

Multi-Level Bin Packing Problem

CSE3000 Research Project Q4 2021/2022

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Multi-Level Bin Packing Problem

Problem Description

Given

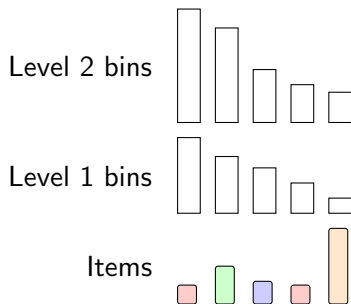
- set of items $\mathcal{I} = \{\mathcal{I}_1, \mathcal{I}_2, \dots, \mathcal{I}_{n^0}\}$
 - with size $s(\mathcal{I}_j) \in \mathbb{N}_{>0}$, $1 \leq j \leq n^0$
- set of bins $\mathcal{B}^i = \{\mathcal{B}_1^i, \mathcal{B}_2^i, \dots, \mathcal{B}_{n^i}^i\}$, $1 \leq i \leq m$, m levels
 - with size $s(\mathcal{B}_k^i) \in \mathbb{N}_{>0}$, capacity $w(\mathcal{B}_k^i) \in \mathbb{N}_{>0}$, and cost $c(\mathcal{B}_k^i) \in \mathbb{N}_{>0}$
 - $1 \leq k \leq n^i$, $1 \leq i \leq m$

Task

- insert each item into a bin of level 1
- insert each used bin of level i , $1 \leq i < m$ into a bin of level $i + 1$
- the capacity of a bin must not exceeded
- minimise total cost of used bins

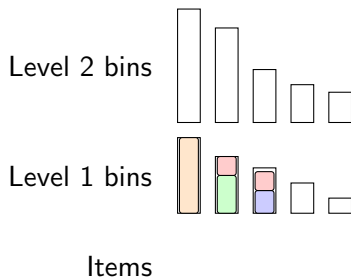
Multi-Level Bin Packing Problem

Example with $m = 2$ levels



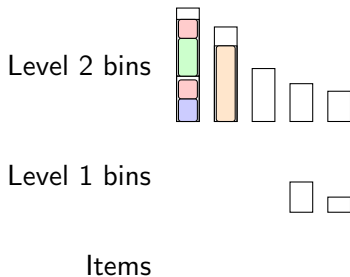
Multi-Level Bin Packing Problem

Example with $m = 2$ levels



Multi-Level Bin Packing Problem

Example with $m = 2$ levels



Class Constrained MLBP (CCMLBP)

In addition to MLBP

- Each item $\mathcal{I}_j \in \mathcal{I}$ belongs to a class $\kappa(\mathcal{I}_j) \in \{1, 2, \dots, q\}$ where q is a positive integer.
- Each level $i \in M$ is associated with a bound $Q^i \in \mathbb{N}_{\geq 0}$ for the number of different classes within a bin $\mathcal{B}_k^i \in \mathcal{B}^i$.

Task

Find a packing that minimises the total cost of the used bins such that every bin at level $i \in M$ satisfies the **capacity constraints** and contains items from no more than Q^i classes.

MLBP with Conflict Constraints (MLBPCC)

In addition to MLBP

- Given: a conflict graph $G = (\mathcal{I}, E)$ where each node corresponds to an item
- An edge $(\mathcal{I}_j, \mathcal{I}_q) \in E$ indicates a **conflict** between items \mathcal{I}_j and \mathcal{I}_q

Task

Find a packing that minimises the total cost of the used bins such that every bin satisfies the **capacity constraints** and each bin **contains no items that are in conflict with each other**.

MLBP with Partial Orders (MLBPPO)

In addition to MLBP

- Given: precedence constraints, denoted by \prec , among items \mathcal{I}
- Precedence refers to the relative ordering of top level bins in \mathcal{B}^m
- if there is a precedence $\mathcal{I}_j \prec \mathcal{I}_q$ between items $\mathcal{I}_j, \mathcal{I}_q \in \mathcal{I}$
 - then if \mathcal{I}_j and \mathcal{I}_q are packed into \mathcal{B}_k^m and \mathcal{B}_l^m , respectively $\Rightarrow k < l$

Task

Find a packing that minimises the total cost of the used bins such that every bin satisfies the **capacity constraints** and all **items satisfy their precedence constraints**.

MLBP with Time Windows (MLBPTW)

In addition to MLBP

- Given: Time window (TW) $t_j = [e_j, l_j]$, $\mathcal{I}_j \in \mathcal{I}$ and penalty factor p
- Item \mathcal{I}_j must be **packed as early as possible** within its earliest packing time e_j and its latest packing time l_j
- Items can only be packed into the same bin if their TWs overlap

Task

- let $u(\mathcal{I}_j)$ be the earliest time item \mathcal{I}_j is packed into a bin
- in addition to the total bin costs, a penalty value $(u(\mathcal{I}_j) - e_j) p$ is added for each item \mathcal{I}_j

Find a packing that minimises the changed objective function such that every bin satisfies the **capacity constraints** and contains **no items with non-overlapping time windows**.

MLBP with Fragmentation Constraints (MLBPFC)

In addition to MLBP

- Each item $\mathcal{I}_j \in \mathcal{I}$ belongs to a group $g(\mathcal{I}_j) \in G = \{1, 2, \dots, q\}$, $q \in \mathbb{N}_{>0}$
- Items of the same group should be packed into the same top level bin \mathcal{B}^m

Task

- let n_r be the number of top level bins that contains items of group $r \in G$
- in addition to the total bin costs, a penalty value $(n_r - 1)p$ is added for each group

Find a packing that minimises the changed objective function such that every bin satisfies the [capacity constraints](#).

C++ Framework

- Provides
 - C++ classes to read instances
 - simple command line argument parser
 - MIP solver class
 - uses CPLEX to solve MIPs
 - add your decision variables, constraints, and objective function
 - sample case for the standard bin packing problem
- It is not mandatory to use the C++ Framework!
 - You can use other MIP solvers and/or other programming languages.

- download and install CPLEX Optimization Studio
 - <https://www.ibm.com/academic/topic/data-science>
 - you will find Optimization Studio under software
 - free for students, but you have to register
- set system environment variable CPLEXDIR=<directory of cplex>
 - e.g., export CPLEXDIR=/opt/ibm/ILOG/CPLEX_Studio201
- run make

C++ Framework

Setup Framework - Windows

- download and install CPLEX Optimization Studio
 - <https://www.ibm.com/academic/topic/data-science>
 - you will find Optimization Studio under software
 - free for students, but you have to register
- download and install Visual Studio with C++
 - <https://visualstudio.microsoft.com/vs/features/cplusplus/>

- click on the project name in solution explorer tab and select properties
- under C/C++ → general → additional include directories → add:
 - <cplex directory>\CPLEX_StudioXXX\concert\include
 - <cplex directory>\CPLEX_StudioXXX\cplex\include
- under C/C++ → preprocessors → preprocessor definitions → add:
 - NDEBUG
 - _CONSOLE
 - IL_STD

- under C/C++ → code generation → runtime library
 - set to “multi-threaded DLL (/MD)”
- under Linker → Input → additional dependencies → add:
 - <cplex directory>\CPLEX_StudioXXX\cplex\lib\x64_windows_vsXXXX\stat_mda\cplexXXX.lib
 - <cplex directory>\CPLEX_StudioXXX\cplex\lib\x64_windows_vsXXXX\stat_mda\ilocplex.lib
 - <cplex directory>\CPLEX_StudioXXX\concert\lib\x64_windows_vsXXXX\stat_mda\concert.lib
- under general → C++ language standard
 - set to “ISO C++ 17 Standard” or to “ISO C++ 20 Standard”