```
02:57:49.920864 IP 192.168.56.101.36286 > 192.168.56.103.sip-tls: Flags [P.], se q 34904:35597, ack 35219, win 1205, options [nop,nop,TS val 1514174486 ecr 15723 01], length 693  
0x0000: 4500 02e9 dda3 4000 4006 684e c0a8 3865 E....@.@.hN..8e
```

0x00000: 4500 02e9 dda3 4000 4006 684e c0a8 3865 E.....@.@.hN..8e 0x0010: c0a8 3867 8dbe 13c5 49bf 2142 37dd f161 ..8g....I.!B7..a 0x0020: 8018 04b5 d869 0000 0101 080a 5a40 7816i....Z@x.

>Attacking encrypted VoIP protocols

```
0x0060:
         dbe7 4a64 5737 b4e7 a3c6 1d5a 6c84 4fc3
                                                   ..JdW7....Zl.O.
         8daf 0422 31fc 8177 ddab 1835 5256 e08d
                                                   ..."1..w...5RV...
0x0070:
                                                   P.:....H.~n....
0x0080:
         50a0 3aa7 94ca 1c48 bd7e 6ed7 f907 ba95
                                                   y...:....%......
         791e f63a 9bb9 f084 9e25 a9db c01f 1da2
0x0090:
0x00a0:
         0814 b053 3587 6e07 5903 21e5 96be d44b
                                                   ...S5.n.Y.!....K
         56cb 1662 8a0c 7d2b f952 5933 ae51 71b0
                                                   V..b..}+.RY3.Qq.
0x00b0:
0x00c0:
         ef53 ad93 1653 f6bf eb04 19f1 55a6 abc5
                                                   .S...S.....U...
                                                   ].....sy....5
0x00d0:
         5d90 d49f cf19 1cc9 053c 79ad e019 ab35
                                                   ......P...*..W...
0x00e0:
         a9a3 abbe d1a4 c550 1211 2abb 9977 9b80
                                                   .{....S ..0.-a..
0x00f0:
         aa7b f38f f2da 535f 02a2 3020 2d67 0ae9
HAXPO Amsterdam, May 2019
                                                       Ivica Stipovic
```

Biography

- 1. Name>Ivica [Eeveetsa] Stipovic
- 2. Work>Ward Solutions, Dublin, Ireland
- 3. Job>Information Security Consultant
- 4. Contact> lvica.Stipovic@ward.ie
- 5. >EOF

What is VoIP protocol?

- VOIP stands for (Voice Over IP) Voice / Video/messaging that uses IP-based transport protocols for transmission
- What is SIP?

SIP stands for Session Initiation Protocol (it is a voice control protocol), developed by IETF

SIP is one of the predominant VOIP control protocols



Source: https://en.wikipedia.org/wiki/Session_Initiation_Protocol

SIP,MGCP,H.323,XMPP...

There are other VoIP protocols as well

- MGCP Media Gateway Control Protocol (MGCP), connection management for media gateways
- H.323 one of the first VoIP call signaling and control protocols that found widespread implementation
- XMPP Extensible Messaging and Presence Protocol, instant messaging, presence information, and contact list maintenance
- Skype protocol, proprietary Internet telephony protocol suite based on peerto-peer architecture

...etc

What applications are using VoIP?

This is a small snapshot of the most popular VoIP applications.

Not all of them use SIP, though: Skype uses proprietary protocol

Viber does use SIP, so as Blink, Cisco IP Communication, Jitsi, Bria, Empathy, Linphone, X-Lite, Zoiper, Yahoo! Messenger...



Best VolP Apps for Your Desktop

- Stay in Touch with These Voice over IP Apps. VoIP, or Voice over Intern Protocol, is the hip way to talk on the phone. ...
- Skype. Skype boasts a huge user base, with more than 300 million subscribers. ...
- · Google Voice. ...
- · Vonage. ...
- · 00V00. ...
- · Viber. ...

Let's talk about SIP

- SIP by default runs as a UDP service over the port 5060
- By default, it is not encrypted
- Full description of SIP protocol is given in the https://tools.ietf.org/html/rfc3261
- SIP structure is very similar to HTTP session structure (both request and response paradigm)

Structure of a SIP session

Similar to HTTP

Request (REGISTER vs GET)

Response (Unathorised | Return code 401)

Request (REGISTER vs GET)

```
REGISTER sip:192.168.0.116:5061 SIP/2.0
Call-ID: 2b0a036d62188707e8c1fa3f99d442eb@0:0:0:0:0:0:0:0
From: "demo-103" <sip:demo-103@192.168.0.116>;tag=fe23e8cf
To: "demo-103" <sip:demo-103@192.168.0.116>
Via: SIP/2.0/TES 192.168.0.164:3565;branch=z9hG4bK-373232-13573379b4dc2c69a
Max-Forwards: 70
User-Agent: Jitsi2.10.5550Windows 10
Expires: 600
Contact: "demo-103" <sip:demo-103@192.168.0.164:3565;transport=tls;register
Content-Length: 0
SIP/2.0 401 Unauthorized
Via: SIP/2.0/TLS
192.168.0.164:3565;branch=z9hG4bK-373232-13573379b4dc2c69aa6994b76a28d8af;r
From: "demo-103" <sip:demo-103@192.168.0.116>;tag=fe23e8cf
To: "demo-103" <sip:demo-103@192.168.0.116>;tag=as174e6496
Call-ID: 2b0a036d62188707e8c1fa3f99d442eb@0:0:0:0:0:0:0:0
CSea: 1 REGISTER
Server: Asterisk PBX 15.7.2
Allow: INVITE, ACK, CANCEL, OPTIONS, BYE, REFER, SUBSCRIBE, NOTIFY, INFO, F
Supported: replaces, timer
WWW-Authenticate: Digest algorithm=MD5, realm="asterisk", nonce="08e37204"
Content-Length: 0
REGISTER sip:192.168.0.116:5061 SIP/2.0
Call-ID: 2b0a036d62188707e8c1fa3f99d442eb@0:0:0:0:0:0:0:0
From: "demo-103" <sip:demo-103@192.168.0.116>;tag=fe23e8cf
To: "demo-103" <sip:demo-103@192.168.0.116>
Max-Forwards: 70
User-Agent: Jitsi2.10.5550Windows 10
```

Attributes of a SIP session

We will focus on attributes required to facilitate password attacks. These are: algorithm, method, username, realm, uri,response,nonce,{nc},{cnonce} and {qop}

```
SUBSCRIBE sip:demo-103@192.168.0.116 SIP/2.0
Call-ID: 7f8aa645e0eaa29432637916eacf4fac@0:0:0:0:0:0:0:0:0
CSeq: 2 SUBSCRIBE
From: "demo-103" <sip:demo-103@192.168.0.116>;tag=27f22122
To: "demo-103" <sip:demo-103@192.168.0.116>
Max-Forwards: 70
Contact: "demo-103" <sip:demo-103@192.168.0.164:3565;transport=tls;registering_acc=192_168_0_116>
User-Agent: Jitsi2.10.5550Windows 10
Event: presence.winfo
Accept: application/watcherinfo+xml
Expires: 3600
Via: SIP/2.0/TLS 192.168.0.164:3565;branch=z9hG4bK-373232-b62e3236d4729306c94af25ea7b34831
Authorization: Digest
username="demo-103",realm="asterisk",nonce="5ef0fb06",uri="sip:demo-103@192.168.0.116",response="a3
Content-Length: 0
```

Why TLS/encryption?

- wrapping SIP into TLS makes it more secure (HTTP vs HTTPS, POP3 vs POP3S, LDAP vs LDAPS etc.)
- interception of encrypted SIP will show obfuscated application layer payload
- only ip/tcp/udp header level information intelligible

Why TLS/encryption?

-challenges to pentesters that TLS presents – we want the application layer

255 38.962338	192.168.0.164	192.168.0.116	TPKT	215 Continuation
256 38.963635	192.168.0.116	192.168.0.164	TCP	60 5061 → 2945 [ACK
257 38.966307	192.168.0.116	192.168.0.164	TCP	60 5061 → 2945 [ACK
258 38.966902	192.168.0.116	192.168.0.164	TCP	60 5061 → 2945 [ACK
265 39.041986	192.168.0.116	192.168.0.164	TPKT	636 Continuation
267 39.081645	192.168.0.164	192.168.0.116	TCP	54 2945 → 5061 [ACK
268 39.157116	192.168.0.116	192.168.0.164	TPKT	652 Continuation
269 39.196799	192.168.0.164	192.168.0.116	TCP	54 2945 → 5061 [ACK
294 40.865386	192.168.0.164	192.168.0.116	TPKT	443 Continuation
304 41.068180	192.168.0.116	192.168.0.164	TCP	60 5061 → 2945 [ACK

rame 294: 443 bytes on wire (3544 bits), 443 bytes captured (3544 bits) on interface 0

thernet II, Src: IntelCor_48:77:f9 (3c:6a:a7:48:77:f9), Dst: IntelCor_80:65:d3 (18:56:80:80:0)

Internet Protocol Version 4, Src: 192.168.0.164, Dst: 192.168.0.116

ransmission Control Protocol, Src Port: 2945, Dst Port: 5061, Seq: 13079, Ack: 6387, Len: 38:

TPKT - ISO on TCP - RFC1006

Continuation data: 1703030180ab648c3bc196077d8836337c44874b13a6a99c...

```
0 01 00 f4 9a 00 00 17 03 03 01 80 ab 64 8c 3b c1 ······d·; ·
0 96 07 7d 88 36 33 7c 44 87 4b 13 a6 a9 9c 56 f1 ···}·63|D ·K····V·
0 34 9d 9d ef 1b de 5d bf bc dc 55 d3 2f aa 02 cd 4·····] · ··U·/···
```

Lower OSI layers

Application layer obfuscated

SIP –two aspects of attacks

1st aspect : interception + decryption

- -for unencrypted SIP sessions, one focuses only on interception
- -in our case, we need to do both

2nd aspect : SIP password cracking

- -for unencrypted SIP sessions, off-the-shelf tools available in Kali
- -in our case, we need to develop either:
 - -manual preparation of a file for sipcrack
 - -our own tool to streamline the cracking (<= chosen approach)

Attacking plaintext SIP passwords

- -Kali with sipdump and sipcrack
- -sipdump takes a pcap of a SIP session as input and generates a text file output for sipcrack
- -sipcrack takes the text output from sipdump and performs password dictionary attack

```
root@kali:/projekti/mitm-relay# sipdump -h
SIPdump 0.2
Usage: sipdump [OPTIONS] <dump file>
       <dump file> = file where captured logins will be written to
       Options:
       -i <interface> = interface to listen on
       -p <file> = use pcap data file
-m = enter login data manually
       -f "<filter>" = set libpcap filter
* Invalid arguments
root@kali:/projekti/mitm-relay# sipcrack -h
SIPcrack 0.2
Usage: sipcrack [OPTIONS] [ -s | -w <wordlist> ] <dump file>
       <dump file> = file containing logins sniffed by SIPdump
       Options:
                    = use stdin for passwords
       -w wordlist = file containing all passwords to try
                     = print cracking process every n passwords (for -w)
       -p num
                       (ATTENTION: slows down heavily)
* Invalid arguments
root@kali:/projekti/mitm-relay#
```

Solution design-1st part (interception and decryption)

An idea of MITM occurred as one plausible attack vector

This is what we need to achieve:

- build a mechanism capable of intercepting and decrypting the TLS wrapped session
- search for some kind of protocol-agnostic proxy capable of decrypting TLS
- forward the traffic from this protocol-agnostic proxy to Burp so we can play with packets

Solution design-1st part (interception and decryption)

- Burpsuite does that job No. 1, but only for HTTP(S) it does not speak SIP or any other non-HTTP(S) protocols for that matter
- The Solution: mitm_relay.py (https://github.com/jrmdev/mitm-relay)

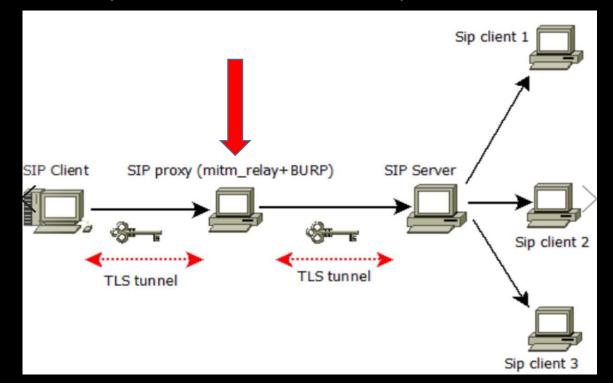
Solution design-1st part (interception and decryption)

-Topology design developed as below

-Key component: mitm_relay.py + Burpsuite running as SIP proxy on Kali

- you can have them on separate VMs too, also, no particular need for Kali – any linux

will do, I guess



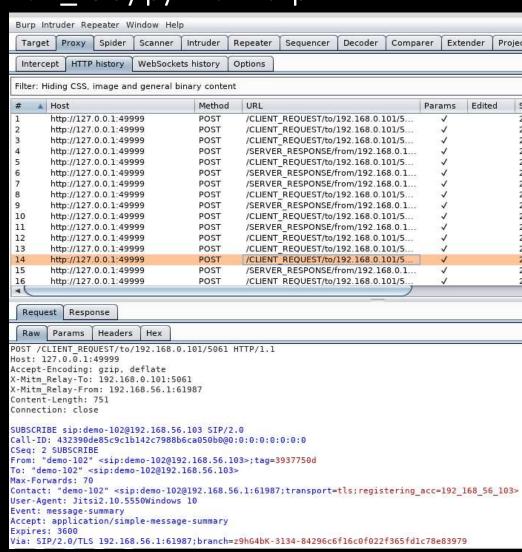
Chaining mitm_relay.py with Burp

- -l 0.0.0.0 mitm_relay.py listens on all interfaces
- -p 127.0.0.1:8080 Burp runs on localhost , port 8080
- -r tcp:5061:192.168.0.101:5061 relay all traffic on tcp 5061 port to final SIP server running on 192.168.0.101, port 5061
- -c mitm.pem digital certificate for mitm_relay.py
- -k privatemitm.key private key for mitm_relay.py

```
root@kali:/projekti/VOIP# !174
python /opt/mitm relay.py -l 0.0.0.0 -p 127.0.0.1:8080 -r tcp:5061:192.168.0.101:5061 -c mitm.pem -k privatemitm.
[+] Webserver listening on ('127.0.0.1', 49999)
[+] Relay listening on tcp 5061 -> 192.168.0.101:5061
[+] New client: ('192.168.56.1', 61984) --> 1892.168.0.101:5061
C >> S [ 192.168.56.1:61984 >> 192.168.0.101:5061 ] [
                                                     Fri 05 Apr 17:55:44 ] [ 267 ]
00000000: 16 03 03 01 06 10 00 01 02 01 00 30 8f ec a7 03
          af 4a a4 ff e5 c4 6c 32 a2 a3 d7 4d 7d b0 14 ac
                                                              .J....l2...M}...
                                                             |Z\rangle \dots h' \dots t[.x]
                                                             ||q/`...)a...o.>.
                                                             |..F).@...pD..".5
                                                             |?..3....0.+...
                                                              .7....K..{.HG.I.
                                                              .~v.M.ob..aT..{Q
                                                              ......lJ..G..a..
                                                             t.-.=...i<...Z.
                                                              [.)[...(..Y...:..
                                                             [.]............[.D]
                                                             |>w(v...j[.....v.
                                                              ....*.5...Dw...i
                                                             T..N.......8...
                                                             .g.HT.*....
00000100: b4 67 cb 48 54 8f 2a 0e ec c5 07
```

Chaining mitm_relay.py with Burp

- If we switch over to Burp, we can see decrypted SIP negotiation
- POST requests have embedded "/CLIENT_REQUEST/to/<IP>" and "/SERVER_RESPONSE/from/<IP>"
- that allows Burp to process SIP Sessions (it only transports stuff, does not really speak SIP)



Chaining mitm_relay.py with Burp

- Mechanics of the interception
- SIP client ->mitm_relay.py->Burpsuite->SIP server->SIP client 2
- NOTE: remember that now that we have Burp reading the SIP, several other attacks can be mounted:
 - -send SIP request to Burp Repeater,
 - -change call destinations,
 - -brute force destination numbers,
 - -change user agent fingerprint,
 - -inject some funky headers/establish covert channel attack,
 - -spoof calling ID...
- Tampering in BURP is interesting, but out of scope of this research

Solution design - 2nd part (coding new app)

- Why is new app required –what about sipdump and sipcrack?

 sipdump will not work here –we have dumped the decrypted session into a non-pcap format
 - -sipdump has no way of importing the private key to decrypt the captured wireshark pcap session
 - -we can decrypt wireshark pcap, remember? we possess the private key for mitm_relay.py
- What is new app doing?
 - -new app will process the mitm_relay.py dump which is more-less text based file and will extract all authentication attributes and perform password dictionary-based attack

Solution design - 2nd part (coding new app)

- Functional mapping of sipdump+sipcrack into sipcrack2
- Streamline mitm_relay dump parsing with password cracking
- How do we crack digest algorithm?
 - -https://www.ietf.org/rfc/rfc2069.txt later ammended by RFC 2617
 - -{qop},{nc},{cnonce} increase the variability of hashing but will still not protect against our attack and all these attributes can be intercepted and decrypted

Solution design - 2nd part (coding new app)

- Digest authentication algorithm
- If {qop},{cnonce} and {nc} are not defined (RFC2069) then H1=MD5(username:realm:password) H2=MD5(method:uri) Response=MD5 (H1:nonce:H2)

Else if {qop},{cnonce} and {nc } are defined (RFC2617) then

Response=MD5(H1:nonce:nc:cnonce:qop:H2)

https://github.com/adenosine-phosphatase/sipcrack2

Final thoughts

- Development sipcrack2 is for now just a linux version, hope to release Windows version with CUDA/multithreading & parallel processing in a near future
- How realistic/difficult is the attack?
 The attack shown is relatively difficult to implement (ARP poisoning required to redirect the traffic from legitimate proxy to the attacker)
- Recommendations
 - -use strong passwords
 - -do not use self-signed certificates
 - -use client side certificate in addition to server
- Risk rating proposed: Medium

<u>Demo</u>



Questions?

```
Telling INIT to go to single user mode.
init: rc main process (2205) killed by TERM signal
[root@centos-4 /]#
```

Shutting down