

Computer Architecture
ELE 475 / COS 475
Slide Deck 4: Superscalar 1

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Types of Data Hazards

Consider executing a sequence of

$$r_k \leftarrow r_i \text{ op } r_j$$


type of instructions

Data-dependence

$$\begin{array}{l} r_3 \leftarrow r_1 \text{ op } r_2 \\ r_5 \leftarrow r_3 \text{ op } r_4 \end{array}$$


Read-after-Write
(RAW) hazard

Anti-dependence

$$\begin{array}{l} r_3 \leftarrow r_1 \text{ op } r_2 \\ r_1 \leftarrow r_4 \text{ op } r_5 \end{array}$$


Write-after-Read
(WAR) hazard

Output-dependence

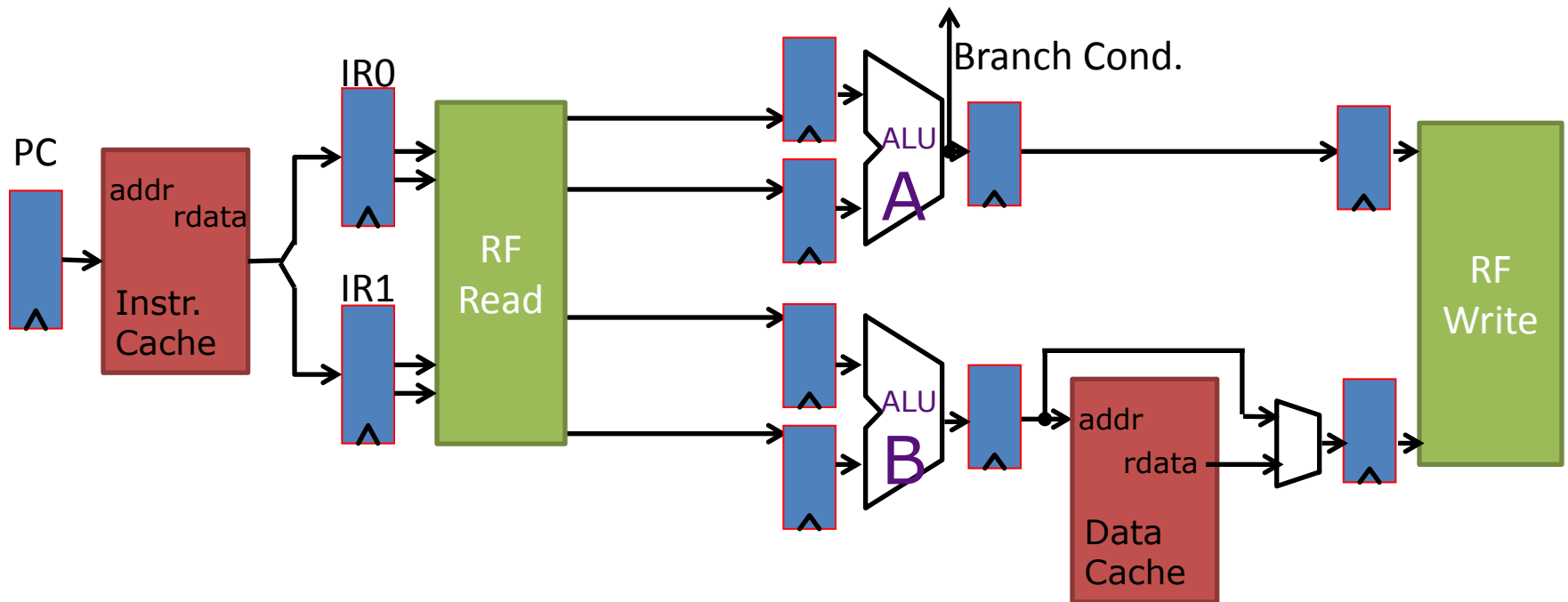
$$\begin{array}{l} r_3 \leftarrow r_1 \text{ op } r_2 \\ r_3 \leftarrow r_6 \text{ op } r_7 \end{array}$$


Write-after-Write
(WAW) hazard

Introduction to Superscalar Processor

- Processors studied so far are fundamentally limited to $CPI \geq 1$
- Superscalar processors enable $CPI < 1$ ($IPC > 1$) by executing multiple instructions in parallel
- Can have both in-order and out-of-order superscalar processors. We will start with in-order.

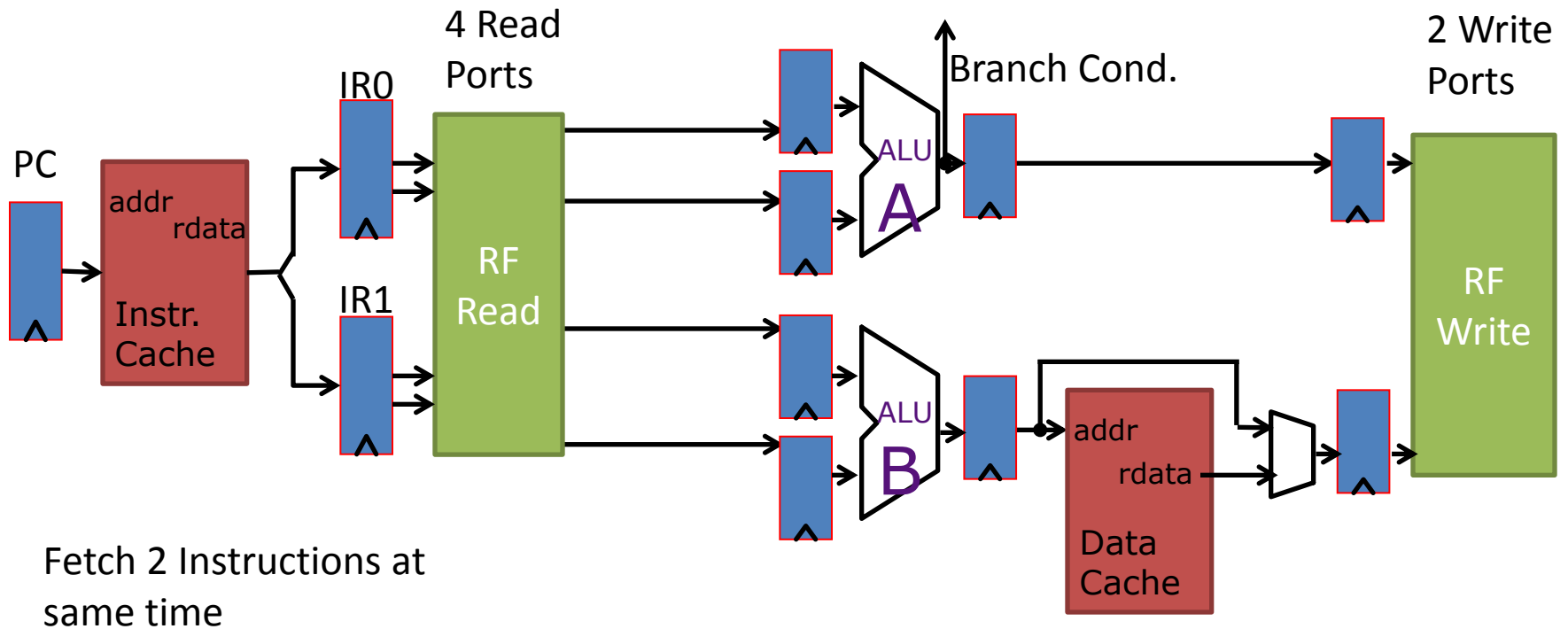
Baseline 2-Way In-Order Superscalar Processor



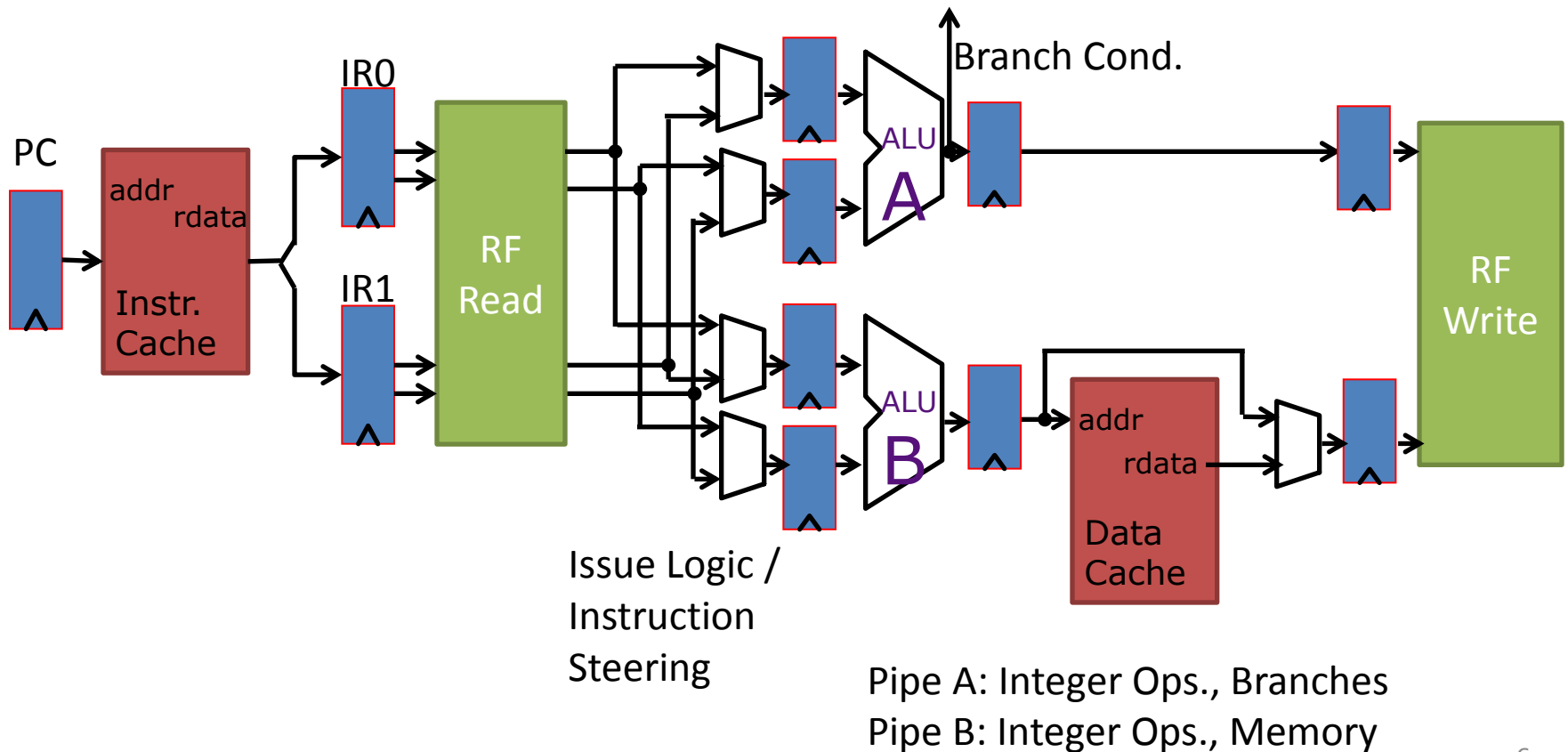
Pipe A: Integer Ops., Branches

Pipe B: Integer Ops., Memory

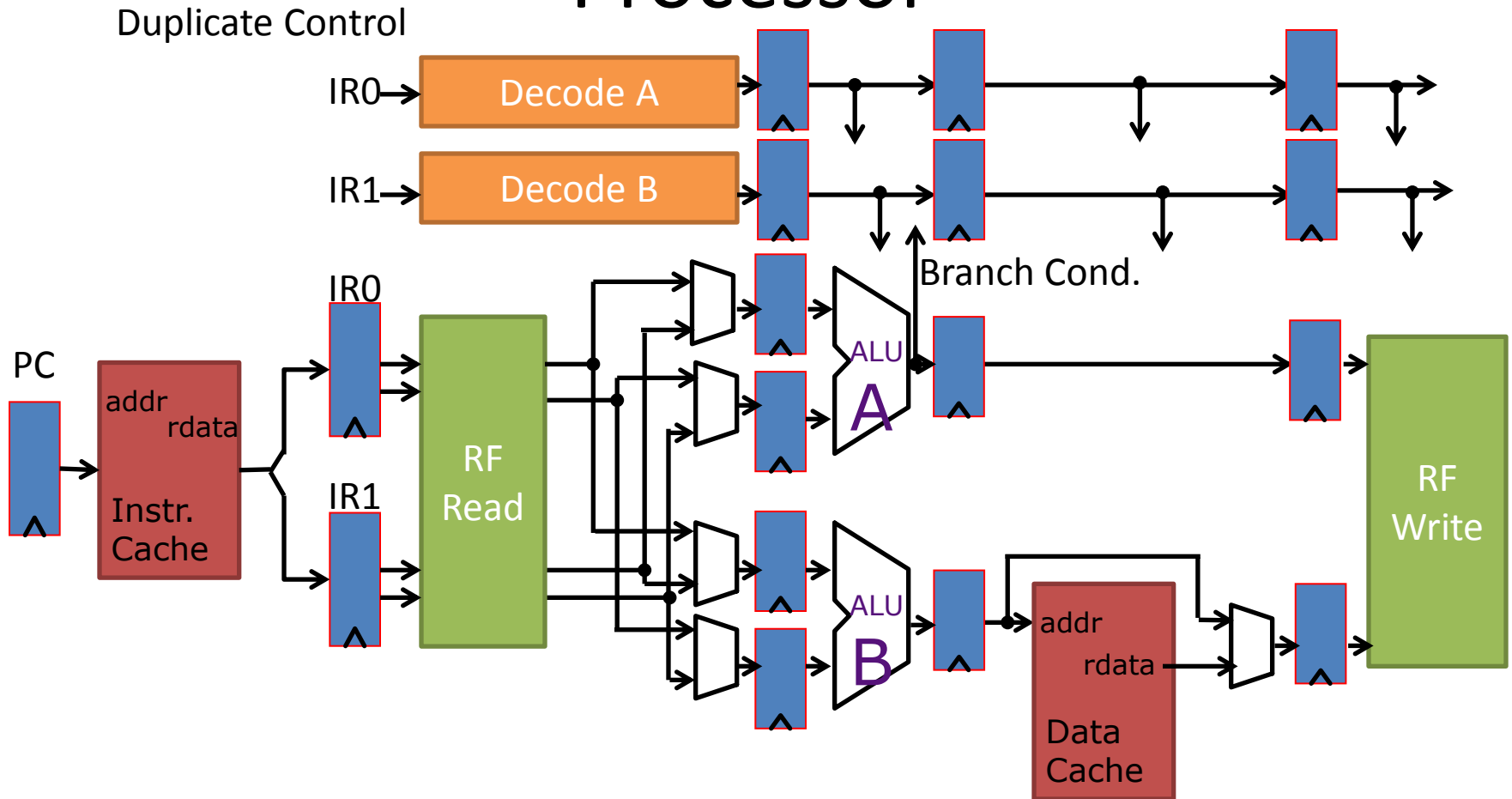
Baseline 2-Way In-Order Superscalar Processor



Baseline 2-Way In-Order Superscalar Processor



Baseline 2-Way In-Order Superscalar Processor



Pipe A: Integer Ops., Branches
Pipe B: Integer Ops., Memory

Issue Logic Pipeline Diagrams

OpA	F	D	A0	A1	W		
OpB	F	D	B0	B1	W		
OpC		F	D	A0	A1	W	
OpD		F	D	B0	B1	W	
OpE			F	D	A0	A1	W
OpF			F	D	B0	B1	W

CPI = 0.5 (IPC = 2)

Double Issue Pipeline
Can have two instructions in same stage at same time

ADDIU	F	D	A0	A1	W		
LW	F	D	B0	B1	W		
LW		F	D	B0	B1	W	
ADDIU		F	D	A0	A1	W	
LW			F	D	B0	B1	W
LW			F	D	D	B0	B

Instruction Issue Logic swaps from natural position

Structural Hazard

Dual Issue Data Hazards

No Bypassing:

ADDIU	R1,R1,1	F	D	A0	A1	W			
ADDIU	R3,R4,1	F	D	B0	B1	W			
ADDIU	R5,R6,1		F	D	A0	A1	W		
ADDIU	R7,R5,1		F	D	D	D	D	A0	A1 W

Full Bypassing:

ADDIU	R1,R1,1	F	D	A0	A1	W			
ADDIU	R3,R4,1	F	D	B0	B1	W			
ADDIU	R5,R6,1		F	D	A0	A1	W		
ADDIU	R7,R5,1		F	D	D	A0	A1	W	

Dual Issue Data Hazards

Order Matters:

ADDIU	R1,R1,1	F	D	A0	A1	W
ADDIU	R3,R4,1	F	D	B0	B1	W
ADDIU	R7,R5,1		F	D	A0	A1 W
ADDIU	R5,R6,1		F	D	B0	B1 W

WAR Hazard Possible?

Fetch Logic and Alignment

Cyc	Addr	Instr
0	0x000	OpA
0	0x004	OpB
1	0x008	OpC
1	0x00C	J 0x100
...		
2	0x100	OpD
2	0x104	J 0x204
...		
3	0x204	OpE
3	0x208	J 0x30C
...		
4	0x30C	OpF
4	0x310	OpG
5	0x314	OpH

0x000	0	0	1	1
...				
0x100	2	2		
...				
0x200		3	3	
...				
0x300				4
0x310	4	5		

Fetching across cache Lines is
very hard. May need extra ports

Fetch Logic and Alignment

Cyc Addr Instr

```

0   0x000 OpA
0   0x004 OpB
1   0x008 OpC
1   0x00C J 0x100
...
2   0x100 OpD
2   0x104 J 0x204
...
3   0x204 OpE
3   0x208 J 0x30C
...
4   0x30C OpF
4   0x310 OpG
5   0x314 OpH
    
```




Ideal, No Alignment Constraints

```

OpA F  D  A0 A1 W
OpB F  D  B0 B1 W
OpC   F  D  B0 B1 W
J     F  D  A0 A1 W
OpD   F  D  B0 B1 W
J     F  D  A0 A1 W
OpE   F  D  B0 B1 W
J     F  D  A0 A1 W
OpF   F  D  A0 A1 W
OpG   F  D  B0 B1 W
OpH   F  D  A0 A1 W
    
```

With Alignment Constraints

Cyc	Addr	Instr
?	0x000	OpA
?	0x004	OpB
?	0x008	OpC
?	0x00C	J 0x100
...		
?	0x100	OpD
?	0x104	J 0x204
...		
?	0x204	OpE
?	0x208	J 0x30C
...		
?	0x30C	OpF
?	0x310	OpG
?	0x314	OpH

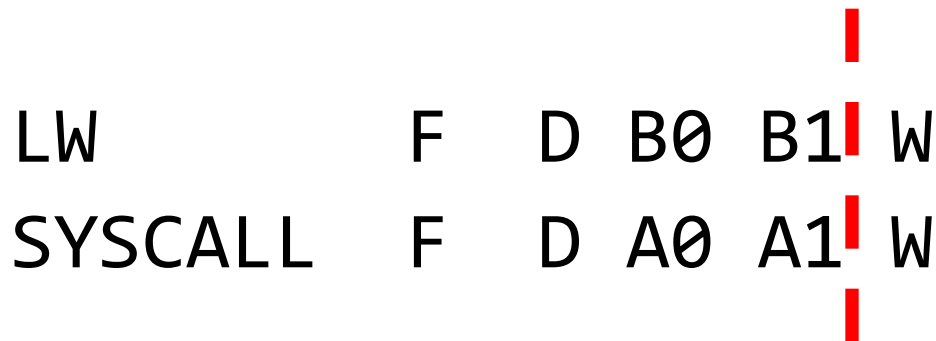
0x000	0	0	1	1
...				
0x100	2	2		
...				
0x200	3 	3	4	4 
...				
0x300			5 	5
0x310	6	6		

With Alignment Constraints

Cyc	Addr	Instr											
1	0x000	OpA	F	D	A0	A1	W						
1	0x004	OpB	F	D	B0	B1	W						
2	0x008	OpC		F	D	B0	B1	W					
2	0x00C	J 0x100		F	D	A0	A1	W					
3	0x100	OpD			F	D	B0	B1	W				
3	0x104	J 0x204			F	D	A0	A1	W				
4	0x200	?				F	-	-	-	-			
4	0x204	OpE				F	D	A0	A1	W			
5	0x208	J 0x30C					F	D	A0	A1	W		
5	0x20C	?					F	-	-	-	-		
6	0x308	?						F	-	-	-	-	
6	0x30C	OpF						F	D	A0	A1	W	
7	0x310	OpG							F	D	A0	A1	W
7	0x314	OpH							F	D	B0	B1	W

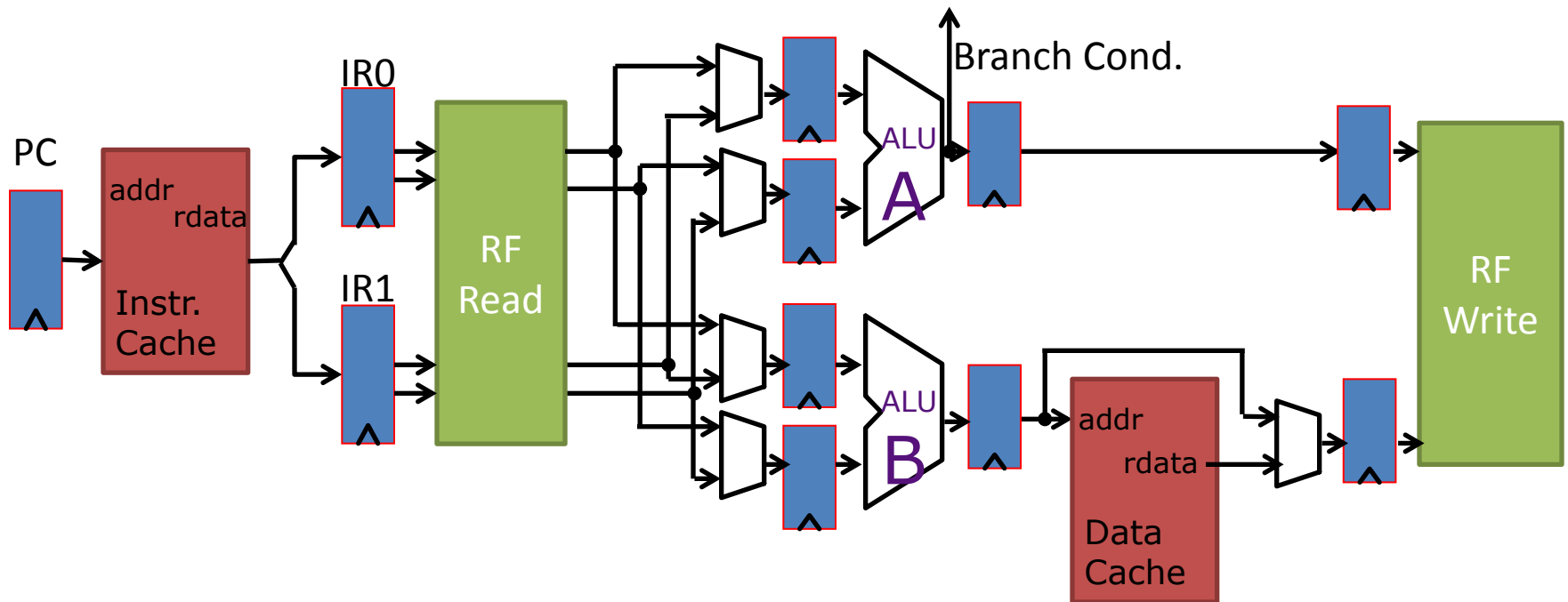
Precise Exceptions and Superscalars

- Similar to tracking program order for data dependencies, we need to track order for exceptions

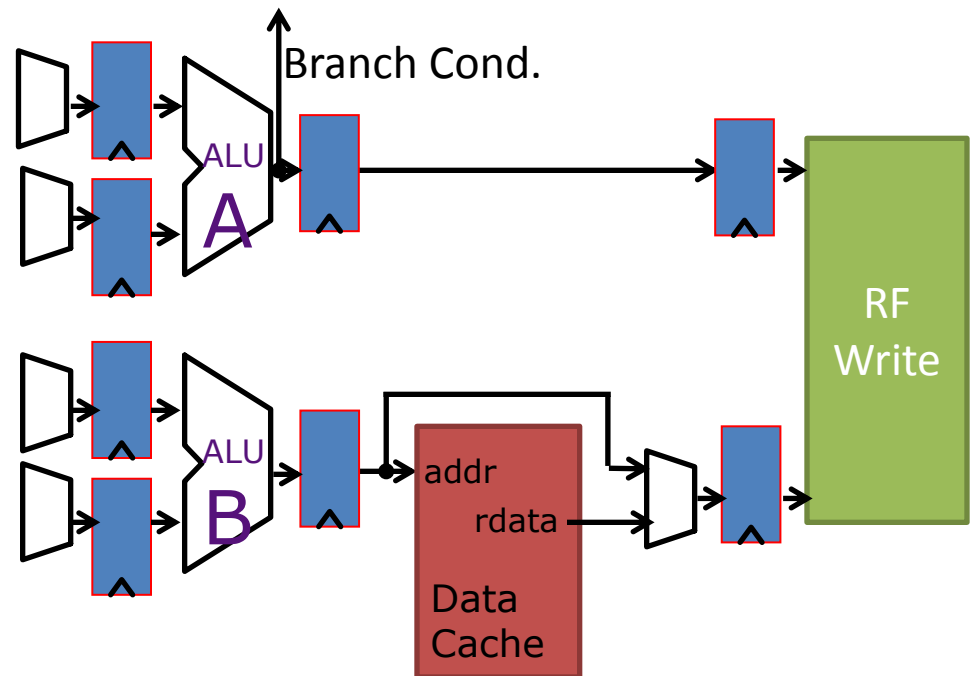


LW is in B pipeline, but commits first in logical order!

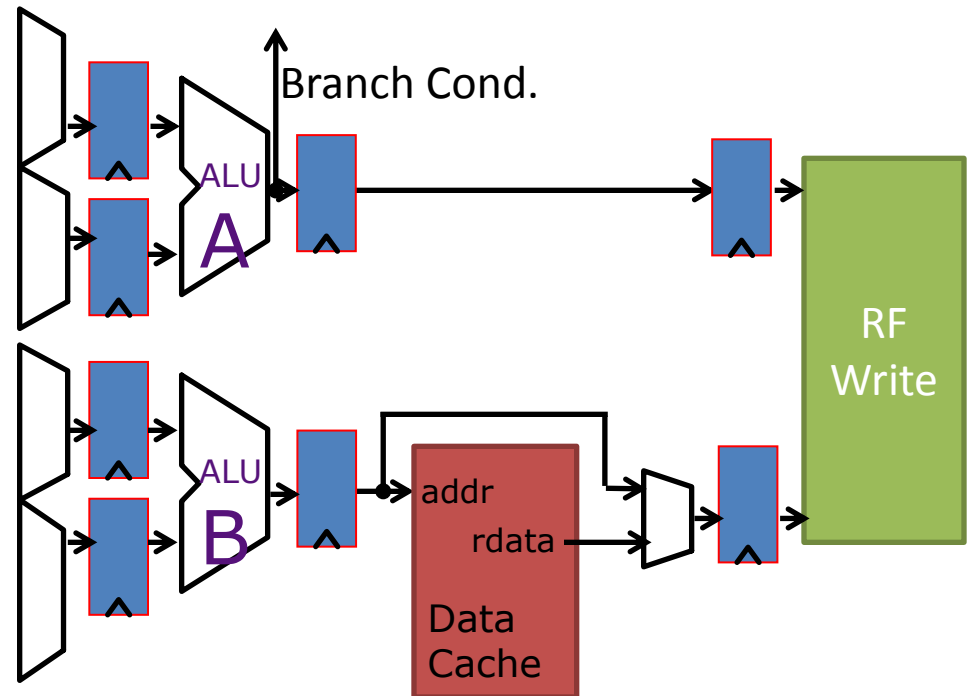
Bypassing in Superscalar Pipelines



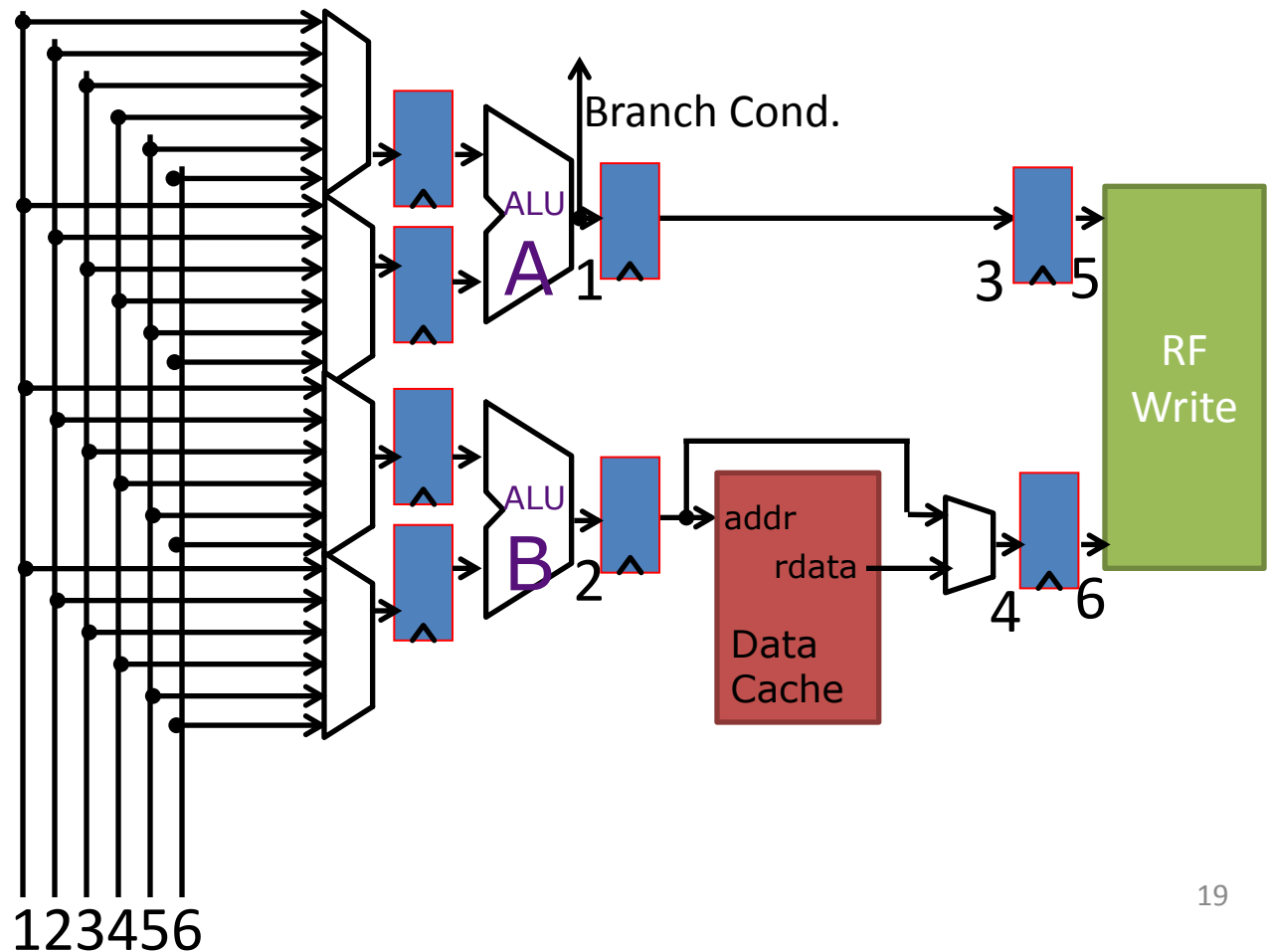
Bypassing in Superscalar Pipelines



Bypassing in Superscalar Pipelines



Bypassing in Superscalar Pipelines



Breaking Decode and Issue Stage

- Bypass Network can become very complex
- Can motivate breaking Decode and Issue Stage

D = Decode, Possibly resolve structural Hazards

I = Register file read, Bypassing, Issue/Steer
Instructions to proper unit

OpA	F	D	I	A0	A1	W	
OpB	F	D	I	B0	B1	W	
OpC		F	D	I	A0	A1	W
OpD		F	D	I	B0	B1	W

Superscalars Multiply Branch Cost

BEQZ	F	D	I	A0	A1	W				
OpA	F	D	I	B0	-	-				
OpB		F	D	I	-	-	-			
OpC		F	D	I	-	-	-			
OpD			F	D	-	-	-	-		
OpE			F	D	-	-	-	-		
OpF				F	-	-	-	-	-	
OpG				F	-	-	-	-	-	
OpH					F	D	I	A0	A1	W
OpI					F	D	I	B0	B1	W

Acknowledgements

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