

Alien Method Call

- ```
class A{
 private B b;
 private ReentrantLock lock;
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock();
 }
}
```
- Class B{  
    private A a;  
    public void b1(){  
        do something;  
    }  
}
- This code is deadlock-prone.
  - Calling an *alien method* with a lock held.
  - Alien methods: methods in other classes and overrideable methods in the same class
- It can cause a deadlock if an alien method (b1())...
  - Runs an infinite loop.
  - Acquires B's lock and A's lock.
  - Potential lock-ordering deadlock
  - Spawns a new thread and does a callback.

## A Warning Sign (or Smell) for Deadlocks

### Lock-ordering Deadlock

```
class A{
 private B b;
 private ReentrantLock lock;
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock();
 }
}

class B{
 private A a;
 private ReentrantLock lock;
 public void b1(){
 lock.lock();
 a.a1();
 lock.unlock();
 }
}
```

- This code is deadlock-prone.
  - Calling an *alien method* with a lock held.
  - Alien methods: methods in other classes and overrideable methods in the same class
- It can cause a deadlock if an alien method (b1())...
  - Acquires B's lock and A's lock.
  - Potential lock-ordering deadlock

### Lock-ordering Deadlock

The diagram shows two threads, Thread #1 and Thread #2, interacting with two classes, A and B. Thread #1 starts by locking class A's lock (red arrow). It then calls class B's b1() method (purple arrow). Inside b1(), it locks class B's lock (purple arrow) and calls class A's a1() method (red arrow). Finally, it unlocks class B's lock (purple arrow) and then class A's lock (red arrow). Thread #2 starts by locking class B's lock (purple arrow). It then calls class A's a1() method (red arrow). Inside a1(), it locks class A's lock (red arrow) and calls class B's b1() method (purple arrow). Finally, it unlocks class A's lock (red arrow) and then class B's lock (purple arrow).

```
class A{
 private B b;
 private ReentrantLock lock;
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock();
 }
}

class B{
 private A a;
 private ReentrantLock lock;
 public void b1(){
 lock.lock();
 a.a1();
 lock.unlock();
 }
}
```

- This code can cause a deadlock if an alien method (b1())...
  - Acquires B's lock and A's lock.
  - Potential lock-ordering deadlock
- Thread #1: acquires A's lock, B's lock and A's lock.
  - Acquires A's lock twice (nested locking): no problem
- Thread #2: acquires B's lock and A's lock.

## Note

- ```

• class A{
    private B b;
    private ReentrantLock lock;
    public void a1(){
        lock.lock();
        b.b1();
        lock.unlock();
    }
}

```
- ```

• Class B{
 private A a;
 public void b1(){
 do something;
 }
}

```
- If you implement A and B AND use them by yourself, be careful NOT to generate a deadlock.
- If you implement A and B as an API designer and allow others to use them, you have NO WAYS to prevent them from causing a lock-ordering deadlock.

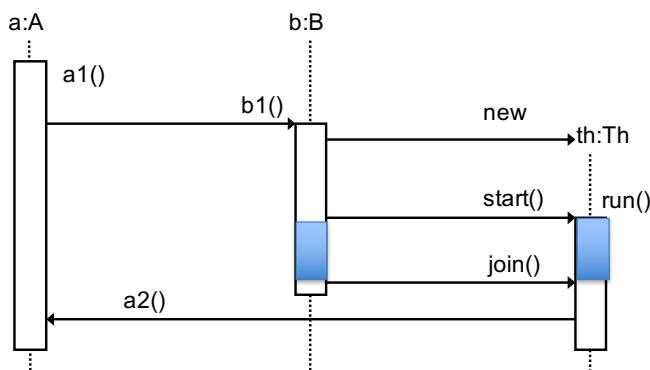
```

• class A{
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock(); }
 public void a2(){
 lock.lock();
 ...
 lock.unlock(); }
}

• class B{
 public void b1(){
 Thread Th = new Thread(
 new Th());
 th.start();
 ...
 th.join(); } }

class Th implements Runnable{
 public void run(){ a.a2(); }
}

```



## Deadlock by a Concurrent Callback

```

• class A{
 private B b;
 private ReentrantLock lock;
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock();
 }
 public void a2(){
 lock.lock();
 ...
 lock.unlock();
 }
}

• Class B{
 private A a;
 public void b1(){

 Thread th = new Thread(
 new Th());
 th.start();
 ...
 th.join();
 }
}

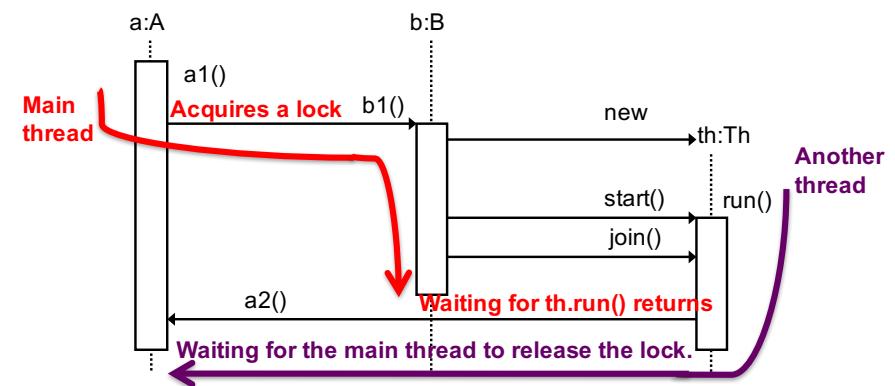
class Th implements Runnable{
 private A a;
 public void run(){
 a.a2();
 }
}

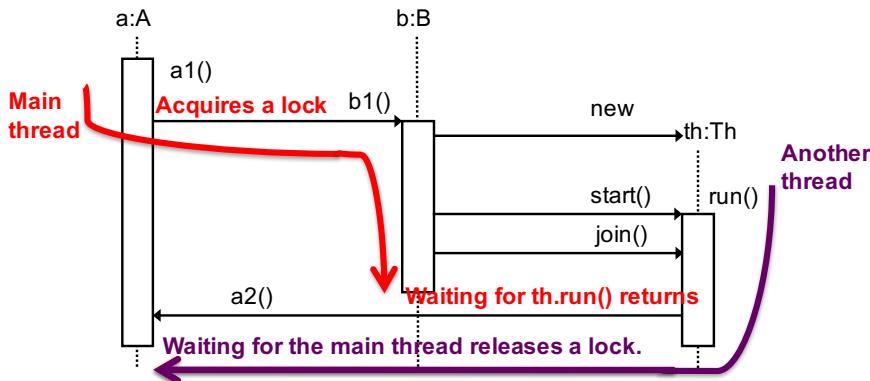
• class A{
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock(); }
 public void a2(){
 lock.lock();
 ...
 lock.unlock(); }
}

• class B{
 public void b1(){
 Thread Th = new Thread(
 new Th());
 th.start();
 ...
 th.join(); } }

class Th implements Runnable{
 public void run(){ a.a2(); }
}

```





- A client class holds a lock and calls an **alien** method.
  - Alien methods = methods in other classes and overrideable methods (neither private nor final) in the same class.
- The alien method spawns a thread that callbacks a method on the client.
  - The callback method requires a lock that the client class already holds.
- The alien method waits for a certain condition that the newly-created thread makes true.

```

• class A{
 private B b;
 private ReentrantLock lock;
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock();
 }
}

• Class B{
 private A a;
 public void b1(){
 do something;
 }
}

```

- If you implement A and B AND use them by yourself, be careful NOT to generate a deadlock.
- If you implement A and B as an API designer and allow others to use them, you have NO WAYS to prevent them from making a callback in `b1()`.

## Another Example

```

• class A{
 private B b;
 private ReentrantLock lock;
 public void a1(){
 lock.lock();
 b.b1();
 lock.unlock();
 }
}

• Class B{
 private A a;
 public abstract void b1(){}
}

```

- `b1()` is an alien method. Deadlock prone.
- If you allow others to define B's subclass and implement `b1()`, you have NO WAYS to prevent them from causing a deadlock.

## One More Example

```

• class A{
 private B b;
 private ReentrantLock lock;
 public void a1(){
 lock.lock();
 a2();
 lock.unlock();
 }
 protected abstract void a2(){}
}

```

- `b1()` is an alien method. Deadlock prone.
- If you allow others to define A's subclass and implement `a2()`, you have NO WAYS to prevent them from causing a deadlock.

## If You Cannot Prevent Your Class (API) from Causing Deadlocks...

- ```
class A{  
    private B b;  
    private ReentrantLock lock;  
    public void a1(){  
        lock.lock();  
        ...  
        lock.unlock();  
        b.b1(); // open call  
    }  
}
```
- Open call
 - Do not call `b1()`, an alien method, in atomic code; move the method call outside the atomic code.

```
Class B{  
    private A a;  
    public void b1(){  
        do something;  
    }  
}
```

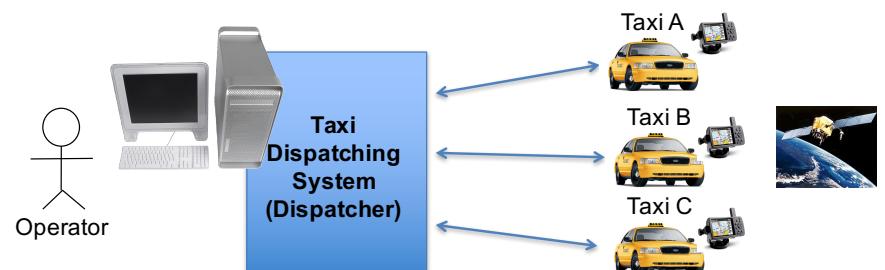
Rule of Thumb

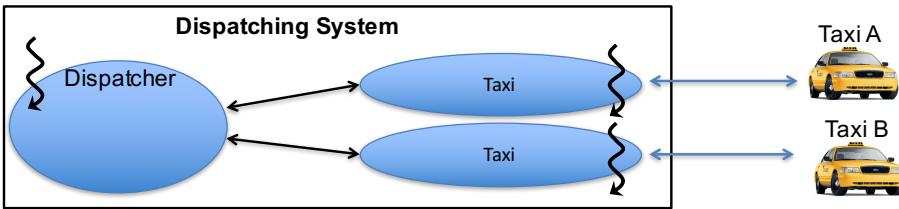
- Try NOT to call alien methods from atomic code.
- Call an alien method(s) outside atomic code.
 - Open call.
- Carefully eliminate race conditions.
- If necessary, compromise to favor being deadlock-free than being free from race conditions.

Exercise: Taxi Dispatching System

Exercise: Taxi Dispatching System

- Dispatcher
 - Periodically obtains each taxi's current location and displays it on the operator's monitor.
 - Allows the operator to set the destination location to each taxi.
- Each taxi's in-car system
 - obtains the current location and notifies the dispatcher if it has arrived at the destination.





```

class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ...
    public void notifyAvailable(Taxi t){
        availableTaxis.add(t);
        ...
    }
    public void displayAvailableTaxis(){
        for(Taxi t: availableTaxis)
            display.draw(t.getLocation());
        ...
    }
}

```

- class Taxi implements Runnable{
 private Point location;
 private Point dest;
 private Dispatcher dispatcher;
- public Point getLocation(){
 return location;
 }
- public setLocation(Point loc){
 location = loc;
 if(location.equals(dest))
 dispatcher.notifyAvailable(this);
 }
- public void run(){
 while(true){
 setLocation(getGPSLoc());
 }
 } }

- What variables are shared by multiple threads?
- Dispatcher's availableTaxis
- Dispatcher could add/remove taxis to/from availableTaxis.

```

class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ...
    public void notifyAvailable(Taxi t){
        availableTaxis.add(t);
        ...
    }
    public void displayAvailableTaxis(){
        for(Taxi t: availableTaxis)
            display.draw(t.getLocation());
        ...
    }
}

```

- class Taxi implements Runnable{
 private Point location;
 private Point dest;
 private Dispatcher dispatcher;
- public Point getLocation(){
 return location;
 }
- public setLocation(Point loc){
 location = loc;
 if(location.equals(dest))
 dispatcher.notifyAvailable(this);
 }
- public void run(){
 while(true){
 setLocation(getGPSLoc());
 }
 } }

- What variables are shared by multiple threads?
- Taxi's location

```

class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ...
    public void notifyAvailable(Taxi t){
        availableTaxis.add(t);
        ...
    }
    public void displayAvailableTaxis(){
        for(Taxi t: availableTaxis)
            display.draw(t.getLocation());
        ...
    }
}

```

- class Taxi implements Runnable{
 private Point location;
 private Point dest;
 private Dispatcher dispatcher;
- public Point getLocation(){
 return location;
 }
- public setLocation(Point loc){
 location = loc;
 if(location.equals(dest))
 dispatcher.notifyAvailable(this);
 }
- public void run(){
 while(true){
 setLocation(getGPSLoc());
 }
 } }

- Guard shared variables with locks.
- Now, what about deadlocks?

```

class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ...
    public void notifyAvailable(Taxi t){
        lockD.lock();
        availableTaxis.add(t);
        ...
        lockD.unlock();
    }
    public void displayAvailableTaxis(){
        lockD.lock();
        for(Taxi t: availableTaxis)
            display.draw(t.getLocation());
        ...
        lockD.unlock();
    }
}

```

- class Taxi implements Runnable{
 private Point location;
 private Point dest;
 private Dispatcher dispatcher;
- public Point getLocation(){
 lockT.lock();
 return location;
 lockT.unlock();
 }
- public setLocation(Point loc){
 lockT.lock();
 location = loc;
 if(location.equals(dest))
 dispatcher.notifyAvailable(this);
 lockT.unlock();
 }
- public void run(){
 while(true){
 setLocation(getGPSLoc());
 }
 } }

- Now, what about deadlocks?

- Trace threads of control and understand the order of lock acquisitions.

```

Class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ...
    public notifyAvailable(Taxi t){
        lockD.lock();
        availableTaxis.add(t);
        ...
        lockD.unlock();
    }
    ...
    public displayAvailableTaxis(){
        lockD.lock();
        for(Taxi t: availableTaxis)
            display.draw(t.getLocation());
        ...
        lockD.unlock();
    }
    ...
}

    • class Taxi implements Runnable{
        private Point location;
        private Point dest;
        private Dispatcher dispatcher;
        ...
        public Point getLocation(){
            lockT.lock();
            return location;
            lockT.unlock();
        }
        public setLocation(Point loc){
            lockT.lock();
            location = loc;
            if(location.equals(dest))
                dispatcher.notifyAvailable(this);
            lockT.unlock();
        }
        public void run(){
            while(true){
                setLocation(getGPSLoc());
            }
            .... }
    }
}

```

lockD → lockT

- Now, what about deadlocks?

- Two threads acquire the same set of locks in order! Be careful!!!
- Lock-ordering deadlocks can occur.

```

Class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ...
    public notifyAvailable(Taxi t){
        lockD.lock(); //blocked
        availableTaxis.add(t);
        ...
        lockD.unlock();
    }
    ...
    public displayAvailableTaxis(){
        lockD.lock();
        for(Taxi t: availableTaxis)
            display.draw(t.getLocation());
        ...
        lockD.unlock();
    }
    ...
}

    • class Taxi implements Runnable{
        private Point location;
        private Point dest;
        private Dispatcher dispatcher;
        ...
        public Point getLocation(){
            lockT.lock(); //blocked
            return location;
            lockT.unlock();
        }
        public setLocation(Point loc){
            lockT.lock();
            location = loc;
            if(location.equals(dest))
                dispatcher.notifyAvailable(this);
            lockT.unlock();
        }
        public void run(){
            while(true){
                setLocation(getGPSLoc());
            }
            .... }
    }
}

```

(1) (2) (3) (4)

- Now, what about deadlocks?

- Trace threads of control and understand the order of lock acquisitions.

```

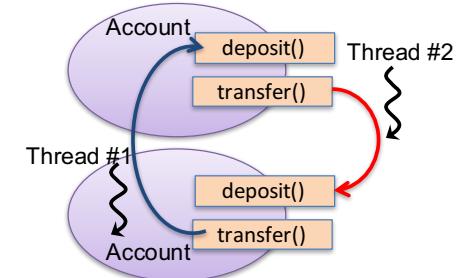
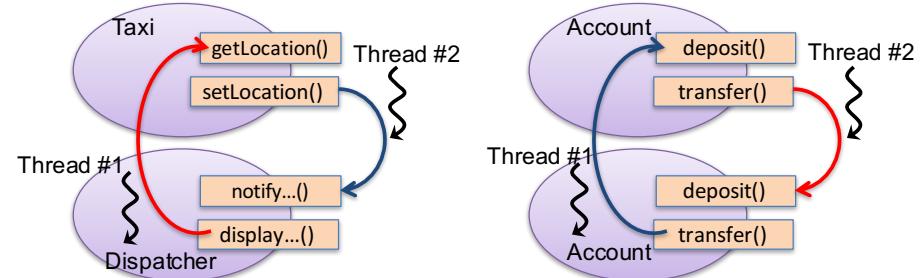
Class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ...
    public notifyAvailable(Taxi t){
        lockD.lock();
        availableTaxis.add(t);
        ...
        lockD.unlock();
    }
    ...
    public displayAvailableTaxis(){
        lockD.lock();
        for(Taxi t: availableTaxis)
            display.draw(t.getLocation());
        ...
        lockD.unlock();
    }
    ...
}

    • class Taxi implements Runnable{
        private Point location;
        private Point dest;
        private Dispatcher dispatcher;
        ...
        public Point getLocation(){
            lockT.lock();
            return location;
            lockT.unlock();
        }
        public setLocation(Point loc){
            lockT.lock();
            location = loc;
            if(location.equals(dest))
                dispatcher.notifyAvailable(this);
            lockT.unlock();
        }
        public void run(){
            while(true){
                setLocation(getGPSLoc());
            }
            .... }
    }
}

```

lockD → lockT **lockT → lockD**

Lock-ordering Deadlocks



Common Solutions

- Static lock
- Timed locking
- Ordered locking
- Nested tryLock()

- Now, getLocation() and notifyAvailable() are open calls.
 - No lock-ordering deadlocks occur.

```
Class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ....
    public void notifyAvailable(Taxi t){
        lockD.lock();
        availableTaxis.add(t);
        ....
        lockD.unlock();
    }

    public void displayAvailableTaxis(){
        HashSet<Taxi> availableTaxisLocal;
        lockD.lock();
        availableTaxisLocal =
            new HashSet<Taxi>(availableTaxis);
        lockD.unlock();
        for(Taxi t: availableTaxisLocal)
            display.draw(t.getLocation());
            // Open call
        ....
    }
}
```

```
    • class Taxi implements Runnable{
        private Point location;
        private Point dest;
        private Dispatcher dispatcher;

        public Point getLocation(){
            lockT.lock();
            return location;
            lockT.unlock();
        }

        public void setLocation(Point loc){
            lockT.lock();
            try{
                location = loc;
                if(!location.equals(dest))
                    return; // Balking
            }finally{
                lockT.unlock(); }
            dispatcher.notifyAvailable(this);
            // Open call
        }

        public void run(){...} }
```

An Alternative Solution

- Open method call
 - Calling a method with no locks held.

Note 1

```
Class Dispatcher{
    private HashSet<Taxi> taxis;
    private HashSet<Taxi> availableTaxis;
    ....
    public void notifyAvailable(Taxi t){
        lockD.lock();
        availableTaxis.add(t);
        ....
        lockD.unlock();
    }

    public Point getLocation(){
        lockT.lock();
        return location;
        lockT.unlock();
    }

    public void setLocation(Point loc){
        lockT.lock();
        try{
            location = loc;
            if(!location.equals(dest))
                return; // Balking
        }finally{
            lockT.unlock(); }
        dispatcher.notifyAvailable(this);
        // Open call
    }

    public void run(){...} }
```

lockT has been released when calling notifyAvailable().

A context switch can occur here. The state of "location==dest" never changes in the rest of this method.

Calling notifyAvailable() can be outside atomic code.

Note 2

```
Class Dispatcher{  
    private HashSet<Taxi> taxis;  
    private HashSet<Taxi> availableTaxis;  
    ....  
    public notifyAvailable(Taxi t){  
        lockD.lock();  
        availableTaxis.add(t);  
        ....  
        lockD.unlock();  
    }  
  
    public displayAvailableTaxis(){  
        HashSet<Taxi> availableTaxisLocal;  
        lockD.lock();  
        availableTaxisLocal =  
            new HashSet<Taxi>(availableTaxis);  
        lockD.unlock();  
        for(Taxi t: availableTaxisLocal)  
            display.draw(t.getLocation());  
            // Open call  
        ....  
    } ...}
```

- class Taxi implements Runnable{
 private Point location;
 private Point dest;
 private Dispatcher dispatcher;

 public Point getLocation(){
 lockT.lock();
 return location;
 lockT.unlock();
 }
 A local variable is not shared
 by multiple threads.
 HashSet is not thread-safe.

 A copy to a local variable must
 be in atomic code. notifyAvailable()
 should not be called during this
 copy operation.

Note 2

```
Class Dispatcher{  
    private HashSet<Taxi> taxis;  
    private HashSet<Taxi> availableTaxis;  
    ....  
    public notifyAvailable(Taxi t){  
        lockD.lock();  
        availableTaxis.add(t);  
        ....  
        lockD.unlock();  
    }  
  
    public displayAvailableTaxis(){  
        HashSet<Taxi> availableTaxisLocal;  
        lockD.lock();  
        availableTaxisLocal =  
            new HashSet<Taxi>(availableTaxis);  
        lockD.unlock();  
        for(Taxi t: availableTaxisLocal)  
            display.draw(t.getLocation());  
            // Open call  
        ....  
    } ...}
```

- class Taxi implements Runnable{
 private Point location;
 private Point dest;
 private Dispatcher dispatcher;

 public Point getLocation(){
 lockT.lock();
 return location;
 lockT.unlock();
 }
 A context switch can occur here.
 notifyAvailable() may be called.
 A race condition can occur in fact.

 availableTaxisLocal may not be in
 synch with availableTaxis.

Note 3

```
Class Dispatcher{  
    private HashSet<Taxi> taxis;  
    private HashSet<Taxi> availableTaxis;  
    ....  
    public notifyAvailable(Taxi t){  
        lockD.lock();  
        availableTaxis.add(t);  
        ....  
        lockD.unlock();  
    }  
  
    public displayAvailableTaxis(){  
        lockD.lock();  
        for(Taxi t: availableTaxis)  
            display.draw(t.getLocation());  
            // NOT OPEN CALL  
        ....  
        lockD.unlock();  
    } ...}  
  
A deadlock can still occur!  
Need two open calls.
```

- class Taxi implements Runnable{
 private Point location;
 private Point dest;
 private Dispatcher dispatcher;

 public Point getLocation(){
 lockT.lock();
 return location;
 lockT.unlock();
 }
 public setLocation(Point loc){
 lockT.lock();
 try{
 location = loc;
 if(!location.equals(dest))
 return; // Balking
 }finally{
 lockT.unlock();
 }
 dispatcher.notifyAvailable(this);
 // Open call
 }
 public void run(){...} }

Rule of Thumb

- Try NOT to call alien methods from atomic code.
- Call alien methods outside atomic code.
 - Open call.
- Carefully eliminate race conditions.
- If necessary, compromise to favor being deadlock-free than being race condition free.

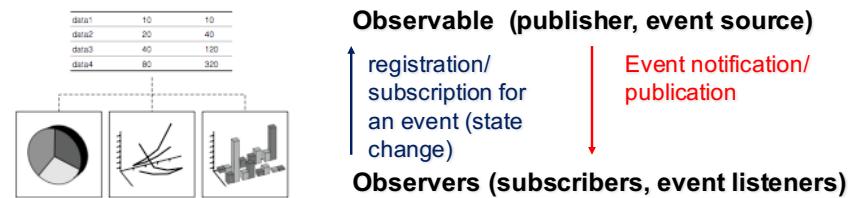
HW 22 [Optional]

- Complete thread-safe code for the taxi application
 - Use open calls

Exercise: Concurrent Observer Design Pattern

Observer Design Pattern (Recap)

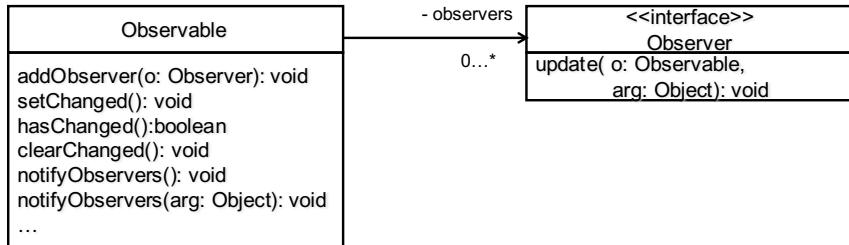
- Intent
 - Event notification
 - Define a one-to-many dependency between objects so that when one object changes its state, all its dependents are notified and updated automatically
- a.k.a
 - Publish-Subscribe (pub/sub)
 - Event source - event listener
- Two key participants (classes/interfaces)
 - Observable (model, publisher or subject)
 - Propagates an event to its dependents (observers) when its state changes.
 - Observer (view and subscriber)
 - Receives events from an observable object.



- Separate data processing from data management.
 - Data management: Observable
 - Visualization/data processing: Observers
 - e.g., Data calculation, GUI display, etc.

Class Structure

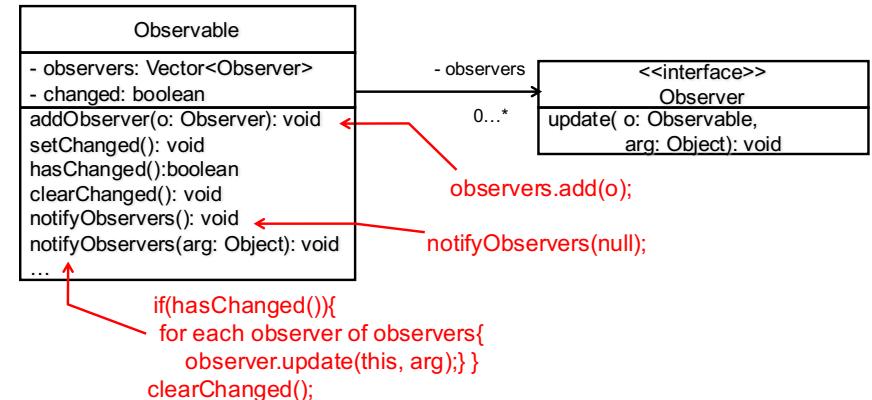
- java.util.Observable
- java.util.Observer



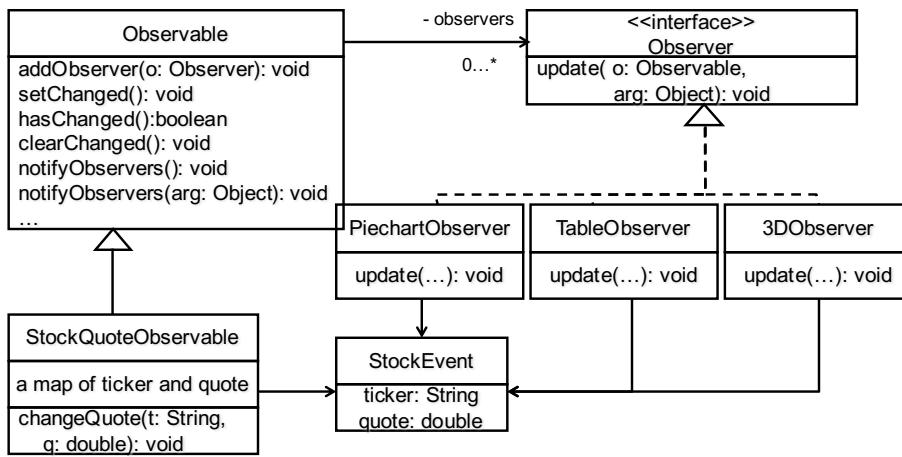
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Class Structure

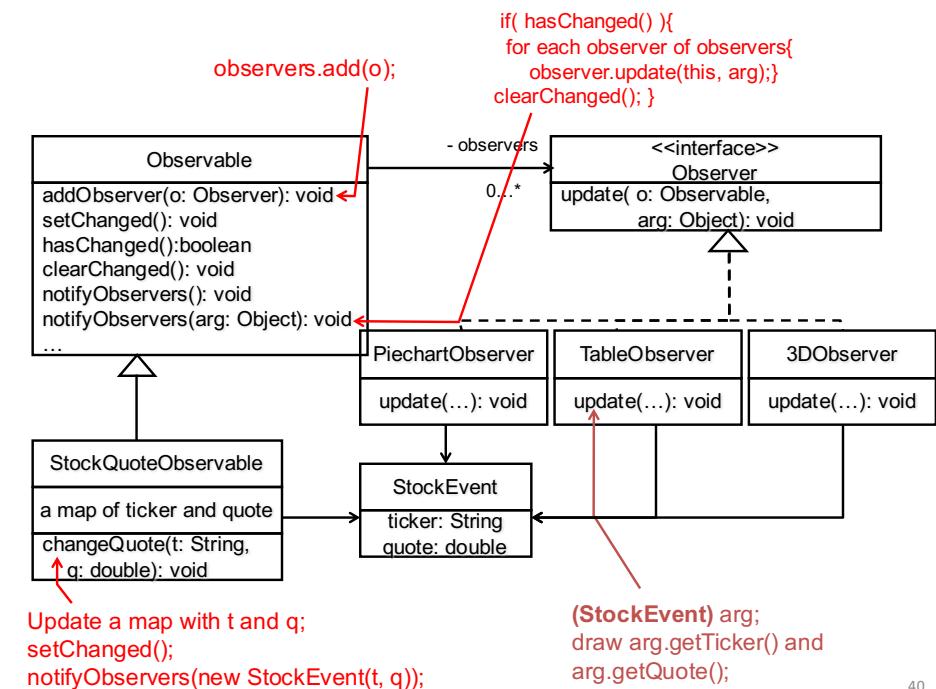
- java.util.Observable
- java.util.Observer



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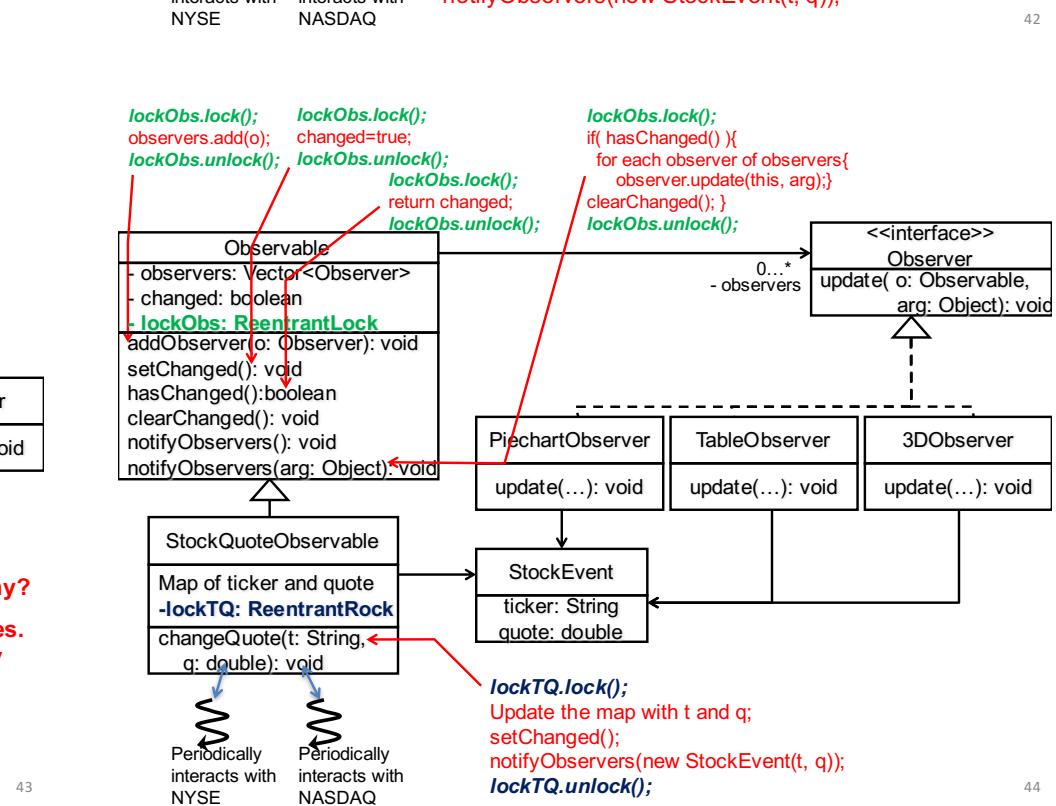
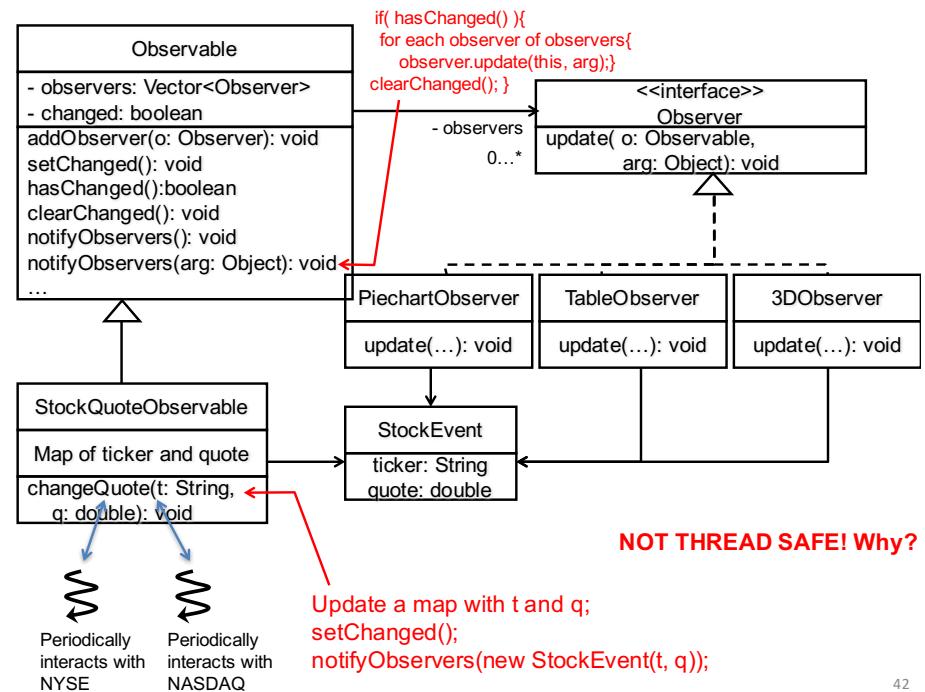
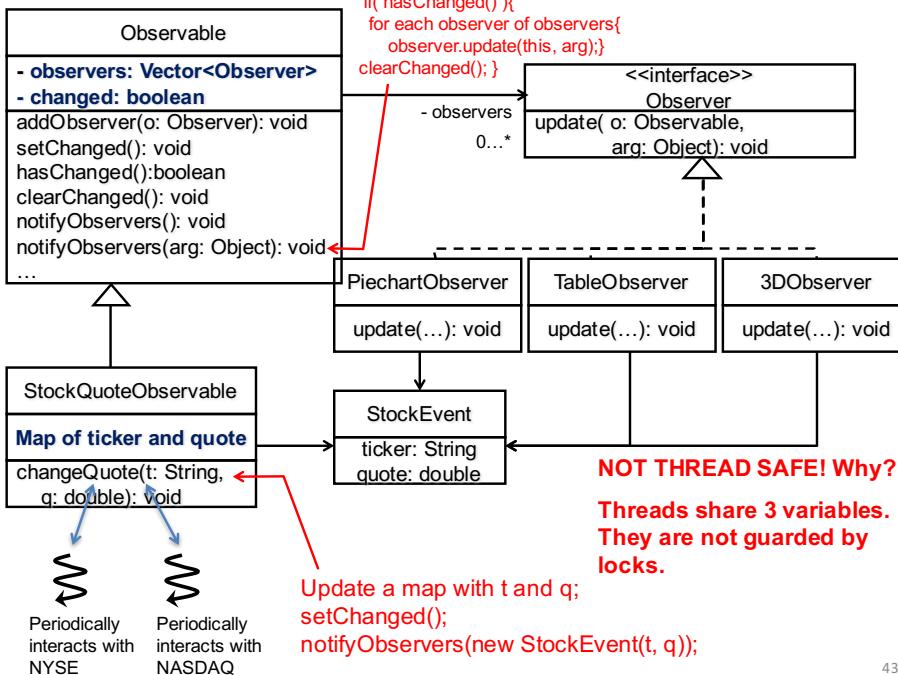
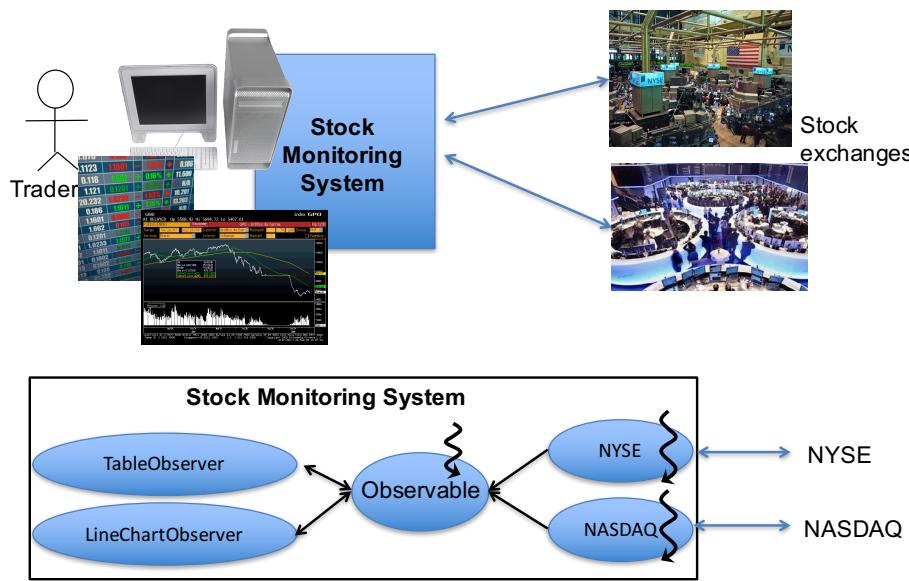


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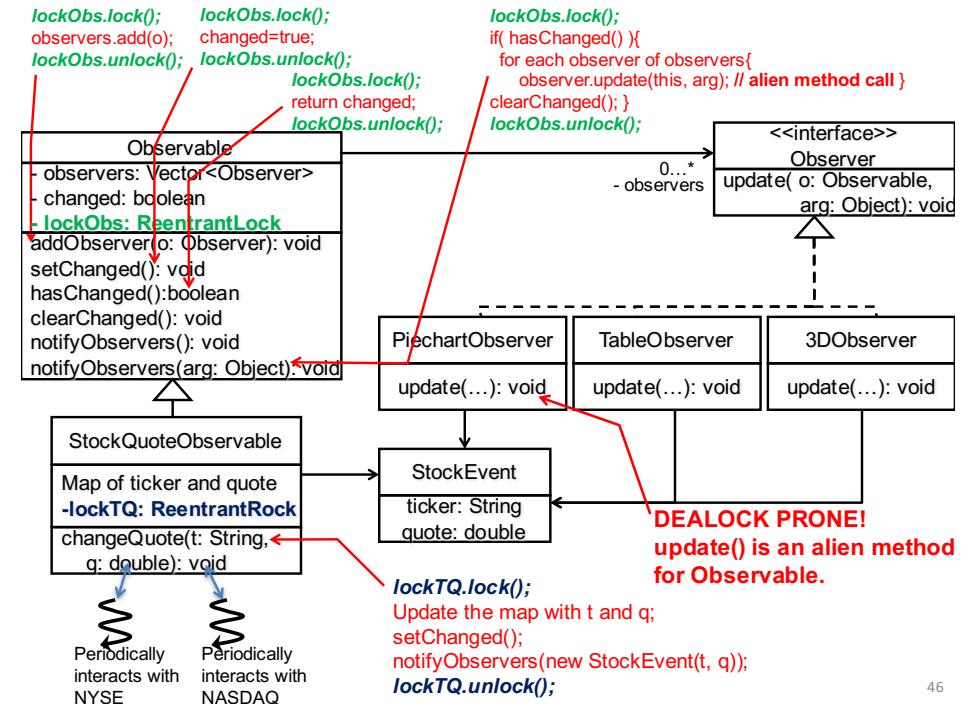
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Concurrent Observer



Notes

- Two threads try to acquire two locks in order in `changeQuote()`:
 - `lockTQ` and `lockObs`
 - Be careful for potential lock-ordering deadlocks.
 - Are there any threads that try to acquire `lockObs` and then `lockTQ`?
- Observable guards two shared variables (“observers” and “changed”) with a single lock.
 - No problem in terms of thread safety
 - Simple and thread-safe implementation strategy
 - Not the most efficient implementation strategy
 - Could be more efficient by guarding two shared variables with two locks.
 - One for “observers” and the other for “changed”



Potential Deadlocks due to Alien Method Calls

- `update()` (an alien method) may...
 - run infinite loop.
 - obtain an observer’s lock and spawn a new thread that requires `lockObs` (i.e. the lock in Observable)
 - Lock-ordering deadlocks could occur.
 - spawn a new thread and perform a callback toward Observable.
- If you are an API designer for Observable, there are no ways for you to prevent these potential cases.
 - You need to implement Observable in a preventive way.

Open Call in Observable.notifyObservers()

```

• class Observable {
    private Vector obs;           //observers
    private boolean changed = false;

    public void notifyObservers(Object arg) {
        Object[] arrLocal;
        synchronized (this) {
            if (!changed) return; // balking
            arrLocal = obs.toArray(); // observers copied to arrLocal
            changed = false;
        }
        for (int i = arrLocal.length-1; i <=0; i--)
            ((Observer)arrLocal[i]).update(this, arg); // OPEN CALL
    }
    ....
}
  
```

Observable.notifyObservers()

```
• class Observable {  
    private Vector obs;           //observers  
    private boolean changed = false;  
  
    public void notifyObservers(Object arg) {  
        Object[] arrLocal;  
        synchronized(this){          // lock.lock()  
            if (!changed) return;    // balking  
            arrLocal = obs.toArray(); // observers copied to arrLocal  
            changed = false;  
        }                           // lock.unlock()  
        for(int i = arrLocal.length-1; i <= 0; i--)  
            ((Observer)arrLocal[i]).update(this, arg); // OPEN CALL  
    }  
....  
}
```

A context switch can occur here.
addObserver() and removeObserver()
may be called. A race condition can
occur.

arrLocal may not be in synch with obs.

Comments in Observable's Source Code Say...

- The worst result of any potential race-condition here is that:
 - 1) a newly-added Observer will miss a notification in progress
 - 2) a recently unregistered Observer will be wrongly notified when it doesn't care (about state-change notifications anymore).
- To look up Java API's source code...
 - Use src.zip in your JDK installation directory
 - Use online source code browsers
 - <http://www.docjar.com/>
 - <http://grepcode.com/>
 - <http://javasourcecode.org/>

Preventive Concurrent API Design

- If you implement the Observer-Observable API and its application(s) AND use both of them, be careful about potential deadlocks.
 - Deadlock by an infinite loop
 - Lock-ordering deadlock
 - Deadlock by a callback

Rule of Thumb

- Try NOT to call alien methods from atomic code.
- Call alien methods outside atomic code.
 - Open call
- Carefully eliminate race conditions.
- If necessary, compromise to favor being deadlock-free than having no race conditions.

HW 23

- Define your own Observable (class) and Observer (interface).
 - DO NOT reuse java.util.Observable and java.util.Observer.
 - Define Observer as a functional interface.
 - update() as an abstract method.
 - Implement the listed methods for Observable
 - C.f. Java API doc for expected behaviors/responsibilities for the methods
- Note that you have done this in HW 4.

