CHAROTAR UNIVERSITY OF SCIENCE & TECHNOLOGY

Sixth Semester of B. Tech (CE) Examination September 2018

CE306/CE306.01 Theory of Computation (TOC)

Date: 22.09.2018, Saturday Time: 10.00 a.m. To 01.00 p.m. Maximum Marks: 70

Instructions:

- 1. The symbols used carry their usual meanings.
- 2. Section I and II must be attempted in separate answer sheets.
- 3. Make suitable assumptions and draw neat figures wherever required.
- 4. Rough work is to be done in the last page of main supplementary, please don't write anything on the question paper.
- 5. Indicate clearly, the option(s) you attempt along with its respective question number.
- 6. Figures to the right indicate marks.

SECTION - I

O - 1 Do as directed:

a. State True/False with reason: Deterministic Finite Automata can have more than one [02] accepting state.

Consider the Finite Automata M given below & state T/F for the following [02]

b. statements with reasons:



- a. M accepts a Null String
- b. M accepts all strings over {a, b}
- c. M is a deterministic FA
- d. M is an NFA
- c. Define: Equivalence Relation [02]
- d. What are the differences between Moore Machine & Mealy Machine? [02]
- e. Convert the following regular expression to deterministic finite automata. [03]

 $(a^*|b^*)^*$

Q-2 Answer the following questions (Any Three).

[12]

- **a.** Use mathematical induction to prove that 1 + 2 + 3 + ... + n = n (n + 1) / 2 for all positive integers n.
- **b.** State and prove Kleene's theorem part I.
- **c.** Show and explain in detail the procedure to convert NFA -^ to NFA with example.
- **d.** State the pumping lemma for regular language. Prove that $\{0^n1^n \mid n \ge 0\}$ is not a regular language.

Q - 3 Answer the following questions (Any Three).

[12]

- **a.** Write and explain an algorithm to minimize the given DFA in detail with an appropriate example.
- **b.** What is a proof? Prove by Contradiction: For any set A, B and C, if $A \cap B = \emptyset$ and C is a subset of B then $A \cap C = \emptyset$.
- **c.** An NFA with states 1-5 and input alphabet {a, b} has the following transition table.

Q	δ (q, a)	δ (q, b)
1	{1, 2}	{1}
2	{3}	{3}
3	{4}	{4}
4	{5}	Ø
5	Ø	{5}

Find:

- a. Draw Transition Diagram
- b. Calculate $\delta^*(1, ab)$
- **d.** Obtain a Deterministic Finite Automata to accept strings of a's and b's having even number of a's and b's.

SECTION - II

Q-4 Do as directed:

a. Compare Turing Machine with FA and PDA.

- [03]
- **b.** "Non-Determinism adds power to a NFA". Justify the correctness or falsehood of the statement.
- [02]

- **c.** Remove useless symbol from the given CFG:
 - $S \rightarrow AB|a \qquad A \rightarrow BC|b \qquad B \rightarrow aB|C \qquad C \rightarrow aC|B$
- **d.** Give applications of finite automata.

[03]

Q-5 Answer the following questions (Any Three).

[12]

- **a.** Find a regular expression for the set of all strings over $\{a, b\}$.
 - i) The language of all strings containing at least two a's
 - ii) The language of all strings containing at most two b's
- **b.** Design a Turing Machine accepting language of even length palindrome over $\{0, 1\}$.
- **c.** When a grammar is said to be an ambiguous grammar? Check whether the grammar G with production rules

$$X \rightarrow X+X \mid X*X \mid X \mid a$$

is ambiguous or not.

d. Define: "Context Free Grammar". What languages are generated by Context Free Grammar whose production rules are given below:

$$S \rightarrow aSa$$
 , $S \rightarrow bSb$,

$$S \rightarrow aa, S \rightarrow bb$$

Q-6 Answer the following questions (Any Three).

[12]

- **a.** Define: Turing Machine. Mention Input device, Output device and Storage device in Turing Machine. Also design a Turing Machine to accept language a⁺.
- **b.** Show using pumping lemma that the given language is not a CFL. $L = \{a^n b^n c^n \mid n \ge 1\}$
- **c.** Write down a CFG generating the given language.
 - (1) The set of odd-length strings in $\{a, b\}^*$ with middle symbol a
 - (2) The set of even-length strings in $\{a, b\}^*$ with two middle symbols equal
- **d.** Convert the following context free grammar into the Chomsky normal form.

$$S \rightarrow aXbX$$
 $X \rightarrow aY \mid bY \mid^{\wedge} Y \rightarrow X \mid c$