ALGOLAB TUTORIAL #9

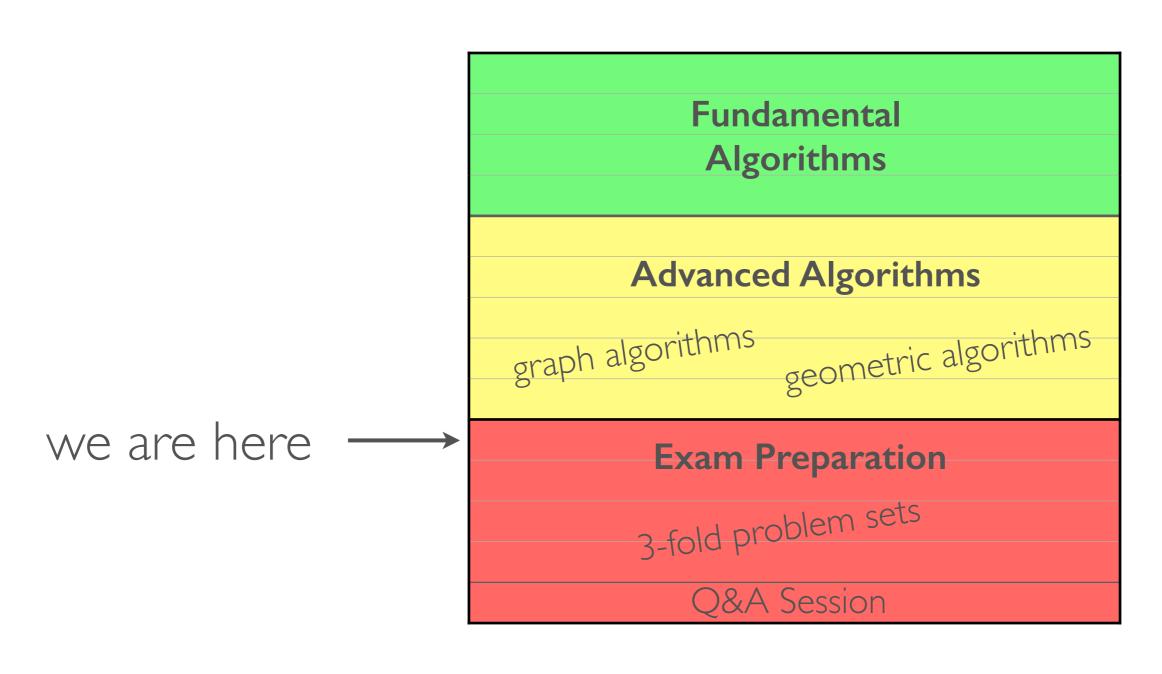
Exam Preparation Week I

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Contents

- Information about the test exam
- Mow to solve Algolab problems (meta-guidelines)
- Discussion of Problems (Firsthit, HINI)

ALGOLAB TIMELINE



Exam (afawk): Jan 27 and Feb 1, 2016, 13-19.

TEST EXAM

Date/time: Tu, Dec 8, 2015, 17:00-19:15, ETH HG.

No PotW on Mon, Dec 7, 2015.

Participation is optional.

Prerequisite: Being registered for Algolab exam.

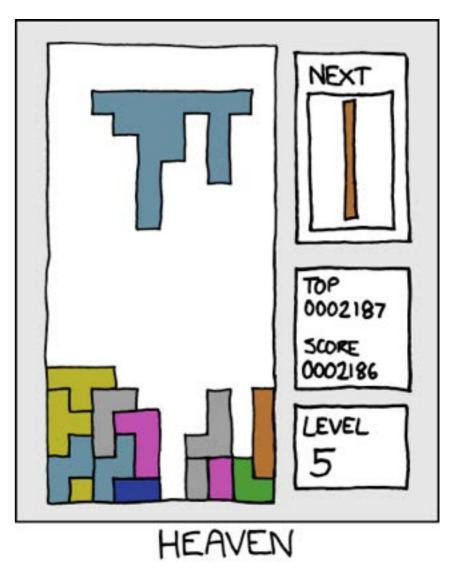
You have to register on the forum until latest

Mon Nov 23, 2015, 10:00 AM Zurich time.

Computer activity (screen) is logged.

HOW TO SOLVE PROBLEMS

- Know what to know
- Understand your task
- Find an appropriate model
- Design an efficient algorithm
- Implement that algorithm
- Avoid "stupid" mistakes



http://xkcd.com/888/

KNOW WHAT TO KNOW

- ▶ Both the material from the tutorials and the collection of problems form the contents of this course.
- Ney concepts, techniques, and skills were covered in the tutorials and/or practiced in a problem.
- Also meta skills such as time management play a role here (practiced in PotWs). THE KEY TO LEADING CHOOSE GOALS, BUILD A SCHEDILE AND HAVE THE SCHEDILE AN







http://xkcd.com/874/

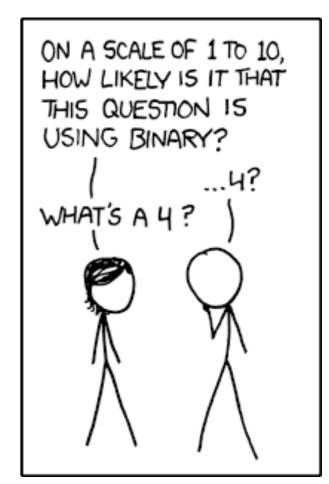
KNOW WHAT TO KNO

We will not ask you to do something drastically different from what you have seen during the semester.

The problems from the exam preparation weeks give you a good idea of how problems in the exam may look like.

If you use a data structure/algorithm/ technique that was not covered, you are most likely not solving the problem You go down a risky road. If that works in a way we intended.

out, kudos to you for the original approach! If not ... you knew the risks ...



http://xkcd.com/953/

UNDERSTANDYOURTASK

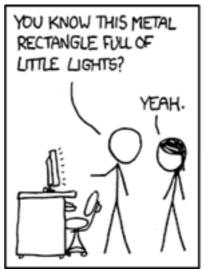
- Read the problem statement carefully.
- Read the problem statement again.

Make sure that you understand what is asked. Do not make any assumptions/ interpretations that are not clearly supported by what is written.

- Check the provided example(s) and if they concur with your understanding. These examples are part of the problem description.
- If (and only if) you think the problem is not clearly stated, ask an assistant.

 | YOU KNOW THIS METAL | I SPEND MOST OF MY LIFE | BUT TOODAY, THE PA

The assistants are not there to confirm your understanding. They will answer: ``Please read the problem statement. It is clearly stated.'' - unless they agree it is not clear.







http://xkcd.com/722/

FIND A MODEL

Rephrase the problem in abstract/ mathematical terms.

- let (using terms like graph, vertex, edge, component, matching, point, line, matrix, relation, inequality, ...) rather than planes, aliens, countries, or antennas.
- Sometimes this task is straightforward and sometimes there are choices to make.
- The goal is to get rid of the story and unveil the algorithmic problem.



Body of Knowledge (Jaume Plensa, 2010)

ALGORITHM DESIGN

- Mow can you attack this problem?
- Do not get caught in the story!

You should let yourself get inspired by what you already know. Not on the story-level, but on the algorithmic level instead!

Try to think about different alternatives: evaluate them briefly, which look promising?

(LP, network flow, maximum matching, dynamic programming, Delaunay/Voronoi, minimum enclosing shapes, greedy, scan, binary search, shortest paths,...)

- If you have no idea, consider another problem.
- If your try gets stuck, consider a different approach.

Performing such meta-searching effectively is an objective of this course.

IMPI FMFNTATION

- Every problem can be solved with no more than
 - ~ 100 lines of well-written code.
- Use suitable data types for input/output processing Unlike for some earlier problems we will not tell you which (precision vs. speed). type to use, because you were taught all the necessary bits...
- Avoid premature optimizations.
- Practice helps a lot...

That's why this is a lab. The more you practice, the less likely it is that you run into a particular issue for the first time during the exam...

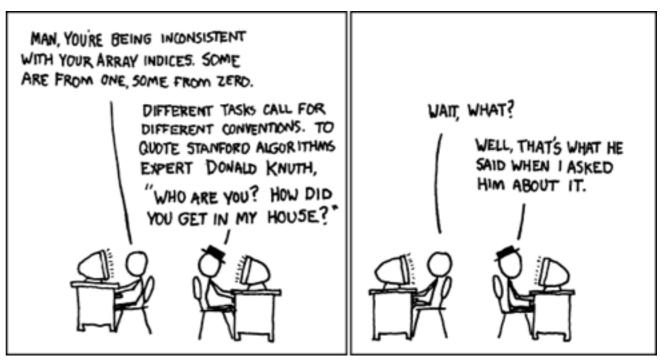


THE REASON I AM SO INEFFICIENT

http://xkcd.com/1445/

AVOID "STUPID" MISTAKES

- Submit to the right problem (correct link).
- Read all input, even if the result is determined half the way along already. Otherwise, you mess up a possibly following problem instance.
- > std::ios_base::sync_with_stdio(false);
- ▶ TMMTL ...
 - practice helps...



WHY ARE THERE MIRRORS ABOVE, BEDS

WHY IS THERE NOT A POKEMON MMO WHY ISTHERE LAUGHING IN TV SHOWS WHY ARE THERE DOORS ON THE FREEWAY

WHY ARE THERE TWO SLASHES AFTER HTTP WHY IS THERE AN ARROW ON AANG'S HEAD THERE, MUSTACHES ON CLOTHES WHY ARE THERE MUSTACHES ON CARS WHY ARE THERE MUSTACHES EVERYWHERE WHY ARE THERE SO MANY BIRDS IN OHIO

WHY ARE THERE SCARY SOUNDS IN MINECRAFT WHY IS THERE. KICKING IN MY STOMACH

SQUIRRELS

WHY ARE THERE

WHY DO TWINS HAVE DIFFERENT FINGERPRINTS \WHY IS HT WHY DO YOUR BOOKS HURT WHY ARE AMERICANS AFRAID OF DRAGONS WHY IS THERE A RED LINE THROUGH HTTPS ON FACEBOOK

GHOSTS

WHY DO AMERICANS CALL

WHY IS SPACE BLACK WHY IS OUTER SPACE SO COLD WHY IS OHIO WEATHER SO WEIRD TO WHY IS NASA SHUTTING DOWN

. THERE MALE AND FEMALE WHY ARE THERE BRIDESMAIDS WHY ARE THERE TINY SPIDERS IN MY HOUSE

UNIV ARE OLD KUNSONS DIFFERENT TO WHY DO SPIDERS COME INSIDE IN MY ARE THERE LOTS OF SPIDERS IN MY HOUSE E WHY ARE THERE SPIDERS IN MY ROOM AWHY ARE THERE SO MANY SPIDERS IN MY ROOM

TI WHY IS THERE NO GPS IN LAPTOPS 🗲 MHY DO KNEES CLICK ₹ WHY IS PROGRAMMING SO HARD WHY AREN'T THERE E GRADES WHY IS ISOLATION BAD WHY DO TREES DIE WHY DO BOYS LIKE ME O

WHY DO TREES DIE WHY DON'T BOYS LIKE ME O

WHY AREN'T POKEHON REAL WHY IS THERE ALLIANS A JAWA UPDATE
WHY AREN'T POKEHON REAL WHY ARE THERE RED DOTS ON HY THISHS JO

WHY AREN'T BULLETS SHARP WHY IS LYING GOOD THE



WHY IS THERE AN OWL ON THE DOLLAR BILL WHY ARE THERE TWO SPOOKS

VESUVIUS THERE MWHY ARE WRESTLERS ALWAYS WET 🔫 WHY ARE OCEANS BECOMING MORE ACIDIC 🗢

WHY AREN'T MY QUAIL LAYING EGGS WHY AREN'T MY QUAIL EGGS HATCHING & WHY IS STEALING WRONG COWHY AREN'T THERE ANY FOREIGN MILITARY BASES IN AMERICA

MHY IS HTTPS IMPORTANT WHY AREN'T MY arms Growing

WHY ARE THERE SO MANY CROWS IN ROCHESTER, MIN

WHY IS THERE AN OWL IN MY BACKYARD

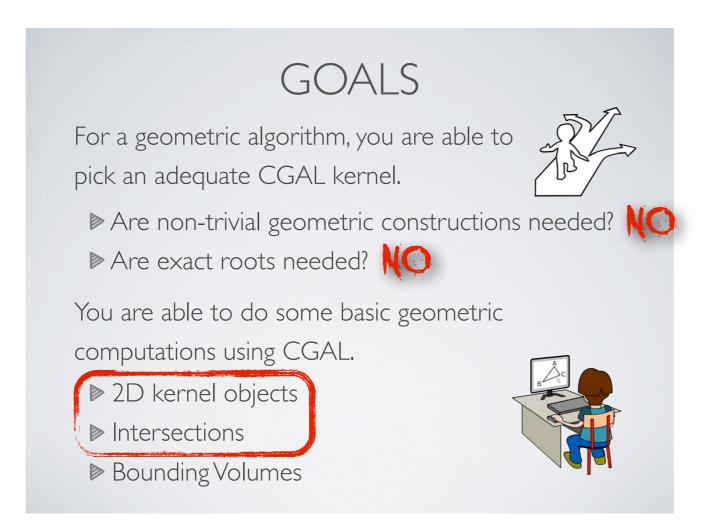
WHY IS THERE AN OWL OUTSIDE MY WINDOW

WHY ARE THERE HELICOPTERS CIRCLING MY HOUSE ARETHERE DUCKS IN MY POOL WHY IS JESUS WHITE UNITY IS THERE LIQUID IN MY EAR WHY DO Q TIPS FEEL GOOD WHY DO GOOD PEOPLE DIE WHY AREN'T



HIT

Problem: Given a ray and some line segments, does the ray intersect any segment?



Do not just guess, make a conscious decision!

Of course, you can just try all three kernels.

But often this choice also impacts the design of the algorithm.

No non-trivial constructions needed

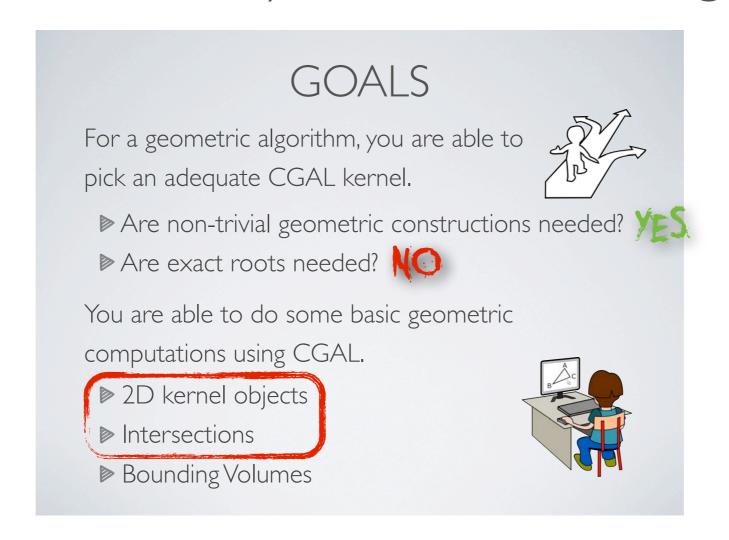


CGAL::Exact_predicates_inexact_constructions_kernel



```
#include <CGAL/Exact_predicates_inexact_constructions_kernel.h>
#include <iostream>
typedef CGAL::Exact_predicates_inexact_constructions_kernel K;
int main()
  std::ios_base::sync_with_stdio(false);
  for (std::size_t n; std::cin >> n && n > 0;) {
    K::Ray_2 r;
    std::cin >> r; ←
                                   With 52-bit numbers (-> double) and EPIC-kernel this is fine here.
    bool found = false;
                                   In general, be careful (-> IO performance slide from CGAL tutorial I)...
    do {
      K::Segment_2 s;
       std::cin >> s; ≰
       if (!found && CGAL::do_intersect(s,r)) found = true;
    } while (--n > 0);
    std::cout << (found ? "yes" : "no") << std::endl;</pre>
```

Problem: Given a ray and some line segments, where does the ray first intersect a segment?



Need intersection point

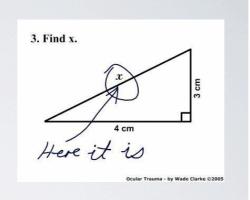


CGAL::Exact_predicates_exact_constructions_kernel

Problem: Given a ray and some line segments, where does the ray first intersect a segment?

PREREQUISITES

You know basic Euclidean geometry (e.g., distance/area/volume, angles, Pythagoras, ...) and can apply this knowledge to describe and analyze problems, to design models and algorithms.



You know basic algorithmic techniques (e.g., D.P., binary search, sorting, line sweep...). You skillfully combine them with the geometric techniques discussed here.

Most expensive operation is intersection construction



aim to minimize the number of constructions

Problem: Given a ray and some line segments, where does the ray first intersect a segment?

aim to minimize the number of constructions How?

- Find a segment s that intersects the ray r.

 If no such segment exists, we are done.
- Clip r at snr to a segment rc.
- Continue to test the remaining segments against rc; clip rc at each intersection.



no intersection construction for segments that intersect r but not rc

HRSTHIT

Problem: Given a ray and some line segments, where does the ray first intersect a segment?

no intersection construction for segments that intersect r but not rc

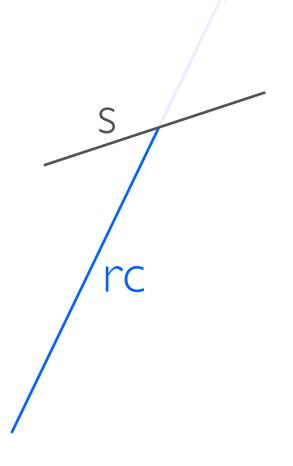
What's the gain?

If we find the first segment hit by rearly, no further intersection is constructed.



randomize segment order.

Claim. Only O(log n) intersections constructed in expectation.



Consider a permutation $x = x_1,...,x_n$ of $\{1,...,n\}$ u.a.r. and the sequence $m = m_0, m_1,..., m_n$ with $m_i = \min \{x_1,...,x_i\}$. How many changes occur in this sequence? $m_i = -\infty$ (indices $i \in \{1,...,n\}$ with $m_i \neq m_{i-1}$)

Example. n = 7 and x = 4,2,6,1,7,3,5 m = 4,2,2,1,1,1,1 with 3 changes

 $\text{Expected \#changes} = E(\Sigma C_i) = \Sigma E(C_i) \qquad c_i = \left\{ \begin{smallmatrix} I_i, \text{ if } m_i \neq m_{i-1} \\ 0_i, \text{ otherwise} \end{smallmatrix} \right.$

$$= 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} = H_n = \Theta(\log n)$$

Same as Firsthit with segment shortening (intersection is only constructed if new segment is hit before)



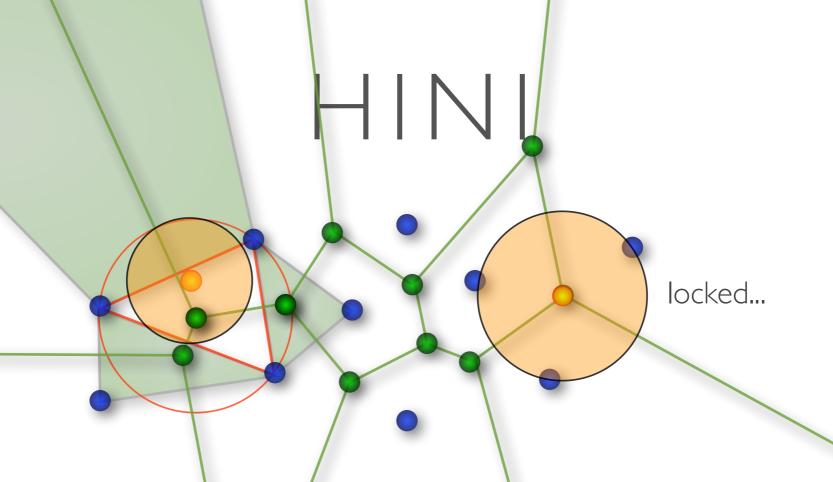
Claim

```
#include <CGAL/Exact_predicates_exact_constructions_kernel.h>
#include <vector>
#include <algorithm>
#include <stdexcept>
typedef CGAL::Exact predicates exact constructions kernel K;
void find_hit(std::size_t n) {
  // read input
  K::Ray_2 r;
  std::cin >> r;
  std::vector<K::Segment_2> segs;
  segs.reserve(n);
  double ix, iy, jx, jy;
  for (std::size t i = 0; i < n; ++i) {
    // as each coordinate can be represented as double, this is
    // significantly faster than K::Segment 2 s; std::cin >> s;
    std::cin \gg ix \gg jy \gg jy;
    segs.push_back(K::Segment_2(K::Point_2(ix,iy),K::Point_2(jx,jy)));
  std::random_shuffle(segs.begin(), segs.end());
                                important to avoid worst case orders
int main()
  std::cin.sync with stdio(false);
  std::cout << std::setiosflags(std::ios::fixed) << std::setprecision(0);</pre>
  for (std::size t n; std::cin >> n && n > 0;)
    find hit(n);
  return 0;
}
```



```
// find some segment hit by r
K::Segment_2 rc(r.source(), r.source()); // clipped ray
std::size_t i = 0;
for (; i < n; ++i)
  if (CGAL::do_intersect(segs[i], r)) {
    shorten_segment(rc, CGAL::intersection(segs[i], r));
    break;
if (i == n) { std::cout << "no\n"; return; }</pre>
// check remaining segments against rc
while (++i < n)
  if (CGAL::do_intersect(segs[i], rc))
    shorten_segment(rc, CGAL::intersection(segs[i], r)); // not rc!
std::cout << floor_to_double(rc.target().x()) << " "</pre>
          << floor_to_double(rc.target().y()) << "\n";
```

```
typedef std::result_of<K::Intersect_2(K::Ray_2,K::Segment_2)>::type IT;
// clip/set target of s to o
void shorten_segment(K::Segment_2& s, const IT& o)
  if (const K::Point_2* p = boost::get<K::Point_2>(&*o))
    s = K::Segment_2(s.source(), *p);
  else if (const K::Segment_2* t = boost::get<K::Segment_2>(&*o))
    // select endpoint of *t closer to s.source()
    if (CGAL::collinear_are_ordered_along_line
      (s.source(), t->source(), t->target()))
      s = K::Segment_2(s.source(), t->source());
    else
      s = K::Segment_2(s.source(), t->target());
  else
    throw std::runtime_error("Strange segment intersection.");
```



Problem. How to move a disk D without colliding with a given point set P?

Approach

- Locate triangle t where we are; is it safe to be there?
- Which of the three neighboring triangles can we move to?

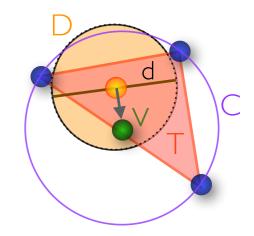


DFS on the Voronoi/Delaunay

HINI

Lemma. If the center of D is inside a triangle T of the DT of P and the placement is safe, then D can be safely moved straight s.t. its center coincides with the circumcenter of T.

- \triangleright c := center of D
- C ≔ circumdisk of T
- V ≔ center of C
- \triangleright d = diametrical segment of D orthogonal to vc.

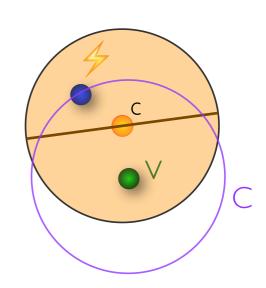


Dovious if d lies inside C.

The disk C is empty and the part of D outside C can only shrink (as a subset of \mathbb{R}^2) during the movement. As the radius of D is at most the radius of C, the final placement is safe.

What if the endpoints of d are outside C?

As the center of D is contained in T, there must be a vertex of T (which is a point from P) "behind" d (as seen from v). But this region lies inside D, in contradiction to D being empty of points from P.



HINI

Challenges

Use an external map where each access costs O(log n) time. Or store it in a face directly to achieve constant access time.

- How to mark faces as visited to avoid cycling?
- These marks have to be initialized. If this is done explicitly for each query, it can be costly if there are many queries.

A query may escape or get stuck quickly, without having to search the whole graph...

Repetitive initialization can be avoided, e.g., using a counter "visited in query i" rather than a bool.

Obs. In the worst case, a query has to search its way through a linear number of faces.

Better: precompute for each triangle/Voronoi vertex the maximum escape radius (Dijkstra in O(n log n)).

HINI

```
typedef CGAL::Triangulation vertex base 2<K>
                                                                  Vb;
typedef CGAL::Triangulation_face_base_with_info_2<int,K>
                                                                  Fb;
typedef CGAL::Triangulation_data_structure_2<Vb,Fb>
                                                                  Tds;
typedef CGAL::Delaunay triangulation 2<K,Tds>
                                                                  Delaunay;
try {
 if (r <= 0) throw true;
                                                               - s = guery position, r = (squared) guery radius
    Face_handle f = t.locate(s);
    if (CGAL::squared_distance(s, t.nearest_vertex(s, f)->point()) < r)</pre>
    throw false:
    // DFS
    std::vector<Face_handle> stack;
    stack.push_back(f);
                                               — query number
    f \rightarrow info() = i; \leftarrow
    while (!stack.empty()) {
    f = stack.back();
      stack.pop_back();
      if (t.is_infinite(f)) throw true;
      for (int j = 0; j < 3; ++j)
        if (f->neighbor(j)->info() < i &&</pre>
            t.segment(Delaunay::Edge(f,j)).squared_length() >= 4 * r) {
        stack.push_back(f->neighbor(j));
           f->neighbor(j)->info() = i;
    throw false;
catch (bool solvable) { std::cout << (solvable ? "y" : "n"); }</pre>
```