

KFC
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1. DATA STRUCTURES

1.1. Fenwick Tree.

```
struct fenwick {
- vi ar;
- fenwick(vi &ar) : ar(_ar.size(), 0) {
- for (int i = 0; i < ar.size(); ++i) {
- ar[i] += _ar[i];
- int j = i | (i+1);
- if (j < ar.size())
- ar[j] += ar[i]; } }
- int sum(int i) {
- int res = 0;
- for (; i >= 0; i = (i & (i+1)) - 1)
- res += ar[i];
- return res; }
- int sum(int i, int j) { return sum(j) - sum(i-1); }
- void add(int i, int val) {
- for (; i < ar.size(); i |= i+1)
- ar[i] += val; }
- int get(int i) {
- int res = ar[i];
- if (i) {
- int lca = (i & (i+1)) - 1;
- for (--i; i != lca; i = (i&(i+1))-1)
- res -= ar[i]; }
- return res; }
- void set(int i, int val) { add(i, -get(i) + val); }
- // range update, point query //
- void add(int i, int j, int val) {
- add(i, val); add(j+1, -val); }
- int getl(int i) { return sum(i); } };
```

1.2. Leq Counter.

```
1.2.1. Leq Counter Array.
#include "segtree.cpp"
struct LeqCounter {
- segtree **roots;
- LeqCounter(int *ar, int n) {
- std::vector<ii> nums;
- for (int i = 0; i < n; ++i)
- nums.push_back({ar[i], i});
- std::sort(nums.begin(), nums.end());
- roots = new segtree*[n];
- roots[0] = new segtree(0, n);
- int prev = 0;
- for (ii &e : nums) {
- for (int i = prev+1; i < e.first; ++i)
- roots[i] = roots[prev];
- roots[e.first] = roots[prev]->update(e.second, 1);
- prev = e.first; }
- for (int i = prev+1; i < n; ++i)
- roots[i] = roots[prev]; }
- int count(int i, int j, int x) {
- return roots[x]->query(i, j); } };
```

1.2.2. Leq Counter Map.

```
struct LeqCounter {
- std::map<int, segtree*> roots;
- std::set<int> neg_nums;
- LeqCounter(int *ar, int n) {
- std::vector<ii> nums;
- for (int i = 0; i < n; ++i) {
- nums.push_back({ar[i], i});
- neg_nums.insert(-ar[i]);
- }
- std::sort(nums.begin(), nums.end());
- roots[0] = new segtree(0, n);
- int prev = 0;
- for (ii &e : nums) {
- roots[e.first] = roots[prev]->update(e.second, 1);
- prev = e.first; } }
- int count(int i, int j, int x) {
- auto it = neg_nums.lower_bound(-x);
- if (it == neg_nums.end()) return 0;
- return roots[-*it]->query(i, j); } };
```

1.3. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest element.

```
#define BITS 15
struct misof_tree {
- int cnt[BITS][1<<BITS];
- misof_tree() { memset(cnt, 0, sizeof(cnt)); }
- void insert(int x) {
- for (int i = 0; i < BITS; cnt[i++][x]++, x >= 1); }
- void erase(int x) {
- for (int i = 0; i < BITS; cnt[i++][x]--, x >= 1); }
- int nth(int n) {
- int res = 0;
- for (int i = BITS-1; i >= 0; i--)
```

```
---- if (cnt[i][res <= 1] <= n) n -= cnt[i][res], res |= 1;
- return res; } };
```

1.4. Mo's Algorithm.

```
struct query {
- int id, l, r; ll hilbert_index;
- query(int id, int l, int r) : id(id), l(l), r(r) {
- hilbert_index = hilbert_order(l, r, LOGN, 0); }
- ll hilbert_order(int x, int y, int pow, int rotate) {
- if (pow == 0) return 0;
- int hpow = 1 << (pow-1);
- int seg = ((x<hpow) ? ((y<hpow)?0:3) : ((y<hpow)?1:2));
- seg = (seg + rotate) & 3;
- const int rotate_delta[4] = {3, 0, 0, 1};
- int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
- int nrot = (rotate + rotate_delta[seg]) & 3;
- ll sub_sq_size = ll(1) << (2*pow - 2);
- ll ans = seg * sub_sq_size;
- ll add = hilbert_order(nx, ny, pow-1, nrot);
- ans += (seg==1 || seg==2) ? add : (sub_sq_size-add-1);
- return ans; }
- bool operator<(const query& other) const {
- return this->hilbert_index < other.hilbert_index; } };
std::vector<query> queries;
for(const query &q : queries) { // [l,r] inclusive
- for(; r > q.r; r--) update(r, -1);
- for(r = r+1; r <= q.r; r++) update(r);
- r--;
- for(; l < q.l; l++) update(l, -1);
- for(l = l-1; l >= q.l; l--) update(l);
- l++; }
```

1.5. Ordered Statistics Tree.

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename T>
using index_set = tree<T, null_type, std::less<T>,
splay_tree_tag, tree_order_statistics_node_update>;
// indexed_set<int> t; t.insert(...);
// t.find_by_order(index); // 0-based
// t.order_of_key(key);
```

1.6. Segment Tree.

1.6.1. Recursive, Point-update Segment Tree.

1.6.2. Iterative, Point-update Segment Tree.

```
struct segtree {
- int n;
- int *vals;
- segtree(vi &ar, int n) {
- this->n = n;
- vals = new int[2*n];
- for (int i = 0; i < n; ++i)
- vals[i+n] = ar[i];
- for (int i = n-1; i > 0; --i)
- vals[i] = vals[i<<1] + vals[i<<1|1]; }
```

```
- void update(int i, int v) {
-- for (vals[i += n] += v; i > 1; i >>= 1)
--   vals[i>>1] = vals[i] + vals[i^1]; }
- int query(int l, int r) {
-- int res = 0;
-- for (l += n, r += n+1; l < r; l >>= 1, r >>= 1) {
--   if (l&1) res += vals[l++];
--   if (r&1) res += vals[--r]; }
-- return res; } };
```

1.6.3. Pointer-based, Range-update Segment Tree.

```
struct segtree {
- int i, j, val, temp_val = 0;
- segtree *l, *r;
- segtree(vi &ar, int _i, int _j) : i(_i), j(_j) {
-- if (i == j) {
--   val = ar[i];
--   l = r = NULL;
-- } else {
--   int k = (i + j) >> 1;
--   l = new segtree(ar, i, k);
--   r = new segtree(ar, k+1, j);
--   val = l->val + r->val; } }
- void visit() {
-- if (temp_val) {
--   val += (j-i+1) * temp_val;
--   if (l) {
--     l->temp_val += temp_val;
--     r->temp_val += temp_val; }
--   temp_val = 0; } }
- void increase(int _i, int _j, int _inc) {
-- visit();
-- if (_i <= i && j <= _j) {
--   temp_val += _inc;
--   visit();
-- } else if (_j < i or j < _i) {
--   // do nothing
-- } else {
--   l->increase(_i, _j, _inc);
--   r->increase(_i, _j, _inc);
--   val = l->val + r->val; } }
- int query(int _i, int _j) {
-- visit();
-- if (_i <= i and j <= _j)
--   return val;
-- else if (_j < i || j < _i)
--   return 0;
-- else
--   return l->query(_i, _j) + r->query(_i, _j);
} };
```

1.6.4. Array-based, Range-update Segment Tree -.

```
struct segtree {
- int n, *vals, *deltas;
- segtree(vi &ar) {
-- n = ar.size();
-- vals = new int[4*n];
```

```
-- deltas = new int[4*n];
-- build(ar, 1, 0, n-1); }
- void build(vi &ar, int p, int i, int j) {
-- deltas[p] = 0;
-- if (i == j)
--   vals[p] = ar[i];
-- else {
--   int k = (i + j) / 2;
--   build(ar, p<<1, i, k);
--   build(ar, p<<1|1, k+1, j);
--   pull(p); } }
- void pull(int p) { vals[p] = vals[p<<1] + vals[p<<1|1]; }
- void push(int p, int i, int j) {
-- if (deltas[p]) {
--   vals[p] += (j - i + 1) * deltas[p];
--   if (i != j) {
--     deltas[p<<1] += deltas[p];
--     deltas[p<<1|1] += deltas[p]; }
--   deltas[p] = 0; } }
- void update(int _i, int _j, int v, int p, int i, int j) {
-- push(p, i, j);
-- if (_i <= i && j <= _j) {
--   deltas[p] += v;
--   push(p, i, j);
-- } else if (_j < i || j < _i) {
--   // do nothing
-- } else {
--   int k = (i + j) / 2;
--   update(_i, _j, v, p<<1, i, k);
--   update(_i, _j, v, p<<1|1, k+1, j);
--   pull(p); } }
- int query(int _i, int _j, int p, int i, int j) {
-- push(p, i, j);
-- if (_i <= i and j <= _j)
--   return vals[p];
-- else if (_j < i || j < _i)
--   return 0;
-- else {
--   int k = (i + j) / 2;
--   return query(_i, _j, p<<1, i, k) +
--     query(_i, _j, p<<1|1, k+1, j); } } };
```

1.6.5. 2D Segment Tree.

```
struct segtree_2d {
- int n, m, **ar;
- segtree_2d(int n, int m) {
-- this->n = n; this->m = m;
-- ar = new int[n];
-- for (int i = 0; i < n; ++i) {
--   ar[i] = new int[m];
--   for (int j = 0; j < m; ++j)
--     ar[i][j] = 0; } }
- void update(int x, int y, int v) {
-- ar[x + n][y + m] = v;
-- for (int i = x + n; i > 0; i >>= 1) {
--   for (int j = y + m; j > 0; j >>= 1) {
```

```
--   ar[i>>1][j] = min(ar[i][j], ar[i^1][j]);
--   ar[i][j>>1] = min(ar[i][j], ar[i][j^1]);
-- } } // just call update one by one to build
- int query(int x1, int x2, int y1, int y2) {
-- int s = INF;
-- if (~x2) for (int a=x1+n, b=x2+n+1; a<b; a>>=1, b>>=1) {
--   if (a & 1) s = min(s, query(a++, -1, y1, y2));
--   if (b & 1) s = min(s, query(--b, -1, y1, y2));
-- } else for (int a=y1+m, b=y2+m+1; a<b; a>>=1, b>>=1) {
--   if (a & 1) s = min(s, ar[x1][a++]);
--   if (b & 1) s = min(s, ar[x1][--b]);
-- } return s; } };
```

1.6.6. Persistent Segment Tree.

```
struct segtree {
- int i, j, val;
- segtree *l, *r;
- segtree(vi &ar, int _i, int _j) : i(_i), j(_j) {
-- if (i == j) {
--   val = ar[i];
--   l = r = NULL;
-- } else {
--   int k = (i+j) >> 1;
--   l = new segtree(ar, i, k);
--   r = new segtree(ar, k+1, j);
--   val = l->val + r->val;
-- } }
- segtree(int i, int j, segtree *l, segtree *r, int val) :
-- i(i), j(j), l(l), r(r), val(val) {}
- segtree* update(int _i, int _val) {
-- if (_i <= i and j <= _i)
--   return new segtree(i, j, l, r, val + _val);
-- else if (_i < i or j < _i)
--   return this;
-- else {
--   segtree *nl = l->update(_i, _val);
--   segtree *nr = r->update(_i, _val);
--   return new segtree(i, j, nl, nr, nl->val + nr->val); } }
- int query(int _i, int _j) {
-- if (_i <= i and j <= _j)
--   return val;
-- else if (_j < i or j < _i)
--   return 0;
-- else
--   return l->query(_i, _j) + r->query(_i, _j); } };
```

1.7. Sparse Table.

1.7.1. 1D Sparse table.

```
int lg[MAXN+1], spt[20][MAXN];
void build(vi &arr, int n) {
- lg[0] = lg[1] = 0;
- for (int i = 2; i <= n; ++i) lg[i] = lg[i>>1] + 1;
- for (int i = 0; i < n; ++i) spt[0][i] = arr[i];
- for (int j = 0; (2 <= j) <= n; ++j)
-- for (int i = 0; i + (2 <= j) <= n; ++i)
--   spt[j+1][i] = std::min(spt[j][i], spt[j][i+(1<<j)]); }
```

```
int query(int a, int b) {
- int k = lg[b-a+1], ab = b - (1<<k) + 1;
- return std::min(spt[k][a], spt[k][ab]); }
```

1.7.2. 2D Sparse Table.

```
const int N = 100, LGN = 20;
int lg[N], A[N][N], st[LGN][LGN][N][N];
void build(int n, int m) {
- for(int k=2; k<=std::max(n,m); ++k) lg[k] = lg[k>1]+1;
- for(int i = 0; i < n; ++i)
- for(int j = 0; j < m; ++j)
- st[0][0][i][j] = A[i][j];
- for(int bj = 0; (2 << bj) <= m; ++bj)
- for(int j = 0; j + (2 << bj) <= m; ++j)
- for(int i = 0; i < n; ++i)
- st[0][bj+1][i][j] =
- std::max(st[0][bj][i][j],
- st[0][bj][i][j + (1 << bj)]);
- for(int bi = 0; (2 << bi) <= n; ++bi)
- for(int i = 0; i + (2 << bi) <= n; ++i)
- for(int j = 0; j < m; ++j)
- st[bi+1][0][i][j] =
- std::max(st[bi][0][i][j],
- st[bi][0][i + (1 << bi)][j]);
- for(int bi = 0; (2 << bi) <= n; ++bi)
- for(int i = 0; i + (2 << bi) <= n; ++i)
- for(int bj = 0; (2 << bj) <= m; ++bj)
- for(int j = 0; j + (2 << bj) <= m; ++j) {
- int ik = i + (1 << bi);
- int jk = j + (1 << bj);
- st[bi+1][bj+1][i][j] =
- std::max(std::max(st[bi][bj][i][j],
- st[bi][bj][ik][j]),
- std::max(st[bi][bj][i][jk],
- st[bi][bj][ik][jk])); } }
int query(int x1, int x2, int y1, int y2) {
- int kx = lg[x2 - x1 + 1], ky = lg[y2 - y1 + 1];
- int x12 = x2 - (1<<kx) + 1, y12 = y2 - (1<<ky) + 1;
- return std::max(std::max(st[kx][ky][x1][y1],
- st[kx][ky][x1][y12]),
- std::max(st[kx][ky][x12][y1],
- st[kx][ky][x12][y12])); }
```

1.8. Splay Tree.

```
struct node *null;
struct node {
- node *left, *right, *parent;
- bool reverse; int size, value;
- node& get(int d) {return d == 0 ? left : right;}
- node(int v=0): reverse(0), size(0), value(v) {
- left = right = parent = null ? null : this; } };
struct SplayTree {
- node *root;
- SplayTree(int arr[] = NULL, int n = 0) {
- if (!null) null = new node();
- root = build(arr, n); }
- node* build(int arr[], int n) {
```

```
if (n == 0) return null;
int mid = n >> 1;
node *p = new node(arr ? arr[mid] : 0);
link(p, build(arr, mid), 0);
link(p, build(arr? arr+mid+1 : NULL, n-mid-1), 1);
pull(p); return p; }
void pull(node *p) {
- p->size = p->left->size + p->right->size + 1; }
void push(node *p) {
if (p != null && p->reverse) {
- swap(p->left, p->right);
- p->left->reverse ^= 1;
- p->right->reverse ^= 1;
- p->reverse ^= 1; } }
void link(node *p, node *son, int d) {
- p->get(d) = son;
- son->parent = p; }
int dir(node *p, node *son) {
- return p->left == son ? 0 : 1; }
void rotate(node *x, int d) {
- node *y = x->get(d), *z = x->parent;
- link(x, y->get(d ^ 1), d);
- link(y, x, d ^ 1);
- link(z, y, dir(z, x));
- pull(x); pull(y); }
node* splay(node *p) {
- while (p->parent != null) {
- node *m = p->parent, *g = m->parent;
- push(g); push(m); push(p);
- int dm = dir(m, p), dg = dir(g, m);
- if (g == null) rotate(m, dm);
- else if (dm == dg) rotate(g, dg), rotate(m, dm);
- else rotate(m, dm), rotate(g, dg);
- } return root = p; }
node* get(int k) {
- node *p = root;
- while (push(p), p->left->size != k) {
- if (k < p->left->size) p = p->left;
- else k -= p->left->size + 1, p = p->right; }
- return p == null ? null : splay(p); }
void split(node *&r, int k) {
if (k == 0) { r = root; root = null; return; }
- r = get(k - 1)->right;
- root->right = r->parent = null;
- pull(root); }
void merge(node *r) {
if (root == null) {root = r; return;}
- link(get(root->size - 1), r, 1);
- pull(root); }
void assign(int k, int val) {
- get(k)->value = val; pull(root); }
void reverse(int L, int R) {
- node *m, *r; split(r, R + 1); split(m, L);
- m->reverse ^= 1; push(m); merge(m); merge(r); }
node* insert(int k, int v) {
- node *r; split(r, k);
```

```
node *p = new node(v); p->size = 1;
link(root, p, 1); merge(r);
return p; }
void erase(int k) {
- node *r, *m;
- split(r, k + 1); split(m, k);
- merge(r); delete m; } }
```

1.9. Treap.

1.9.1. Implicit Treap.

```
struct cartree {
- typedef struct _Node {
- int node_val, subtree_val, delta, prio, size;
- _Node *l, *r;
- _Node(int val) : node_val(val), subtree_val(val),
- delta(0), prio((rand()<<16)^rand()), size(1),
- l(NULL), r(NULL) {}
- ~_Node() { delete l; delete r; }
- } *Node;
- int get_subtree_val(Node v) {
- return v ? v->subtree_val : 0; }
- int get_size(Node v) { return v ? v->size : 0; }
- void apply_delta(Node v, int delta) {
- if (!v) return;
- v->delta += delta;
- v->node_val += delta;
- v->subtree_val += delta * get_size(v); }
- void push_delta(Node v) {
- if (!v) return;
- apply_delta(v->l, v->delta);
- apply_delta(v->r, v->delta);
- v->delta = 0; }
- void update(Node v) {
- if (!v) return;
- v->subtree_val = get_subtree_val(v->l) + v->node_val
- + get_subtree_val(v->r);
- v->size = get_size(v->l) + 1 + get_size(v->r); }
- Node merge(Node l, Node r) {
- push_delta(l); push_delta(r);
- if (!l || !r) return l ? l : r;
- if (l->size <= r->size) {
- l->r = merge(l->r, r);
- update(l);
- return l; }
- else {
- r->l = merge(l, r->l);
- update(r);
- return r; } }
- void split(Node v, int key, Node &l, Node &r) {
- push_delta(v);
- l = r = NULL;
- if (!v) return;
- if (key <= get_size(v->l)) {
- split(v->l, key, l, v->l);
- r = v;
- } else {
```

```
----- split(v->r, key - get_size(v->l) - 1, v->r, r); -----
----- l = v; } -----
----- update(v); } -----
- Node root; -----
public:
- cartree() : root(NULL) {} -----
- ~cartree() { delete root; } -----
- int get(Node v, int key) { -----
----- push_delta(v); -----
----- if (key < get_size(v->l)) -----
-----     return get(v->l, key); -----
----- else if (key > get_size(v->l)) -----
-----     return get(v->r, key - get_size(v->l) - 1); -----
----- return v->node_val; } -----
- int get(int key) { return get(root, key); } -----
- void insert(Node item, int key) { -----
----- Node l, r; -----
----- split(root, key, l, r); -----
----- root = merge(merge(l, item), r); } -----
- void insert(int key, int val) { -----
----- insert(new _Node(val), key); } -----
- void erase(int key) { -----
----- Node l, m, r; -----
----- split(root, key + 1, m, r); -----
----- split(m, key, l, m); -----
----- delete m; -----
----- root = merge(l, r); } -----
- int query(int a, int b) { -----
----- Node l1, r1; -----
----- split(root, b+1, l1, r1); -----
----- Node l2, r2; -----
----- split(l1, a, l2, r2); -----
----- int res = get_subtree_val(r2); -----
----- l1 = merge(l2, r2); -----
----- root = merge(l1, r1); -----
----- return res; } -----
- void update(int a, int b, int delta) { -----
----- Node l1, r1; -----
----- split(root, b+1, l1, r1); -----
----- Node l2, r2; -----
----- split(l1, a, l2, r2); -----
----- apply_delta(r2, delta); -----
----- l1 = merge(l2, r2); -----
----- root = merge(l1, r1); } -----
- int size() { return get_size(root); } };
```

1.9.2. Persistent Treap .

1.10. Union Find.

```
struct union_find { -----
- vi p; union_find(int n) : p(n, -1) {} -----
- int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); } -----
- bool unite(int x, int y) { -----
----- int xp = find(x), yp = find(y); -----
----- if (xp == yp) return false; -----
----- if (p[xp] > p[yp]) std::swap(xp,yp); -----
```

```
--- p[xp] += p[yp], p[yp] = xp; return true; } -----
- int size(int x) { return -p[find(x)]; } };
```

1.11. Unique Counter.

```
struct UniqueCounter { -----
- int *B; std::map<int, int> last; LeqCounter *leq_cnt; -----
- UniqueCounter(int *ar, int n) { // 0-index A[i] -----
----- B = new int[n+1]; -----
----- B[0] = 0; -----
----- for (int i = 1; i <= n; ++i) { -----
-----     B[i] = last[ar[i-1]]; -----
-----     last[ar[i-1]] = i; } -----
----- leq_cnt = new LeqCounter(B, n+1); } -----
- int count(int l, int r) { -----
----- return leq_cnt->count(l+1, r+1, l); } };
```

2. DYNAMIC PROGRAMMING

2.1. Dynamic Convex Hull Trick.

```
// USAGE: hull.insert_line(m, b); hull.gety(x); -----
bool UPPER_HULL = true; // you can edit this -----
bool IS_QUERY = false, SPECIAL = false; -----
struct line { -----
- ll m, b; line(ll m=0, ll b=0): m(m), b(b) {} -----
- mutable std::multiset<line>::iterator it; -----
- const line *see(std::multiset<line>::iterator it)const; -----
- bool operator < (const line& k) const { -----
----- if (!IS_QUERY) return m < k.m; -----
----- if (!SPECIAL) { -----
-----     ll x = k.m; const line *s = see(it); -----
-----     if (!s) return 0; -----
-----     return (b - s->b) < (x) * (s->m - m); -----
----- } else { -----
-----     ll y = k.m; const line *s = see(it); -----
-----     if (!s) return 0; -----
-----     ll n1 = y - b, d1 = m; -----
-----     ll n2 = b - s->b, d2 = s->m - m; -----
-----     if (d1 < 0) n1 *= -1, d1 *= -1; -----
-----     if (d2 < 0) n2 *= -1, d2 *= -1; -----
-----     return (n1) * d2 > (n2) * d1; } } } -----
struct dynamic_hull : std::multiset<line> { -----
- bool bad(iterator y) { -----
----- iterator z = next(y); -----
----- if (y == begin()) { -----
-----     if (z == end()) return 0; -----
-----     return y->m == z->m && y->b <= z->b; } -----
----- iterator x = prev(y); -----
----- if (z == end()) return y->m == x->m && y->b <= x->b; -----
----- return (x->b - y->b)*(z->m - y->m)>= -----
----- (y->b - z->b)*(y->m - x->m); } -----
- iterator next(iterator y) {return ++y;} -----
- iterator prev(iterator y) {return --y;} -----
- void insert_line(ll m, ll b) { -----
----- IS_QUERY = false; -----
----- if (!UPPER_HULL) m *= -1; -----
----- iterator y = insert(line(m, b)); -----
----- y->it = y; if (bad(y)) {erase(y); return;} -----
```

```
--- while (next(y) != end() && bad(next(y))) -----
----- erase(next(y)); -----
--- while (y != begin() && bad(prev(y))) -----
----- erase(prev(y)); } -----
- ll gety(ll x) { -----
----- IS_QUERY = true; SPECIAL = false; -----
----- const line& L = *lower_bound(line(x, 0)); -----
----- ll y = (L.m) * x + L.b; -----
----- return UPPER_HULL ? y : -y; } -----
- ll getx(ll y) { -----
----- IS_QUERY = true; SPECIAL = true; -----
----- const line& l = *lower_bound(line(y, 0)); -----
----- return /*floor*/ ((y - l.b + l.m - 1) / l.m); } -----
} hull; -----
const line* line::see(std::multiset<line>::iterator it) -----
const {return ++it == hull.end() ? NULL : &*it;} -----
```

2.2. Divide and Conquer Optimization. For DP problems of the form

$$dp(i,j) = \min_{k \leq j} \{ dp(i-1,k) + C(k,j) \}$$

where $C(k,j)$ is some cost function.

```
ll dp[G+1][N+1]; -----
void solve_dp(int g, int k_L, int k_R, int n_L, int n_R) { ---
- int n_M = (n_L+n_R)/2; -----
- dp[g][n_M] = INF; -----
- int best_k = -1; -----
- for (int k = k_L; k <= n_M && k <= k_R; k++) -----
----- if (dp[g-1][k]+cost(k+1,n_M) < dp[g][n_M]) { -----
-----     dp[g][n_M] = dp[g-1][k]+cost(k+1,n_M); -----
-----     best_k = k; } -----
- if (n_L <= n_M-1) -----
----- solve_dp(g, k_L, best_k, n_L, n_M-1); -----
- if (n_M+1 <= n_R) -----
----- solve_dp(g, best_k, k_R, n_M+1, n_R); } -----
```

3. GEOMETRY

```
#include <complex> -----
#define x real() -----
#define y imag() -----
typedef std::complex<double> point; // 2D point only -----
const double PI = acos(-1.0), EPS = 1e-7; -----
```

3.1. Dots and Cross Products.

```
double dot(point a, point b) { -----
- return a.x * b.x + a.y * b.y; } // + a.z * b.z; -----
double cross(point a, point b) { -----
- return a.x * b.y - a.y * b.x; } -----
double cross(point a, point b, point c) { -----
- return cross(a, b) + cross(b, c) + cross(c, a); } -----
double cross3D(point a, point b) { -----
- return point(a.x*b.y - a.y*b.x, a.y*b.z - -----
----- a.z*b.y, a.z*b.x - a.x*b.z); } -----
```

3.2. Angles and Rotations.

```
double angle(point a, point b, point c) { -----
- // angle formed by abc in radians: PI < x <= PI -----
- return abs(remainder(arg(a-b) - arg(c-b), 2*PI)); } -----
point rotate(point p, point a, double d) { -----
- //rotate point a about pivot p CCW at d radians -----
- return p + (a - p) * point(cos(d), sin(d)); } -----
```

3.3. Spherical Coordinates.

$$\begin{aligned} x &= r \cos \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \cos \theta \sin \phi & \theta &= \cos^{-1} x/r \\ z &= r \sin \theta & \phi &= \text{atan2}(y, x) \end{aligned}$$

3.4. Point Projection.

```
point proj(point p, point v) { -----
- // project point p onto a vector v (2D & 3D) -----
- return dot(p, v) / norm(v) * v; } -----
point projLine(point p, point a, point b) { -----
- // project point p onto line ab (2D & 3D) -----
- return a + dot(p-a, b-a) / norm(b-a) * (b-a); } -----
point projSeg(point p, point a, point b) { -----
- // project point p onto segment ab (2D & 3D) -----
- double s = dot(p-a, b-a) / norm(b-a);
- return a + min(1.0, max(0.0, s)) * (b-a); } -----
point projPlane(point p, double a, double b, -----
- double c, double d) { -----
- // project p onto plane ax+by+cz+d=0 (3D) -----
- // same as: o + p - project(p - o, n); -----
- double k = -d / (a*a + b*b + c*c); -----
- point o(a*k, b*k, c*k), n(a, b, c); -----
- point v(p.x-o.x, p.y-o.y, p.z-o.z); -----
- double s = dot(v, n) / dot(n, n); -----
- return point(o.x + p.x + s * n.x, o.y + -----
- p.y +s * n.y, o.z + p.z + s * n.z); } -----
```

3.5. Great Circle Distance.

```
double greatCircleDist(double lat1, double long1, -----
- double lat2, double long2, double R) { -----
- long1 *= PI / 180; lat1 *= PI / 180; // to radians -----
- long2 *= PI / 180; lat2 *= PI / 180; -----
- return R*acos(sin(lat1)*sin(lat2) + -----
- cos(lat1)*cos(lat2)*cos(abs(long1 - long2))); } -----
// another version, using actual (x, y, z) -----
double greatCircleDist(point a, point b) { -----
- return atan2(abs(cross3D(a, b)), dot3D(a, b)); } -----
```

3.6. Point/Line/Plane Distances.

```
double distPtLine(point p, double a, double b, double c) { ---
- // dist from point p to line ax+by+c=0 ---
- return abs(a*p.x + b*p.y + c) / sqrt(a*a + b*b);}
double distPtLine(point p, point a, point b) { -----
- // dist from point p to line ab -----
- return abs((a.y - b.y) * (p.x - a.x) + -----
- (b.x - a.x) * (p.y - a.y)) / -----
- hypot(a.x - b.x, a.y - b.y);}
double distPtPlane(point p, double a, double b, -----
- double c, double d) { -----
```

```
- // distance to 3D plane ax + by + cz + d = 0 -----
- return (a*p.x+b*p.y+c*p.z+d)/sqrt(a*a+b*b+c*c); } -----
/*! // distance between 3D lines AB & CD (untested) -----
double distLine3D(point A,point B,point C,point D){ -----
- point u = B - A, v = D - C, w = A - C; -----
- double a = dot(u, u), b = dot(u, v); -----
- double c = dot(v, v), d = dot(u, w); -----
- double e = dot(v, w), det = a*c - b*b; -----
- double s = det < EPS ? 0.0 : (b*e - c*d) / det; -----
- double t = det < EPS -----
- ? (b > c ? d/b : e/c) // parallel -----
- : (a*e - b*d) / det; -----
- point top = A + u * s, bot = w - A - v * t; -----
- return dist(top, bot); -----
} // dist<EPS: intersection */ -----
```

3.7. Intersections.

```
3.7.1. Line-Segment Intersection. Get intersection points of 2D -----
lines/segments  $\overline{ab}$  and  $\overline{cd}$ .
point null(HUGE_VAL, HUGE_VAL); -----
point line_inter(point a, point b, point c, -----
- point d, bool seg = false) { -----
- point ab(b.x - a.x, b.y - a.y); -----
- point cd(d.x - c.x, d.y - c.y); -----
- point ac(c.x - a.x, c.y - a.y); -----
- double D = -cross(ab, cd); // determinant -----
- double Ds = cross(cd, ac); -----
- double Dt = cross(ab, ac); -----
- if (abs(D) < EPS) { // parallel -----
- if (seg && abs(Ds) < EPS) { // collinear -----
- point p[] = {a, b, c, d}; -----
- sort(p, p + 4, [](point a, point b) { -----
- return a.x < b.x-EPS || -----
- (dist(a,b) < EPS && a.y < b.y-EPS); }); -----
- return dist(p[1], p[2]) < EPS ? p[1] : null; } -----
- return null; } -----
- double s = Ds / D, t = Dt / D; -----
- if (seg && (min(s,t)<-EPS||max(s,t)>1+EPS)) return null; ---
- return point(a.x + s * ab.x, a.y + s * ab.y); } -----
/* double A = cross(d-a, b-a), B = cross(c-a, b-a); -----
return (B*d - A*c)/(B - A); */ -----
```

```
3.7.2. Circle-Line Intersection. Get intersection points of circle at center -----
c, radius r, and line  $\overline{ab}$ .
std::vector<point> CL_inter(point c, double r, -----
- point a, point b) { -----
- point p = projLine(c, a, b); -----
- double d = abs(c - p); vector<point> ans; -----
- if (d > r + EPS); // none -----
- else if (d > r - EPS) ans.push_back(p); // tangent -----
- else if (d < EPS) { // diameter -----
- point v = r * (b - a) / abs(b - a); -----
- ans.push_back(c + v); -----
- ans.push_back(c - v); -----
- } else { -----
- double t = acos(d / r); -----
```

```
--- p = c + (p - c) * r / d; -----
--- ans.push_back(rotate(c, p, t)); -----
--- ans.push_back(rotate(c, p, -t)); -----
- } return ans; } -----
```

3.7.3. Circle-Circle Intersection.

```
std::vector<point> CC_intersection(point c1, -----
- double r1, point c2, double r2) { -----
- double d = dist(c1, c2); -----
- vector<point> ans; -----
- if (d < EPS) { -----
- if (abs(r1-r2) < EPS); // inf intersections -----
- } else if (r1 < EPS) { -----
- if (abs(d - r2) < EPS) ans.push_back(c1); -----
- } else { -----
- double s = (r1*r1 + d*d - r2*r2) / (2*r1*d); -----
- double t = acos(max(-1.0, min(1.0, s))); -----
- point mid = c1 + (c2 - c1) * r1 / d; -----
- ans.push_back(rotate(c1, mid, t)); -----
- if (abs(sin(t)) >= EPS) -----
- ans.push_back(rotate(c2, mid, -t)); -----
- } return ans; } -----
```

3.8. Areas.

3.8.1. Polygon Area. Find the area of any 2D polygon given as points in $O(n)$.

```
double area(point p[], int n) { -----
- double a = 0; -----
- for (int i = 0, j = n - 1; i < n; j = i++) -----
- a += cross(p[i], p[j]); -----
- return abs(a) / 2; } -----
```

3.8.2. Triangle Area. Find the area of a triangle using only their lengths. Lengths must be valid.

```
double area(double a, double b, double c) { -----
- double s = (a + b + c) / 2; -----
- return sqrt(s*(s-a)*(s-b)*(s-c)); } -----
```

3.8.3. Cyclic Quadrilateral Area. Find the area of a cyclic quadrilateral using only their lengths. A quadrilateral is cyclic if its inner angles sum up to 360°.

```
double area(double a, double b, double c, double d) { -----
- double s = (a + b + c + d) / 2; -----
- return sqrt((s-a)*(s-b)*(s-c)*(s-d)); } -----
```

3.9. Polygon Centroid. Get the centroid/center of mass of a polygon in $O(m)$.

```
point centroid(point p[], int n) { -----
- point ans(0, 0); -----
- double z = 0; -----
- for (int i = 0, j = n - 1; i < n; j = i++) { -----
- double cp = cross(p[j], p[i]); -----
- ans += (p[j] + p[i]) * cp; -----
- z += cp; -----
- } return ans / (3 * z); } -----
```

3.10. Convex Hull.

3.10.1. *2D Convex Hull*. Get the convex hull of a set of points using Graham-Andrew's scan. This sorts the points at $O(n \log n)$, then performs the Monotonic Chain Algorithm at $O(n)$.

```
// counterclockwise hull in p[], returns size of hull -----
bool xcmp(const point& a, const point& b) { -----
- return a.x < b.x || (a.x == b.x && a.y < b.y); }
int convex_hull(point p[], int n) { -----
- std::sort(p, p + n, xcmp); if (n <= 1) return n;
- int k = 0; point *h = new point[2 * n];
- double zer = EPS; // -EPS to include collinears -----
- for (int i = 0; i < n; h[k++] = p[i++]) -----
-- while (k >= 2 && cross(h[k-2],h[k-1],p[i]) < zer) -----
-- k; -----
- for(int i = n-2, t = k; i >= 0; h[k++] = p[i--]) -----
-- while (k > t && cross(h[k-2],h[k-1],p[i]) < zer) -----
-- k; -----
- k -= 1 + (h[0].x==h[1].x&&h[0].y==h[1].y ? 1 : 0);
- copy(h, h + k, p); delete[] h; return k; } -----
```

3.10.2. *3D Convex Hull*. Currently $O(N^2)$, but can be optimized to a randomized $O(N \log N)$ using the Clarkson-Shor algorithm. Sauce: Efficient 3D Convex Hull Tutorial on CF.

```
typedef std::vector<bool> vb; -----
struct point3D { -----
- ll x, y, z; -----
- point3D(ll x = 0, ll y = 0, ll z = 0) : x(x), y(y), z(z) {}
- point3D operator-(const point3D &o) const { -----
-- return point3D(x - o.x, y - o.y, z - o.z); } -----
- point3D cross(const point3D &o) const { -----
-- return point3D(y*o.z-z*o.y, z*o.x-x*o.z, x*o.y-y*o.x); }
- ll dot(const point3D &o) const { -----
-- return x*o.x + y*o.y + z*o.z; } -----
- bool operator==(const point3D &o) const { -----
-- return std::tie(x, y, z) == std::tie(o.x, o.y, o.z); } -----
- bool operator<(const point3D &o) const { -----
-- return std::tie(x, y, z) < std::tie(o.x, o.y, o.z); } };
struct face { -----
- std::vector<int> p_idx; -----
- point3D q; };
std::vector<face> convex_hull_3D(std::vector<point3D> &points) { -----
- int n = points.size(); -----
- std::vector<face> faces; -----
- std::vector<vb> dead(points.size(), vb(points.size(), true));
- auto add_face = [&](int a, int b, int c) { -----
-- faces.push_back({a, b, c}, -----
-- (points[b] - points[a]).cross(points[c] - points[a]));
-- dead[a][b] = dead[b][c] = dead[c][a] = false; };
- add_face(0, 1, 2); -----
- add_face(0, 2, 1); -----
- for (int i = 3; i < n; ++i) { -----
-- std::vector<face> faces_inv; -----
-- for(face &f : faces) { -----
---- if ((points[i] - points[f.p_idx[0]]).dot(f.q) > 0) -----
---- for (int j = 0; j < 3; ++j) -----
----- dead[f.p_idx[j]][f.p_idx[(j+1)%3]] = true; -----
----- else -----
```

```
----- faces_inv.push_back(f); } -----
- faces.clear(); -----
- for(face &f : faces_inv) { -----
-- for (int j = 0; j < 3; ++j) { -----
-- int a = f.p_idx[j], b = f.p_idx[(j + 1) % 3]; -----
-- if(dead[b][a]) -----
-- add_face(b, a, i); } } -----
- faces.insert( -----
- faces.end(), faces_inv.begin(), faces_inv.end()); } -----
- return faces; } -----
```

3.11. **Delaunay Triangulation**. Simply map each point (x,y) to (x,y,x^2+y^2) , find the 3d convex hull, and drop the 3rd dimension.

3.12. **Point in Polygon**. Check if a point is strictly inside (or on the border) of a polygon in $O(n)$.

```
bool inPolygon(point q, point p[], int n) { -----
- bool in = false; -----
- for (int i = 0, j = n - 1; i < n; j = i++) -----
-- in ^= (((p[i].y > q.y) != (p[j].y > q.y)) && -----
-- q.x < (p[j].x - p[i].x) * (q.y - p[i].y) / -----
-- (p[j].y - p[i].y) + p[i].x); -----
- return in; } -----
bool onPolygon(point q, point p[], int n) { -----
- for (int i = 0, j = n - 1; i < n; j = i++) -----
- if (abs(dist(p[i], q) + dist(p[j], q) -----
-- dist(p[i], p[j])) < EPS) -----
-- return true; -----
- return false; } -----
```

3.13. **Cut Polygon by a Line**. Cut polygon by line \overline{ab} to its left in $O(n)$, such that $\angle abp$ is counter-clockwise.

```
vector<point> cut(point p[],int n,point a,point b) { -----
- vector<point> poly; -----
- for (int i = 0, j = n - 1; i < n; j = i++) { -----
-- double c1 = cross(a, b, p[j]); -----
-- double c2 = cross(a, b, p[i]); -----
-- if (c1 > -EPS) poly.push_back(p[j]); -----
-- if (c1 * c2 < -EPS) -----
-- poly.push_back(line_inter(p[j], p[i], a, b)); -----
- } return poly; } -----
```

3.14. **Triangle Centers**.

```
point bary(point A, point B, point C, -----
-- double a, double b, double c) { -----
- return (A*a + B*b + C*c) / (a + b + c); }
point trilinear(point A, point B, point C, -----
-- double a, double b, double c) { -----
- return bary(A,B,C,abs(B-C)*a, -----
-- abs(C-A)*b,abs(A-B)*c); }
point centroid(point A, point B, point C) { -----
- return bary(A, B, C, 1, 1, 1); }
point circumcenter(point A, point B, point C) { -----
- double a=norm(B-C), b=norm(C-A), c=norm(A-B); -----
- return bary(A,B,C,a*(b+c-a),b*(c+a-b),c*(a+b-c)); }
point orthocenter(point A, point B, point C) { -----
- return bary(A,B,C, tan(angle(B,A,C)), -----
```

```
----- tan(angle(A,B,C)), tan(angle(A,C,B))); } -----
point incenter(point A, point B, point C) { -----
- return bary(A,B,C,abs(B-C),abs(A-C),abs(A-B)); } -----
// incircle radius given the side lengths a, b, c -----
double inradius(double a, double b, double c) { -----
- double s = (a + b + c) / 2; -----
- return sqrt(s * (s-a) * (s-b) * (s-c)) / s; }
point excenter(point A, point B, point C) { -----
- double a = abs(B-C), b = abs(C-A), c = abs(A-B); -----
- return bary(A, B, C, -a, b, c); } -----
// return bary(A, B, C, a, -b, c); -----
// return bary(A, B, C, a, b, -c); -----
point brocard(point A, point B, point C) { -----
- double a = abs(B-C), b = abs(C-A), c = abs(A-B); -----
- return bary(A,B,C,c/b*a,c/b*a,c/a*b,a/b*c); // CCW -----
// return bary(A,B,C,b/c*a,c/a*b,a/b*c); // CW } -----
point symmedian(point A, point B, point C) { -----
- return bary(A,B,C,norm(B-C),norm(C-A),norm(A-B)); } -----
```

3.15. **Convex Polygon Intersection**. Get the intersection of two convex polygons in $O(n^2)$.

```
std::vector<point> convex_polygon_inter( -----
-- point a[], int an, point b[], int bn) { -----
- point ans[an + bn + an*bn]; -----
- int size = 0; -----
- for (int i = 0; i < an; ++i) -----
-- if (inPolygon(a[i],b,bn) || onPolygon(a[i],b,bn)) -----
-- ans[size++] = a[i]; -----
- for (int i = 0; i < bn; ++i) -----
-- if (inPolygon(b[i],a,an) || onPolygon(b[i],a,an)) -----
-- ans[size++] = b[i]; -----
- for (int i = 0, I = an - 1; i < an; I = i++) -----
-- for (int j = 0, J = bn - 1; j < bn; J = j++) { -----
-- try { -----
-- point p=line_inter(a[i],a[I],b[j],b[J],true); -----
-- ans[size++] = p; -----
-- } catch (exception ex) {} } -----
- size = convex_hull(ans, size); -----
- return vector<point>(ans, ans + size); } -----
```

3.16. **Pick's Theorem for Lattice Points**. Count points with integer coordinates inside and on the boundary of a polygon in $O(n)$ using Pick's theorem: $\text{Area} = I + B/2 - 1$.

```
int interior(point p[], int n) { -----
- return area(p,n) - boundary(p,n) / 2 + 1; }
int boundary(point p[], int n) { -----
- int ans = 0; -----
- for (int i = 0, j = n - 1; i < n; j = i++) -----
-- ans += gcd(p[i].x - p[j].x, p[i].y - p[j].y); -----
- return ans; } -----
```

3.17. **Minimum Enclosing Circle**. Get the minimum bounding ball that encloses a set of points (2D or 3D) in Θn .

```
std::pair<point, double> bounding_ball(point p[], int n){ -----
- std::random_shuffle(p, p + n); -----
- point center(0, 0); double radius = 0; -----
- for (int i = 0; i < n; ++i) { -----
```

```
--- if (dist(center, p[i]) > radius + EPS) { -----
----- center = p[i]; radius = 0; -----
    for (int j = 0; j < i; ++j) -----
----- if (dist(center, p[j]) > radius + EPS) { -----
----- center.x = (p[i].x + p[j].x) / 2; -----
----- center.y = (p[i].y + p[j].y) / 2; -----
----- // center.z = (p[i].z + p[j].z) / 2; -----
----- radius = dist(center, p[i]); // midpoint -----
    for (int k = 0; k < j; ++k) -----
----- if (dist(center, p[k]) > radius + EPS) { -----
----- center = circumcenter(p[i], p[j], p[k]); -----
----- radius = dist(center, p[i]); } } } -----
- return {center, radius}; } -----
```

3.18. **Shamos Algorithm.** Solve for the polygon diameter in $O(n \log n)$.

```
double shamos(point p[], int n) { -----
- point *h = new point[n+1]; copy(p, p + n, h); -----
- int k = convex_hull(h, n); if (k <= 2) return 0; -----
- h[k] = h[0]; double d = HUGE_VAL; -----
- for (int i = 0, j = 1; i < k; ++i) { -----
----- while (distPtLine(h[j+1], h[i], h[i+1]) >= -----
----- distPtLine(h[j], h[i], h[i+1])) { -----
----- j = (j + 1) % k; } -----
- d = min(d, distPtLine(h[j], h[i], h[i+1])); -----
- } return d; } -----
```

3.19. **k D Tree.** Get the k -nearest neighbors of a point within pruned radius in $O(k \log k \log n)$.

```
#define cpoint const point& -----
bool cmpx(cpoint a, cpoint b) {return a.x < b.x;} -----
bool cmpy(cpoint a, cpoint b) {return a.y < b.y;} -----
struct KDTree { -----
- KDTree(point p[],int n): p(p), n(n) {build(0,n);} -----
- priority_queue< pair<double, point*> > pq; -----
- point *p; int n, k; double qx, qy, prune; -----
- void build(int L, int R, bool dvx=false) { -----
----- if (L >= R) return; -----
----- int M = (L + R) / 2; -----
----- nth_element(p + L, p + M, p + R, dvx?cmpx:cmpy); -----
----- build(L, M, !dvx); build(M + 1, R, !dvx); } -----
- void dfs(int L, int R, bool dvx) { -----
----- if (L >= R) return; -----
----- int M = (L + R) / 2; -----
----- double dx = qx - p[M].x, dy = qy - p[M].y; -----
----- double delta = dvx ? dx : dy; -----
----- double D = dx * dx + dy * dy; -----
----- if (D<=prune && (pq.size()<k||D<pq.top().first)) { -----
----- pq.push(make_pair(D, &p[M])); -----
----- if (pq.size() > k) pq.pop(); } -----
----- int nL = L, nR = M, fL = M + 1, fR = R; -----
----- if (delta > 0) {swap(nL, fL); swap(nR, fR);} -----
----- dfs(nL, nR, !dvx); -----
----- D = delta * delta; -----
----- if (D<=prune && (pq.size()<k||D<pq.top().first)) -----
----- dfs(fL, fR, !dvx); } -----
- // returns k nearest neighbors of (x, y) in tree -----
- // usage: vector<point> ans = tree.knn(x, y, 2); -----
```

```
- vector<point> knn(double x, double y, -----
----- int k=1, double r=-1) { -----
----- qx=x; qy=y; this->k=k; prune=r<0?HUGE_VAL:r*r; -----
----- dfs(0, n, false); vector<point> v; -----
----- while (!pq.empty()) { -----
----- v.push_back(*pq.top().second); -----
----- pq.pop(); -----
----- } reverse(v.begin(), v.end()); -----
----- return v; } };
```

3.20. **Line Sweep (Closest Pair).** Get the closest pair distance of a set of points in $O(n \log n)$ by sweeping a line and keeping a bounded rectangle. Modifiable for other metrics such as Minkowski and Manhattan distance. For external point queries, see k D Tree.

```
bool cmpy(const point& a, const point& b) { return a.y < b.y; }
double closest_pair_sweep(point p[], int n) { -----
- if (n <= 1) return HUGE_VAL; -----
- std::sort(p, p + n, cmpy); -----
- std::set<point> box; box.insert(p[0]); -----
- double best = 1e13; // infinity, but not HUGE_VAL -----
- for (int L = 0, i = 1; i < n; ++i) { -----
----- while(L < i && p[i].y - p[L].y > best) -----
----- box.erase(p[L++]); -----
----- point bound(p[i].x - best, p[i].y - best); -----
----- std::set<point>::iterator it = box.lower_bound(bound); -----
----- while (it != box.end() && p[i].x+best >= it->x){ -----
----- double dx = p[i].x - it->x; -----
----- double dy = p[i].y - it->y; -----
----- best = std::min(best, std::sqrt(dx*dx + dy*dy)); -----
----- ++it; } -----
----- box.insert(p[i]); -----
- } return best; } -----
```

3.21. **Line upper/lower envelope.** To find the upper/lower envelope of a collection of lines $a_i + b_ix$, plot the points (b_i, a_i) , add the point $(0, \pm\infty)$ (depending on if upper/lower envelope is desired), and then find the convex hull.

3.22. **Formulas.** Let $a = (a_x, a_y)$ and $b = (b_x, b_y)$ be two-dimensional vectors.

- $a \cdot b = |a||b| \cos \theta$, where θ is the angle between a and b .
- $a \times b = |a||b| \sin \theta$, where θ is the signed angle between a and b .
- $a \times b$ is equal to the area of the parallelogram with two of its sides formed by a and b . Half of that is the area of the triangle formed by a and b .
- The line going through a and b is $Ax + By = C$ where $A = b_y - a_y$, $B = a_x - b_x$, $C = Aa_x + Ba_y$.
- Two lines $A_1x + B_1y = C_1$, $A_2x + B_2y = C_2$ are parallel iff. $D = A_1B_2 - A_2B_1$ is zero. Otherwise their unique intersection is $(B_2C_1 - B_1C_2, A_1C_2 - A_2C_1)/D$.
- **Euler's formula:** $V - E + F = 2$
- Side lengths a, b, c can form a triangle iff. $a + b > c$, $b + c > a$ and $a + c > b$.
- Sum of internal angles of a regular convex n -gon is $(n - 2)\pi$.
- **Law of sines:** $\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C}$
- **Law of cosines:** $b^2 = a^2 + c^2 - 2ac \cos B$

- Internal tangents of circles $(c_1, r_1), (c_2, r_2)$ intersect at $(c_1r_2 + c_2r_1)/(r_1 + r_2)$, external intersect at $(c_1r_2 - c_2r_1)/(r_1 + r_2)$.

4. GRAPHS

4.1. Single-Source Shortest Paths.

4.1.1. *Dijkstra.*

```
#include "graph_template_adjlist.cpp" -----
// insert inside graph; needs n, dist[], and adj[] -----
void dijkstra(int s) { -----
- for (int u = 0; u < n; ++u) -----
----- dist[u] = INF; -----
----- dist[s] = 0; -----
----- std::priority_queue<ii, vii, std::greater<ii> > pq; -----
----- pq.push({0, s}); -----
----- while (!pq.empty()) { -----
----- int u = pq.top().second; -----
----- int d = pq.top().first; -----
----- pq.pop(); -----
----- if (dist[u] < d) -----
----- continue; -----
----- dist[u] = d; -----
----- for (auto &e : adj[u]) { -----
----- int v = e.first; -----
----- int w = e.second; -----
----- if (dist[v] > dist[u] + w) { -----
----- dist[v] = dist[u] + w; -----
----- pq.push({dist[v], v}); } } } -----
```

4.1.2. *Bellman-Ford.*

```
#include "graph_template_adjlist.cpp" -----
// insert inside graph; needs n, dist[], and adj[] -----
void bellman_ford(int s) { -----
- for (int u = 0; u < n; ++u) -----
----- dist[u] = INF; -----
----- dist[s] = 0; -----
----- for (int i = 0; i < n-1; ++i) -----
----- for (int u = 0; u < n; ++u) -----
----- for (auto &e : adj[u]) -----
----- if (dist[u] + e.second < dist[e.first]) -----
----- dist[e.first] = dist[u] + e.second; } -----
// you can call this after running bellman_ford() -----
bool has_neg_cycle() { -----
- for (int u = 0; u < n; ++u) -----
----- for (auto &e : adj[u]) -----
----- if (dist[e.first] > dist[u] + e.second) -----
----- return true; -----
- return false; } -----
```

4.1.3. *Shortest Path Faster Algorithm.*

```
#include "graph_template_adjlist.cpp" -----
// insert inside graph; -----
// needs n, dist[], in_queue[], num_vis[], and adj[] -----
bool spfa(int s) { -----
- for (int u = 0; u < n; ++u) { -----
----- dist[u] = INF; -----
----- in_queue[u] = 0; -----
```

```
-- num_vis[u] = 0; }
-- dist[s] = 0;
-- in_queue[s] = 1;
-- bool has_negative_cycle = false;
-- std::queue<int> q; q.push(s);
-- while (not q.empty()) {
--     int u = q.front(); q.pop(); in_queue[u] = 0;
--     if (++num_vis[u] >= n)
--         dist[u] = -INF, has_negative_cycle = true;
--     for (auto &[v, c] : adj[u])
--         if (dist[v] > dist[u] + c) {
--             dist[v] = dist[u] + c;
--             if (!in_queue[v]) {
--                 q.push(v);
--                 in_queue[v] = 1; } } }
-- return has_negative_cycle; }
```

4.2. All-Pairs Shortest Paths.

4.2.1. Floyd-Washall.

```
#include "graph_template_adjmat.cpp"
// insert inside graph; needs n and mat[][]
void floyd_warshall() {
    for (int k = 0; k < n; ++k)
        for (int i = 0; i < n; ++i)
            for (int j = 0; j < n; ++j)
                if (mat[i][k] + mat[k][j] < mat[i][j])
                    mat[i][j] = mat[i][k] + mat[k][j]; }
```

4.3. Strongly Connected Components.

4.3.1. Kosaraju.

```
struct kosaraju_graph {
    int n, *vis;
    vi **adj;
    std::vector<vi> sccs;
    kosaraju_graph(int n) {
        this->n = n;
        vis = new int[n];
        adj = new vi*[2];
        for (int dir = 0; dir < 2; ++dir)
            adj[dir] = new vi[n];
        void add_edge(int u, int v) {
            adj[0][u].push_back(v);
            adj[1][v].push_back(u);
        }
        void dfs(int u, int p, int dir, vi &topo) {
            vis[u] = 1;
            for (int v : adj[dir][u])
                if (!vis[v] && v != p) dfs(v, u, dir, topo);
            topo.push_back(u);
        }
        void kosaraju() {
            vi topo;
            for (int u = 0; u < n; ++u) vis[u] = 0;
            for (int u = 0; u < n; ++u) if (!vis[u]) dfs(u, -1, 0, topo);
            for (int u = 0; u < n; ++u) vis[u] = 0;
            for (int i = n-1; i >= 0; --i) {
                if (!vis[topo[i]]) {
```

```
                sccs.push_back({});
                dfs(topo[i], -1, 1, sccs.back()); } } };
```

4.3.2. Tarjan's Offline Algorithm.

```
int n, id[N], low[N], st[N], in[N], TOP, ID;
int scc[N], SCC_SIZE; // 0 <= scc[u] < SCC_SIZE
vector<int> adj[N]; // 0-based adjlist
void dfs(int u) {
    id[u] = low[u] = ID++;
    st[TOP++] = u; in[u] = 1;
    for (int v : adj[u]) {
        if (id[v] == -1) {
            dfs(v);
            low[u] = min(low[u], low[v]);
        } else if (in[v] == 1)
            low[u] = min(low[u], id[v]);
        if (id[u] == low[u]) {
            int sid = SCC_SIZE++;
            do {
                int v = st[--TOP];
                in[v] = 0; scc[v] = sid;
            } while (st[TOP] != u);
        }
    }
}
void tarjan() { // call tarjan() to load SCC
    memset(id, -1, sizeof(int) * n);
    SCC_SIZE = ID = TOP = 0;
    for (int i = 0; i < n; ++i)
        if (id[i] == -1) dfs(i); }
```

4.4. Minimum Mean Weight Cycle. Run this for each strongly connected component

```
typedef std::vector<double> vd;
double min_mean_cycle(graph &g) {
    double mn = INF;
    std::vector<vd> dp(g.n+1, vd(g.n, mn));
    dp[0][0] = 0;
    for (int k = 1; k <= g.n; ++k)
        for (int u = 0; u < g.n; ++u)
            for (auto &[v, w] : g.adj[u])
                dp[k][v] = std::min(ar[k][v], dp[k-1][u] + w);
    for (int k = 0; k < g.n; ++k) {
        double mx = -INF;
        for (int u = 0; u < g.n; ++u)
            mx = std::max(mx, (dp[g.n][u] - dp[k][u]) / (g.n - k));
        mn = std::min(mn, mx);
    }
    return mn; }
```

4.5. Biconnected Components.

4.5.1. Bridges and Articulation Points.

```
struct graph {
    int n, *disc, *low, TIME;
    vi *adj, stk, articulation_points;
    std::set<ii> bridges;
    vvi comps;
    graph(int n) {
        adj = new vi[n];
        disc = new int[n];
```

```
        low = new int[n];
        void add_edge(int u, int v) {
            adj[u].push_back(v);
            adj[v].push_back(u);
        }
        void _bridges_artics(int u, int p) {
            disc[u] = low[u] = TIME++;
            stk.push_back(u);
            int children = 0;
            bool has_low_child = false;
            for (int v : adj[u]) {
                if (disc[v] == -1) {
                    _bridges_artics(v, u);
                    children++;
                    if (disc[u] < low[v])
                        bridges.insert({std::min(u, v), std::max(u, v)});
                    if (disc[u] <= low[v]) {
                        has_low_child = true;
                        comps.push_back({u});
                        while (comps.back().back() != v and !stk.empty()) {
                            comps.back().push_back(stk.back());
                            stk.pop_back();
                        }
                        low[u] = std::min(low[u], low[v]);
                    } else if (v != p)
                        low[u] = std::min(low[u], disc[v]);
                }
                if ((p == -1 && children >= 2) || (p != -1 && has_low_child))
                    articulation_points.push_back(u);
            }
            void bridges_artics() {
                for (int u = 0; u < n; ++u) disc[u] = -1;
                stk.clear();
                articulation_points.clear();
                bridges.clear();
                comps.clear();
                TIME = 0;
                for (int u = 0; u < n; ++u) if (disc[u] == -1)
                    _bridges_artics(u, -1);
            }
        }
```

4.5.2. Block Cut Tree.

```
// insert inside code for finding articulation points
graph build_block_cut_tree() {
    int bct_n = articulation_points.size() + comps.size();
    vi block_id(n, is_art(n, 0));
    graph tree(bct_n);
    for (int i = 0; i < articulation_points.size(); ++i) {
        block_id[articulation_points[i]] = i;
        is_art[articulation_points[i]] = 1;
    }
    for (int i = 0; i < comps.size(); ++i) {
        int id = i + articulation_points.size();
        for (int u : comps[i])
            if (is_art[u]) tree.add_edge(block_id[u], id);
        else
            block_id[u] = id;
    }
    return tree; }
```


<div>4.5.3. Bridge Tree.</div> <div>// insert inside code for finding bridges ----- // requires union_find and hasher ----- graph build_bridge_tree() { - union_find uf(n); ----- - for (int u = 0; u < n; ++u) { ----- -- for (int v : adj[u]) { ----- ---- ii uv = { std::min(u, v), std::max(u, v) }; ----- ---- if (bridges.find(uv) == bridges.end()) ----- ----- uf.unite(u, v); } } ----- - hasher h; ----- - for (int u = 0; u < n; ++u) ----- -- if (u == uf.find(u)) h.get_hash(u); ----- - int tn = h.h.size(); ----- - graph tree(tn); ----- - for (int i = 0; i < M; ++i) { ----- -- int ui = h.get_hash(uf.find(u)); ----- -- int vi = h.get_hash(uf.find(v)); ----- -- if (ui != vi) tree.add_edge(ui, vi); } ----- - return tree; } -----</div> <div>4.6. Minimum Spanning Tree.</div> <div>4.6.1. Kruskal.</div> <div>#include "graph_template_edgelist.cpp" ----- #include "union_find.cpp" ----- // insert inside graph; needs n, and edges ----- void kruskal(viii &res) { ----- - viii().swap(res); // or use res.clear(); ----- - std::priority_queue<iii, viii, std::greater<iii> > pq; ----- - for (auto &edge : edges) ----- -- pq.push(edge); ----- - union_find uf(n); ----- - while (!pq.empty()) { ----- -- auto node = pq.top(); pq.pop(); ----- -- int u = node.second.first; ----- -- int v = node.second.second; ----- -- if (uf.unite(u, v)) ----- ---- res.push_back(node); } } -----</div> <div>4.6.2. Prim.</div> <div>#include "graph_template_adjlist.cpp" ----- // insert inside graph; needs n, vis[], and adj[] ----- void prim(viii &res, int s=0) { ----- - viii().swap(res); // or use res.clear(); ----- - std::priority_queue<ii, vii, std::greater<ii> > pq; ----- - pq.push({0, s}); ----- - vis[s] = true; ----- - while (!pq.empty()) { ----- -- int u = pq.top().second; pq.pop(); ----- -- vis[u] = true; ----- -- for (auto &[v, w] : adj[u]) { ----- ---- if (v == u) continue; ----- ---- if (vis[v]) continue; ----- ---- res.push_back({w, {u, v}}); ----- ---- pq.push({w, v}); } } } -----</div>	<div>4.7. Euler Path/Cycle .</div> <div>4.7.1. Euler Path/Cycle in a Directed Graph .</div> <div>#define MAXV 1000 ----- #define MAXE 5000 ----- int indeg[MAXV], outdeg[MAXV], res[MAXE + 1]; ii start_end(graph &g) { - int start = -1, end = -1, any = 0, c = 0; - for (int u = 0; u < n; ++u) { -- if (outdeg[u] > 0) any = u; -- if (indeg[u] + 1 == outdeg[u]) start = u, c++; -- else if (indeg[u] == outdeg[u] + 1) end = u, c++; -- else if (indeg[u] != outdeg[u]) return {-1, -1}; } - if ((start == -1) != (end == -1) (c != 2 && c != 0)) -- return {-1, -1}; - if (start == -1) start = end = any; - return {start, end}; } bool euler_path(graph &g) { - ii se = start_end(g); - int cur = se.first, at = g.edges.size() + 1; - if (cur == -1) return false; - std::stack<int> s; - while (true) { -- if (outdeg[cur] == 0) { --- res[--at] = cur; --- if (s.empty()) break; --- cur = s.top(); s.pop(); -- } else s.push(cur), cur = g.adj[cur][--outdeg[cur]]; } - return at == 0; } -----</div> <div>4.7.2. Euler Path/Cycle in an Undirected Graph .</div> <div>std::multiset<int> adj[1010]; std::list<int> L; std::list<int>::iterator euler(- int at, int to, std::list<int>::iterator it) { - if (at == to) return it; - L.insert(it, at), --it; - while (!adj[at].empty()) { -- int nxt = *adj[at].begin(); -- adj[at].erase(adj[at].find(nxt)); -- adj[nxt].erase(adj[nxt].find(at)); -- if (to == -1) { --- it = euler(nxt, at, it); --- L.insert(it, at); --- --it; -- } else { --- it = euler(nxt, to, it); --- to = -1; } } - return it; } // euler(0, -1, L.begin()) -----</div> <div>4.8. Bipartite Matching .</div> <div>4.8.1. Alternating Paths Algorithm .</div>	<div>vi* adj; ----- bool* done; // initially all false ----- int* owner; // initially all -1 ----- int alternating_path(int left) { ----- - if (done[left]) return 0; ----- - done[left] = true; ----- - for (int right : adj[left]) { ----- -- if (owner[right] == -1 alternating_path(owner[right])) { --- owner[right] = left; return 1; } } ----- - return 0; } -----</div> <div>4.8.2. Hopcroft-Karp Algorithm .</div> <div>#define MAXN 5000 ----- int dist[MAXN+1], q[MAXN+1]; ----- #define dist(v) dist[v == -1 ? MAXN : v] ----- struct bipartite_graph { ----- - int n, m, *L, *R; vi *adj; ----- - bipartite_graph(int n, int m) : n(n), m(m), ----- -- L(new int[n]), R(new int[m]), adj(new vi[n]) {} ----- - ~bipartite_graph() { delete[] adj; delete[] L; delete[] R; } ----- - void add_edge(int u, int v) { adj[u].push_back(v); } ----- - bool bfs() { ----- -- int l = 0, r = 0; ----- -- for (int v = 0; v < n; ++v) ----- --- if (L[v] == -1) dist(v) = 0, q[r++] = v; ----- --- else dist(v) = INF; ----- -- dist(-1) = INF; ----- -- while (l < r) { ----- --- int v = q[l++]; ----- --- if (dist(v) < dist(-1)) ----- --- for (int u : adj[v]) ----- --- if (dist(R[u]) == INF) { ----- ----- dist(R[u]) = dist(v) + 1; ----- ----- q[r++] = R[u]; } } ----- -- return dist(-1) != INF; } ----- - bool dfs(int v) { ----- - if (v != -1) { ----- -- for (int u : adj[v]) ----- -- if (dist(R[u]) == dist(v) + 1) ----- -- if (dfs(R[u])) { R[u] = v; L[v] = u; return true; } ----- -- dist(v) = INF; ----- -- return false; } ----- - return true; } ----- - int maximum_matching() { ----- -- int matching = 0; ----- -- for (int u = 0; u < n; ++u) ----- -- L[u] = R[u] = -1; ----- -- while (bfs()) ----- -- for (int u = 0; u < n; ++u) ----- -- matching += L[u] == -1 && dfs(u); ----- -- return matching; } };</div> <div>4.8.3. Minimum Vertex Cover in Bipartite Graphs .</div> <div>#include "hopcroft_karp.cpp" ----- std::vector<bool> alt; ----- void dfs(bipartite_graph &g, int u) { -----</div>
-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------	----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------	-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------

```
- alt[u] = true;
- for (int v: g.adj[u]) {
-   alt[v + g.n] = true;
-   if (g.R[v] != -1 && !alt[g.R[v]])
-     dfs(g, g.R[v]); } }
vi mvc_bipartite(bipartite_graph &g) {
- vi res; g.maximum_matching();
- alt.assign(g.n + g.m, false);
- for (int i = 0; i < g.n; ++i) if (g.L[i] == -1) dfs(g, i);
- for (int i = 0; i < g.n; ++i) if (!alt[i]) res.push_back(i);
- for (int i = 0; i < g.m; ++i)
-   if (alt[g.n + i]) res.push_back(g.n + i);
- return res; }
```

4.9. Maximum Flow.

4.9.1. **Edmonds-Karp**. $O(VE^2)$

4.9.2. *Dinic*. $O(V^2E)$

```
struct flow_network_dinic {
- struct edge {
-   int u, v; ll c, f;
-   edge(int u, int v, ll c) : u(u), v(v), c(c), f(0) {} };
- int n;
- std::vector<int> adj_ptr, par, dist;
- std::vector<std::vector<int>> adj;
- std::vector<edge> edges;
- flow_network_dinic(int n) : n(n) {
-   std::vector<std::vector<int>>(n).swap(adj);
-   reset(); }
- void reset() {
-   std::vector<int>(n).swap(adj_ptr);
-   std::vector<int>(n).swap(par);
-   std::vector<int>(n).swap(dist);
-   for (edge &e : edges) e.f = 0; }
- void add_edge(int u, int v, ll c, bool bi = false) {
-   adj[u].push_back(edges.size());
-   edges.push_back(edge(u, v, c));
-   adj[v].push_back(edges.size());
-   edges.push_back(edge(v, u, (bi ? c : 0LL))); }
- ll res(const edge &e) { return e.c - e.f; }
- bool make_level_graph(int s, int t) {
-   for (int u = 0; u < n; ++u) dist[u] = -1;
-   dist[s] = 0;
-   std::queue<int> q; q.push(s);
-   while (!q.empty()) {
-     int u = q.front(); q.pop();
-     for (int i: adj[u]) {
-       edge &e = edges[i];
-       if (dist[e.v] < 0 and res(e)) {
-         dist[e.v] = dist[u] + 1;
-         q.push(e.v); } } }
-   return dist[t] != -1; }
- bool is_next(int u, int v) {
-   return dist[v] == dist[u] + 1; }
- bool dfs(int u, int t) {
-   if (u == t) return true;
```

```
for (int &ii = adj_ptr[u]; ii < adj[u].size(); ++ii) {
  int i = adj[u][ii];
  edge &e = edges[i];
  if (is_next(u, e.v) and res(e) > 0 and dfs(e.v, t)) {
    par[e.v] = i;
    return true; } }
return false; }
bool aug_path(int s, int t) {
  for (int u = 0; u < n; ++u) par[u] = -1;
  return dfs(s, t); }
ll calc_max_flow(int s, int t) {
  ll total_flow = 0;
  while (make_level_graph(s, t)) {
    for (int u = 0; u < n; ++u) adj_ptr[u] = 0;
    while (aug_path(s, t)) {
      ll flow = pvl::LL_INF;
      for (int i = par[t]; i != -1; i = par[edges[i].u])
        flow = std::min(flow, res(edges[i]));
      for (int i = par[t]; i != -1; i = par[edges[i].u]) {
        edges[i].f += flow;
        edges[i^1].f -= flow; }
      total_flow += flow; } }
  return total_flow; }
std::vector<bool> min_cut(int s, int t) {
  calc_max_flow(s, t);
  assert(!make_level_graph(s, t));
  std::vector<bool> cut_mem(n);
  for (int u = 0; u < n; ++u)
    cut_mem[u] = (dist[u] != -1);
  return cut_mem; } }
```

4.9.3. *Push-relabel*. $\omega(VE + V^2\sqrt{E})$, $O(V^3)$

```
int n;
std::vector<vi> capacity, flow;
vi height, excess;
void push(int u, int v) {
  int d = min(excess[u], capacity[u][v] - flow[u][v]);
  flow[u][v] += d; flow[v][u] -= d;
  excess[u] -= d; excess[v] += d; }
void relabel(int u) {
  int d = INF;
  for (int i = 0; i < n; ++i)
    if (capacity[u][i] - flow[u][i] > 0)
      d = min(d, height[i]);
  if (d < INF) height[u] = d + 1; }
vi find_max_height_vertices(int s, int t) {
  vi max_height;
  for (int i = 0; i < n; ++i) {
    if (i != s && i != t && excess[i] > 0) {
      if (!max_height.empty() && height[i] > height[max_height[0]])
        max_height.clear();
      if (max_height.empty() || height[i] == height[max_height[0]])
        max_height.push_back(i); } }
  return max_height; }
int max_flow(int s, int t) {
  flow.assign(n, vi(n, 0));
```

```
height.assign(n, 0); height[s] = n;
excess.assign(n, 0); excess[s] = INF;
for (int i = 0; i < n; ++i) if (i != s) push(s, i);
vi current;
while (!(current = find_max_height_vertices(s, t)).empty()) {
  for (int i: current) {
    bool pushed = false;
    for (int j = 0; j < n && excess[i]; ++j) {
      if (capacity[i][j] - flow[i][j] > 0 &&
          height[i] == height[j] + 1) {
        push(i, j);
        pushed = true; } }
    if (!pushed) relabel(i), break; } }
  int max_flow = 0;
  for (int i = 0; i < n; ++i) max_flow += flow[i][t];
  return max_flow; }
```

4.9.4. *Gomory-Hu* (*All-pairs Maximum Flow*). $O(V^3E)$, possibly amortized $O(V^2E)$ with a big constant factor.

```
#include "dinic.cpp"
struct gomory_hu_tree {
  int n;
  std::vector<int> dep;
  std::vector<std::pair<int, ll>> par;
  explicit gomory_hu_tree(flow_network_dinic &g) : n(g.n) {
    std::vector<std::pair<int, ll>>(n, {0, 0LL}).swap(par);
    std::vector<int>(n, 0).swap(dep);
    std::vector<int> temp_par(n, 0);
    for (int u = 1; u < n; ++u) {
      g.reset();
      ll flow = g.calc_max_flow(u, temp_par[u]);
      std::vector<bool> cut_mem = g.min_cut(u, temp_par[u]);
      for (int v = u+1; v < n; ++v)
        if (cut_mem[u] == cut_mem[v]
            and temp_par[u] == temp_par[v])
          add_edge(temp_par[u], u, flow); } }
  void add_edge(int u, int v, ll w) {
    par[v] = {u, w}; dep[v] = dep[u] + 1; }
  ll calc_max_flow(int s, int t) {
    ll ans = pvl::LL_INF;
    while (dep[s] > dep[t]) {
      ans = std::min(ans, par[s].second); s = par[s].first; }
    while (dep[s] < dep[t]) {
      ans = std::min(ans, par[t].second); t = par[t].first; }
    while (s != t) {
      ans = std::min(ans, par[s].second); s = par[s].first;
      ans = std::min(ans, par[t].second); t = par[t].first; }
    return ans; } }
```

4.10. Minimum Cost Maximum Flow.

```
struct edge {
  int u, v; ll cost, cap, flow;
  edge(int u, int v, ll cap, ll cost) :
    u(u), v(v), cap(cap), cost(cost), flow(0) {} };
struct flow_network {
  int n, s, t, *par, *in_queue, *num_vis; ll *dist, *pot;
```

```
std::vector<edge> edges;
std::vector<int> *adj;
std::map<std::pair<int, int>, std::vector<int> > edge_idx;
flow_network(int n, int s, int t) : n(n), s(s), t(t) {
    adj = new std::vector<int>[n];
    par = new int[n];
    in_queue = new int[n];
    num_vis = new int[n];
    dist = new ll[n];
    pot = new ll[n];
    for (int u = 0; u < n; ++u) pot[u] = 0; }
void add_edge(int u, int v, ll cap, ll cost) {
    adj[u].push_back(edges.size());
    edge_idx[{u, v}].push_back(edges.size());
    edges.push_back(edge(u, v, cap, cost));
    adj[v].push_back(edges.size());
    edge_idx[{v, u}].push_back(edges.size());
    edges.push_back(edge(v, u, 0LL, -cost)); }
ll get_flow(int u, int v) {
    ll f = 0;
    for (int i : edge_idx[{u, v}]) f += edges[i].flow;
    return f; }
ll res(edge &e) { return e.cap - e.flow; }
void bellman_ford() {
    for (int u = 0; u < n; ++u) pot[u] = INF;
    pot[s] = 0;
    for (int it = 0; it < n-1; ++it)
        for (auto e : edges)
            if (res(e) > 0)
                pot[e.v] = std::min(pot[e.v], pot[e.u] + e.cost); }
bool spfa () {
    std::queue<int> q; q.push(s);
    while (not q.empty()) {
        int u = q.front(); q.pop(); in_queue[u] = 0;
        if (++num_vis[u] >= n) {
            dist[u] = -INF;
            return false; }
        for (int i : adj[u]) {
            edge e = edges[i];
            if (res(e) <= 0) continue;
            ll nd = dist[u] + e.cost + pot[u] - pot[e.v];
            if (dist[e.v] > nd) {
                dist[e.v] = nd;
                par[e.v] = i;
                if (not in_queue[e.v]) {
                    q.push(e.v);
                    in_queue[e.v] = 1; } } } }
    return dist[t] != INF; }
bool aug_path() {
    for (int u = 0; u < n; ++u) {
        par[u] = -1;
        in_queue[u] = 0;
        num_vis[u] = 0;
        dist[u] = INF; }
    dist[s] = 0;
    in_queue[s] = 1;
```

```
return spfa();
}
pll calc_max_flow(bool do_bellman_ford=false) {
    ll total_cost = 0, total_flow = 0;
    if (do_bellman_ford)
        bellman_ford();
    while (aug_path()) {
        ll f = INF;
        for (int i = par[t]; i != -1; i = par[edges[i].u])
            f = std::min(f, res(edges[i]));
        for (int i = par[t]; i != -1; i = par[edges[i].u]) {
            edges[i].flow += f;
            edges[i^1].flow -= f; }
        total_cost += f * (dist[t] + pot[t] - pot[s]);
        total_flow += f;
        for (int u = 0; u < n; ++u)
            if (par[u] != -1) pot[u] += dist[u]; }
    return {total_cost, total_flow}; } }
```

4.10.1. Hungarian Algorithm.

```
int n, m; // size of A, size of B
int cost[N+1][N+1]; // input cost matrix, 1-indexed
int way[N+1], p[N+1]; // p[i]=j: Ai is matched to Bj
int minv[N+1], A[N+1], B[N+1]; bool used[N+1];
int hungarian() {
    for (int i = 0; i <= N; ++i)
        A[i] = B[i] = p[i] = way[i] = 0; // init
    for (int i = 1; i <= n; ++i) {
        p[0] = i; int R = 0;
        for (int j = 0; j <= m; ++j)
            minv[j] = INF, used[j] = false;
        do {
            int L = p[R], dR = 0;
            int delta = INF;
            used[R] = true;
            for (int j = 1; j <= m; ++j)
                if (!used[j]) {
                    int c = cost[L][j] - A[L] - B[j];
                    if (c < minv[j]) minv[j] = c, way[j] = R;
                    if (minv[j] < delta) delta = minv[j], dR = j;
                }
            for (int j = 0; j <= m; ++j)
                if (used[j]) A[p[j]] += delta, B[j] -= delta;
            else minv[j] -= delta;
            R = dR;
        } while (p[R] != 0);
        for (; R != 0; R = way[R])
            p[R] = p[way[R]]; }
    return -B[0]; }
```

4.11. **Minimum Arborescence**. Given a weighted directed graph, finds a subset of edges of minimum total weight so that there is a unique path from the root r to each vertex. Returns a vector of size n , where the i th element is the edge for the i th vertex. The answer for the root is undefined!

```
#include "../data-structures/union_find.cpp"
struct arborescence {
```

```
int n; union_find uf;
vector<vector<pair<ii,int> > > adj;
arborescence(int _n) : n(_n), uf(n), adj(n) { }
void add_edge(int a, int b, int c) {
    adj[b].push_back(make_pair(ii(a,b),c)); }
vii find_min(int r) {
    vi vis(n,-1), mn(n,INF); vii par(n);
    rep(i,0,n) {
        if (uf.find(i) != i) continue;
        int at = i;
        while (at != r && vis[at] == -1) {
            vis[at] = i;
            iter(it,adj[at]) if (it->second < mn[at] &&
                uf.find(it->first.first) != at)
                mn[at] = it->second, par[at] = it->first;
            if (par[at] == ii(0,0)) return vii();
            at = uf.find(par[at].first); }
        if (at == r || vis[at] != i) continue;
        union_find tmp = uf; vi seq;
        do { seq.push_back(at); at = uf.find(par[at].first);
        } while (at != seq.front());
        iter(it,seq) uf.unite(*it,seq[0]);
        int c = uf.find(seq[0]);
        vector<pair<ii,int> > nw;
        iter(it,seq) iter(jt,adj[*it])
            nw.push_back(make_pair(jt->first,
                jt->second - mn[*it]));
        adj[c] = nw;
        vii rest = find_min(r);
        if (size(rest) == 0) return rest;
        ii use = rest[c];
        rest[at = tmp.find(use.second)] = use;
        iter(it,seq) if (*it != at)
            rest[*it] = par[*it];
        return rest; }
    return par; } }
```

4.12. **Blossom algorithm**. Finds a maximum matching in an arbitrary graph in $O(|V|^4)$ time. Be vary of loop edges.

```
#define MAXV 300
bool marked[MAXV], emarked[MAXV][MAXV];
int S[MAXV];
vi find_augmenting_path(const vector<vi> &adj, const vi &m){
    int n = size(adj), s = 0;
    vi par(n,-1), height(n), root(n,-1), q, a, b;
    memset(marked,0,sizeof(marked));
    memset(emarked,0,sizeof(emarked));
    rep(i,0,n) if (m[i] >= 0) emarked[i][m[i]] = true;
    else root[i] = i, S[s++] = i;
    while (s) {
        int v = S[--s];
        iter(wt,adj[v]) {
            int w = *wt;
            if (emarked[v][w]) continue;
            if (root[w] == -1) {
                int x = S[s++] = m[w];
```

```
----- par[w]=v, root[w]=root[v], height[w]=height[v]+1; -----
----- par[x]=w, root[x]=root[w], height[x]=height[w]+1; -----
----- } else if (height[w] % 2 == 0) { -----
----- if (root[v] != root[w]) { -----
----- while (v != -1) q.push_back(v), v = par[v]; -----
----- reverse(q.begin(), q.end()); -----
----- while (w != -1) q.push_back(w), w = par[w]; -----
----- return q; -----
----- } else { -----
----- int c = v; -----
----- while (c != -1) a.push_back(c), c = par[c]; -----
----- c = w; -----
----- while (c != -1) b.push_back(c), c = par[c]; -----
----- while (!a.empty() && !b.empty() && a.back() == b.back()) -----
----- c = a.back(), a.pop_back(), b.pop_back(); -----
----- memset(marked, 0, sizeof(marked)); -----
----- fill(par.begin(), par.end(), 0); -----
----- iter(it, a) par[*it] = 1; iter(it, b) par[*it] = 1; -----
----- par[c] = s = 1; -----
----- rep(i, 0, n) root[par[i]] = par[i] ? 0 : s++; i; -----
----- vector<vi> adj2(s); -----
----- rep(i, 0, n) iter(it, adj[i]) { -----
----- if (par[*it] == 0) continue; -----
----- if (par[i] == 0) { -----
----- if (!marked[par[*it]]) { -----
----- adj2[par[i]].push_back(par[*it]); -----
----- adj2[par[*it]].push_back(par[i]); -----
----- marked[par[*it]] = true; } -----
----- } else adj2[par[i]].push_back(par[*it]); } -----
----- vi m2(s, -1); -----
----- if (m[c] != -1) m2[m2[par[m[c]]] = 0] = par[m[c]]; -----
----- rep(i, 0, n) if (par[i] != 0 && m[i] != -1 && par[m[i]] != 0) -----
----- m2[par[i]] = par[m[i]]; -----
----- vi p = find_augmenting_path(adj2, m2); -----
----- int t = 0; -----
----- while (t < size(p) && p[t]) t++; -----
----- if (t == size(p)) { -----
----- rep(i, 0, size(p)) p[i] = root[p[i]]; -----
----- return p; } -----
----- if (!p[0] || (m[c] != -1 && p[t+1] != par[m[c]])) -----
----- reverse(p.begin(), p.end()), t=(int)size(p)-t-1; -----
----- rep(i, 0, t) q.push_back(root[p[i]]); -----
----- iter(it, adj[root[p[t-1]]) { -----
----- if (par[*it] != (s = 0)) continue; -----
----- a.push_back(c), reverse(a.begin(), a.end()); -----
----- iter(jt, b) a.push_back(*jt); -----
----- while (a[s] != *it) s++; -----
----- if ((height[*it] & 1) ^ (s < (int)size(a) - (int)size(b))) -----
----- reverse(a.begin(), a.end()), s=(int)size(a)-s-1; -----
----- while (a[s] != c) q.push_back(a[s]), s=(s+1)%size(a); -----
----- q.push_back(c); -----
----- rep(i, t+1, size(p)) q.push_back(root[p[i]]); -----
----- return q; } } } -----
----- emarked[v][w] = emarked[w][v] = true; } -----
----- marked[v] = true; } return q; } -----
vii max_matching(const vector<vi> &adj) {
```

```
----- vi m(size(adj), -1), ap; vii res, es; -----
----- rep(i, 0, size(adj)) iter(it, adj[i]) es.emplace_back(i, *it); -----
----- random_shuffle(es.begin(), es.end()); -----
----- iter(it, es) if (m[it->first] == -1 && m[it->second] == -1) -----
----- m[it->first] = it->second, m[it->second] = it->first; -----
----- do { ap = find_augmenting_path(adj, m); -----
----- rep(i, 0, size(ap)) m[m[ap[i^1]] = ap[i]] = ap[i^1]; -----
----- } while (!ap.empty()); -----
----- rep(i, 0, size(m)) if (i < m[i]) res.emplace_back(i, m[i]); -----
----- return res; } -----
```

4.13. **Maximum Density Subgraph.** Given (weighted) undirected graph G . Binary search density. If g is current density, construct flow network: $(S, u, m), (u, T, m + 2g - d_u), (u, v, 1), (v, T, 1)$, where m is a large constant (larger than sum of edge weights). Run floating-point max-flow. If minimum cut has empty S -component, then maximum density is smaller than g , otherwise it's larger. Distance between valid densities is at least $1/(n(n-1))$. Edge case when density is 0. This also works for weighted graphs by replacing d_u by the weighted degree, and doing more iterations (if weights are not integers).

4.14. **Maximum-Weight Closure.** Given a vertex-weighted directed graph G . Turn the graph into a flow network, adding weight ∞ to each edge. Add vertices S, T . For each vertex v of weight w , add edge (S, v, w) if $w \geq 0$, or edge $(v, T, -w)$ if $w < 0$. Sum of positive weights minus minimum $S - T$ cut is the answer. Vertices reachable from S are in the closure. The maximum-weight closure is the same as the complement of the minimum-weight closure on the graph with edges reversed.

4.15. **Maximum Weighted Ind. Set in a Bipartite Graph.** This is the same as the minimum weighted vertex cover. Solve this by constructing a flow network with edges $(S, u, w(u))$ for $u \in L, (v, T, w(v))$ for $v \in R$ and (u, v, ∞) for $(u, v) \in E$. The minimum S, T -cut is the answer. Vertices adjacent to a cut edge are in the vertex cover.

4.16. **Synchronizing word problem.** A DFA has a synchronizing word (an input sequence that moves all states to the same state) iff. each pair of states has a synchronizing word. That can be checked using reverse DFS over pairs of states. Finding the shortest synchronizing word is NP-complete.

4.17. **Max flow with lower bounds on edges.** Change edge $(u, v, l \leq f \leq c)$ to $(u, v, f \leq c - l)$. Add edge (t, s, ∞) . Create super-nodes S, T . Let $M(u) = \sum_v l(v, u) - \sum_v l(u, v)$. If $M(u) < 0$, add edge $(u, T, -M(u))$, else add edge $(S, u, M(u))$. Max flow from S to T . If all edges from S are saturated, then we have a feasible flow. Continue running max flow from s to t in original graph.

4.18. **Tutte matrix for general matching.** Create an $n \times n$ matrix A . For each edge $(i, j), i < j$, let $A_{ij} = x_{ij}$ and $A_{ji} = -x_{ij}$. All other entries are 0. The determinant of A is zero iff. the graph has a perfect matching. A randomized algorithm uses the Schwartz-Zippel lemma to check if it is zero.

4.19. **Heavy Light Decomposition.**

```
#include "segment_tree.cpp" -----
struct heavy_light_tree { -----
- int n, *par, *heavy, *dep, *path_root, *pos; -----
- std::vector<int> *adj; -----
} -----
```

```
----- segtree *segment_tree; -----
----- heavy_light_tree(int n) : n(n) { -----
----- this->adj = new std::vector<int>[n]; -----
----- segment_tree = new segtree(0, n-1); -----
----- par = new int[n]; -----
----- heavy = new int[n]; -----
----- dep = new int[n]; -----
----- path_root = new int[n]; -----
----- pos = new int[n]; } -----
----- void add_edge(int u, int v) { -----
----- adj[u].push_back(v); -----
----- adj[v].push_back(u); } -----
----- void build(int root) { -----
----- for (int u = 0; u < n; ++u) -----
----- heavy[u] = -1; -----
----- par[root] = root; -----
----- dep[root] = 0; -----
----- dfs(root); -----
----- for (int u = 0, p = 0; u < n; ++u) { -----
----- if (par[u] == -1 or heavy[par[u]] != u) { -----
----- for (int v = u; v != -1; v = heavy[v]) { -----
----- path_root[v] = u; -----
----- pos[v] = p++; } } } -----
----- int dfs(int u) { -----
----- int sz = 1; -----
----- int max_subtree_sz = 0; -----
----- for (int v : adj[u]) { -----
----- if (v != par[u]) { -----
----- par[v] = u; -----
----- dep[v] = dep[u] + 1; -----
----- int subtree_sz = dfs(v); -----
----- if (max_subtree_sz < subtree_sz) { -----
----- max_subtree_sz = subtree_sz; -----
----- heavy[u] = v; } -----
----- sz += subtree_sz; } } -----
----- return sz; } -----
----- int query(int u, int v) { -----
----- int res = 0; -----
----- while (path_root[u] != path_root[v]) { -----
----- if (dep[path_root[u]] > dep[path_root[v]]) -----
----- std::swap(u, v); -----
----- res += segment_tree->sum(pos[path_root[v]], pos[v]); -----
----- v = par[path_root[v]]; } -----
----- res += segment_tree->sum(pos[u], pos[v]); -----
----- return res; } -----
----- void update(int u, int v, int c) { -----
----- for (; path_root[u] != path_root[v]; v = par[path_root[v]]) { -----
----- if (dep[path_root[u]] > dep[path_root[v]]) -----
----- std::swap(u, v); -----
----- segment_tree->increase(pos[path_root[v]], pos[v], c); } -----
----- segment_tree->increase(pos[u], pos[v], c); } } } -----
```

4.20. **Centroid Decomposition**

```
#define MAXV 100100 -----
#define LGMAXV 20 -----
int jmp[MAXV][LGMAXV], -----
```

```
- path[MAXV][LGMAXV],
- sz[MAXV], seph[MAXV],
- shortest[MAXV];
struct centroid_decomposition {
- int n; vvi adj;
- centroid_decomposition(int _n) : n(_n), adj(n) {}
- void add_edge(int a, int b) {
-- adj[a].push_back(b); adj[b].push_back(a);
- int dfs(int u, int p) {
-- sz[u] = 1;
-- for (int i = 0; i < adj[u].size())
-- if (adj[u][i] != p) sz[u] += dfs(adj[u][i], u);
-- return sz[u];
- void makepaths(int sep, int u, int p, int len) {
-- jmp[u][seph[sep]] = sep, path[u][seph[sep]] = len;
-- int bad = -1;
-- for (int i = 0; i < adj[u].size()) {
-- if (adj[u][i] == p) bad = i;
-- else makepaths(sep, adj[u][i], u, len + 1);
-- if (p == sep)
-- swap(adj[u][bad], adj[u].back()), adj[u].pop_back();
- void separate(int h=0, int u=0) {
-- dfs(u, -1); int sep = u;
-- down:
-- for (int nxt : adj[sep])
-- if (sz[nxt] < sz[sep] && sz[nxt] > sz[u]/2)
-- sep = nxt, goto down;
-- seph[sep] = h, makepaths(sep, sep, -1, 0);
-- for (int i = 0; i < adj[sep].size())
-- separate(h+1, adj[sep][i]);
- void paint(int u) {
-- for (int h = 0; h < seph[u] + 1)
-- shortest[jmp[u][h]] =
-- std::min(shortest[jmp[u][h]], path[u][h]);
- int closest(int u) {
-- int mn = INF/2;
-- for (int h = 0; h < seph[u] + 1)
-- mn = std::min(mn, path[u][h] + shortest[jmp[u][h]]);
-- return mn; } };
```

4.21. Least Common Ancestor.

4.21.1. Binary Lifting.

```
struct graph {
- int n, logn, *dep, **par;
- std::vector<int> *adj;
- graph(int n, int logn=20) : n(n), logn(logn) {
-- adj = new std::vector<int>[n];
-- dep = new int[n];
-- par = new int*[n];
-- for (int i = 0; i < n; ++i) par[i] = new int[logn];
- void dfs(int u, int p, int d) {
-- dep[u] = d;
-- par[u][0] = p;
-- for (int v : adj[u])
-- if (v != p) dfs(v, u, d+1);
- int ascend(int u, int k) {
```

```
-- for (int i = 0; i < logn; ++i)
-- if (k & (1 << i)) u = par[u][i];
-- return u; }
- int lca(int u, int v) {
-- if (dep[u] > dep[v]) u = ascend(u, dep[u] - dep[v]);
-- if (dep[v] > dep[u]) v = ascend(v, dep[v] - dep[u]);
-- if (u == v) return u;
-- for (int k = logn-1; k >= 0; --k) {
-- if (par[u][k] != par[v][k]) {
-- u = par[u][k]; v = par[v][k]; }
-- return par[u][0];
- bool is_anc(int u, int v) {
-- if (dep[u] < dep[v]) std::swap(u, v);
-- return ascend(u, dep[u] - dep[v]) == v;
- void prep_lca(int root=0) {
-- dfs(root, root, 0);
-- for (int k = 1; k < logn; ++k)
-- for (int u = 0; u < n; ++u)
-- par[u][k] = par[par[u][k-1]][k-1]; } };
```

4.21.2. Euler Tour Sparse Table.

```
struct graph {
- int n, logn, *par, *dep, *first, *lg, **spt;
- vi *adj, euler; // spt size should be ~ 2n
- graph(int n, int logn=20) : n(n), logn(logn) {
-- adj = new vi[n];
-- par = new int[n];
-- dep = new int[n];
-- first = new int[n];
- void add_edge(int u, int v) {
-- adj[u].push_back(v); adj[v].push_back(u);
- void dfs(int u, int p, int d) {
-- dep[u] = d; par[u] = p;
-- first[u] = euler.size();
-- euler.push_back(u);
-- for (int v : adj[u])
-- if (v != p) {
-- dfs(v, u, d+1);
-- euler.push_back(u); }
- void prep_lca(int root=0) {
-- dfs(root, root, 0);
-- int en = euler.size();
-- lg = new int[en+1];
-- lg[0] = lg[1] = 0;
-- for (int i = 2; i <= en; ++i)
-- lg[i] = lg[i >> 1] + 1;
-- spt = new int*[en];
-- for (int i = 0; i < en; ++i) {
-- spt[i] = new int[lg[en]];
-- spt[i][0] = euler[i];
-- for (int k = 0; (2 << k) <= en; ++k)
-- for (int i = 0; i + (2 << k) <= en; ++i)
-- if (dep[spt[i][k]] < dep[spt[i+(1<<k)]] [k]))
-- spt[i][k+1] = spt[i][k];
-- else
-- spt[i][k+1] = spt[i+(1<<k)]] [k]; }
```

```
- int lca(int u, int v) {
-- int a = first[u], b = first[v];
-- if (a > b) std::swap(a, b);
-- int k = lg[b-a+1], ba = b - (1 << k) + 1;
-- if (dep[spt[a][k]] < dep[spt[ba][k]]) return spt[a][k];
-- return spt[ba][k]; } };
```

4.21.3. Tarjan Off-line LCA.

```
#include "data-structures/union-find.cpp"
struct tarjan_lca {
- vi ancestor, answers;
- vvi adj;
- vvi queries;
- std::vector<bool> colored;
- union_find uf;
- tarjan_lca(int n, vvi &adj) : adj(adj), uf(n) {
-- vi(n).swap(ancestor);
-- vvi(n).swap(queries);
-- std::vector<bool>(n, false).swap(colored);
- void query(int x, int y) {
-- queries[x].push_back(ii(y, size(answers)));
-- queries[y].push_back(ii(x, size(answers)));
-- answers.push_back(-1);
- void process(int u) {
-- ancestor[u] = u;
-- for (int v : adj[u]) {
-- process(v);
-- uf.unite(u,v);
-- ancestor[uf.find(u)] = u;
-- colored[u] = true;
-- for (auto &[a, b]: queries[u])
-- if (colored[a]) answers[b] = ancestor[uf.find(a)];
} };
```

4.22. Counting Spanning Trees. Kirchoff's Theorem: The number of spanning trees of any graph is the determinant of any cofactor of the Laplacian matrix in $O(n^3)$.

- (1) Let A be the adjacency matrix.
- (2) Let D be the degree matrix (matrix with vertex degrees on the diagonal).
- (3) Get $D - A$ and delete exactly one row and column. Any row and column will do. This will be the cofactor matrix.
- (4) Get the determinant of this cofactor matrix using Gauss-Jordan.
- (5) Spanning Trees = |cofactor($D - A$)|

4.23. Erdős-Gallai Theorem. A sequence of non-negative integers $d_1 \geq \dots \geq d_n$ can be represented as the degree sequence of finite simple graph on n vertices if and only if $d_1 + \dots + d_n$ is even and the following holds for $1 \leq k \leq n$:

$$\sum_{i=1}^n d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

4.24. Tree Isomorphism.

```
// REQUIREMENT: list of primes pr[], see prime sieve
typedef long long LL;
int pre[N], q[N], path[N]; bool vis[N];
// perform BFS and return the last node visited
```



```
int bfs(int u, vector<int> adj[]) {
    memset(vis, 0, sizeof(vis));
    int head = 0, tail = 0;
    q[tail++] = u; vis[u] = true; pre[u] = -1;
    while (head != tail) {
        u = q[head]; if (++head == N) head = 0;
        for (int i = 0; i < adj[u].size(); ++i) {
            int v = adj[u][i];
            if (!vis[v]) {
                vis[v] = true; pre[v] = u;
                q[tail++] = v; if (tail == N) tail = 0;
            }
        }
    }
    return u;
} // returns the list of tree centers
vector<int> tree_centers(int r, vector<int> adj[]) {
    int size = 0;
    for (int u=bfs(bfs(r, adj), adj); u!=-1; u=pre[u])
        path[size++] = u;
    vector<int> med(1, path[size/2]);
    if (size % 2 == 0) med.push_back(path[size/2-1]);
    return med;
} // returns "unique hashcode" for tree with root u
LL rootcode(int u, vector<int> adj[], int p=-1, int d=15){
    vector<LL> k; int nd = (d + 1) % primes;
    for (int i = 0; i < adj[u].size(); ++i)
        if (adj[u][i] != p)
            k.push_back(rootcode(adj[u][i], adj, u, nd));
    sort(k.begin(), k.end());
    LL h = k.size() + 1;
    for (int i = 0; i < k.size(); ++i)
        h = h * pr[d] + k[i];
    return h;
} // returns "unique hashcode" for the whole tree
LL treecode(int root, vector<int> adj[]) {
    vector<int> c = tree_centers(root, adj);
    if (c.size()==1)
        return (rootcode(c[0], adj) << 1) | 1;
    return (rootcode(c[0],adj)*rootcode(c[1],adj))<<1;
} // checks if two trees are isomorphic
bool isomorphic(int r1, vector<int> adj1[], int r2,
    vector<int> adj2[], bool rooted = false) {
    if (rooted)
        return rootcode(r1, adj1) == rootcode(r2, adj2);
    return treecode(r1, adj1) == treecode(r2, adj2);
}

5. MATH I - ALGEBRA

5.1. Generating Function Manager.
const int DEPTH = 19;
const int ARR_DEPTH = (1<<DEPTH); //approx 5x10^5
const int SZ = 12;
ll temp[SZ][ARR_DEPTH+1];
const ll MOD = 998244353;
struct GF_Manager {
    int tC = 0;
    std::stack<int> to_be_freed;
    const static ll DEPTH = 23;
    ll prim[DEPTH+1], prim_inv[DEPTH+1], two_inv[DEPTH+1];
    ll mod_pow(ll base, ll exp) {
        if(exp==0) return 1;
        if(exp&1) return (base*mod_pow(base,exp-1))%MOD;
        else return mod_pow((base*base)%MOD, exp/2);
    }
    void set_up_primitives() {
        prim[DEPTH] = 31;
        prim_inv[DEPTH] = mod_pow(prim[DEPTH], MOD-2);
        two_inv[DEPTH] = mod_pow(1<<DEPTH,MOD-2);
        for(int n = DEPTH-1; n >= 0; n--) {
            prim[n] = (prim[n+1]*prim[n+1])%MOD;
            prim_inv[n] = mod_pow(prim[n],MOD-2);
            two_inv[n] = mod_pow(1<<n,MOD-2);
        }
    }
    GF_Manager(){ set_up_primitives(); }
    void start_claiming(){ to_be_freed.push(0); }
    ll* claim(){
        ++to_be_freed.top(); assert(tC < SZ); return temp[tC++];
    }
    void end_claiming(){tC-=to_be_freed.top(); to_be_freed.pop();}
    void NTT(ll A[], int n, ll t[],
        bool is_inverse=false, int offset=0) {
        if (n==0) return;
        //Put the evens first, then the odds
        for (int i = 0; i < (1<<(n-1)); i++) {
            t[i] = A[offset+2*i];
            t[i+(1<<(n-1))] = A[offset+2*i+1];
        }
        for(int i = 0; i < (1<<n); i++)
            A[offset+i] = t[i];
        NTT(A, n-1, t, is_inverse, offset);
        NTT(A, n-1, t, is_inverse, offset+(1<<(n-1)));
        ll w1 = (is_inverse ? prim_inv[n] : prim[n]), w = 1;
        for (int i = 0; i < (1<<(n-1)); i++, w=(w*w1)%MOD) {
            t[i] = (A[offset+i] + w*A[offset+(1<<(n-1))+i])%MOD;
            t[i+(1<<(n-1))] = (A[offset+i] -
                w*A[offset+(1<<(n-1))+i])%MOD;
        }
        for (int i = 0; i < (1<<n); i++) A[offset+i] = t[i];
    }
    int add(ll A[], int an, ll B[], int bn, ll C[]) {
        int cn = 0;
        for(int i = 0; i < max(an,bn); i++) {
            C[i] = A[i]+B[i];
            if(C[i] <= -MOD) C[i] += MOD;
            if(MOD <= C[i]) C[i] -= MOD;
            if(C[i]!=0) cn = i;
        }
        return cn;
    }
    int subtract(ll A[], int an, ll B[], int bn, ll C[]) {
        int cn = 0;
        for(int i = 0; i < max(an,bn); i++) {
            C[i] = A[i]-B[i];
            if(C[i] <= -MOD) C[i] += MOD;
            if(MOD <= C[i]) C[i] -= MOD;
            if(C[i]!=0) cn = i;
        }
        return cn+1;
    }
    int scalar_mult(ll v, ll A[], int an, ll C[]) {
        for(int i = 0; i < an; i++) C[i] = (v*A[i])%MOD;
        return v==0 ? 0 : an;
    }
    int mult(ll A[], int an, ll B[], int bn, ll C[]) {
        start_claiming();
        // make sure you've called setup prim first
        // note: an and bn refer to the *number of items in
        // each array*, NOT the degree of the largest term
        int n, degree = an+bn-1;
        for(n=0; (1<<n) < degree; n++);
        ll *tA = claim(), *tB = claim(), *t = claim();
        copy(A,A+an,tA); fill(tA+an,tA+(1<<n),0);
        copy(B,B+bn,tB); fill(tB+bn,tB+(1<<n),0);
        NTT(tA,n,t);
        NTT(tB,n,t);
        for(int i = 0; i < (1<<n); i++)
            tA[i] = (tA[i]*tB[i])%MOD;
        NTT(tA,n,t,true);
        scalar_mult(two_inv[n],tA,degree,C);
        end_claiming();
        return degree;
    }
    int reciprocal(ll F[], int fn, ll R[]) {
        start_claiming();
        ll *tR = claim(), *tempR = claim();
        int n; for(n=0; (1<<n) < fn; n++);
        fill(tempR,tempR+(1<<n),0);
        tempR[0] = mod_pow(F[0],MOD-2);
        for (int i = 1; i <= n; i++) {
            mult(tempR,1<<i,F,1<<i,tR);
            tR[0] -= 2;
            scalar_mult(-1,tR,1<<i,tR);
            mult(tempR,1<<i,tR,1<<i,tempR);
        }
        copy(tempR,tempR+fn,R);
        end_claiming();
        return n;
    }
    int quotient(ll F[], int fn, ll G[], int gn, ll Q[]) {
        start_claiming();
        ll* revF = claim();
        ll* revG = claim();
        ll* tempQ = claim();
        copy(F,F+fn,revF); reverse(revF,revF+fn);
        copy(G,G+gn,revG); reverse(revG,revG+gn);
        int qn = fn-gn+1;
        reciprocal(revG,qn,revG);
        mult(revF,qn,revG,qn,tempQ);
        reverse(tempQ,tempQ+qn);
        copy(tempQ,tempQ+qn,Q);
        end_claiming();
        return qn;
    }
    int mod(ll F[], int fn, ll G[], int gn, ll R[]) {
        start_claiming();
        ll *Q = claim(), *GQ = claim();
        int qn = quotient(F, fn, G, gn, Q);
        int gqn = mult(G, gn, Q, qn, GQ);
        int rn = subtract(F, fn, GQ, gqn, R);
        end_claiming();
        return rn;
    }
    ll horners(ll F[], int fn, ll xi) {
        ll ans = 0;
        for(int i = fn-1; i >= 0; i--)
```

```

---- ans = (ans*xi+F[i]) % MOD;
-- return ans; } };
GF_Manager gfManager;
ll split[DEPTH+1][2*(ARR_DEPTH)+1];
ll Fi[DEPTH+1][2*(ARR_DEPTH)+1];
int bin_splitting(ll a[], int l, int r, int s=0, int offset=0) {
- if(l == r) {
-- split[s][offset] = -a[l]; //x^0
-- split[s][offset+1] = 1; //x^1
-- return 2; }
- int m = (l+r)/2;
- int sz = m-l+1;
- int da = bin_splitting(a, l, m, s+1, offset);
- int db = bin_splitting(a, m+1, r, s+1, offset+(sz<<1));
- return gfManager.mult(split[s+1]+offset, da,
-- split[s+1]+offset+(sz<<1), db, split[s]+offset); }
void multipoint_eval(ll a[], int l, int r, ll F[], int fn,
- ll ans[], int s=0, int offset=0) {
-- if(l == r) {
---- ans[l] = gfManager.horners(F,fn,a[l]);
-- return; }
- int m = (l+r)/2;
- int sz = m-l+1;
- int da = gfManager.mod(F, fn, split[s+1]+offset,
-- sz+1, Fi[s]+offset);
- int db = gfManager.mod(F, fn, split[s+1]+offset+(sz<<1),
-- r-m+1, Fi[s]+offset+(sz<<1));
-- multipoint_eval(a,l,m,Fi[s]+offset,da,ans,s+1,offset);
-- multipoint_eval(a,m+1,r,Fi[s]+offset+(sz<<1),
-- db,ans,s+1,offset+(sz<<1));
}

```

5.2. **Fast Fourier Transform.** Compute the Discrete Fourier Transform (DFT) of a polynomial in $O(n \log n)$ time.

```

struct poly {
-- double a, b;
-- poly(double a=0, double b=0): a(a), b(b) {}
-- poly operator+(const poly& p) const {
---- return poly(a + p.a, b + p.b);}
-- poly operator-(const poly& p) const {
---- return poly(a - p.a, b - p.b);}
-- poly operator*(const poly& p) const {
---- return poly(a*p.a - b*p.b, a*p.b + b*p.a);}
};
void fft(poly in[], poly p[], int n, int s) {
-- if (n < 1) return;
-- if (n == 1) {p[0] = in[0]; return;}
-- n >= 1; fft(in, p, n, s << 1);
-- fft(in + s, p + n, n, s << 1);
-- poly w(1), wn(cos(M_PI/n), sin(M_PI/n));
-- for (int i = 0; i < n; ++i) {
---- poly even = p[i], odd = p[i + n];
---- p[i] = even + w * odd;
---- p[i + n] = even - w * odd;
---- w = w * wn;
-- }
}

```

```

}
void fft(poly p[], int n) {
-- poly *f = new poly[n]; fft(p, f, n, 1);
-- copy(f, f + n, p); delete[] f;
}
void inverse_fft(poly p[], int n) {
-- for(int i=0; i<n; i++) {p[i].b *= -1;} fft(p, n);
-- for(int i=0; i<n; i++) {p[i].a/=n; p[i].b/= -1*n;}
}

```

5.3. **FFT Polynomial Multiplication.** Multiply integer polynomials a, b of size an, bn using FFT in $O(n \log n)$. Stores answer in an array c , rounded to the nearest integer (or double).

```

// note: c[] should have size of at least (an+bn)
int mult(int a[],int an,int b[],int bn,int c[]) {
-- int n, degree = an + bn - 1;
-- for (n = 1; n < degree; n <= 1); // power of 2
-- poly *A = new poly[n], *B = new poly[n];
-- copy(a, a + an, A); fill(A + an, A + n, 0);
-- copy(b, b + bn, B); fill(B + bn, B + n, 0);
-- fft(A, n); fft(B, n);
-- for (int i = 0; i < n; i++) A[i] = A[i] * B[i];
-- inverse_fft(A, n);
-- for (int i = 0; i < degree; i++)
-- c[i] = int(A[i].a + 0.5); // same as round(A[i].a)
-- delete[] A, B; return degree;
}

```

5.4. **Number Theoretic Transform.** Other possible moduli: 2113929217(2^{25}), 2013265920268435457(2^{28} , with $g = 5$)

```

#include "../mathematics/primitive_root.cpp"
int mod = 998244353, g = primitive_root(mod),
- ginv = mod_pow<ll>(g, mod-2, mod),
- inv2 = mod_pow<ll>(2, mod-2, mod);
#define MAXN (1<<22)
struct Num {
-- int x;
-- Num(ll _x=0) { x = (_x%mod+mod)%mod; }
-- Num operator +(const Num &b) { return x + b.x; }
-- Num operator -(const Num &b) const { return x - b.x; }
-- Num operator *(const Num &b) const { return (ll)x * b.x; }
-- Num operator /(const Num &b) const {
-- return (ll)x * b.inv().x; }
-- Num inv() const { return mod_pow<ll>((ll)x, mod-2, mod); }
-- Num pow(int p) const { return mod_pow<ll>((ll)x, p, mod); }
} T1[MAXN], T2[MAXN];
void ntt(Num x[], int n, bool inv = false) {
-- Num z = inv ? ginv : g;
-- z = z.pow((mod - 1) / n);
-- for (ll i = 0, j = 0; i < n; i++) {
-- if (i < j) swap(x[i], x[j]);
-- ll k = n>>1;
-- while (1 <= k && k <= j) j -= k, k >= 1;
-- j += k; }
-- for (int mx = 1, p = n/2; mx < n; mx <= 1, p >= 1) {
-- Num wp = z.pow(p), w = 1;
-- for (int k = 0; k < mx; k++, w = w*wp) {

```

```

-- for (int i = k; i < n; i += mx << 1) {
---- Num t = x[i + mx] * w;
---- x[i + mx] = x[i] - t;
---- x[i] = x[i] + t; } } }
- if (inv) {
-- Num ni = Num(n).inv();
-- rep(i,0,n) { x[i] = x[i] * ni; } } }
void inv(Num x[], Num y[], int l) {
- if (l == 1) { y[0] = x[0].inv(); return; }
- inv(x, y, l>>1);
- // NOTE: maybe l<<2 instead of l<<1
- rep(i,l>>1,l<<1) T1[i] = y[i] = 0;
- rep(i,0,l) T1[i] = x[i];
- ntt(T1, l<<1); ntt(y, l<<1);
- rep(i,0,l<<1) y[i] = y[i]*2 - T1[i] * y[i] * y[i];
- ntt(y, l<<1, true); }
void sqrt(Num x[], Num y[], int l) {
- if (l == 1) { assert(x[0].x == 1); y[0] = 1; return; }
- sqrt(x, y, l>>1);
- inv(y, T2, l>>1);
- rep(i,l>>1,l<<1) T1[i] = T2[i] = 0;
- rep(i,0,l) T1[i] = x[i];
- ntt(T2, l<<1); ntt(T1, l<<1);
- rep(i,0,l<<1) T2[i] = T1[i] * T2[i];
- ntt(T2, l<<1, true);
- rep(i,0,l) y[i] = (y[i] + T2[i]) * inv2; }
// vim: cc=60 ts=2 sts=2 sw=2:

```

5.5. **Polynomial Long Division.** Divide two polynomials A and B to get Q and R , where $\frac{A}{B} = Q + \frac{R}{B}$.

```

typedef vector<double> Poly;
Poly Q, R; // quotient and remainder
void trim(Poly& A) { // remove trailing zeroes
-- while (!A.empty() && abs(A.back()) < EPS)
-- A.pop_back();
}
void divide(Poly A, Poly B) {
-- if (B.size() == 0) throw exception();
-- if (A.size() < B.size()) {Q.clear(); R=A; return;}
-- Q.assign(A.size() - B.size() + 1, 0);
-- Poly part;
-- while (A.size() >= B.size()) {
---- int As = A.size(), Bs = B.size();
---- part.assign(As, 0);
---- for (int i = 0; i < Bs; i++)
---- part[As-Bs+i] = B[i];
---- double scale = Q[As-Bs] = A[As-1] / part[As-1];
---- for (int i = 0; i < As; i++)
---- A[i] -= part[i] * scale;
---- trim(A);
-- } R = A; trim(Q); }

```

5.6. **Matrix Multiplication.** Multiplies matrices $A_{p \times q}$ and $B_{q \times r}$ in $O(n^3)$ time, modulo MOD.

```

long[][] multiply(long A[][], long B[][]) {
-- int p = A.length, q = A[0].length, r = B[0].length;
-- // if(q != B.length) throw new Exception("((");

```

```
--- long AB[][] = new long[p][r]; -----
--- for (int i = 0; i < p; i++) -----
--- for (int j = 0; j < q; j++) -----
--- for (int k = 0; k < r; k++) -----
--- (AB[i][k] += A[i][j] * B[j][k]) %= MOD;
--- return AB; }
```

5.7. **Matrix Power.** Computes for B^e in $O(n^3 \log e)$ time. Refer to Matrix Multiplication.

```
long[][] power(long B[], long e) {
int n = B.length;
long ans[][]= new long[n][n];
for (int i = 0; i < n; i++) ans[i][i] = 1;
while (e > 0) {
if (e % 2 == 1) ans = multiply(ans, b);
b = multiply(b, b); e /= 2;
} return ans;}
```

5.8. **Fibonacci Matrix.** Fast computation for n th Fibonacci $\{F_1, F_2, \dots, F_n\}$ in $O(\log n)$:

$$\begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^n \times \begin{bmatrix} F_2 \\ F_1 \end{bmatrix}$$

5.9. **Gauss-Jordan/Matrix Determinant.** Row reduce matrix A in $O(n^3)$ time. Returns true if a solution exists.

```
boolean gaussJordan(double A[][]) {
int n = A.length, m = A[0].length;
boolean singular = false;
// double determinant = 1;
for (int i=0, p=0; i<n && p<m; i++, p++) {
for (int k = i + 1; k < n; k++) {
if (Math.abs(A[k][p]) > EPS) { // swap
// determinant *= -1;
double t[]=A[i]; A[i]=A[k]; A[k]=t;
break;
}
}
// determinant *= A[i][p];
if (Math.abs(A[i][p]) < EPS)
{ singular = true; i--; continue; }
for (int j = m-1; j >= p; j--) A[i][j]/= A[i][p];
for (int k = 0; k < n; k++) {
if (i == k) continue;
for (int j = m-1; j >= p; j--)
A[k][j] -= A[k][p] * A[i][j];
}
} return !singular; }
```

6. MATH II - COMBINATORICS

6.1. **Lucas Theorem.** Compute $\binom{n}{k} \bmod p$ in $O(p + \log_p n)$ time, where p is a prime.

```
LL f[P], lid; // P: biggest prime
LL lucas(LL n, LL k, int p) {
if (k == 0) return 1;
if (n < p && k < p) {
if (lid != p) {
lid = p; f[0] = 1;
```

```
for (int i = 0; i < p; ++i) f[i]=f[i-1]*i%p;
}
return f[n] * modpow(f[n-k]*f[k]%p, p-2, p) % p;}
return lucas(n/p,k/p,p) * lucas(n%p,k%p,p) % p; }
```

6.2. **Granville's Theorem.** Compute $\binom{n}{k} \bmod m$ (for any m) in $O(m^2 \log^2 n)$ time.

```
def fprime(n, p):
# counts the number of prime divisors of n!
pk, ans = p, 0
while pk <= n:
ans += n // pk
pk *= p
return ans
def granville(n, k, p, E):
# n choose k (mod p^E)
prime_pow = fprime(n,p)-fprime(k,p)-fprime(n-k,p)
if prime_pow >= E: return 0
e = E - prime_pow
pe = p ** e
r, f = n - k, [1]*pe
for i in range(1, pe):
x = i
if x % p == 0:
x = 1
f[i] = f[i-1] * x % pe
numer, denom, negate, ptr = 1, 1, 0, 0
while n:
if f[-1] != 1 and ptr >= e:
negate ^= (n&1) ^ (k&1) ^ (r&1)
numer = numer * f[n%pe] % pe
denom = denom * f[k%pe] % pe * f[r%pe] % pe
n, k, r = n//p, k//p, r//p
ptr += 1
ans = numer * modinv(denom, pe) % pe
if negate and (p != 2 or e < 3):
ans = (pe - ans) % pe
return mod(ans * p**prime_pow, p**E)
```

```
def choose(n, k, m): # generalized (n choose k) mod m
factors, x, p = [], m, 2
while p*p <= x:
e = 0
while x % p == 0:
e += 1
x /= p
if e: factors.append((p, e))
p += 1
if x > 1: factors.append((x, 1))
crt_array = [granville(n,k,p,e) for p, e in factors]
mod_array = [p**e for p, e in factors]
return chinese_remainder(crt_array, mod_array)[0]
```

6.3. **Derangements.** Compute the number of permutations with n elements such that no element is at their original position:

$$D(n) = (n - 1) (D(n - 1) + D(n - 2)) = nD(n - 1) + (-1)^n$$

6.4. **Factoradics.** Convert a permutation of n items to factoradics and vice versa in $O(n \log n)$.

```
// use fenwick tree add, sum, and low code
typedef long long LL;
void factoradic(int arr[], int n) { // 0 to n-1
for (int i = 0; i <=n; i++) fen[i] = 0;
for (int i = 1; i < n; i++) add(i, 1);
for (int i = 0; i < n; i++) {
int s = sum(arr[i]);
add(arr[i], -1); arr[i] = s;
}
void permute(int arr[], int n) { // factoradic to perm
for (int i = 0; i <=n; i++) fen[i] = 0;
for (int i = 1; i < n; i++) add(i, 1);
for (int i = 0; i < n; i++) {
arr[i] = low(arr[i] - 1);
add(arr[i], -1);
}
}
```

6.5. **k th Permutation.** Get the next k th permutation of n items, if exists, using factoradics. All values should be from 0 to $n - 1$. Use factoradics methods as discussed above.

```
std::vector<int> nth_permutation(int cnt, int n) {
std::vector<int> idx(cnt), per(cnt), fac(cnt);
for (int i = 0; i < cnt; ++i) idx[i] = i;
for (int i = 1; i < cnt+1; ++i) fac[i - 1] = n % i, n /= i;
for (int i = cnt - 1; i >= 0; --i)
per[cnt - i - 1] = idx[fac[i]],
idx.erase(idx.begin() + fac[i]);
return per; }
```

6.6. **Catalan Numbers.**

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1}$$

- (1) The number of non-crossing partitions of an n -element set
- (2) The number of expressions with n pairs of parentheses
- (3) The number of ways $n + 1$ factors can be parenthesized
- (4) The number of full binary trees with $n + 1$ leaves
- (5) The number of monotonic lattice paths of an $n \times n$ grid (5-SAT problem)
- (6) The number of triangulations of a convex polygon with $n + 2$ sides (non-rotational)
- (7) The number of permutations $\{1, \dots, n\}$ without a 3-term increasing subsequence
- (8) The number of ways to form a mountain range with n ups and n downs

6.7. **Stirling Numbers.** s_1 : Count the number of permutations of n elements with k disjoint cycles

s_2 : Count the ways to partition a set of n elements into k nonempty subsets

$$s_1(n, k) = \begin{cases} 1 & n = k = 0 \\ s_1(n - 1, k - 1) - (n - 1)s_1(n - 1, k) & n, k > 0 \\ 0 & \text{elsewhere} \end{cases}$$

$$s_2(n,k)=\begin{cases}1&n=k=0\\s_2(n-1,k-1)+ks_2(n-1,k)&n,k>0\\0&\text{elsewhere}\end{cases}$$

6.8. **Partition Function.** Pregenerate the number of partitions of positive integer n with n positive addends.

$$p(n,k)=\begin{cases}1&n=k=0\\0&n<k\\p(n-1,k-1)+p(n-k,k)&n\geq k\end{cases}$$

7. MATH III - NUMBER THEORY

7.1. **Number/Sum of Divisors.** If a number n is prime factorized where $n = p_1^{e_1} \times p_2^{e_2} \times \cdots \times p_k^{e_k}$, where σ_0 is the number of divisors while σ_1 is the sum of divisors:

$$\sum_{d|n} d^k = \sigma_k(n) = \prod \frac{p_i^{k(e_i)+1} - 1}{p_i - 1}$$

Product: $\prod_{d|n} d = n^{\frac{\sigma_1(n)}{2}}$

7.2. **Möbius Sieve.** The Möbius function μ is the Möbius inverse of e such that $e(n) = \sum_{d|n} \mu(d)$.

```
std::bitset<N> is; int mu[N];
void mobiusSieve() {
- for (int i = 1; i < N; ++i) mu[i] = 1;
- for (int i = 2; i < N; ++i) if (!is[i]) {
--- for (int j = i; j < N; j += i) is[j] = 1, mu[j] *= -1;
--- for (ll j = 1LL*i*i; j < N; j += i*i) mu[j] = 0; } }
```

7.3. **Möbius Inversion.** Given arithmetic functions f and g :

$$g(n) = \sum_{d|n} f(d) \iff f(n) = \sum_{d|n} \mu(d) g\left(\frac{n}{d}\right)$$

7.4. **GCD Subset Counting.** Count number of subsets $S \subseteq A$ such that $\gcd(S) = g$ (modifiable).

```
int f[MX+1]; // MX is maximum number of array
long long gcnt[MX+1]; // gcnt[G]: answer when gcd==G
long long C(int f) {return (1LL << f) - 1;}
// f: frequency count
// C(f): # of subsets of f elements (YOU CAN EDIT)
// Usage: int subsets_with_gcd_1 = gcnt[1];
void gcd_counter(int a[], int n) {
- memset(f, 0, sizeof f);
- memset(gcnt, 0, sizeof gcnt);
- int mx = 0;
- for (int i = 0; i < n; ++i) {
----- f[a[i]] += 1;
----- mx = max(mx, a[i]); }
- for (int i = mx; i >= 1; --i) {
--- int add = f[i];
--- long long sub = 0;
--- for (int j = 2*i; j <= mx; j += i) {
----- add += f[j];
----- sub += gcnt[j]; }
--- gcnt[i] = C(add) - sub; }
```

7.5. **Euler Totient.** Counts all integers from 1 to n that are relatively prime to n in $O(\sqrt{n})$ time.

```
ll totient(ll n) {
- if (n <= 1) return 1;
- ll tot = n;
- for (int i = 2; i * i <= n; i++) {
--- if (n % i == 0) tot -= tot / i;
--- while (n % i == 0) n /= i; }
- if (n > 1) tot -= tot / n;
- return tot; }
```

7.6. **Extended Euclidean.** Assigns x, y such that $ax + by = \gcd(a, b)$ and returns $\gcd(a, b)$.

```
ll mod(ll x, ll m) { // use this instead of x % m
- if (m == 0) return 0;
- if (m < 0) m *= -1;
- return (x%m + m) % m; // always nonnegative }
ll extended_euclid(ll a, ll b, ll &x, ll &y) {
- if (b==0) {x = 1; y = 0; return a;}
- ll g = extended_euclid(b, a%b, x, y);
- ll z = x - a/b*y;
- x = y; y = z; return g; }
}
```

7.7. **Modular Exponentiation.** Find $b^e \pmod m$ in $O(\log e)$ time.

```
template <class T>
T mod_pow(T b, T e, T m) {
- T res = T(1);
- while (e) {
--- if (e & T(1)) res = smod(res * b, m);
--- b = smod(b * b, m), e >>= T(1); }
- return res; }
```

7.8. **Modular Inverse.** Find unique x such that $ax \equiv 1 \pmod m$. Returns 0 if no unique solution is found. Please use modulo solver for the non-unique case.

```
ll modinv(ll a, ll m) {
- ll x, y; ll g = extended_euclid(a, m, x, y);
- if (g == 1 || g == -1) return mod(x * g, m);
- return 0; // 0 if invalid }
```

7.9. **Modulo Solver.** Solve for values of x for $ax \equiv b \pmod m$. Returns $(-1, -1)$ if there is no solution. Returns a pair (x, M) where solution is $x \pmod M$.

```
pll modsolver(ll a, ll b, ll m) {
- ll x, y; ll g = extended_euclid(a, m, x, y);
- if (b % g != 0) return {-1, -1};
- return {mod(x*b/g, m/g), abs(m/g)}; }
```

7.10. **Linear Diophantine.** Computes integers x and y such that $ax + by = c$, returns $(-1, -1)$ if no solution. Tries to return positive integer answers for x and y if possible.

```
pll null(-1, -1); // needs extended euclidean
pll diophantine(ll a, ll b, ll c) {
- if (!a && !b) return c ? null : {0, 0};
- if (!a) return c % b ? null : {0, c / b};
- if (!b) return c % a ? null : {c / a, 0}; }
```

```
- ll x, y; ll g = extended_euclid(a, b, x, y);
- if (c % g) return null;
- y = mod(y * (c/g), a/g);
- if (y == 0) y += abs(a/g); // prefer positive sol.
- return {(c - b*y)/a, y}; }
```

7.11. **Chinese Remainder Theorem.** Solves linear congruence $x \equiv b_i \pmod{m_i}$. Returns $(-1, -1)$ if there is no solution. Returns a pair (x, M) where solution is $x \pmod M$.

```
pll chinese(ll b1, ll m1, ll b2, ll m2) {
- ll x, y; ll g = extended_euclid(m1, m2, x, y);
- if (b1 % g != b2 % g) return ii(-1, -1);
- ll M = abs(m1 / g * m2);
- return {mod(mod(x*b2*m1+y*b1*m2, M*g)/g,M), M}; }
ii chinese_remainder(ll b[], ll m[], int n) {
- ii ans(0, 1);
- for (int i = 0; i < n; ++i) {
--- ans = chinese(b[i],m[i],ans.first,ans.second);
--- if (ans.second == -1) break; }
- return ans; }
```

7.11.1. *Super Chinese Remainder.* Solves linear congruence $a_i x \equiv b_i \pmod{m_i}$. Returns $(-1, -1)$ if there is no solution.

```
pll super_chinese(ll a[], ll b[], ll m[], int n) {
- pll ans(0, 1);
- for (int i = 0; i < n; ++i) {
--- pll two = modsolver(a[i], b[i], m[i]);
--- if (two.second == -1) return two;
--- ans = chinese(ans.first, ans.second,
--- two.first, two.second);
--- if (ans.second == -1) break; }
- return ans; }
```

7.12. **Primitive Root.**

```
#include "mod_pow.cpp"
ll primitive_root(ll m) {
- std::vector<ll> div;
- for (ll i = 1; i*i <= m-1; i++) {
--- if ((m-1) % i == 0) {
----- if (i < m) div.push_back(i);
----- if (m/i < m) div.push_back(m/i); } }
- for (int x = 2; x < m; ++x) {
--- bool ok = true;
--- for (int d : div) if (mod_pow<ll>(x, d, m) == 1) {
----- ok = false; break; }
--- if (ok) return x; }
- return -1; }
```

7.13. **Josephus.** Last man standing out of n if every k th is killed. Zero-based, and does not kill 0 on first pass.

```
int J(int n, int k) {
- if (n == 1) return 0;
- if (k == 1) return n-1;
- if (n < k) return (J(n-1,k)+k)%n;
- int np = n - n/k;
- return k*((J(np,k)+np-n%k*np)%np) / (k-1); }
```

7.14. **Number of Integer Points under a Lines.** Count the number of integer solutions to $Ax + By \leq C, 0 \leq x \leq n, 0 \leq y$. In other words, evaluate the sum $\sum_{x=0}^n \left\lfloor \frac{C - Ax}{B} + 1 \right\rfloor$. To count all solutions, let $n = \left\lfloor \frac{C}{A} \right\rfloor$. In any case, it must hold that $C - nA \geq 0$. Be very careful about overflows.

8. MATH IV - NUMERICAL METHODS

8.1. **Fast Square Testing.** An optimized test for square integers.

```
long long M;
void init_is_square() {
- for (int i = 0; i < 64; ++i) M |= 1ULL << (63-(i*i)%64); }
inline bool is_square(ll x) {
- if (x == 0) return true; // XXX
- if ((M << x) >= 0) return false;
- int c = std::__builtin_ctz(x);
- if (c & 1) return false;
- x >>= c;
- if ((x&7) - 1) return false;
- ll r = std::sqrt(x);
- return r*r == x; }
```

8.2. **Simpson Integration.** Use to numerically calculate integrals

```
const int N = 1000 * 1000; // number of steps
double simpson_integration(double a, double b){
- double h = (b - a) / N;
- double s = f(a) + f(b); // a = x_0 and b = x_2n
- for (int i = 1; i <= N - 1; ++i) {
- double x = a + h * i;
- s += f(x) * ((i & 1) ? 4 : 2); }
- s *= h / 3;
- return s; }
```

9. STRINGS

9.1. **Knuth-Morris-Pratt**. Count and find all matches of string f in string s in $O(n)$ time.

```
int par[N]; // parent table
void buildKMP(string& f) {
- par[0] = -1, par[1] = 0;
- int i = 2, j = 0;
- while (i <= f.length()) {
- if (f[i-1] == f[j]) par[i++] = ++j;
- else if (j > 0) j = par[j];
- else par[i++] = 0; } }
std::vector<int> KMP(string& s, string& f) {
- buildKMP(f); // call once if f is the same
- int i = 0, j = 0; vector<int> ans;
- while (i + j < s.length()) {
- if (s[i + j] == f[j]) {
- if (++j == f.length()) {
- ans.push_back(i);
- i += j - par[j];
- if (j > 0) j = par[j]; }
- } else {
- i += j - par[j]; }
```

```
if (j > 0) j = par[j]; }
- } return ans; }
```

9.2. Trie.

```
template <class T>
struct trie {
- struct node {
- map<T, node*> children;
- int prefixes, words;
- node() { prefixes = words = 0; } };
- node* root;
- trie() : root(new node()) { }
- template <class I>
- void insert(I begin, I end) {
- node* cur = root;
- while (true) {
- cur->prefixes++;
- if (begin == end) { cur->words++; break; }
- else {
- T head = *begin;
- typename map<T, node*>::const_iterator it;
- it = cur->children.find(head);
- if (it == cur->children.end()) {
- pair<T, node*> nw(head, new node());
- it = cur->children.insert(nw).first;
- } begin++, cur = it->second; } } }
- template<class I>
- int countMatches(I begin, I end) {
- node* cur = root;
- while (true) {
- if (begin == end) return cur->words;
- else {
- T head = *begin;
- typename map<T, node*>::const_iterator it;
- it = cur->children.find(head);
- if (it == cur->children.end()) return 0;
- begin++, cur = it->second; } } }
- template<class I>
- int countPrefixes(I begin, I end) {
- node* cur = root;
- while (true) {
- if (begin == end) return cur->prefixes;
- else {
- T head = *begin;
- typename map<T, node*>::const_iterator it;
- it = cur->children.find(head);
- if (it == cur->children.end()) return 0;
- begin++, cur = it->second; } } } }
```

9.2.1. Persistent Trie.

```
const int MAX_KIDS = 2;
const char BASE = '0'; // 'a' or 'A'
struct trie {
- int val, cnt;
- std::vector<trie*> kids;
- trie () : val(-1), cnt(0), kids(MAX_KIDS, NULL) {}
- trie (int val) : val(val), cnt(0), kids(MAX_KIDS, NULL) { }
```

```
- trie (int val, int cnt, std::vector<trie*> &n_kids) :
- val(val), cnt(cnt), kids(n_kids) {}
- trie *insert(std::string &s, int i, int n) {
- trie *n_node = new trie(val, cnt+1, kids);
- if (i == n) return n_node;
- if (!n_node->kids[s[i]-BASE])
- n_node->kids[s[i]-BASE] = new trie(s[i]);
- n_node->kids[s[i]-BASE] =
- n_node->kids[s[i]-BASE]->insert(s, i+1, n);
- return n_node; } };
// max xor on a binary trie from version 'a+1' to 'b' (b > a):
int get_max_xor(trie *a, trie *b, int x) {
- int ans = 0;
- for (int i = MAX_BITS; i >= 0; --i) {
- // don't flip the bit for min xor
- int u = ((x & (1 << i)) > 0) ^ 1;
- int res_cnt = (b and b->kids[u] ? b->kids[u]->cnt : 0)
- (a and a->kids[u] ? a->kids[u]->cnt : 0);
- if (res_cnt == 0) u ^= 1;
- ans ^= (u << i);
- if (a) a = a->kids[u];
- if (b) b = b->kids[u]; }
- return ans; }
```

9.3. **Suffix Array**. Construct a sorted catalog of all substrings of s in $O(n \log n)$ time using counting sort.

```
int n, equiv[N+1], suffix[N+1];
ii equiv_pair[N+1];
string T;
void make_suffix_array(string& s) {
- if (s.back()!='$') s += '$';
- n = s.length();
- for (int i = 0; i < n; i++)
- suffix[i] = i;
- sort(suffix, suffix+n, [&s](int i, int j){return s[i] < s[j]});
- int sz = 0;
- for(int i = 0; i < n; i++){
- if(i==0 || s[suffix[i]]!=s[suffix[i-1]])
- ++sz;
- equiv[suffix[i]] = sz; }
- for (int t = 1; t < n; t<=1) {
- for (int i = 0; i < n; i++)
- equiv_pair[i] = {equiv[i],equiv[(i+t)%n]};
- sort(suffix, suffix+n, [](int i, int j) {
- return equiv_pair[i] < equiv_pair[j]});
- int sz = 0;
- for (int i = 0; i < n; i++) {
- if(i==0 || equiv_pair[suffix[i]]!=equiv_pair[suffix[i-1]])
- ++sz;
- equiv[suffix[i]] = sz; } } }
int count_occurences(string& G) { // in string T
- int L = 0, R = n-1;
- for (int i = 0; i < G.length(); i++){
- // lower/upper = first/last time G[i] is
- // the ith character in suffixes from [L,R]
- std::tie(L,R) = {lower(G[i],i,L,R), upper(G[i],i,L,R)}; }
```



```

-- if (L==1 && R==1) return 0; } -----
- return R-L+1; } -----

```

9.4. **Longest Common Prefix**. Find the length of the longest common prefix for every substring in $O(n)$.

```

int lcp[N]; // lcp[i] = LCP(s[sa[i]:], s[sa[i+1]:]) -----
void buildLCP(std::string s) { // build suffix array first -----
- for (int i = 0, k = 0; i < n; i++) { -----
-- if (pos[i] != n - 1) { -----
---- for(int j = sa[pos[i]+1]; s[i+k]==s[j+k];k++); -----
---- lcp[pos[i]] = k; if (k > 0) k--; -----
-- } else { lcp[pos[i]] = 0; } } } -----

```

9.5. **Aho-Corasick Trie**. Find all multiple pattern matches in $O(n)$ time. This is KMP for multiple strings.

```

class Node { -----
- HashMap<Character, Node> next = new HashMap<>(); -----
- Node fail = null; -----
- long count = 0; -----
- public void add(String s) { // adds string to trie -----
-- Node node = this; -----
-- for (char c : s.toCharArray()) { -----
--- if (!node.contains(c)) -----
---- node.next.put(c, new Node()); -----
---- node = node.get(c); -----
-- } node.count++; } -----
- public void prepare() { -----
-- // prepares fail links of Aho-Corasick Trie -----
-- Node root = this; root.fail = null; -----
-- Queue<Node> q = new ArrayDeque<Node>(); -----
-- for (Node child : next.values()) // BFS -----
--- { child.fail = root; q.offer(child); } -----
-- while (!q.isEmpty()) { -----
--- Node head = q.poll(); -----
--- for (Character letter : head.next.keySet()) { -----
---- // traverse upwards to get nearest fail link -----
--- Node p = head; -----
--- Node nextNode = head.get(letter); -----
--- do { p = p.fail; } -----
--- while(p != root && !p.contains(letter)); -----
--- if (p.contains(letter)) { // fail link found -----
--- p = p.get(letter); -----
--- nextNode.fail = p; -----
--- nextNode.count += p.count; -----
--- } else { nextNode.fail = root; } -----
--- q.offer(nextNode); } } } -----
- public BigInteger search(String s) { -----
-- // counts the words added in trie present in s -----
-- Node root = this, p = this; -----
-- BigInteger ans = BigInteger.ZERO; -----
-- for (char c : s.toCharArray()) { -----
--- while (p != root && !p.contains(c)) p = p.fail; -----
--- if (p.contains(c)) { -----
--- p = p.get(c); -----
--- ans = ans.add(BigInteger.valueOf(p.count)); -----
--- } return ans; } -----

```

```

- private Node get(char c) { return next.get(c); } -----
- private boolean contains(char c) { -----
-- return next.containsKey(c); } } -----
// Usage: Node trie = new Node(); -----
// for (String s : dictionary) trie.add(s); -----
// trie.prepare(); BigInteger m = trie.search(str); -----

```

9.6. **Palindromes**.

9.6.1. **Palindromic Tree**. Find lengths and frequencies of all palindromic substrings of a string in $O(n)$ time.

Theorem: there can only be up to n unique palindromic substrings for any string.

```

int par[N*2+1], child[N*2+1][28]; -----
int len[N*2+1], node[N*2+1], cs[N*2+1], size; -----
long long cnt[N + 2]; // count can be very large -----
int newNode(int p = -1) { -----
- cnt[size] = 0; par[size] = p; -----
- len[size] = (p == -1 ? 0 : len[p] + 2); -----
- memset(child[size], -1, sizeof child[size]); -----
- return size++; } -----
int get(int i, char c) { -----
- if (child[i][c] == -1) child[i][c] = newNode(i); -----
- return child[i][c]; } -----
void manachers(char s[]) { -----
- int n = strlen(s), cn = n * 2 + 1; -----
- for (int i = 0; i < n; i++) { -----
-- cs[i * 2] = -1; cs[i * 2 + 1] = s[i]; } -----
- size = n * 2; -----
- int odd = newNode(), even = newNode(); -----
- int cen = 0, rad = 0, L = 0, R = 0; -----
- size = 0; len[odd] = -1; -----
- for (int i = 0; i < cn; i++) -----
-- node[i] = (i % 2 == 0 ? even : get(odd, cs[i])); -----
- for (int i = 1; i < cn; i++) { -----
-- if (i > rad) { L = i - 1; R = i + 1; } -----
-- else { -----
--- int M = cen * 2 - i; // retrieve from mirror -----
--- node[i] = node[M]; -----
--- if (len[node[M]] < rad - i) L = -1; -----
--- else { -----
--- R = rad + 1; L = i * 2 - R; -----
--- while (len[node[i]] > rad - i) -----
--- node[i] = par[node[i]]; } } // expand palindrome -----
-- while (L >= 0 && R < cn && cs[L] == cs[R]) { -----
--- if (cs[L] != -1) node[i] = get(node[i], cs[L]); -----
--- L--, R++; } -----
-- cnt[node[i]]++; -----
-- if (i + len[node[i]] > rad) { -----
--- rad = i + len[node[i]]; cen = i; } } -----
- for (int i = size - 1; i >= 0; --i) -----
- cnt[par[i]] += cnt[i]; // update parent count } -----
int countUniquePalindromes(char s[]) { -----
- manachers(s); return size; } -----
int countAllPalindromes(char s[]) { -----
- manachers(s); int total = 0; -----
- for (int i = 0; i < size; i++) total += cnt[i]; -----

```

```

- return total; } -----
// longest palindrome substring of s -----
std::string longestPalindrome(char s[]) { -----
- manachers(s); -----
- int n = strlen(s), cn = n * 2 + 1, mx = 0; -----
- for (int i = 1; i < cn; i++) -----
-- if (len[node[mx]] < len[node[i]]) -----
--- mx = i; -----
- int pos = (mx - len[node[mx]]) / 2; -----
- return std::string(s + pos, s + pos + len[node[mx]]); } -----

```

9.6.2. **Eertree**.

```

struct node { -----
- int start, end, len, back_edge, *adj; -----
- node() { -----
-- adj = new int[26]; -----
-- for (int i = 0; i < 26; ++i) adj[i] = 0; } -----
- node(int start, int end, int len, int back_edge) : -----
-- start(start), end(end), len(len), back_edge(back_edge) { -----
-- adj = new int[26]; -----
-- for (int i = 0; i < 26; ++i) adj[i] = 0; } } } -----
struct eertree { -----
- int ptr, cur_node; -----
- std::vector<node> tree; -----
- eertree() { -----
-- tree.push_back(node()); -----
-- tree.push_back(node(0, 0, -1, 1)); -----
-- tree.push_back(node(0, 0, 0, 1)); -----
-- cur_node = 1; -----
-- ptr = 2; } -----
- int get_link(int temp, std::string &s, int i) { -----
-- while (true) { -----
--- int cur_len = tree[temp].len; -----
--- // don't return immediately if you want to -----
--- // get all palindromes; not recommended -----
--- if (i-cur_len-1 >= 0 and s[i] == s[i-cur_len-1]) -----
--- return temp; -----
--- temp = tree[temp].back_edge; } -----
-- return temp; } -----
- void insert(std::string &s, int i) { -----
-- int temp = cur_node; -----
-- temp = get_link(temp, s, i); -----
-- if (tree[temp].adj[s[i] - 'a'] != 0) { -----
--- cur_node = tree[temp].adj[s[i] - 'a']; -----
--- return; } -----
-- ptr++; -----
-- tree[temp].adj[s[i] - 'a'] = ptr; -----
-- int len = tree[temp].len + 2; -----
-- tree.push_back(node(i-len+1, i, len, 0)); -----
-- temp = tree[temp].back_edge; -----
-- cur_node = ptr; -----
-- if (tree[cur_node].len == 1) { -----
--- tree[cur_node].back_edge = 2; -----
--- return; } -----
-- temp = get_link(temp, s, i); -----
-- tree[cur_node].back_edge = tree[temp].adj[s[i] - 'a']; } -----

```

```
- void insert(std::string &s) { -----
--- for (int i = 0; i < s.size(); ++i) -----
----- insert(s, i); } }; -----
----- for (int j = 0; j < s.size(); ++j) -----
----- h_ans[i][j+1] = (h_ans[i][j] + -----
----- s[j] * p_pow[i][j]) % MOD; } } }; ---
```

9.7. **Z Algorithm**. Find the longest common prefix of all substrings of s with itself in $O(n)$ time.

```
int z[N]; // z[i] = lcp(s, s[i:]) -----
void computeZ(string s) { -----
- int n = s.length(), L = 0, R = 0; z[0] = n; -----
- for (int i = 1; i < n; i++) { -----
--- if (i > R) { -----
----- L = R = i; -----
----- while (R < n && s[R - L] == s[R]) R++; -----
----- z[i] = R - L; R--; -----
--- } else { -----
----- int k = i - L; -----
----- if (z[k] < R - i + 1) z[i] = z[k]; -----
----- else { -----
----- L = i; -----
----- while (R < n && s[R - L] == s[R]) R++; -----
----- z[i] = R - L; R--; } } } } -----
```

9.8. **Booth's Minimum String Rotation**. Booth's Algo: Find the index of the lexicographically least string rotation in $O(n)$ time.

```
int f[N * 2]; -----
int booth(string S) { -----
- S.append(S); // concatenate itself -----
- int n = S.length(), i, j, k = 0; -----
- memset(f, -1, sizeof(int) * n); -----
- for (j = 1; j < n; j++) { -----
--- i = f[j-k-1]; -----
--- while (i != -1 && S[j] != S[k + i + 1]) { -----
----- if (S[j] < S[k + i + 1]) k = j - i - 1; -----
----- i = f[i]; -----
--- } if (i == -1 && S[j] != S[k + i + 1]) { -----
----- if (S[j] < S[k + i + 1]) k = j; -----
----- f[j - k] = -1; -----
--- } else f[j - k] = i + 1; -----
- } return k; } -----
```

9.9. Hashing.

9.9.1. Rolling Hash.

```
int MAXN = 1e5+1, MOD = 1e9+7; -----
struct hasher { -----
- int n; -----
- std::vector<ll> *p_pow, *h_ans; -----
- hash(vi &s, vi primes) : n(primes.size()) { -----
--- p_pow = new std::vector<ll>[n]; -----
--- h_ans = new std::vector<ll>[n]; -----
--- for (int i = 0; i < n; ++i) { -----
----- p_pow[i] = std::vector<ll>(MAXN); -----
----- p_pow[i][0] = 1; -----
----- for (int j = 0; j+1 < MAXN; ++j) -----
----- p_pow[i][j+1] = (p_pow[i][j] * primes[i]) % MOD; -----
----- h_ans[i] = std::vector<ll>(MAXN); -----
----- h_ans[i][0] = 0; -----
```