

# **CS 3101 Assignment 1**

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## Problem 1

### Question 1 :

a)  $T(n) = O(n \log_2 n)$ , for an array of size  $n$ .

b) Since a temporary array is created  $S(n) = O(n)$ .

### Question 2 :

$$n < \frac{Z}{L}$$

### Question 3 :

$Q(n) = \tau + 2Q(\frac{n}{2}) + \frac{\tau n}{Z}$ , Where  $\tau$  is the cost of a cache miss.

### Question 4 :

$$Q(n) = \tau + 2Q(\frac{n}{2}) + \frac{\tau n}{Z}$$

$$Q(n) = \tau + 2(\tau + 2Q(\frac{n}{2}) + \frac{\tau n}{Z}) + \frac{\tau n}{Z}$$

$Q(n) = \tau(2^i + 1 + \frac{in}{Z})$ , Where  $\tau$  is the cost of a cache miss and  $i$  is s.t

$$2^i = n.$$

$$Q(n) = O(n)$$

### Question 5 :

Merge Sort is not an optimal sorting method for modern desktop/laptop computers since for large data sets the algorithm has high memory usage and high cache complexity, as a result negating any advantage gained from the  $O(n \log_2 n)$  time complexity.

## Problem 2

### Question 1:

Access time:  $100 + S - 1$

Locality: Temporal

Miss Type: Cold

### Question 2:

Access time:  $\frac{100S}{16} + \frac{15S}{16}$

Locality: Spacial

Miss Type: Cold

### Question 3:

Access time:  $100S$

Locality: None

Miss Type: Cold & Capacity

### Question 4:

Access time:  $200S$

Locality: None

Miss Type: Cold & Conflict

### **Problem 3**

#### **Question 1:**

$Q(m, n) \approx \frac{2nmZ\tau}{L}$ , where  $\tau$  is the cost of a cache miss.

#### **Question 2:**

For large values of m,n there is a massive amount of capacity related cache misses which makes this a very slow algorithm, thus making it unsuitable for general use.

#### **Question 3:**

$Q(n, m) \approx \frac{Z}{L}((j - i) + (l - k) + (q - p))$  for  $\alpha > \frac{j - i + 1}{Z}$

#### **Question 4:**

I ran both methods for  $2^{17}$  and found that the second method was about 30 percent faster. I have included the results in results.txt.