# Spectral Graph Theory

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### 1 Introduction

#### 1.1 Adjacency Matrix Representation of A Graph

We represent a Graph G with a square adjacency matrix A, where entry

$$a_{ij} = \begin{cases} 1 & \text{if (i, j) is an edge in } G \\ 0 & \text{otherwise} \end{cases}$$

Evidently, the matrix is symmetric, or  $a_{ij} = a_{ji}$ . Therefore by spectral theorem, we have a orthonormal basis of eigenvectors, with real eigenvalues associated.

## 2 Counting Paths with Adjacency Matrix

**Theorem L.** et G be a graph on labeled vertices, let A be its adjacency matrix, and let k be a positive integer. Then  $A_{i,j}^k$  is equal to the number of walks from i to j that are of length k.

*Proof.* The proof is fairly simple, and we will do it by induction.

When k = 1, we look at the original adjacency matrix, and  $A_{i,j}$  indicates whether there's an edge between i, j, which is a path of length 1.

Now assume that the statement is true for k, and prove it for k+1.

Let's first think about it intuitively,  $A^k$  gives the number of paths walks from i to all other points. If one such point is v, then we just need to determine if there's and edge from v to j, if so then we just add the number of walks from i to k.

Let z be any vertex of G. If there are  $b_{i,z}$  walks of length k from i to z, and there are  $a_{z,j}$  walks of length one (in other words, edges) from z to j, then there are  $b_{i,z}a_{z,j}$  walks of

length k+1 from i to j whose next-to-last vertex is z. Therefore, the number of all walks of length k+1 from i to j is:

$$c(i,j) = \sum_{z \in G} b_{i,z} a_{z,j}$$

Since  $b_{i,z}$  correspond to an entry in  $A^k$ , the formula above is basically a matrix multiplication.

At this point, it's a good habit to check our proof again. Ask ourselves this question:

We know  $A_{i,j}^k$  represents some number of walks from i to j, but does it count all of them? We'll leave it as a quick thought exercise.

#### 2.1 Connectivity

**Theorem L.** et G be a simple graph on n vertices, and let A be the adjacency matrix of G. Then G is connected iff  $(I + A)^{n-1}$  consists of strictly positive entries.

*Proof.* We will only give the central idea of the proof here

A path from one point to another consists at most n vertices, or n-1 edges. Therefore, if we cannot find a path between two points within n-1 edges, the graph is not connected

Using the previous theorem, we know that  $A_{n-1}$  gives the number of paths if length n-1 between any two points, however, this is **not enough**.

Notice that the theorem indicates it's  $(I + A)^{n-1}$  instead of  $A^{n-1}$ . So what does the I do?

Graphically, it means that we assume any vertex is connected to itself. What is  $(I + A)^{n-1}$  then?

For simplicity, call (I + A) = A'. With  $A'_{i,i} = 1$ ,  $A'^k$  no longer counts the number of paths of length **strictly** k, but the number of paths of length  $\leq k$ . This is because we don't have to move to another point every time we multiply A', we can choose to stay at the same point since  $A'_{i,i} = 1$ .

Therefore  $(I+A)^{n-1}$  counts the number of paths of length  $\leq n-1$ , and if there still exist an 0 entry  $A_{i,j}^{m-1}$ , then it means if we cannot find a path between two points within n-1 edges, the graph is not connected.