



**SRM Institute of Science and Technology**  
**College of Engineering and Technology**  
**SCHOOL OF COMPUTING**

SRM Nagar, Kattankulathur – 603203, Chengalpattu District, Tamilnadu

**Academic Year: 2022-23 (EVEN)**

**SET-B**

**Test: CLAT-2**

**Course Code & Title: 18CSC304J -COMPILER DESIGN**

**Year & Sem: III & VI**

**Date : 04.04.2023**

**Duration : 2 Periods**

**Max. Marks: 50**

**Course Articulation Matrix:**

S.No.	Course Outcome	PO1	PO2	PO3	PO4	PO5	PO6	PO7	PO8	PO9	PO10	PO11	PO12
1	CO3	3	3	3									

**Part – A (10 X 1 = 10 Marks)**

**Answer ALL Questions**

Q. No	Question	Marks	BL	CO	PO	PI Code
1	The grammar $A \rightarrow Ax \mid (A) \mid \epsilon$ is not suitable for predictive-parsing because the grammar is? a) Left factoring <b>b) Left recursive</b> c) Right recursive d) An operator grammar  <b>Answer: B</b>	1	1	3	1	1.7.1
2	For the grammar, $E \rightarrow EE \mid (E) \mid \epsilon$ , number of parse trees to produce empty string is? a) One b) Two c) Three <b>d) Infinite</b>  <b>Answer: D</b>	1	1	3	1	1.7.1
3	Which grammar rules violate the requirements of an operator grammar? 1. $E \rightarrow FG$ 2. $F \rightarrow E s F$ 3. $G \rightarrow F t H p$ 4. $H \rightarrow \epsilon$  a) 1 only b) 1 and 3 only <b>c) 1 and 4 only</b> d) 1, 3 and 4 only  <b>Answer: C</b>	1	2	3	2	2.8.2
4	A form of recursive descent parsing that does not require any back-tracking is known as? a) recursive parsing b) non-recursive parsing <b>c) predictive parsing</b> d) non-predictive parsing  <b>Answer: C</b>	1	1	3	1	1.7.1

5	<p>For the grammar given below, find FIRST(X)</p> $X \rightarrow Ya \mid bZ$ $Y \rightarrow c \mid \varepsilon$ <p>a) {a, b}  b) {c, <math>\varepsilon</math>}  c) {a, b, c}  <b>d) {a, b, c, <math>\varepsilon</math>}</b></p> <p><b>Answer: D</b></p>	1	1	3	1	1.7.1
6	<p>In shift reduce parsing handle is at _____</p> <p><b>a) Top of the stack</b>  b) Bottom of the stack  c) Anywhere in the stack  d) Nowhere in the stack</p> <p><b>Answer: A</b></p>	1	2	3	2	2.8.2
7	<p>Choose the correct precedence relations in operator precedence parsing:  if operator <math>O_1</math> has higher precedence than operator <math>O_2</math></p> <p>a) <math>O_1 \succ O_2</math>  b) <math>O_1 = O_2</math>  c) <math>O_2 \prec O_1</math>  <b>d) <math>O_1 \succ O_2</math> and <math>O_2 \prec O_1</math></b></p> <p><b>Answer: D</b></p>	1	2	3	4	4.6.2
8	<p>In SLR, CLR and LALR parser, which have same number of states?</p> <p>a) SLR and CLR  <b>b) SLR and LALR</b>  c) CLR and LALR  d) SLR, CLR and LALR</p> <p><b>Answer: B</b></p>	1	2	3	1	1.7.1
9	<p>What is the LEADING(X) for the following grammar?</p> $X \rightarrow X - B \mid B$ $B \rightarrow B * A \mid A$ $A \rightarrow (X) \mid id$ <p>a) LEADING(X)={-,*,(,)}  <b>b) LEADING(X)={-,*,),id}</b>  c) LEADING(X)={-,*,(,id}  d) LEADING(X)={-,*,(}</p> <p><b>Answer: B</b></p>	1	2	3	4	4.6.2
10	<p>Construction of parsing table in which strategies do not need the FOLLOW set?</p> <p>a) SLR and CLR  b) SLR and LALR  <b>c) CLR and LALR</b>  d) SLR, CLR and LALR</p> <p><b>Answer: C</b></p>	1	1	3	3	3.8.2



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**SET-B**

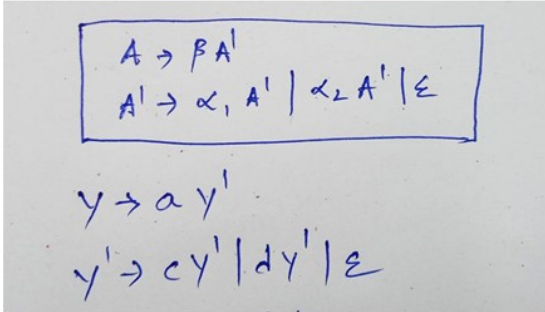
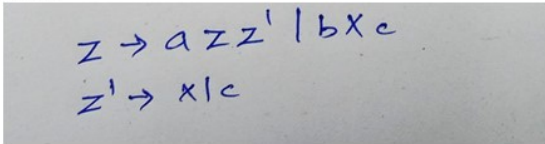
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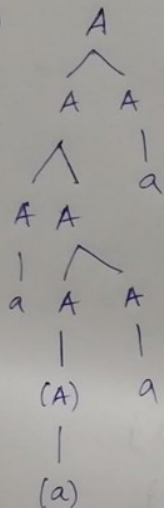
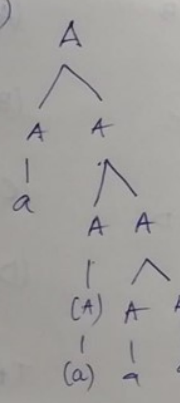
Date : 04.04.2023  
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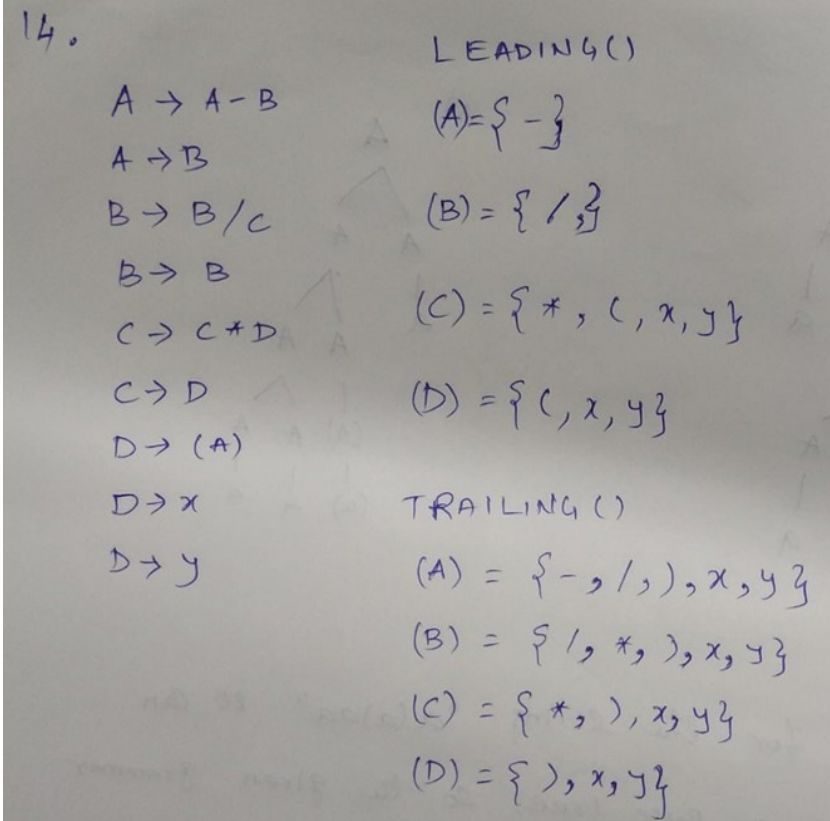
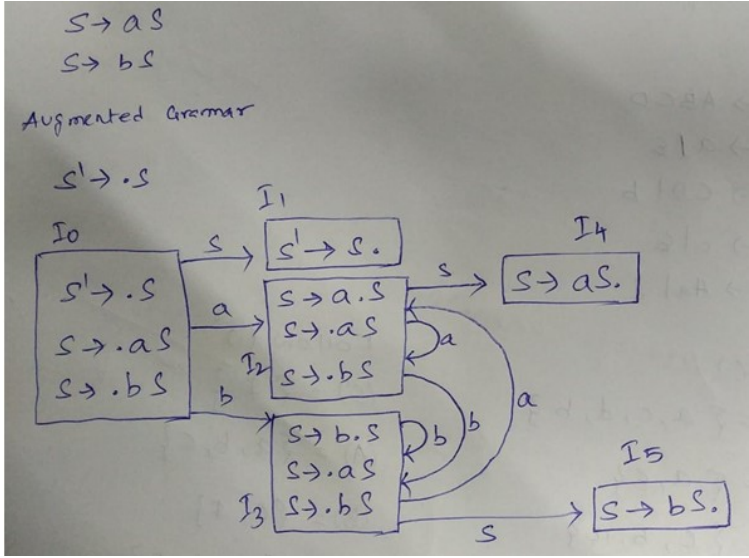
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1	CO3	3	3	3									

**PART B (4 x 4 = 16) ANSWER ANY FOUR**

11	<p>Eliminate left recursion and left factoring in the following grammar:</p> $X \rightarrow Ya \mid b \mid c$ $Y \rightarrow Yc \mid Yd \mid a$ $Z \rightarrow aZX \mid bXc \mid aZc$ <p><b>Left recursion – 2 marks</b></p>  <p><b>Left factoring – 2 marks</b></p> 	4	2	3	3	3.8.2
12	<p>Check the following grammar is ambiguous or not by parsing the input string “a(a)aa”:</p> $A \rightarrow AA$ $A \rightarrow (A)$ $A \rightarrow a$	4	2	3	4	4.6.2

	<p>12.</p> <div style="display: flex; justify-content: space-around;"> <div style="text-align: center;"> <p>a)</p>  </div> <div style="text-align: center;"> <p>b)</p>  </div> </div> <p>for the string "a(a)aa" it can generate 2 parse tree, so the given grammar is Ambiguous.</p>					
13	<p>Compute FIRST() and FOLLOW() for the grammar:</p> <p><math>S \rightarrow ABCD</math>  <math>A \rightarrow a \mid \epsilon</math>  <math>B \rightarrow CD \mid b</math>  <math>C \rightarrow c \mid \epsilon</math>  <math>D \rightarrow Aa \mid d</math></p> <p>Compute FIRST() – 2 marks          Compute FOLLOW() – 2 marks</p> <div style="border: 1px solid black; padding: 10px; margin-top: 10px;"> <p>13. <math>S \rightarrow ABCD</math>  <math>A \rightarrow a \mid \epsilon</math>  <math>B \rightarrow CD \mid b</math>  <math>C \rightarrow c \mid \epsilon</math>  <math>D \rightarrow Aa \mid d</math></p> <div style="display: flex; justify-content: space-between;"> <div style="width: 45%;"> <p>FIRST()</p> <p><math>(S) = \{a, c, d, b\}</math>  <math>(A) = \{a, \epsilon\}</math>  <math>(B) = \{c, b, a\}</math>  <math>(C) = \{c, \epsilon\}</math>  <math>(D) = \{a, d, \epsilon\}</math></p> </div> <div style="width: 45%;"> <p>FOLLOW()</p> <p><math>(S) = \{\\$ \}</math>  <math>(A) = \{a, b, c\}</math>  <math>(B) = \{c, \\$ \}</math>  <math>(C) = \{a, d, \\$ \}</math>  <math>(D) = \{c, \\$ \}</math></p> </div> </div> </div>	4	3	3	3	3.8.2
14	<p>Find LEADING() and TRAILING() for all the non-terminals in the following grammar:</p> <p><math>A \rightarrow A - B \mid B</math></p>	4	3	3	3	3.8.2

	$B \rightarrow B / C \mid B$ $C \rightarrow C * D \mid D$ $D \rightarrow (A) \mid x \mid y$  Compute LEADING() – 2 marks Compute TRAILING() – 2 marks  					
15	Find the canonical collection of LR(0) items for the following grammar: $S \rightarrow aS \mid bS$  	4	3	3	2	2.6.4
Part – C (2 x 12 = 24 Marks) Answer ALL Questions						
16.	Consider the grammar: $A \rightarrow pqC \mid pBs \mid pAD$ $B \rightarrow qB \mid \epsilon$ $C \rightarrow s \mid \epsilon$	12	2	3	3	3.8.2

$D \rightarrow p | q | \epsilon$

Check whether the following inputs are accepted or not by the grammar using recursive decent parsing:

i) pqqp - 6 marks

ii) ppqqss - 6 marks

1b

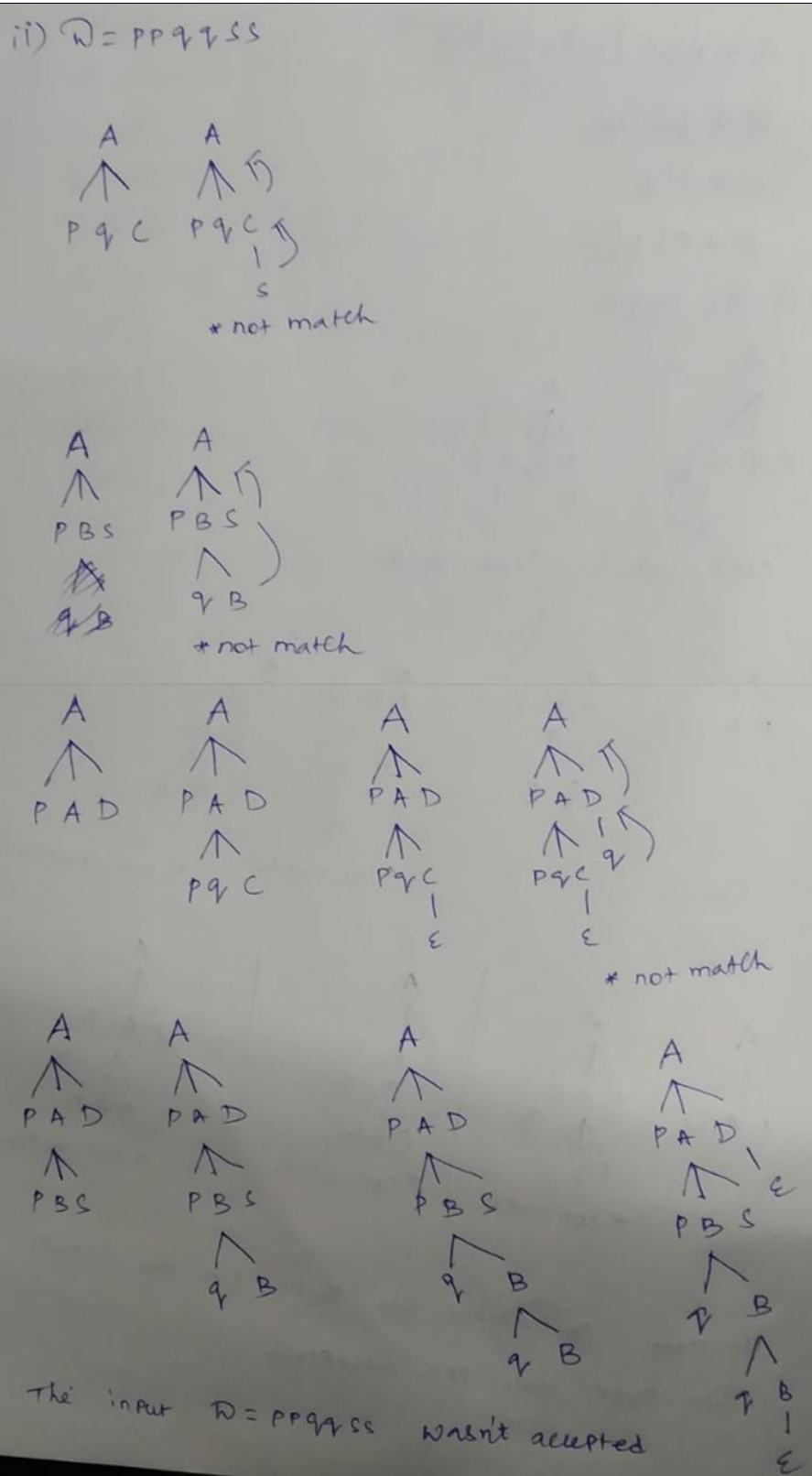
$$A \rightarrow p q C \mid p B s \mid p A D$$
$$B \rightarrow q B \mid \epsilon$$
$$C \rightarrow s \mid \epsilon$$
$$D \rightarrow p \mid q \mid \epsilon$$

i)  $w = p q q p$

Parse trees for  $w = p q q p$ :

- Tree 1:  $A \rightarrow p q C$ ,  $C \rightarrow s$ . \*not match
- Tree 2:  $A \rightarrow p q C$ ,  $C \rightarrow \epsilon$ . \*not match
- Tree 3:  $A \rightarrow p B s$ ,  $B \rightarrow q B$ ,  $B \rightarrow \epsilon$ . \*not match
- Tree 4:  $A \rightarrow p B s$ ,  $B \rightarrow q B$ ,  $B \rightarrow \epsilon$ . \*not match
- Tree 5:  $A \rightarrow p A D$ ,  $A \rightarrow p q C$ . \*not match
- Tree 6:  $A \rightarrow p A D$ ,  $A \rightarrow p B s$ . \*not match

No more productions are there, so the input  $w = p q q p$  was not accepted



(OR)

17.	Election commission has announced the MLA election for Kanchipuram constituency. In view of this, applications are invited for the MLA election nomination. A candidate should produce proof for Age, Qualification, and any Work experience. The basic criteria for age limit is Age>20 and Age<50. The academic qualification can be UG or PG or Diploma or no qualification. Then, it includes whether the candidate has a work experience or not. Construct CFG for	12	3	3	4	4.6.2
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	<p>the given scenario and parse the below input strings using Shift Reduce parsing.</p> <p>Input:</p> <p>i) 31 dip no</p> <p>ii) 20 ug yes</p> <p><b>Grammar:</b> - 6 marks</p> <p><math>M \rightarrow AQE</math></p> <p><math>A \rightarrow 2X \mid 3Y \mid 4Y</math></p> <p><math>X \rightarrow 1 \mid 2 \mid 3 \mid \dots \mid 9</math></p> <p><math>Y \rightarrow 0 \mid X</math></p> <p><math>Q \rightarrow ug \mid pg \mid dip \mid \epsilon</math></p> <p><math>E \rightarrow yes \mid no</math></p> <p><b>Parsing the inputs:</b></p> <p>i) 31 dip no - 3 marks</p> <table><tr><th>Stack</th><th>Input</th><th>Action</th></tr><tr><td>\$</td><td>31 dip no \$</td><td>Shift</td></tr><tr><td>\$ 3</td><td>1 dip no \$</td><td>Shift</td></tr><tr><td>\$ 31</td><td>dip no \$</td><td>Reduce <math>X \rightarrow 1</math></td></tr><tr><td>\$ 3X</td><td>dip no \$</td><td>Reduce <math>Y \rightarrow X</math></td></tr><tr><td>\$ 3Y</td><td>dip no \$</td><td>Reduce <math>A \rightarrow 3Y</math></td></tr><tr><td>\$ A</td><td>dip no \$</td><td>Shift</td></tr><tr><td>\$ A dip</td><td>no \$</td><td>Reduce <math>Q \rightarrow dip</math></td></tr><tr><td>\$ AQ</td><td>no \$</td><td>Shift</td></tr><tr><td>\$ AQ no</td><td>\$</td><td>Reduce <math>E \rightarrow no</math></td></tr><tr><td>\$ AQE</td><td>\$</td><td>Reduce <math>M \rightarrow AQE</math></td></tr><tr><td>\$ M</td><td>\$</td><td>Accept</td></tr></table> <p>ii) 20 ug yes - 3 marks</p> <table><tr><th>Stack</th><th>Input</th><th>Action</th></tr><tr><td>\$</td><td>20 ug yes \$</td><td>Shift</td></tr><tr><td>\$ 2</td><td>0 ug yes \$</td><td>Shift</td></tr><tr><td>\$ 20</td><td>ug yes \$</td><td>Reduce <math>Y \rightarrow 0</math></td></tr><tr><td>\$ 2Y</td><td>ug yes \$</td><td>Shift</td></tr><tr><td>\$ 2Y ug</td><td>yes \$</td><td>Reduce <math>Q \rightarrow ug</math></td></tr><tr><td>\$ 2YQ</td><td>yes \$</td><td>Shift</td></tr><tr><td>\$ 2YQ yes</td><td>\$</td><td>Reduce <math>E \rightarrow yes</math></td></tr><tr><td>\$ 2YQE</td><td>\$</td><td>Not Accepted</td></tr></table>	Stack	Input	Action	\$	31 dip no \$	Shift	\$ 3	1 dip no \$	Shift	\$ 31	dip no \$	Reduce $X \rightarrow 1$	\$ 3X	dip no \$	Reduce $Y \rightarrow X$	\$ 3Y	dip no \$	Reduce $A \rightarrow 3Y$	\$ A	dip no \$	Shift	\$ A dip	no \$	Reduce $Q \rightarrow dip$	\$ AQ	no \$	Shift	\$ AQ no	\$	Reduce $E \rightarrow no$	\$ AQE	\$	Reduce $M \rightarrow AQE$	\$ M	\$	Accept	Stack	Input	Action	\$	20 ug yes \$	Shift	\$ 2	0 ug yes \$	Shift	\$ 20	ug yes \$	Reduce $Y \rightarrow 0$	\$ 2Y	ug yes \$	Shift	\$ 2Y ug	yes \$	Reduce $Q \rightarrow ug$	\$ 2YQ	yes \$	Shift	\$ 2YQ yes	\$	Reduce $E \rightarrow yes$	\$ 2YQE	\$	Not Accepted					
Stack	Input	Action																																																																			
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18.	<p>Consider the following grammar:</p> <p><math>S \rightarrow (L) \mid a</math></p> <p><math>L \rightarrow L,S \mid S</math></p> <p>Construct operator precedence parsing table using Leading and Trailing, precedence graph and precedence function</p>	12	2	3	3	3.8.2																																																															



18.

$S \rightarrow (L)$

$S \rightarrow a$

$L \rightarrow L, S$

$L \rightarrow S$

LEADING()

$(S) = \{ (, a \}$

$(L) = \{ , (, a \}$

TRAILING()

$(S) = \{ ), a \}$

$(L) = \{ , , ), a \}$

Three types of precedence =  $\{ <, >, = \}$

Rule ①

set  $\$ < a$

set  $b > \$$

$\$ < \text{Lead}(S)$

$\$ < \{ (, a \}$

; for all 'a'  $\text{Lead}(S)$

; for all 'b'  $\text{Trail}(S)$

$\text{Trail}(S) > \$$

$\{ ), a \} > \$$

Rule ②

$A \rightarrow xy$  ; if  $xy$  are terminal set  $x=y$

NOT APPLICABLE

Rule ③

$A \rightarrow xBy$  ; if 'xy' are terminal 'B' is nonTerminal  
set  $x=y$

$S \rightarrow (L)$

$( = )$

Rule ④

$A \rightarrow \alpha x B \beta$  ; if 'x' is terminal and 'B' non Terminal  
set  $x < \text{Lead}(B)$

$S \rightarrow \frac{(}{x} \frac{L}{B})$

$L \rightarrow L, \frac{S}{x B}$

$< \{ , , (, a \}$

$> \{ (, a \}$

Rule ⑤

$A \rightarrow \alpha B x \beta$  ; if 'x' is terminal and 'B' non Terminal  
set  $\text{Trail}(B) > x$

$S \rightarrow \frac{(}{B} \frac{L}{x})$

$L \rightarrow \frac{L}{B} , \frac{S}{x}$

$\{ , , , a \} > )$

$\{ , , , a \} > ,$

Rule ⑥

$\$, \$ = \text{Accepted}$

	a	,	(	)	\$
a		>		>	>
,	<	>	<	>	
(	<	<	<	=	
)		>		>	>
\$	<		<		Accepted

(OR)

19.

Consider the following grammar:

$S \rightarrow L = R \mid R$

$L \rightarrow *R \mid \text{id}$

$R \rightarrow L$

Construct the SLR parser table for the grammar. Show action of the parser, for the input string **\*id=id**

12

3

3

3

3.8.2

**Given Grammar:** $S \rightarrow L = R$  $S \rightarrow R$  $L \rightarrow *R$  $L \rightarrow \text{id}$  $R \rightarrow L$ **Step-1: Augmented grammar - 1 mark** $S' \rightarrow S$  $S \rightarrow L = R$  $S \rightarrow R$  $L \rightarrow *R$  $L \rightarrow \text{id}$  $R \rightarrow L$ **Step-2: Find LR(0) collections - 3 marks** $I_0:$  $S' \rightarrow .S$  $S \rightarrow .L = R$  $S \rightarrow .R$  $L \rightarrow .*R$  $L \rightarrow .\text{id}$  $R \rightarrow .L$  $I_1: \text{goto}(I_0, S)$  $S' \rightarrow S.$  $I_2: \text{goto}(I_0, L)$  $S \rightarrow L. = R$  $R \rightarrow L.$  $I_3: \text{goto}(I_0, R)$  $S \rightarrow R.$  $I_4: \text{goto}(I_0, *)$  $L \rightarrow *.R$  $R \rightarrow .L$  $L \rightarrow .*R$  $L \rightarrow .\text{id}$  $I_5: \text{goto}(I_0, \text{id})$  $L \rightarrow \text{id}.$  $I_6: \text{goto}(I_2, =)$  $S \rightarrow L =. R$  $R \rightarrow .L$  $L \rightarrow .*R$  $L \rightarrow .\text{id}$  $I_7: \text{goto}(I_4, R)$  $L \rightarrow *R.$  $I_8: \text{goto}(I_4, L)$  $R \rightarrow L.$

I<sub>4</sub>: goto(I<sub>4</sub>, \*)  
 $L \rightarrow *.R$   
 $R \rightarrow .L$   
 $L \rightarrow .*R$   
 $L \rightarrow .id$

I<sub>5</sub>: goto(I<sub>4</sub>, id)  
 $L \rightarrow id.$

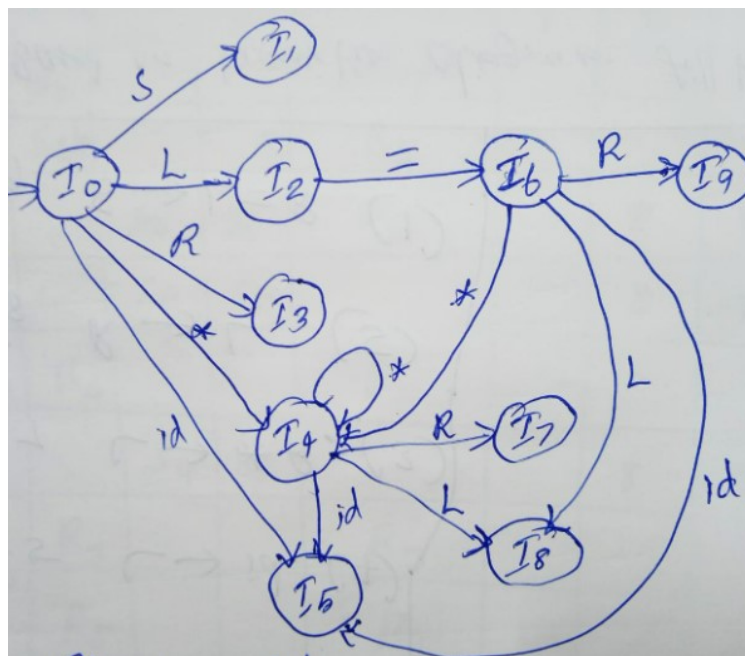
I<sub>9</sub>: goto(I<sub>6</sub>, R)  
 $S \rightarrow L=R.$

I<sub>8</sub>: goto(I<sub>6</sub>, L)  
 $R \rightarrow L.$

I<sub>4</sub>: goto(I<sub>6</sub>, \*)  
 $L \rightarrow *.R$   
 $R \rightarrow .L$   
 $L \rightarrow .*R$   
 $L \rightarrow .id$

I<sub>5</sub>: goto(I<sub>6</sub>, id)  
 $L \rightarrow id.$

**Transition Diagram of goto function - 2 marks**



**Step-3: Find FOLLOW() - 1 marks**

$\text{FOLLOW}(S') = \{ \$ \}$

$\text{FOLLOW}(S) = \{ \$ \}$

$\text{FOLLOW}(L) = \{ =, \text{FOLLOW}(R) \} = \{ =, \$ \}$

$\text{FOLLOW}(R) = \{ \text{FOLLOW}(S), \text{FOLLOW}(L) \} = \{ \$, = \}$

**Step-4: Construct the parsing table - 3 marks**

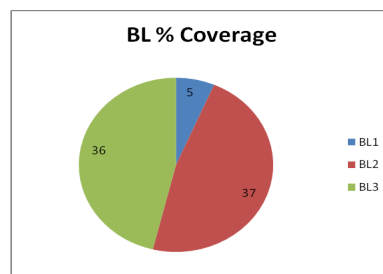
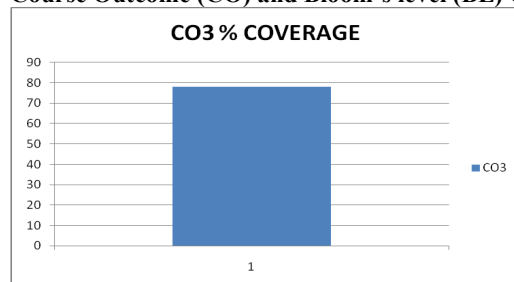
	=	*	id	\$		S	L	R
0		S4	S5			1	2	3
1				Accept				
2	S6/R5			R5				
3				R2				
4		S4	S5				8	7
5	R4			R4				
6		S4	S5				8	9
7	R3			R3				
8	R5			R5				
9				R1				

**Parsing the input: \*id = id - 2 marks**

Stack	Input	Action
0	*id = id \$	Shift 4
0 * 4	id = id \$	Shift 5
0 * 4 id 5	= id \$	Reduce by L → id
0 * 4 L 8	= id \$	Reduce by R → L
0 * 4 R 7	= id \$	Reduce by L → * R
0 L 2	= id \$	Shift 6
0 L 2 = 6	id \$	Shift 5
0 L 2 = 6 id 5	\$	Reduce by L → id
0 L 2 = 6 L 8	\$	Reduce by R → L
0 L 2 = 6 R 9	\$	Reduce by S → L = R
0 S 1	\$	Accept

\*Performance Indicators are available separately for Computer Science and Engineering in AICTE examination reforms policy

**Course Outcome (CO) and Bloom's level (BL) Coverage in Questions**



Approved by the Audit Professor/Course Coordinator