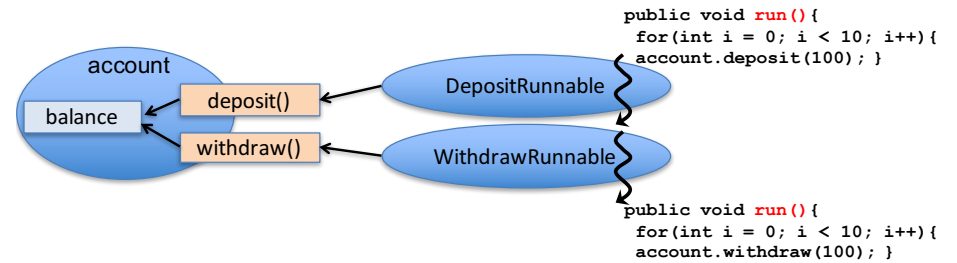
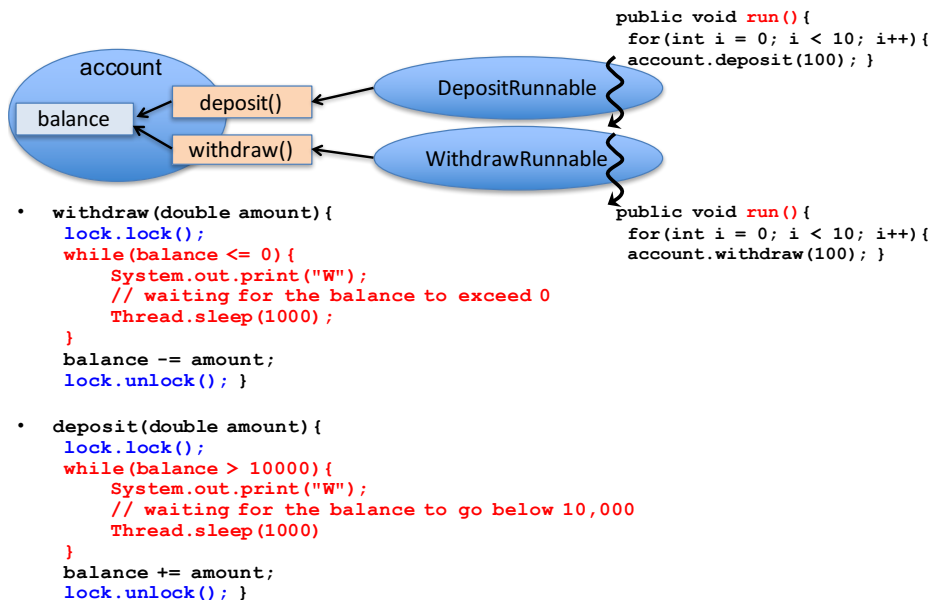


Deadlock

DeadlockedBankAccount.java

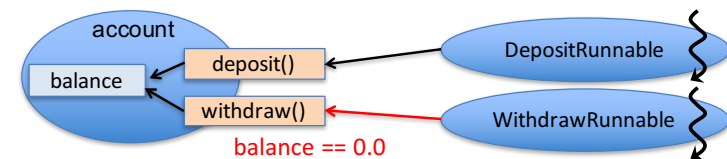


DeadlockedBankAccount.java



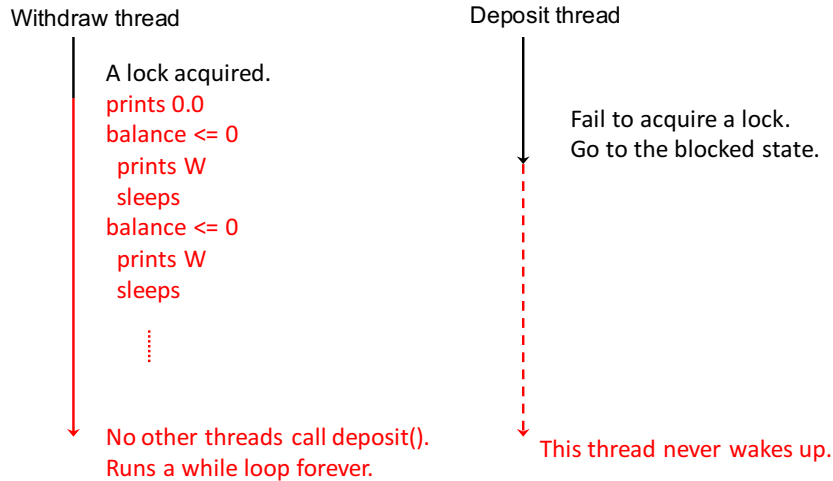
Deadlock

- Assume the withdrawal thread goes ahead and runs first.
- Output
 - Lock obtained
 - Current balance (w): 0.0



DeadlockedBankAccount2.java

- Current balance (w): 0.0WWWWWWWW



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- Previous version

```

- withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        System.out.print("W");
        // waiting for the
        // balance to exceed 0
        Thread.sleep(1000);
    }
    balance -= amount;
    lock.unlock();
}

- deposit(double amount){
    lock.lock();
    while(balance > 10000){
        System.out.print("W");
        // waiting for the balance
        // to go below 10,000
        Thread.sleep(1000);
    }
    balance += amount;
    lock.unlock();
}

```

- New version

```

- withdraw(double amount){
    while( balance <= 0 ){
        System.out.print("W");
        Thread.sleep(2);
    }
    lock.lock();
    balance -= amount;
    lock.unlock();
}

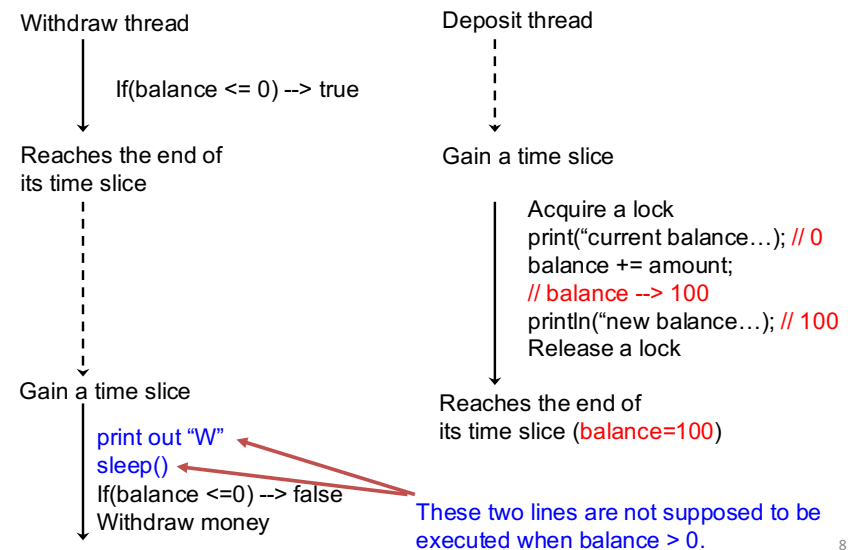
- deposit(double amount){
    while( balance > 10000 ){
        System.out.print("W");
        Thread.sleep(2);
    }
    lock.lock();
    balance += amount;
    lock.unlock();
}

```

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- Has no deadlock problems.
- Can generate race conditions.

A Potential Race Condition in DeadlockedBankAccount2

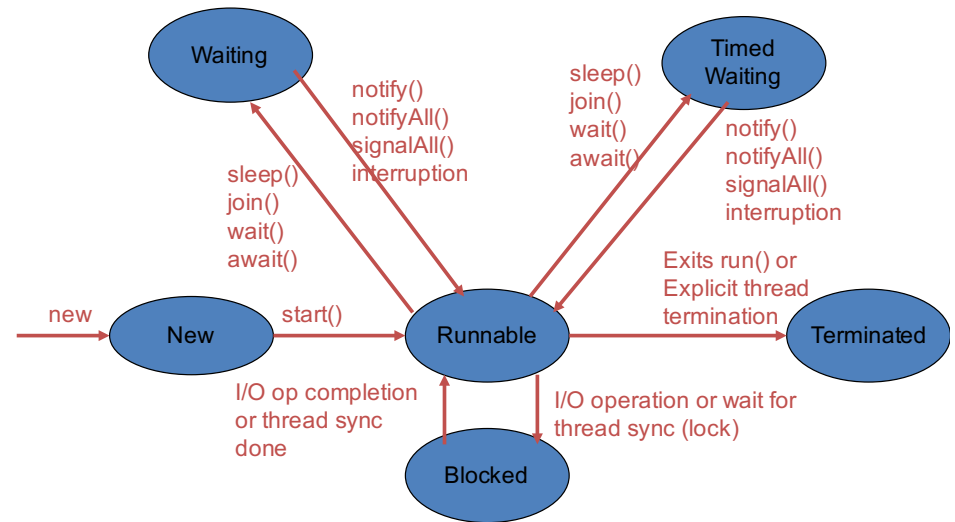


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Avoiding Deadlocks and Race Conditions

- Use a Condition object(s).
 - Allow a thread to
 - temporarily release a lock so that another thread can proceed
 - The thread goes to the Waiting state from the Runnable state.
 - re-acquire the lock later.
- java.util.concurrent.locks.Condition
 - Obtain its instance from a lock object
 - `ReentrantLock lock = new ReentrantLock();`
`Condition condition = lock.newCondition(); //factory method`
`condition.await();`



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ThreadSafeBankAccount2

- `Condition sufficientFundsCondition = lock.newCondition();`
`Condition belowUpperLimitFundsCondition = lock.newCondition();`
- `withdraw(double amount){`
`lock.lock();`
`while(balance <= 0){`
`// waiting for the balance to exceed 0`
`sufficientFundsCondition.await(); }`
`balance -= amount;`
`belowUpperLimitFundsCondition.signalAll();`
`lock.unlock(); }`
- `deposit(double amount){`
`lock.lock();`
`while(balance >= 300){`
`// waiting for the balance to go below 10000.`
`belowUpperLimitFundsCondition.await(); }`
`balance += amount;`
`sufficientFundsCondition.signalAll();`
`lock.unlock(); }`

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ThreadSafeBankAccount2

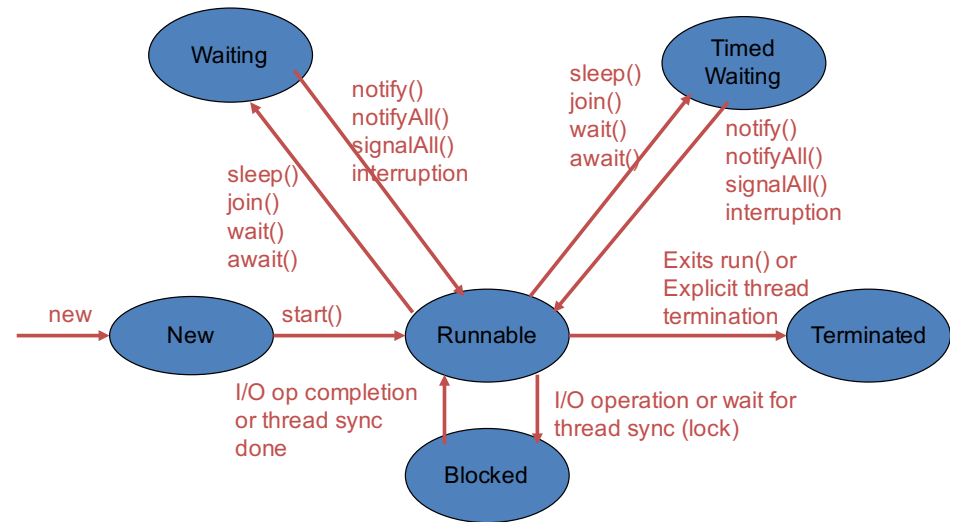
- Output
 - Lock obtained
 - 7 (d): current balance: 0.0
 - 7 (d): new balance: 100.0
 - Lock released
 - Lock obtained
 - 8 (d): current balance: 100.0
 - 8 (d): new balance: 200.0
 - Lock released
 - Lock obtained
 - 9 (d): current balance: 200.0
 - 9 (d): new balance: 300.0
 - Lock released
 - Lock obtained
 - 10 (d): current balance: 300.0
 - 10 (d): await(): Balance has reached the upper limit.
 - Lock obtained
 - 11 (d): current balance: 300.0
 - 11 (d): await(): Balance has reached the upper limit.
 - Lock obtained
 - 12 (w): current balance: 300.0
 - 12 (w): new balance: 200.0
 - Lock released
 - 10 (d): new balance: 300.0
 - Lock released
 - 11 (d): await(): Balance has reached the upper limit.

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Condition

- `await()`
 - Will be waiting until it is signaled or interrupted
 - Will be waiting until it is signaled or interrupted, or until a specified waiting time (relative time) elapsed.
 - Will be waiting until it is signaled or interrupted, or until a specified deadline (absolute time).
 - If signaled, goes to the Runnable state and re-acquires a lock.
 - Will be “blocked” if fails to re-acquire the lock.
 - Throws an `InterruptedException` if interrupted.
 - c.f. previous lecture note that explains `InterruptedException`
- `signalAll()`
 - Wakes up all waiting threads on a condition object.
 - All of them go to the “runnable” state.
 - One of them will re-acquire a lock.

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- When a thread calls `await()`, `signal()` or `signalAll()` on a Condition object,
 - the thread is assumed to hold the lock associated with the Condition object.
- If the thread does not,
 - an `IllegalMonitorStateException` is thrown.

SignalAll() Before or After a State Change?

```

• withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
    balance -= amount;
    belowUpperLimitFundsCondition.signalAll();
    lock.unlock(); }

• deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
    balance += amount;
    sufficientFundsCondition.signalAll();
    lock.unlock(); }

```

- What happens if you call `signalAll()` first and then update the balance? Any problems?

- ```

withdraw(double amount){
 lock.lock();
 while(balance <= 0){
 // waiting for the balance to exceed 0
 sufficientFundsCondition.await(); }
 belowUpperLimitFundsCondition.signalAll();
 balance -= amount;
 lock.unlock(); }

```
- ```

deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
        sufficientFundsCondition.signalAll();
        balance += amount;
        lock.unlock(); }

```
- Need to worry about race conditions in this case?

(1) W thread: "waiting" temporarily releases the lock

```

withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
        belowUpperLimitFundsCondition.signalAll();
        balance -= amount;
        lock.unlock(); }

```

(2) D thread: signalAll(). Ctx switch

```

deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
        sufficientFundsCondition.signalAll();
        balance += amount;
        lock.unlock(); }

```

- A context switch can occur in between signalAll() and state change?
- A "W" thread can withdraw money before a "D" thread deposits money?
 - Can the balance variable have a negative value?

Two Important Things (1)

(1) W thread: "waiting" temporarily releases the lock

```

withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await(); }
        belowUpperLimitFundsCondition.signalAll();
        balance -= amount;
        lock.unlock(); }

```

(3) W thread: "runnable" Tries to acquire the lock again and fails. Goes to "blocked."

```

deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await(); }
        sufficientFundsCondition.signalAll();
        balance += amount;
        lock.unlock(); }

```

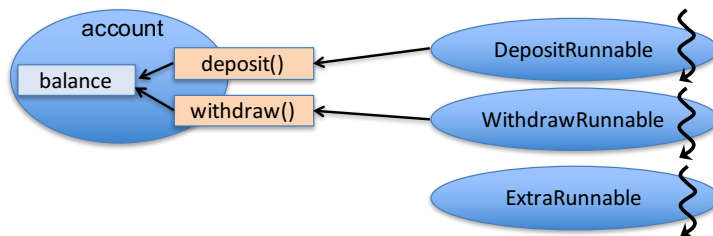
(2) D thread: signalAll(). Ctx switch

- A "W" thread CANNOT withdraw money before a "D" thread deposits money.
- A "D" thread CANNOT deposit money before a "W" thread withdraws money.

- A state change (or value change) can be made on a shared variable safely after calling signalAll()
 - AS FAR AS the state changes in atomic code
- Common programming convention/practice:
 - A state change first, followed by signalAll().

Two Important Things (2)

- A JVM DOES context switches even when a thread runs in atomic code.
 - A lock guarantees that only one thread exclusively runs atomic code at a time.



Note that...

- Some books and online materials explicitly/implicitly say that context switches never occur when a thread runs in atomic code.
- It is wrong!

signal() and signalAll()

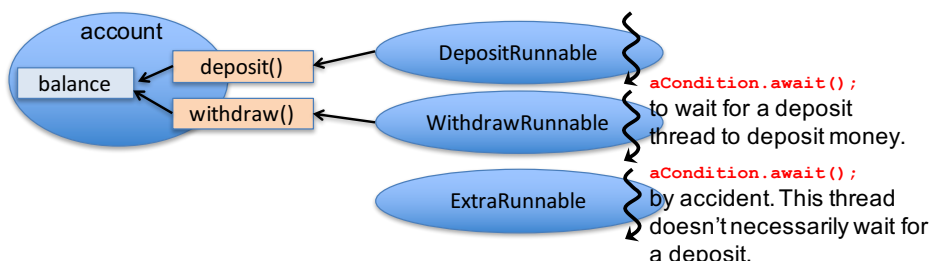
- signalAll()
 - Wakes up all waiting threads on a condition object.
 - All of them go to the “runnable” state.
 - One of them will re-acquire a lock.
- signal()
 - Wakes up one of waiting threads on a condition object.
 - One of them goes to the “runnable” state. The others stay at the “waiting” state.
 - JVM’s thread scheduler selects one of them. Assume a random selection.
 - Not predictable which waiting thread to be selected.

ThreadSafeBankAccount2

- ```
Condition sufficientFundsCondition = lock.newCondition();
Condition belowUpperLimitFundsCondition = lock.newCondition();
```
- ```
withdraw(double amount){
    lock.lock();
    while(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await();
    }
    balance -= amount;
    belowUpperLimitFundsCondition.signalAll();
    lock.unlock();
}
```
- ```
deposit(double amount){
 lock.lock();
 while(balance >= 300){
 // waiting for the balance to go below 300.
 belowUpperLimitFundsCondition.await();
 }
 balance += amount;
 sufficientFundsCondition.signalAll();
 lock.unlock();
}
```

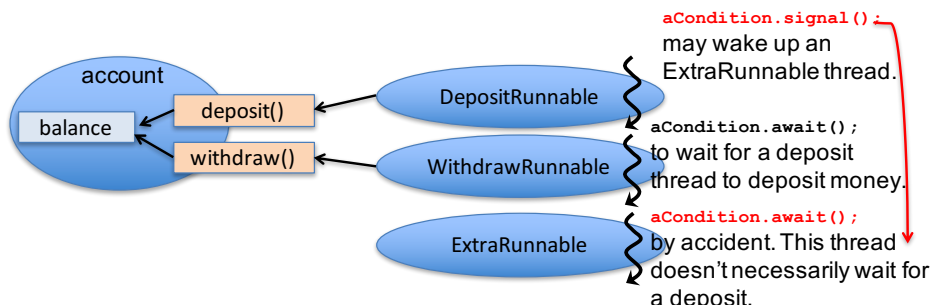
## signal() or signalAll()?

- `signal()` is more lightweight than `signalAll()`.
  - Waking up a waiting thread is computationally expensive.
  - `signal()` wakes up only one thread.
- `signalAll()` is more *protective*.
  - `signalAll()` should be favored (at least in my personal taste).



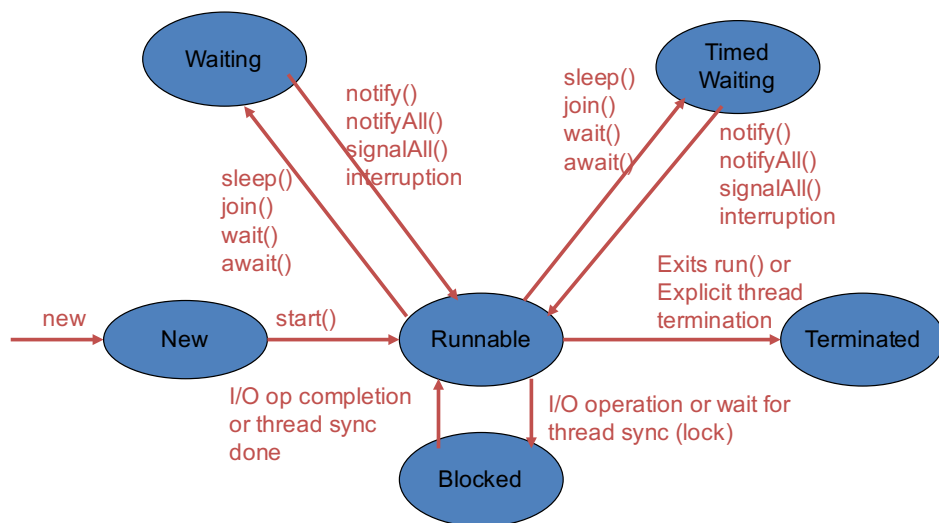
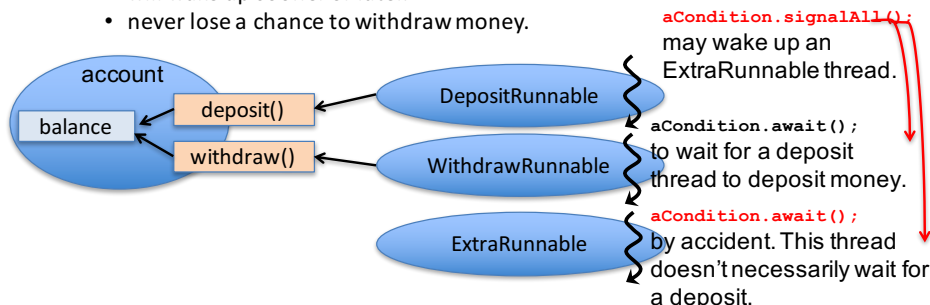
## signal()

- A deposit thread may wake up an `ExtraRunnable` thread by calling `signal()`.
  - A withdraw thread
    - loses a chance to withdraw money even if some money in the account.
    - may never be waked up if no threads call deposit afterward.

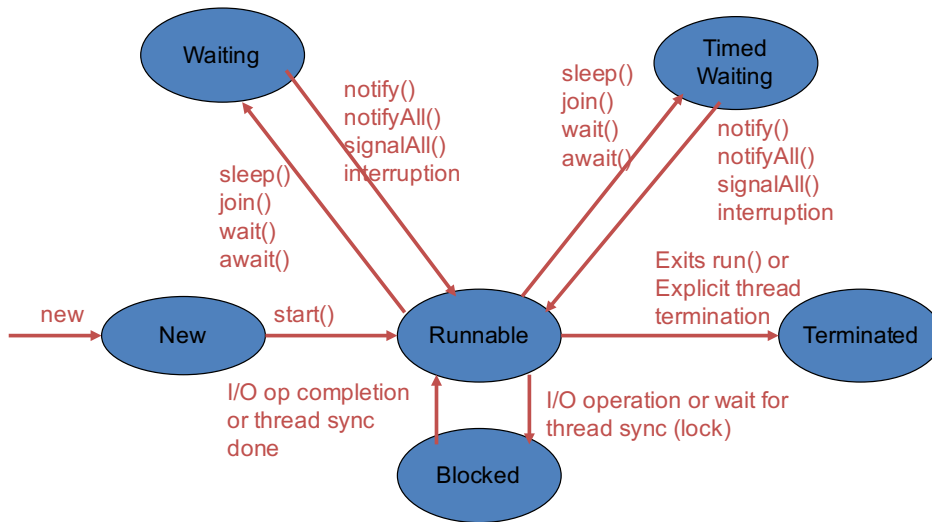


## signalAll()

- A deposit thread wake up both a withdraw thread and an `ExtraRunnable` thread.
  - Both threads go into the “Runnable” state.
- If the `ExtraRunnable` thread acquires the lock associated with `aCondition`,
  - the withdraw thread
    - goes to the “blocked” state.
    - acquires the target lock when the `ExtraRunnable` releases it.
    - will wake up sooner or later.
    - never lose a chance to withdraw money.



# ThreadSafeBankAccount2



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- ```

Condition sufficientFundsCondition = lock.newCondition();
Condition belowUpperLimitFundsCondition = lock.newCondition();

```
- ```

withdraw(double amount){
 lock.lock();
 while(balance <= 0){
 // waiting for the balance to exceed 0
 sufficientFundsCondition.await();
 }
 balance -= amount;
 belowUpperLimitFundsCondition.signalAll();
 lock.unlock();
}

```
- ```

deposit(double amount){
    lock.lock();
    while(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await();
    }
    balance += amount;
    sufficientFundsCondition.signalAll();
    lock.unlock();
}

```

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“while” or “if” to Surround await()?

- ```

withdraw(double amount){
 lock.lock();
 while(balance <= 0){
 // waiting for the balance to exceed 0
 sufficientFundsCondition.await();
 }
 balance -= amount;
 belowUpperLimitFundsCondition.signalAll();
 lock.unlock();
}

```
- ```

withdraw(double amount){
    lock.lock();
    if(balance <= 0){
        // waiting for the balance to exceed 0
        sufficientFundsCondition.await();
    }
    balance -= amount;
    belowUpperLimitFundsCondition.signalAll();
    lock.unlock();
}

```
- “while” should be used rather than “if” when multiple threads call withdraw() concurrently. Why?

Problem

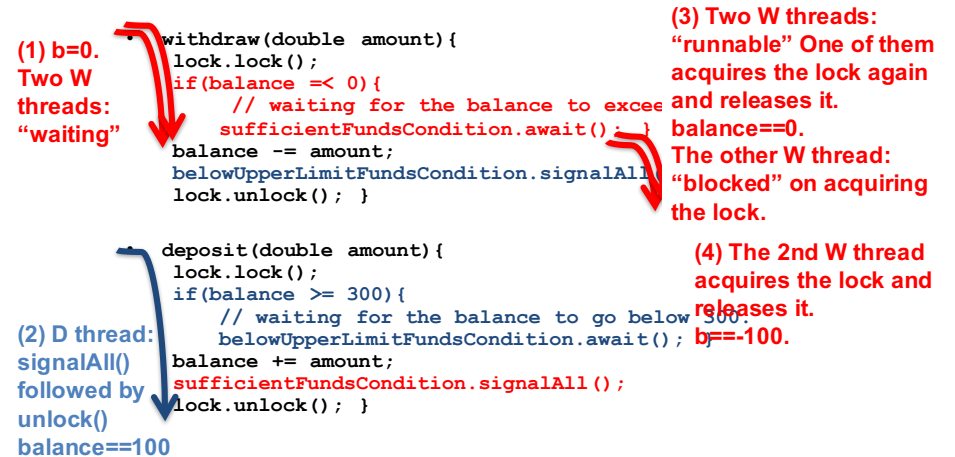
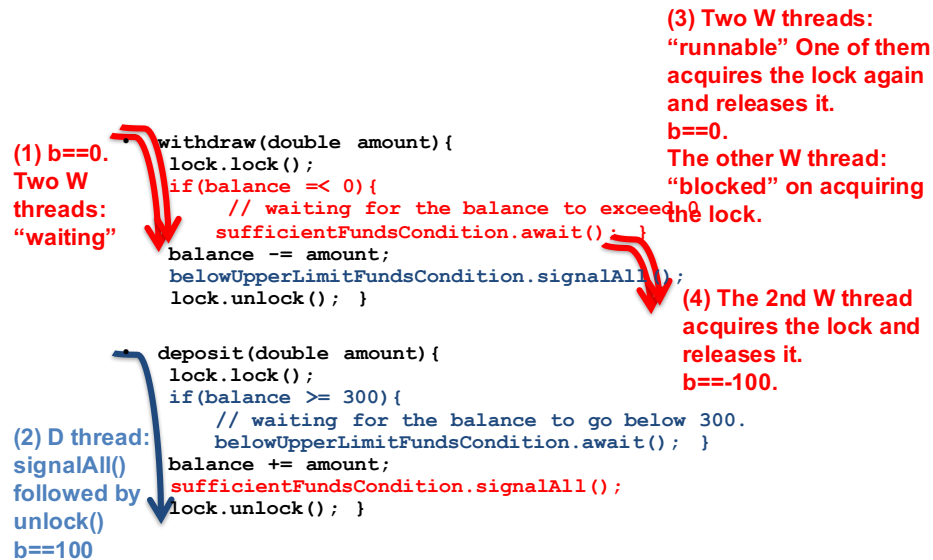
- (1) $b == 0$. Two W threads: “waiting”
- ```

withdraw(double amount){
 lock.lock();
 if(balance <= 0){
 // waiting for the balance to exceed 0
 sufficientFundsCondition.await();
 }
 balance -= amount;
 belowUpperLimitFundsCondition.signalAll();
 lock.unlock();
}

```
- (2) D thread: signalAll() followed by unlock()  $b == 100$
- ```

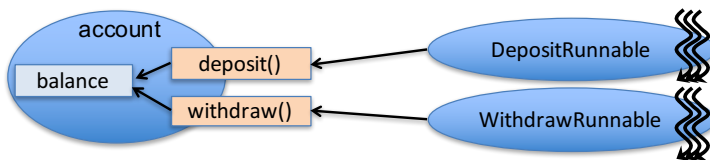
deposit(double amount){
    lock.lock();
    if(balance >= 300){
        // waiting for the balance to go below 300.
        belowUpperLimitFundsCondition.await();
    }
    balance += amount;
    sufficientFundsCondition.signalAll();
    lock.unlock();
}

```
- (3) Two W threads: “runnable” One of them acquires the lock again and releases it. $b == 0$. The other W thread: “blocked” on acquiring the lock.



- The 2nd thread should have made sure " $balance > 0$."
- If only one "W" thread runs, this problem does not occur.
- Always use a while loop.

ThreadSafeBankAccount2



"if" or "while" in Atomic Code?

- You can use "if" rather than "while" for a conditional checking
 - if you use signal(), not signalAll().
- However, in practice, the while-signalAll pair is more common than the if-signal pair.

InterruptedException

- Some methods in Java API can throw InterruptedException.
 - Thread.sleep()
 - Thread.join()
 - Condition.await()
 - ReentrantLock.tryLock()
- These methods can be long-running and cancellable.

Condition.await()

- await() lets the currently-executed thread to wait/sleep until another thread wakes it up with signal()/signalAll().
- interrupt() can interrupt a waiting/sleeping thread on a condition object.
 - The waiting/sleeping thread re-acquires a lock and throws an InterruptedException.
 - It does NOT immediately throw an InterruptedException.

```
withdraw(double amount){  
    lock.lock();  
    while(balance <= 0){  
        try{  
            // waiting for the balance to exceed 0  
            sufficientFundsCondition.await();  
        }catch(InterruptedException e){  
            //Do something  
        }  
    }  
    belowUpperLimitFundsCondition.signalAll();  
    balance -= amount;  
    lock.unlock(); }  

```

