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1 Templates

1.1 Start

```
// .vimrc
syn on
set mouse=a sw=4 ts=4 ai si nu wrap
nnoremap;:

// Terminal: comparing generated output to sample output
./my_program < sample.in | diff sample.out -
```

1.2 Template - C++

```
#include<bits/stdc++.h>
using namespace std;
typedef long long l1;
static bool DBG = 1;

ll mod(ll a, ll b) { return ((a%b)+b)%b; }

int main() {
   ios_base::sync_with_stdio(0);
   cout << fixed << setprecision(15);
   int n;
   cin >> n;
   cout << n << endl;
   return 0;
}</pre>
```

1.3 Template - Java

```
import java.util.*;
import java.math.*;
import java.io.*;

class modelo {
    static final double EPS = 1.e-10;
    static final boolean DBG = true;

    private static int cmp(double x, double y = 0, double tol = EPS) {
        return (x <= y + tol)? (x + tol < y)? -1 : 0 : 1;
    }

    public static void main(String[] argv) {
        Scanner s = new Scanner(System.in);
    }
}</pre>
```

2 Data Structures

2.1 BIT

```
template<typename T> struct BIT{
   int S;
   vector<T> v;
   BIT<T>(int _S){
       S = _S;
       v.resize(S+1):
   void update(int i, T k){
       for(i++; i<=S; i+=i&-i)</pre>
          v[i] = v[i] + k;
   T read(int i){
       T sum = 0;
       for(i++; i; i-=i&-i)
          sum = sum + v[i];
       return sum;
   T read(int 1, int r){
       return read(r) - read(l-1);
```

2.2 KD Tree

```
// A straightforward, but probably sub-optimal KD-tree implmentation
// that's probably good enough for most things (current it's a
// 2D-tree)
// - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
// distributed
// - worst case for nearest-neighbor may be linear in pathological
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
   ntype x, y;
   point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
```

```
bool operator==(const point &a, const point &b) {
   return a.x == b.x && a.y == b.y;
// sorts points on x-coordinate
bool on_x(const point &a, const point &b) {
   return a.x < b.x;</pre>
// sorts points on y-coordinate
bool on_y(const point &a, const point &b) {
   return a.y < b.y;</pre>
// squared distance between points
ntype pdist2(const point &a, const point &b) {
   ntype dx = a.x-b.x, dy = a.y-b.y;
   return dx*dx + dy*dy;
// bounding box for a set of points
struct bbox {
   ntype x0, x1, y0, y1;
   bbox(): x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
   // computes bounding box from a bunch of points
   void compute(const vector<point> &v) {
       for (int i = 0; i < v.size(); ++i) {</pre>
           x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
           y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
       }
   // squared distance between a point and this bbox, 0 if inside
   ntype distance(const point &p) {
       if (p.x < x0) {
           if (p.y < y0) return pdist2(point(x0, y0), p);</pre>
           else if (p.y > y1) return pdist2(point(x0, y1), p);
                     return pdist2(point(x0, p.y), p);
       else if (p.x > x1) {
           if (p.y < y0) return pdist2(point(x1, y0), p);</pre>
           else if (p.y > y1) return pdist2(point(x1, y1), p);
                     return pdist2(point(x1, p.y), p);
       }
       else {
           if (p.y < y0) return pdist2(point(p.x, y0), p);</pre>
           else if (p.y > y1) return pdist2(point(p.x, y1), p);
                     return 0;
};
// stores a single node of the kd-tree, either internal or leaf
struct kdnode {
   bool leaf; // true if this is a leaf node (has one point)
   point pt; // the single point of this is a leaf
   bbox bound; // bounding box for set of points in children
   kdnode *first, *second; // two children of this kd-node
   kdnode() : leaf(false), first(0), second(0) {}
   "kdnode() { if (first) delete first; if (second) delete second; }
   // intersect a point with this node (returns squared distance)
   ntype intersect(const point &p) {
```

```
return bound.distance(p);
   }
   // recursively builds a kd-tree from a given cloud of points
   void construct(vector<point> &vp) {
       // compute bounding box for points at this node
       bound.compute(vp):
       // if we're down to one point, then we're a leaf node
       if (vp.size() == 1) {
           leaf = true;
           pt = vp[0];
       } else {
           // split on x if the bbox is wider than high (not best
               heuristic...)
           if (bound.x1-bound.x0 >= bound.y1-bound.y0)
              sort(vp.begin(), vp.end(), on_x);
           // otherwise split on y-coordinate
           else sort(vp.begin(), vp.end(), on_y);
           // divide by taking half the array for each child
           // (not best performance if many duplicates in the middle)
           int half = vp.size()/2;
           vector<point> vl(vp.begin(), vp.begin()+half);
           vector<point> vr(vp.begin()+half, vp.end());
           first = new kdnode(); first->construct(vl);
           second = new kdnode(); second->construct(vr);
   }
};
// simple kd-tree class to hold the tree and handle queries
struct kdtree {
   kdnode *root;
   // constructs a kd-tree from a points (copied here, as it sorts them)
   kdtree(const vector<point> &vp) {
       vector<point> v(vp.begin(), vp.end());
       root = new kdnode():
       root->construct(v);
    "kdtree() { delete root; }
   // recursive search method returns squared distance to nearest point
   ntype search(kdnode *node, const point &p) {
       if (node->leaf) {
           // commented special case tells a point not to find itself
           //
                  if (p == node->pt) return sentry;
           //
                  else
           return pdist2(p, node->pt);
       ntype bfirst = node->first->intersect(p);
       ntype bsecond = node->second->intersect(p);
       // choose the side with the closest bounding box to search first
       // (note that the other side is also searched if needed)
       if (bfirst < bsecond) {</pre>
           ntype best = search(node->first, p);
           if (bsecond < best)</pre>
              best = min(best, search(node->second, p));
           return best;
       else {
           ntype best = search(node->second, p);
           if (bfirst < best)</pre>
```

```
best = min(best, search(node->first, p));
           return best;
       }
   }
   // squared distance to the nearest
   ntype nearest(const point &p) {
       return search(root, p);
};
// some basic test code here
int main() {
   // generate some random points for a kd-tree
   vector<point> vp;
   for (int i = 0; i < 100000; ++i) {</pre>
       vp.push_back(point(rand()%100000, rand()%100000));
   kdtree tree(vp);
   // query some points
   for (int i = 0; i < 10; ++i) {
       point q(rand()%100000, rand()%100000);
       cout << "Closest squared distance to (" << q.x << ", " << q.y <<</pre>
           << " is " << tree.nearest(q) << endl;</pre>
   }
```

2.3 LCA

```
struct lca {
   int L, N;
   vector<int> depth, size, link;
   lca(const vvi &graph, int root = 0) {
       N = graph.size();
       for (L = 0; (1 << L) <= N; L++);
       depth.resize(N);
       size.resize(N):
       link.resize(L*N);
       init(root, root, graph);
   void init(int loc, int par, const vvi &graph) {
       link[loc] = par;
       for (int 1 = 1; 1 < L; 1++)
          link[1*N + loc] = link[(1-1)*N + link[(1-1)*N + loc]];
       for (int nbr : graph[loc]) {
           if (nbr == par) continue;
           depth[nbr] = depth[loc] + 1;
           init(nbr, loc, graph);
          size[loc] += size[nbr];
       }
       size[loc]++;
   int above(int loc, int dist) {
       for (int 1 = 0; 1 < L; 1++)
          if ((dist >> 1)&1)
              loc = link[l*N + loc];
```

2.4 Lazy Segment Tree

```
// Modular lazy segment tree
// It takes a type T for vertex values and a type U for update
// operations. Type T should have an operator '+' specifying how to
// combine vertices. Type U should have an operator '()' specifying how
// to apply updates to vertices and an operator '+' for combining two
// updates. Example below.
template<typename T, typename U> struct seg_tree_lazy {
   int S, H;
   T zero;
   vector<T> value:
   U noop;
   vector<bool> dirty;
   vector<U> prop;
   seg_tree_lazy<T, U>(int _S, T _zero = T(), U _noop = U()) {
       zero = _zero, noop = _noop;
       for (S = 1, H = 1; S < S;) S *= 2, H++;
       value.resize(2*S, zero);
       dirty.resize(2*S, false);
       prop.resize(2*S, noop);
   void set_leaves(vector<T> &leaves) {
       copy(leaves.begin(), leaves.end(), value.begin() + S);
       for (int i = S - 1; i > 0; i--)
          value[i] = value[2 * i] + value[2 * i + 1];
   void apply(int i, U &update) {
       value[i] = update(value[i]);
       if(i < S) {
          prop[i] = prop[i] + update;
          dirty[i] = true;
   void rebuild(int i) {
       for (int 1 = i/2; 1; 1 /= 2) {
          T combined = value[2*1] + value[2*1+1];
           value[1] = prop[1](combined);
   }
   void propagate(int i) {
       for (int h = H; h > 0; h--) {
          int 1 = i \gg h;
```

```
if (dirty[1]) {
               apply(2*1, prop[1]);
              apply(2*1+1, prop[1]);
              prop[1] = noop;
               dirty[l] = false;
       }
    void upd(int i, int j, U update) {
       i += S, j += S;
       propagate(i), propagate(j);
       for (int 1 = i, r = j; 1 <= r; 1 /= 2, r /= 2) {
           if((1&1) == 1) apply(1++, update);
           if((r\&1) == 0) apply(r--, update);
       rebuild(i), rebuild(j);
    T query(int i, int j){
       i += S, j += S;
       propagate(i), propagate(j);
       T res_left = zero, res_right = zero;
       for(; i <= j; i /= 2, j /= 2){
           if((i&1) == 1) res_left = res_left + value[i++];
           if((j&1) == 0) res_right = value[j--] + res_right;
       return res_left + res_right;
};
// Example that supports following operations:
// 1: Add amount V to the values in range [L,R].
// 2: Reset the values in range [L,R] to value V.
// 3: Query for the sum of the values in range [L,R].
// Here's what T looks like:
struct node {
    int sum, width;
    node operator+(const node &n) {
       return { sum + n.sum, width + n.width }
// Here's what U looks like:
struct update {
    bool type; // 0 for add, 1 for reset
    int value:
    node operator()(const node &n) {
       if(type) return { n.width * value, n.width };
       else return { n.sum + n.width * value, n.width };
    update operator+(const update &u) {
       if(u.type) return u;
       return { type, value + u.value };
}:
int main() {
    int N = 100;
    node zero = \{0,0\};
    update noop = {false, 0};
    vector<node> leaves(N, {0,1});
```

```
seg_tree_lazy<node, update> st(N, zero, noop);
st.set_leaves(leaves);
```

2.5 Segment Tree 2D

```
#define Max 506
#define INF (1 << 30)
int P[Max][Max]; // container for 2D grid
/* 2D Segment Tree node */
struct Point {
   int x, y, mx;
   Point() {}
   Point(int x, int y, int mx) : x(x), y(y), mx(mx) {}
   bool operator < (const Point& other) const {</pre>
       return mx < other.mx;</pre>
};
struct Segtree2d {
   Point T[2 * Max * Max];
   int n, m;
   // initialize and construct segment tree
   void init(int n. int m) {
       this \rightarrow n = n;
       this -> m = m;
       build(1, 1, 1, n, m);
   // build a 2D segment tree from data [ (a1, b1), (a2, b2) ]
   // Time: O(n logn)
   Point build(int node, int a1, int b1, int a2, int b2) {
       // out of range
       if (a1 > a2 \text{ or } b1 > b2)
           return def();
       // if it is only a single index, assign value to node
       if (a1 == a2 \text{ and } b1 == b2)
           return T[node] = Point(a1, b1, P[a1][b1]);
       // split the tree into four segments
       T[node] = def();
       T[node] = maxNode(T[node], build(4 * node - 2, a1, b1, (a1 + a2) /
            2, (b1 + b2) / 2);
       T[node] = maxNode(T[node], build(4 * node - 1, (a1 + a2) / 2 + 1,
            b1, a2, (b1 + b2) / 2));
       T[node] = maxNode(T[node], build(4 * node + 0, a1, (b1 + b2) / 2 +
            1, (a1 + a2) / 2, b2));
       T[node] = maxNode(T[node], build(4 * node + 1, (a1 + a2) / 2 + 1,
            (b1 + b2) / 2 + 1, a2, b2);
       return T[node];
   // helper function for query(int, int, int, int);
   Point query(int node, int a1, int b1, int a2, int b2, int x1, int y1,
        int x2, int y2) {
       // if we out of range, return dummy
       if (x1 > a2 or y1 > b2 or x2 < a1 or y2 < b1 or a1 > a2 or b1 > b2)
```

```
return def():
               // if it is within range, return the node
               if (x1 <= a1 and y1 <= b1 and a2 <= x2 and b2 <= y2)
                       return T[node];
               // split into four segments
               Point mx = def():
               mx = maxNode(mx, query(4 * node - 2, a1, b1, (a1 + a2) / 2, (b1 + a2)) / (a2 + a2)) / (a3 + a2)) / (a4 + a2)) / (a5 + a2
                         b2) / 2, x1, y1, x2, y2) );
               mx = maxNode(mx, query(4 * node - 1, (a1 + a2) / 2 + 1, b1, a2, (b1 + b2) / 2, x1, y1, x2, y2));
               mx = maxNode(mx, query(4 * node + 0, a1, (b1 + b2) / 2 + 1, (a1 + b2))
                         a2) / 2, b2, x1, y1, x2, y2) );
               mx = maxNode(mx, query(4 * node + 1, (a1 + a2) / 2 + 1, (b1 + b2))
                         /2 + 1, a2, b2, x1, y1, x2, y2));
               // return the maximum value
               return mx;
        // query from range [ (x1, y1), (x2, y2) ]
        // Time: O(logn)
        Point query(int x1, int y1, int x2, int y2) {
               return query(1, 1, 1, n, m, x1, y1, x2, y2);
        // helper function for update(int, int, int);
        Point update(int node, int a1, int b1, int a2, int b2, int x, int y,
                 int value) {
               if (a1 > a2 \text{ or } b1 > b2)
                       return def();
               if (x > a2 \text{ or } y > b2 \text{ or } x < a1 \text{ or } y < b1)
                       return T[node];
               if (x == a1 \text{ and } y == b1 \text{ and } x == a2 \text{ and } y == b2)
                       return T[node] = Point(x, y, value);
               Point mx = def();
               mx = maxNode(mx, update(4 * node - 2, a1, b1, (a1 + a2) / 2, (b1 + a2))
                         b2) / 2, x, y, value) );
               mx = maxNode(mx, update(4 * node - 1, (a1 + a2) / 2 + 1, b1, a2,
                         (b1 + b2) / 2, x, y, value));
               mx = maxNode(mx, update(4 * node + 0, a1, (b1 + b2) / 2 + 1, (a1 + b2))
                         a2) / 2, b2, x, y, value));
               mx = maxNode(mx, update(4 * node + 1, (a1 + a2) / 2 + 1, (b1 + b2))
                         / 2 + 1, a2, b2, x, y, value) );
               return T[node] = mx;
        // update the value of (x, y) index to 'value'
        // Time: O(logn)
        Point update(int x, int y, int value) {
               return update(1, 1, 1, n, m, x, y, value);
        // utility functions; these functions are virtual because they will
                  be overridden in child class
        virtual Point maxNode(Point a, Point b) {
               return max(a, b);
        // dummy node
        virtual Point def() {
               return Point(0, 0, -INF);
};
```

```
/* 2D Segment Tree for range minimum query; a override of Segtree2d class
    */
struct Segtree2dMin : Segtree2d {
   // overload maxNode() function to return minimum value
   Point maxNode(Point a, Point b) { return min(a, b); }
   Point def() { return Point(0, 0, INF); }
};
// initialize class objects
Segtree2d Tmax;
Segtree2dMin Tmin;
/* Drier program */
int main(void) {
   int n, m;
   // input
   scanf("%d %d", &n, &m);
   for(int i = 1; i <= n; i++)
       for(int j = 1; j \le m; j++)
           scanf("%d", &P[i][j]);
   // initialize
   Tmax.init(n, m);
   Tmin.init(n, m);
   // query
   int x1, y1, x2, y2;
   scanf("%d %d %d %d", &x1, &y1, &x2, &y2);
   Tmax.query(x1, y1, x2, y2).mx;
   Tmin.query(x1, y1, x2, y2).mx;
   // update
   int x, y, v;
   scanf("%d %d %d", &x, &y, &v);
   Tmax.update(x, y, v);
   Tmin.update(x, y, v);
   return 0;
```

2.6 Segment Tree

```
template<typename T> struct seg_tree {
   int S;
   T zero:
   vector<T> value;
   seg_tree<T>(int _S, T _zero = T()) {
       S = _S, zero = _zero;
       value.resize(2*S+1, zero);
   void set_leaves(vector<T> &leaves) {
       copy(leaves.begin(), leaves.end(), value.begin() + S);
       for (int i = S - 1; i > 0; i--)
          value[i] = value[2 * i] + value[2 * i + 1];
   void upd(int i, T v) {
       i += S;
       value[i] = v:
       while(i>1){
          i/=2;
          value[i] = value[2*i] + value[2*i+1];
```

```
}
}
T query(int i, int j) {
    T res_left = zero, res_right = zero;
    for(i += S, j += S; i <= j; i /= 2, j /= 2){
        if((i&1) == 1) res_left = res_left + value[i++];
        if((j&1) == 0) res_right = value[j--] + res_right;
    }
    return res_left + res_right;
}
</pre>
```

2.7 Union Find

```
// (struct) also keeps track of sizes
struct union find {
   vector<int> P,S;
   union_find(int N) {
       P.resize(N), S.resize(N, 1);
       for(int i = 0; i < N; i++) P[i] = i;</pre>
   int rep(int i) {return (P[i] == i) ? i : P[i] = rep(P[i]);}
   bool unio(int a, int b) {
       a = rep(a), b = rep(b);
       if(a == b) return false;
       P[b] = a;
       S[a] += S[b];
       return true;
// (Shorter) union-find set: the vector/array contains the parent of each
int find(vector <int>& C, int x){return (C[x]==x) ? x : C[x]=find(C,
    int find(int x){return (C[x]==x)?x:C[x]=find(C[x]);} //C
```

3 Graph

3.1 2-SAT

```
struct two_sat {
   int N;
   vector<vector<int>> impl;

   two_sat(int _N) {
      N = _N;
      impl.resize(2 * N);
   }

   void add_impl(int var1, bool neg1, int var2, bool neg2) {
      impl[2 * var1 + neg1].push_back(2 * var2 + neg2);
      impl[2 * var2 + !neg2].push_back(2 * var1 + !neg1);
   }

   void add_clause(int var1, bool neg1, int var2, bool neg2) {
      add_impl(var1, !neg1, var2, neg2);
   }
}
```

```
void add_clause(int var1, bool neg1) {
       add_clause(var1, neg1, var1, neg1);
   int V, L, C;
   stack<int> view;
   int dfs(int loc) {
       visit[loc] = V:
       label[loc] = L++;
       int low = label[loc];
       view.push(loc);
       in_view[loc] = true;
       for (int nbr : impl[loc]) {
           if(!visit[nbr]) low = min(low, dfs(nbr));
           else if(in_view[nbr]) low = min(low, label[nbr]);
       if(low == label[loc]) {
           while (true) {
              int mem = view.top();
              comp[mem] = C;
              in_view[mem] = false;
              view.pop();
              if(mem == loc) break;
           }
           C++;
       return low;
   vector<int> visit, label, comp, in_view;
   void reset(vector<int> &v) {
       v.resize(2 * N);
       fill(v.begin(), v.end(), 0);
   bool consistent() {
       V = 0, L = 0, C = 0;
       reset(visit), reset(label), reset(comp), reset(in_view);
       for (int i = 0; i < 2 * N; i++) {</pre>
           if(!visit[i]) {
              V++:
              dfs(i);
           }
       for (int i = 0; i < N; i++)</pre>
           if(comp[2 * i] == comp[2 * i + 1]) return false;
       return true;
   }
};
```

3.2 Dense Dijkstra

```
void Dijkstra (const VVT &w, VT &dist, VI &prev, int start) {
   int n = w.size();
   VI found (n);
   prev = VI(n, -1);
   dist = VT(n, 1000000000);
```

```
dist[start] = 0;
while (start != -1){
   found[start] = true;
   int best = -1;
   for (int k = 0; k < n; k++) if (!found[k]) {
      if (dist[k] > dist[start] + w[start][k]) {
         dist[k] = dist[start] + w[start][k];
        prev[k] = start;
      }
      if (best == -1 || dist[k] < dist[best]) best = k;
   }
   start = best;
}</pre>
```

3.3 Dijkstra

```
// Implementation of Dijkstra's algorithm using adjacency lists
// and priority queue for efficiency.
// Running time: O(|E| log |V|)
typedef pair<int,int> PII;
const int INF = 2000000000;
int main(){
   int N, s, t;
   scanf ("%d%d%d", &N, &s, &t);
   vector<vector<PII> > edges(N);
   for (int i = 0; i < N; i++){
       int M:
       scanf ("%d", &M);
      for (int j = 0; j < M; j++){
          int vertex, dist;
          scanf ("%d%d", &vertex, &dist);
          edges[i].push_back (make_pair (dist, vertex)); // note order
               of arguments here
       }
   // use priority queue in which top element has the "smallest" priority
   priority_queue<PII, vector<PII>, greater<PII> > Q;
   vector<int> dist(N, INF), dad(N, -1);
   Q.push (make_pair (0, s));
   dist[s] = 0;
   while (!Q.empty()){
       PII p = Q.top();
       if (p.second == t) break;
       Q.pop();
       int here = p.second;
       for (vector<PII>::iterator it=edges[here].begin();
           it!=edges[here].end(); it++){
           if (dist[here] + it->first < dist[it->second]){
              dist[it->second] = dist[here] + it->first;
              dad[it->second] = here;
              Q.push (make_pair (dist[it->second], it->second));
```

```
printf ("%d\n", dist[t]);
if (dist[t] < INF)
    for(int i=t;i!=-1;i=dad[i])
        printf ("%d%c", i, (i==s?'\n':' '));
return 0;
}</pre>
```

3.4 Eulerian Path

```
struct Edge:
typedef list<Edge>::iterator iter;
struct Edge {
   int next_vertex;
   iter reverse_edge;
   Edge(int next_vertex) : next_vertex(next_vertex) {}
};
const int max_vertices = 100;
int num vertices:
list<Edge> adj[max_vertices]; // adjacency list
vector<int> path;
void find_path(int v) {
   while(adj[v].size() > 0) {
       int vn = adj[v].front().next_vertex;
       adj[vn].erase(adj[v].front().reverse_edge);
       adj[v].pop_front();
       find_path(vn);
   path.push_back(v);
void add_edge(int a, int b) {
   adj[a].push_front(Edge(b));
   iter ita = adj[a].begin();
   adj[b].push_front(Edge(a));
   iter itb = adj[b].begin();
   ita->reverse_edge = itb;
   itb->reverse_edge = ita;
```

3.5 Heavy Light

```
template<typename T> struct heavy_light {
    lca links;
    seg_tree<T> st;
    vector<int> preorder, index, jump;

    heavy_light(const vvi &graph, int root) {
        links = lca(graph, 0);
        st = seg_tree<T>(graph.size());
        index.resize(graph.size()), jump.resize(graph.size());
        dfs(root, root, root, graph);
    }
    void dfs(int loc, int par, int lhv, const vvi &graph) {
        jump[loc] = lhv;
        index[loc] = preorder.size();
    }
}
```

```
preorder.push_back(loc);
       vector<int> ch = graph[loc];
       sort(ch.begin(), ch.end(), [&](int i, int j) {
              return links.size[i] > links.size[j]; });
       if (loc != par) ch.erase(ch.begin());
       for (int c = 0; c < ch.size(); c++)</pre>
           dfs(ch[c], loc, c ? ch[c] : lhv, graph);
   void assign(int loc, T value) {
       st.upd(index[loc], value);
   T __sum(int u, int r) {
       T res;
       while (u != r) {
           int go = max(index[r] + 1, index[jump[u]]);
           res = res + st.query(go, index[u]);
          u = links.link[preorder[go]];
       }
       return res;
   T sum(int u, int v) {
       int r = links.find(u, v);
       return st.query(index[r], index[r]) + __sum(u, r) + __sum(v, r);
};
```

3.6 Poset Width

```
// requires bipartite graph (4.1)
vector<int> width(vector<vector<int>> poset) {
   int N = poset.size();
   bipartite_graph g(N, N);
   for (int i = 0; i < N; i++) {
       for (int j : poset[i])
          g.edge(i, i);
   g.matching();
   vector<bool> vis[2];
   vis[false].resize(2 * N, false);
   vis[true].resize(2 * N, false);
   for (int i = 0; i < N; i++) {</pre>
       if (g.match[i] != -1) continue;
       if (vis[false][i]) continue;
       queue<pair<bool, int>> bfs;
       bfs.push(make_pair(false, i));
       vis[false][i] = true;
       while (!bfs.empty()) {
           bool inm = bfs.front().first;
           int loc = bfs.front().second;
          bfs.pop();
          for (int nbr : g.adj[loc]) {
              if (vis[!inm][nbr]) continue;
              if ((g.match[loc] == nbr) ^ inm) continue;
              vis[!inm][nbr] = true;
              bfs.push(make_pair(!inm, nbr));
```

```
vector<bool> inz(2 * N, false);
for (int i = 0; i < 2 * N; i++)
   inz[i] = vis[true][i] || vis[false][i];
vector<bool> ink(N, false);
for (int i = 0; i < N; i++)</pre>
   if (!inz[i])
       ink[i] = true;
for (int i = N; i < 2 * N; i++)
   if (inz[i])
       ink[i - N] = true;
vector<int> res;
for (int i = 0; i < N; i++) {</pre>
   if (!ink[i])
       res.push_back(i);
}
return res;
```

3.7 SCC

```
struct edge{int e, nxt;};
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x) {
   int i;
   v[x]=true;
   for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
   stk[++stk[0]]=x:
void fill_backward(int x) {
   int i:
   v[x]=false;
   group_num[x]=group_cnt;
   for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e]) fill_backward(er[i].e);
void add_edge(int v1, int v2) { //add edge v1->v2
   e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
   er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC() {
   int i:
   stk[0]=0;
   memset(v, false, sizeof(v));
   for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);</pre>
   for(i=stk[0];i>=1;i--) if(v[stk[i]]){group_cnt++;
        fill_backward(stk[i]);}
```

3.8 Topological Sort

```
// This function uses performs a non-recursive topological sort. //
```

```
// Running time: O(|V|^2). If you use adjacency lists (vector<map<int> >),
//
         the running time is reduced to O(|E|).
//
// INPUT: w[i][j] = 1 if i should come before j, 0 otherwise
    OUTPUT: a permutation of 0,...,n-1 (stored in a vector)
       which represents an ordering of the nodes which
//
       is consistent with w
// If no ordering is possible, false is returned.
typedef double TYPE;
typedef vector<TYPE> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool TopologicalSort (const VVI &w, VI &order){
   int n = w.size();
   VI parents (n);
   queue<int> q;
   order.clear();
   for (int i = 0; i < n; i++){
       for (int j = 0; j < n; j++)
           if (w[j][i]) parents[i]++;
       if (parents[i] == 0) q.push (i);
   while (q.size() > 0){
       int i = q.front();
       q.pop();
       order.push_back (i);
       for (int j = 0; j < n; j++) if (w[i][j]){
           parents[j]--:
           if (parents[j] == 0) q.push (j);
   return (order.size() == n);
```

4 4 Combinatorial Optimization

4.1 Bipartite Graph

```
struct bipartite_graph {
   int A, B;
   vector<vector<int>> adj;

   bipartite_graph(int _A, int _B) {
        A = _A, B = _B;
        adj.resize(A + B);
   }

   void edge(int i, int j) {
        adj[i].push_back(A+j);
        adj[A+j].push_back(i);
   }

   vector<int> visit, match;
```

```
bool augment(int loc, int run) {
       if(visit[loc] == run) return false;
       visit[loc] = run;
       for (int nbr : adj[loc]) {
          if (match[nbr] == -1 || augment(match[nbr], run)) {
              match[loc] = nbr, match[nbr] = loc;
              return true;
       }
       return false;
   int matching() {
       visit = vector<int>(A+B, -1);
       match = vector<int>(A+B, -1);
       int ans = 0;
       for (int i = 0; i < A; i++)</pre>
          ans += augment(i, i);
       return ans:
   }
   vector<bool> vertex_cover() {
       vector<bool> res(A + B, false);
       queue<int> bfs;
       for (int i = 0; i < A; i++) {
           if (match[i] == -1) bfs.push(i);
           else res[i] = true;
       while (!bfs.empty()) {
          int loc = bfs.front();
           bfs.pop();
           for (int nbr : adj[loc]) {
              if (res[nbr]) continue;
              res[nbr] = true;
              int loc2 = match[nbr];
              if (loc2 == -1) continue;
              res[loc2] = false:
              bfs.push(loc2);
          }
       }
       return res;
   }
};
```

4.2 Max Flow - Dinic

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
//
// Running time:
// O(|V|^2 |E|)
//
// INPUT:
// - graph, constructed using AddEdge()
// - source and sink
//
// OUTPUT:
// - maximum flow value
// - To obtain actual flow values, look at edges with capacity > 0
// (zero capacity edges are residual edges).
```

```
struct Edge {
    int from, to, cap, flow, index;
    Edge(int from, int to, int cap, int flow, int index) :
       from(from), to(to), cap(cap), flow(flow), index(index) {}
    11 rcap() { return cap - flow; }
};
struct Dinic {
    int N:
    vector<vector<Edge> > G;
    vector<vector<Edge *> > Lf:
    vector<int> layer;
    vector<int> Q;
    Dinic(int N) : N(N), G(N), Q(N) 
    void AddEdge(int from, int to, int cap) {
       if (from == to) return;
       G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
       G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
    11 BlockingFlow(int s, int t) {
       layer.clear(); layer.resize(N, -1);
       layer[s] = 0;
       Lf.clear(); Lf.resize(N);
       int head = 0, tail = 0;
       Q[tail++] = s;
       while (head < tail) {</pre>
           int x = Q[head++];
           for (int i = 0; i < G[x].size(); i++) {</pre>
               Edge &e = G[x][i]; if (e.rcap() <= 0) continue;
               if (layer[e.to] == -1) {
                   layer[e.to] = layer[e.from] + 1;
                   Q[tail++] = e.to;
               if (layer[e.to] > layer[e.from]) {
                  Lf[e.from].push_back(&e);
           }
       }
       if (layer[t] == -1) return 0;
       11 totflow = 0;
       vector<Edge *> P;
       while (!Lf[s].empty()) {
           int curr = P.empty() ? s : P.back()->to;
if (curr == t) { // Augment
               11 amt = P.front()->rcap();
               for (int i = 0; i < P.size(); ++i) {</pre>
                   amt = min(amt, P[i]->rcap());
               totflow += amt;
               for (int i = P.size() - 1; i >= 0; --i) {
                  P[i]->flow += amt:
                   G[P[i]->to][P[i]->index].flow -= amt;
                   if (P[i]->rcap() <= 0) {</pre>
                      Lf[P[i]->from].pop_back();
                      P.resize(i);
```

```
} else if (Lf[curr].empty()) { // Retreat
               P.pop_back();
               for (int i = 0; i < N; ++i)</pre>
                   for (int j = 0; j < Lf[i].size(); ++j)</pre>
                       if (Lf[i][j]->to == curr)
                           Lf[i].erase(Lf[i].begin() + j);
           } else { // Advance
               P.push_back(Lf[curr].back());
       }
       return totflow;
    11 GetMaxFlow(int s, int t) {
       11 \text{ totflow} = 0;
       while (ll flow = BlockingFlow(s, t))
           totflow += flow;
       return totflow;
    }
};
```

4.3 Min Cost Matching

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
// second.
// cost[i][j] = cost for pairing left node i with right node j
   Lmate[i] = index of right node that left node i pairs with
   Rmate[j] = index of left node that right node j pairs with
//
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
   int n = int(cost.size());
   // construct dual feasible solution
   VD u(n);
   VD v(n);
   for (int i = 0; i < n; i++) {</pre>
      u[i] = cost[i][0];
      for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
   }
   for (int j = 0; j < n; j++) {
      v[j] = cost[0][j] - u[0];
      for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
   }
```

```
// construct primal solution satisfying complementary slackness
Lmate = VI(n, -1);
Rmate = VI(n, -1);
int mated = 0;
for (int i = 0; i < n; i++) {</pre>
   for (int j = 0; j < n; j++) {
       if (Rmate[j] != -1) continue;
       if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
           Lmate[i] = j;
           Rmate[j] = i;
           mated++;
           break;
       }
   }
VD dist(n);
VI dad(n);
VI seen(n);
// repeat until primal solution is feasible
while (mated < n) {</pre>
   // find an unmatched left node
   int s = 0;
   while (Lmate[s] !=-1) s++;
   // initialize Dijkstra
   fill(dad.begin(), dad.end(), -1);
   fill(seen.begin(), seen.end(), 0);
   for (int k = 0; k < n; k++)
       dist[k] = cost[s][k] - u[s] - v[k];
   int j = 0;
   while (true) {
       // find closest
       j = -1;
       for (int k = 0; k < n; k++) {
           if (seen[k]) continue;
           if (j == -1 || dist[k] < dist[j]) j = k;</pre>
       seen[i] = 1;
       // termination condition
       if (Rmate[j] == -1) break;
       // relax neighbors
       const int i = Rmate[j];
       for (int k = 0; k < n; k++) {
           if (seen[k]) continue;
           const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
           if (dist[k] > new_dist) {
               dist[k] = new_dist;
               dad[k] = j;
       }
   }
   // update dual variables
   for (int k = 0; k < n; k++) {
       if (k == j || !seen[k]) continue;
       const int i = Rmate[k];
       v[k] += dist[k] - dist[j];
       u[i] -= dist[k] - dist[j];
   u[s] += dist[j];
   // augment along path
```

```
while (dad[j] >= 0) {
    const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
    j = d;
}
Rmate[j] = s;
Lmate[s] = j;
    mated++;
}
double value = 0;
for (int i = 0; i < n; i++)
    value += cost[i][Lmate[i]];
return value;</pre>
```

4.4 Min Cost Max Flow

```
// Min cost max flow algorithm using an adjacency matrix. If you
// want just regular max flow, setting all edge costs to 1 gives
// running time O(|E|^2 |V|).
//
// Running time: O(\min(|V|^2 * totflow, |V|^3 * totcost))
// INPUT: cap -- a matrix such that cap[i][j] is the capacity of
         a directed edge from node i to node j
     cost -- a matrix such that cost[i][j] is the (positive)
         cost of sending one unit of flow along a
//
         directed edge from node i to node j
     source -- starting node
     sink -- ending node
//
// OUTPUT: max flow and min cost; the matrix flow will contain
//
      the actual flow values (note that unlike in the MaxFlow
      code, you don't need to ignore negative flow values -- there
//
      shouldn't be any)
typedef vector<ll> vll;
typedef vector<vll> vvll;
const 11 INF = 1LL << 60;</pre>
struct MCMF {
   int N:
   vll found, dad, dist, pi;
   vvll cap, flow, cost;
   MCMF(int N) : N(N), cap(N, vll(N)), flow(cap), cost(cap),
   dad(N), found(N), pi(N), dist(N+1) {};
   void add_edge(int from, int to, ll ca, ll co) {
       cap[from][to] = ca; cost[from][to] = co; }
   bool search(int source, int sink) {
       fill(found.begin(), found.end(), 0);
       fill(dist.begin(), dist.end(), INF);
       dist[source] = 0;
       while(source != N) {
           int best = N;
           found[source] = 1;
           for(int k = 0; k < N; k++) {
```

```
if(found[k]) continue;
              if(flow[k][source]) {
                  ll val = dist[source] + pi[source] - pi[k] -
                      cost[k][source];
                  if(dist[k] > val) {
                      dist[k] = val:
                      dad[k] = source;
              if(flow[source][k] < cap[source][k]) {</pre>
                  11 val = dist[source] + pi[source] - pi[k] +
                      cost[source][k];
                  if(dist[k] > val) {
                      dist[k] = val;
                      dad[k] = source;
              if(dist[k] < dist[best]) best = k;</pre>
           }
           source = best;
       for(int k = 0; k < N; k++)
           pi[k] = min((11)(pi[k] + dist[k]), INF);
       return found[sink];
   pair<11.11> mcmf(int source, int sink) {
       11 totflow = 0, totcost = 0;
       while(search(source, sink)) {
          11 amt = INF;
           for(int x = sink; x != source; x = dad[x])
              amt = min(amt, (11)(flow[x][dad[x]] != 0 ?
                         flow[x][dad[x]] : cap[dad[x]][x] -
                              flow[dad[x]][x]));
           for(int x = sink; x != source; x = dad[x]) {
              if(flow[x][dad[x]] != 0) {
                  flow[x][dad[x]] -= amt;
                  totcost -= amt * cost[x][dad[x]];
              } else {
                  flow[dad[x]][x] += amt;
                  totcost += amt * cost[dad[x]][x];
           }
           totflow += amt;
       return {totflow, totcost};
};
```

4.5 Min Cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
//
// Running time:
// O(|V|^3)
//
// INPUT:
// - graph, constructed using AddEdge()
//
```

```
// OUTPUT:
// - (min cut value, nodes in half of min cut)
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
   int N = weights.size();
   VI used(N), cut, best_cut;
   int best_weight = -1;
   for (int phase = N-1; phase >= 0; phase--) {
       VI w = weights[0];
       VI added = used;
       int prev, last = 0;
       for (int i = 0; i < phase; i++) {</pre>
           prev = last;
           last = -1;
           for (int j = 1; j < N; j++)
              if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
           if (i == phase-1) {
              for (int j = 0; j < N; j++) weights[prev][j] +=</pre>
                   weights[last][j];
              for (int j = 0; j < N; j++) weights[j][prev] =</pre>
                   weights[prev][j];
              used[last] = true;
              cut.push_back(last);
              if (best_weight == -1 || w[last] < best_weight) {</pre>
                  best_cut = cut;
                  best_weight = w[last];
           } else {
              for (int j = 0; j < N; j++)
                  w[j] += weights[last][j];
              added[last] = true;
           }
   }
   return make_pair(best_weight, best_cut);
```

5 5 Geometry

5.1 Convex Hull

```
typedef double T;
const T EPS = 1e-7;
struct PT {
   T x, y;
   PT() {}
   PT(T x, T y) : x(x), y(y) {}
   bool operator<(const PT &rhs) const { return make_pair(y,x) <</pre>
        make_pair(rhs.y,rhs.x); }
   bool operator == (const PT &rhs) const { return make_pair(y,x) ==
        make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE_REDUNDANT
bool between(const PT &a, const PT &b, const PT &c) {
   return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 &&
        (a.y-b.y)*(c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
   sort(pts.begin(), pts.end());
   pts.erase(unique(pts.begin(), pts.end()), pts.end());
   vector<PT> up, dn;
   for (int i = 0; i < pts.size(); i++) {</pre>
       while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i])
            >= 0) up.pop_back();
       while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i])
            <= 0) dn.pop_back();
       up.push_back(pts[i]);
       dn.push_back(pts[i]);
   for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
   if (pts.size() <= 2) return;</pre>
   dn.clear();
   dn.push_back(pts[0]);
   dn.push_back(pts[1]);
   for (int i = 2; i < pts.size(); i++) {</pre>
       if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i]))
            dn.pop_back();
       dn.push_back(pts[i]);
   if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
       dn[0] = dn.back();
       dn.pop_back();
   pts = dn;
#endif
```

5.2 Delaunay

```
// Slow but simple Delaunay triangulation. Does not handle // degenerate cases (from O'Rourke, Computational Geometry in C) // Running time: O(n^4)
```

```
// INPUT: x[] = x-coordinates
       y[] = y-coordinates
//
// OUTPUT: triples = a vector containing m triples of indices
            corresponding to triangle vertices
typedef double T;
struct triple {
   int i, j, k;
   triple() {}
   triple(int i, int j, int k) : i(i), j(j), k(k) {}
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
   int n = x.size();
   vector<T> z(n):
   vector<triple> ret;
   for (int i = 0; i < n; i++)
       z[i] = x[i] * x[i] + y[i] * y[i];
   for (int i = 0; i < n-2; i++) {
       for (int j = i+1; j < n; j++) {
           for (int k = i+1; k < n; k++) {
              if (i == k) continue;
              double xn = (y[j]-y[i])*(z[k]-z[i]) -
                   (y[k]-y[i])*(z[j]-z[i]);
              double yn = (x[k]-x[i])*(z[j]-z[i]) -
                   (x[j]-x[i])*(z[k]-z[i]);
              double zn = (x[j]-x[i])*(y[k]-y[i]) -
                   (x[k]-x[i])*(y[j]-y[i]);
              bool flag = zn < 0;
              for (int m = 0; flag && m < n; m++)</pre>
                  flag = flag && ((x[m]-x[i])*xn +
                          (y[m]-y[i])*yn +
                          (z[m]-z[i])*zn <= 0);
              if (flag) ret.push_back(triple(i, j, k));
          }
       }
   }
   return ret;
int main() {
   T xs[]={0, 0, 1, 0.9};
   T ys[]={0, 1, 0, 0.9};
   vector<T> x(\&xs[0], \&xs[4]), y(\&ys[0], \&ys[4]);
   vector<triple> tri = delaunayTriangulation(x, y);
   //expected: 0 1 3
   //
          0 3 2
   int i;
   for(i = 0; i < tri.size(); i++)</pre>
       printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
   return 0;
```

5.3 Geometry

```
// C++ routines for computational geometry.
double INF = 1e100;
double EPS = 1e-12;
```

```
struct PT {
   double x, y;
   PT() {}
   PT(double x, double y) : x(x), y(y) {}
   PT(const PT &p) : x(p.x), y(p.y) {}
   PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
   PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
   PT operator * (double c) const { return PT(x*c, y*c ); }
   PT operator / (double c) const { return PT(x/c, y/c); }
double dot(PT p, PT q) { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
   os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
   return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
   return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
   double r = dot(b-a,b-a);
   if (fabs(r) < EPS) return a;</pre>
   r = dot(c-a, b-a)/r;
   if (r < 0) return a;
   if (r > 1) return b;
   return a + (b-a)*r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
   return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
       double a, double b, double c, double d) {
   return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
   return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
   return LinesParallel(a, b, c, d)
```

```
&& fabs(cross(a-b, a-c)) < EPS
       && fabs(cross(c-d, c-a)) < EPS;
}
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
   if (LinesCollinear(a, b, c, d)) {
       if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
              dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
       if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
           return false:
       return true;
   if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
   if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
   return true;
}
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
   b=\bar{b}-a; d=c-d; c=c-a;
   assert(dot(b, b) > EPS \&\& dot(d, d) > EPS);
   return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
   b=(a+b)/2:
   c=(a+c)/2:
   return ComputeLineIntersection(b, b+RotateCW90(a-b), c,
        c+RotateCW90(a-c));
}
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
   bool c = 0;
   for (int i = 0; i < p.size(); i++){</pre>
       int j = (i+1)%p.size();
       if ((p[i].y <= q.y && q.y < p[j].y ||
                  p[j].y \le q.y && q.y < p[i].y) &&
              q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) /
                   (p[j].y - p[i].y)
           c = !c:
   }
   return c;
}
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
```

```
for (int i = 0; i < p.size(); i++)</pre>
       if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) <</pre>
           return true;
   return false;
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
   vector<PT> ret;
   b = b-a:
   a = a-c;
   double A = dot(b, b);
   double B = dot(a, b);
   double C = dot(a, a) - r*r;
   double D = B*B - A*C;
   if (D < -EPS) return ret:
   ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
   if (D > EPS)
       ret.push_back(c+a+b*(-B-sqrt(D))/A);
   return ret;
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a. PT b. double r. double R) {
   vector<PT> ret;
   double d = sqrt(dist2(a, b));
   if (d > r+R \mid d+min(r, R) < max(r, R)) return ret;
   double x = (d*d-R*R+r*r)/(2*d);
   double y = sqrt(r*r-x*x);
   PT v = (b-a)/d;
   ret.push_back(a+v*x + RotateCCW90(v)*y);
   if (y > 0)
       ret.push_back(a+v*x - RotateCCW90(v)*y);
   return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
   double area = 0:
   for(int i = 0; i < p.size(); i++) {</pre>
       int j = (i+1) % p.size();
       area += p[i].x*p[j].y - p[j].x*p[i].y;
   return area / 2.0;
}
double ComputeArea(const vector<PT> &p) {
   return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
   PT c(0.0):
   double scale = 6.0 * ComputeSignedArea(p);
```

```
for (int i = 0; i < p.size(); i++){</pre>
       int j = (i+1) % p.size();
       c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
   return c / scale;
}
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
   for (int i = 0; i < p.size(); i++) {</pre>
       for (int k = i+1; k < p.size(); k++) {</pre>
           int j = (i+1) % p.size();
           int 1 = (k+1) % p.size();
           if (i == 1 || j == k) continue;
           if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
              return false;
       }
   }
   return true;
// Plane distance between parallel planes aX + bY + cZ + d1 = 0 and
// aX + bY + cZ + d2 = 0 is abs(d1 - d2) / sqrt(a*a + b*b + c*c)
// distance from point (px, py, pz) to line (x1, y1, z1)-(x2, y2, z2)
// (or ray, or segment; in the case of the ray, the endpoint is the
// first point)
public static final int LINE = 0;
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1, double y1, double z1,
       double x2, double y2, double z2, double px, double py, double pz,
       int type) {
   double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-y2) + (z1-z2)*(z1-z2);
   double x, y, z;
   if (pd2 == 0) {
       x = x1;
       y = y1;
       z = z1;
   else {
       double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (pz-z1)*(z2-z1)) /
           pd2;
       x = x1 + u * (x2 - x1);
       y = y1 + u * (y2 - y1);
       z = z1 + u * (z2 - z1):
       if (type != LINE && u < 0) {</pre>
           x = x1;
           y = y1;
           z = z1;
       if (type == SEGMENT && u > 1.0) {
           x = x2;
           y = y2;
           z = z2;
   return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)*(z-pz);
}
int main() {
   // expected: (-5,2)
```

```
cerr << RotateCCW90(PT(2,5)) << endl;</pre>
// expected: (5,-2)
cerr << RotateCW90(PT(2,5)) << endl;</pre>
// expected: (-5,2)
cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
// expected: (5,2)
cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
// expected: (5,2) (7.5,3) (2.5,1)
cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
    << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "</pre>
   << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;</pre>
// expected: 6.78903
cerr \lt\lt DistancePointPlane(4,-4,3,2,-2,5,-8) \lt\lt endl;
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
   << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
   << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
// expected: 0 0 1
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
    << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "</pre>
    << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
// expected: 1 1 1 0
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
    << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "</pre>
    << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
    << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;</pre>
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3))
    << endl;
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v;
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
    << PointInPolygon(v, PT(2,0)) << " "
    << PointInPolygon(v, PT(0,2)) << " "
    << PointInPolygon(v, PT(5,2)) << " "</pre>
    << PointInPolygon(v, PT(2,5)) << endl;</pre>
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
    << PointOnPolygon(v, PT(2,0)) << " "
   << PointOnPolygon(v, PT(0,2)) << " "</pre>
    << PointOnPolygon(v, PT(5,2)) << " "
    << PointOnPolygon(v, PT(2,5)) << endl;</pre>
// expected: (1,6)
        (5,4)(4,5)
//
//
        blank line
//
        (4,5)(5,4)
//
        blank line
        (4,5)(5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
```

```
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;</pre>
```

6 6 Numerics

6.1 Euclid

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
typedef vector<int> VI;
typedef pair<int,int> PII;
// return a % b (positive value)
int mod(int a, int b) {
   return ((a%b)+b)%b;
// computes gcd(a,b)
int gcd(int a, int b) {
   int tmp;
   while(b){a%=b; tmp=a; a=b; b=tmp;}
   return a;
// computes lcm(a,b)
int lcm(int a, int b) {
   return a/gcd(a,b)*b;
// returns d = gcd(a,b); finds x,y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
   int xx = y = 0;
   int yy = x = 1;
   while (b) {
       int q = a/b;
       int \bar{t} = b; b = a\%b; a = t;
       t = xx; xx = x-q*xx; x = t;
       t = yy; yy = y-q*yy; y = t;
   }
   return a;
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
   int x, y;
   VI solutions;
   int d = extended_euclid(a, n, x, y);
```

```
if (!(b%d)) {
       x = mod(x*(b/d), n);
       for (int i = 0; i < d; i++)</pre>
           solutions.push_back(mod(x + i*(n/d), n));
   return solutions:
// computes b such that ab = 1 (mod n), returns -1 on failure
int mod_inverse(int a, int n) {
   int x, y;
   int d = extended_euclid(a, n, x, y);
   if (d > 1) return -1;
   return mod(x,n);
// Chinese remainder theorem (special case): find z such that
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y).
// Return (z,M). On failure, M = -1.
PII chinese_remainder_theorem(int x, int a, int y, int b) {
   int s, t;
   int d = extended_euclid(x, y, s, t);
   if (a%d != b%d) return make_pair(0, -1);
   return make_pair(mod(s*b*x+t*a*y,x*y)/d, x*y/d);
// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm_i (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &x, const VI &a) {
   PII ret = make_pair(a[0], x[0]);
   for (int i = 1; i < x.size(); i++) {</pre>
       ret = chinese_remainder_theorem(ret.second, ret.first, x[i], a[i]);
       if (ret.second == -1) break;
   return ret:
// computes x and y such that ax + by = c; on failure, x = y = -1
void linear_diophantine(int a, int b, int c, int &x, int &y) {
   int d = gcd(a,b);
   if (c%d) {
       x = y = -1;
   } else {
       x = c/d * mod_inverse(a/d, b/d);
       y = (c-a*x)/b;
}
int main() {
   // expected: 2
   cout << gcd(14, 30) << endl;</pre>
   // expected: 2 -2 1
   int x, y;
   int d = extended_euclid(14, 30, x, y);
   cout << d << " " << x << " " << y << endl;
   // expected: 95 45
   VI sols = modular_linear_equation_solver(14, 30, 100);
   for (int i = 0; i < (int) sols.size(); i++) cout << sols[i] << " ";</pre>
   cout << endl;</pre>
   // expected: 8
```

6.2 FFT

```
namespace fft {
   struct cnum {
       double a, b;
       cnum operator+(const cnum &c) { return { a + c.a, b + c.b }; }
       cnum operator-(const cnum &c) { return { a - c.a, b - c.b }; }
       cnum operator*(const cnum &c) { return { a*c.a - b*c.b, a*c.b +
       cnum operator/(double d) { return { a / d, b / d }; }
   };
   const double PI = 2 * atan2(1, 0);
   int deg;
   vector<int> rev;
   void set_degree(int _deg) {
       assert(__builtin_popcount(_deg) == 1);
       deg = _deg;
       rev.resize(deg);
       for (int i = 1, j = 0; i < deg; i++) {
           int bit = deg / 2;
          for (; j >= b̄it; bit /= 2)
             j -= bit;
          j += bit;
          rev[i] = j;
   void transform(vector<cnum> &poly, bool invert) {
       if(deg != poly.size()) set_degree(poly.size());
       for (int i = 1; i < deg; i++)
          if(rev[i] > i)
              swap(poly[i], poly[rev[i]]);
       for (int len = 2; len <= deg; len *= 2) {
           double ang = 2 * PI / len * (invert ? -1 : 1);
           cnum base = { cos(ang), sin(ang) };
           for (int i = 0; i < deg; i += len) {
              cnum w = \{1, 0\};
              for (int j = 0; j < len / 2; j++) {
                  cnum u = poly[i+j];
                  cnum v = w * poly[i+j+len/2];
                  poly[i+j] = u + v;
                  polv[i+j+len/2] = u - v;
                  w = w * base;
```

6.3 Gauss-Jordan

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
    (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
// (3) computing determinants of square matrices
// Running time: O(n^3)
//
// INPUT: a[][] = an nxn matrix
        b[][] = an nxm matrix
//
// OUTPUT: X = an nxm matrix (stored in b[][])
        A^{-1} = an nxn matrix (stored in a[][])
        returns determinant of a[][]
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
const double EPS = 1e-10;
T GaussJordan(VVT &a, VVT &b) {
   const int n = a.size();
   const int m = b[0].size();
   VI irow(n), icol(n), ipiv(n);
   T \det = 1;
   for (int i = 0; i < n; i++) {
       int pj = -1, pk = -1;
       for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
           for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
              if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j;}
                   pk = k; 
       if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." <<</pre>
            endl; exit(0); }
       ipiv[pk]++;
       swap(a[pj], a[pk]);
       swap(b[pj], b[pk]);
       if (pj != pk) det *= -1;
       irow[i] = pj;
       icol[i] = pk;
       T c = 1.0 / a[pk][pk];
       det *= a[pk][pk];
       a[pk][pk] = 1.0;
       for (int p = 0; p < n; p++) a[pk][p] *= c;</pre>
       for (int p = 0; p < m; p++) b[pk][p] *= c;
```

```
for (int p = 0; p < n; p++) if (p != pk) {
           c = a[p][pk];
           a[p][pk] = 0;
           for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
           for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
   }
   for (int p = n-1; p \ge 0; p--) if (irow[p] != icol[p]) {
       for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
   return det;
int main() {
   const int n = 4;
   const int m = 2;
   double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
   double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
   VVT a(n), b(n);
   for (int i = 0; i < n; i++) {
       a[i] = VT(A[i], A[i] + n);
       b[i] = VT(B[i], B[i] + m);
   }
   double det = GaussJordan(a, b);
   // expected: 60
   cout << "Determinant: " << det << endl;</pre>
   // expected: -0.233333 0.166667 0.133333 0.0666667
   //
            0.166667 0.166667 0.333333 -0.333333
   //
            //
            0.05 -0.75 -0.1 0.2
   cout << "Inverse: " << endl;</pre>
   for (int i = 0; i < n; i++) {
       for (int j = 0; j < n; j++)
           cout << a[i][i] << ' ';
       cout << endl:</pre>
   }
   // expected: 1.63333 1.3
            -0.166667 0.5
   //
            2.36667 1.7
            -1.85 -1.35
   cout << "Solution: " << endl;</pre>
   for (int i = 0; i < n; i++) {</pre>
       for (int j = 0; j < m; j++)
           cout << b[i][j] << ' ';
       cout << endl;</pre>
   }
}
```

6.4 Matrix

```
template<typename T> struct matrix {
  int N;
  vector<T> dat;

matrix<T> (int _N, T fill = T(0), T diag = T(0)) {
    N = _N;
    dat.resize(N * N, fill);
```

```
for (int i = 0; i < N; i++)</pre>
           (*this)(i, i) = diag;
    T& operator()(int i, int j) {
       return dat[N * i + j];
    matrix<T> operator *(matrix<T> &b){
       matrix<T> r(N);
       for(int i=0; i<N; i++)</pre>
           for(int j=0; j<N; j++)</pre>
               for(int k=0; k<N; k++)</pre>
                   r(i, j) = r(i, j) + (*this)(i, k) * b(k, j);
       return r;
   matrix<T> pow(ll expo){
       if(!expo) return matrix<T>(N, T(0), T(1));
       matrix<T> r = (*this * *this).pow(expo/2);
       return expo&1 ? r * *this : r;
    friend ostream& operator<<(ostream &os, matrix<T> &m){
       os << "{";
       for(int i=0; i<m.N; i++){</pre>
           if(i) os << "},\n ";
           os << "{";
           for(int j=0; j<m.N; j++){</pre>
               if(j) os << ", ":
               os << setw(10) << m(i, j) << setw(0);
           }
       }
       return os << "}}";
};
struct mll {
    const int MOD:
    ll val;
    mll(ll _val = 0) {
       val = _val % MOD;
       if (val < 0) val += MOD;</pre>
   mll operator+(const mll &o) {
       return mll((val + o.val) % MOD);
   mll operator*(const mll &o) {
       return mll((val * o.val) % MOD);
    friend ostream& operator<<(ostream &os, mll &m) {</pre>
       return os << m.val;</pre>
};
```

6.5 Primes

```
// O(sqrt(x)) Exhaustive Primality Test
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x) {
   if(x<=1) return false;</pre>
```

```
if(x<=3) return true;
   if (!(x\%2) || !(x\%3)) return false;
   LL s=(LL)(sqrt((double)(x))+EPS);
   for(LL i=5;i<=s;i+=6)</pre>
      if (!(x%i) || !(x%(i+2))) return false;
   return true:
// Primes less than 1000:
    2 3 5 7 11 13 17 19 23
    41 43 47 53 59 61 67 71 73 79
                                   83 89
       101 103
                 107
                      109
                          113 127
                                    131 137
                                               139
                                                   149
                                                        151
// 157
       163
            167
                 173
                           181
                                     193
                                              199
                                                        223
                                191
                                          197
                                                   211
// 227
        229
            233
                 239
                           251
                                257
                                     263
                                          269
                                                   277
                                                        281
                      241
                                               271
// 283
            307
                 311
                           317
                                331
                                     337
                                                   353
        293
                      313
                                          347
                                               349
                                                        359
// 367
        373
            379
                 383
                      389
                           397
                                401
                                     409
                                          419
                                               421
                                                   431
                                                        433
// 439
        443
            449
                 457
                      461
                           463
                                467
                                     479
                                          487
                                               491
                                                   499
                                                        503
// 509
        521
            523
                 541
                      547
                           557
                                563
                                     569
                                          571
                                              577
                                                   587
                                                        593
// 599
        601
            607
                 613
                      617
                           619
                                631
                                     641
                                          643
                                              647
                                                        659
                                                   653
// 661
        673
            677
                 683
                      691
                           701
                                709
                                     719
                                          727
                                              733
                                                   739
                                                        743
                      773
                           787
                                797
                                               821
                                                        827
// 751
        757
            761
                 769
                                     809
                                          811
                                                   823
// 829
        839
            853
                 857
                      859
                           863
                                     881
                                                   907
                                                        911
                                877
                                          883
                                               887
// 919
       929
            937
                 941
                      947
                           953
                                967
                                     971
                                          977
                                                   991
                                                        997
// Other primes:
// The largest prime smaller than 10^x:
// 7 97 997 9973 99991 999983 9999991 99999989 999999937 9999999967
```

6.6 Reduced Row Echelon Form

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxn matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
        returns rank of a[][]
typedef vector<double> VD;
typedef vector<VD> VVD;
const double EPSILON = 1e-7;
// returns rank
int rref (VVD &a){
   int i,j,r,c;
   int n = a.size();
   int m = a[0].size();
   for (r=c=0:c<m:c++){</pre>
       j=r;
       for (i=r+1;i< n;i++) if (fabs(a[i][c])>fabs(a[j][c])) j = i;
       if (fabs(a[j][c]) < EPSILON) continue;</pre>
       for (i=0;i<m;i++) swap(a[j][i],a[r][i]);</pre>
       double s = a[r][c]:
       for (j=0;j<m;j++) a[r][j] /= s;</pre>
       for (i=0;i<n;i++) if (i != r){</pre>
           double t = a[i][c];
```

6.7 Simplex

```
// Two-phase simplex algorithm for solving linear programs of the form
//
    maximize c^T x
    subject to Ax <= b
//
          x >= 0
//
// INPUT: A -- an m x n matrix
     b -- an m-dimensional vector
     c -- an n-dimensional vector
     x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
11
      above, nan if infeasible)
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
   int m, n;
   VI B, N;
   VVD D:
   LPSolver(const VVD &A, const VD &b, const VD &c):
       m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
           for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
               D[i][j] = A[i][j];
           for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1;
               D[i][n + 1] = b[i]; }
           for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
           N[n] = -1; D[m + 1][n] = 1;
   void Pivot(int r, int s) {
       for (int i = 0; i < m + 2; i++) if (i != r)
           for (int j = 0; j < n + 2; j++) if (j != s)
              D[i][j] -= D[r][j] * D[i][s] / D[r][s];
       for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] /= D[r][s];
       for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] /= -D[r][s];
       D[r][s] = 1.0 / D[r][s];
       swap(B[r], N[s]);
   bool Simplex(int phase) {
       int x = phase == 1 ? m + 1 : m;
       while (true) {
           int s = -1;
```

```
for (int j = 0; j <= n; j++) {
              if (phase == 2 && N[j] == -1) continue;
              if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D[x][s] &&
                   N[i] < N[s]) s = i;
          }
           if (D[x][s] > -EPS) return true;
           int r = -1;
           for (int i = 0; i < m; i++) {</pre>
              if (D[i][s] < EPS) continue;</pre>
              if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] /
                   D[r][s] ||
                      (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s])
                          && B[i] < B[r]) r = i;
           if (r == -1) return false;
           Pivot(r, s);
   DOUBLE Solve(VD &x) {
       int r = 0:
       for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
       if (D[r][n + 1] < -EPS) {
           Pivot(r, n);
           if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return
               -numeric_limits<DOUBLE>::infinity();
           for (int i = 0: i < m: i++) if (B[i] == -1) {
              int s = -1;
              for (int j = 0; j \le n; j++)
                  if (s == -1 | D[i][i] < D[i][s] | D[i][j] == D[i][s]
                       && N[i] < N[s]) s = i;
              Pivot(i, s);
           }
       if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
       x = VD(n):
       for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
       return D[m][n + 1];
   }
};
int main() {
   const int m = 4;
   const int n = 3;
   DOUBLE A[m][n] = {
       \{6, -1, 0\},\
       \{-1, -5, 0\},\
       { 1, 5, 1 },
       \{-1, -5, -1\}
   DOUBLE _b[m] = \{ 10, -4, 5, -5 \};
   DOUBLE _c[n] = \{ 1, -1, 0 \};
   VVD A(m);
   VD b(_b, _b + m);
   VD c(_c, _c + n);
   for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);
   LPSolver solver(A, b, c);
   VD x;
   DOUBLE value = solver.Solve(x):
```

```
cerr << "VALUE: " << value << endl; // VALUE: 1.29032
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];
cerr << endl;
return 0;
}</pre>
```

7 7 String

7.1 Aho-Corasick

```
namespace aho_corasick {
   const int SIGMA = 2;
   const int TOTL = 1e7 + 100;
   struct node {
       int link[SIGMA];
       int suff, dict, patt;
       node() {
          suff = 0, dict = 0, patt = -1;
          memset(link, 0, sizeof(link));
       // link[]: contains trie links + failure links
       // suff: link to longest proper suffix that exists in the trie
       // dict: link to longest suffix that exists in the dictionary
       // patt: index of this node's word in the dictionary
   };
   int tail = 1;
   vector<node> trie(TOTL);
   vector<string> patterns;
   void add_pattern(string &s) {
       int loc = 0;
       for (char c : s) {
          int &nloc = trie[loc].link[c-'a'];
           if (!nloc) nloc = tail++;
          loc = nloc:
       trie[loc].dict = loc;
       trie[loc].patt = patterns.size();
       patterns.push_back(s);
   void calc_links() {
       queue<int> bfs({0});
       while (!bfs.empty()) {
          int loc = bfs.front(); bfs.pop();
          int fail = trie[loc].suff;
          if (!trie[loc].dict)
              trie[loc].dict = trie[fail].dict;
          for (int c = 0; c < SIGMA; c++) {
              int &succ = trie[loc].link[c];
              if (succ) {
                  trie[succ].suff = loc ? trie[fail].link[c] : 0:
                  bfs.push(succ);
              } else succ = trie[fail].link[c];
```

7.2 KMP

```
template<typename T> struct kmp {
   int M:
   vector<T> needle;
   vector<int> succ;
   kmp(vector<T> _needle) {
       needle = _needle;
       M = needle.size();
       succ.resize(M + 1);
       succ[0] = -1, succ[1] = 0;
       int cur = 0;
       for (int i = 2; i <= M; ) {</pre>
           if (needle[i-1] == needle[cur]) succ[i++] = ++cur;
           else if (cur) cur = succ[cur];
           else succ[i++] = 0;
       }
   }
   vector<bool> find(vector<T> &haystack) {
       int N = haystack.size(), i = 0;
       vector<bool> res(N);
       for (int m = 0; m + i < N; ) {</pre>
           if (i < M && needle[i] == haystack[m + i]) {</pre>
               if (i == M - 1) res[m] = true;
              i++:
           } else if (succ[i] != -1) {
              m = m + i - succ[i];
              i = succ[i];
           } else {
              i = 0;
               m++;
           }
       return res;
   }
};
```

7.3 Suffix Arrays

```
// Suffix array construction in O(L \log^2 L) time. Routine for // computing the length of the longest common prefix of any two
```

```
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
       of substring s[i...L-1] in the list of sorted suffixes.
//
       That is, if we take the inverse of the permutation suffix[],
//
       we get the actual suffix array.
struct SuffixArray {
 const int L:
 string s;
 vector<vector<int> > P;
 vector<pair<int,int>,int> > M;
  SuffixArray(const string &s): L(s.length()), s(s), P(1, vector<int>(L,
      0)), M(L) {
   for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
   for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {</pre>
     P.push_back(vector<int>(L, 0));
     for (int i = 0; i < L; i++)</pre>
       M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ?</pre>
           P[level-1][i + skip] : -1000), i);
     sort(M.begin(), M.end());
     for (int i = 0; i < L; i++)
       P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ?
           P[level][M[i-1].second] : i;
 vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and
      s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
   int len = 0;
   if (i == j) return L - i;
   for (int k = P.size() - 1; k \ge 0 && i < L && j < L; k--) {
     if (P[k][i] == P[k][j]) {
       i += 1 << k;
       j += 1 << k;
       len += 1 << k;
   return len;
};
int main() {
 // bobocel is the 0'th suffix
 // obocel is the 5'th suffix
 // bocel is the 1'st suffix
 // ocel is the 6'th suffix
 // cel is the 2'nd suffix
 // el is the 3'rd suffix
 // l is the 4'th suffix
 SuffixArray suffix("bobocel");
 vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
```

```
for (int i = 0; i < v.size(); i++) cout << v[i] << " ";
cout << endl;
cout << suffix.LongestCommonPrefix(0, 2) << endl;
}</pre>
```

8 8 Misc

8.1 IO

```
int main() {
 // Ouput a specific number of digits past the decimal point,
 // in this case 5
 cout.setf(ios::fixed); cout << setprecision(5);</pre>
 cout << 100.0/7.0 << endl;
 cout.unsetf(ios::fixed);
 // Output the decimal point and trailing zeros
 cout.setf(ios::showpoint);
 cout << 100.0 << endl;
 cout.unsetf(ios::showpoint);
 // Output a '+' before positive values
 cout.setf(ios::showpos);
 cout << 100 << " " << -100 << endl;
 cout.unsetf(ios::showpos);
 // Output numerical values in hexadecimal
 cout << hex << 100 << " " << 1000 << " " << 10000 << dec << endl;
```

8.2 Longest Increasing Subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
    INPUT: a vector of integers
// OUTPUT: a vector containing the longest increasing subsequence
typedef vector<int> VI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
 VPII best;
 VI dad(v.size(), -1);
 for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
   PII item = make_pair(v[i], 0);
   VPII::iterator it = lower_bound(best.begin(), best.end(), item);
   item.second = i;
#else
   PII item = make_pair(v[i], i);
   VPII::iterator it = upper_bound(best.begin(), best.end(), item);
   if (it == best.end()) {
     dad[i] = (best.size() == 0 ? -1 : best.back().second);
     best.push_back(item);
```

```
} else {
    dad[i] = dad[it->second];
    *it = item;
}

VI ret;
for (int i = best.back().second; i >= 0; i = dad[i])
    ret.push_back(v[i]);
    reverse(ret.begin(), ret.end());
    return ret;
}
```

8.3 Regular Expressions - Java

```
// Code which demonstrates the use of Java's regular expression libraries.
// This is a solution for
// Loglan: a logical language
// http://acm.uva.es/p/v1/134.html
// In this problem, we are given a regular language, whose rules can be
// inferred directly from the code. For each sentence in the input, we
// determine whether the sentence matches the regular expression or not.
// code consists of (1) building the regular expression (which is fairly
// complex) and (2) using the regex to match sentences.
import java.util.*;
import java.util.regex.*;
public class LogLan {
  public static String BuildRegex (){
   String space = " +";
   String A = "([aeiou])";
   String C = "([a-z\&\&[^aeiou]])";
   String MOD = "(g" + A + ")";
   String BA = (b^{\dagger} + A + )^{\dagger};
   String DA = (d' + A + )';
   String LA = "(1" + A + ")";
   String NAM = "([a-z]*" + C + ")";
   String PREDA = "(" + C + C + A + C + A + "|" + C + A + C + C + A +
   String predstring = "(" + PREDA + "(" + space + PREDA + ")*)";
   String predname = "(" + LA + space + predstring + "|" + NAM + ")";
   String preds = "(" + predstring + "(" + space + A + space +
        predstring + ")*)";
   String predclaim = "(" + predname + space + BA + space + preds + "|"
       + DA + space +
     preds + ")";
   String verbpred = "(" + MOD + space + predstring + ")";
   String statement = "(" + predname + space + verbpred + space +
       predname + "|" +
     predname + space + verbpred + ")";
   String sentence = "(" + statement + "|" + predclaim + ")";
```

```
return "^" + sentence + "$";
}
public static void main (String args[]){
  String regex = BuildRegex();
  Pattern pattern = Pattern.compile (regex);
  Scanner s = new Scanner(System.in);
  while (true) {
    // In this problem, each sentence consists of multiple lines, where
    // line is terminated by a period. The code below reads lines until
    // encountering a line whose final character is a '.'. Note the use
    // s.length() to get length of string
    // s.charAt() to extract characters from a Java string
    // s.trim() to remove whitespace from the beginning and end of Java
        string
    // Other useful String manipulation methods include
    // s.compareTo(t) < 0 if s < t, lexicographically</pre>
    // s.indexOf("apple") returns index of first occurrence of "apple"
    // s.lastIndexOf("apple") returns index of last occurrence of
        "apple" in s
    // s.replace(c,d) replaces occurrences of character c with d
    // s.startsWith("apple) returns (s.indexOf("apple") == 0)
    // s.toLowerCase() / s.toUpperCase() returns a new lower/uppercased
        string
    //
    // Integer.parseInt(s) converts s to an integer (32-bit)
    // Long.parseLong(s) converts s to a long (64-bit)
    // Double.parseDouble(s) converts s to a double
    String sentence = "";
    while (true){
     sentence = (sentence + " " + s.nextLine()).trim();
     if (sentence.equals("#")) return;
     if (sentence.charAt(sentence.length()-1) == '.') break;
    // now, we remove the period, and match the regular expression
    String removed_period = sentence.substring(0,
        sentence.length()-1).trim();
    if (pattern.matcher (removed_period).find()){
     System.out.println ("Good");
    } else {
     System.out.println ("Bad!");
}
```

8.4 Tokenizer

```
// char tokens
vector<string> &split(const string &s, char delim, vector<string> &elems){
   stringstream ss(s);
   string item;
   while (getline(ss, item, delim)) {
       elems.push_back(item);
   return elems;
vector<string> split(const string &s, char delim) {
   vector<string> elems;
   split(s, delim, elems);
   return elems;
// string tokens
vector<string> split(string str, string token) {
   vector<string> result;
   int next;
   while((next = str.find(token)) != string::npos) {
       result.push_back(str.substr(0, next));
       str = str.substr(next + token.size());
   return result;
int main() {
   string test = "Hello, this is a string that we might, like, like to
       parse! ";
   auto v_space = split(test, ' ');
   auto v_comma = split(test, ',');
   cout << "|"; for(auto s : v_space) cout << s << "|"; cout << endl;</pre>
   cout << "|"; for(auto s : v_comma) cout << s << "|"; cout << endl;</pre>
   // |Hello,|this|is|a|string|that|we|might,|like,|like|to|parse!|
   // |Hello| this is a string that we might | like | like to parse! |
   return 0;
```