

CRYPTO Q

Using quantum in classical crypto

TEAM 2

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Problems we were trying to tackle

Symmetric key distribution

- Diffie-Helman: Key exchanged using same channel as data
- Several interactions to agree about the key

Blockchain consensus algorithms

- Proof-of-work? Expensive and slow
- Proof-of-stake? Problem of nothing-at-stake
- Proof-of-elapsed time? I can cheat selecting my random timer.

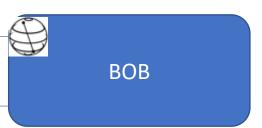
Quantum Key Distribution

- Use quantum entanglement to share the symmetric key
- Encrypt using classical cryptography and send data using classical link.

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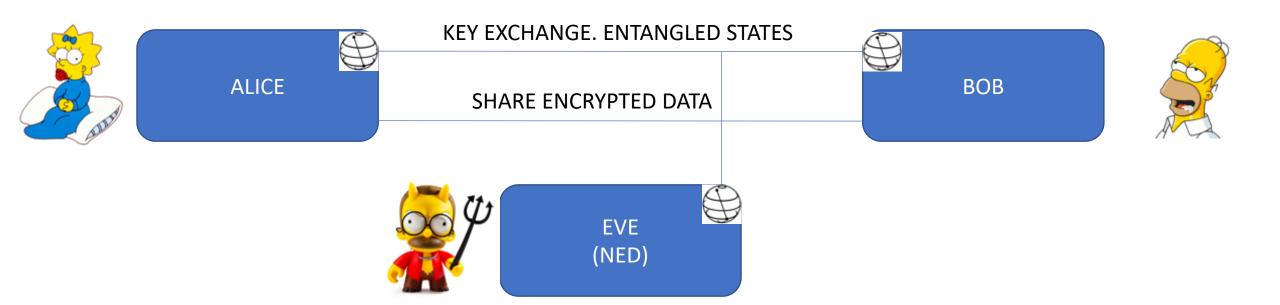






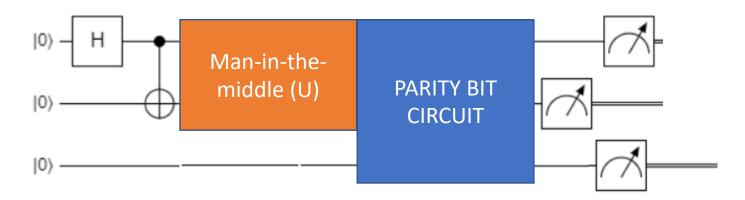
Quantum Key Distribution

- BUT WAIT! Our key exchange could be eavesdropped or sabotaged!
- Simulating Eve's effect and adding a parity bit to detect this matter.



How did we do it? Our parity circuit!

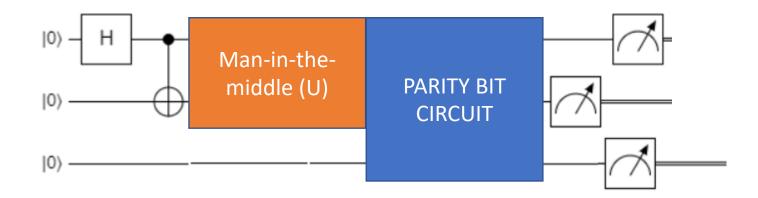
• We implemented a basic circuit cell to be reused all over the project.



$$U = R(\theta_A) \otimes R(\theta_B)$$

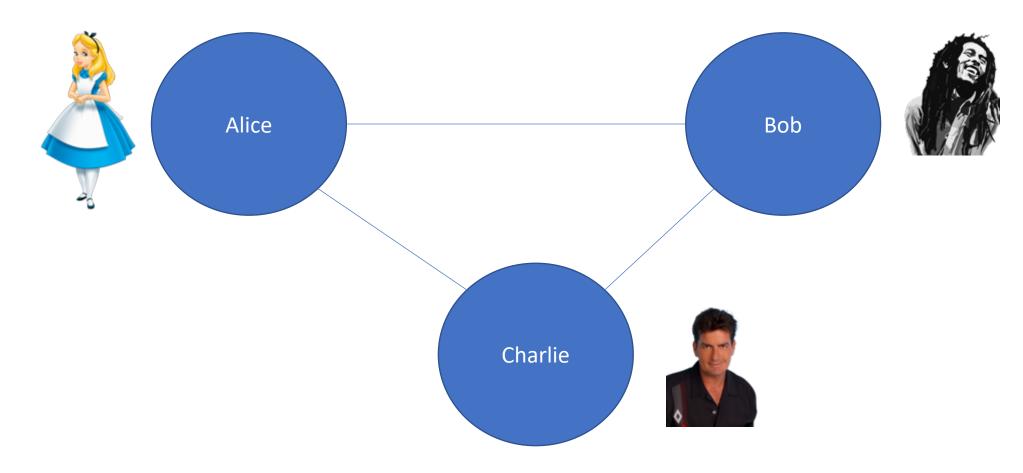
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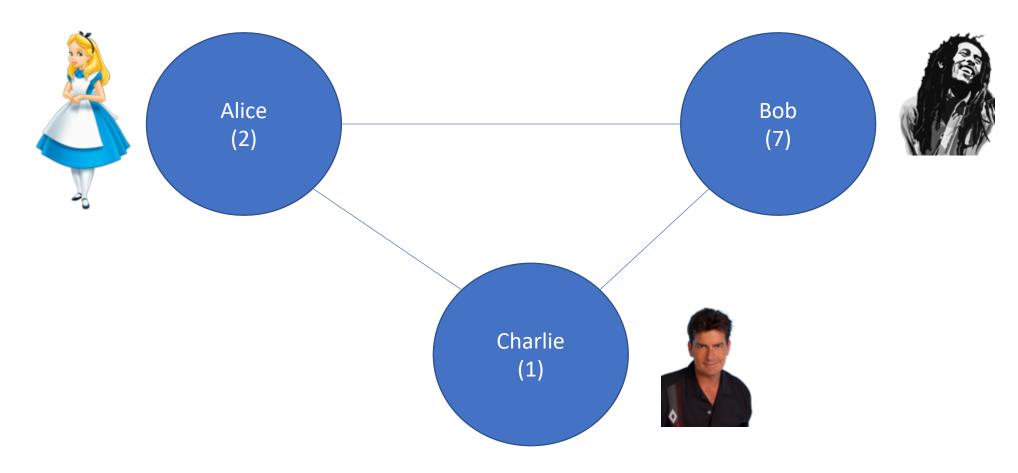


- It worked like a charm in simulation! We detected man-in-the-middle effects and fixed it. Security level of keys may be set.
- However... things started breaking in a real device. Noice affected our parity bit (we could fix it with a classical processing after measurement)

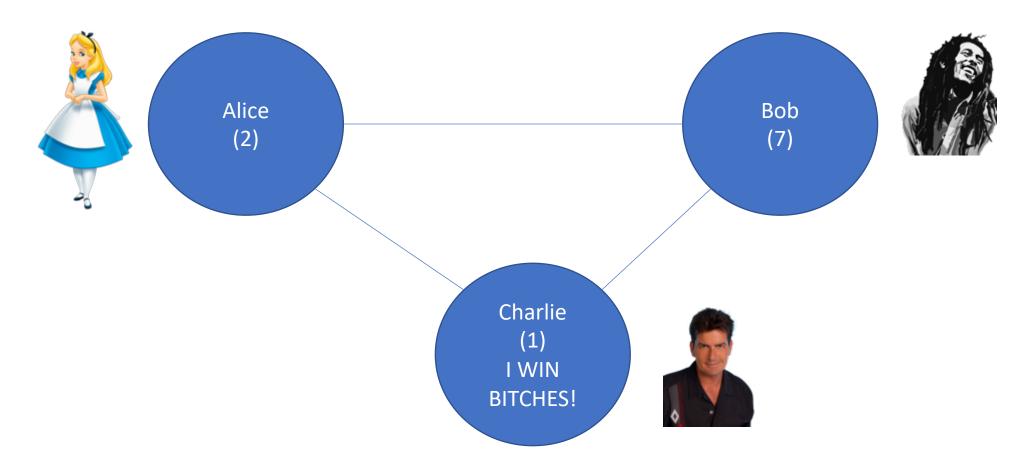
Quantum Distributed Consensus

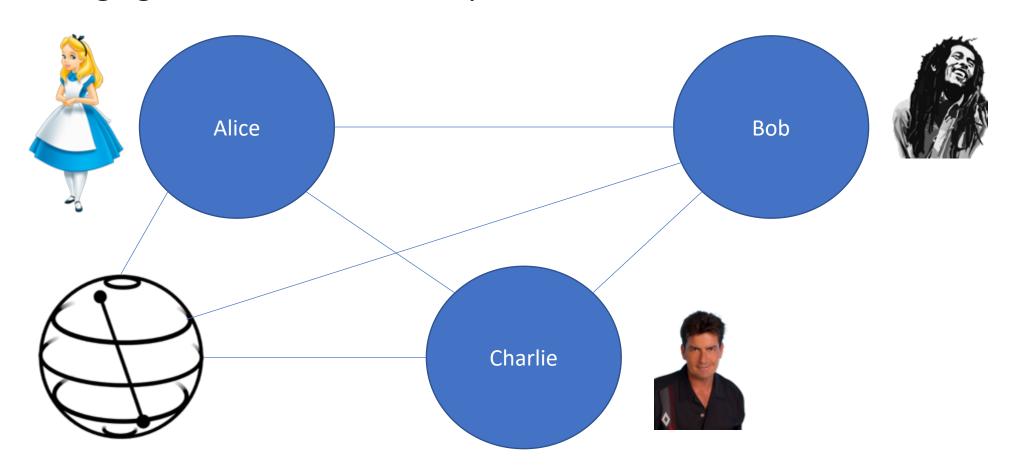


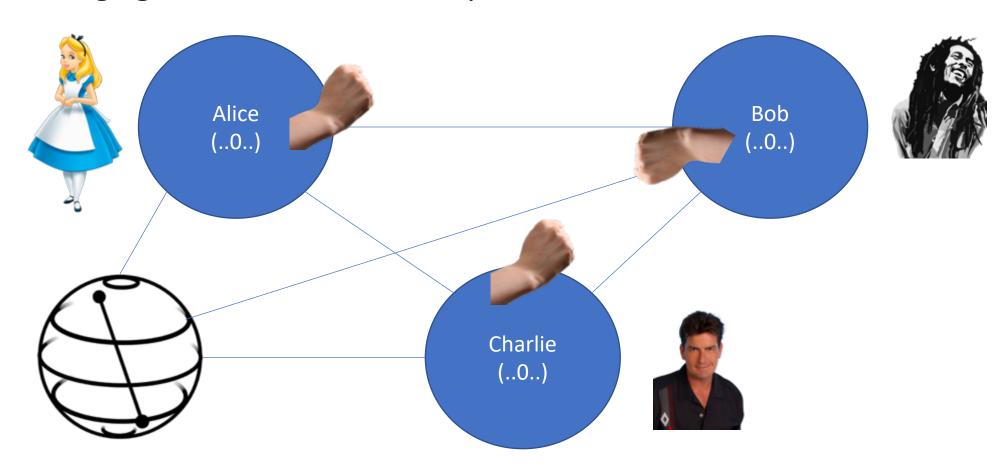
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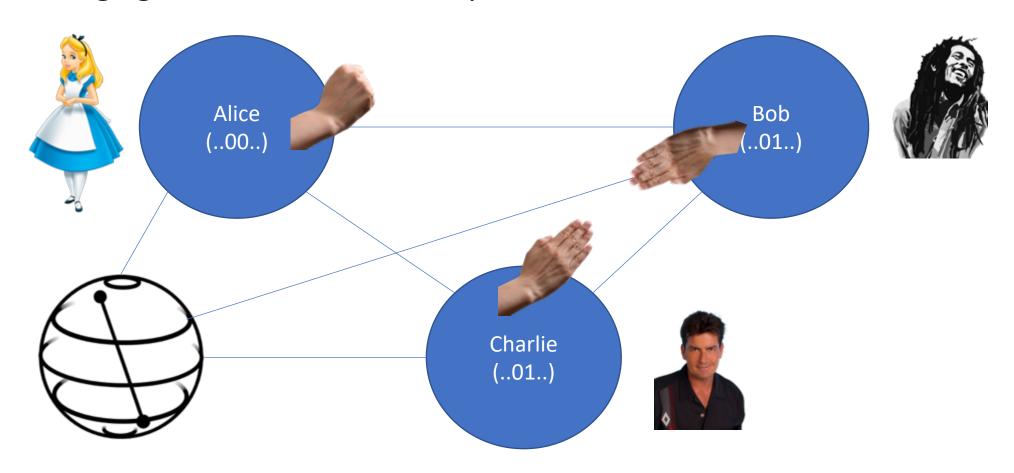


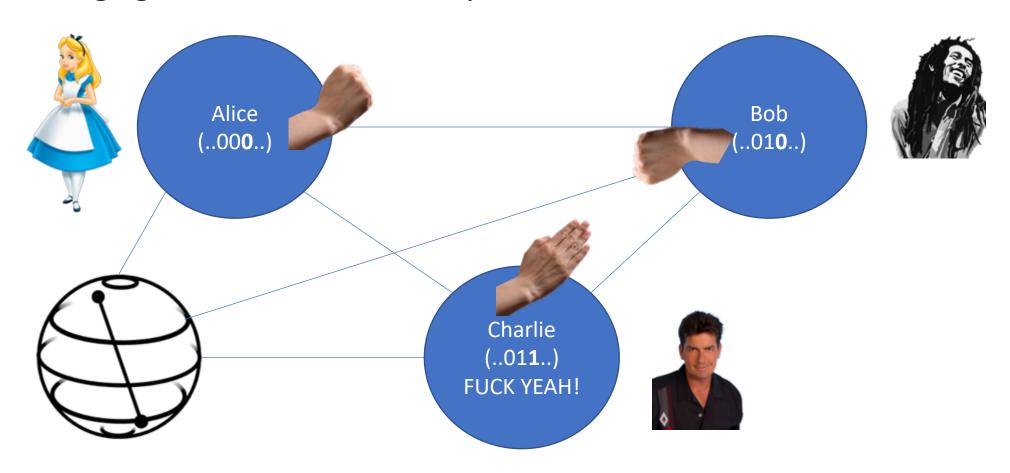
Quantum Distributed Consensus











- Bitwise entanglement so cheating may be detected (I know my bit and someone else's)
- But we still can cheat...
- Simulate cheating with "cheating matrices"

$$Ch_1|A,B,C\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix} = |1\rangle$$

$$Ch_2|A_b, B_c, C_a\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix} \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} = |0\rangle$$

• We are aiming a **Byzantine Fault Tolerant Consensus** (1/3 can cheat)

Cheating allowed	Consensus
No cheating	Always
JUST cheat 1 JUST cheat 2	No one cheats → Consensus One player cheats → Consensus Two players cheat → 1/4 prob. consensus Everyone cheats → 1/8 prob. consensus
Cheat 1 AND Cheat 2	No one cheats → Consensus One player cheats → ¼ prob. consensus Two players cheat → No consensus Everyone cheats → No consensus

What else? Future work.

Quantum Key Distribution in real devices with parity bit.

• Enhanced QRP consensus where cheating is penalized. Test it in a real setup.

• QRPS (Rock-Paper-Scissors) consensus with encoded 3-D Qudits

 Learn more about quantum and qiskit! And publish a paper? We'll see.

Quantum Physics