

# Algolab: quick ref

## Table of contents

<b>1 misc</b>	<b>1</b>
<b>2 STL</b>	<b>2</b>
2.1 array	2
2.2 vector, queue, stack	2
2.3 sort, pq	2
2.4 set, map	2
<b>3 CGAL</b>	<b>3</b>
3.1 CGAL basic	3
3.2 Triangulation	3
3.2.1 range search, circumcircles	4
3.3 LP QP	4
<b>4 BGL</b>	<b>5</b>
4.1 BGL basics	5
4.1.1 connected component	5
4.1.2 strong component <i>-bgl_tutorial_code.cpp</i>	5
4.1.3 biconnected components	5
4.1.4 MST	5
4.1.5 dijkstra <i>-bgl_tutorial_code.cpp</i>	6
4.1.6 topological sort	6
4.1.7 bfs/dfs	6
4.1.8 max (cardinality) matching	6
4.2 König's theorem	6
4.3 Max flow	7
4.3.1 flow modelling	7
4.3.2 circulation	7
4.3.3 edge lower bounds	7
4.3.4 mincut	8
4.3.5 edge-disjoint paths	8
4.3.6 vertex cover (bipartite matching)	8
4.3.7 mincost maxflow	9
<b>5 DP</b>	<b>9</b>

## 1 misc

- keyboard, capslock

```
setxkbmap us
```

```
setxkbmap -option ctrl:nocaps
```

geany keybinding:

complete word, delete cur line, toggle line comment, find next/prev selection

- basrc-cgal...

```
gencgal_cmake_eclipse(){#~/.bashrc
```

```
  cgal_create_cmake_script
```

```
  echo 'set(CMAKE_CXX_FLAGS "${CMAKE_CXX_FLAGS} -std=c++11")' >> CMakeLists.txt
```

```
  cmake -G "Eclipse CDT4 - Unix Makefiles" . }
```

- io config

```
ios_base::sync_with_stdio(false);
```

```
std::cout << std::setiosflags(std::ios::fixed);
```

```
std::cout << std::setprecision(0);
```

- rounding up/down

```
int round_up(const CGAL::Quotient<ET>& e) {
    double d = std::ceil(CGAL::to_double(e)); // #include <cmath>
    while(d < e) d++; while(d - 1 >= e) d--; return d; }
int round_down(const CGAL::Quotient<ET>& e) {
    double d = std::floor(CGAL::to_double(e));
    while(d > e) d--; while(d + 1 <= e) d++; return d; }
```

- def

```
#define forloop(i,lo,hi) for(int i = (lo); i <= (hi); ++i)
#define rep(i,N) forloop(i,0,(int)N-1)
```

## 2 STL

### 2.1 array

```
void f(int ** a){...} // function that takes 2d array
int **arr = new int *[n];
rep(i,n) arr[i]=new int[n]; f(arr);
rep(i,n) delete arr[i]; delete arr;
```

### 2.2 vector, queue, stack

```
#include <vector>/<queue>/<algorithm>
vector<int> v3(5,3); // v3 initialed as [3,3,3,3,3]
v.push_back(i);
for(vector<int>::iterator it=v1.begin(); it!=v1.end(); it++) cout << *it << " ";
sort(v1.begin(), v1.end()); // simple sort (ascending order)
reverse(v5.begin(), v5.end()); // reverse elements in v5
std::queue<int> q; vector<bool> visited(n,false);
q.push(s); visited[s]=true; // do a bfs
while(!q.empty()){const int u=q.top(); q.pop();
    for(v : out_vertices(u)) if(visited[v]==false)
        {visited[v]=true; q.push(v);} }
```

### 2.3 sort, pq

```
bool my_cmp(const pair<int, int> &lhs, const pair<int, int> &rhs)
    {return (lhs.first < rhs.first) || (lhs.first==rhs.first && lhs.second < rhs.second);}
sort(vp.begin(), vp.end(), my_cmp); // put the function name as 3rd argument
struct MyFooStruct { int x,y;
    MyFooStruct(int xx, int yy) {x = xx;y = yy;}
    bool operator < ( const MyFooStruct & other ) const
        {return (x<other.x) || (x==other.x && y<other.y);} };
vector<MyFooStruct> vs; sort(vs.begin(), vs.end());
priority_queue<MyFooStruct> spq;
rep(i,4) spq.push( MyFooStruct(a1[i], a2[i]) );
cout << spq.top().x << " " << spq.top().y << endl; spq.pop();
```

### 2.4 set, map

```
set<int> s; rep(i, 10) s.insert(i); s.size();
for(set<int>::iterator it=s.begin(); it!=s.end(); it++) cout << *it << " ";
cout << (s.find(10)!=s.end()) << " " << (s.count(9)==1) << endl;
s.erase(100); s.erase(s.begin()); s.erase(--s.end());
set<int> s1,s2, s_union, s_intersect, s_diff; // #include <algorithm>
set_union( s1.begin(), s1.end(), s2.begin(), s2.end(), inserter(s_union, s_union.end()) );
set_intersection( s1.begin(),s1.end(),s2.begin(),s2.end(),inserter(s_intersect,s_intersect.begin()) );
set_difference( s1.begin(), s1.end(), s2.begin(), s2.end(), inserter(s_diff,s_diff.end()) );
map<string,int> m;
rep(i, 3) m.insert( make_pair(wds[i], cnts[i]) );// insert by making <k,v> pair
m["aa"] = 3; // update/insert by assignment
for(map<string,int>::iterator it=m.begin(); it!=m.end(); it++)
```

```
cout << it->first << ":" << it->second << ", ";
```

### 3 CGAL

#### 3.1 CGAL basic

*intersect.cpp*, *hello-really-exact.cpp*, *two-kernels.cpp*, *minball.cpp*

- functions

```
typedef CGAL::Exact_predicates_exact_constructions_kernel K; // or other kernels
K::FT CGAL::squared_distance (Type1<K> obj1, Type2<K> obj2){} //distance^2 between 2 geometric obj
std::sqrt(CGAL::to_double(CGAL::squared_distance(r,l)))/ //obtain an approximation of the real dist
bool CGAL::left_turn (Point_2 &p, Point_2 &q, Point_2 &r){}
int CGAL::orientation (const Point_2 &p1, const Point_2 &p2, const Point_2 &p3){}
auto CGAL::intersection (Type1< Kernel > obj1, Type2< Kernel > obj2){}
```

- intersection –*intersect.cpp*
- min circle –*minball.cpp*

```
typedef CGAL::Exact_predicates_exact_constructions_kernel_with_sqrt K;
typedef CGAL::Min_circle_2_traits_2<K> Traits;
typedef CGAL::Min_circle_2<Traits> Min_circle;
typedef K::Point_2 P;
P points[n]; // almost-antenna pb
Min_circle mc1( points, points+n, false); // very slow
Min_circle mc(points, points + n, true);
Traits::Circle c = mc.circle();
K::FT min_r = c.squared_radius();
if (mc.is_degenerate() == false)
    for (int i = 0; i < mc.number_of_support_points(); i++) P sp = mc.support_point(i);
```

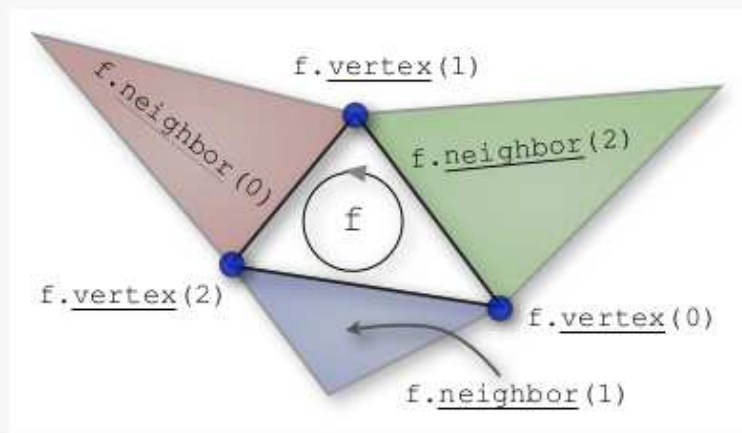
#### 3.2 Triangulation

- find nearest vertex
- iterate through all edges
- for a vertex, iterate through its incident edges
- do BFS/DFS on faces
- adding additional information

Proximity.pdf [delanay.cpp](#)  $O(n \log n)$

- in CGAL documentation, see *Triangulation\_2* and *Delaunay\_triangulation\_2* for member functions of a triangulation.
- see *TriangulationDataStructure\_2* for info wrt. Edge/Face/Vertex
- 

*CGAL's triangulation data structure is vertex/face based.  
Edges are represented implicitly only.*



Edges in `CGAL::Triangulation_data_structure_2` are represented as a `std::pair<Face_handle,int>`: A pair  $(f,i)$  represents the  $i$ -th edge along the boundary of  $*f$ . The edge connects the *vertices*  $(i+1)\%3$  and  $(i+2)\%3$  of  $*f$ .

```
//A safe strategy is to let the triangulation choose a suitable insertion order: Instead of
inserting points one by one using t.insert(p) insert a whole (iterator) range [b,e) of points using
t.insert(b,e)
rep(i,n) {
    K::Point_2 p; std::cin >> p;
    pts.push_back(p);}
// construct triangulation
Triangulation t;
t.insert(pts.begin(), pts.end());
for (Edge_iterator e = t.finite_edges_begin(); e != t.finite_edges_end(); ++e){
    //~ if(t.segment(e).vertex(1)!=pts[0] && t.segment(e).vertex(0)!=pts[0]) continue;
    best = min(best, t.segment(e).squared_length() );
}
```

### 3.2.1 range search, circumcircles

cf. *radiation2*

## 3.3 LP QP

A model of `QuadraticProgram` describes a convex quadratic program of the form.

$$\begin{aligned} \text{(QP) minimize } & \mathbf{x}^T D \mathbf{x} + \mathbf{c}^T \mathbf{x} + c_0 \\ \text{subject to } & A \mathbf{x} \begin{matrix} \geq \\ \leq \end{matrix} \mathbf{b}, \\ & \mathbf{l} \leq \mathbf{x} \leq \mathbf{u} \end{aligned}$$

in  $n$  real variables  $\mathbf{x} = (x_0, \dots, x_{n-1})$ .

Here,

- $A$  is an  $m \times n$  matrix (the constraint matrix),
- $\mathbf{b}$  is an  $m$ -dimensional vector (the right-hand side),
- $\begin{matrix} \geq \\ \leq \end{matrix}$  is an  $m$ -dimensional vector of relations from  $\{\leq, =, \geq\}$ ,
- $\mathbf{l}$  is an  $n$ -dimensional vector of lower bounds for  $\mathbf{x}$ , where  $l_j \in \mathbb{R} \cup \{-\infty\}$  for all  $j$
- $\mathbf{u}$  is an  $n$ -dimensional vector of upper bounds for  $\mathbf{x}$ , where  $u_j \in \mathbb{R} \cup \{\infty\}$  for all  $j$
- $D$  is a symmetric positive-semidefinite  $n \times n$  matrix (the quadratic objective function),
- $\mathbf{c}$  is an  $n$ -dimensional vector (the linear objective function), and
- $c_0$  is a constant.

portfolio.cpp LP\_QP.pdf

```
// general QP: ** min x'Dx + c'x + c0 st. Ax<=b, l<=x<=u **
Program qp (CGAL::LARGER, false, 0, false, 0);
qp.set_u(X,true); qp.set_u(Y,true); // x,y <= 0
qp.set_l(Z2,true); // z^2>=0
qp.set_a(X,0,1);qp.set_a(Y,0,1);qp.set_b(0,-4); // constraint-0: x + y >= -4
qp.set_a(X,1,4);qp.set_a(Y,1,2);qp.set_a(Z2,1,1);qp.set_b(1,-a*b); // constraint-1: 4x + 2y + z2 >= -ab
qp.set_a(X,2,-1);qp.set_a(Y,2,1);qp.set_b(2,-1); // constraint-2: -x + y >= -1
// obj: min ax^2 + by + z^4
qp.set_d(Z2,Z2,2); // z^4
qp.set_d(X,X,a*2); // a*x^2 -- need to *2!
qp.set_c(Y,b); // b*y --need to *2!
Solution s = CGAL::solve_quadratic_program(qp, ET());
assert (s.solves_quadratic_program(qp));
if (s.status() == CGAL::QP_INFEASIBLE) {cout << "no" << endl;continue;}
```

```

else if (s.status() == CGAL::QP_UNBOUNDED) {cout<<"unbounded"<<endl; continue;}
double obj = CGAL::to_double(s.objective_value());
cout.precision(0); cout << fixed << ceil(obj) << endl;

```

## 4 BGL

### 4.1 BGL basics

*bgl\_tutorial\_code.cpp, bgl\_handout\_1.pdf*

#### 4.1.1 connected component

```

int V = num_vertices(G); vector<int> comp(V); // stores index of the vertices' component
int ncomp = connected_components(G, &comp[0]);
cout << ncomp << " connected components in G. " << endl;
multimap<int, int> mm; // map scc id to vertices
rep(i, V) mm.insert( make_pair(comp[i], i) );
rep(i, ncomp){
    cout << "component-" << i << " have vertices: "; multimap<int, int>::iterator ibeg, iend;
    for( tie(ibeg,iend) = mm.equal_range(i); ibeg!=iend; ibeg++) cout << ibeg->second << ", "; }

```

#### 4.1.2 strong component –*bgl\_tutorial\_code.cpp*

cf. *bgl-tutorial monkey island*

```

vector<int> scc(V); // `scc` stores vertices' strong component id (ie. one PARTITION of all vertices)
int nscc = strong_components(G,
    make_iterator_property_map(scc.begin(),get(vertex_index, G)) ); // nscc = nb of scc in G

```

#### 4.1.3 biconnected components

```

WeightMap wm = get(edge_weight, G); // biconnected component is a partition of *edges*, so we need a
property_map (for edges) to store the component id for each edge
/* we just use the WeightMap (which is just a property_map for edges) to store the components
* else we can use the edge_component_t as above... ( need to redefine the 'Graph' type, see:
http://www.boost.org/doc/libs/1_59_0/libs/graph/example/biconnected_components.cpp )*/
int ncomp = biconnected_components(G, wm); // put wm as argument to store edge's component
cout << ncomp << " biconnected components" << endl;
EdgeIt ei, ei_end;
for (tie(ei, ei_end) = edges(G); ei != ei_end; ++ei)
    cout << source(*ei, G) << "-" << target(*ei, G) << ", is in biconnected-component-" <<wm[*ei];
vector<Vertex> art_pts; // get articulation points
articulation_points(G, back_inserter(art_pts));
cout << art_pts.size() << " articulation points, they are: " << endl;
for(vector<Vertex>::iterator it=art_pts.begin(); it!=art_pts.end(); it++) cout << *it << ", ";

```

#### 4.1.4 MST

- Kruskal

```

vector<Edge> mst; // edge vector to store mst: a list of V-1 edges
kruskal_minimum_spanning_tree(G, back_inserter(mst));
int mst_weight = 0;
for (vector<Edge>::iterator ebegin = mst.begin(); ebegin != mst.end(); ++ebegin) {
    int u = source(*ebegin, G), v = target(*ebegin, G); int c = wm[*ebegin];
    cout << u<<'-'<<v<<" , weight="<<c<<endl; mst_weight += c;
} cout << "MST total weight is: " << mst_weight << endl;

```

- Prim

```

int V = num_vertices(G);
vector<int> pred(V); // predecessor vector
prim_minimum_spanning_tree(G, &pred[0]);
rep(j, V){
    Edge e; bool success;
    tie(e, success) = edge(j, pred[j], G);

```

```

    if(success){// because the mst root do not have pred...
        int u = source(e, G), v = target(e, G); int c = wm[e];
        cout << u<<'-'<<v<<"", weight="<<c<<endl; }
    }
}

```

#### 4.1.5 dijkstra -bgl\_tutorial\_code.cpp

```

int V = num_vertices(G);
vector<int> dist(V), pred(V);
dijkstra_shortest_paths(
    G, 0, // dijkstra with source=0
    predecessor_map( make_iterator_property_map(pred.begin(), get(vertex_index, G)) ),
    distance_map( make_iterator_property_map(dist.begin(), get(vertex_index, G)) )
);
rep(i,V) cout << "0-" << i << ", dist = " << dist[i] << endl;

```

#### 4.1.6 topological sort

```

vector<Vertex> c; // container
topological_sort(G, back_inserter(c));
cout << "A topological ordering: ";
for ( vector<Vertex>::reverse_iterator ii=c.rbegin(); ii!=c.rend(); ++ii) cout << (*ii) << " ";

```

#### 4.1.7 bfs/dfs

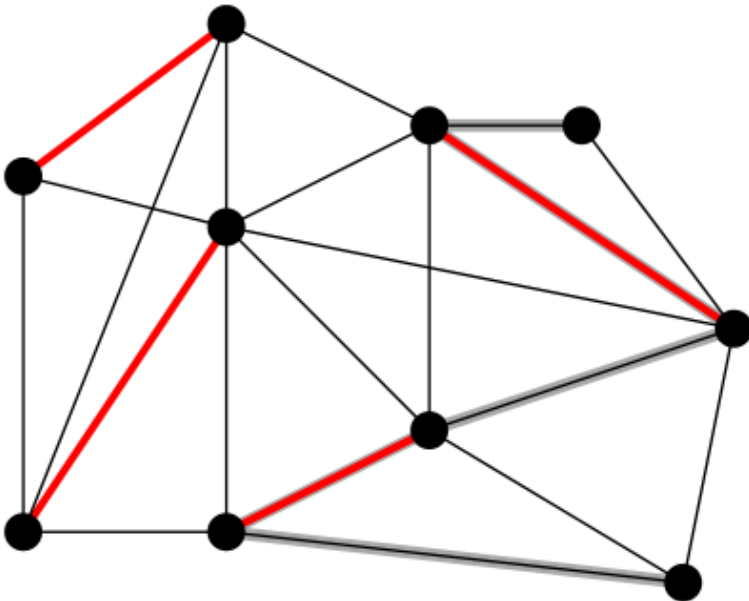
```

class custom_bfs_visitor : public boost::default_bfs_visitor
{ public:
    template < typename Vertex, typename Graph >
    void discover_vertex(Vertex u, const Graph & g) const { cout << u << ", "; }
};
custom_bfs_visitor vis;
breadth_first_search(G, vertex(0, G), visitor(vis)); // almost the same for DFS

```

#### 4.1.8 max (cardinality) matching

cf. *buddy selection*



- $G = (V, E)$
- $M \subseteq E$  is a matching if and only if no two edges of  $M$  are adjacent.
- In an unweighted graph, a maximum matching is a matching of maximum cardinality.
- In a weighted graph, a maximum matching is a matching such that the weight sum over the included edges is maximum.
- BGL does not provide weighted matching algorithms.

cf. *knight's*

#### 4.2 König's theorem

A graph is bipartite if and only if it does not contain an odd cycle.

In bipartite graphs, the size of [minimum vertex cover](#) is equal to the size of the [maximum matching](#); this is König's theorem.<sup>[16][17]</sup>

An alternative and equivalent form of this theorem is that the size of the [maximum independent set](#) plus the size of the maximum matching is equal to the number of vertices.

### 4.3 Max flow

[bgl\\_flows.cpp](#) [residualBFS.cpp](#) [mincost\\_maxflow.cpp](#)

#### 4.3.1 flow modelling

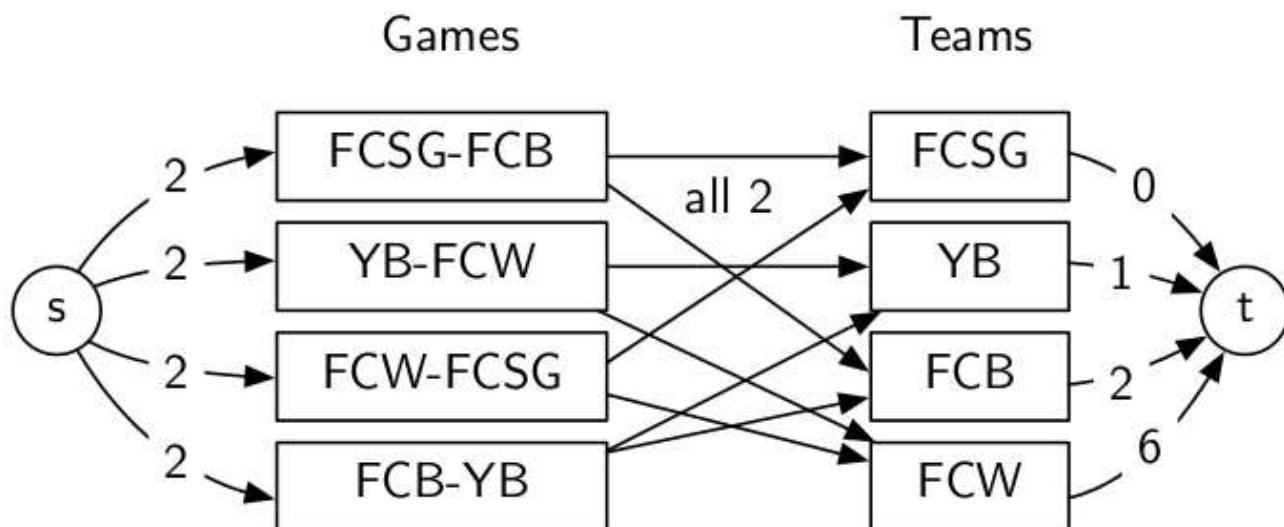


Figure 1.

#### 4.3.2 circulation

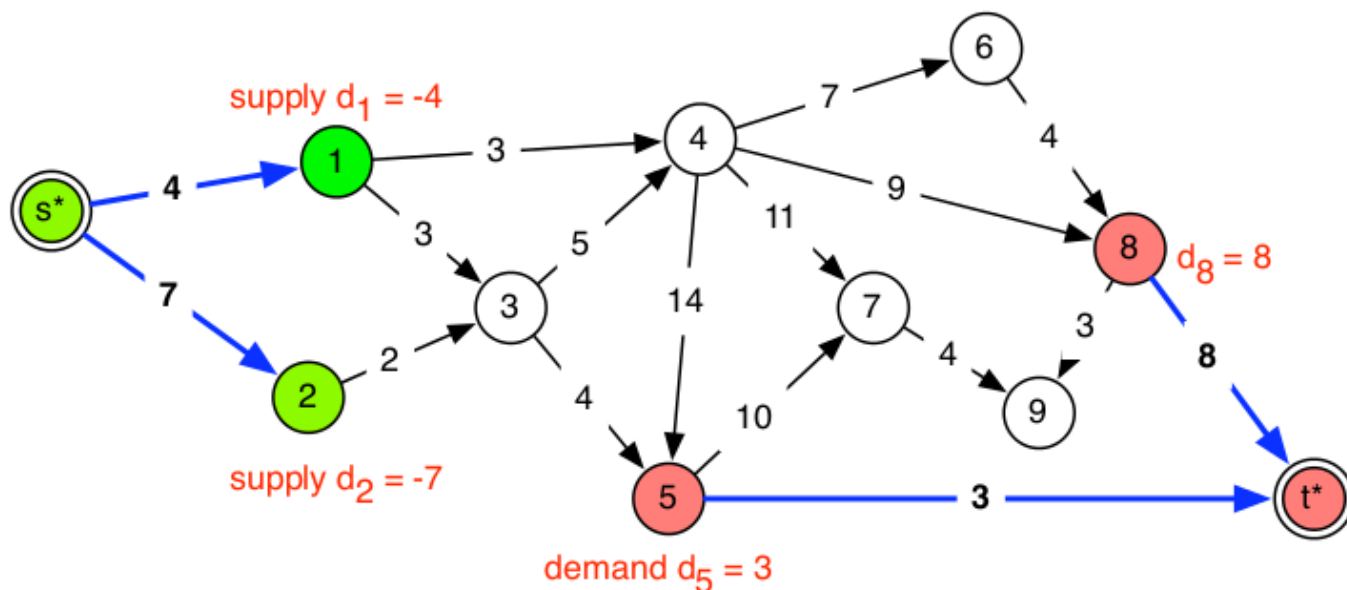


Figure 2.

#### 4.3.3 edge lower bounds

first let lower bound flow, adjust supply/demand of each vertex  $\rightarrow$  to a circulation pb.



### New demand constraints:

$$f^{\text{in}}(v) - f^{\text{out}}(v) = d_v - L_v$$

Also,  $f_0$  uses some of the edge capacities already, so we have:

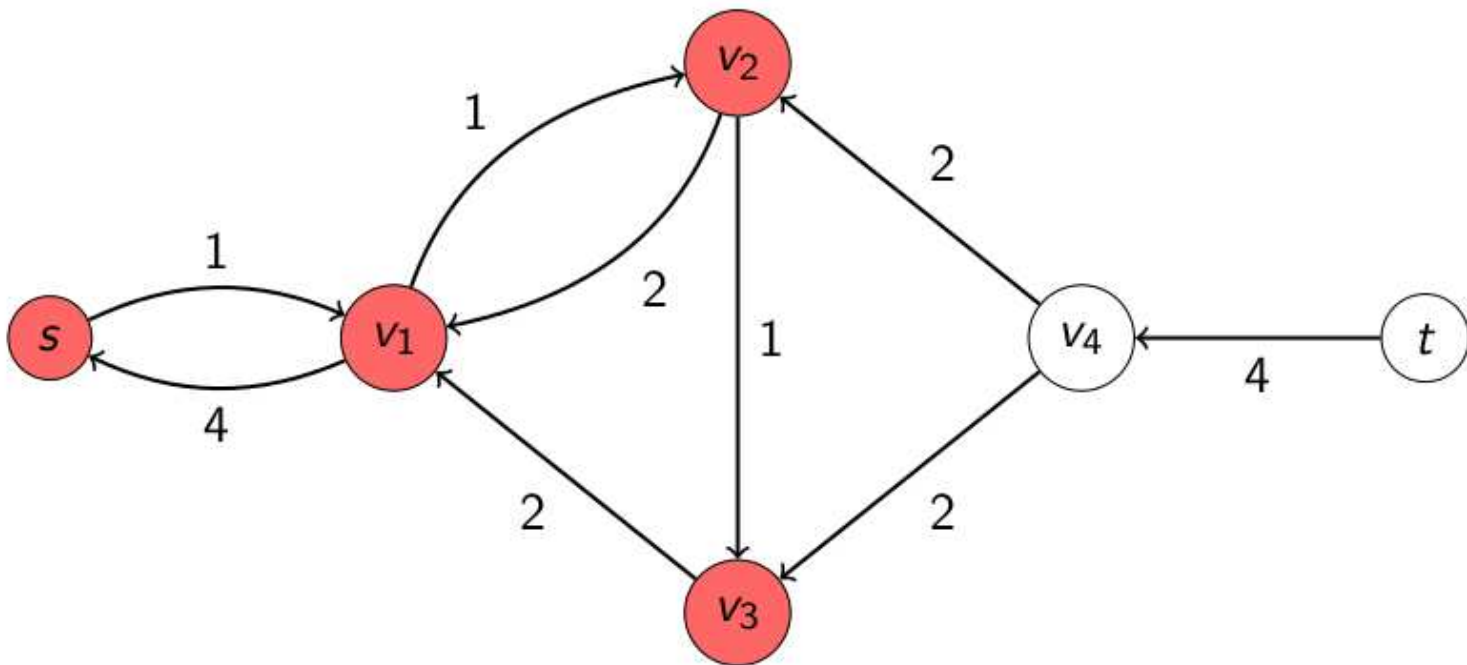
### New capacity constraints:

$$0 \leq f(e) \leq c_e - \ell_e$$

These constraints give a standard instance of the circulation problem.

#### 4.3.4 mincut

$V = S \cup T$ ,  $S$ =reachable vertices from  $s$  in residual graph. *residualBFS.cpp*



#### 4.3.5 edge-disjoint paths

setting each edge of cap=1, maxflow from  $s$  to  $t$  = nb of edge-disjoint paths from  $s$  to  $t$ .

cf. *Phantom menace*

cf. *tetris*

#### 4.3.6 vertex cover (bipartite matching)

*residualBFS.cpp*, *bgl\_tutorial\_3\_moreflows.pdf*



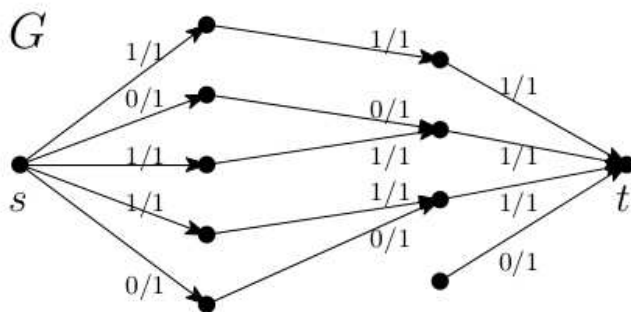
## ■ Maximum independent set

Largest  $T \subseteq V$ , such that  
 $\nexists u, v \in T : (u, v) \in E$ .

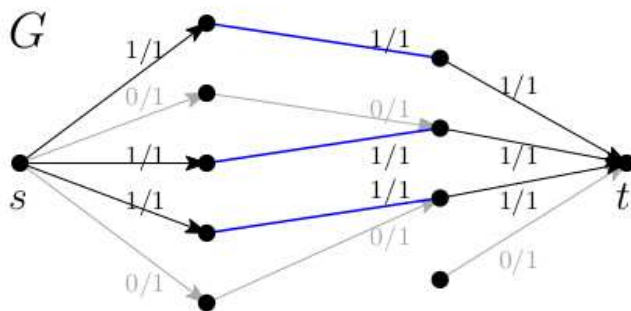
## ■ Minimum vertex cover

Smallest  $S \subseteq V$ , such that  
 $\forall (u, v) \in E : u \in S \vee v \in S$ .

Compute the flow:

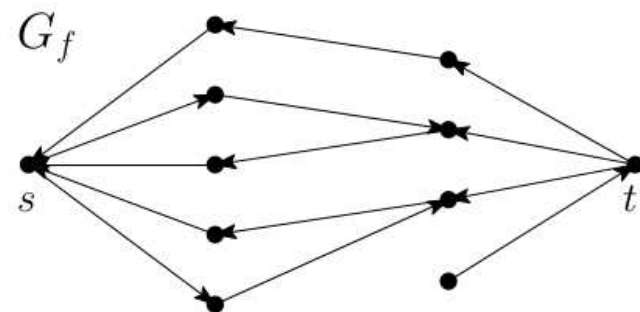


Find the matching:



⇒ result is the unvisited vertices in L, and visited vertices in R.

Compute the residual graph  $G_f$ :



Find visited vertices with BFS from

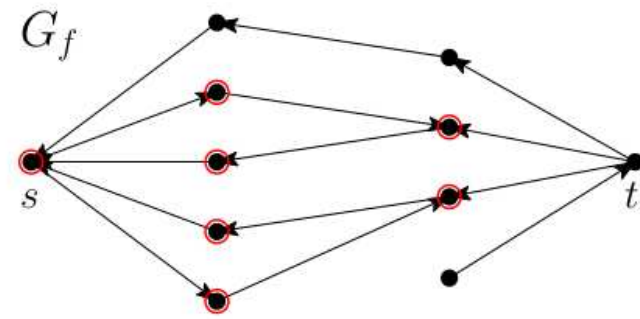


Figure 3.

### 4.3.7 mincost maxflow

canceling negative weights → cf *canteen*, *carsharing*

## 5 DP

compressing state!

cf. *bonus level*, *poker chips*