

ENGG5781 Matrix Analysis and Computations

Lecture 6: Least Squares Revisited

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Lecture 6: Least Squares Revisited

- Part I: regularization
- Part II: sparsity
 - ℓ_0 minimization
 - greedy pursuit, ℓ_1 minimization, and variations
 - majorization-minimization for ℓ_2 - ℓ_1 minimization
 - dictionary learning
- Part III: LS with errors in \mathbf{A}
 - total LS
 - robust LS, and its equivalence to regularization

Part I: Regularization

Sensitivity to Noise

- **Question:** how sensitive is the LS solution when there is noise?
- **Model:**

$$\mathbf{y} = \mathbf{A}\bar{\mathbf{x}} + \boldsymbol{\nu},$$

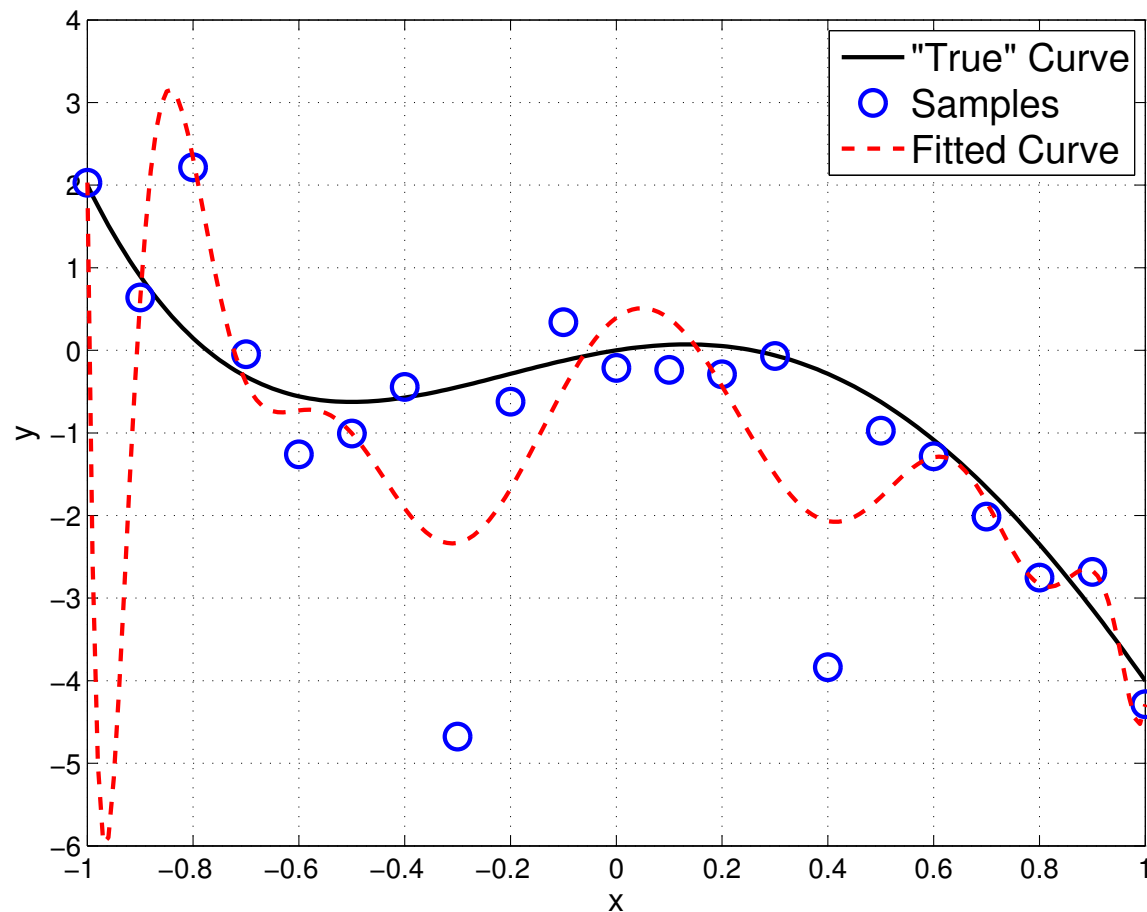
where $\bar{\mathbf{x}}$ is the true result; $\mathbf{A} \in \mathbb{R}^{m \times n}$ has full column rank; $\boldsymbol{\nu}$ is noise, modeled as a random vector with mean zero and covariance $\gamma^2 \mathbf{I}$.

- **Mean square error (MSE) analysis:** from $\mathbf{x}_{\text{LS}} = \mathbf{A}^\dagger \mathbf{y} = \bar{\mathbf{x}} + \mathbf{A}^\dagger \boldsymbol{\nu}$ we get

$$\begin{aligned} \mathbb{E}[\|\mathbf{x}_{\text{LS}} - \bar{\mathbf{x}}\|_2^2] &= \mathbb{E}[\|\mathbf{A}^\dagger \boldsymbol{\nu}\|_2^2] = \mathbb{E}[\text{tr}(\mathbf{A}^\dagger \boldsymbol{\nu} \boldsymbol{\nu}^T (\mathbf{A}^\dagger)^T)] = \text{tr}(\mathbf{A}^\dagger \mathbb{E}[\boldsymbol{\nu} \boldsymbol{\nu}^T] (\mathbf{A}^\dagger)^T) \\ &= \gamma^2 \text{tr}(\mathbf{A}^\dagger (\mathbf{A}^\dagger)^T) = \gamma^2 \text{tr}((\mathbf{A}^T \mathbf{A})^{-1}) \\ &= \gamma^2 \sum_{i=1}^n \frac{1}{\sigma_i^2(\mathbf{A})} \end{aligned}$$

- **Observation:** the MSE becomes very large if some $\sigma_i(\mathbf{A})$'s are close to zero.

Toy Demonstration: Curve Fitting



The same curve fitting example in Lecture 2. The “true” curve is the true $f(x)$ with model order $n = 4$. In practice, the model order may not be known and we may have to guess. The fitted curve above is done by LS with a guessed model order $n = 16$.

ℓ_2 -Regularized LS

- **Intuition:** replace $\mathbf{x}_{LS} = (\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T \mathbf{y}$ by

$$\mathbf{x}_{RLS} = (\mathbf{A}^T \mathbf{A} + \lambda \mathbf{I})^{-1} \mathbf{A}^T \mathbf{y},$$

for some $\lambda > 0$, where the term $\lambda \mathbf{I}$ is added to improve the system conditioning, thereby attempting to reduce noise sensitivity

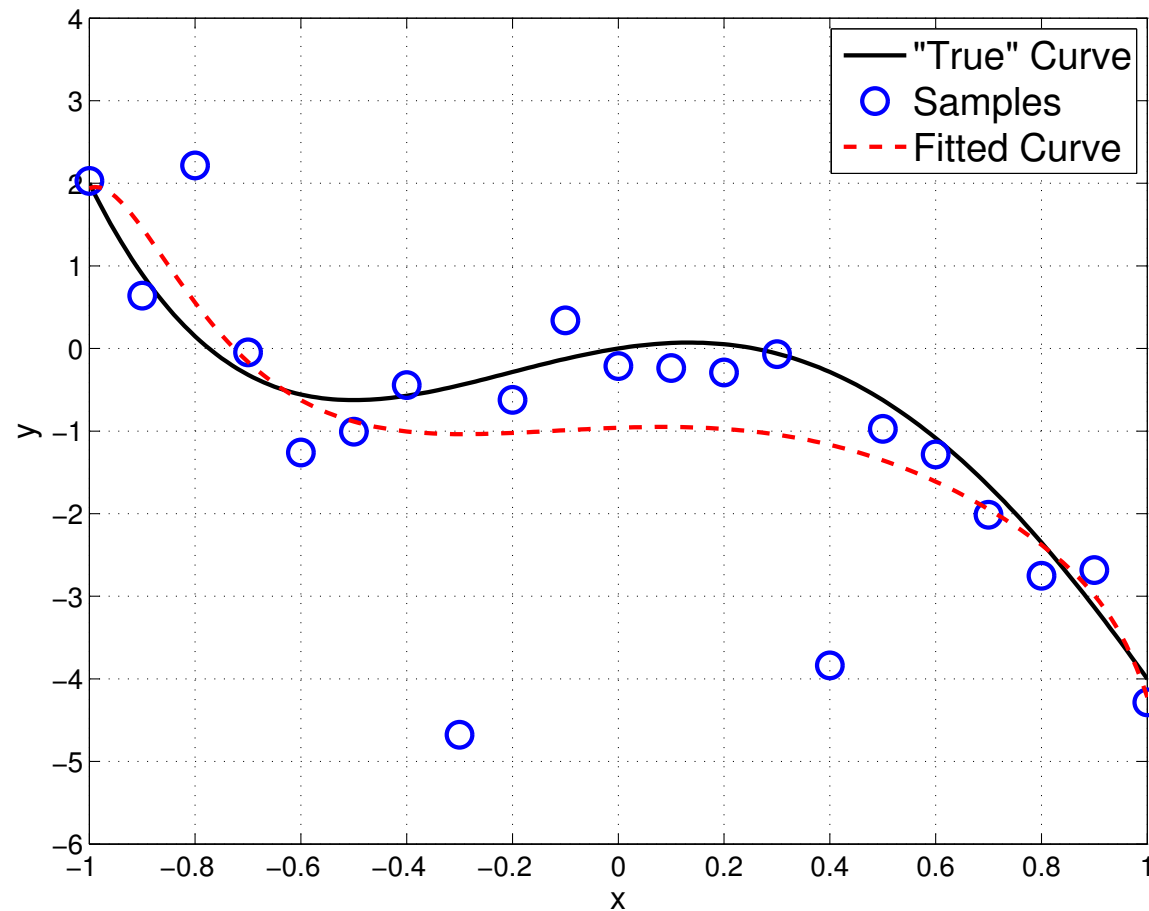
- how may we make sense out of such a modification?
- ℓ_2 -regularized LS: find an \mathbf{x} that solves

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{A}\mathbf{x} - \mathbf{y}\|_2^2 + \lambda \|\mathbf{x}\|_2^2$$

for some pre-determined $\lambda > 0$.

- the solution is uniquely given by $\mathbf{x}_{RLS} = (\mathbf{A}^T \mathbf{A} + \lambda \mathbf{I})^{-1} \mathbf{A}^T \mathbf{y}$
- the formulation says that we try to minimize both $\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2$ and $\|\mathbf{x}\|_2^2$, and λ controls which one should be more emphasized in the minimization

Toy Demonstration: Curve Fitting



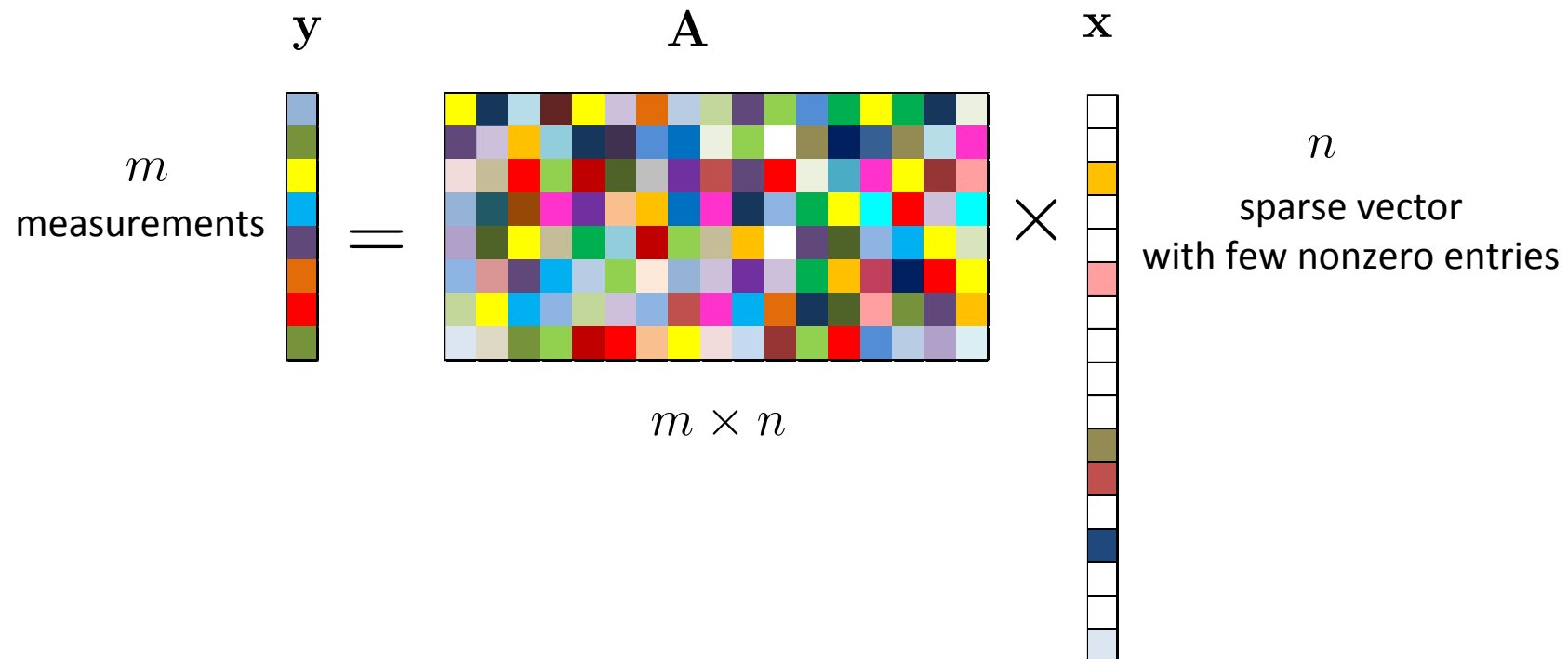
The fitted curve is done by ℓ_2 -regularized LS with a guessed model order $n = 18$ and with $\lambda = 0.1$.

Part II: Sparsity

The Sparse Recovery Problem

Problem: given $\mathbf{y} \in \mathbb{R}^m$, $\mathbf{A} \in \mathbb{R}^{m \times n}$, $m < n$, find a **sparsest** $\mathbf{x} \in \mathbb{R}^n$ such that

$$\mathbf{y} = \mathbf{A}\mathbf{x}.$$



- by sparsest, we mean that \mathbf{x} should have as many zero elements as possible.

A Sparsity Optimization Formulation

- let

$$\|\mathbf{x}\|_0 = \sum_{i=1}^n \mathbb{1}\{x_i \neq 0\}$$

denote the cardinality function

– commonly called the “ ℓ_0 -norm”, though it is not a norm.

- Minimum ℓ_0 -norm formulation:

$$\begin{aligned} \min_{\mathbf{x} \in \mathbb{R}^n} \quad & \|\mathbf{x}\|_0 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}. \end{aligned}$$

- **Question:** suppose that $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$, where $\bar{\mathbf{x}}$ is the vector we seek to recover. Can the min. ℓ_0 -norm problem recover $\bar{\mathbf{x}}$ in an exact and unique fashion?
 - an answer lies in the notion of **spark**, which may be seen as a strong definition of rank

Spark

Spark: the spark of \mathbf{A} , denoted by $\text{spark}(\mathbf{A})$, is the **smallest** number of **linearly dependent** columns of \mathbf{A} .

- let $\text{spark}(\mathbf{A}) = k$. Then, k is the smallest number such that there exists a linearly dependent $\{\mathbf{a}_{i_1}, \dots, \mathbf{a}_{i_k}\}$ for some $\{i_1, \dots, i_k\} \subseteq \{1, \dots, n\}$ ¹.
- let $\text{spark}(\mathbf{A}) = r + 1$. Then, $\{\mathbf{a}_{i_1}, \dots, \mathbf{a}_{i_r}\}$ is linearly independent for any $\{i_1, \dots, i_r\} \subseteq \{1, \dots, n\}$
 - any collection of r columns of \mathbf{A} is linearly independent, simply stated
- **Comparison with rank:** Let $\text{rank}(\mathbf{A}) = r$ (not the same r above). Then, there **exists** a linearly independent $\{\mathbf{a}_{i_1}, \dots, \mathbf{a}_{i_r}\}$ for some $\{i_1, \dots, i_r\} \subseteq \{1, \dots, n\}$.
- **Kruskal rank:** this is an alternative definition of rank. The Kruskal rank of \mathbf{A} , denoted by $\text{krank}(\mathbf{A})$, has its definition equivalent to $\text{krank}(\mathbf{A}) = \text{spark}(\mathbf{A}) - 1$.

¹We leave it implicit that $i_k \neq i_j$ for any $k \neq j$.

Spark

- if any collection of m vectors in $\{\mathbf{a}_1, \dots, \mathbf{a}_n\} \subseteq \mathbb{R}^m$, with $n > m$, is linearly independent, then

$$\text{spark}(\mathbf{A}) = m + 1, \quad \text{rank}(\mathbf{A}) = m.$$

- an example is Vandemonde matrices with distinct roots
- some specifically designed bases also have this property
- but there also exist instances in which rank and spark are very different
 - let $\{\mathbf{v}_1, \dots, \mathbf{v}_r\} \in \mathbb{R}^m$ be linearly independent, and let $\mathbf{A} = [\mathbf{v}_1, \dots, \mathbf{v}_r, \mathbf{v}_1]$.
 - we have $\text{rank}(\mathbf{A}) = r$, but $\text{spark}(\mathbf{A}) = 2$
- to conclude, spark may be seen as a stronger definition of rank, and

$$\text{spark}(\mathbf{A}) - 1 \leq \text{rank}(\mathbf{A})$$

Perfect Recovery Guarantee of the Min. ℓ_0 -Norm Problem

Theorem 6.1. Suppose that $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$. Then, $\bar{\mathbf{x}}$ is the unique solution to the minimum ℓ_0 -norm problem if

$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}\text{spark}(\mathbf{A}).$$

- **Implication:** if $\bar{\mathbf{x}}$ is sufficiently sparse, then the minimum ℓ_0 -norm problem perfectly recovers $\bar{\mathbf{x}}$
- **Proof sketch:**
 1. let \mathbf{x}^* be a solution to the min. ℓ_0 -norm problem. Let $\mathbf{e} = \bar{\mathbf{x}} - \mathbf{x}^*$.
 2. $\mathbf{0} = \mathbf{A}\bar{\mathbf{x}} - \mathbf{A}\mathbf{x}^* = \mathbf{A}\mathbf{e}$; $\|\mathbf{e}\|_0 \leq \|\bar{\mathbf{x}}\|_0 + \|\mathbf{x}^*\|_0 \leq 2\|\bar{\mathbf{x}}\|_0$.
 3. suppose $\mathbf{e} \neq \mathbf{0}$. Then, $\mathbf{A}\mathbf{e} = \mathbf{0}$, $\|\mathbf{e}\|_0 \leq 2\|\bar{\mathbf{x}}\|_0 \implies \text{spark}(\mathbf{A}) \leq 2\|\bar{\mathbf{x}}\|_0$

Perfect Recovery Guarantee of the Min. ℓ_0 -Norm Problem

- **coherence:** the coherence of \mathbf{A} is defined as

$$\mu(\mathbf{A}) = \max_{j \neq k} \frac{|\mathbf{a}_j^T \mathbf{a}_k|}{\|\mathbf{a}_j\|_2 \|\mathbf{a}_k\|_2}.$$

– measures how similar the columns of \mathbf{A} are in the worst-case sense.

- a weak version of Theorem 6.1:

Corollary 6.1. Suppose that $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$. Then, $\bar{\mathbf{x}}$ is the unique solution to the minimum ℓ_0 -norm problem if

$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}(1 + \mu(\mathbf{A})^{-1}).$$

- **Implication:** perfect recovery may depend on how incoherent \mathbf{A} is.
- proof idea: show that $\text{spark}(\mathbf{A}) \geq 1 + \mu(\mathbf{A})^{-1}$

On Solving the Minimum ℓ_0 -Norm Problem

Question: How should we solve the minimum ℓ_0 -norm problem

$$\begin{aligned} \min_{\mathbf{x}} \quad & \|\mathbf{x}\|_0 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}, \end{aligned}$$

or can it be efficiently solved?

- ℓ_0 -norm minimization does not lead to a simple solution as in 2-norm min.
- the minimum ℓ_0 -norm problem is **NP-hard** in general
 - what does that mean?
 - * given any \mathbf{y} , \mathbf{A} , the problem is unlikely to be exactly solvable in polynomial time (i.e., in a complexity of $\mathcal{O}(n^p)$ for any $p > 0$)

Brute Force Search for the Minimum ℓ_0 -Norm Problem

- notation: $\mathbf{A}_{\mathcal{I}}$ denotes a submatrix of \mathbf{A} obtained by keeping the columns indicated by \mathcal{I}
- we may solve the ℓ_0 -norm minimization problem via brute force search:

input: \mathbf{A}, \mathbf{y}
for all $\mathcal{I} \subseteq \{1, 2, \dots, n\}$ do
 if $\mathbf{y} = \mathbf{A}_{\mathcal{I}}\tilde{\mathbf{x}}$ has a solution for some $\tilde{\mathbf{x}} \in \mathbb{R}^{|\mathcal{I}|}$
 record $(\tilde{\mathbf{x}}, \mathcal{I})$ as one of candidate solutions
end
output: a candidate solution $(\tilde{\mathbf{x}}, \mathcal{I})$ whose $|\mathcal{I}|$ is the smallest

- example: for $n = 3$, we test $\mathcal{I} = \{1\}, \mathcal{I} = \{2\}, \mathcal{I} = \{3\}, \mathcal{I} = \{1, 2\}, \mathcal{I} = \{2, 3\}, \mathcal{I} = \{1, 3\}, \mathcal{I} = \{1, 2, 3\}$
- manageable for very small n , too expensive even for moderate n
- how about a greedy search that searches less?

Greedy Pursuit

- consider a greedy search called the **orthogonal matching pursuit (OMP)**

Algorithm: OMP

input: \mathbf{A}, \mathbf{y}

set $\mathcal{I} = \emptyset, \hat{\mathbf{x}} = \mathbf{0}$

repeat

$$\mathbf{r} = \mathbf{y} - \mathbf{A}\hat{\mathbf{x}}$$

$$k = \arg \max_{j \in \{1, \dots, n\}} |\mathbf{a}_j^T \mathbf{r}| / \|\mathbf{a}_j\|_2$$

$$\mathcal{I} := \mathcal{I} \cup \{k\}$$

$$\hat{\mathbf{x}} := \arg \min_{\mathbf{x} \in \mathbb{R}^n, x_i=0 \ \forall i \notin \mathcal{I}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2$$

until a stopping rule is satisfied, e.g., $\|\mathbf{y} - \mathbf{A}\hat{\mathbf{x}}\|_2$ is sufficiently small

output: $\hat{\mathbf{x}}$

- note: there are many other greedy search strategies

Perfect Recovery Guarantee of Greedy Pursuit

- again, a key question is conditions under which OMP admits perfect recovery
- there are many such theoretical conditions, not only for OMP but also for other greedy algorithms
- one such result is as follows:

Theorem 6.2. Suppose that $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$. Then, OMP recovers $\bar{\mathbf{x}}$ if

$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}(1 + \mu(\mathbf{A})^{-1}).$$

- proof idea: show that OMP is guaranteed to pick a correct column at every stage.

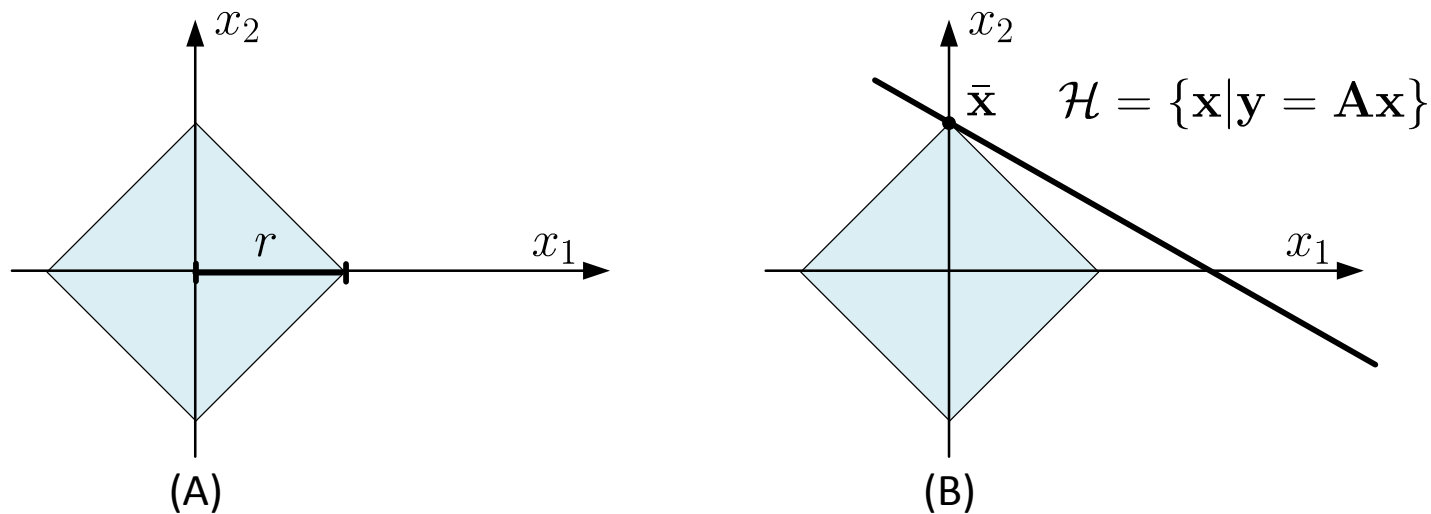
Convex Relaxation

Another approximation approach is to replace $\|\mathbf{x}\|_0$ by a convex function:

$$\begin{aligned} \min_{\mathbf{x}} \quad & \|\mathbf{x}\|_1 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}. \end{aligned}$$

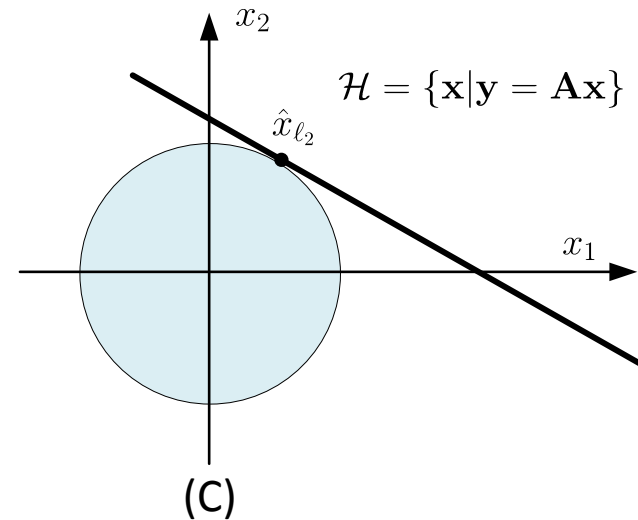
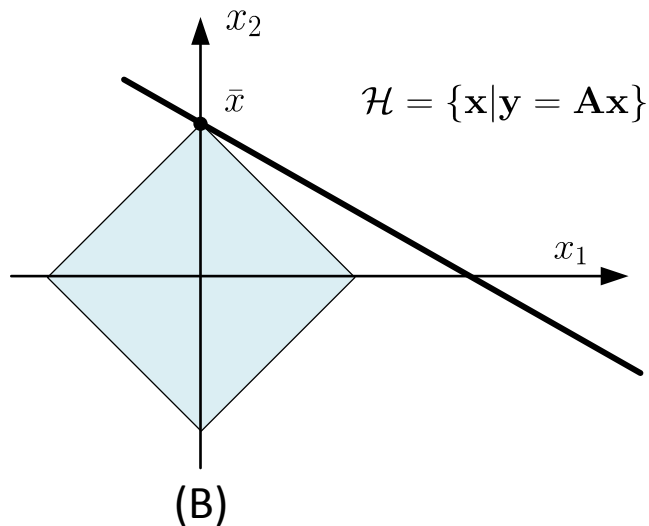
- also known as basis pursuit in the literature
- convex, a linear program
- no closed-form solution (while the minimum 2-norm problem has)
- but the success of this minimum 1-norm problem, both in theory and practice, has motivated a large body of work on computationally efficient algorithms for it

Illustration of 1-Norm Geometry



- Fig. A shows the 1-norm ball of radius r in \mathbb{R}^2 . Note that the 1-norm ball is “pointy” along the axes.
- Fig. B shows the 1-norm recovery solution. The point $\bar{\mathbf{x}}$ is a “sparse” vector; the line \mathcal{H} is the set of all \mathbf{x} that satisfy $\mathbf{y} = \mathbf{A}\mathbf{x}$.

Illustration of 1-Norm Geometry



- The 1-norm recovery problem is to pick out a point in \mathcal{H} that has the minimum 1-norm. We can see that $\bar{\mathbf{x}}$ is such a point.
- Fig. C shows the geometry when 2-norm is used. We can see that the solution $\hat{\mathbf{x}}$ may not be sparse.

Perfect Recovery Guarantee of the Min. 1-Norm Problem

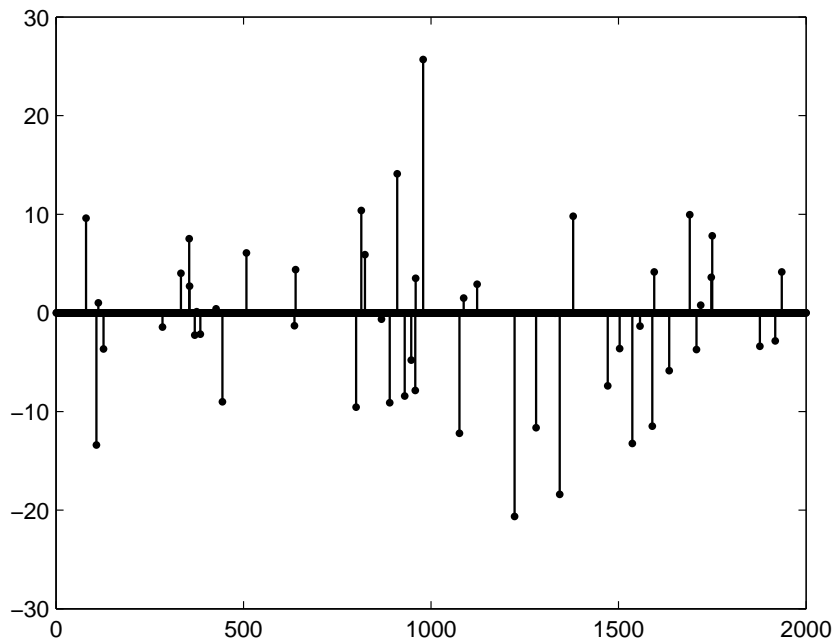
- again, researchers studied conditions under which the minimum 1-norm problem admits perfect recovery
- this has been an exciting topic, with many provable conditions such as the restricted isometry property (RIP), the nullspace property (NSP), ...
 - see the literature for details, and here is one: [\[Yin'13\]](#)
- a simple one is as follows:

Theorem 6.3. Suppose that $\mathbf{y} = \mathbf{A}\bar{\mathbf{x}}$. Then, $\bar{\mathbf{x}}$ is the unique solution to the minimum 1-norm problem if

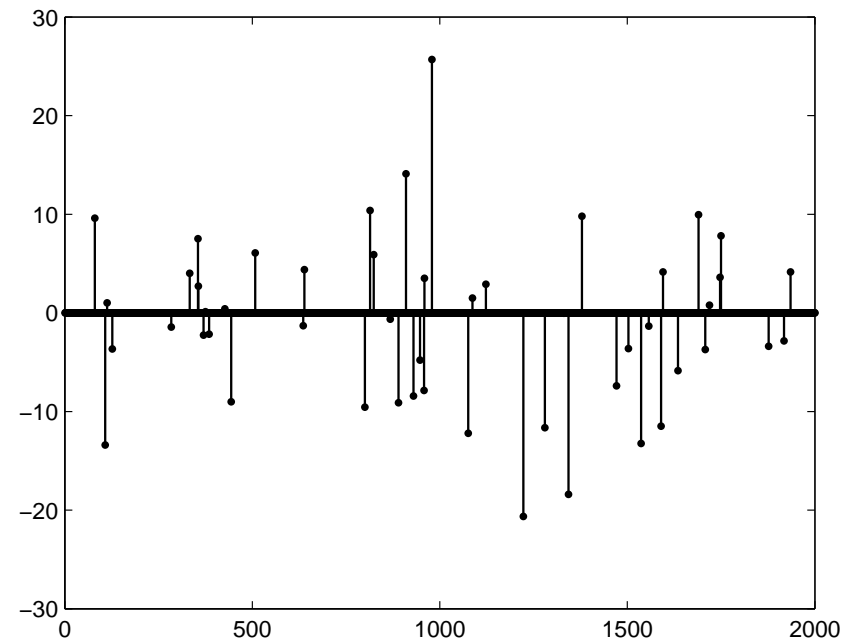
$$\|\bar{\mathbf{x}}\|_0 < \frac{1}{2}(1 + \mu(\mathbf{A})^{-1}).$$

Toy Demonstration: Sparse Signal Reconstruction

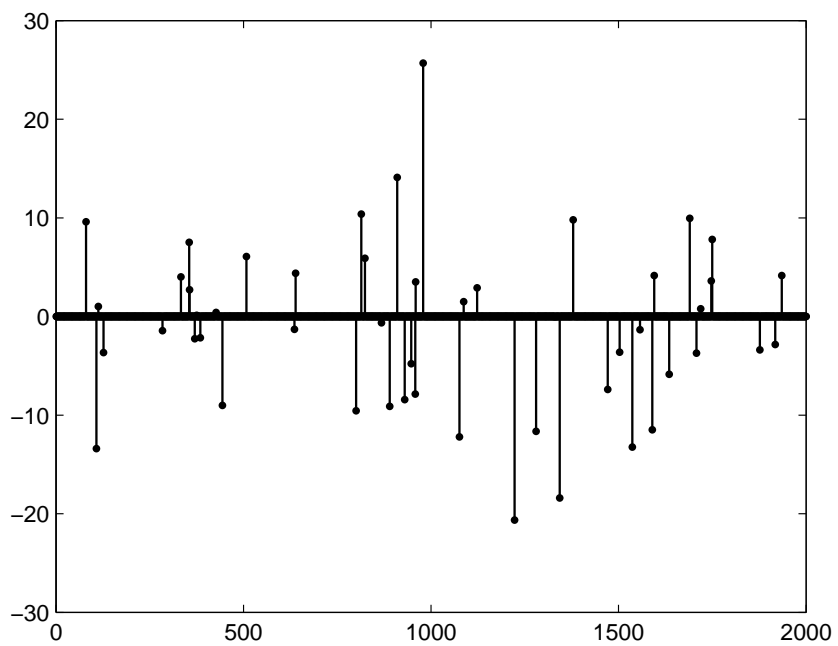
- Sparse vector $\mathbf{x} \in \mathbb{R}^n$ with $n = 2000$ and $\|\mathbf{x}\|_0 = 50$.
- $m = 400$ noise-free observations of $\mathbf{y} = \mathbf{A}\mathbf{x}$, a_{ij} is randomly generated.



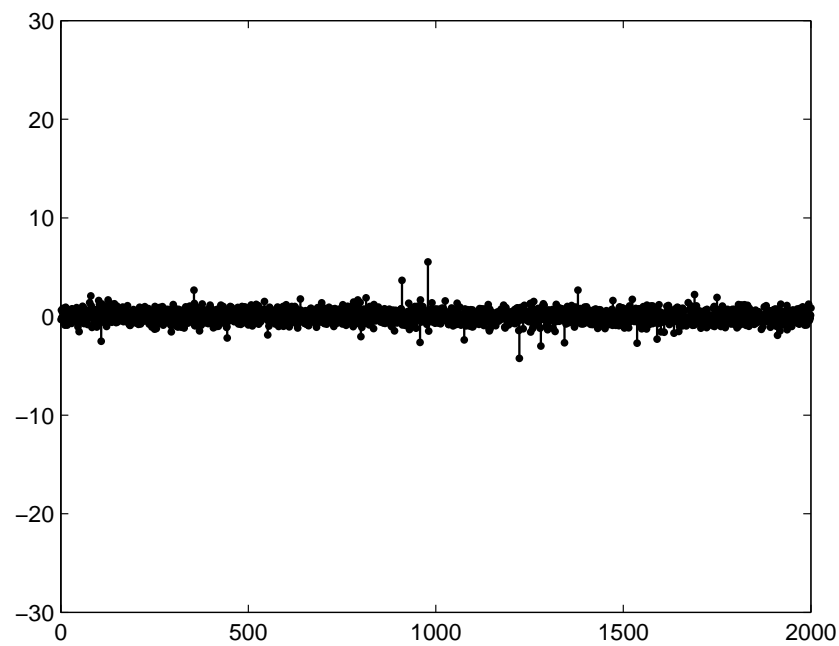
(a) Sparse source signal



(b) Recovery by 1-norm minimization



(c) Sparse source signal



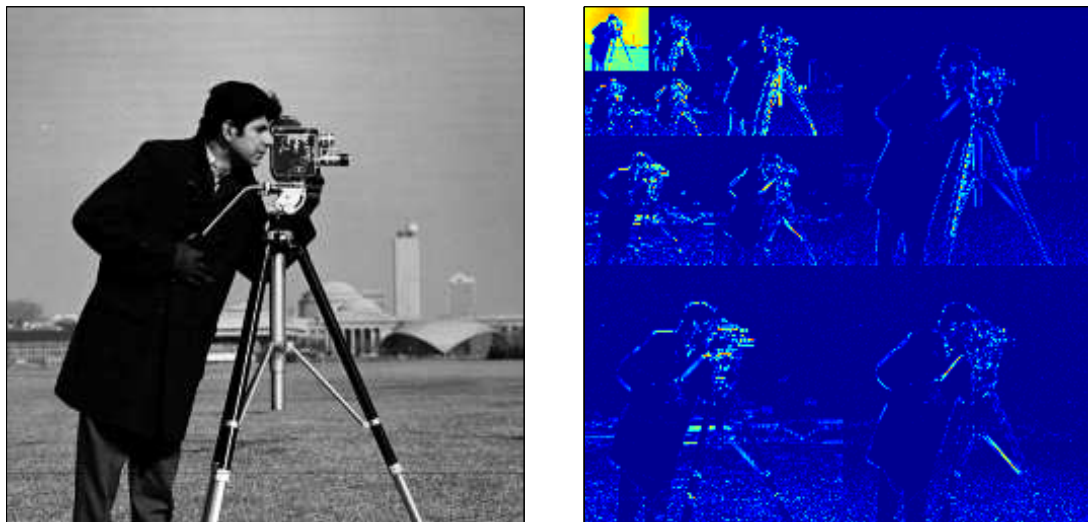
(d) Recovery by 2-norm minimization

Application: Compressive sensing (CS)

- Consider a signal $\tilde{\mathbf{x}} \in \mathbb{R}^n$ that has a sparse representation $\mathbf{x} \in \mathbb{R}^n$ in the domain of $\Psi \in \mathbb{R}^{n \times n}$ (e.g. DCT or wavelet), i.e.,

$$\tilde{\mathbf{x}} = \Psi \mathbf{x},$$

where \mathbf{x} is sparse.



Left: the original image $\tilde{\mathbf{x}}$. Right: the corresponding coefficient \mathbf{x} in the wavelet domain, which is sparse. Source: [\[Romberg-Wakin'07\]](#)

Application: CS

- To acquire \mathbf{x} , we use a sensing matrix $\Phi \in \mathbb{R}^{m \times n}$ to observe \mathbf{x}

$$\mathbf{y} = \Phi \tilde{\mathbf{x}} = \Phi \Psi \mathbf{x}.$$

Here, we have $m \ll n$, i.e., much fewer observations than the no. of unknowns

- Such a \mathbf{y} will be good for compression, transmission and storage.
- $\tilde{\mathbf{x}}$ is recovered by recovering \mathbf{x} :

$$\begin{aligned} \min \quad & \|\mathbf{x}\|_0 \\ \text{s.t.} \quad & \mathbf{y} = \mathbf{A}\mathbf{x}, \end{aligned}$$

where $\mathbf{A} = \Phi \Psi$

- how to choose Φ ? CS research suggests that i.i.d. random Φ will work well!

Application: CS

$$\begin{aligned}
 y_1 &= \left\langle \begin{array}{c} \text{img} \\ \text{noise} \end{array} \right\rangle, \\
 y_2 &= \left\langle \begin{array}{c} \text{img} \\ \text{noise} \end{array} \right\rangle, \\
 y_3 &= \left\langle \begin{array}{c} \text{img} \\ \text{noise} \end{array} \right\rangle, \\
 &\vdots \\
 y_M &= \left\langle \begin{array}{c} \text{img} \\ \text{noise} \end{array} \right\rangle,
 \end{aligned}$$

(a) measurements via i.i.d. random Φ



original (25k wavelets)

(b) original image



perfect recovery

(c) ℓ_1 recovery

Source: [\[Romberg-Wakin'07\]](#)

Variations

- when \mathbf{y} is contaminated by noise, or when $\mathbf{y} = \mathbf{A}\mathbf{x}$ does not exactly hold, some variants of the previous min. 1-norm formulation may be considered:

- **basis pursuit denoising**: given $\epsilon > 0$, solve

$$\min_{\mathbf{x}} \|\mathbf{x}\|_1 \quad \text{s.t.} \quad \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 \leq \epsilon$$

- **ℓ_1 -regularized LS**: given $\lambda > 0$, solve

$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \lambda \|\mathbf{x}\|_1$$

- **Lasso**: given $\tau > 0$, solve

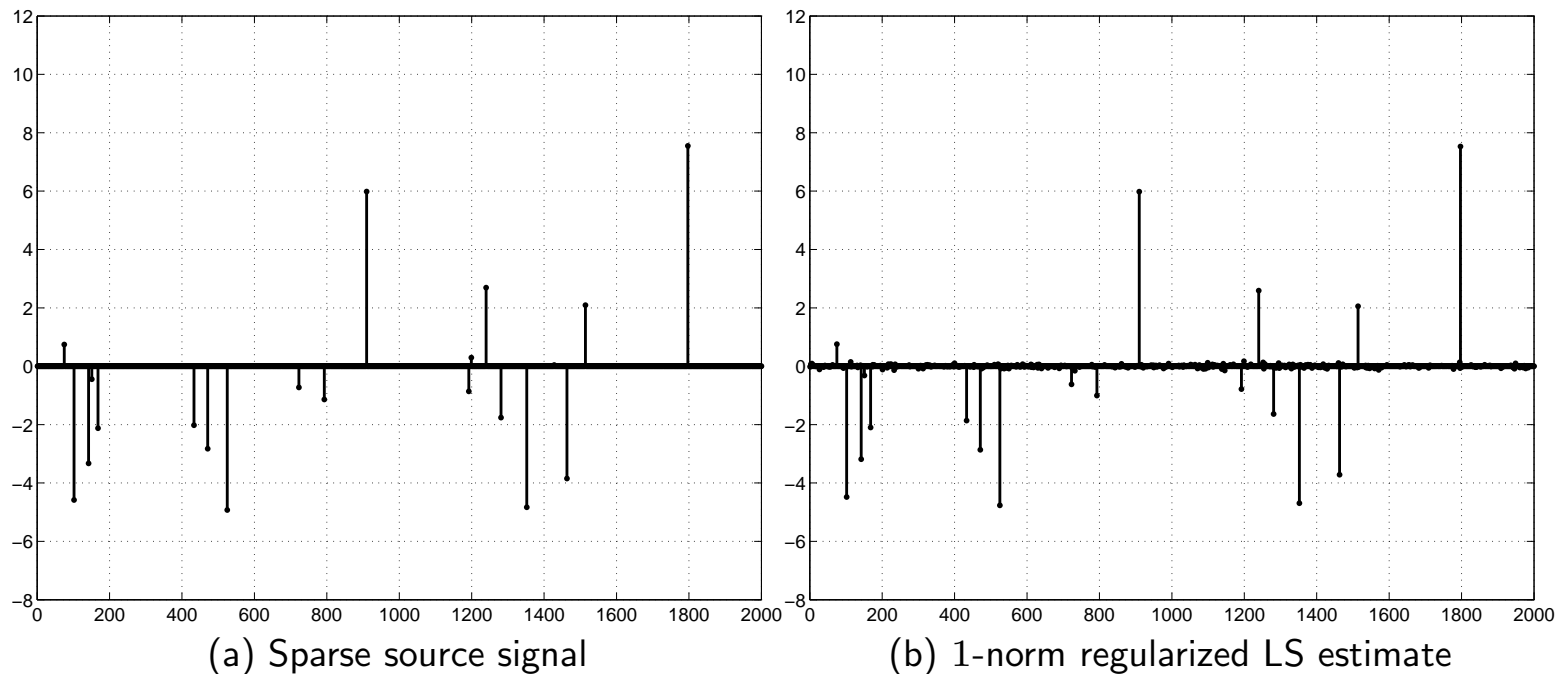
$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 \quad \text{s.t.} \quad \|\mathbf{x}\|_1 \leq \tau$$

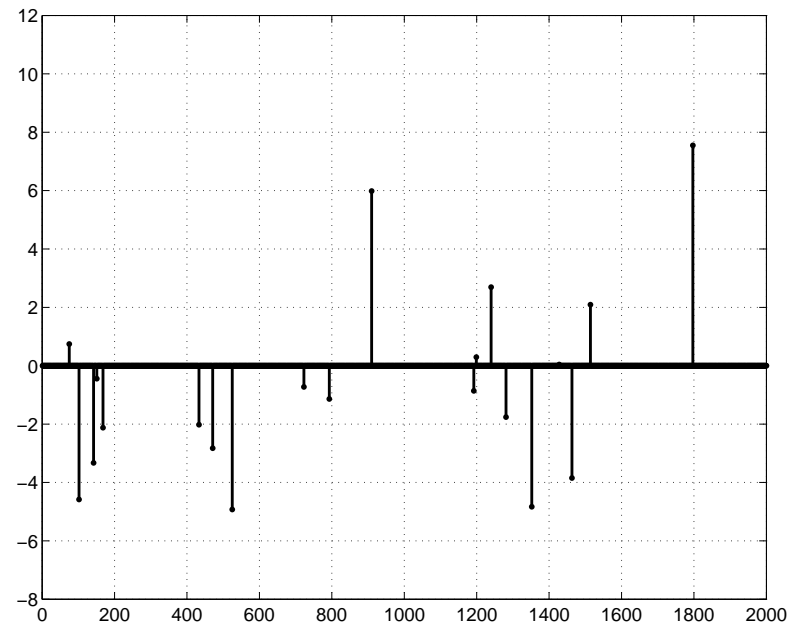
- when outliers exist in \mathbf{y} (i.e., some elements of \mathbf{y} are badly corrupted), we also want the residual $\mathbf{r} = \mathbf{y} - \mathbf{A}\mathbf{x}$ to be sparse; so,

$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_1 + \lambda \|\mathbf{x}\|_1.$$

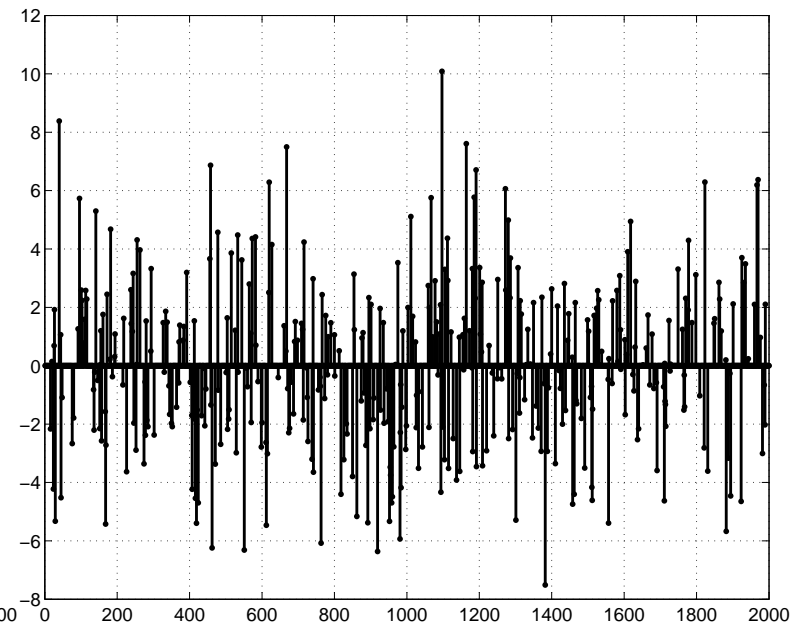
Toy Demonstration: Noisy Sparse Signal Reconstruction

- Sparse signal $\mathbf{x} \in \mathbb{R}^n$ with $n = 2000$ and $\|\mathbf{x}\|_0 = 20$.
- $m = 400$ noisy observations of $\mathbf{y} = \mathbf{A}\mathbf{x} + \boldsymbol{\nu}$, both a_{ij} and ν_i are randomly generated.
- 1-norm regularized LS $\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \lambda \|\mathbf{x}\|_1$ is used. $\lambda = 0.1$.



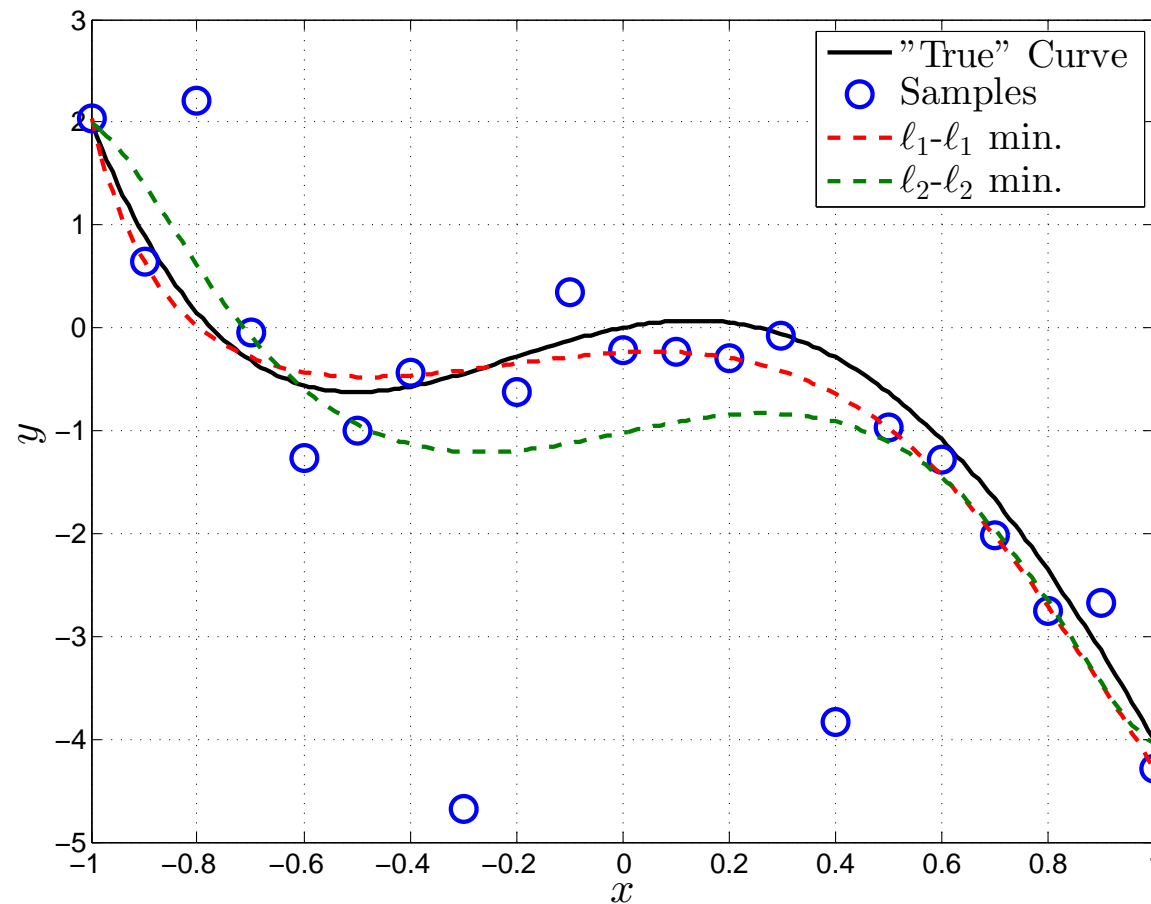


(c) Sparse source signal



(d) LS estimate

Toy Demonstration: Curve Fitting



The same curve fitting problem in Lecture 2. The guessed model order is $n = 18$.

$$\ell_2\text{-}\ell_2 \text{ min.: } \min \|\mathbf{y} - \mathbf{Ax}\|_2^2 + \lambda \|\mathbf{x}\|_2^2$$

$$\ell_1\text{-}\ell_1 \text{ min.: } \min \|\mathbf{y} - \mathbf{Ax}\|_1 + \lambda \|\mathbf{x}\|_1$$

Total Variation (TV) Denoising

- Scenario:

- estimate $\mathbf{x} \in \mathbb{R}^n$ from a noisy measurement $\mathbf{x}_{\text{cor}} = \mathbf{x} + \boldsymbol{\nu}$.
- \mathbf{x} is known to be piecewise linear, i.e., for most i we have

$$x_i - x_{i-1} = x_{i+1} - x_i \iff -x_{i+1} + 2x_i - x_{i-1} = 0.$$

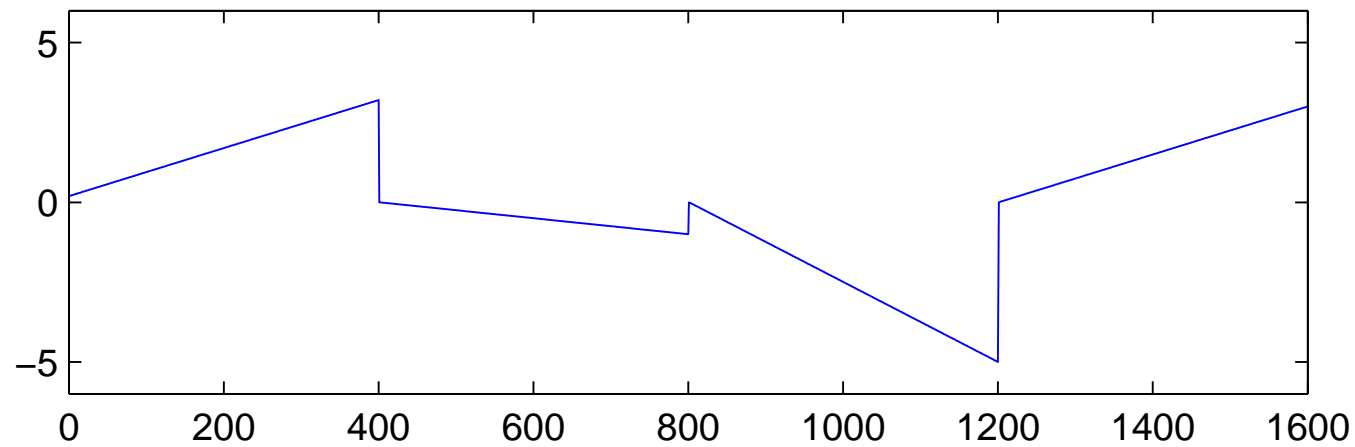
- equivalently, $\mathbf{D}\mathbf{x}$ is sparse, where

$$\mathbf{D} = \begin{bmatrix} -1 & 2 & 1 & 0 & \dots \\ 0 & -1 & 2 & 1 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \dots & \dots & -1 & 2 & 1 \end{bmatrix}.$$

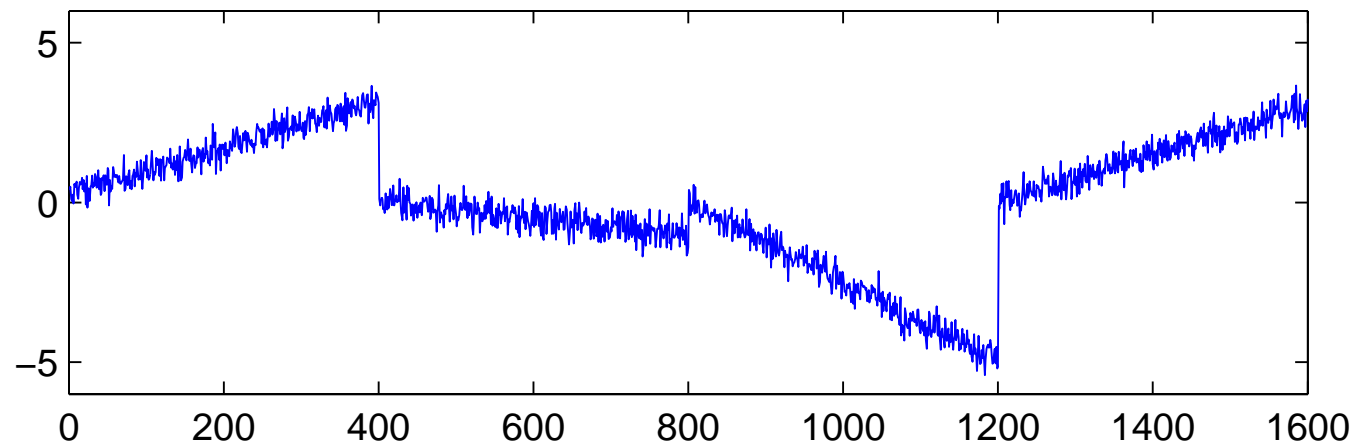
- TV denoising: estimate \mathbf{x} by solving

$$\min_{\mathbf{x}} \|\mathbf{x}_{\text{cor}} - \mathbf{x}\|_2^2 + \lambda \|\mathbf{D}\mathbf{x}\|_1$$

Source

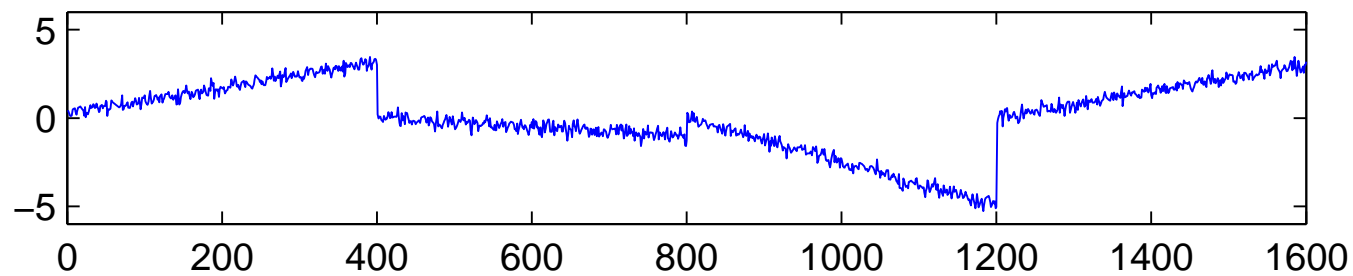


Corrupted by noise

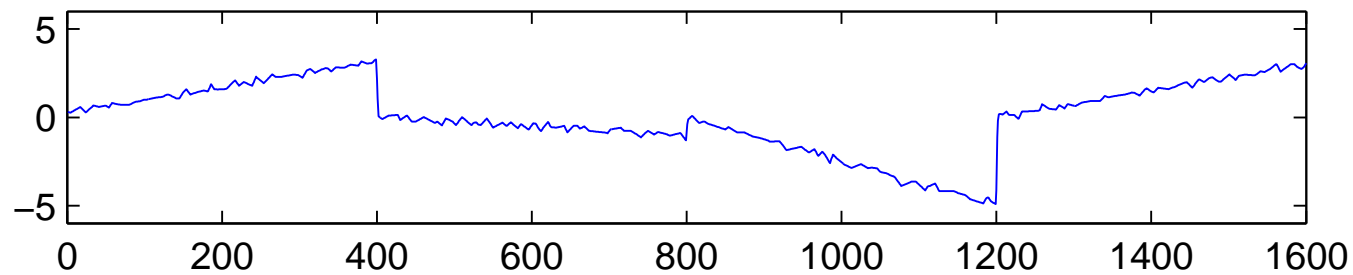


Original \mathbf{x} and corrupted \mathbf{x}_{cor}

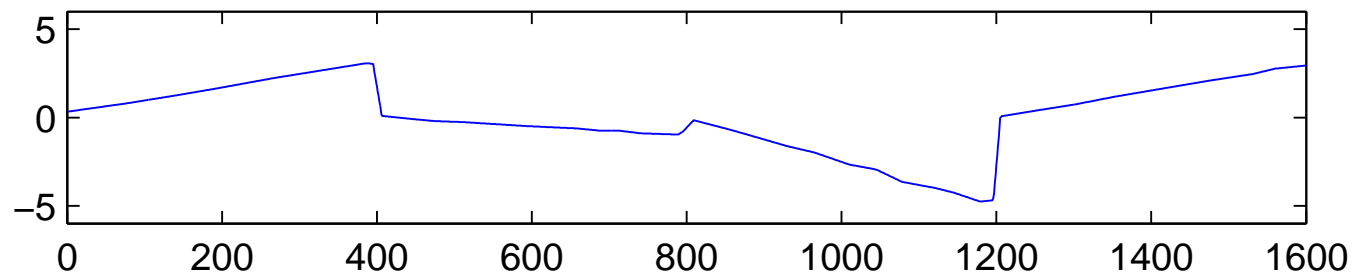
\hat{x} with $\lambda = 0.1$



\hat{x} with $\lambda = 1$

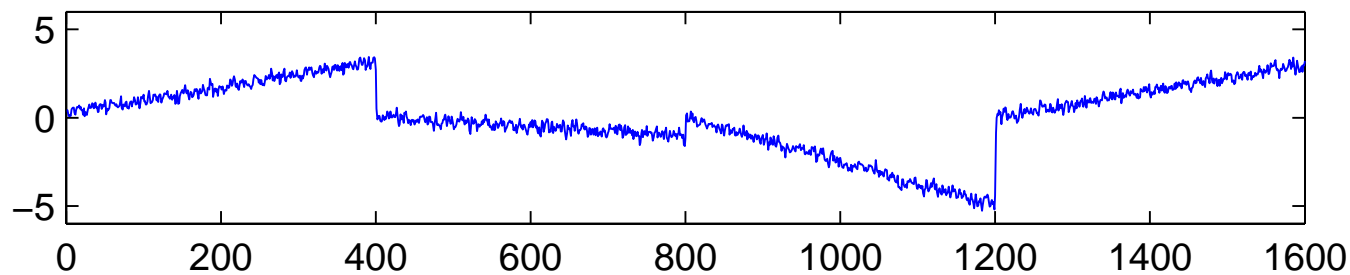


\hat{x} with $\lambda = 10$

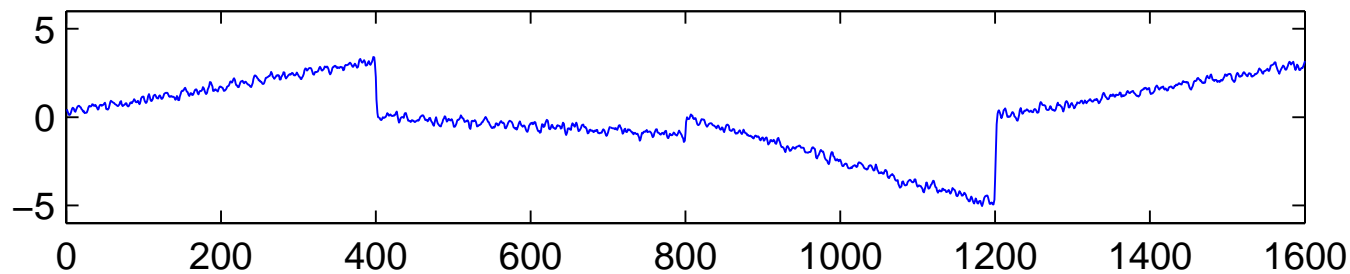


TV denoised signals for various λ 's.

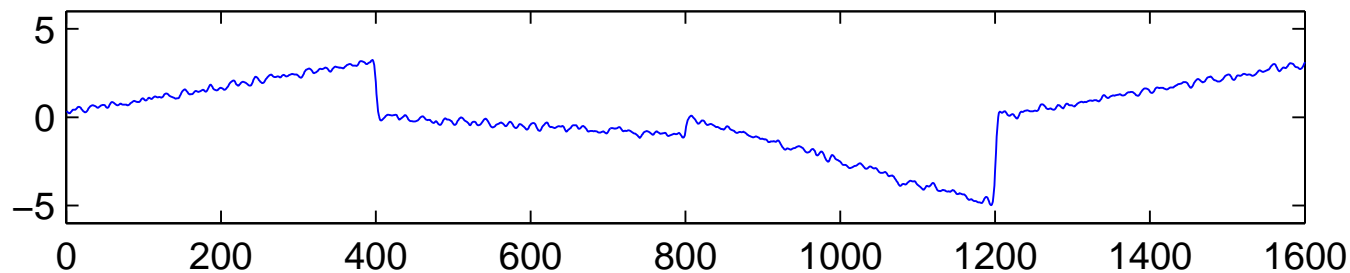
\hat{x} with $\lambda = 0.1$



\hat{x} with $\lambda = 1$



\hat{x} with $\lambda = 10$



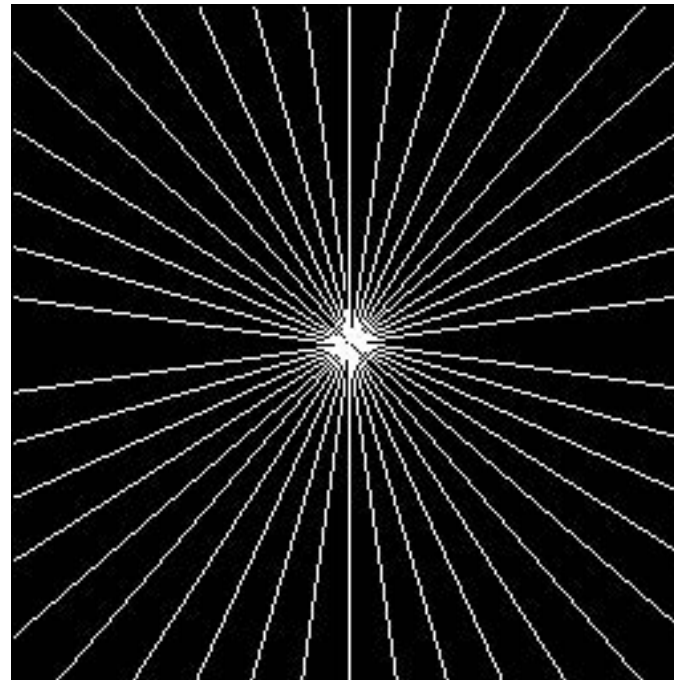
TV denoised signals via ℓ_2 regularization and for various λ 's.

Application: Magnetic Resonance Imaging (MRI)

Problem: MRI image reconstruction.



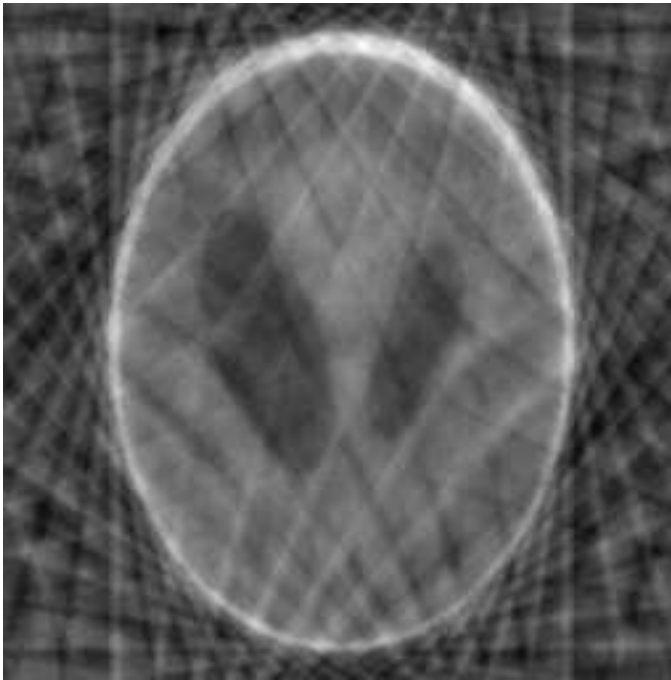
(a)



(b)

Fig. a shows the original test image. Fig. b shows the sampling region in the frequency domain. Fourier coefficients are sampled along 22 approximately radial lines. Source: [\[Candès-Romberg-Tao'06\]](#)

Application: MRI



(c)



(d)

Fig. c is the recovery by filling the unobserved Fourier coefficients to zero. Fig. d is the recovery by a TV minimization problem. Source: [\[Candès-Romberg-Tao'06\]](#)

Efficient Computations of the $\ell_2 - \ell_1$ Minimization Solution

- consider the $\ell_2 - \ell_1$ minimization problem

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{y} - \mathbf{Ax}\|_2^2 + \lambda \|\mathbf{x}\|_1.$$

- as mentioned, the problem is convex and there are many optimization algorithms custom-designed for it
 - some keywords for such algorithms: majorization-minimization (MM), ADMM, fast proximal gradient (or the so-called FISTA), Frank-Wolfe,...
- Aim: get some flavor of one particular algorithm, namely, MM, that is sufficiently “matrix” and is suitable for large-scale problems

MM for $\ell_2 - \ell_1$ Minimization: LS as an Example

- to see the insight of MM, we start with the plain old LS

$$\min_{\mathbf{x}} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2.$$

- observe that for a given $\bar{\mathbf{x}}$, one has

$$\begin{aligned} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 &= \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}} - \mathbf{A}(\mathbf{x} - \bar{\mathbf{x}})\|_2^2 \\ &= \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + \|\mathbf{A}(\mathbf{x} - \bar{\mathbf{x}})\|_2^2 \\ &\leq \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + c\|\mathbf{x} - \bar{\mathbf{x}}\|_2^2 \end{aligned}$$

for any $\mathbf{x} \in \mathbb{R}^n$ and for any $c \geq \sigma_{\max}^2(\mathbf{A})$

MM for $\ell_2 - \ell_1$ Minimization: LS as an Example

- let $c \geq \sigma_{\max}^2(\mathbf{A})$, and let

$$g(\mathbf{x}, \bar{\mathbf{x}}) = \|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + c\|\mathbf{x} - \bar{\mathbf{x}}\|_2^2$$

- we have

$$\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 \leq g(\mathbf{x}, \bar{\mathbf{x}}), \quad \text{for any } \mathbf{x}, \bar{\mathbf{x}} \in \mathbb{R}^n$$

$$\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 = g(\mathbf{x}, \mathbf{x}), \quad \text{for any } \mathbf{x} \in \mathbb{R}^n$$

- also,

$$\arg \min_{\mathbf{x} \in \mathbb{R}^n} g(\mathbf{x}, \bar{\mathbf{x}}) = \frac{1}{c} \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + \bar{\mathbf{x}}$$

- **Idea:** given an initial point $\mathbf{x}^{(0)}$, do

$$\mathbf{x}^{(k+1)} = \arg \min_{\mathbf{x} \in \mathbb{R}^n} g(\mathbf{x}, \mathbf{x}^{(k)}) = \frac{1}{c} \mathbf{A}^T (\mathbf{y} - \mathbf{A}\mathbf{x}^{(k)}) + \mathbf{x}^{(k)}, \quad k = 1, 2, \dots$$

- note: not very interesting at this moment as the above iteration is the same as gradient descent with step size $1/c$

MM for $\ell_2 - \ell_1$ Minimization: General MM Principle

- the example shown above is an instance of MM
- general MM principle:
 - consider a general optimization problem

$$\min_{\mathbf{x} \in \mathcal{C}} f(\mathbf{x})$$

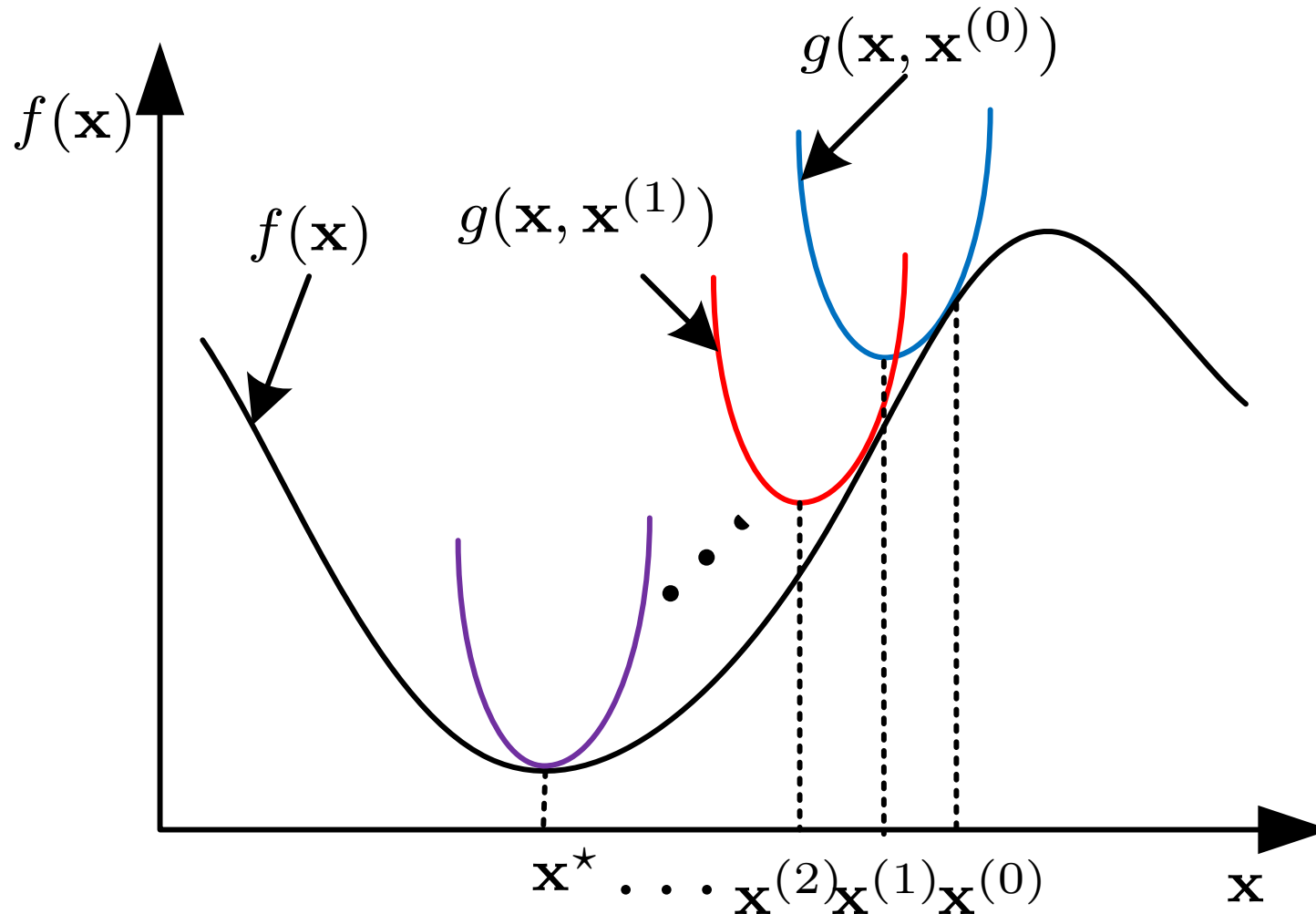
and suppose that f is hard to minimize directly

- let $g(\mathbf{x}, \bar{\mathbf{x}})$ be a **surrogate function** that is easy to minimize and satisfies

$$f(\mathbf{x}) \leq g(\mathbf{x}, \bar{\mathbf{x}}) \text{ for all } \mathbf{x}, \bar{\mathbf{x}}, \quad f(\mathbf{x}) = g(\mathbf{x}, \mathbf{x}) \text{ for all } \mathbf{x}$$

- MM algorithm: $\mathbf{x}^{(k+1)} = \arg \min_{\mathbf{x} \in \mathcal{C}} g(\mathbf{x}, \mathbf{x}^{(k)}), k = 1, 2, \dots$
- as a basic result, $f(\mathbf{x}^{(0)}) \geq f(\mathbf{x}^{(1)}) \geq f(\mathbf{x}^{(2)}) \dots$
- suppose that f is convex and \mathcal{C} is convex. MM is guaranteed to converge to an optimal solution under some mild assumption **[Razaviyayn-Hong-Luo'13]**

MM for $\ell_2 - \ell_1$ Minimization: General MM Principle



MM for $\ell_2 - \ell_1$ Minimization

- now consider applying MM to the $\ell_2 - \ell_1$ minimization problem

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2 + \lambda \|\mathbf{x}\|_1.$$

- let $c \geq \sigma_{\max}^2(\mathbf{A})$, and let

$$g(\mathbf{x}, \bar{\mathbf{x}}) = \frac{1}{2} (\|\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}\|_2^2 - 2(\mathbf{x} - \bar{\mathbf{x}})^T \mathbf{A}^T (\mathbf{y} - \mathbf{A}\bar{\mathbf{x}}) + c\|\mathbf{x} - \bar{\mathbf{x}}\|_2^2) + \lambda \|\mathbf{x}\|_1$$

– simply plug the same surrogate for $\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2^2$ we saw previously

- it can be shown that

$$\mathbf{x}^{(k+1)} = \text{soft} \left(\frac{1}{c} \mathbf{A}^T (\mathbf{y} - \mathbf{A}\mathbf{x}^{(k)}) + \mathbf{x}^{(k)}, \lambda/c \right)$$

where soft is called the soft-thresholding operator and is defined as follows: if $\mathbf{z} = \text{soft}(\mathbf{x}, \delta)$ then $z_i = \text{sign}(x_i) \max\{|x_i| - \delta, 0\}$

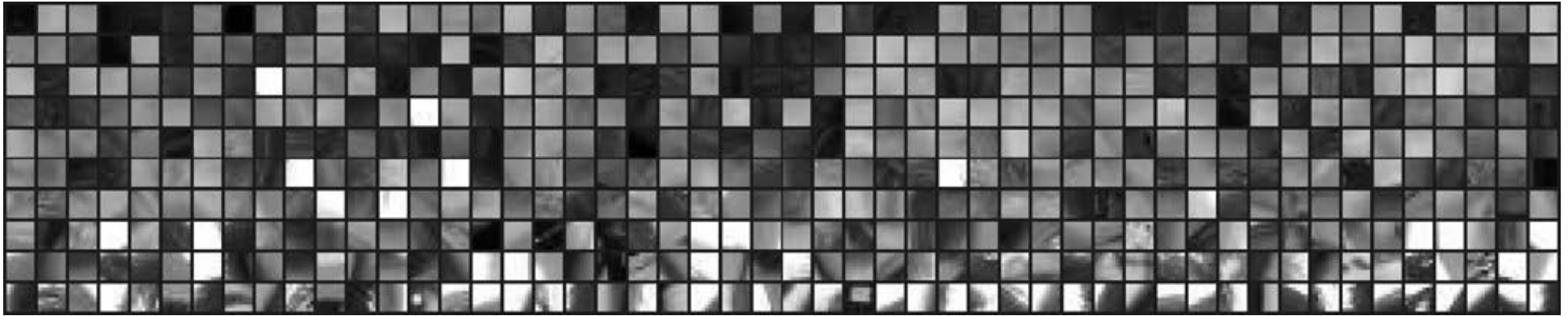
Dictionary Learning

- previously \mathbf{A} is assumed to be given
- how about learning a fat \mathbf{A} from data, as in matrix factorization?
- Dictionary learning (DL): given $\tau > 0$ and $\mathbf{Y} \in \mathbb{R}^{m \times n}$, solve

$$\begin{aligned} \min_{\mathbf{A} \in \mathbb{R}^{m \times k}, \mathbf{B} \in \mathbb{R}^{k \times n}} \quad & \sum_{i=1}^n \|\mathbf{y}_i - \mathbf{A}\mathbf{b}_i\|_2^2 \\ \text{s.t.} \quad & \|\mathbf{b}_i\|_0 \leq \tau, \quad i = 1, \dots, n \end{aligned}$$

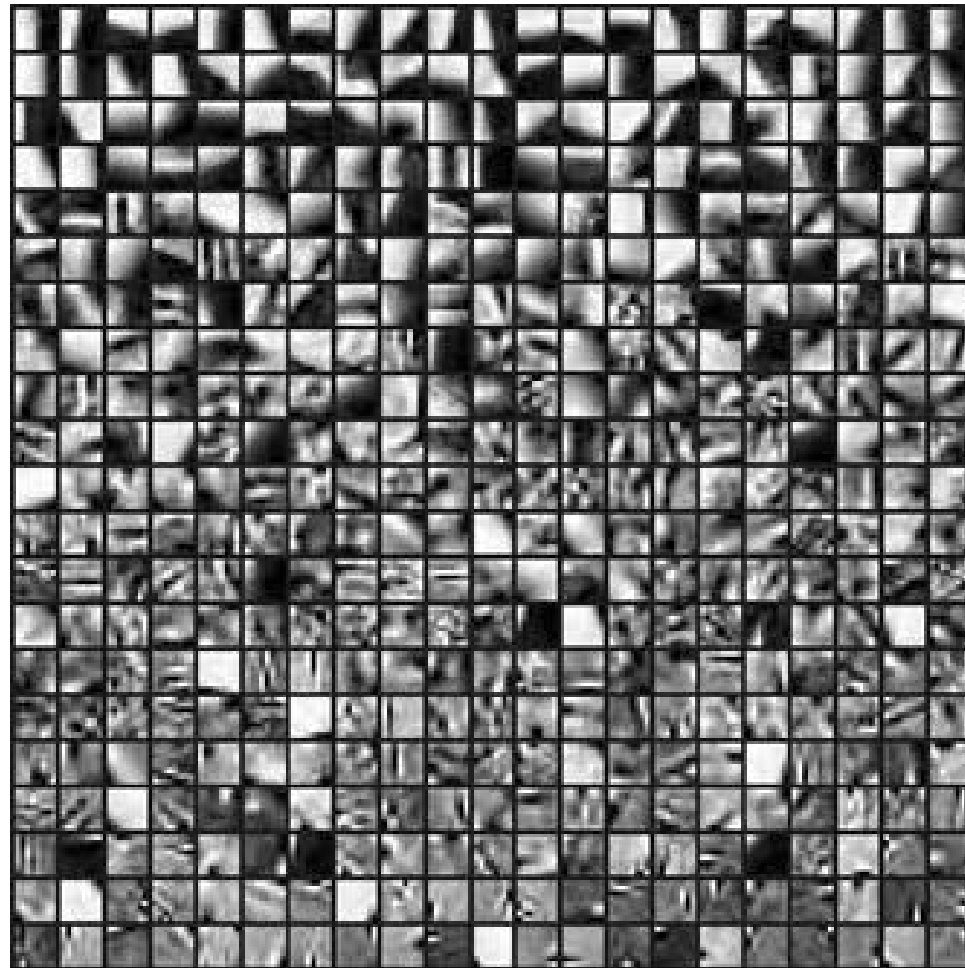
- DL considers $k \geq m$, and \mathbf{A} is called an overcomplete dictionary
- DL is handled by alternating optimization—the same approach in matrix fac.

Dictionary Learning



A collection of 500 random image blocks. Source: [\[Aharon-Elad-Bruckstein'06\]](#).

Dictionary Learning



The learned dictionary. Source: [\[Aharon-Elad-Bruckstein'06\]](#).

Part III: LS with Errors in A

LS with Errors in \mathbf{A}

- **Scenario:** errors exist in the system matrix \mathbf{A}
- **Aim:** mitigate the effects of the system matrix errors on the LS solution
- there are many ways to do so, and we look at two
- **Total LS (TLS):**

$$\min_{\mathbf{x} \in \mathbb{R}^n, \mathbf{\Delta} \in \mathbb{R}^{m \times n}} \|\mathbf{y} - (\mathbf{A} + \mathbf{\Delta})\mathbf{x}\|_2^2 + \|\mathbf{\Delta}\|_F^2$$

– minimally perturb the system matrix for best fitting in the Euclidean sense

- **Robust LS :**

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\mathbf{\Delta} \in \mathcal{U}} \|\mathbf{y} - (\mathbf{A} + \mathbf{\Delta})\mathbf{x}\|_2^2$$

for some pre-determined uncertainty set $\mathcal{U} \subset \mathbb{R}^{m \times n}$

– robustify the LS via a worst-case means

Total LS

$$\min_{\mathbf{x} \in \mathbb{R}^n, \mathbf{\Delta} \in \mathbb{R}^{m \times n}} \|\mathbf{y} - (\mathbf{A} + \mathbf{\Delta})\mathbf{x}\|_2^2 + \|\mathbf{\Delta}\|_F^2$$

- does not seem to have a closed-form solution at first sight
- turns out to have a closed-form solution under some mild assumptions
- assume \mathbf{A} to be of full column rank with $m \geq n + 1$
- let $\mathbf{C} = [\mathbf{A} \ \mathbf{y}]$, and let v_{n+1} be the $(n + 1)$ th right singular value of \mathbf{C} . If

$$\text{rank}(\mathbf{C}) = n + 1, \quad v_{n+1,n+1} \neq 0,$$

then

$$\mathbf{x}_{\text{TLS}} = -\frac{1}{v_{n+1,n+1}} \begin{bmatrix} v_{1,n+1} \\ \vdots \\ v_{n,n+1} \end{bmatrix}$$

is a TLS solution

- see [\[Golub-Van Loan'12\]](#) for further discussion on issues like $v_{n+1,n+1} \neq 0$

Proof Sketch of the TLS Solution

- idea: turn the TLS problem to a low-rank matrix approximation problem
- by a change of variables

$$\mathbf{C} = [\mathbf{A} \ \mathbf{y}] \in \mathbb{R}^{m \times (n+1)}, \quad \mathbf{D} = [\mathbf{\Delta} \ (\mathbf{A} + \mathbf{\Delta})\mathbf{x}] \in \mathbb{R}^{m \times (n+1)},$$

the TLS problem can be formulated as

$$\min_{\mathbf{x}, \mathbf{D}} \|\mathbf{C} - \mathbf{D}\|_F^2 \quad \text{s.t.} \quad \mathbf{D} \begin{bmatrix} \mathbf{x} \\ -1 \end{bmatrix} = \mathbf{0} \quad (\dagger)$$

- the constraint in (\dagger) , together with $m \geq n + 1$, implies $\text{rank}(\mathbf{D}) \leq n$
- or, we can equivalently rewrite (\dagger) as

$$\min_{\mathbf{x}, \mathbf{D}} \|\mathbf{C} - \mathbf{D}\|_F^2 \quad \text{s.t.} \quad \text{rank}(\mathbf{D}) \leq n, \quad \mathbf{D} \begin{bmatrix} \mathbf{x} \\ -1 \end{bmatrix} = \mathbf{0}$$

Proof Sketch of the TLS Solution

- consider a *relaxation* of (\dagger) :

$$\min_{\mathbf{D}} \|\mathbf{C} - \mathbf{D}\|_F^2 \quad \text{s.t. } \text{rank}(\mathbf{D}) \leq n, \quad (\ddagger)$$

where we drop the constraint $\mathbf{D} \begin{bmatrix} \mathbf{x} \\ -1 \end{bmatrix} = \mathbf{0}$

- let \mathbf{D}^* be a solution to (\ddagger) . If there exists an \mathbf{x} such that $\mathbf{D}^* \begin{bmatrix} \mathbf{x} \\ -1 \end{bmatrix} = \mathbf{0}$, \mathbf{D}^* is also a solution to (\dagger) and \mathbf{x} is a TLS solution
- let $\mathbf{C} = \sum_{i=1}^{n+1} \sigma_i \mathbf{u}_i \mathbf{v}_i^T$ be the SVD
- by the Eckart-Young-Mirsky theorem, a solution to (\ddagger) is $\mathbf{D}^* = \sum_{i=1}^n \sigma_i \mathbf{u}_i \mathbf{v}_i^T$.
- as a basic fact of SVD, we have $\mathbf{D}^* \mathbf{v}_{n+1} = \mathbf{0}$.
- thus, if $v_{n+1,n+1} \neq 0$, we have the desired TLS solution

Robust LS

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\Delta \in \mathcal{U}} \|\mathbf{y} - (\mathbf{A} + \Delta)\mathbf{x}\|_2$$

- consider the case of $\mathcal{U} = \{\Delta \in \mathbb{R}^{m \times n} \mid \|\Delta\|_2 \leq \lambda\}$ for some $\lambda > 0$
- the robust LS problem can be shown to be equivalent to

$$\min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2 + \lambda \|\mathbf{x}\|_2$$

- **Observations and Implications:**
 - the equivalent form of the robust LS is very similar to (but not exactly the same as) the previous ℓ_2 -regularized LS
 - robustification is equivalent to regularization
- it can be shown that the same equivalence holds if we replace the uncertainty set by $\mathcal{U} = \{\Delta \in \mathbb{R}^{m \times n} \mid \|\Delta\|_F \leq \lambda\}$

Proof Sketch of the Robust LS Equivalence Result

- by the definition of induced norms, we have

$$\|\Delta\|_2 \leq \lambda \iff \|\Delta\mathbf{x}\|_2 \leq \lambda\|\mathbf{x}\|_2 \quad \text{for all } \mathbf{x} \in \mathbb{R}^n$$

- then, for any $\mathbf{x} \in \mathbb{R}^n$ and for any $\Delta \in \mathcal{U}$,

$$\begin{aligned} \|\mathbf{y} - (\mathbf{A} + \Delta)\mathbf{x}\|_2 &\leq \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2 + \|\Delta\mathbf{x}\|_2 \\ &\leq \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2 + \lambda\|\mathbf{x}\|_2, \end{aligned} \quad (*)$$

and note that the 1st equality above holds if $\mathbf{y} - \mathbf{A}\mathbf{x} = -\alpha\Delta\mathbf{x}$ for some $\alpha \geq 0$, and the 2nd equality above holds if \mathbf{x} is the 1st right singular vector of Δ

- consider the case of $\mathbf{x} \neq \mathbf{0}$, $\mathbf{y} - \mathbf{A}\mathbf{x} \neq \mathbf{0}$. It can be verified that

$$\Delta = -\frac{\lambda}{\|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2\|\mathbf{x}\|_2}(\mathbf{y} - \mathbf{A}\mathbf{x})\mathbf{x}^T$$

attains the equalities in $(*)$ and lies in \mathcal{U}

- the other cases of \mathbf{x} are handled in a similar fashion

More Robust LS Equivalences

- denote $\mathcal{U}_{q,p} = \{\Delta \in \mathbb{R}^{m \times n} \mid \|\Delta \mathbf{x}\|_p \leq \lambda \|\mathbf{x}\|_q \ \forall \mathbf{x}\}$, where $p, q \geq 1$. We have

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\Delta \in \mathcal{U}_{q,p}} \|\mathbf{y} - (\mathbf{A} + \Delta)\mathbf{x}\|_p = \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_p + \lambda \|\mathbf{x}\|_q$$

- proof: almost the same as the previous case
- some interesting special cases:

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\Delta \in \mathcal{U}_{2,1}} \|\mathbf{y} - (\mathbf{A} + \Delta)\mathbf{x}\|_2 = \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_2 + \lambda \|\mathbf{x}\|_1$$

$$\min_{\mathbf{x} \in \mathbb{R}^n} \max_{\substack{\Delta \in \mathbb{R}^{m \times n} \\ \|\delta_i\|_1 \leq \lambda \ \forall i}} \|\mathbf{y} - (\mathbf{A} + \Delta)\mathbf{x}\|_1 = \min_{\mathbf{x} \in \mathbb{R}^n} \|\mathbf{y} - \mathbf{A}\mathbf{x}\|_1 + \lambda \|\mathbf{x}\|_1$$

- Implication:** ℓ_1 regularization may also be seen as an act of robustification
- suggested reading: **[Bertsimas-Copenhaver'17]**, including extension to PCA

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