COMP 431

Internet Services & Protocols

Applications & Application-Layer Protocols:

The Domain Name System

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Application-Layer Protocols Outline The architecture of distributed systems Client/Server computing Example client/server systems and their application-level protocols The World-Wide Web (HTTP) Reliable file transfer (FTP) E-mail (SMTP & POP) Internet Domain Name System (DNS) The programming model used in constructing distributed systems Socket programming

Application-Layer Protocols

The Domain Name System (DNS)

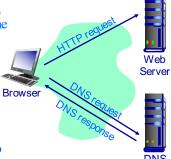
- Computers (hosts, routers) connected to the Internet have two forms of names:
 - » IP address a 32 bit identifier used for addressing hosts and routing data to them
 - » Hostname an ASCII string used by applications
- The DNS is an Internet-wide *service* that provides mappings between IP addresses and hostnames
 - » The DNS is a distributed database implemented in a hierarchy of name servers
 - » The DNS is also an application-layer protocol
- Hosts and routers use name servers to resolve names (address/name translation)
 - » Name resolution is an *essential* Internet function implemented as application-layer protocol

32 bit numbers are not easy to remember

The Domain Name System

Web browsing (HTTP) example

- The DNS is mainly used by applications, not end-users
 - » And virtually all applications use the DNS for every request they generate
- Web browsing: User enters URL www.someSchool.edu
 - » In order to create the socket to www.someSchool.edu, the OS (TCP) must resolve the hostname to an IP address
 - » The OS contacts a DNS name server to learn the web server's IP address
 - » The IP address is then used by TCP to create the socket to the server
 - » All this happens transparently to the user and the browser!

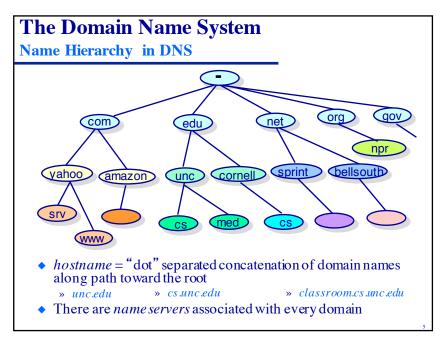


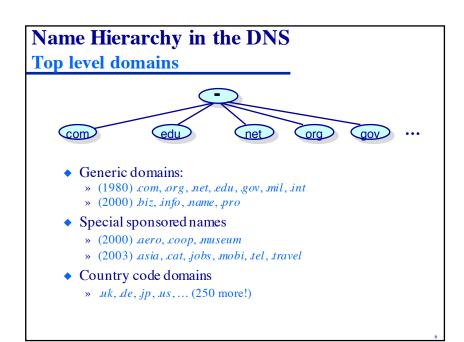
DNS - domain name system

used to map domain names to ip addresses

decoupling allows the ip address and server to be changed whenever

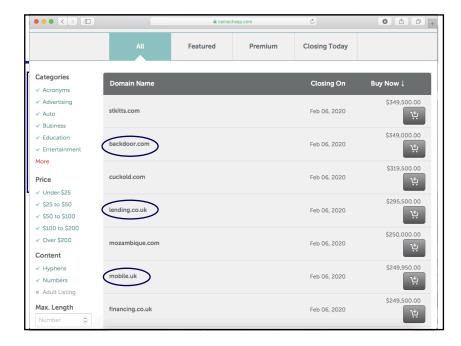
whenever it is changed UNC or the entity would only have to go to its dns database and change the mappings

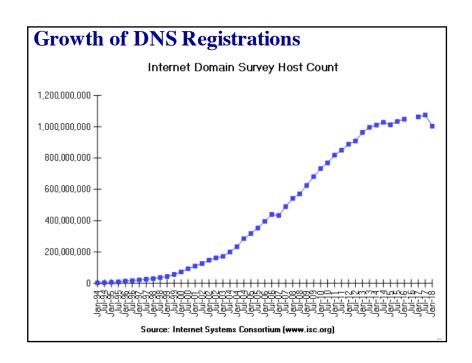




Can be the same machine but just two qualitatively different servers







Designing a distributed service



- » A server process on a big, well connected supercomputer?
- Centralized systems do not scale!
 - » Poor reliability: centralized = single point of failure
 - » Poor performance: centralized = "remote access" for most users
 - » Difficult to manage: centralized = all customer traffic goes to one location, a large staff has to be present to handle registrations
- A centralized system is not politically feasible in an international network

Designing a Distributed Service

DNS Name Servers

- No server has every hostname-to-IP address mapping
- Authoritative name server:
 - » Every host is registered with at least one authoritative server that stores that host's IP address and name
 - » The authoritative name server can perform name/address translation for that host's name/address
- *Local* authoritative name servers:
 - » Each ISP, university, company, has a local (default) name server authoritative for its own hosts
 - » Resolvers always query a name server local to it to resolve any host name

Local name server bristol.cs.unc.edu



Name resolution: Query and Reply



Local host

classroom.cs.unc.edu

What if the name is not a local host (e.g.,

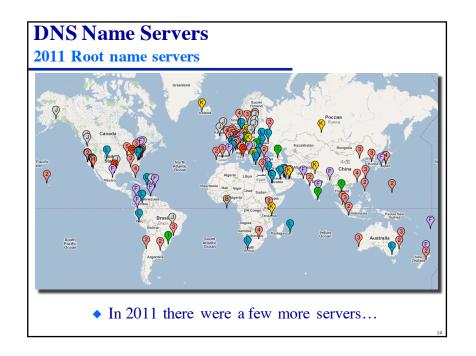
www.yahoo.com)?

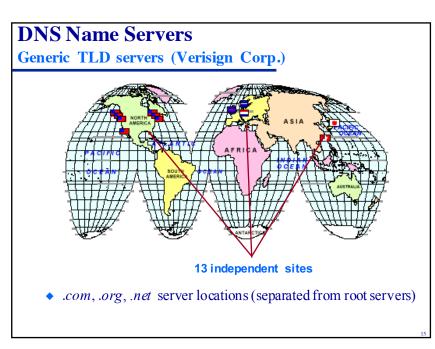
DNS Name Servers

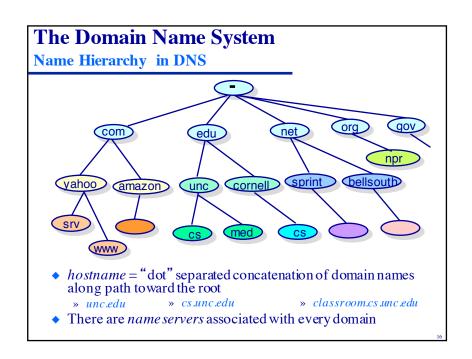
Root name servers

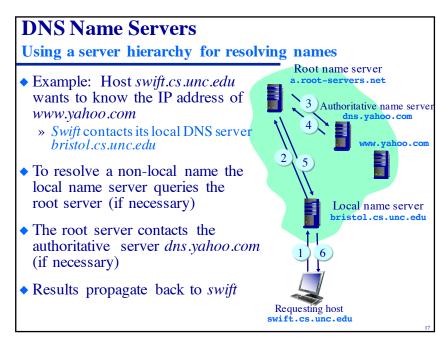


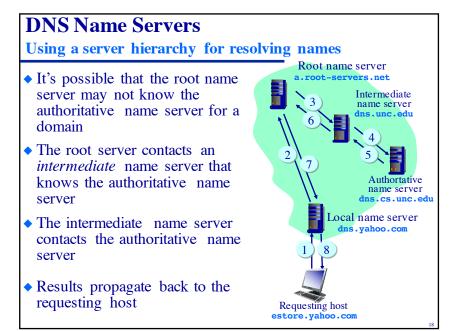
- A root name server is contacted when a local name server that can' resolve a name
 - » The root server either resolves the name or provides pointers to authoritative servers at lower level of name hierarchy
- In 1998, there were a dozen root name servers worldwide

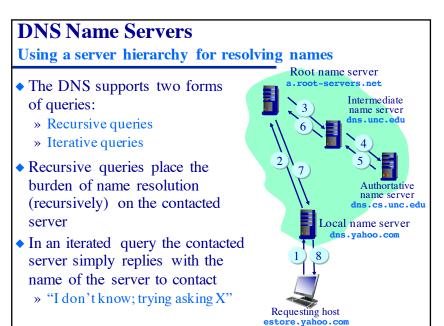


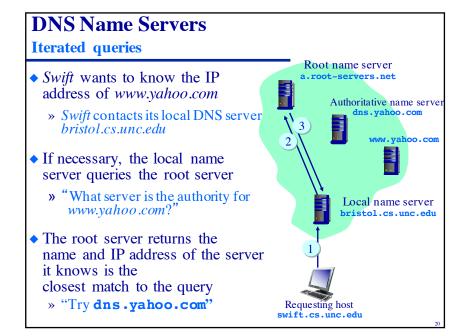


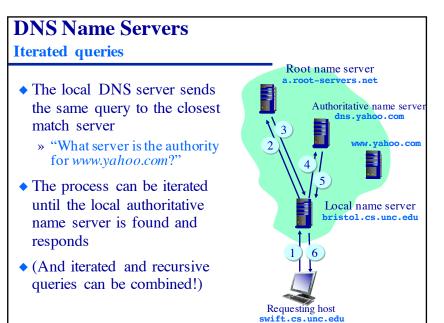








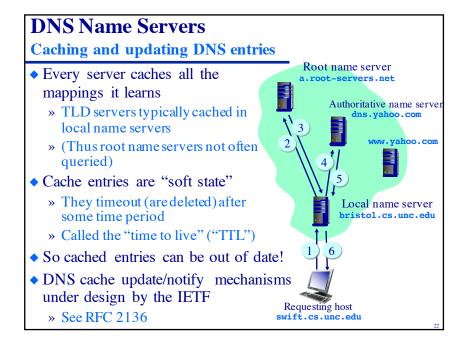




More load on local server

Local server can cache

response time includes network latency and server load



DNS Name Servers

DNS resource records

RR format: <name, value, type, time_to_live>

- The DNS is a distributed database storing resource records (RRs)
- Type = A
 - » name is a hostname
 - » value is hostname's IP address
- Type = CNAME
 - » name is an alias name for some "canonical" (the real) name
 - » value is canonical name

- Type = NS
 - » name is a domain
 - » value is name of authoritative name server for this domain
- Type = MX
 - » value is name of mail server host associated with name

Inserting records into the DNS

- Example: New startup "Network Utopia"
- Register name networkuptopia.com at DNS registrar (e.g., Network Solutions)
 - » You provide names & IP addresses of authoritative name server (primary and secondary)
 - » The registrar inserts two RRs into .com TLD server:
 - networkutopia.com, dns1.networkutopia.com, NS
 - * dns1.networkutopia.com, 212.212.212.1, A
- ◆ You stand up *dns1 networkutopia.com* running BIND and create:
 - » Authoritative server type A record for www.networkuptopia.com
 - » MX record for networkutopia.com



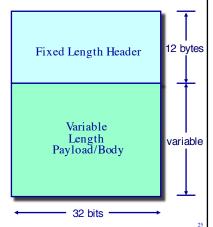
dns1.networkutopia.com

www.networkutopia.com

The Domain Name System

The DNS protocol

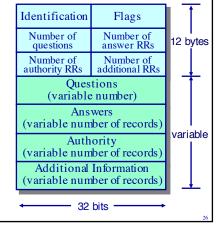
- ◆ The DNS service is implemented by the DNS protocol
- ◆ A request/response protocol run on top of UDP
 - » Uses port 53
- ◆ Why UDP?!
 - » Doesn't reliability matter?!



DNS Protocol

DNS query and reply messages

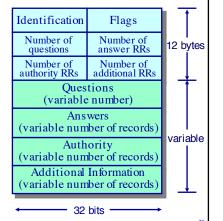
- DNS *query* and *reply* messages both have the same message format
- Messages have a fixed length message header
 - » Identification 16 bit query/reply identifier used to match replies to queries
 - » Flags:
 - Query/Reply bit
 - * "Reply is authoritative" bit
 - * "Recursion desired" bit
 - **....**



DNS Protocol

DNS query and reply messages

- Messages have a variable-length "question & answer" body
- Questions:
 - » The name and type fields (type A or MX) for a query hotmail.com MX
- Answers:
 - » One RR for each IP address answering query
- Authority:
 - » Resource records of other authoritative servers



DNS Resource Records

nslookup query/reply message example

```
(parris) 101> nslookup
> set debug
> www.yahoo.com
Server: bristol.cs.unc.edu
Address: 152.2.131.228

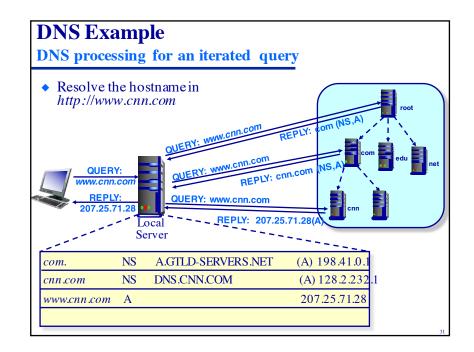
QUESTIONS:
    www.yahoo.com, type = A, class = IN
```

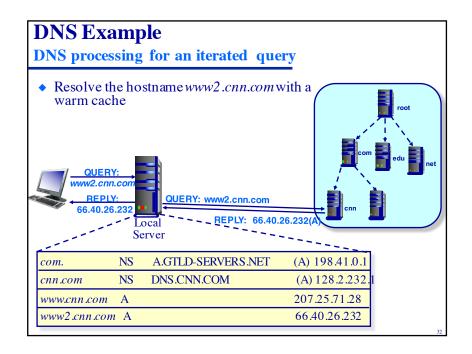
DNS Resource Records

nslookup query/reply message example

```
ANSWERS:
-> www.yahoo.com
canonical name = www.yahoo-ht3.akadns.net
-> www.yahoo-ht3.akadns.net
    internet address = 69.147.114.210
AUTHORITY RECORDS:
-> akadns.net
    nameserver = zc.akadns.org.
-> akadns.net
    nameserver = zd.akadns.org.
-> akadns.net
    nameserver = eurl.akadns.net.
-> akadns.net
    nameserver = use3.akadns.net.
-> akadns.net
    nameserver = use4.akadns.net.
-> akadns.net
    nameserver = usw2.akadns.net.
-> akadns.net
    nameserver = asia9.akadns.net.
-> akadns.net
    nameserver = za.akadns.org.
-> akadns.net
    nameserver = zb.akadns.org.
```

DNS Resource Records nslookup query/reply message example ADDITIONAL RECORDS: -> za.akadns.org internet address = 195.219.3.169 -> zb.akadns.org internet address = 206.132.100.105 -> zc.akadns.org internet address = 124.211.40.4 -> zd.akadns.org internet address = 63.209.3.132-> eurl.akadns.net internet address = 213.254.204.197 -> use3.akadns.net internet address = 204.2.178.133 -> use4.akadns.net internet address = 208.44.108.137 -> usw2.akadns.net internet address = 63.209.3.132 -> asia9.akadns.net internet address = 220.73.220.4 Non-authoritative answer: www.yahoo.com canonical name = www.yahoo-ht3.akadns.net. Name: www.yahoo-ht3.akadns.net Address: 69.147.114.210





www.networkutopia.com

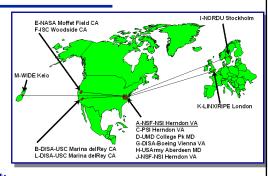
Attacking the DNS

- DDoS attacks: Bombard root servers with requests
 - » Not successful to date(!)
 - » Defeated by traffic filtering
 - » Local DNS servers cache IPs of TLD servers, allowing root server bypass
 - » Bombard TLD servers Potentially more dangerous
- ◆ Redirect attacks
 - » "Man-in-middle" (Intercept queries)
 - » DNS poisoning: Send bogus replies to a DNS server, which will cache them & return to others
- ◆ Exploit DNS for DDoS
 - » Send queries with spoofed source address!
 - » (Requires amplification)



Summary

- F gets 270,000,000+ hits per day
 - » Other servers have comparable load
- ◆ The Verisign TLD servers answer 5,000,000,000 queries per day
- Clearly the DNS would collapse without:
 - » Hierarchy
 - » Distributed processing
 - » Caching



• If DNS fails, Internet services stop working!