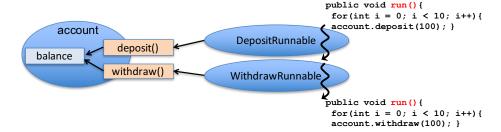
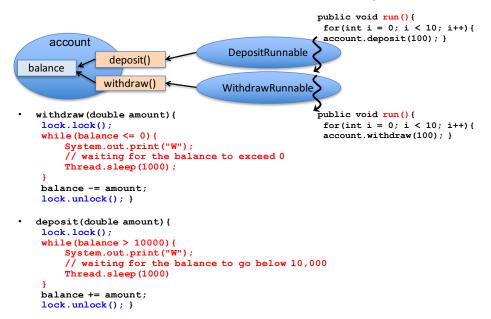
DeadlockedBankAccount.java



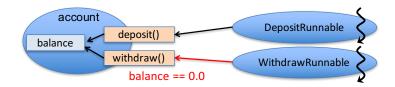
Deadlock

DeadlockedBankAccount.java

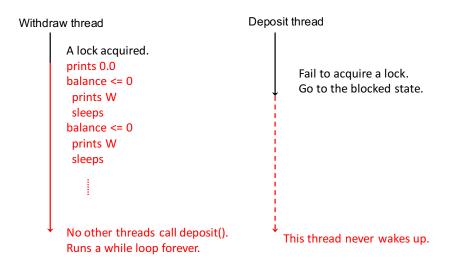


Deadlock

- Assume the withdrawal thread goes ahead and runs first.
- Output
 - Lock obtainedCurrent balance (w): 0.0WWWWWWWW



• Current balance (w): 0.0WWWWWWW



- Has no deadlock problems.
- Can generate race conditions.

DeadlockedBankAccount2.java

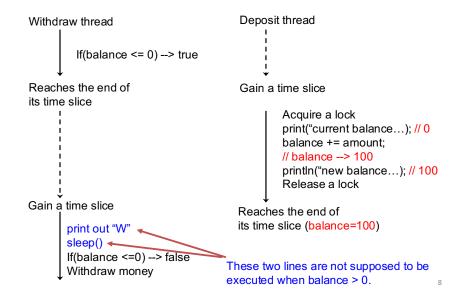
```
· Previous version
    - withdraw(double amount) {
       lock.lock();
       while (balance <= 0) {
           System.out.print("W");
           // waiting for the
           // balance to exceed 0
           Thread.sleep(1000);
       balance -= amount;
       lock.unlock();
   - deposit(double amount) {
       lock.lock();
       while(balance > 10000){
           System.out.print("W");
           // waiting for the balance
      to go below 10,000
           Thread.sleep (1000)
       balance += amount;
       lock.unlock(); }
```

· New version

```
- withdraw(double amount) {
    while( balance <= 0 ) {
        System.out.print("W");
        Thread.sleep(2);
    }
    lock.lock();
    balance -= amount;
    lock.unlock();
}

- deposit(double amount) {
    while( balance > 10000 ) {
        System.out.print("W");
        Thread.sleep(2);
    }
    lock.lock();
    balance += amount;
    lock.unlock();
}
```

A Potential Race Condition in DeadlockedBankAccount2



Avoiding Deadlocks and Race Conditions

- Use a Condition object(s).
 - Allow a thread to
 - temporarily release a lock so that another thread can proceed
 The thread goes to the Waiting state from the Runnable state.
 - re-acquire the lock later.
- java.util.concurrent.locks.Condition
 - Obtain its instance from a lock object

```
• ReentrantLock lock = new ReentrantLock();
Condition condition = lock.newCondition(); //factory method
condition.await();
```

Timed Waiting sleep() Waiting notify() join() notifyAll() wait() notify() sianalAll() notifyAll() await() interruption sleep(signalAll() join() interruption wait() await() Exits run() or **Explicit thread** start() termination new Runnable **Terminated** New I/O op completion I/O operation or wait for or thread sync thread sync (lock) done **Blocked**

ThreadSafeBankAccount2

```
Condition sufficientFundsCondition = lock.newCondition();
Condition belowUpperLimitFundsCondition = lock.newCondition();
withdraw(double amount) {
 lock.lock();
 while (balance <= 0) {
     // waiting for the balance to exceed 0
    sufficientFundsCondition.await(); }
 balance -= amount;
 belowUpperLimitFundsCondition.signalAll();
 lock.unlock(); }
deposit(double amount) {
 lock.lock();
 while (balance >= 300) {
    // waiting for the balance to go below 10000.
    belowUpperLimitFundsCondition.await(); }
 balance += amount;
 sufficientFundsCondition.signalAll();
 lock.unlock(); }
```

ThreadSafeBankAccount2

```
0utput
    Lock obtained
 - 7 (d): current balance: 0.0
 - 7 (d): new balance: 100.0
 - Lock released

    Lock obtained

    8 (d): current balance: 100.0
    8 (d): new balance: 200.0
    Lock released
    Lock obtained
    9 (d): current balance: 200.0
    9 (d): new balance: 300.0

    Lock released

    10 (d): current balance: 300.0
 - 10 (d): await(): Balance has reached the upper limit.

    Lock obtained

    11 (d): current balance: 300.0
    11 (d): await(): Balance has reached the upper limit.

    Lock obtained

 - 12 (w): current balance: 300.0
 - 12 (w): new balance: 200.0

    Lock released

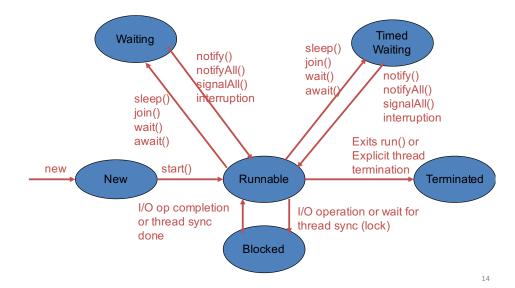
 - 10 (d): new balance: 300.0
 - Lock released
 - 11 (d): await(): Balance has reached the upper limit.
```

11

Condition

- await()
 - Will be waiting until it is signaled or interrupted
 - Will be waiting until it is signaled or interrupted, or until a specified waiting time (relative time) elapsed.
 - Will be waiting until it is signaled or interrupted, or until a specified deadline (absolute time).
 - If signaled, goes to the Runnable state and re-acquires a lock.
 - Will be "blocked" if fails to re-acquire the lock.
 - Throws an InterruptedException if interrupted.
 - c.f. previous lecture note that explains InterruptedException
- signalAll()
 - Wakes up all waiting threads on a condition object.
 - All of them go to the "runnable" state.
 - One of them will re-acquire a lock.

- When a thread calls await(), signal() or signalAll() on a Condition object,
 - the thread is assumed to hold the lock associated with the Condition object.
- If the thread does not,
 - an IllegalMonitorStateException is thrown.



SignalAll() Before or After a State Change?

```
withdraw(double amount) {
 lock.lock();
 while (balance =< 0) {
     // waiting for the balance to exceed 0
    sufficientFundsCondition.await(); }
 balance -= amount;
 belowUpperLimitFundsCondition.signalAll();
 lock.unlock(); }
deposit(double amount) {
 lock.lock();
 while(balance >= 300) {
    // waiting for the balance to go below 300.
    belowUpperLimitFundsCondition.await(); }
 balance += amount;
 sufficientFundsCondition.signalAll();
 lock.unlock(); }
```

 What happens if you call signalAll() first and then update the balance? Any problems?

```
withdraw(double amount) {
 lock.lock();
 while (balance =< 0) {
     // waiting for the balance to exceed 0
    sufficientFundsCondition.await(); }
belowUpperLimitFundsCondition.signalAll();
balance -= amount;
 lock.unlock(); }
deposit(double amount) {
 lock.lock();
 while (balance >= 300) {
    // waiting for the balance to go below 300.
    belowUpperLimitFundsCondition.await(); }
 sufficientFundsCondition.signalAll();
balance += amount;
 lock.unlock(); }
```

Need to worry about race conditions in this case?

```
(1) W thread:
             vithdraw(double amount) {
              lock.lock();
"waiting"
              hile(balance =< 0){
                  // waiting for the balance to exceed _{0} (3) W thread:
temporarily
                                                            "runnable"
releases the
                 sufficientFundsCondition.await(); }
             belowUpperLimitFundsCondition.signalAll();
                                                           Tries to acquire the
lock
             balance -= amount;
                                                            lock again and fails.
             lock.unlock(); }
                                                            Goes to "blocked."
            deposit(double amount) {
             lock.lock();
             while(balance >= 300) {
                 // waiting for the balance to go below 300.
(2) D thread:
                 belowUpperLimitFundsCondition.await(); }
signalAll().
             sufficientFundsCondition.signalAll();
Ctx switch
             balance += amount;
             lock.unlock(); }
```

- A "W" thread CANNOT withdraw money before a "D" thread deposits money.
- A "D" thread CANNOT deposit money before a "W" thread withdraws money.

```
vithdraw(double amount) {
             lock.lock();
"waiting"
              hile(balance =< 0){
temporarily
                  // waiting for the balance to exceed 0
releases the
                 sufficientFundsCondition.await(); }
             belowUpperLimitFundsCondition.signalAll();
lock
             balance -= amount;
             lock.unlock(); }
            deposit(double amount) {
             lock.lock();
             while(balance >= 300) {
                 // waiting for the balance to go below 300.
(2) D thread:
                belowUpperLimitFundsCondition.await(); }
signalAll().
             sufficientFundsCondition.signalAll();
Ctx switch
            balance += amount;
             lock.unlock(); }
```

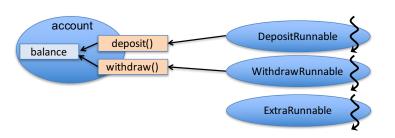
- A context switch can occur in between signalAll() and state change?
- A "W" thread can withdraw money before a "D" thread deposits money?
 - Can the balance variable have a negative value?

Two Important Things (1)

- A state change (or value change) can be made on a shared variable safely after calling signalAll()
 - AS FAR AS the state changes in atomic code
- Common programming convention/practice:
 - A state change first, followed by signalAll().

Two Important Things (2)

- A JVM <u>DOES</u> context switches even when a thread runs in atomic code.
 - A lock guarantees that only one thread exclusively runs atomic code at a time.



signal() and signalAll()

- signalAll()
 - Wakes up all waiting threads on a condition object.
 - All of them go to the "runnable" state.
 - One of them will re-acquire a lock.
- signal()
 - Wakes up one of waiting threads on a condition object.
 - One of them goes to the "runnable" state. The others stay at the "waiting" state.
 - JVM's thread scheduler selects one of them. Assume a random selection.
 - Not predictable which waiting thread to be selected.

Note that...

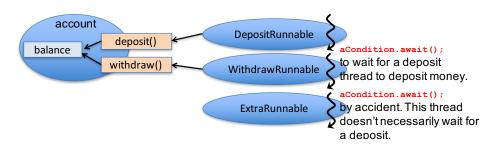
- Some books and online materials explicitly/implicitly say that context switches never occur when a thread runs in atomic code.
- It is wrong!

ThreadSafeBankAccount2

```
Condition sufficientFundsCondition = lock.newCondition();
Condition belowUpperLimitFundsCondition = lock.newCondition();
withdraw(double amount) {
 lock.lock();
 while (balance =< 0) {
     // waiting for the balance to exceed 0
    sufficientFundsCondition.await(); }
 balance -= amount;
 belowUpperLimitFundsCondition.signalAll();
 lock.unlock(); }
deposit(double amount) {
 lock.lock();
 while (balance >= 300) {
    // waiting for the balance to go below 300.
    belowUpperLimitFundsCondition.await(); }
 balance += amount;
 sufficientFundsCondition.signalAll();
 lock.unlock(); }
```

signal() or signalAll()?

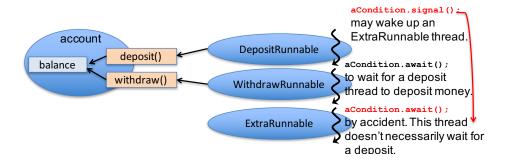
- signal() is more lightweight than signalAll().
 - Waking up a waiting thread is computationally expensive.
 - signal() wakes up only one thread.
- signalAll() is more protective.
 - signalAll() should be favored (at least in my personal taste).



Timed Waiting sleep() Waiting notify() join() notifyAll() notify() wait() signalAll() notifyAll() await(sleep(interruption signalAll() join() interruption wait() await() Exits run() or Explicit thread start() new termination Runnable New **Terminated** I/O op completion I/O operation or wait for or thread sync thread sync (lock) done Blocked

signal()

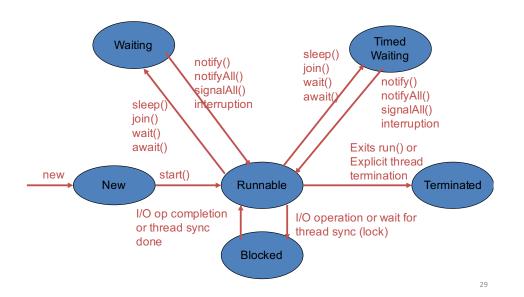
- A deposit thread may wake up an ExtraRunnable thread by calling signal().
 - A withdraw thread
 - loses a chance to withdraw money even if some money in the account.
 - may never be waked up if no threads call deposit afterward.



signalAll()

- A deposit thread wake up both a withdraw thread and an ExtraRunnalbe thread.
 - Both threads go into the "Runnable" state.
- If the ExtraRunnable thread acquires the lock associated with acondition,
 - the withdraw thread
 - · goes to the "blocked" state.
 - acquires the target lock when the ExtraRunnable releases it.
 - will wake up sooner or later.
- aCondition.signalAl never lose a chance to withdraw money. may wake up an ExtraRunnable thread. account DepositRunnable deposit() aCondition.await(); balance to wait for a deposit withdraw() WithdrawRunnable thread to deposit money. aCondition.await(); ExtraRunnable by accident. This thread doesn't necessarily wait for a deposit.

ThreadSafeBankAccount2



```
Condition sufficientFundsCondition = lock.newCondition();
  Condition belowUpperLimitFundsCondition = lock.newCondition();

    withdraw(double amount) {

   lock.lock();
   while (balance =< 0) {
       // waiting for the balance to exceed 0
       sufficientFundsCondition.await(); }
   balance -= amount;
   belowUpperLimitFundsCondition.signalAll();
   lock.unlock(); }
  deposit(double amount) {
   lock.lock();
   while (balance >= 300) {
      // waiting for the balance to go below 300.
      belowUpperLimitFundsCondition.await(); }
   balance += amount;
   sufficientFundsCondition.signalAll();
   lock.unlock(); }
```

"while" or "if" to Surround await()?

```
withdraw(double amount) {
 lock.lock();
 while (balance =< 0) {
     // waiting for the balance to exceed 0
    sufficientFundsCondition.await(); }
 balance -= amount;
 belowUpperLimitFundsCondition.signalAll();
 lock.unlock(); }
withdraw(double amount) {
lock.lock();
 if(balance =< 0){
     // waiting for the balance to exceed 0
    sufficientFundsCondition.await(); }
 balance -= amount;
belowUpperLimitFundsCondition.signalAll();
 lock.unlock(); }
```

 "while" should be used rather than "if" when multiple threads call withdraw() concurrently. Why?

Problem

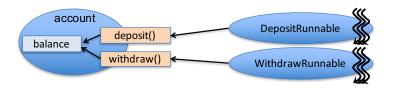
```
"runnable" One of them
            withdraw(double amount){
                                                       acquires the lock again
(1) b==0.
             lock.lock();
                                                        and releases it.
Two W
             if(balance =< 0){
                  // waiting for the balance to excee b==0.
threads:
                 sufficientFundsCondition.await(); } The other W thread:
"waiting"
             balance -= amount;
                                                       "blocked" on acquiring
             belowUpperLimitFundsCondition.signalAl the lock.
             lock.unlock(); }
            deposit(double amount) {
             lock.lock();
             if(balance >= 300){
                 // waiting for the balance to go below 300.
(2) D thread:
                 belowUpperLimitFundsCondition.await(); }
             balance += amount;
signalAll()
             sufficientFundsCondition.signalAll();
followed by .
             lock.unlock(); }
unlock()
b = = 100
```

30

(3) Two W threads:

```
(3) Two W threads:
                                                         "runnable" One of them
                                                         acquires the lock again
                                                         and releases it.
                                                         b==0.
              ithdraw(double amount){
                                                         The other W thread:
              Lock.lock();
Two W
                                                         "blocked" on acquiring
              f(balance = < 0) {
                  // waiting for the balance to exceed he lock.
threads:
                 sufficientFundsCondition.await()
"waiting"
             balance -= amount;
             belowUpperLimitFundsCondition.signalAl
                                                          (4) The 2nd W thread
             lock.unlock(); }
                                                          acquires the lock and
            deposit(double amount) {
                                                          releases it.
             lock.lock();
                                                          b = -100.
             if(balance >= 300){
                 // waiting for the balance to go below 300.
(2) D thread:
                 belowUpperLimitFundsCondition.await(); }
signalAll()
             balance += amount;
             sufficientFundsCondition.signalAll();
followed by .
             lock.unlock(); }
unlock()
b==100
```

ThreadSafeBankAccount2



```
(3) Two W threads:
             vithdraw(double amount){
                                                        "runnable" One of them
(1) b=0.
              lock.lock();
                                                        acquires the lock again
Two W
              if(balance =< 0){
                  // waiting for the balance to excee and releases it.
threads:
                                                        balance==0.
                 sufficientFundsCondition.await()
"waiting"
             balance -= amount;
                                                        The other W thread:
             belowUpperLimitFundsCondition.signalAll
                                                         "blocked" on acquiring
             lock.unlock(); }
                                                        the lock.
            deposit(double amount) {
                                                          (4) The 2nd W thread
             lock.lock();
                                                          acquires the lock and
             if(balance >= 300) {
                 // waiting for the balance to go below releases it.
(2) D thread:
                 belowUpperLimitFundsCondition.await(); b==-100.
signalAll()
             balance += amount;
             sufficientFundsCondition.signalAll();
followed by .
             lock.unlock(); }
unlock()
balance==100
```

- The 2nd thread should have made sure "balance>0."
- If only one "W" thread runs, this problem does not occur.
- Always use a while loop.

"if" or "while" in Atomic Code?

- You can use "if" rather than "while" for a conditional checking
 - if you use signal(), not signalAll().
- However, in practice, the while-signalAll pair is more common than the if-signal pair.

InterruptedException

- Some methods in Java API can throw InterruptedException.
 - Thread.sleep()
 - Thread.join()
 - Condition.await()
 - ReentrantLock.tryLock()
 - These methods can be long-running and cancellable.

Deposit thread Withdraw thread sufficientFundsCond.await() Interrupted. Re-acquires a lock. Goes to the catch clause. withdraw(double amount) { lock.lock(); while(balance =< 0) { try{ // waiting for the balance to exceed 0 sufficientFundsCondition.await(); } catch (InterruptedException e) { //Do something } } ...}

Condition.await()

- await() lets the currently-executed thread to wait/sleep until another thread wakes it up with signal()/signalAll().
- interrupt() can interrupt a waiting/sleeping thread on a condition object.
 - The waiting/sleeping thread re-acquires a lock and throws an InterruptedException.
 - It does NOT immediately throw an Interrupted Exception.

```
withdraw(double amount) {
  lock.lock();
  while(balance =< 0) {
      try{
            // waiting for the balance to exceed 0
            sufficientFundsCondition.await();
      } catch (InterruptedException e) {
            //Do something
      }
  }
  belowUpperLimitFundsCondition.signalAll();
  balance -= amount;
  lock.unlock(); }</pre>
```