

KFC  
AdMU Progar

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1. DATA STRUCTURES

1.1. Fenwick Tree.

```
struct fenwick {
- vi ar;
- fenwick(vi &ar) : ar(ar.size(), 0) {
- for (int i = 0; i < ar.size(); ++i) {
- ar[i] += ar[i];
- int j = i | (i+1);
- if (j < ar.size())
- ar[j] += ar[i]; } }
- int sum(int i) {
- int res = 0;
- for (; i >= 0; i = (i & (i+1)) - 1)
- res += ar[i];
- return res; }
- int sum(int i, int j) { return sum(j) - sum(i-1); }
- void add(int i, int val) {
- for (; i < ar.size(); i |= i+1)
- ar[i] += val; }
- int get(int i) {
- int res = ar[i];
- if (i) {
- int lca = (i & (i+1)) - 1;
- for (--i; i != lca; i = (i&(i+1))-1)
- res -= ar[i]; }
- return res; }
- void set(int i, int val) { add(i, -get(i) + val); }
- // range update, point query //
- void add(int i, int j, int val) {
- add(i, val); add(j+1, -val); }
- int get1(int i) { return sum(i); } };
```

1.2. Leq Counter.

1.2.1. Leq Counter Array.

```
#include "segtree.cpp"
struct LeqCounter {
- segtree **roots;
- LeqCounter(int *ar, int n) {
- std::vector<ii> nums;
- for (int i = 0; i < n; ++i)
- nums.push_back({ar[i], i});
- std::sort(nums.begin(), nums.end());
- roots = new segtree*[n];
- roots[0] = new segtree(0, n);
- int prev = 0;
- for (ii &e : nums) {
- for (int i = prev+1; i < e.first; ++i)
- roots[i] = roots[prev];
- roots[e.first] = roots[prev]->update(e.second, 1);
```

```
prev = e.first; }
- for (int i = prev+1; i < n; ++i)
- roots[i] = roots[prev]; }
- int count(int i, int j, int x) {
- return roots[x]->query(i, j); } };
```

1.2.2. Leq Counter Map.

```
struct LeqCounter {
- std::map<int, segtree*> roots;
- std::set<int> neg_nums;
- LeqCounter(int *ar, int n) {
- std::vector<ii> nums;
- for (int i = 0; i < n; ++i) {
- nums.push_back({ar[i], i});
- neg_nums.insert(-ar[i]);
- }
- std::sort(nums.begin(), nums.end());
- roots[0] = new segtree(0, n);
- int prev = 0;
- for (ii &e : nums) {
- roots[e.first] = roots[prev]->update(e.second, 1);
- prev = e.first; } }
- int count(int i, int j, int x) {
- auto it = neg_nums.lower_bound(-x);
- if (it == neg_nums.end()) return 0;
- return roots[-*it]->query(i, j); } };
```

1.3. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest element.

```
#define BITS 15
struct misof_tree {
- int cnt[BITS][1<<BITS];
- misof_tree() { memset(cnt, 0, sizeof(cnt)); }
- void insert(int x) {
- for (int i = 0; i < BITS; cnt[i++][x]++, x >= 1); }
- void erase(int x) {
- for (int i = 0; i < BITS; cnt[i++][x]--, x >= 1); }
- int nth(int n) {
- int res = 0;
- for (int i = BITS-1; i >= 0; i--)
- if (cnt[i][res <= 1] <= n) n -= cnt[i][res], res |= 1;
- return res; } };
```

1.4. Mo's Algorithm.

```
struct query {
- int id, l, r; ll hilbert_index;
- query(int id, int l, int r) : id(id), l(l), r(r) {
- hilbert_index = hilbert_order(l, r, LOGN, 0); }
- ll hilbert_order(int x, int y, int pow, int rotate) {
- if (pow == 0) return 0;
- int hpow = 1 << (pow-1);
- int seg = ((x<hpow) ? ((y<hpow)?0:3) : ((y<hpow)?1:2));
- seg = (seg + rotate) & 3;
- const int rotate_delta[4] = {3, 0, 0, 1};
- int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
- int nrot = (rotate + rotate_delta[seg]) & 3;
- ll sub_sq_size = ll(1) << (2*pow - 2);
```

```
ll ans = seg * sub_sq_size;
- ll add = hilbert_order(nx, ny, pow-1, nrot);
- ans += (seg==1 || seg==2) ? add : (sub_sq_size-add-1);
- return ans; }
- bool operator<(const query& other) const {
- return this->hilbert_index < other.hilbert_index; } };
std::vector<query> queries;
for(const query &q : queries) { // [l,r] inclusive
- for(; r > q.r; r--) update(r, -1);
- for(r = r+1; r <= q.r; r++) update(r);
- r--;
- for(; l < q.l; l++) update(l, -1);
- for(l = l-1; l >= q.l; l--) update(l);
- l++; }
```

1.5. Ordered Statistics Tree.

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <typename T>
using index_set = tree<T, null_type, std::less<T>,
splay_tree_tag, tree_order_statistics_node_update>;
// indexed_set<int> t; t.insert(...);
// t.find_by_order(index); // 0-based
// t.order_of_key(key);
```

1.6. Segment Tree.

1.6.1. Recursive, Point-update Segment Tree.

1.6.2. Iterative, Point-update Segment Tree.

```
struct segtree {
- int n;
- int *vals;
- segtree(vi &ar, int n) {
- this->n = n;
- vals = new int[2*n];
- for (int i = 0; i < n; ++i)
- vals[i+n] = ar[i];
- for (int i = n-1; i > 0; --i)
- vals[i] = vals[i<<1] + vals[i<<1|1]; }
- void update(int i, int v) {
- for (vals[i += n] += v; i > 1; i >= 1)
- vals[i>>1] = vals[i] + vals[i^1]; }
- int query(int l, int r) {
- int res = 0;
- for (l += n, r += n+1; l < r; l >= 1, r >= 1) {
- if (l&1) res += vals[l++];
- if (r&1) res += vals[--r]; }
- return res; } };
```

1.6.3. *Pointer-based, Range-update Segment Tree.*

```

struct segtree {
-   int i, j, val, temp_val = 0;
-   segtree *l, *r;
-   segtree(vi &ar, int _i, int _j) : i(_i), j(_j) {
-   if (i == j) {
-       val = ar[i];
-       l = r = NULL;
-   } else {
-       int k = (i + j) >> 1;
-       l = new segtree(ar, i, k);
-       r = new segtree(ar, k+1, j);
-       val = l->val + r->val; } }
-   void visit() {
-   if (temp_val) {
-       val += (j-i+1) * temp_val;
-       if (l) {
-           l->temp_val += temp_val;
-           r->temp_val += temp_val; }
-       temp_val = 0; } }
-   void increase(int _i, int _j, int _inc) {
-   visit();
-   if (_i <= i && j <= _j) {
-       temp_val += _inc;
-       visit();
-   } else if (_j < i or j < _i) {
-       // do nothing
-   } else {
-       l->increase(_i, _j, _inc);
-       r->increase(_i, _j, _inc);
-       val = l->val + r->val; } }
-   int query(int _i, int _j) {
-   visit();
-   if (_i <= i and j <= _j)
-       return val;
-   else if (_j < i || j < _i)
-       return 0;
-   else
-       return l->query(_i, _j) + r->query(_i, _j);
} };

```

1.6.4. *Array-based, Range-update Segment Tree* -.

```

struct segtree {
-   int n, *vals, *deltas;
-   segtree(vi &ar) {
-   n = ar.size();
-   vals = new int[4*n];
-   deltas = new int[4*n];
-   build(ar, 1, 0, n-1); }
-   void build(vi &ar, int p, int i, int j) {
-   deltas[p] = 0;
-   if (i == j)
-       vals[p] = ar[i];
-   else {
-       int k = (i + j) / 2;
-       build(ar, p<<1, i, k);

```

```

-       build(ar, p<<1|1, k+1, j);
-       pull(p); } }
-   void pull(int p) { vals[p] = vals[p<<1] + vals[p<<1|1]; }
-   void push(int p, int i, int j) {
-   if (deltas[p]) {
-       vals[p] += (j - i + 1) * deltas[p];
-       if (i != j) {
-           deltas[p<<1] += deltas[p];
-           deltas[p<<1|1] += deltas[p]; }
-       deltas[p] = 0; } }
-   void update(int _i, int _j, int v, int p, int i, int j) {
-   push(p, i, j);
-   if (_i <= i && j <= _j) {
-       deltas[p] += v;
-       push(p, i, j);
-   } else if (_j < i || j < _i) {
-       // do nothing
-   } else {
-       int k = (i + j) / 2;
-       update(_i, _j, v, p<<1, i, k);
-       update(_i, _j, v, p<<1|1, k+1, j);
-       pull(p); } }
-   int query(int _i, int _j, int p, int i, int j) {
-   push(p, i, j);
-   if (_i <= i and j <= _j)
-       return vals[p];
-   else if (_j < i || j < _i)
-       return 0;
-   else {
-       int k = (i + j) / 2;
-       return query(_i, _j, p<<1, i, k) +
-          query(_i, _j, p<<1|1, k+1, j); } } };

```

1.6.5. *2D Segment Tree.*

```

struct segtree_2d {
-   int n, m, **ar;
-   segtree_2d(int n, int m) {
-   this->n = n;   this->m = m;
-   ar = new int[n];
-   for (int i = 0; i < n; ++i) {
-       ar[i] = new int[m];
-       for (int j = 0; j < m; ++j)
-           ar[i][j] = 0; } }
-   void update(int x, int y, int v) {
-   ar[x + n][y + m] = v;
-   for (int i = x + n; i > 0; i >= 1) {
-       for (int j = y + m; j > 0; j >= 1) {
-           ar[i>>1][j] = min(ar[i][j], ar[i^1][j]);
-           ar[i][j>>1] = min(ar[i][j], ar[i][j^1]);
-       } } // just call update one by one to build
-   int query(int x1, int x2, int y1, int y2) {
-   int s = INF;
-   if (~x2) for (int a=x1+n, b=x2+n+1; a<b; a>=1, b>=1) {
-       if (a & 1) s = min(s, query(a++, -1, y1, y2));
-       if (b & 1) s = min(s, query(--b, -1, y1, y2));
-   } else for (int a=y1+m, b=y2+m+1; a<b; a>=1, b>=1) {

```

```

-       if (a & 1) s = min(s, ar[x1][a++]);
-       if (b & 1) s = min(s, ar[x1][--b]);
-   } return s; } };

```

1.6.6. *Persistent Segment Tree.*

```

struct segtree {
-   int i, j, val;
-   segtree *l, *r;
-   segtree(vi &ar, int _i, int _j) : i(_i), j(_j) {
-   if (i == j) {
-       val = ar[i];
-       l = r = NULL;
-   } else {
-       int k = (i+j) >> 1;
-       l = new segtree(ar, i, k);
-       r = new segtree(ar, k+1, j);
-       val = l->val + r->val;
-   } }
-   segtree(int i, int j, segtree *l, segtree *r, int val) :
-   i(i), j(j), l(l), r(r), val(val) {
-   segtree* update(int _i, int _val) {
-   if (_i <= i and j <= _i)
-       return new segtree(i, j, l, r, val + _val);
-   else if (_i < i or j < _i)
-       return this;
-   else {
-       segtree *nl = l->update(_i, _val);
-       segtree *nr = r->update(_i, _val);
-       return new segtree(i, j, nl, nr, nl->val + nr->val); } }
-   int query(int _i, int _j) {
-   if (_i <= i and j <= _j)
-       return val;
-   else if (_j < i or j < _i)
-       return 0;
-   else
-       return l->query(_i, _j) + r->query(_i, _j); } };

```

1.7. *Sparse Table.*1.7.1. *1D Sparse table.*

```

int lg[MAXN+1], spt[20][MAXN];
void build(vi &arr, int n) {
-   lg[0] = lg[1] = 0;
-   for (int i = 2; i <= n; ++i) lg[i] = lg[i>>1] + 1;
-   for (int i = 0; i < n; ++i) spt[0][i] = arr[i];
-   for (int j = 0; (2 <= j) <= n; ++j)
-       for (int i = 0; i + (2 <= j) <= n; ++i)
-           spt[j+1][i] = std::min(spt[j][i], spt[j][i+(1<<j)]); }
int query(int a, int b) {
-   int k = lg[b-a+1], ab = b - (1<<k) + 1;
-   return std::min(spt[k][a], spt[k][ab]); }

```

1.7.2. **2D Sparse Table**.

```

const int N = 100, LGN = 20;
int lg[N], A[N][N], st[LGN][LGN][N][N];
void build(int n, int m) {
-   for (int k=2; k<=std::max(n,m); ++k) lg[k] = lg[k>>1]+1;

```

```

- for(int i = 0; i < n; ++i)
- - for(int j = 0; j < m; ++j)
- - - st[0][0][i][j] = A[i][j];
- for(int bj = 0; (2 << bj) <= m; ++bj)
- - for(int j = 0; j + (2 << bj) <= m; ++j)
- - - for(int i = 0; i < n; ++i)
- - - - st[0][bj+1][i][j] =
- - - - std::max(st[0][bj][i][j],
- - - - st[0][bj][i][j + (1 << bj)]);
- for(int bi = 0; (2 << bi) <= n; ++bi)
- - for(int i = 0; i + (2 << bi) <= n; ++i)
- - - for(int j = 0; j < m; ++j)
- - - - st[bi+1][0][i][j] =
- - - - std::max(st[bi][0][i][j],
- - - - st[bi][0][i + (1 << bi)][j]);
- for(int bi = 0; (2 << bi) <= n; ++bi)
- - for(int i = 0; i + (2 << bi) <= n; ++i)
- - - for(int bj = 0; (2 << bj) <= m; ++bj)
- - - - for(int j = 0; j + (2 << bj) <= m; ++j) {
- - - - - int ik = i + (1 << bi);
- - - - - int jk = j + (1 << bj);
- - - - - st[bi+1][bj+1][i][j] =
- - - - - std::max(std::max(st[bi][bj][i][j],
- - - - - st[bi][bj][ik][j]),
- - - - - std::max(st[bi][bj][i][jk],
- - - - - st[bi][bj][ik][jk])); } }
int query(int x1, int x2, int y1, int y2) {
- int kx = lg[x2 - x1 + 1], ky = lg[y2 - y1 + 1];
- int x12 = x2 - (1<<kx) + 1, y12 = y2 - (1<<ky) + 1;
- return std::max(std::max(st[kx][ky][x1][y1],
- - st[kx][ky][x1][y12]),
- - std::max(st[kx][ky][x12][y1],
- - st[kx][ky][x12][y12])); }

```

## 1.8. Splay Tree.

```

struct node *null;
struct node {
- node *left, *right, *parent;
- bool reverse; int size, value;
- node& get(int d) {return d == 0 ? left : right;}
- node(int v=0): reverse(0), size(0), value(v) {
- left = right = parent = null ? null : this; } }
struct SplayTree {
- node *root;
- SplayTree(int arr[] = NULL, int n = 0) {
- if (!null) null = new node();
- root = build(arr, n); }
- node* build(int arr[], int n) {
- if (n == 0) return null;
- int mid = n >> 1;
- node *p = new node(arr ? arr[mid] : 0);
- link(p, build(arr, mid), 0);
- link(p, build(arr? arr+mid+1 : NULL, n-mid-1), 1);
- pull(p); return p; }
- void pull(node *p) {
- p->size = p->left->size + p->right->size + 1; }

```

```

- void push(node *p) {
- if (p != null && p->reverse) {
- swap(p->left, p->right);
- p->left->reverse ^= 1;
- p->right->reverse ^= 1;
- p->reverse ^= 1; } }
- void link(node *p, node *son, int d) {
- p->get(d) = son;
- son->parent = p; }
- int dir(node *p, node *son) {
- return p->left == son ? 0 : 1; }
- void rotate(node *x, int d) {
- node *y = x->get(d), *z = x->parent;
- link(x, y->get(d ^ 1), d);
- link(y, x, d ^ 1);
- link(z, y, dir(z, x));
- pull(x); pull(y); }
- node* splay(node *p) {
- while (p->parent != null) {
- node *m = p->parent, *g = m->parent;
- push(g); push(m); push(p);
- int dm = dir(m, p), dg = dir(g, m);
- if (g == null) rotate(m, dm);
- else if (dm == dg) rotate(g, dg), rotate(m, dm);
- else rotate(m, dm), rotate(g, dg);
- } return root = p; }
- node* get(int k) {
- node *p = root;
- while (push(p), p->left->size != k) {
- if (k < p->left->size) p = p->left;
- else k -= p->left->size + 1, p = p->right; }
- return p == null ? null : splay(p); }
- void split(node *&r, int k) {
- if (k == 0) { r = root; root = null; return; }
- r = get(k - 1)->right;
- root->right = r->parent = null;
- pull(root); }
- void merge(node *r) {
- if (root == null) { root = r; return; }
- link(get(root->size - 1), r, 1);
- pull(root); }
- void assign(int k, int val) {
- get(k)->value = val; pull(root); }
- void reverse(int L, int R) {
- node *m, *r; split(r, R + 1); split(m, L);
- m->reverse ^= 1; push(m); merge(r); }
- node* insert(int k, int v) {
- node *r; split(r, k);
- node *p = new node(v); p->size = 1;
- link(root, p, 1); merge(r);
- return p; }
- void erase(int k) {
- node *r, *m;
- split(r, k + 1); split(m, k);
- merge(r); delete m; } }

```

## 1.9. Treap.

### 1.9.1. Implicit Treap.

```

struct cartree {
- typedef struct _Node {
- int node_val, subtree_val, delta, prio, size;
- _Node *l, *r;
- _Node(int val) : node_val(val), subtree_val(val),
- delta(0), prio((rand()<<16)^rand()), size(1),
- l(NULL), r(NULL) {}
- ~_Node() { delete l; delete r; }
- } *Node;
- int get_subtree_val(Node v) {
- return v ? v->subtree_val : 0; }
- int get_size(Node v) { return v ? v->size : 0; }
- void apply_delta(Node v, int delta) {
- if (!v) return;
- v->delta += delta;
- v->node_val += delta;
- v->subtree_val += delta * get_size(v); }
- void push_delta(Node v) {
- if (!v) return;
- apply_delta(v->l, v->delta);
- apply_delta(v->r, v->delta);
- v->delta = 0; }
- void update(Node v) {
- if (!v) return;
- v->subtree_val = get_subtree_val(v->l) + v->node_val
- + get_subtree_val(v->r);
- v->size = get_size(v->l) + 1 + get_size(v->r); }
- Node merge(Node l, Node r) {
- push_delta(l); push_delta(r);
- if (!l || !r) return l ? l : r;
- if (l->size <= r->size) {
- l->r = merge(l->r, r);
- update(l);
- return l; }
- else {
- r->l = merge(l, r->l);
- update(r);
- return r; } }
- void split(Node v, int key, Node &l, Node &r) {
- push_delta(v);
- l = r = NULL;
- if (!v) return;
- if (key <= get_size(v->l)) {
- split(v->l, key, l, v->l);
- r = v;
- } else {
- split(v->r, key - get_size(v->l) - 1, v->r, r);
- l = v; }
- update(v); }
- Node root;
public:
- cartree() : root(NULL) {}
- ~cartree() { delete root; }
- int get(Node v, int key) {
- push_delta(v);

```

```
--- if (key < get_size(v->l)) -----
---     return get(v->l, key);
--- else if (key > get_size(v->l)) -----
---     return get(v->r, key - get_size(v->l) - 1);
--- return v->node_val; }
- int get(int key) { return get(root, key); } -----
- void insert(Node item, int key) { -----
-     Node l, r;
-     split(root, key, l, r);
-     root = merge(merge(l, item), r); }
- void insert(int key, int val) { -----
-     insert(new _Node(val), key); }
- void erase(int key) { -----
-     Node l, m, r;
-     split(root, key + 1, m, r);
-     split(m, key, l, m);
-     delete m;
-     root = merge(l, r); }
- int query(int a, int b) { -----
-     Node l1, r1;
-     split(root, b+1, l1, r1);
-     Node l2, r2;
-     split(l1, a, l2, r2);
-     int res = get_subtree_val(r2);
-     l1 = merge(l2, r2);
-     root = merge(l1, r1);
-     return res; }
- void update(int a, int b, int delta) { -----
-     Node l1, r1;
-     split(root, b+1, l1, r1);
-     Node l2, r2;
-     split(l1, a, l2, r2);
-     apply_delta(r2, delta);
-     l1 = merge(l2, r2);
-     root = merge(l1, r1); }
- int size() { return get_size(root); } };
```

1.9.2. **Persistent Treap.**

1.10. Union Find.

```
struct union_find { -----
- vi p; union_find(int n) : p(n, -1) { }
- int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); }
- bool unite(int x, int y) { -----
-     int xp = find(x), yp = find(y);
-     if (xp == yp) return false;
-     if (p[xp] > p[yp]) std::swap(xp,yp);
-     p[xp] += p[yp], p[yp] = xp; return true; }
- int size(int x) { return -p[find(x)]; } };
```

1.11. Unique Counter.

```
struct UniqueCounter { -----
- int *B; std::map<int, int> last; LeqCounter *leq_cnt;
- UniqueCounter(int *ar, int n) { // 0-index A[i] -----
-     B = new int[n+1];
-     B[0] = 0;
-     for (int i = 1; i <= n; ++i) { -----
```

```
-----     B[i] = last[ar[i-1]]; -----
-----     last[ar[i-1]] = i; }
- leq_cnt = new LeqCounter(B, n+1); } -----
- int count(int l, int r) { -----
-     return leq_cnt->count(l+1, r+1, l); } };
```

2. DYNAMIC PROGRAMMING

2.1. Dynamic Convex Hull Trick.

```
// USAGE: hull.insert_line(m, b); hull.gety(x); -----
bool UPPER_HULL = true; // you can edit this -----
bool IS_QUERY = false, SPECIAL = false;
struct line { -----
- ll m, b; line(ll m=0, ll b=0): m(m), b(b) {}
- mutable std::multiset<line>::iterator it;
- const line *see(std::multiset<line>::iterator it)const;
- bool operator < (const line& k) const { -----
-     if (!IS_QUERY) return m < k.m;
-     if (!SPECIAL) { -----
-         ll x = k.m; const line *s = see(it);
-         if (!s) return 0;
-         return (b - s->b) < (x) * (s->m - m);
-     } else { -----
-         ll y = k.m; const line *s = see(it);
-         if (!s) return 0;
-         ll n1 = y - b, d1 = m;
-         ll n2 = b - s->b, d2 = s->m - m;
-         if (d1 < 0) n1 *= -1, d1 *= -1;
-         if (d2 < 0) n2 *= -1, d2 *= -1;
-         return (n1) * d2 > (n2) * d1; } } };
struct dynamic_hull : std::multiset<line> { -----
- bool bad(iterator y) { -----
-     iterator z = next(y);
-     if (y == begin()) { -----
-         if (z == end()) return 0;
-         return y->m == z->m && y->b <= z->b; }
-     iterator x = prev(y); -----
-     if (z == end()) return y->m == x->m && y->b <= x->b;
-     return (x->b - y->b)*(z->m - y->m)>=
-         (y->b - z->b)*(y->m - x->m); }
- iterator next(iterator y) {return ++y;}
- iterator prev(iterator y) {return --y;}
- void insert_line(ll m, ll b) { -----
-     IS_QUERY = false;
-     if (!UPPER_HULL) m *= -1;
-     iterator y = insert(line(m, b));
-     y->it = y; if (bad(y)) {erase(y); return;}
-     while (next(y) != end() && bad(next(y))) -----
-         erase(next(y));
-     while (y != begin() && bad(prev(y))) -----
-         erase(prev(y)); }
- ll gety(ll x) { -----
-     IS_QUERY = true; SPECIAL = false;
-     const line& L = *lower_bound(line(x, 0));
-     ll y = (L.m) * x + L.b;
-     return UPPER_HULL ? y : -y; }
- ll getx(ll y) { -----
```

```
-----     IS_QUERY = true; SPECIAL = true; -----
-----     const line& l = *lower_bound(line(y, 0)); -----
-----     return /*floor*/ ((y - l.b + l.m - 1) / l.m); } -----
} hull;
const line* line::see(std::multiset<line>::iterator it) -----
const {return ++it == hull.end() ? NULL : &it;} -----
```

2.2. Divide and Conquer Optimization. For DP problems of the form

$$dp(i,j) = \min_{k \leq j} \{ dp(i-1,k) + C(k,j) \}$$

where  $C(k,j)$  is some cost function.

```
ll dp[G+1][N+1]; -----
void solve_dp(int g, int k_L, int k_R, int n_L, int n_R) { ---
- int n_M = (n_L+n_R)/2; -----
- dp[g][n_M] = INF; -----
- int best_k = -1;
- for (int k = k_L; k <= n_M && k <= k_R; k++) -----
-     if (dp[g-1][k]+cost(k+1,n_M) < dp[g][n_M]) { -----
-         dp[g][n_M] = dp[g-1][k]+cost(k+1,n_M);
-         best_k = k; } -----
- if (n_L <= n_M-1) -----
-     solve_dp(g, k_L, best_k, n_L, n_M-1);
- if (n_M+1 <= n_R) -----
-     solve_dp(g, best_k, k_R, n_M+1, n_R); }
```

3. GEOMETRY

```
#include <complex> -----
#define x real() -----
#define y imag() -----
typedef std::complex<double> point; // 2D point only -----
const double PI = acos(-1.0), EPS = 1e-7; -----
```

3.1. Dots and Cross Products.

```
double dot(point a, point b) { -----
- return a.x * b.x + a.y * b.y; } // + a.z * b.z; -----
double cross(point a, point b) { -----
- return a.x * b.y - a.y * b.x; } -----
double cross(point a, point b, point c) { -----
- return cross(a, b) + cross(b, c) + cross(c, a); } -----
double cross3D(point a, point b) { -----
- return point(a.x*b.y - a.y*b.x, a.y*b.z -
-         a.z*b.y, a.z*b.x - a.x*b.z); }
```

3.2. Angles and Rotations.

```
double angle(point a, point b, point c) { -----
- // angle formed by abc in radians: PI < x <= PI -----
- return abs(remainder(arg(a-b) - arg(c-b), 2*PI)); } -----
point rotate(point p, point a, double d) { -----
- //rotate point a about pivot p CCW at d radians -----
- return p + (a - p) * point(cos(d), sin(d)); }
```

3.3. Spherical Coordinates.

$$\begin{aligned} x &= r \cos \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \cos \theta \sin \phi & \theta &= \cos^{-1} x/r \\ z &= r \sin \theta & \phi &= \text{atan2}(y,x) \end{aligned}$$

3.4. Point Projection.

```
point proj(point p, point v) {
- // project point p onto a vector v (2D & 3D)
- return dot(p, v) / norm(v) * v; }
point projLine(point p, point a, point b) {
- // project point p onto line ab (2D & 3D)
- return a + dot(p-a, b-a) / norm(b-a) * (b-a); }
point projSeg(point p, point a, point b) {
- // project point p onto segment ab (2D & 3D)
- double s = dot(p-a, b-a) / norm(b-a);
- return a + min(1.0, max(0.0, s)) * (b-a); }
point projPlane(point p, double a, double b,
- double c, double d) {
- // project p onto plane ax+by+cz+d=0 (3D)
- // same as: o + p - project(p - o, n);
- double k = -d / (a*a + b*b + c*c);
- point o(a*k, b*k, c*k), n(a, b, c);
- point v(p.x-o.x, p.y-o.y, p.z-o.z);
- double s = dot(v, n) / dot(n, n);
- return point(o.x + p.x + s * n.x, o.y +
- p.y + s * n.y, o.z + p.z + s * n.z); }
```

3.5. Great Circle Distance.

```
double greatCircleDist(double lat1, double long1,
- double lat2, double long2, double R) {
- long1 *= PI / 180; lat1 *= PI / 180; // to radians
- long2 *= PI / 180; lat2 *= PI / 180;
- return R*acos(sin(lat1)*sin(lat2) +
- cos(lat1)*cos(lat2)*cos(abs(long1 - long2))); }
// another version, using actual (x, y, z)
double greatCircleDist(point a, point b) {
- return atan2(abs(cross3D(a, b)), dot3D(a, b)); }
```

3.6. Point/Line/Plane Distances.

```
double distPtLine(point p, double a, double b, double c) {
- // dist from point p to line ax+by+c=0
- return abs(a*p.x + b*p.y + c) / sqrt(a*a + b*b); }
double distPtLine(point p, point a, point b) {
- // dist from point p to line ab
- return abs((a.y - b.y) * (p.x - a.x) +
- (b.x - a.x) * (p.y - a.y)) /
- hypot(a.x - b.x, a.y - b.y); }
double distPtPlane(point p, double a, double b,
- double c, double d) {
- // distance to 3D plane ax + by + cz + d = 0
- return (a*p.x+b*p.y+c*p.z+d)/sqrt(a*a+b*b+c*c); }
/*! // distance between 3D lines AB & CD (untested)
double distLine3D(point A,point B,point C,point D){
- point u = B - A, v = D - C, w = A - C;
- double a = dot(u, u), b = dot(u, v);
- double c = dot(v, v), d = dot(u, w);
- double e = dot(v, w), det = a*c - b*b;
- double s = det < EPS ? 0.0 : (b*e - c*d) / det;
- double t = det < EPS
- ? (b > c ? d/b : e/c) // parallel
- : (a*e - b*d) / det;
- point top = A + u * s, bot = w - A - v * t;
```

```
- return dist(top, bot);
} // dist<EPS: intersection */
```

3.7. Intersections.

3.7.1. Line-Segment Intersection. Get intersection points of 2D lines/segments *ab* and *cd*.

```
point null(HUGE_VAL, HUGE_VAL);
point line_inter(point a, point b, point c,
- point d, bool seg = false) {
- point ab(b.x - a.x, b.y - a.y);
- point cd(d.x - c.x, d.y - c.y);
- point ac(c.x - a.x, c.y - a.y);
- double D = -cross(ab, cd); // determinant
- double Ds = cross(cd, ac);
- double Dt = cross(ab, ac);
- if (abs(D) < EPS) { // parallel
- if (seg && abs(Ds) < EPS) { // collinear
- point p[] = {a, b, c, d};
- sort(p, p + 4, [](point a, point b) {
- return a.x < b.x-EPS ||
- (dist(a,b) < EPS && a.y < b.y-EPS); });
- return dist(p[1], p[2]) < EPS ? p[1] : null; }
- return null; }
- double s = Ds / D, t = Dt / D;
- if (seg && (min(s,t)<-EPS||max(s,t)>1+EPS)) return null;
- return point(a.x + s * ab.x, a.y + s * ab.y); }
/* double A = cross(d-a, b-a), B = cross(c-a, b-a);
return (B*d - A*c)/(B - A); */
```

3.7.2. Circle-Line Intersection. Get intersection points of circle at center *c*, radius *r*, and line *ab*.

```
std::vector<point> CL_inter(point c, double r,
- point a, point b) {
- point p = projLine(c, a, b);
- double d = abs(c - p); vector<point> ans;
- if (d > r + EPS); // none
- else if (d > r - EPS) ans.push_back(p); // tangent
- else if (d < EPS) { // diameter
- point v = r * (b - a) / abs(b - a);
- ans.push_back(c + v);
- ans.push_back(c - v);
- } else {
- double t = acos(d / r);
- p = c + (p - c) * r / d;
- ans.push_back(rotate(c, p, t));
- ans.push_back(rotate(c, p, -t));
- } return ans; }
```

3.7.3. Circle-Circle Intersection.

```
std::vector<point> CC_intersection(point c1,
- double r1, point c2, double r2) {
- double d = dist(c1, c2);
- vector<point> ans;
- if (d < EPS) {
- if (abs(r1-r2) < EPS); // inf intersections
- } else if (r1 < EPS) {
- if (abs(d - r2) < EPS) ans.push_back(c1);
```

```
- } else {
- double s = (r1*r1 + d*d - r2*r2) / (2*r1*d);
- double t = acos(max(-1.0, min(1.0, s)));
- point mid = c1 + (c2 - c1) * r1 / d;
- ans.push_back(rotate(c1, mid, t));
- if (abs(sin(t)) >= EPS)
- ans.push_back(rotate(c2, mid, -t));
- } return ans; }
```

3.8. Polygon Areas. Find the area of any 2D polygon given as points in *O(n)*.

```
double area(point p[], int n) {
- double a = 0;
- for (int i = 0, j = n - 1; i < n; j = i++)
- a += cross(p[i], p[j]);
- return abs(a) / 2; }
```

3.8.1. Triangle Area. Find the area of a triangle using only their lengths. Lengths must be valid.

```
double area(double a, double b, double c) {
- double s = (a + b + c) / 2;
- return sqrt(s*(s-a)*(s-b)*(s-c)); }
```

Cyclic Quadrilateral Area. Find the area of a cyclic quadrilateral using only their lengths. A quadrilateral is cyclic if its inner angles sum up to 360°.

```
double area(double a, double b, double c, double d) {
- double s = (a + b + c + d) / 2;
- return sqrt((s-a)*(s-b)*(s-c)*(s-d)); }
```

3.9. Polygon Centroid. Get the centroid/center of mass of a polygon in *O(m)*.

```
point centroid(point p[], int n) {
- point ans(0, 0);
- double z = 0;
- for (int i = 0, j = n - 1; i < n; j = i++) {
- double cp = cross(p[j], p[i]);
- ans += (p[j] + p[i]) * cp;
- z += cp;
- } return ans / (3 * z); }
```

3.10. Convex Hull.

3.10.1. 2D Convex Hull. Get the convex hull of a set of points using Graham-Andrew's scan. This sorts the points at *O(n log n)*, then performs the Monotonic Chain Algorithm at *O(n)*.

```
// counterclockwise hull in p[], returns size of hull
bool xcmp(const point& a, const point& b) {
- return a.x < b.x || (a.x == b.x && a.y < b.y); }
int convex_hull(point p[], int n) {
- std::sort(p, p + n, xcmp); if (n <= 1) return n;
- int k = 0; point *h = new point[2 * n];
- double zer = EPS; // -EPS to include collinears
- for (int i = 0; i < n; h[k++] = p[i++])
- while (k >= 2 && cross(h[k-2],h[k-1],p[i]) < zer)
- --k;
- for(int i = n-2, t = k; i >= 0; h[k++] = p[i--])
- while (k > t && cross(h[k-2],h[k-1],p[i]) < zer)
```



```
----- -k;
- k -= 1 + (h[0].x==h[1].x&&h[0].y==h[1].y ? 1 : 0);
- copy(h, h + k, p); delete[] h; return k; }

3.10.2. 3D Convex Hull. Currently  $O(N^2)$ , but can be optimized to a randomized  $O(N \log N)$  using the Clarkson-Shor algorithm. Sauce: Efficient 3D Convex Hull Tutorial on CF.

typedef std::vector<bool> vb;
struct point3D {
    ll x, y, z;
    point3D(ll x = 0, ll y = 0, ll z = 0) : x(x), y(y), z(z) {}
    point3D operator-(const point3D &o) const {
        return point3D(x - o.x, y - o.y, z - o.z); }
    point3D cross(const point3D &o) const {
        return point3D(y*o.z-z*o.y, z*o.x-x*o.z, x*o.y-y*o.x); }
    ll dot(const point3D &o) const {
        return x*o.x + y*o.y + z*o.z; }
    bool operator==(const point3D &o) const {
        return std::tie(x, y, z) == std::tie(o.x, o.y, o.z); }
    bool operator<(const point3D &o) const {
        return std::tie(x, y, z) < std::tie(o.x, o.y, o.z); } };
struct face {
    std::vector<int> p_idx;
    point3D q; };
std::vector<face> convex_hull_3D(std::vector<point3D> &points) {
    int n = points.size();
    std::vector<face> faces;
    std::vector<vb> dead(points.size(), vb(points.size(), true));
    auto add_face = [&](int a, int b, int c) {
        faces.push_back({a, b, c},
            (points[b] - points[a]).cross(points[c] - points[a]));
        dead[a][b] = dead[b][c] = dead[c][a] = false; };
    add_face(0, 1, 2);
    add_face(0, 2, 1);
    for (int i = 3; i < n; ++i) {
        std::vector<face> faces_inv;
        for(face &f : faces) {
            if ((points[i] - points[f.p_idx[0]]).dot(f.q) > 0)
                for (int j = 0; j < 3; ++j)
                    dead[f.p_idx[j]][f.p_idx[(j+1)%3]] = true;
            else
                faces_inv.push_back(f); }
        faces.clear();
        for(face &f : faces_inv) {
            for (int j = 0; j < 3; ++j) {
                int a = f.p_idx[j], b = f.p_idx[(j + 1) % 3];
                if(dead[b][a])
                    add_face(b, a, i); } }
        faces.insert(
            faces.end(), faces_inv.begin(), faces_inv.end()); }
    return faces; }
```

3.11. **Delaunay Triangulation**. Simply map each point  $(x,y)$  to  $(x,y,x^2+y^2)$ , find the 3d convex hull, and drop the 3rd dimension.

3.12. **Point in Polygon**. Check if a point is strictly inside (or on the border) of a polygon in  $O(n)$ .

```
bool inPolygon(point q, point p[], int n) {
    bool in = false;
    for (int i = 0, j = n - 1; i < n; j = i++)
        in ^= (((p[i].y > q.y) != (p[j].y > q.y)) &&
            q.x < (p[j].x - p[i].x) * (q.y - p[i].y) /
            (p[j].y - p[i].y) + p[i].x);
    return in; }

bool onPolygon(point q, point p[], int n) {
    for (int i = 0, j = n - 1; i < n; j = i++)
        if (abs(dist(p[i], q) + dist(p[j], q) -
            dist(p[i], p[j])) < EPS)
            return true;
    return false; }

3.13. Cut Polygon by a Line. Cut polygon by line  $\overline{ab}$  to its left in  $O(n)$ , such that  $\angle abp$  is counter-clockwise.
vector<point> cut(point p[],int n,point a,point b) {
    vector<point> poly;
    for (int i = 0, j = n - 1; i < n; j = i++) {
        double c1 = cross(a, b, p[j]);
        double c2 = cross(a, b, p[i]);
        if (c1 > -EPS) poly.push_back(p[j]);
        if (c1 * c2 < -EPS)
            poly.push_back(line_inter(p[j], p[i], a, b));
    return poly; }
```

3.14. **Triangle Centers**.

```
point bary(point A, point B, point C,
    double a, double b, double c) {
    return (A*a + B*b + C*c) / (a + b + c); }
point trilinear(point A, point B, point C,
    double a, double b, double c) {
    return bary(A,B,C,abs(B-C)*a,
        abs(C-A)*b,abs(A-B)*c); }
point centroid(point A, point B, point C) {
    return bary(A, B, C, 1, 1, 1); }
point circumcenter(point A, point B, point C) {
    double a=norm(B-C), b=norm(C-A), c=norm(A-B);
    return bary(A,B,C,a*(b+c-a),b*(c+a-b),c*(a+b-c)); }
point orthocenter(point A, point B, point C) {
    return bary(A,B,C, tan(angle(B,A,C)),
        tan(angle(A,B,C)), tan(angle(A,C,B))); }
point incenter(point A, point B, point C) {
    return bary(A,B,C,abs(B-C),abs(A-C),abs(A-B)); }
// incircle radius given the side lengths a, b, c
double inradius(double a, double b, double c) {
    double s = (a + b + c) / 2;
    return sqrt(s * (s-a) * (s-b) * (s-c)) / s; }
point excenter(point A, point B, point C) {
    double a = abs(B-C), b = abs(C-A), c = abs(A-B);
    return bary(A, B, C, -a, b, c); }
// return bary(A, B, C, a, -b, c);
// return bary(A, B, C, a, b, -c);
point brocard(point A, point B, point C) {
    double a = abs(B-C), b = abs(C-A), c = abs(A-B);
    return bary(A,B,C,c/b*a,c/b*a,c/b*a); // CCW
    // return bary(A,B,C,b/c*a,b/c*a,b/c*a); // CW }
```

```
point symmedian(point A, point B, point C) {
    return bary(A,B,C,norm(B-C),norm(C-A),norm(A-B)); }
```

3.15. **Convex Polygon Intersection**. Get the intersection of two convex polygons in  $O(n^2)$ .

```
std::vector<point> convex_polygon_inter(
    point a[], int an, point b[], int bn) {
    point ans[an + bn + an*bn];
    int size = 0;
    for (int i = 0; i < an; ++i)
        if (inPolygon(a[i],b,bn) || onPolygon(a[i],b,bn))
            ans[size++] = a[i];
    for (int i = 0; i < bn; ++i)
        if (inPolygon(b[i],a,an) || onPolygon(b[i],a,an))
            ans[size++] = b[i];
    for (int i = 0, I = an - 1; i < an; I = i++)
        for (int j = 0, J = bn - 1; j < bn; J = j++) {
            try {
                point p=line_inter(a[i],a[I],b[j],b[J],true);
                ans[size++] = p;
            } catch (exception ex) {} }
    size = convex_hull(ans, size);
    return vector<point>(ans, ans + size); }
```

3.16. **Pick's Theorem for Lattice Points**. Count points with integer coordinates inside and on the boundary of a polygon in  $O(n)$  using Pick's theorem:  $\text{Area} = I + B/2 - 1$ .

```
int interior(point p[], int n) {
    return area(p,n) - boundary(p,n) / 2 + 1; }
int boundary(point p[], int n) {
    int ans = 0;
    for (int i = 0, j = n - 1; i < n; j = i++)
        ans += gcd(p[i].x - p[j].x, p[i].y - p[j].y);
    return ans; }
```

3.17. **Minimum Enclosing Circle**. Get the minimum bounding ball that encloses a set of points (2D or 3D) in  $\Theta n$ .

```
std::pair<point, double> bounding_ball(point p[], int n){
    std::random_shuffle(p, p + n);
    point center(0, 0); double radius = 0;
    for (int i = 0; i < n; ++i) {
        if (dist(center, p[i]) > radius + EPS) {
            center = p[i]; radius = 0;
            for (int j = 0; j < i; ++j)
                if (dist(center, p[j]) > radius + EPS) {
                    center.x = (p[i].x + p[j].x) / 2;
                    center.y = (p[i].y + p[j].y) / 2;
                    // center.z = (p[i].z + p[j].z) / 2;
                    radius = dist(center, p[i]); // midpoint
                    for (int k = 0; k < j; ++k)
                        if (dist(center, p[k]) > radius + EPS) {
                            center = circumcenter(p[i], p[j], p[k]);
                            radius = dist(center, p[i]); } } } }
    return {center, radius}; }
```

3.18. **Shamos Algorithm.** Solve for the polygon diameter in  $O(n \log n)$ .

```
double shamos(point p[], int n) {
- point *h = new point[n+1]; copy(p, p + n, h);
- int k = convex_hull(h, n); if (k <= 2) return 0;
- h[k] = h[0]; double d = HUGE_VAL;
- for (int i = 0, j = 1; i < k; ++i) {
-   while (distPtLine(h[j+1], h[i], h[i+1]) >=
-     distPtLine(h[j], h[i], h[i+1])) {
-     j = (j + 1) % k; }
-   d = min(d, distPtLine(h[j], h[i], h[i+1]));
- } return d; }
```

3.19. **kD Tree.** Get the  $k$ -nearest neighbors of a point within pruned radius in  $O(k \log k \log n)$ .

```
#define cpoint const point&
bool cmpx(cpoint a, cpoint b) {return a.x < b.x;}
bool cmpy(cpoint a, cpoint b) {return a.y < b.y;}
struct KDTree {
- KDTree(point p[],int n): p(p), n(n) {build(0,n);}
- priority_queue< pair<double, point*> > pq;
- point *p; int n, k; double qx, qy, prune;
- void build(int L, int R, bool dvx=false) {
-   if (L >= R) return;
-   int M = (L + R) / 2;
-   nth_element(p + L, p + M, p + R, dvx?cmpx:cmPy);
-   build(L, M, !dvx); build(M + 1, R, !dvx); }
- void dfs(int L, int R, bool dvx) {
-   if (L >= R) return;
-   int M = (L + R) / 2;
-   double dx = qx - p[M].x, dy = qy - p[M].y;
-   double delta = dvx ? dx : dy;
-   double D = dx * dx + dy * dy;
-   if (D<=prune && (pq.size()<k || D<pq.top().first)) {
-     pq.push(make_pair(D, &p[M]));
-     if (pq.size() > k) pq.pop(); }
-   int nL = L, nR = M, fL = M + 1, fR = R;
-   if (delta > 0) {swap(nL, fL); swap(nR, fR);}
-   dfs(nL, nR, !dvx);
-   D = delta * delta;
-   if (D<=prune && (pq.size()<k || D<pq.top().first))
-   dfs(fL, fR, !dvx); }
- // returns k nearest neighbors of (x, y) in tree
- // usage: vector<point> ans = tree.knn(x, y, 2);
- vector<point> knn(double x, double y,
-   int k=1, double r=-1) {
-   qx=x; qy=y; this->k=k; prune=r<0?HUGE_VAL:r*r;
-   dfs(0, n, false); vector<point> v;
-   while (!pq.empty()) {
-     v.push_back(*pq.top().second);
-     pq.pop();
-   } reverse(v.begin(), v.end());
-   return v; } };
```

3.20. **Line Sweep (Closest Pair).** Get the closest pair distance of a set of points in  $O(n \log n)$  by sweeping a line and keeping a bounded rectangle. Modifiable for other metrics such as Minkowski and Manhattan distance. For external point queries, see  $kD$  Tree.

```
bool cmpy(const point& a, const point& b) { return a.y < b.y; }
double closest_pair_sweep(point p[], int n) {
- if (n <= 1) return HUGE_VAL;
- std::sort(p, p + n, cmpy);
- std::set<point> box; box.insert(p[0]);
- double best = 1e13; // infinity, but not HUGE_VAL
- for (int L = 0, i = 1; i < n; ++i) {
-   while(L < i && p[i].y - p[L].y > best)
-     box.erase(p[L++]);
-   point bound(p[i].x - best, p[i].y - best);
-   std::set<point>::iterator it = box.lower_bound(bound);
-   while (it != box.end() && p[i].x+best >= it->x){
-     double dx = p[i].x - it->x;
-     double dy = p[i].y - it->y;
-     best = std::min(best, std::sqrt(dx*dx + dy*dy));
-     ++it; }
-   box.insert(p[i]);
- } return best; }
```

3.21. **Line upper/lower envelope.** To find the upper/lower envelope of a collection of lines  $a_i + b_ix$ , plot the points  $(b_i, a_i)$ , add the point  $(0, \pm\infty)$  (depending on if upper/lower envelope is desired), and then find the convex hull.

3.22. **Formulas.** Let  $a = (a_x, a_y)$  and  $b = (b_x, b_y)$  be two-dimensional vectors.

- $a \cdot b = |a||b| \cos \theta$ , where  $\theta$  is the angle between  $a$  and  $b$ .
- $a \times b = |a||b| \sin \theta$ , where  $\theta$  is the signed angle between  $a$  and  $b$ .
- $a \times b$  is equal to the area of the parallelogram with two of its sides formed by  $a$  and  $b$ . Half of that is the area of the triangle formed by  $a$  and  $b$ .
- The line going through  $a$  and  $b$  is  $Ax + By = C$  where  $A = b_y - a_y$ ,  $B = a_x - b_x$ ,  $C = Aa_x + Ba_y$ .
- Two lines  $A_1x + B_1y = C_1$ ,  $A_2x + B_2y = C_2$  are parallel iff.  $D = A_1B_2 - A_2B_1$  is zero. Otherwise their unique intersection is  $(B_2C_1 - B_1C_2, A_1C_2 - A_2C_1)/D$ .
- **Euler's formula:**  $V - E + F = 2$
- Side lengths  $a, b, c$  can form a triangle iff.  $a + b > c$ ,  $b + c > a$  and  $a + c > b$ .
- Sum of internal angles of a regular convex  $n$ -gon is  $(n - 2)\pi$ .
- **Law of sines:**  $\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C}$
- **Law of cosines:**  $b^2 = a^2 + c^2 - 2ac \cos B$
- Internal tangents of circles  $(c_1, r_1), (c_2, r_2)$  intersect at  $(c_1r_2 + c_2r_1)/(r_1 + r_2)$ , external intersect at  $(c_1r_2 - c_2r_1)/(r_1 + r_2)$ .

4. GRAPHS

4.1. Single-Source Shortest Paths.

4.1.1. *Dijkstra.*

```
#include "graph_template_adjlist.cpp"
// insert inside graph; needs n, dist[], and adj[]
void dijkstra(int s) {
- for (int u = 0; u < n; ++u)
-   dist[u] = INF;
- dist[s] = 0;
- std::priority_queue<ii, vii, std::greater<ii> > pq;
- pq.push({0, s});
```

```
- while (!pq.empty()) {
-   int u = pq.top().second;
-   int d = pq.top().first;
-   pq.pop();
-   if (dist[u] < d) continue;
-   dist[u] = d;
-   for (auto &e : adj[u]) {
-     int v = e.first;
-     int w = e.second;
-     if (dist[v] > dist[u] + w) {
-       dist[v] = dist[u] + w;
-       pq.push({dist[v], v}); } } }
```

4.1.2. *Bellman-Ford.*

```
#include "graph_template_adjlist.cpp"
// insert inside graph; needs n, dist[], and adj[]
void bellman_ford(int s) {
- for (int u = 0; u < n; ++u)
-   dist[u] = INF;
- dist[s] = 0;
- for (int i = 0; i < n-1; ++i)
-   for (int u = 0; u < n; ++u)
-     for (auto &e : adj[u])
-       if (dist[u] + e.second < dist[e.first])
-         dist[e.first] = dist[u] + e.second; }
// you can call this after running bellman_ford()
bool has_neg_cycle() {
- for (int u = 0; u < n; ++u)
-   for (auto &e : adj[u])
-     if (dist[e.first] > dist[u] + e.second)
-       return true;
- return false; }
```

4.1.3. *Shortest Path Faster Algorithm.*

```
#include "graph_template_adjlist.cpp"
// insert inside graph;
// needs n, dist[], in_queue[], num_vis[], and adj[]
bool spfa(int s) {
- for (int u = 0; u < n; ++u) {
-   dist[u] = INF;
-   in_queue[u] = 0;
-   num_vis[u] = 0; }
- dist[s] = 0;
- in_queue[s] = 1;
- bool has_negative_cycle = false;
- std::queue<int> q; q.push(s);
- while (not q.empty()) {
-   int u = q.front(); q.pop(); in_queue[u] = 0;
-   if (++num_vis[u] >= n)
-     dist[u] = -INF, has_negative_cycle = true;
-   for (auto &[v, c] : adj[u])
-     if (dist[v] > dist[u] + c) {
-       dist[v] = dist[u] + c;
-       if (!in_queue[v]) {
-         q.push(v);
```

```
----- in_queue[v] = 1; } } } -----
- return has_negative_cycle; } -----
```

4.2. All-Pairs Shortest Paths.

4.2.1. Floyd-Washall.

```
#include "graph_template_adjmat.cpp" -----
// insert inside graph; needs n and mat[][] -----
void floyd_warshall() { -----
- for (int k = 0; k < n; ++k) -----
- - for (int i = 0; i < n; ++i) -----
- - - for (int j = 0; j < n; ++j) -----
- - - - if (mat[i][k] + mat[k][j] < mat[i][j]) -----
- - - - - mat[i][j] = mat[i][k] + mat[k][j]; } -----
```

4.3. Strongly Connected Components.

4.3.1. Kosaraju.

```
struct kosaraju_graph { -----
- int n, *vis; -----
- vi **adj; -----
- std::vector<vi> sccs; -----
- kosaraju_graph(int n) { -----
- - this->n = n; -----
- - vis = new int[n]; -----
- - adj = new vi*[2]; -----
- - for (int dir = 0; dir < 2; ++dir) -----
- - - adj[dir] = new vi[n]; } -----
- void add_edge(int u, int v) { -----
- - adj[0][u].push_back(v); -----
- - adj[1][v].push_back(u); } -----
- void dfs(int u, int p, int dir, vi &topo) { -----
- - vis[u] = 1; -----
- - for (int v : adj[dir][u]) -----
- - - if (!vis[v] && v != p) dfs(v, u, dir, topo); -----
- - topo.push_back(u); } -----
- void kosaraju() { -----
- - vi topo; -----
- - for (int u = 0; u < n; ++u) vis[u] = 0; -----
- - for (int u = 0; u < n; ++u) if(!vis[u]) dfs(u, -1, 0, topo); -----
- - for (int u = 0; u < n; ++u) vis[u] = 0; -----
- - for (int i = n-1; i >= 0; --i) { -----
- - - if (!vis[topo[i]]) { -----
- - - - sccs.push_back({}); -----
- - - - dfs(topo[i], -1, 1, sccs.back()); } } } } -----
```

4.3.2. Tarjan's Offline Algorithm .

```
int n, id[N], low[N], st[N], in[N], TOP, ID; -----
int scc[N], SCC_SIZE; // 0 <= scc[u] < SCC_SIZE -----
vector<int> adj[N]; // 0-based adjlist -----
void dfs(int u) { -----
- id[u] = low[u] = ID++; -----
- st[TOP++] = u; in[u] = 1; -----
- for (int v : adj[u]) { -----
- - if (id[v] == -1) { -----
- - - dfs(v); -----
- - - low[u] = min(low[u], low[v]); -----
- - } else if (in[v] == 1) -----
```

```
----- low[u] = min(low[u], id[v]); } -----
- if (id[u] == low[u]) { -----
- - int sid = SCC_SIZE++; -----
- - do { -----
- - - int v = st[--TOP]; -----
- - - in[v] = 0; scc[v] = sid; -----
- - } while (st[TOP] != u); } } -----
void tarjan() { // call tarjan() to load SCC -----
- memset(id, -1, sizeof(int) * n); -----
- SCC_SIZE = ID = TOP = 0; -----
- for (int i = 0; i < n; ++i) -----
- - if (id[i] == -1) dfs(i); } -----
```

4.4. Minimum Mean Weight Cycle . Run this for each strongly connected component

```
typedef std::vector<double> vd; -----
double min_mean_cycle(graph &g) { -----
- double mn = INF; -----
- std::vector<vd> dp(g.n+1, vd(g.n, mn)); -----
- dp[0][0] = 0; -----
- for (int k = 1; k <= g.n; ++k) -----
- - for (int u = 0; u < g.n; ++u) -----
- - - for (auto &[v, w]: g.adj[u]) -----
- - - - dp[k][v] = std::min(ar[k][v], dp[k-1][u] + w); -----
- - for (int k = 0; k < g.n; ++k) { -----
- - - double mx = -INF; -----
- - - for (int u = 0; u < g.n; ++u) -----
- - - - mx = std::max(mx, (dp[g.n][u] - dp[k][u]) / (g.n - k)); -----
- - - mn = std::min(mn, mx); } -----
- return mn; } -----
```

4.5. Biconnected Components.

4.5.1. Bridges and Articulation Points.

```
struct graph { -----
- int n, *disc, *low, TIME; -----
- vi *adj, stk, articulation_points; -----
- std::set<ii> bridges; -----
- vvi comps; -----
- graph(int n) : n(n) { -----
- - adj = new vi[n]; -----
- - disc = new int[n]; -----
- - low = new int[n]; } -----
- void add_edge(int u, int v) { -----
- - adj[u].push_back(v); -----
- - adj[v].push_back(u); } -----
- void _bridges_artics(int u, int p) { -----
- - disc[u] = low[u] = TIME++; -----
- - stk.push_back(u); -----
- - int children = 0; -----
- - bool has_low_child = false; -----
- - for (int v : adj[u]) { -----
- - - if (disc[v] == -1) { -----
- - - - _bridges_artics(v, u); -----
- - - - children++; -----
- - - - if (disc[u] < low[v]) -----
- - - - - bridges.insert({std::min(u, v), std::max(u, v)}); -----
```

```
----- if (disc[u] <= low[v]) { -----
----- - has_low_child = true; -----
----- - comps.push_back({u}); -----
----- - while (comps.back().back() != v and !stk.empty()) { -----
----- - - comps.back().push_back(stk.back()); -----
----- - - stk.pop_back(); } } -----
----- low[u] = std::min(low[u], low[v]); -----
----- } else if (v != p) -----
----- - low[u] = std::min(low[u], disc[v]); } -----
----- if ((p == -1 && children >= 2) || -----
----- - (p != -1 && has_low_child)) -----
----- - articulation_points.push_back(u); } -----
- void bridges_artics() { -----
- - for (int u = 0; u < n; ++u) disc[u] = -1; -----
- - stk.clear(); -----
- - articulation_points.clear(); -----
- - bridges.clear(); -----
- - comps.clear(); -----
- - TIME = 0; -----
- - for (int u = 0; u < n; ++u) if (disc[u] == -1) -----
- - - _bridges_artics(u, -1); } } -----
```

4.5.2. Block Cut Tree.

```
// insert inside code for finding articulation points -----
graph build_block_cut_tree() { -----
- int bct_n = articulation_points.size() + comps.size(); -----
- vi block_id(n, is_art(n, 0)); -----
- graph tree(bct_n); -----
- for (int i = 0; i < articulation_points.size(); ++i) { -----
- - block_id[articulation_points[i]] = i; -----
- - is_art[articulation_points[i]] = 1; } -----
- for (int i = 0; i < comps.size(); ++i) { -----
- - int id = i + articulation_points.size(); -----
- - for (int u : comps[i]) -----
- - - if (is_art[u]) tree.add_edge(block_id[u], id); -----
- - - else block_id[u] = id; } -----
- return tree; } -----
```

4.5.3. Bridge Tree.

```
// insert inside code for finding bridges -----
// requires union_find and hasher -----
graph build_bridge_tree() { -----
- union_find uf(n); -----
- for (int u = 0; u < n; ++u) { -----
- - for (int v : adj[u]) { -----
- - - ii uv = { std::min(u, v), std::max(u, v) }; -----
- - - if (bridges.find(uv) == bridges.end()) -----
- - - - uf.unite(u, v); } } -----
- hasher h; -----
- for (int u = 0; u < n; ++u) -----
- - if (u == uf.find(u)) h.get_hash(u); -----
- - int tn = h.h.size(); -----
- - graph tree(tn); -----
- - for (int i = 0; i < M; ++i) { -----
- - - int ui = h.get_hash(uf.find(u)); -----
- - - int vi = h.get_hash(uf.find(v)); -----
```



```
--- if (ui != vi) tree.add_edge(ui, vi); } -----
- return tree; } -----
```

4.6. Minimum Spanning Tree.

4.6.1. Kruskal.

```
#include "graph_template_edgelist.cpp" -----
#include "union_find.cpp" -----
// insert inside graph; needs n, and edges -----
void kruskal(viii &res) { -----
- viii().swap(res); // or use res.clear(); -----
- std::priority_queue<iii, viii, std::greater<iii> > pq; -----
- for (auto &edge : edges) -----
- pq.push(edge); -----
- union_find uf(n); -----
- while (!pq.empty()) { -----
- auto node = pq.top(); pq.pop(); -----
- int u = node.second.first; -----
- int v = node.second.second; -----
- if (uf.unite(u, v)) -----
- res.push_back(node); } } -----
```

4.6.2. Prim.

```
#include "graph_template_adjlist.cpp" -----
// insert inside graph; needs n, vis[], and adj[] -----
void prim(viii &res, int s=0) { -----
- viii().swap(res); // or use res.clear(); -----
- std::priority_queue<ii, vii, std::greater<ii> > pq; -----
- pq.push({0, s}); -----
- vis[s] = true; -----
- while (!pq.empty()) { -----
- int u = pq.top().second; pq.pop(); -----
- vis[u] = true; -----
- for (auto &[v, w] : adj[u]) { -----
- if (v == u) continue; -----
- if (vis[v]) continue; -----
- res.push_back({w, {u, v}}); -----
- pq.push({w, v}); } } } -----
```

4.7. Euler Path/Cycle .

4.7.1. Euler Path/Cycle in a Directed Graph .

```
#define MAXV 1000 -----
#define MAXE 5000 -----
int indeg[MAXV], outdeg[MAXV], res[MAXE + 1];
ii start_end(graph &g) { -----
- int start = -1, end = -1, any = 0, c = 0; -----
- for (int u = 0; u < n; ++u) { -----
- if (outdeg[u] > 0) any = u; -----
- if (indeg[u] + 1 == outdeg[u]) start = u, c++; -----
- else if (indeg[u] == outdeg[u] + 1) end = u, c++; -----
- else if (indeg[u] != outdeg[u]) return {-1, -1}; } -----
- if ((start == -1) != (end == -1) || (c != 2 && c != 0)) -----
- return {-1, -1}; -----
- if (start == -1) start = end = any; -----
- return {start, end}; } -----
bool euler_path(graph &g) { -----
```

```
- ii se = start_end(g); -----
- int cur = se.first, at = g.edges.size() + 1; -----
- if (cur == -1) return false; -----
- std::stack<int> s; -----
- while (true) { -----
- if (outdeg[cur] == 0) { -----
- res[at] = cur; -----
- if (s.empty()) break; -----
- cur = s.top(); s.pop(); -----
- } else s.push(cur), cur = g.adj[cur][--outdeg[cur]]; } -----
- return at == 0; } -----
```

4.7.2. Euler Path/Cycle in an Undirected Graph .

```
std::multiset<int> adj[1010]; -----
std::list<int> L; -----
std::list<int>::iterator euler( -----
- int at, int to, std::list<int>::iterator it -----
) { -----
- if (at == to) return it; -----
- L.insert(it, at), --it; -----
- while (!adj[at].empty()) { -----
- int nxt = *adj[at].begin(); -----
- adj[at].erase(adj[at].find(nxt)); -----
- adj[nxt].erase(adj[nxt].find(at)); -----
- if (to == -1) { -----
- it = euler(nxt, at, it); -----
- L.insert(it, at); -----
- --it; -----
- } else { -----
- it = euler(nxt, to, it); -----
- to = -1; } } -----
- return it; } -----
// euler(0, -1, L.begin()) -----
```

4.8. Bipartite Matching .

4.8.1. Alternating Paths Algorithm .

```
vi* adj; -----
bool* done; // initially all false -----
int* owner; // initially all -1 -----
int alternating_path(int left) { -----
- if (done[left]) return 0; -----
- done[left] = true; -----
- for (int right : adj[left]) { -----
- if (owner[right] == -1 || alternating_path(owner[right])) { -----
- owner[right] = left; return 1; } } -----
- return 0; } -----
```

4.8.2. Hopcroft-Karp Algorithm .

```
#define MAXN 5000 -----
int dist[MAXN+1], q[MAXN+1]; -----
#define dist(v) dist[v == -1 ? MAXN : v] -----
struct bipartite_graph { -----
- int n, m, *L, *R; vi *adj; -----
- bipartite_graph(int n, int m) : n(n), m(m), -----
- L(new int[n]), R(new int[m]), adj(new vi[n]) {} -----
- ~bipartite_graph() { delete[] adj; delete[] L; delete[] R; } -----
```

```
- void add_edge(int u, int v) { adj[u].push_back(v); } -----
- bool bfs() { -----
- int l = 0, r = 0; -----
- for (int v = 0; v < n; ++v) -----
- if(L[v] == -1) dist(v) = 0, q[r++] = v; -----
- else dist(v) = INF; -----
- dist(-1) = INF; -----
- while(l < r) { -----
- int v = q[l++]; -----
- if(dist(v) < dist(-1)) -----
- for (int u : adj[v]) -----
- if(dist(R[u]) == INF) { -----
- dist(R[u]) = dist(v) + 1; -----
- q[r++] = R[u]; } } -----
- return dist(-1) != INF; } -----
- bool dfs(int v) { -----
- if(v != -1) { -----
- for (int u : adj[v]) -----
- if(dist(R[u]) == dist(v) + 1) -----
- if(dfs(R[u])) { R[u] = v; L[v] = u; return true; } -----
- dist(v) = INF; -----
- return false; } -----
- return true; } -----
- int maximum_matching() { -----
- int matching = 0; -----
- for (int u = 0; u < n; ++u) -----
- L[u] = R[u] = -1; -----
- while(bfs()) -----
- for (int u = 0; u < n; ++u) -----
- matching += L[u] == -1 && dfs(u); -----
- return matching; } };
```

4.8.3. Minimum Vertex Cover in Bipartite Graphs .

```
#include "hopcroft_karp.cpp" -----
std::vector<bool> alt; -----
void dfs(bipartite_graph &g, int u) { -----
- alt[u] = true; -----
- for (int v : g.adj[u]) { -----
- alt[v + g.n] = true; -----
- if (g.R[v] != -1 && !alt[g.R[v]]) -----
- dfs(g, g.R[v]); } } -----
vi mvc_bipartite(bipartite_graph &g) { -----
- vi res; g.maximum_matching(); -----
- alt.assign(g.n + g.m, false); -----
- for (int i = 0; i < g.n; ++i) if (g.L[i] == -1) dfs(g, i); -----
- for (int i = 0; i < g.n; ++i) if (!alt[i]) res.push_back(i); -----
- for (int i = 0; i < g.m; ++i) -----
- if (alt[g.n + i]) res.push_back(g.n + i); -----
- return res; } -----
```

4.9. Maximum Flow.

4.9.1. Edmonds-Karp .  $O(VE^2)$

4.9.2. Dinic.  $O(V^2E)$ 

```

struct edge {
- int u, v;
- ll c, f;
- edge(int u, int v, ll c) : u(u), v(v), c(c), f(0) {} };
struct flow_network {
- int n, s, t, *adj_ptr, *par, *dist;
- std::vector<edge> edges;
- std::vector<int> *adj;
- flow_network(int n, int s, int t) : n(n), s(s), t(t) {
-- adj = new std::vector<int>[n];
-- adj_ptr = new int[n];
-- par = new int[n];
-- dist = new int[n]; }
- void add_edge(int u, int v, ll c, bool bi=false) {
-- adj[u].push_back(edges.size());
-- edges.push_back(edge(u, v, c));
-- adj[v].push_back(edges.size());
-- edges.push_back(edge(v, u, (bi ? c : 0LL))); }
- ll res(edge &e) { return e.c - e.f; }
- bool make_level_graph() {
-- for (int u = 0; u < n; ++u) dist[u] = -1;
-- dist[s] = 0;
-- std::queue<int> q; q.push(s);
-- while (!q.empty()) {
-- int u = q.front(); q.pop();
-- for (int i : adj[u]) {
-- edge &e = edges[i];
-- if (dist[e.v] < 0 and res(e)) {
-- dist[e.v] = dist[u] + 1;
-- q.push(e.v); } } }
-- return dist[t] != -1; }
- bool is_next(int u, int v) {
-- return dist[v] == dist[u] + 1; }
- bool dfs(int u) {
-- if (u == t) return true;
-- for (int &ii = adj_ptr[u]; ii < adj[u].size(); ++ii) {
-- int i = adj[u][ii];
-- edge &e = edges[i];
-- if (is_next(u, e.v) and res(e) > 0 and dfs(e.v)) {
-- par[e.v] = i;
-- return true; } }
-- return false; }
- bool aug_path() {
-- for (int u = 0; u < n; ++u) par[u] = -1;
-- return dfs(s); }
- ll calc_max_flow() {
-- ll total_flow = 0;
-- while (make_level_graph()) {
-- for (int u = 0; u < n; ++u) adj_ptr[u] = 0;
-- while (aug_path()) {
-- ll flow = INF;
-- for (int i = par[t]; i != -1; i = par[edges[i].u])
-- flow = std::min(flow, res(edges[i]));
-- for (int i = par[t]; i != -1; i = par[edges[i].u]) {
-- edges[i].f += flow;

```

```

----- edges[i^1].f -= flow; } -----
----- total_flow += flow; } } -----
-- return total_flow; } };
```

4.9.3. Push-relabel.  $\omega(VE + V^2\sqrt{E})$ ,  $O(V^3)$ 

```

int n;
std::vector<vi> capacity, flow;
vi height, excess;
void push(int u, int v) {
- int d = min(excess[u], capacity[u][v] - flow[u][v]);
- flow[u][v] += d; flow[v][u] -= d;
- excess[u] -= d; excess[v] += d; }
void relabel(int u) {
- int d = INF;
- for (int i = 0; i < n; i++)
-- if (capacity[u][i] - flow[u][i] > 0)
-- d = min(d, height[i]);
- if (d < INF) height[u] = d + 1; }
vi find_max_height_vertices(int s, int t) {
- vi max_height;
- for (int i = 0; i < n; i++) {
-- if (i != s && i != t && excess[i] > 0) {
-- if (!max_height.empty() && height[i] > height[max_height[0]])
-- max_height.clear();
-- if (max_height.empty() || height[i] == height[max_height[0]])
-- max_height.push_back(i); } }
- return max_height; }
int max_flow(int s, int t) {
- flow.assign(n, vi(n, 0));
- height.assign(n, 0); height[s] = n;
- excess.assign(n, 0); excess[s] = INF;
- for (int i = 0; i < n; i++) if (i != s) push(s, i);
- vi current;
- while (!current.empty()) {
-- for (int i : current) {
-- bool pushed = false;
-- for (int j = 0; j < n && excess[i]; j++) {
-- if (capacity[i][j] - flow[i][j] > 0 &&
-- height[i] == height[j] + 1) {
-- push(i, j);
-- pushed = true; } }
-- if (!pushed) relabel(i), break; } }
- int max_flow = 0;
- for (int i = 0; i < n; i++) max_flow += flow[i][t];
- return max_flow; }
```

## 4.9.4. Gomory-Hu (All-pairs Maximum Flow).

```

#define MAXV 2000
int q[MAXV], d[MAXV];
struct flow_network {
- struct edge { int v, nxt, cap;
-- edge(int _v, int _cap, int _nxt)
-- : v(_v), nxt(_nxt), cap(_cap) {} } };
- int n, *head, *curh; vector<edge> e, e_store;
- flow_network(int _n) : n(_n) {
-- curh = new int[n];
-- memset(head = new int[n], -1, n*sizeof(int)); }

```

```

- void reset() { e = e_store; }
- void add_edge(int u, int v, int uv, int vu=0) {
-- e.push_back(edge(v,uv,head[u])); head[u]=(int)size(e)-1;
-- e.push_back(edge(u,vu,head[v])); head[v]=(int)size(e)-1; }
- int augment(int v, int t, int f) {
-- if (v == t) return f;
-- for (int &i = curh[v], ret; i != -1; i = e[i].nxt)
-- if (e[i].cap > 0 && d[e[i].v] + 1 == d[v])
-- if ((ret = augment(e[i].v, t, min(f, e[i].cap))) > 0)
-- return (e[i].cap -= ret, e[i^1].cap += ret, ret);
-- return 0; }
- int max_flow(int s, int t, bool res=true) {
-- e_store = e;
-- int l, r, f = 0, x;
-- while (true) {
-- memset(d, -1, n*sizeof(int));
-- l = r = 0, d[q[r++]] = t = 0;
-- while (l < r)
-- for (int v = q[l++], i = head[v]; i != -1; i = e[i].nxt)
-- if (e[i^1].cap > 0 && d[e[i].v] == -1)
-- d[q[r++]] = e[i].v = d[v]+1;
-- if (d[s] == -1) break;
-- memcpy(curh, head, n * sizeof(int));
-- while ((x = augment(s, t, INF)) != 0) f += x; }
-- if (res) reset();
-- return f; } };
bool same[MAXV];
pair<vii, vvi> construct_gh_tree(flow_network &g) {
- int n = g.n, v;
- vii par(n, ii(0, 0)); vvi cap(n, vi(n, -1));
- rep(s,1,n) {
-- int l = 0, r = 0;
-- par[s].second = g.max_flow(s, par[s].first, false);
-- memset(d, 0, n * sizeof(int));
-- memset(same, 0, n * sizeof(bool));
-- d[q[r++]] = s = 1;
-- while (l < r) {
-- same[v = q[l++]] = true;
-- for (int i = g.head[v]; i != -1; i = g.e[i].nxt)
-- if (g.e[i].cap > 0 && d[g.e[i].v] == 0)
-- d[q[r++]] = g.e[i].v = 1; }
-- rep(i,s+1,n)
-- if (par[i].first == par[s].first && same[i])
-- par[i].first = s;
-- g.reset(); }
- rep(i,0,n) {
-- int mn = INF, cur = i;
-- while (true) {
-- cap[cur][i] = mn;
-- if (cur == 0) break;
-- mn = min(mn, par[cur].second, cur = par[cur].first; } }
- return make_pair(par, cap); }
int compute_max_flow(int s, int t, const pair<vii, vvi> &gh) {
- int cur = INF, at = s;
- while (gh.second[at][t] == -1)
-- cur = min(cur, gh.first[at].second),

```

```

-- at = gh.first[at].first;
-- return min(cur, gh.second[at][t]); }

```

#### 4.10. Minimum Cost Maximum Flow.

```

struct edge {
- int u, v; ll cost, cap, flow;
- edge(int u, int v, ll cap, ll cost) :
- u(u), v(v), cap(cap), cost(cost), flow(0) {} };
struct flow_network {
- int n, s, t, *par, *in_queue, *num_vis; ll *dist, *pot;
- std::vector<edge> edges;
- std::vector<int> *adj;
- std::map<std::pair<int, int>, std::vector<int>> edge_idx;
- flow_network(int n, int s, int t) : n(n), s(s), t(t) {
- adj = new std::vector<int>[n];
- par = new int[n];
- in_queue = new int[n];
- num_vis = new int[n];
- dist = new ll[n];
- pot = new ll[n];
- for (int u = 0; u < n; ++u) pot[u] = 0; }
- void add_edge(int u, int v, ll cap, ll cost) {
- adj[u].push_back(edges.size());
- edge_idx[{u, v}].push_back(edges.size());
- edges.push_back(edge(u, v, cap, cost));
- adj[v].push_back(edges.size());
- edge_idx[{v, u}].push_back(edges.size());
- edges.push_back(edge(v, u, 0LL, -cost)); }
- ll get_flow(int u, int v) {
- ll f = 0;
- for (int i : edge_idx[{u, v}]) f += edges[i].flow;
- return f; }
- ll res(edge &e) { return e.cap - e.flow; }
- void bellman_ford() {
- for (int u = 0; u < n; ++u) pot[u] = INF;
- pot[s] = 0;
- for (int it = 0; it < n-1; ++it)
- for (auto e : edges)
- if (res(e) > 0)
- pot[e.v] = std::min(pot[e.v], pot[e.u] + e.cost); }
- bool spfa () {
- std::queue<int> q; q.push(s);
- while (not q.empty()) {
- int u = q.front(); q.pop(); in_queue[u] = 0;
- if (++num_vis[u] >= n) {
- dist[u] = -INF;
- return false; }
- for (int i : adj[u]) {
- edge e = edges[i];
- if (res(e) <= 0) continue;
- ll nd = dist[u] + e.cost + pot[u] - pot[e.v];
- if (dist[e.v] > nd) {
- dist[e.v] = nd;
- par[e.v] = i;
- if (not in_queue[e.v]) {
- q.push(e.v);

```

```

- in_queue[e.v] = 1; } } } }
- return dist[t] != INF; }
- bool aug_path() {
- for (int u = 0; u < n; ++u) {
- par[u] = -1;
- in_queue[u] = 0;
- num_vis[u] = 0;
- dist[u] = INF; }
- dist[s] = 0;
- in_queue[s] = 1;
- return spfa();
- }
- pll calc_max_flow(bool do_bellman_ford=false) {
- ll total_cost = 0, total_flow = 0;
- if (do_bellman_ford)
- bellman_ford();
- while (aug_path()) {
- ll f = INF;
- for (int i = par[t]; i != -1; i = par[edges[i].u])
- f = std::min(f, res(edges[i]));
- for (int i = par[t]; i != -1; i = par[edges[i].u]) {
- edges[i].flow += f;
- edges[i^1].flow -= f; }
- total_cost += f * (dist[t] + pot[t] - pot[s]);
- total_flow += f;
- for (int u = 0; u < n; ++u)
- if (par[u] != -1) pot[u] += dist[u]; }
- return {total_cost, total_flow}; } };

```

##### 4.10.1. Hungarian Algorithm.

```

int n, m; // size of A, size of B
int cost[N+1][N+1]; // input cost matrix, 1-indexed
int way[N+1], p[N+1]; // p[i]=j: Ai is matched to Bj
int minv[N+1], A[N+1], B[N+1]; bool used[N+1];
int hungarian() {
- for (int i = 0; i <= N; ++i)
- A[i] = B[i] = p[i] = way[i] = 0; // init
- for (int i = 1; i <= n; ++i) {
- p[0] = i; int R = 0;
- for (int j = 0; j <= m; ++j)
- minv[j] = INF, used[j] = false;
- do {
- int L = p[R], dR = 0;
- int delta = INF;
- used[R] = true;
- for (int j = 1; j <= m; ++j)
- if (!used[j]) {
- int c = cost[L][j] - A[L] - B[j];
- if (c < minv[j]) minv[j] = c, way[j] = R;
- if (minv[j] < delta) delta = minv[j], dR = j;
- }
- for (int j = 0; j <= m; ++j)
- if (used[j]) A[p[j]] += delta, B[j] -= delta;
- else minv[j] -= delta;
- R = dR;
- } while (p[R] != 0);

```

```

- for (; R != 0; R = way[R])
- p[R] = p[way[R]]; }
- return -B[0]; }

```

4.11. **Minimum Arborescence**. Given a weighted directed graph, finds a subset of edges of minimum total weight so that there is a unique path from the root  $r$  to each vertex. Returns a vector of size  $n$ , where the  $i$ th element is the edge for the  $i$ th vertex. The answer for the root is undefined!

```

#include "../data-structures/union_find.cpp"
struct arborescence {
- int n; union_find uf;
- vector<vector<pair<ii,int>>> adj;
- arborescence(int _n) : n(_n), uf(n), adj(n) {}
- void add_edge(int a, int b, int c) {
- adj[b].push_back(make_pair(ii(a,b),c)); }
- vii find_min(int r) {
- vi vis(n,-1), mn(n,INF); vii par(n);
- rep(i,0,n) {
- if (uf.find(i) != i) continue;
- int at = i;
- while (at != r && vis[at] == -1) {
- vis[at] = i;
- iter(it,adj[at]) if (it->second < mn[at] &&
- uf.find(it->first.first) != at)
- mn[at] = it->second, par[at] = it->first;
- if (par[at] == ii(0,0)) return vii();
- at = uf.find(par[at].first); }
- if (at == r || vis[at] != i) continue;
- union_find tmp = uf; vi seq;
- do { seq.push_back(at); at = uf.find(par[at].first);
- } while (at != seq.front());
- iter(it,seq) uf.unite(*it,seq[0]);
- int c = uf.find(seq[0]);
- vector<pair<ii,int>> > nw;
- iter(it,seq) iter(jt,adj[*it])
- nw.push_back(make_pair(jt->first,
- jt->second - mn[*it]));
- adj[c] = nw;
- vii rest = find_min(r);
- if (size(rest) == 0) return rest;
- ii use = rest[c];
- rest[at = tmp.find(use.second)] = use;
- iter(it,seq) if (*it != at)
- rest[*it] = par[*it];
- return rest; }
- return par; } };

```

4.12. **Blossom algorithm**. Finds a maximum matching in an arbitrary graph in  $O(|V|^4)$  time. Be vary of loop edges.

```

#define MAXV 300
bool marked[MAXV], emarked[MAXV][MAXV];
int S[MAXV];
vi find_augmenting_path(const vector<vi> &adj, const vi &m){
- int n = size(adj), s = 0;
- vi par(n,-1), height(n), root(n,-1), q, a, b;

```

```

- memset(marked,0,sizeof(marked));
- memset(emarked,0,sizeof(emarked));
- rep(i,0,n) if (m[i] >= 0) emarked[i][m[i]] = true;
-       else root[i] = i, S[s++] = i;
- while (s) {
-   int v = S[--s];
-   iter(wt,adj[v]) {
-     int w = *wt;
-     if (emarked[v][w]) continue;
-     if (root[w] == -1) {
-       int x = S[s++] = m[w];
-       par[w]=v, root[w]=root[v], height[w]=height[v]+1;
-       par[x]=w, root[x]=root[w], height[x]=height[w]+1;
-     } else if (height[w] % 2 == 0) {
-       if (root[v] != root[w]) {
-         while (v != -1) q.push_back(v), v = par[v];
-         reverse(q.begin(), q.end());
-         while (w != -1) q.push_back(w), w = par[w];
-         return q;
-       } else {
-         int c = v;
-         while (c != -1) a.push_back(c), c = par[c];
-         c = w;
-         while (c != -1) b.push_back(c), c = par[c];
-         while (!a.empty() && !b.empty() && a.back() == b.back())
-           c = a.back(), a.pop_back(), b.pop_back();
-         memset(marked,0,sizeof(marked));
-         fill(par.begin(), par.end(), 0);
-         iter(it,a) par[*it] = 1; iter(it,b) par[*it] = 1;
-         par[c] = s = 1;
-         rep(i,0,n) root[par[i] = par[i] ? 0 : s++] = i;
-         vector<vi> adj2(s);
-         rep(i,0,n) iter(it,adj[i]) {
-           if (par[*it] == 0) continue;
-           if (par[i] == 0) {
-             if (!marked[par[*it]]) {
-               adj2[par[i]].push_back(par[*it]);
-               adj2[par[*it]].push_back(par[i]);
-               marked[par[*it]] = true;
-             }
-             } else adj2[par[i]].push_back(par[*it]);
-         }
-         vi m2(s, -1);
-         if (m[c] != -1) m2[m2[par[m[c]]] = 0] = par[m[c]];
-         rep(i,0,n) if (par[i] != 0 && m[i] != -1 && par[m[i]] != 0)
-           m2[par[i]] = par[m[i]];
-         vi p = find_augmenting_path(adj2, m2);
-         int t = 0;
-         while (t < size(p) && p[t]) t++;
-         if (t == size(p)) {
-           rep(i,0,size(p)) p[i] = root[p[i]];
-           return p;
-         }
-         if (!p[0] || (m[c] != -1 && p[t+1] != par[m[c]]))
-           reverse(p.begin(), p.end()), t=(int)size(p)-t-1;
-         rep(i,0,t) q.push_back(root[p[i]]);
-         iter(it,adj[root[p[t-1]]) {
-           if (par[*it] != (s = 0)) continue;
-           a.push_back(c), reverse(a.begin(), a.end());

```

```

-       iter(jt,b) a.push_back(*jt);
-       while (a[s] != *it) s++;
-       if ((height[*it]&1)^(s<(int)size(a)-(int)size(b)))
-         reverse(a.begin(),a.end()), s=(int)size(a)-s-1;
-       while(a[s]!=c)q.push_back(a[s]),s=(s+1)%size(a);
-       q.push_back(c);
-       rep(i,t+1,size(p)) q.push_back(root[p[i]]);
-       return q; } }
-   emarked[v][w] = emarked[w][v] = true;
-   marked[v] = true; } return q; }
vii max_matching(const vector<vi> &adj) {
-   vi m(size(adj), -1), ap; vii res, es;
-   rep(i,0,size(adj)) iter(it,adj[i]) es.emplace_back(i,*it);
-   random_shuffle(es.begin(), es.end());
-   iter(it,es) if (m[it->first] == -1 && m[it->second] == -1)
-     m[it->first] = it->second, m[it->second] = it->first;
-   do { ap = find_augmenting_path(adj, m);
-     rep(i,0,size(ap)) m[m[ap[i^1]] = ap[i]] = ap[i^1];
-   } while (!ap.empty());
-   rep(i,0,size(m)) if (i < m[i]) res.emplace_back(i, m[i]);
-   return res; }

```

4.13. **Maximum Density Subgraph.** Given (weighted) undirected graph  $G$ . Binary search density. If  $g$  is current density, construct flow network:  $(S, u, m), (u, T, m + 2g - d_u), (u, v, 1)$ , where  $m$  is a large constant (larger than sum of edge weights). Run floating-point max-flow. If minimum cut has empty  $S$ -component, then maximum density is smaller than  $g$ , otherwise it's larger. Distance between valid densities is at least  $1/(n(n-1))$ . Edge case when density is 0. This also works for weighted graphs by replacing  $d_u$  by the weighted degree, and doing more iterations (if weights are not integers).

4.14. **Maximum-Weight Closure.** Given a vertex-weighted directed graph  $G$ . Turn the graph into a flow network, adding weight  $\infty$  to each edge. Add vertices  $S, T$ . For each vertex  $v$  of weight  $w$ , add edge  $(S, v, w)$  if  $w \geq 0$ , or edge  $(v, T, -w)$  if  $w < 0$ . Sum of positive weights minus minimum  $S - T$  cut is the answer. Vertices reachable from  $S$  are in the closure. The maximum-weight closure is the same as the complement of the minimum-weight closure on the graph with edges reversed.

4.15. **Maximum Weighted Ind. Set in a Bipartite Graph.** This is the same as the minimum weighted vertex cover. Solve this by constructing a flow network with edges  $(S, u, w(u))$  for  $u \in L, (v, T, w(v))$  for  $v \in R$  and  $(u, v, \infty)$  for  $(u, v) \in E$ . The minimum  $S, T$ -cut is the answer. Vertices adjacent to a cut edge are in the vertex cover.

4.16. **Synchronizing word problem.** A DFA has a synchronizing word (an input sequence that moves all states to the same state) iff. each pair of states has a synchronizing word. That can be checked using reverse DFS over pairs of states. Finding the shortest synchronizing word is NP-complete.

4.17. **Max flow with lower bounds on edges.** Change edge  $(u, v, l \leq f \leq c)$  to  $(u, v, f \leq c - l)$ . Add edge  $(t, s, \infty)$ . Create super-nodes  $S, T$ . Let  $M(u) = \sum_v l(v, u) - \sum_v l(u, v)$ . If  $M(u) < 0$ , add edge  $(u, T, -M(u))$ , else add edge  $(S, u, M(u))$ . Max flow from  $S$  to  $T$ . If all edges from  $S$  are saturated, then we have a feasible flow. Continue running max flow from  $s$  to  $t$  in original graph.

4.18. **Tutte matrix for general matching.** Create an  $n \times n$  matrix  $A$ . For each edge  $(i, j)$ ,  $i < j$ , let  $A_{ij} = x_{ij}$  and  $A_{ji} = -x_{ij}$ . All other entries are 0. The determinant of  $A$  is zero iff. the graph has a perfect matching. A randomized algorithm uses the Schwartz-Zippel lemma to check if it is zero.

4.19. **Heavy Light Decomposition.**

```

#include "segment_tree.cpp"
struct heavy_light_tree {
-   int n, *par, *heavy, *dep, *path_root, *pos;
-   std::vector<int> adj;
-   segtree segment_tree;
-   heavy_light_tree(int n) : n(n) {
-     this->adj = new std::vector<int>[n];
-     segment_tree = new segtree(0, n-1);
-     par = new int[n];
-     heavy = new int[n];
-     dep = new int[n];
-     path_root = new int[n];
-     pos = new int[n];
-   }
-   void add_edge(int u, int v) {
-     adj[u].push_back(v);
-     adj[v].push_back(u);
-   }
-   void build(int root) {
-     for (int u = 0; u < n; ++u)
-       heavy[u] = -1;
-     par[root] = root;
-     dep[root] = 0;
-     dfs(root);
-     for (int u = 0, p = 0; u < n; ++u) {
-       if (par[u] == -1 or heavy[par[u]] != u) {
-         for (int v = u; v != -1; v = heavy[v]) {
-           path_root[v] = u;
-           pos[v] = p++;
-         }
-       }
-     }
-     int dfs(int u) {
-       int sz = 1;
-       int max_subtree_sz = 0;
-       for (int v : adj[u]) {
-         if (v != par[u]) {
-           par[v] = u;
-           dep[v] = dep[u] + 1;
-           int subtree_sz = dfs(v);
-           if (max_subtree_sz < subtree_sz) {
-             max_subtree_sz = subtree_sz;
-             heavy[u] = v;
-           }
-           sz += subtree_sz;
-         }
-       }
-       return sz;
-     }
-     int query(int u, int v) {
-       int res = 0;
-       while (path_root[u] != path_root[v]) {
-         if (dep[path_root[u]] > dep[path_root[v]])
-           std::swap(u, v);
-         res += segment_tree->sum(pos[path_root[v]], pos[v]);
-         v = par[path_root[v]];
-       }
-       res += segment_tree->sum(pos[u], pos[v]);
-       return res;
-     }

```

```
- void update(int u, int v, int c) {
-- for (; path_root[u] != path_root[v]; v = par[path_root[v]]){
----- if (dep[path_root[u]] > dep[path_root[v]])
----- std::swap(u, v);
----- segment_tree->increase(pos[path_root[v]], pos[v], c); }
-- segment_tree->increase(pos[u], pos[v], c); } };
```

4.20. Centroid Decomposition .

```
#define MAXV 100100
#define LGMAXV 20
int jmp[MAXV][LGMAXV],
  path[MAXV][LGMAXV],
  sz[MAXV], seph[MAXV],
  shortest[MAXV];
struct centroid_decomposition {
- int n; vvi adj;
- centroid_decomposition(int _n) : n(_n), adj(n) { }
- void add_edge(int a, int b) {
-- adj[a].push_back(b); adj[b].push_back(a); }
- int dfs(int u, int p) {
-- sz[u] = 1;
-- rep(i,0,size(adj[u]))
---- if (adj[u][i] != p) sz[u] += dfs(adj[u][i], u);
-- return sz[u]; }
- void makepaths(int sep, int u, int p, int len) {
-- jmp[u][seph[sep]] = sep, path[u][seph[sep]] = len;
-- int bad = -1;
-- rep(i,0,size(adj[u])) {
---- if (adj[u][i] == p) bad = i;
---- else makepaths(sep, adj[u][i], u, len + 1); }
-- if (p == sep)
---- swap(adj[u][bad], adj[u].back()), adj[u].pop_back(); }
- void separate(int h=0, int u=0) {
-- dfs(u,-1); int sep = u;
-- down: iter(nxt,adj[sep])
---- if (sz[*nxt] < sz[sep] && sz[*nxt] > sz[u]/2) {
----- sep = *nxt; goto down; }
-- seph[sep] = h, makepaths(sep, sep, -1, 0);
-- rep(i,0,size(adj[sep])) separate(h+1, adj[sep][i]); }
- void paint(int u) {
-- rep(h,0,seph[u]+1)
---- shortest[jmp[u][h]] = min(shortest[jmp[u][h]],
----- path[u][h]); }

- int closest(int u) {
-- int mn = INF/2;
-- rep(h,0,seph[u]+1)
---- mn = min(mn, path[u][h] + shortest[jmp[u][h]]);
-- return mn; } };
```

4.21. Least Common Ancestor.

4.21.1. Binary Lifting.

```
struct graph {
- int n, logn, *dep, **par;
- std::vector<int> *adj;
- graph(int n, int logn=20) : n(n), logn(logn) {
-- adj = new std::vector<int>[n];
```

```
dep = new int[n];
par = new int*[n];
for (int i = 0; i < n; ++i) par[i] = new int[logn]; }
void dfs(int u, int p, int d) {
-- dep[u] = d;
-- par[u][0] = p;
-- for (int v : adj[u])
---- if (v != p) dfs(v, u, d+1); }
- int ascend(int u, int k) {
-- for (int i = 0; i < logn; ++i)
---- if (k & (1 << i)) u = par[u][i];
-- return u; }
- int lca(int u, int v) {
-- if (dep[u] > dep[v]) u = ascend(u, dep[u] - dep[v]);
-- if (dep[v] > dep[u]) v = ascend(v, dep[v] - dep[u]);
-- if (u == v) return u;
-- for (int k = logn-1; k >= 0; --k) {
---- if (par[u][k] != par[v][k]) {
----- u = par[u][k]; v = par[v][k]; } }
-- return par[u][0]; }
- bool is_anc(int u, int v) {
-- if (dep[u] < dep[v]) std::swap(u, v);
-- return ascend(u, dep[u] - dep[v]) == v; }
- void prep_lca(int root=0) {
-- dfs(root, root, 0);
-- for (int k = 1; k < logn; ++k)
---- for (int u = 0; u < n; ++u)
----- par[u][k] = par[par[u][k-1]][k-1]; } }
```

4.21.2. Euler Tour Sparse Table.

```
struct graph {
- int n, logn, *par, *dep, *first, *lg, **spt;
- vi *adj, euler; // spt size should be ~ 2n
- graph(int n, int logn=20) : n(n), logn(logn) {
-- adj = new vi[n];
-- par = new int[n];
-- dep = new int[n];
-- first = new int[n]; }
- void add_edge(int u, int v) {
-- adj[u].push_back(v); adj[v].push_back(u); }
- void dfs(int u, int p, int d) {
-- dep[u] = d; par[u] = p;
-- first[u] = euler.size();
-- euler.push_back(u);
-- for (int v : adj[u])
---- if (v != p) {
----- dfs(v, u, d+1);
----- euler.push_back(u); } }
- void prep_lca(int root=0) {
-- dfs(root, root, 0);
-- int en = euler.size();
-- lg = new int[en+1];
-- lg[0] = lg[1] = 0;
-- for (int i = 2; i <= en; ++i)
---- lg[i] = lg[i >> 1] + 1;
-- spt = new int*[en];
```

```
for (int i = 0; i < en; ++i) {
-- spt[i] = new int[lg[en]];
-- spt[i][0] = euler[i]; }
for (int k = 0; (2 << k) <= en; ++k)
-- for (int i = 0; i + (2 << k) <= en; ++i)
---- if (dep[spt[i][k]] < dep[spt[i+(1<<k)][k]])
----- spt[i][k+1] = spt[i][k];
-- else
---- spt[i][k+1] = spt[i+(1<<k)][k]; }
- int lca(int u, int v) {
-- int a = first[u], b = first[v];
-- if (a > b) std::swap(a, b);
-- int k = lg[b-a+1], ba = b - (1 << k) + 1;
-- if (dep[spt[a][k]] < dep[spt[ba][k]]) return spt[a][k];
-- return spt[ba][k]; } };
```

4.21.3. Tarjan Off-line LCA.

```
#include "../data-structures/union_find.cpp"
struct tarjan_olca {
- int *ancestor;
- vi *adj, answers;
- vii *queries;
- bool *colored;
- union_find uf;
- tarjan_olca(int n, vi *_adj) : adj(_adj), uf(n) {
-- colored = new bool[n];
-- ancestor = new int[n];
-- queries = new vii[n];
-- memset(colored, 0, n); }
- void query(int x, int y) {
-- queries[x].push_back(ii(y, size(answers)));
-- queries[y].push_back(ii(x, size(answers)));
-- answers.push_back(-1); }
- void process(int u) {
-- ancestor[u] = u;
-- rep(i,0,size(adj[u])) {
---- int v = adj[u][i];
---- process(v);
---- uf.unite(u,v);
---- ancestor[uf.find(u)] = u; }
-- colored[u] = true;
-- rep(i,0,size(queries[u])) {
---- int v = queries[u][i].first;
---- if (colored[v])
----- answers[queries[u][i].second] = ancestor[uf.find(v)];
-- } } };
```

4.22. Counting Spanning Trees. Kirchoff's Theorem: The number of spanning trees of any graph is the determinant of any cofactor of the Laplacian matrix in  $O(n^3)$ .

- (1) Let  $A$  be the adjacency matrix.
- (2) Let  $D$  be the degree matrix (matrix with vertex degrees on the diagonal).
- (3) Get  $D - A$  and delete exactly one row and column. Any row and column will do. This will be the cofactor matrix.
- (4) Get the determinant of this cofactor matrix using Gauss-Jordan.
- (5) Spanning Trees = |cofactor( $D - A$ )|



4.23. **Erdős-Gallai Theorem.** A sequence of non-negative integers  $d_1 \geq \cdots \geq d_n$  can be represented as the degree sequence of finite simple graph on  $n$  vertices if and only if  $d_1 + \cdots + d_n$  is even and the following holds for  $1 \leq k \leq n$ :

$$\sum_{i=1}^n d_i \leq k(k-1) + \sum_{i=k+1}^n \min(d_i, k)$$

4.24. **Tree Isomorphism**.

```
// REQUIREMENT: list of primes pr[], see prime sieve -----
typedef long long LL;
int pre[N], q[N], path[N]; bool vis[N];
// perform BFS and return the last node visited -----
int bfs(int u, vector<int> adj[]) {
    memset(vis, 0, sizeof(vis));
    int head = 0, tail = 0;
    q[tail++] = u; vis[u] = true; pre[u] = -1;
    while (head != tail) {
        u = q[head]; if (++head == N) head = 0;
        for (int i = 0; i < adj[u].size(); ++i) {
            int v = adj[u][i];
            if (!vis[v]) {
                vis[v] = true; pre[v] = u;
                q[tail++] = v; if (tail == N) tail = 0;
            }
        }
    }
    return u;
} // returns the list of tree centers -----
vector<int> tree_centers(int r, vector<int> adj[]) {
    int size = 0;
    for (int u=bfs(bfs(r, adj), adj); u!=-1; u=pre[u])
        path[size++] = u;
    vector<int> med(1, path[size/2]);
    if (size % 2 == 0) med.push_back(path[size/2-1]);
    return med;
} // returns "unique hashcode" for tree with root u -----
LL rootcode(int u, vector<int> adj[], int p=-1, int d=15){
    vector<LL> k; int nd = (d + 1) % primes;
    for (int i = 0; i < adj[u].size(); ++i)
        if (adj[u][i] != p)
            k.push_back(rootcode(adj[u][i], adj, u, nd));
    sort(k.begin(), k.end());
    LL h = k.size() + 1;
    for (int i = 0; i < k.size(); ++i)
        h = h * pr[d] + k[i];
    return h;
} // returns "unique hashcode" for the whole tree -----
LL treecode(int root, vector<int> adj[]) {
    vector<int> c = tree_centers(root, adj);
    if (c.size()==1)
        return (rootcode(c[0], adj) << 1) | 1;
    return (rootcode(c[0], adj)*rootcode(c[1], adj))<<1;
} // checks if two trees are isomorphic -----
bool isomorphic(int r1, vector<int> adj1[], int r2,
                vector<int> adj2[], bool rooted = false) {
    if (rooted)
```

```
        return rootcode(r1, adj1) == rootcode(r2, adj2);
    return treecode(r1, adj1) == treecode(r2, adj2); }

5. MATH I - ALGEBRA

5.1. Generating Function Manager.
const int DEPTH = 19;
const int ARR_DEPTH = (1<<DEPTH); //approx 5x10^5
const int SZ = 12;
ll temp[SZ][ARR_DEPTH+1];
const ll MOD = 998244353;
struct GF_Manager {
    int tC = 0;
    std::stack<int> to_be_freed;
    const static ll DEPTH = 23;
    ll prim[DEPTH+1], prim_inv[DEPTH+1], two_inv[DEPTH+1];
    ll mod_pow(ll base, ll exp) {
        if(exp==0) return 1;
        if(exp&1) return (base*mod_pow(base,exp-1))%MOD;
        else return mod_pow((base*base)%MOD, exp/2); }
    void set_up_primitives() {
        prim[DEPTH] = 31;
        prim_inv[DEPTH] = mod_pow(prim[DEPTH], MOD-2);
        two_inv[DEPTH] = mod_pow(1<<DEPTH,MOD-2);
        for(int n = DEPTH-1; n >= 0; n--) {
            prim[n] = (prim[n+1]*prim[n+1])%MOD;
            prim_inv[n] = mod_pow(prim[n],MOD-2);
            two_inv[n] = mod_pow(1<<n,MOD-2); } }
    GF_Manager(){ set_up_primitives(); }
    void start_claiming(){ to_be_freed.push(0); }
    ll* claim(){
        ++to_be_freed.top(); assert(tC < SZ); return temp[tC++]; }
    void end_claiming(){tC=to_be_freed.top(); to_be_freed.pop();}
    void NTT(ll A[], int n, ll t[],
             bool is_inverse=false, int offset=0) {
        if (n==0) return;
        //Put the evens first, then the odds
        for (int i = 0; i < (1<<(n-1)); i++) {
            t[i] = A[offset+2*i];
            t[i+(1<<(n-1))] = A[offset+2*i+1]; }
        for(int i = 0; i < (1<<n); i++)
            A[offset+i] = t[i];
        NTT(A, n-1, t, is_inverse, offset);
        NTT(A, n-1, t, is_inverse, offset+(1<<(n-1)));
        ll w1 = (is_inverse ? prim_inv[n] : prim[n]), w = 1;
        for (int i = 0; i < (1<<(n-1)); i++, w=(w*w1)%MOD) {
            t[i] = (A[offset+i] + w*A[offset+(1<<(n-1))+i])%MOD;
            t[i+(1<<(n-1))] = (A[offset+i]-
                               w*A[offset+(1<<(n-1))+i])%MOD; }
        for (int i = 0; i < (1<<n); i++) A[offset+i] = t[i];
    }
    int add(ll A[], int an, ll B[], int bn, ll C[]) {
        int cn = 0;
        for(int i = 0; i < max(an,bn); i++) {
            C[i] = A[i]+B[i];
            if(C[i] <= -MOD) C[i] += MOD;
            if(MOD <= C[i]) C[i] -= MOD;
        }
    }
    int subtract(ll A[], int an, ll B[], int bn, ll C[]) {
        int cn = 0;
        for(int i = 0; i < max(an,bn); i++) {
            C[i] = A[i]-B[i];
            if(C[i] <= -MOD) C[i] += MOD;
            if(MOD <= C[i]) C[i] -= MOD;
        }
        return cn+1; }
    int scalar_mult(ll v, ll A[], int an, ll C[]) {
        for(int i = 0; i < an; i++) C[i] = (v*A[i])%MOD;
        return v==0 ? 0 : an; }
    int mult(ll A[], int an, ll B[], int bn, ll C[]) {
        start_claiming();
        // make sure you've called setup prim first
        // note: an and bn refer to the *number of items in
        // each array*, NOT the degree of the largest term
        int n, degree = an+bn-1;
        for(n=0; (1<<n) < degree; n++);
        ll *tA = claim(), *tB = claim(), *t = claim();
        copy(A,A+an,tA); fill(tA+an,tA+(1<<n),0);
        copy(B,B+bn,tB); fill(tB+bn,tB+(1<<n),0);
        NTT(tA,n,t);
        NTT(tB,n,t);
        for(int i = 0; i < (1<<n); i++)
            tA[i] = (tA[i]*tB[i])%MOD;
        NTT(tA,n,t,true);
        scalar_mult(two_inv[n],tA,degree,C);
        end_claiming();
        return degree; }
    int reciprocal(ll F[], int fn, ll R[]) {
        start_claiming();
        ll *tR = claim(), *tempR = claim();
        int n; for(n=0; (1<<n) < fn; n++);
        fill(tempR,tempR+(1<<n),0);
        tempR[0] = mod_pow(F[0],MOD-2);
        for (int i = 1; i <= n; i++) {
            mult(tempR,1<<i,F,1<<i,tR);
            tR[0] -= 2;
            scalar_mult(-1,tR,1<<i,tR);
            mult(tempR,1<<i,tR,1<<i,tempR); }
        copy(tempR,tempR+fn,R);
        end_claiming();
        return n; }
    int quotient(ll F[], int fn, ll G[], int gn, ll Q[]) {
        start_claiming();
        ll* revF = claim();
        ll* revG = claim();
        ll* tempQ = claim();
        copy(F,F+fn,revF); reverse(revF,revF+fn);
        copy(G,G+gn,revG); reverse(revG,revG+gn);
        int qn = fn-gn+1;
        reciprocal(revG,qn,revG);
        mult(revF,qn,revG,qn,tempQ);
        reverse(tempQ,tempQ+qn);
```

```
-- copy(tempQ,tempQ+qn,Q);
-- end_claiming();
-- return qn; }
int mod(ll F[], int fn, ll G[], int gn, ll R[]) {
start_claiming();
ll *Q = claim(), *GQ = claim();
int qn = quotient(F, fn, G, gn, Q);
int gqn = mult(G, gn, Q, qn, GQ);
int rn = subtract(F, fn, GQ, gqn, R);
end_claiming();
return rn; }
ll horners(ll F[], int fn, ll xi) {
ll ans = 0;
for(int i = fn-1; i >= 0; i--)
ans = (ans*xi+F[i]) % MOD;
return ans; } };
GF_Manager gfManager;
ll split[DEPTH+1][2*(ARR_DEPTH)+1];
ll Fi[DEPTH+1][2*(ARR_DEPTH)+1];
int bin_splitting(ll a[], int l, int r, int s=0, int offset=0) {
if(l == r) {
split[s][offset] = a[l]; //x^0
split[s][offset+1] = 1; //x^1
return 2; }
int m = (l+r)/2;
int sz = m-l+1;
int da = bin_splitting(a, l, m, s+1, offset);
int db = bin_splitting(a, m+1, r, s+1, offset+(sz<<1));
return gfManager.mult(split[s+1]+offset, da,
split[s+1]+offset+(sz<<1), db, split[s]+offset); }
void multipoint_eval(ll a[], int l, int r, ll F[], int fn,
ll ans[], int s=0, int offset=0) {
if(l == r) {
ans[l] = gfManager.horners(F,fn,a[l]);
return; }
int m = (l+r)/2;
int sz = m-l+1;
int da = gfManager.mod(F, fn, split[s+1]+offset,
sz+1, Fi[s]+offset);
int db = gfManager.mod(F, fn, split[s+1]+offset+(sz<<1),
r-m+1, Fi[s]+offset+(sz<<1));
multipoint_eval(a,l,m,Fi[s]+offset,da,ans,s+1,offset);
multipoint_eval(a,m+1,r,Fi[s]+offset+(sz<<1),
db,ans,s+1,offset+(sz<<1));
}
```

5.2. **Fast Fourier Transform.** Compute the Discrete Fourier Transform (DFT) of a polynomial in  $O(n \log n)$  time.

```
struct poly {
double a, b;
poly(double a=0, double b=0): a(a), b(b) {}
poly operator+(const poly& p) const {
return poly(a + p.a, b + p.b);}
poly operator-(const poly& p) const {
return poly(a - p.a, b - p.b);}
poly operator*(const poly& p) const {
```

```
-- return poly(a*p.a - b*p.b, a*p.b + b*p.a);}
};
void fft(poly in[], poly p[], int n, int s) {
if (n < 1) return;
if (n == 1) {p[0] = in[0]; return;}
n >>= 1; fft(in, p, n, s << 1);
fft(in + s, p + n, n, s << 1);
poly w(1), wn(cos(M_PI/n), sin(M_PI/n));
for (int i = 0; i < n; ++i) {
poly even = p[i], odd = p[i + n];
p[i] = even + w * odd;
p[i + n] = even - w * odd;
w = w * wn;
}
}
void fft(poly p[], int n) {
poly *f = new poly[n]; fft(p, f, n, 1);
copy(f, f + n, p); delete[] f;
}
void inverse_fft(poly p[], int n) {
for(int i=0; i<n; i++) {p[i].b *= -1;} fft(p, n);
for(int i=0; i<n; i++) {p[i].a/=n; p[i].b/= -1*n;}
```

5.3. **FFT Polynomial Multiplication.** Multiply integer polynomials  $a, b$  of size  $an, bn$  using FFT in  $O(n \log n)$ . Stores answer in an array  $c$ , rounded to the nearest integer (or double).

```
// note: c[] should have size of at least (an+bn)
int mult(int a[],int an,int b[],int bn,int c[]) {
int n, degree = an + bn - 1;
for (n = 1; n < degree; n <= 1); // power of 2
poly *A = new poly[n], *B = new poly[n];
copy(a, a + an, A); fill(A + an, A + n, 0);
copy(b, b + bn, B); fill(B + bn, B + n, 0);
fft(A, n); fft(B, n);
for (int i = 0; i < n; i++) A[i] = A[i] * B[i];
inverse_fft(A, n);
for (int i = 0; i < degree; i++)
c[i] = int(A[i].a + 0.5); // same as round(A[i].a)
delete[] A, B; return degree;
}
```

5.4. **Number Theoretic Transform.** Other possible moduli: 2113929217( $2^{25}$ ), 2013265920268435457( $2^{28}$ ,  $withg = 5$ )

```
#include "../mathematics/primitive_root.cpp"
int mod = 998244353, g = primitive_root(mod),
ginv = mod_pow<ll>(g, mod-2, mod),
inv2 = mod_pow<ll>(2, mod-2, mod);
#define MAXN (1<<22)
struct Num {
int x;
Num(ll _x=0) { x = (_x%mod+mod)%mod; }
Num operator +(const Num &b) { return x + b.x; }
Num operator -(const Num &b) const { return x - b.x; }
Num operator *(const Num &b) const { return (ll)x * b.x; }
Num operator /(const Num &b) const {
return (ll)x * b.inv().x; }
```

```
- Num inv() const { return mod_pow<ll>((ll)x, mod-2, mod); }
- Num pow(int p) const { return mod_pow<ll>((ll)x, p, mod); }
} T1[MAXN], T2[MAXN];
void ntt(Num x[], int n, bool inv = false) {
Num z = inv ? ginv : g;
z = z.pow((mod - 1) / n);
for (ll i = 0, j = 0; i < n; i++) {
if (i < j) swap(x[i], x[j]);
ll k = n>>1;
while (1 <= k && k <= j) j -= k, k >>= 1;
j += k; }
for (int mx = 1, p = n/2; mx < n; mx <= 1, p >>= 1) {
Num wp = z.pow(p), w = 1;
for (int k = 0; k < mx; k++, w = w*wp) {
for (int i = k; i < n; i += mx << 1) {
Num t = x[i + mx] * w;
x[i + mx] = x[i] - t;
x[i] = x[i] + t; } } }
if (inv) {
Num ni = Num(n).inv();
rep(i,0,n) { x[i] = x[i] * ni; } }
void inv(Num x[], Num y[], int l) {
if (l == 1) { y[0] = x[0].inv(); return; }
inv(x, y, l>>1);
// NOTE: maybe l<<2 instead of l<<1
rep(i,l>>1,l<<1) T1[i] = y[i] = 0;
rep(i,0,l) T1[i] = x[i];
ntt(T1, l<<1); ntt(y, l<<1);
rep(i,0,l<<1) y[i] = y[i]*2 - T1[i] * y[i] * y[i];
ntt(y, l<<1, true); }
void sqrt(Num x[], Num y[], int l) {
if (l == 1) { assert(x[0].x == 1); y[0] = 1; return; }
sqrt(x, y, l>>1);
inv(y, T2, l>>1);
rep(i,l>>1,l<<1) T1[i] = T2[i] = 0;
rep(i,0,l) T1[i] = x[i];
ntt(T2, l<<1); ntt(T1, l<<1);
rep(i,0,l<<1) T2[i] = T1[i] * T2[i];
ntt(T2, l<<1, true);
rep(i,0,l) y[i] = (y[i] + T2[i]) * inv2; }
// vim: cc=60 ts=2 sts=2 sw=2:
```

5.5. **Polynomial Long Division.** Divide two polynomials  $A$  and  $B$  to get  $Q$  and  $R$ , where  $\frac{A}{B} = Q + \frac{R}{B}$ .

```
typedef vector<double> Poly;
Poly Q, R; // quotient and remainder
void trim(Poly& A) { // remove trailing zeroes
while (!A.empty() && abs(A.back()) < EPS)
A.pop_back(); }
void divide(Poly A, Poly B) {
if (B.size() == 0) throw exception();
if (A.size() < B.size()) {Q.clear(); R=A; return;}
Q.assign(A.size() - B.size() + 1, 0);
Poly part;
while (A.size() >= B.size()) {
```

## 6. MATH II - COMBINATORICS

- (1) The number of non-crossing partitions of an  $n$ -element set
- (2) The number of expressions with  $n$  pairs of parentheses
- (3) The number of ways  $n + 1$  factors can be parenthesized
- (4) The number of full binary trees with  $n + 1$  leaves
- (5) The number of monotonic lattice paths of an  $n \times n$  grid (5-SAT problem)
- (6) The number of triangulations of a convex polygon with  $n + 2$  sides (non-rotational)

- (7) The number of permutations  $\{1, \dots, n\}$  without a 3-term increasing subsequence

(8) The number of ways to form a mountain range with  $n$  ups and  $n$  downs

6.7. **Stirling Numbers.**  $s_1$ : Count the number of permutations of  $n$  elements with  $k$  disjoint cycles

$s_2$ : Count the ways to partition a set of  $n$  elements into  $k$  nonempty subsets

$$s_1(n,k) = \begin{cases} 1 & n = k = 0 \\ s_1(n-1,k-1) - (n-1)s_1(n-1,k) & n, k > 0 \\ 0 & \text{elsewhere} \end{cases}$$
$$s_2(n,k) = \begin{cases} 1 & n = k = 0 \\ s_2(n-1,k-1) + ks_2(n-1,k) & n, k > 0 \\ 0 & \text{elsewhere} \end{cases}$$

6.8. **Partition Function.** Pregenerate the number of partitions of positive integer  $n$  with  $n$  positive addends.

$$p(n,k) = \begin{cases} 1 & n = k = 0 \\ 0 & n < k \\ p(n-1,k-1) + p(n-k,k) & n \geq k \end{cases}$$

7. MATH III - NUMBER THEORY

7.1. **Number/Sum of Divisors.** If a number  $n$  is prime factorized where  $n = p_1^{e_1} \times p_2^{e_2} \times \dots \times p_k^{e_k}$ , where  $\sigma_0$  is the number of divisors while  $\sigma_1$  is the sum of divisors:

$$\sum_{d|n} d^k = \sigma_k(n) = \prod \frac{p_i^{k(e_i)+1} - 1}{p_i - 1}$$
$$\text{Product: } \prod_{d|n} d = n^{\frac{\sigma_1(n)}{2}}$$

7.2. **Möbius Sieve.** The Möbius function  $\mu$  is the Möbius inverse of  $e$  such that  $e(n) = \sum_{d|n} \mu(d)$ .

```
std::bitset<N> is; int mu[N];
void mobiusSieve() {
- for (int i = 1; i < N; ++i) mu[i] = 1;
- for (int i = 2; i < N; ++i) if (!is[i]) {
--- for (int j = i; j < N; j += i) { is[j] = 1; mu[j] *= -1; }
--- for (long long j = 1LL*i*i; j < N; j += i*i) mu[j] = 0; } }
```

7.3. **Möbius Inversion.** Given arithmetic functions  $f$  and  $g$ :

$$g(n) = \sum_{d|n} f(d) \quad \Leftrightarrow \quad f(n) = \sum_{d|n} \mu(d) g\left(\frac{n}{d}\right)$$

7.4. **GCD Subset Counting.** Count number of subsets  $S \subseteq A$  such that  $\gcd(S) = g$  (modifiable).

```
int f[MX+1]; // MX is maximum number of array
long long gcnt[MX+1]; // gcnt[G]: answer when gcd==G
long long C(int f) {return (1ll << f) - 1;}
// f: frequency count
// C(f): # of subsets of f elements (YOU CAN EDIT)
// Usage: int subsets_with_gcd_1 = gcnt[1];
void gcd_counter(int a[], int n) {
```

```
- memset(f, 0, sizeof f);
- memset(gcnt, 0, sizeof gcnt);
- int mx = 0;
- for (int i = 0; i < n; ++i) {
---- f[a[i]] += 1;
---- mx = max(mx, a[i]); }
- for (int i = mx; i >= 1; --i) {
--- int add = f[i];
--- long long sub = 0;
--- for (int j = 2*i; j <= mx; j += i) {
---- add += f[j];
---- sub += gcnt[j]; }
--- gcnt[i] = C(add) - sub; }
```

7.5. **Euler Totient.** Counts all integers from 1 to  $n$  that are relatively prime to  $n$  in  $O(\sqrt{n})$  time.

```
ll totient(ll n) {
- if (n <= 1) return 1;
- ll tot = n;
- for (int i = 2; i * i <= n; i++) {
--- if (n % i == 0) tot -= tot / i;
--- while (n % i == 0) n /= i; }
- if (n > 1) tot -= tot / n;
- return tot; }
```

7.6. **Extended Euclidean.** Assigns  $x, y$  such that  $ax + by = \gcd(a, b)$  and returns  $\gcd(a, b)$ .

```
ll mod(ll x, ll m) { // use this instead of x % m
- if (m == 0) return 0;
- if (m < 0) m *= -1;
- return (x%m + m) % m; // always nonnegative
}
ll extended_euclid(ll a, ll b, ll &x, ll &y) {
- if (b==0) {x = 1; y = 0; return a;}
- ll g = extended_euclid(b, a%b, x, y);
- ll z = x - a/b*y;
- x = y; y = z; return g;
}
```

7.7. **Modular Exponentiation.** Find  $b^e \pmod m$  in  $O(\log e)$  time.

```
template <class T>
T mod_pow(T b, T e, T m) {
- T res = T(1);
- while (e) {
--- if (e & T(1)) res = smod(res * b, m);
--- b = smod(b * b, m), e >>= T(1); }
- return res; }
```

7.8. **Modular Inverse.** Find unique  $x$  such that  $ax \equiv 1 \pmod m$ . Returns 0 if no unique solution is found. Please use modulo solver for the non-unique case.

```
ll modinv(ll a, ll m) {
- ll x, y; ll g = extended_euclid(a, m, x, y);
- if (g == 1 || g == -1) return mod(x * g, m);
- return 0; // 0 if invalid }
```

7.9. **Modulo Solver.** Solve for values of  $x$  for  $ax \equiv b \pmod m$ . Returns  $(-1, -1)$  if there is no solution. Returns a pair  $(x, M)$  where solution is  $x \bmod M$ .

```
pll modsolver(ll a, ll b, ll m) {
- ll x, y; ll g = extended_euclid(a, m, x, y);
- if (b % g != 0) return {-1, -1};
- return {mod(x*b/g, m/g), abs(m/g)}; }
```

7.10. **Linear Diophantine.** Computes integers  $x$  and  $y$  such that  $ax + by = c$ , returns  $(-1, -1)$  if no solution. Tries to return positive integer answers for  $x$  and  $y$  if possible.

```
pll null(-1, -1); // needs extended euclidean
pll diophantine(ll a, ll b, ll c) {
- if (!a && !b) return c ? null : {0, 0};
- if (!a) return c % b ? null : {0, c / b};
- if (!b) return c % a ? null : {c / a, 0};
- ll x, y; ll g = extended_euclid(a, b, x, y);
- if (c % g) return null;
- y = mod(y * (c/g), a/g);
- if (y == 0) y += abs(a/g); // prefer positive sol.
- return {(c - b*y)/a, y}; }
```

7.11. **Chinese Remainder Theorem.** Solves linear congruence  $x \equiv b_i \pmod{m_i}$ . Returns  $(-1, -1)$  if there is no solution. Returns a pair  $(x, M)$  where solution is  $x \bmod M$ .

```
pll chinese(ll b1, ll m1, ll b2, ll m2) {
- ll x, y; ll g = extended_euclid(m1, m2, x, y);
- if (b1 % g != b2 % g) return ii(-1, -1);
- ll M = abs(m1 / g * m2);
- return {mod(mod(x*b2*m1+y*b1*m2, M*g)/g,M), M}; }
ii chinese_remainder(ll b[], ll m[], int n) {
- ii ans(0, 1);
- for (int i = 0; i < n; ++i) {
--- ans = chinese(b[i],m[i],ans.first,ans.second);
--- if (ans.second == -1) break; }
- return ans; }
```

7.11.1. *Super Chinese Remainder.* Solves linear congruence  $a_ix \equiv b_i \pmod{m_i}$ . Returns  $(-1, -1)$  if there is no solution.

```
pll super_chinese(ll a[], ll b[], ll m[], int n) {
- pll ans(0, 1);
- for (int i = 0; i < n; ++i) {
--- pll two = modsolver(a[i], b[i], m[i]);
--- if (two.second == -1) return two;
--- ans = chinese(ans.first, ans.second, two.first, two.second);
--- if (ans.second == -1) break; }
- return ans; }
```

7.12. **Primitive Root.**

```
#include "mod_pow.cpp"
ll primitive_root(ll m) {
- vector<ll> div;
- for (ll i = 1; i*i <= m-1; i++) {
--- if ((m-1) % i == 0) {
----- if (i < m) div.push_back(i);
----- if (m/i < m) div.push_back(m/i); } }
```

```
- rep(x,2,m) { -----
-- bool ok = true; -----
-- iter(it,div) if (mod_pow<ll>(x, *it, m) == 1) { -----
---- ok = false; break; } -----
-- if (ok) return x; } -----
- return -1; } -----
```

7.13. **Josephus.** Last man standing out of  $n$  if every  $k$ th is killed. Zero-based, and does not kill 0 on first pass.

```
int J(int n, int k) { -----
- if (n == 1) return 0; -----
- if (k == 1) return n-1; -----
- if (n < k) return (J(n-1,k)+k)%n; -----
- int np = n - n/k; -----
- return k*((J(np,k)+np-n%k*np)%np) / (k-1); } -----
```

7.14. **Number of Integer Points under a Lines.** Count the number of integer solutions to  $Ax + By \leq C, 0 \leq x \leq n, 0 \leq y$ . In other words, evaluate the sum  $\sum_{x=0}^n \left\lfloor \frac{C-Ax}{B} + 1 \right\rfloor$ . To count all solutions, let  $n = \left\lfloor \frac{C}{a} \right\rfloor$ . In any case, it must hold that  $C - nA \geq 0$ . Be very careful about overflows.

8. MATH IV - NUMERICAL METHODS

8.1. **Fast Square Testing.** An optimized test for square integers.

```
long long M; -----
void init_is_square() { -----
- rep(i,0,64) M |= 1ULL << (63-(i*i)%64); } -----
inline bool is_square(ll x) { -----
- if (x == 0) return true; // XXX -----
- if ((M << x) >= 0) return false; -----
- int c = __builtin_ctz(x); -----
- if (c & 1) return false; -----
- x >>= c; -----
- if ((x&7) - 1) return false; -----
- ll r = sqrt(x); -----
- return r*r == x; } -----
```

8.2. **Simpson Integration.** Use to numerically calculate integrals

```
const int N = 1000 * 1000; // number of steps -----
double simpson_integration(double a, double b){ -----
- double h = (b - a) / N; -----
- double s = f(a) + f(b); // a = x_0 and b = x_2n -----
- for (int i = 1; i <= N - 1; ++i) { -----
-- double x = a + h * i; -----
-- s += f(x) * ((i & 1) ? 4 : 2); } -----
- s *= h / 3; -----
- return s; } -----
```

9. STRINGS

9.1. **Knuth-Morris-Pratt.** Count and find all matches of string  $f$  in string  $s$  in  $O(n)$  time.

```
int par[N]; // parent table -----
void buildKMP(string& f) { -----
- par[0] = -1, par[1] = 0; -----
- int i = 2, j = 0; -----
```

```
- while (i <= f.length()) { -----
-- if (f[i-1] == f[j]) par[i++] = ++j; -----
-- else if (j > 0) j = par[j]; -----
-- else par[i++] = 0; } } -----
std::vector<int> KMP(string& s, string& f) { -----
- buildKMP(f); // call once if f is the same -----
- int i = 0, j = 0; vector<int> ans; -----
- while (i + j < s.length()) { -----
-- if (s[i + j] == f[j]) { -----
--- if (++j == f.length()) { -----
---- ans.push_back(i); -----
--- i += j - par[j]; -----
--- if (j > 0) j = par[j]; } -----
-- } else { -----
--- i += j - par[j]; -----
--- if (j > 0) j = par[j]; } -----
- } return ans; } -----
```

9.2. Trie.

```
template <class T> -----
struct trie { -----
- struct node { -----
-- map<T, node*> children; -----
-- int prefixes, words; -----
-- node() { prefixes = words = 0; } }; -----
- node* root; -----
- trie() : root(new node()) { } -----
- template <class I> -----
- void insert(I begin, I end) { -----
-- node* cur = root; -----
-- while (true) { -----
--- cur->prefixes++; -----
--- if (begin == end) { cur->words++; break; } -----
--- else { -----
---- T head = *begin; -----
---- typename map<T, node*>::const_iterator it; -----
---- it = cur->children.find(head); -----
---- if (it == cur->children.end()) { -----
----- pair<T, node*> nw(head, new node()); -----
----- it = cur->children.insert(nw).first; -----
----- } begin++, cur = it->second; } } } -----
- template<class I> -----
- int countMatches(I begin, I end) { -----
-- node* cur = root; -----
-- while (true) { -----
--- if (begin == end) return cur->words; -----
--- else { -----
---- T head = *begin; -----
---- typename map<T, node*>::const_iterator it; -----
---- it = cur->children.find(head); -----
---- if (it == cur->children.end()) return 0; -----
---- begin++, cur = it->second; } } } -----
- template<class I> -----
- int countPrefixes(I begin, I end) { -----
-- node* cur = root; -----
-- while (true) { -----
```

```
- if (begin == end) return cur->prefixes; -----
- else { -----
-- T head = *begin; -----
-- typename map<T, node*>::const_iterator it; -----
-- it = cur->children.find(head); -----
-- if (it == cur->children.end()) return 0; -----
-- begin++, cur = it->second; } } } } -----
```

9.2.1. *Persistent Trie.*

```
const int MAX_KIDS = 2; -----
const char BASE = '0'; // 'a' or 'A' -----
struct trie { -----
- int val, cnt; -----
- std::vector<trie*> kids; -----
- trie () : val(-1), cnt(0), kids(MAX_KIDS, NULL) {} -----
- trie (int val) : val(val), cnt(0), kids(MAX_KIDS, NULL) {} -----
- trie (int val, int cnt, std::vector<trie*> &n_kids) : -----
-- val(val), cnt(cnt), kids(n_kids) {} -----
- trie *insert(std::string &s, int i, int n) { -----
-- trie *n_node = new trie(val, cnt+1, kids); -----
-- if (i == n) return n_node; -----
-- if (!n_node->kids[s[i]-BASE]) -----
--- n_node->kids[s[i]-BASE] = new trie(s[i]); -----
-- n_node->kids[s[i]-BASE] = -----
--- n_node->kids[s[i]-BASE]->insert(s, i+1, n); -----
-- return n_node; } }; -----
// max xor on a binary trie from version 'a+1' to 'b' (b > a):
int get_max_xor(trie *a, trie *b, int x) { -----
- int ans = 0; -----
- for (int i = MAX_BITS; i >= 0; --i) { -----
-- // don't flip the bit for min xor -----
-- int u = ((x & (1 << i)) > 0) ^ 1; -----
-- int res_cnt = (b and b->kids[u] ? b->kids[u]->cnt : 0) - -----
-- (a and a->kids[u] ? a->kids[u]->cnt : 0); -----
-- if (res_cnt == 0) u ^= 1; -----
-- ans ^= (u << i); -----
-- if (a) a = a->kids[u]; -----
-- if (b) b = b->kids[u]; } -----
- return ans; } -----
```

9.3. **Suffix Array.** Construct a sorted catalog of all substrings of  $s$  in  $O(n \log n)$  time using counting sort.

```
int n, equiv[N+1], suffix[N+1]; -----
ii equiv_pair[N+1]; -----
string T; -----
void make_suffix_array(string& s) { -----
- if (s.back()!='$') s += '$'; -----
- n = s.length(); -----
- for (int i = 0; i < n; i++) -----
-- suffix[i] = i; -----
- sort(suffix,suffix+n,[&s](int i, int j){return s[i] < s[j]}); -----
- int sz = 0; -----
- for(int i = 0; i < n; i++){ -----
-- if(i==0 || s[suffix[i]]!=s[suffix[i-1]]) -----
--- ++sz; -----
-- equiv[suffix[i]] = sz; } -----
- for (int t = 1; t < n; t<=1) { -----
```



```
for (int i = 0; i < n; i++)
    equiv_pair[i] = {equiv[i],equiv[(i+t)%n]};
sort(suffix, suffix+n, [](int i, int j) {
    return equiv_pair[i] < equiv_pair[j];});
int sz = 0;
for (int i = 0; i < n; i++) {
    if(i==0 || equiv_pair[suffix[i]]!=equiv_pair[suffix[i-1]])
        ++sz;
    equiv[suffix[i]] = sz; } }
int count_occurrences(string& G) { // in string T
    int L = 0, R = n-1;
    for (int i = 0; i < G.length(); i++){
        // lower/upper = first/last time G[i] is
        // the ith character in suffixes from [L,R]
        std::tie(L,R) = {lower(G[i],i,L,R), upper(G[i],i,L,R)};
        if (L==-1 && R==-1) return 0; }
    return R-L+1; }
```

9.4. **Longest Common Prefix**. Find the length of the longest common prefix for every substring in  $O(n)$ .

```
int lcp[N]; // lcp[i] = LCP(s[sa[i]:], s[sa[i+1]:])
void buildLCP(std::string s) { // build suffix array first
    for (int i = 0, k = 0; i < n; i++) {
        if (pos[i] != n - 1) {
            for(int j = sa[pos[i]+1]; s[i+k]==s[j+k];k++);
            lcp[pos[i]] = k; if (k > 0) k--;
        } else { lcp[pos[i]] = 0; } } }
```

9.5. **Aho-Corasick Trie**. Find all multiple pattern matches in  $O(n)$  time. This is KMP for multiple strings.

```
class Node {
    HashMap<Character, Node> next = new HashMap<>();
    Node fail = null;
    long count = 0;
    public void add(String s) { // adds string to trie
        Node node = this;
        for (char c : s.toCharArray()) {
            if (!node.contains(c))
                node.next.put(c, new Node());
            node = node.get(c);
        } node.count++;
        public void prepare() {
            // prepares fail links of Aho-Corasick Trie
            Node root = this; root.fail = null;
            Queue<Node> q = new ArrayDeque<Node>();
            for (Node child : next.values()) // BFS
                { child.fail = root; q.offer(child); }
            while (!q.isEmpty()) {
                Node head = q.poll();
                for (Character letter : head.next.keySet()) {
                    // traverse upwards to get nearest fail link
                    Node p = head;
                    Node nextNode = head.get(letter);
                    do { p = p.fail; }
                    while(p != root && !p.contains(letter));
                    if (p.contains(letter)) { // fail link found
```

```
        p = p.get(letter);
        nextNode.fail = p;
        nextNode.count += p.count;
        } else { nextNode.fail = root; }
        q.offer(nextNode); } } }
    public BigInteger search(String s) {
        // counts the words added in trie present in s
        Node root = this, p = this;
        BigInteger ans = BigInteger.ZERO;
        for (char c : s.toCharArray()) {
            while (p != root && !p.contains(c)) p = p.fail;
            if (p.contains(c)) {
                p = p.get(c);
                ans = ans.add(BigInteger.valueOf(p.count));
            } return ans; }
        private Node get(char c) { return next.get(c); }
        private boolean contains(char c) {
            return next.containsKey(c); }
        // Usage: Node trie = new Node();
        // for (String s : dictionary) trie.add(s);
        // trie.prepare(); BigInteger m = trie.search(str); }
```

9.6. Palindromes.

9.6.1. **Palindromic Tree**. Find lengths and frequencies of all palindromic substrings of a string in  $O(n)$  time.

Theorem: there can only be up to  $n$  unique palindromic substrings for any string.

```
int par[N*2+1], child[N*2+1][26];
int len[N*2+1], node[N*2+1], cs[N*2+1], size;
long long cnt[N+2]; // count can be very large
int newNode(int p = -1) {
    cnt[size] = 0; par[size] = p;
    len[size] = (p == -1 ? 0 : len[p] + 2);
    memset(child[size], -1, sizeof child[size]);
    return size++;
}
int get(int i, char c) {
    if (child[i][c] == -1) child[i][c] = newNode(i);
    return child[i][c];
}
void manachers(char s[]) {
    int n = strlen(s), cn = n * 2 + 1;
    for (int i = 0; i < n; i++) {
        cs[i * 2] = -1; cs[i * 2 + 1] = s[i];
        size = n * 2;
        int odd = newNode(), even = newNode();
        int cen = 0, rad = 0, L = 0, R = 0;
        size = 0; len[odd] = -1;
        for (int i = 0; i < cn; i++)
            node[i] = (i % 2 == 0 ? even : get(odd, cs[i]));
        for (int i = 1; i < cn; i++) {
            if (i > rad) { L = i - 1; R = i + 1; }
            else {
                int M = cen * 2 - i; // retrieve from mirror
                node[i] = node[M];
                if (len[node[M]] < rad - i) L = -1;
                else {
                    R = rad + 1; L = i * 2 - R;
```

```
                while (len[node[i]] > rad - i)
                    node[i] = par[node[i]]; } } // expand palindrome
            while (L >= 0 && R < cn && cs[L] == cs[R]) {
                if (cs[L] != -1) node[i] = get(node[i],cs[L]);
                L--, R++; }
            cnt[node[i]]++;
            if (i + len[node[i]] > rad) {
                rad = i + len[node[i]]; cen = i; }
            for (int i = size - 1; i >= 0; --i)
                cnt[par[i]] += cnt[i]; // update parent count
        int countUniquePalindromes(char s[]) {
            manachers(s); return size; }
        int countAllPalindromes(char s[]) {
            manachers(s); int total = 0;
            for (int i = 0; i < size; i++) total += cnt[i];
            return total; }
        // longest palindrome substring of s
        std::string longestPalindrome(char s[]) {
            manachers(s);
            int n = strlen(s), cn = n * 2 + 1, mx = 0;
            for (int i = 1; i < cn; i++)
                if (len[node[mx]] < len[node[i]])
                    mx = i;
            int pos = (mx - len[node[mx]]) / 2;
            return std::string(s + pos, s + pos + len[node[mx]]); }
```

9.6.2. Eertree.

```
struct node {
    int start, end, len, back_edge, *adj;
    node() {
        adj = new int[26];
        for (int i = 0; i < 26; ++i) adj[i] = 0; }
    node(int start, int end, int len, int back_edge) :
        start(start), end(end), len(len), back_edge(back_edge) {
        adj = new int[26];
        for (int i = 0; i < 26; ++i) adj[i] = 0; } };
struct eertree {
    int ptr, cur_node;
    std::vector<node> tree;
    eertree() {
        tree.push_back(node());
        tree.push_back(node(0, 0, -1, 1));
        tree.push_back(node(0, 0, 0, 1));
        cur_node = 1;
        ptr = 2; }
    int get_link(int temp, std::string &s, int i) {
        while (true) {
            int cur_len = tree[temp].len;
            // don't return immediately if you want to
            // get all palindromes; not recommended
            if (i-cur_len-1 >= 0 and s[i] == s[i-cur_len-1])
                return temp;
            temp = tree[temp].back_edge; }
        return temp; }
    void insert(std::string &s, int i) {
        int temp = cur_node;
```

```
--- temp = get_link(temp, s, i); -----
--- if (tree[temp].adj[s[i] - 'a'] != 0) { -----
----- cur_node = tree[temp].adj[s[i] - 'a']; -----
----- return; } -----
--- ptr++; -----
--- tree[temp].adj[s[i] - 'a'] = ptr; -----
--- int len = tree[temp].len + 2; -----
--- tree.push_back(node(i-len+1, i, len, 0)); -----
--- temp = tree[temp].back_edge; -----
--- cur_node = ptr; -----
--- if (tree[cur_node].len == 1) { -----
----- tree[cur_node].back_edge = 2; -----
----- return; } -----
--- temp = get_link(temp, s, i); -----
--- tree[cur_node].back_edge = tree[temp].adj[s[i] - 'a']; } ---
- void insert(std::string &s) { -----
--- for (int i = 0; i < s.size(); ++i) -----
----- insert(s, i); } };
```

9.7. **Z Algorithm**. Find the longest common prefix of all substrings of  $s$  with itself in  $O(n)$  time.

```
int z[N]; // z[i] = lcp(s, s[i:]) -----
void computeZ(string s) { -----
- int n = s.length(), L = 0, R = 0; z[0] = n; -----
- for (int i = 1; i < n; i++) { -----
--- if (i > R) { -----
----- L = R = i; -----
----- while (R < n && s[R - L] == s[R]) R++; -----
----- z[i] = R - L; R--; -----
--- } else { -----
----- int k = i - L; -----
----- if (z[k] < R - i + 1) z[i] = z[k]; -----
----- else { -----
----- L = i; -----
----- while (R < n && s[R - L] == s[R]) R++; -----
----- z[i] = R - L; R--; } } } }
```

9.8. **Booth's Minimum String Rotation**. Booth's Algo: Find the index of the lexicographically least string rotation in  $O(n)$  time.

```
int f[N * 2]; -----
int booth(string S) { -----
- S.append(S); // concatenate itself -----
- int n = S.length(), i, j, k = 0; -----
- memset(f, -1, sizeof(int) * n); -----
- for (j = 1; j < n; j++) { -----
--- i = f[j-k-1]; -----
--- while (i != -1 && S[j] != S[k + i + 1]) { -----
----- if (S[j] < S[k + i + 1]) k = j - i - 1; -----
----- i = f[i]; -----
--- } if (i == -1 && S[j] != S[k + i + 1]) { -----
----- if (S[j] < S[k + i + 1]) k = j; -----
----- f[j - k] = -1; -----
--- } else f[j - k] = i + 1; -----
- } return k; }
```

9.9. Hashing.

```
9.9.1. Rolling Hash.
int MAXN = 1e5+1, MOD = 1e9+7; -----
struct hasher { -----
- int n; -----
- std::vector<ll> *p_pow, *h_ans; -----
- hash(vi &s, vi primes) : n(primes.size()) { -----
--- p_pow = new std::vector<ll>[n]; -----
--- h_ans = new std::vector<ll>[n]; -----
--- for (int i = 0; i < n; ++i) { -----
----- p_pow[i] = std::vector<ll>(MAXN); -----
----- p_pow[i][0] = 1; -----
----- for (int j = 0; j+1 < MAXN; ++j) -----
----- p_pow[i][j+1] = (p_pow[i][j] * primes[i]) % MOD; -----
----- h_ans[i] = std::vector<ll>(MAXN); -----
----- h_ans[i][0] = 0; -----
----- for (int j = 0; j < s.size(); ++j) -----
----- h_ans[i][j+1] = (h_ans[i][j] + -----
----- s[j] * p_pow[i][j]) % MOD; } } };
```