

## Introduction

A segment tree is a heap-like data structure that can be used for making update/query operations upon array intervals in logarithmical time. We define the segment trees for the interval  $[i, j]$  in the following recursive manner.

- The root node (first node in the array) node will hold the information for the interval  $[i, j]$ .
- If  $i < j$  the left and right children will hold the information for the interval  $[i, (i + j)/2]$  and  $[(i+j)/2 + 1, j]$ .

A segment tree is a cool data structure, primarily used for a range query. It is a height-balanced binary tree with a static structure. The nodes of the segment tree correspond to various intervals and can be augmented with appropriate information pertaining to those intervals.

## Need for Segment Trees

Consider the following problem -

Given an array A consisting of n integers. You need to support 2 types of queries.

1. Find the sum of all elements of the array from 1 to i for a given index i.
2. Update the ith element of the array from  $a[i]$  to x.

There can be two possible approaches to this problem -

1. Maintain the array A dynamically, so that updates can be done in  $O(1)$  but every query of the first type will take  $O(n)$  time in the worst case as it will require finding the sum of n elements. So if there are q queries and majority of them are of the first type, then the worst case of this algorithm will be  $O(nq)$ .
2. Maintain the array B dynamically, where  $B[i]$  stores the sum of all elements from  $A[0]$  to  $A[i]$ . Now each query of the first type can be done in  $O(1)$  but every query of the second type will take  $O(n)$  time in the worst case as it will require updating all elements of B from i to n. So if there are q queries and the majority of them are of the second type, then the worst case of this algorithm will be  $O(nq)$ .

We will now discuss how we can use segment trees to support both these type of queries with the total run time of  $O(q \log n)$  where q is the number of queries and each query takes  $O(\log n)$  time.

In short problems where we need to maintain some information with updates, segment trees come to the rescue.