

G and P and K, Oh My!

- Most fast adders “G”enerate, “P”ropagate, or “K”ill the carry
 - Usually only G and P are used; K only appears in some carry chains
- When does a bit Generate a carry out?
 - $G_i = A_i \text{ AND } B_i$
 - If G_i is true, then $\text{Cout}_i = C_{i+1}$ is forced to be true
- When does a bit Propagate a carry in to the carry out?
 - $P_i = A_i \text{ XOR } B_i$
 - If P_i is true, then $\text{Cout}_i (=C_{i+1})$ follows C_i
 - Usually implemented as $P_i = A_i \text{ OR } B_i$
 - OR is cheaper/faster than an XOR
 - If you are doing logic, Cout is still equal to $G_i + P_i C_i$
 - Just beware that $\text{Sum}_i \neq P_i \text{ XOR } C_i$

Using G and P

- We can combine G_i and P_i into larger blocks
 - Call these “group generate” and “group propagate” terms
- When does a group Generate a carry out? (e.g., 4 bits)
 - $G_{3:0} = G_3 + P_3G_2 + P_3P_2G_1 + P_3P_2P_1G_0$
- When does a group Propagate a carry in to the carry out?
 - $P_{3:0} = P_3P_2P_1P_0$
- We can also combine groups of groups
 - $G_{i:j} = G_{i:k} + P_{i:k}G_{k-1:j}$
 - $P_{i:j} = P_{i:k}P_{k-1:j}$



Faster Carry Bypass (or Carry Skip) Adders

- We see the basic idea is to form multi-level carry chains
 - Break the bits into groups
 - Ripple the carry in each group, in parallel
 - Ripple the global carry across the groups
- How big should each group be? (N bits total, k bits per group)
 - If ripple time equals block skip time then delay = $2(N-1) + (N/k - 2)$
- Would groups of different sizes be faster? (yes)
 - Middle groups have longer to generate carry outs; should be larger
 - Early and late groups have ripples in critical path; should be shorter
 - Called “Variable Block Adders”