Pumping Lemma for RE

Pumping Lemma (Fox Resular Languages)

- >> Pumping Lemma is used to prove that a longuage is NOT REGULAR.
- >> It cannot be used to prove that a Language is Regular.

If A is a Regular Language, then A has a pumping Length 'p' such that any stowing 's' where $|S| \ge P$ may be divided into 3 parts S = xyz such that the following Conditions must be force.

- 1) x y iz E A for every i>0
- 2) 141>0
- 3> 1xy1 < P

Pumping Lemma (Post Regular Expression)

> To prove that a language is not Resular wring pumping Lemma, follow the below steps:

(me prove using Contradiction)

- -> Assume that A is Regular
- > It has to have a pumping Length (say p)
- -> All strings longer than p can be pumped 151>P
- > Now find a string 'S' in A such that ISI>P
- > Dévide S into XYZ
- > Show that xy'z & A for some i
- > Then Consider all consider that 5 can be divided
- Show that none of these can satisfy all the 3 pamping conditions at the same time
- > S can not be pumped = Contradiction

Pumping Lemma (Fox Regular Languages)-Example => Unite pumping Lemma priore that the language A = Lory 1 nzof is not pegular priot -> Assume that Ais Regular -> It need to have a pumping length a, song p > Choose stants, S = app > Por example, P=7, >> S = aaaaaaa bbbbbbbb -> All the possible ways to devide S are as follows. Case1: The y is in the 'a' paul 7xy'z=> xy'z > xy'z agaaaaabbbbbbbb Case 2: The y is in the 'b' paul -> xy'z >> xy'z case 3: The y in in the a and b' pant > xy'z => xy'z Check IXYISP, P=7

Pumping Lemma (For RL) - Example

=> Using pumping Lemma priore that the language $A = \{yy \mid y \in \{0,1\}^{\#}\} \text{ is Not Resular}$

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Proof

> Assume that A i's Regular

-> pumping length say P

$$\Rightarrow S = 0^{p} \cdot 0^{p} \cdot 1$$

P=7

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Therefore: As not Resular