

# Introduction to Automata Theory

# Purpose and motivation

- What are the mathematical properties of computer hardware and software?
- What is a **computation** and what is an **algorithm**? Can we give rigorous mathematical definitions of these **notions**?
- What are the limitations of computers? Can “everything” be computed?

**Purpose:** Develop formal mathematical models of computation that reflect real-world computers.

# Alan Turing (1912-1954)

- Father of Modern Computer Science
- English mathematician
- Studied abstract machines called **Turing machines** even before computers existed



# Theory of Computation: A historical perspective

1930s	<ul style="list-style-type: none"><li>• Alan Turing studies <b>Turing machines</b></li><li>• <b>Decidability</b></li><li>• <b>Halting problem</b></li></ul>
1940-1950s	<ul style="list-style-type: none"><li>• “<b>Finite automata</b>” machines studied</li><li>• Noam Chomsky proposes the “<b>Chomsky Hierarchy</b>” for formal languages</li></ul>
1969	Cook introduces “intractable” problems or “ <b>NP-Hard</b> ” problems
1970-	Modern computer science: <b>compilers</b> , <b>computational &amp; complexity theory</b> evolve

# Theory of Computation, nowadays

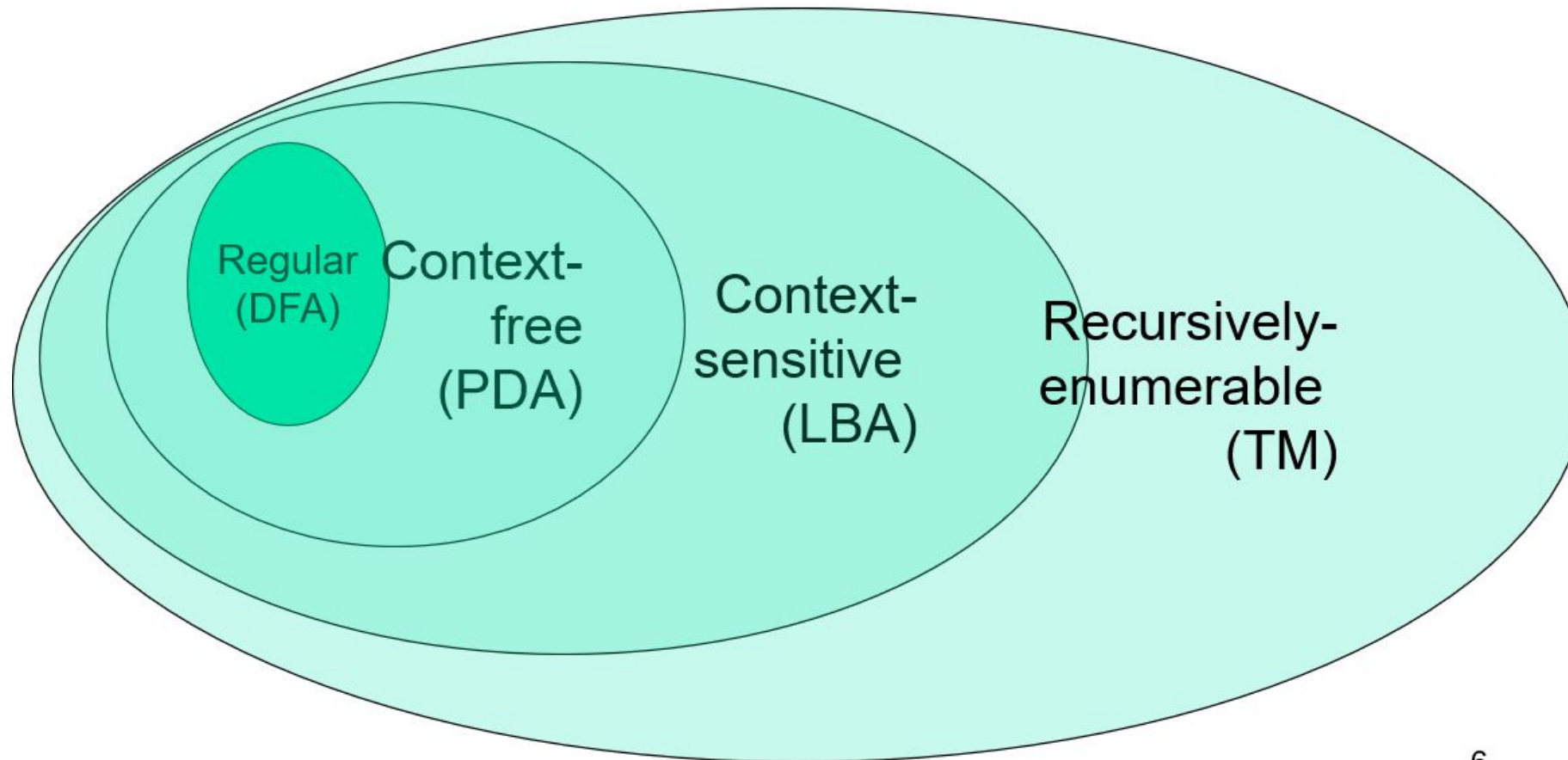
- Nowadays, the Theory of Computation can be divided into the following three areas:
- **Complexity theory:** What makes some problems computationally **hard** and other problems **easy**?
- **Computability theory:** Classify problems as being **solvable** or **unsolvable**.
- **Automata theory:** It deals with definitions and properties of different types of **computation models**.

# Computation models

- **Finite Automata:** These are used in text processing, compilers, and hardware design.
- **Context-Free Grammars:** These are used to define programming languages and in Artificial Intelligence.
- **Turing Machines:** These form a simple abstract model of a “real” computer, such as our PC at home.

# The Chomsky hierarchy

- A hierarchy of classes of formal languages



# Languages & Grammars

An **alphabet** is a set of symbols:

$\{0,1\}$

**Sentences** are strings of symbols:

0,1,00,01,10,1,...

A **language** is a set of sentences:

$L = \{000,0100,0010,..\}$

A **grammar** is a finite list of rules defining a language.

$S \longrightarrow 0A$

$B \longrightarrow 1B$

$A \longrightarrow 1A$

$B \longrightarrow 0F$

$A \longrightarrow 0B$

$F \longrightarrow \epsilon$

- Languages: "A language is a collection of sentences of finite length all constructed from a finite alphabet of symbols"
- Grammars: "A grammar can be regarded as a device that enumerates the sentences of a language" - nothing more, nothing less

*Reference: N. Chomsky,  
Information and Control, Vol 2,  
1959*



# Rudimentary Elements of Automata Theory

# Alphabet

*An alphabet is a finite, non-empty set of symbols*

- We use the symbol  $\Sigma$  (sigma) to denote an alphabet
- Examples:
  - Binary:  $\Sigma = \{0,1\}$
  - All lower case letters:  $\Sigma = \{a, b, c, ..z\}$
  - Alphanumeric:  $\Sigma = \{a-z, A-Z, 0-9\}$
  - DNA molecule letters:  $\Sigma = \{a, c, g, t\}$
  - ...

# Strings

*A string or word is a finite sequence of symbols chosen from  $\Sigma$*

- **Empty string is  $\epsilon$  (or “epsilon”)**
- Length of a string  $w$ , denoted by “ $|w|$ ”, is equal to the number of (non-  $\epsilon$ ) characters in the string
  - *E.g.,  $x = 010100$   $|x| = 6$*
  - *$x = 01 \epsilon 0 \epsilon 1 \epsilon 00 \epsilon$   $|x| = ?$*
  - *$xy = \text{concatenation}$  of two strings  $x$  and  $y$*

# Powers of an alphabet

Let  $\Sigma$  be an alphabet.

- $\Sigma^k$  = the set of all strings of length  $k$
- $\Sigma^* = \Sigma^0 \cup \Sigma^1 \cup \Sigma^2 \cup \dots$
- $\Sigma^+ = \Sigma^1 \cup \Sigma^2 \cup \Sigma^3 \cup \dots$

# Languages

***L is said to be a language over alphabet  $\Sigma$ , only if  $L \subseteq \Sigma^*$***

- this is because  $\Sigma^*$  is the set of all strings (of all possible length including 0) over the given alphabet  $\Sigma$

Examples:

1. Let L be *the* language of all strings consisting of  $n$  0's followed by  $n$  1's:  
 $L = \{\epsilon, 01, 0011, 000111, \dots\}$
2. Let L be *the* language of all strings of with equal number of 0's and 1's:  
 $L = \{\epsilon, 01, 10, 0011, 1100, 0101, 1010, 1001, \dots\}$

Canonical ordering of strings in the language

**Definition:  $\emptyset$  denotes the Empty language**

- Let  $L = \{\epsilon\}$ ; Is  $L = \emptyset$ ?

# The membership problem

*Given a string  $w \in \Sigma^*$  and a language  $L$  over  $\Sigma$ , decide whether or not  $w \in L$ .*

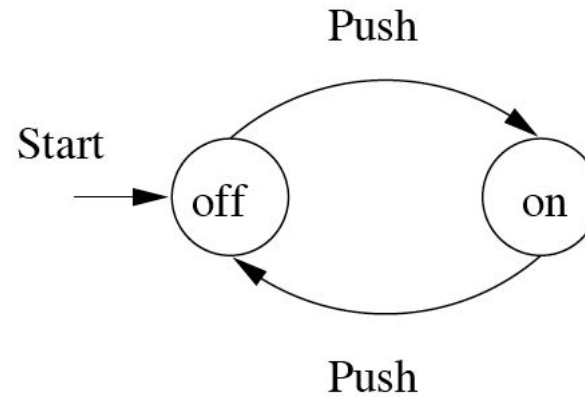
Example:

Let  $w = 100011$

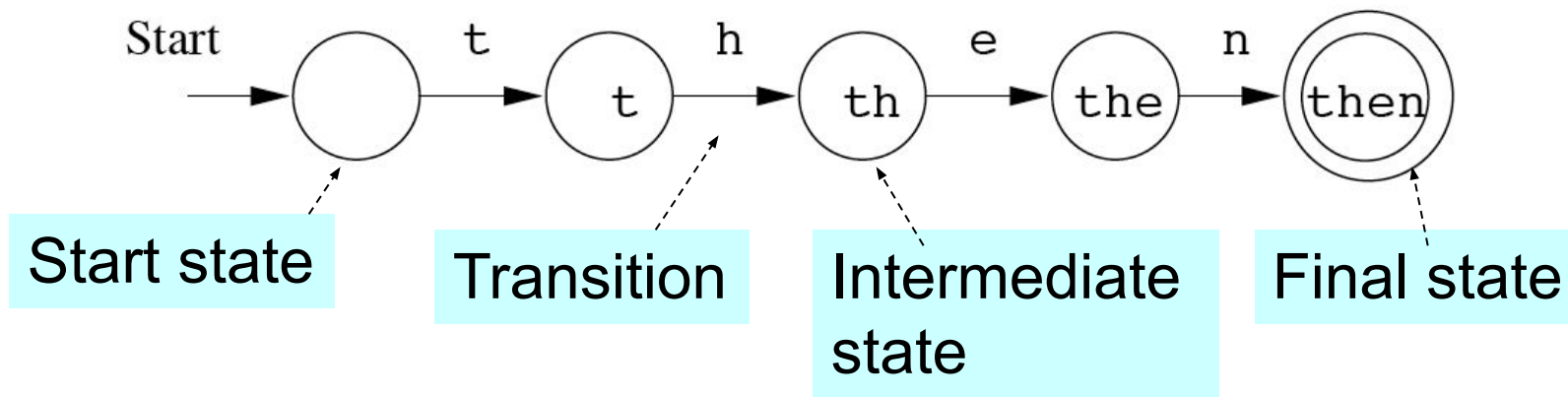
Q) Is  $w \in$  the language of strings with equal number of 0s and 1s?

# Finite Automata : Examples

- On/Off switch



- Modeling recognition of the word “then”



# Finite Automata

- Some Applications
  - Software for designing and checking the behavior of digital circuits.
  - Lexical analyzer of a typical compiler.
  - Software for scanning large bodies of text (e.g., web pages) for pattern finding.
  - Software for verifying systems of all types that have a finite number of states (e.g., stock market transaction, communication/network protocol).