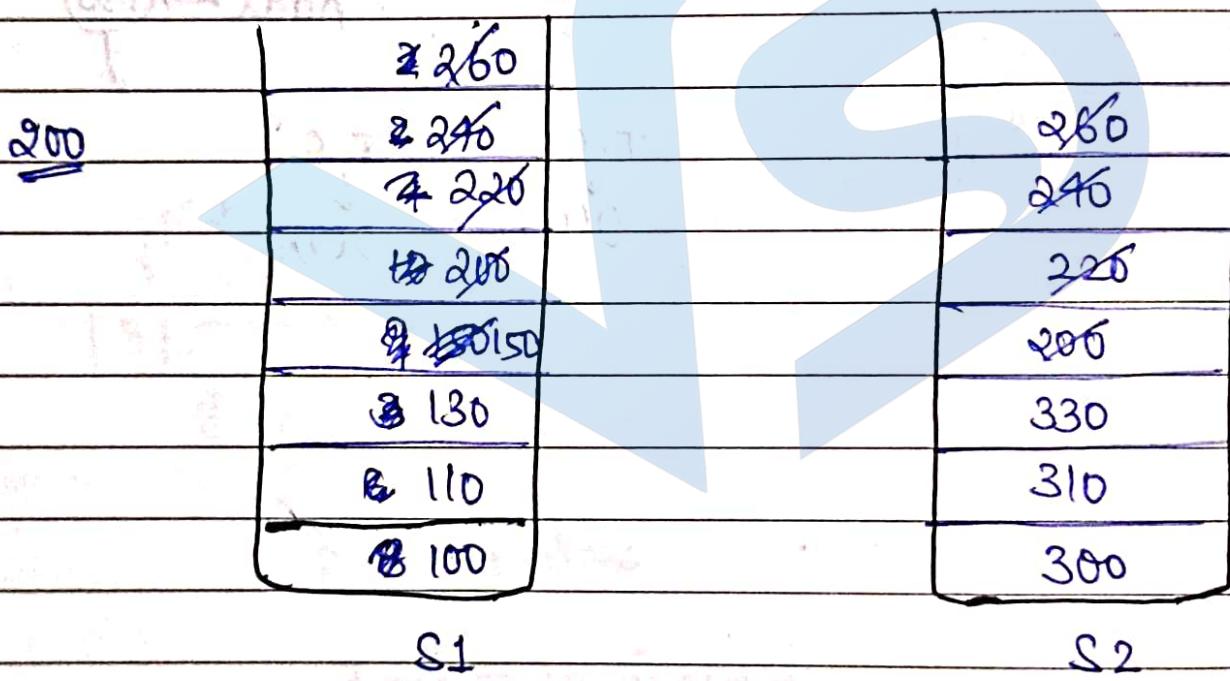
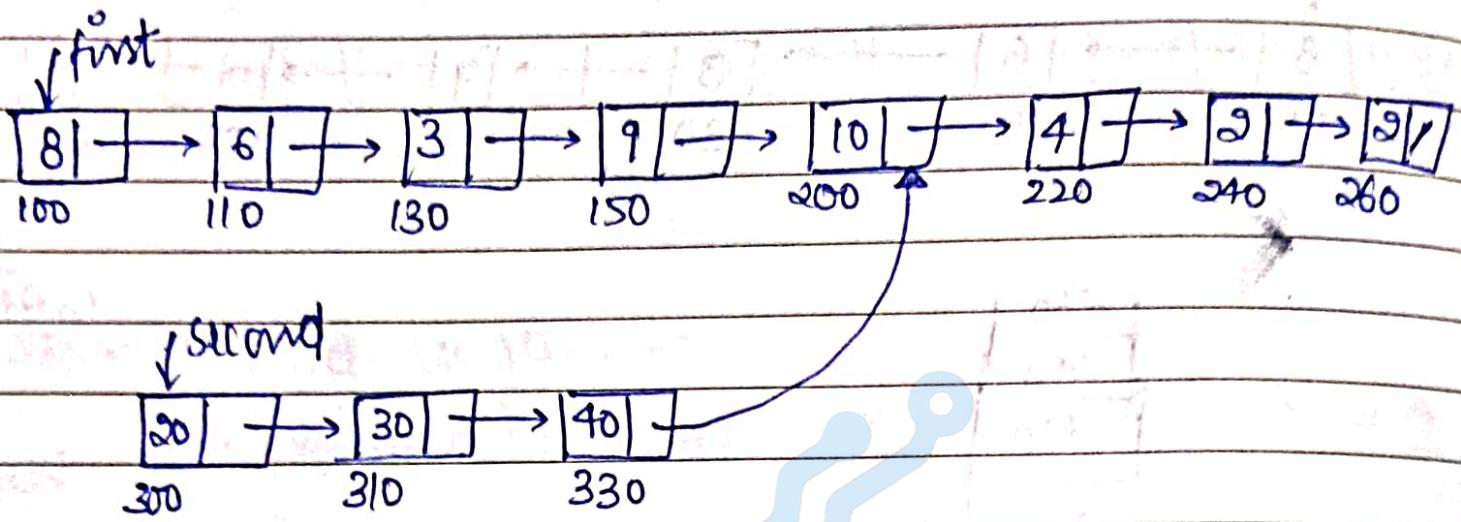
$\rightarrow 8/2 \rightarrow 4$

Addr $\rightarrow 150$

T.C
 $O(n)$

S.C
 $O(n)$

Student challenge: Finding intersecting point of Two Linked list :-



Method 1: Using stack

$T.C \rightarrow O(n+n) + n$
 $O(3n)$

$S.C \rightarrow O(m)$

$p = \text{first};$

$\{\text{else block}\}$

$\text{while}(p \neq \text{NULL})$

{

$\text{push}(\&\text{stk1}, p);$

$p = \text{second};$

$\text{while}(p \neq \text{NULL})$

{

$\text{push}(\&\text{stk2}, p)$

5

$\text{while}(\text{stackTop(stk1)} == \text{stackTop(stk2)})$

{

~~pop~~ $p = \text{pop}(\&\text{stk1});$

$\text{pop}(\&\text{stk2});$

}

$\text{cout} \ll p \rightarrow \text{data};$

intersecting point