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# 1 1 Templates

### 1.1 Start

```
// .vimrc
syn on
set mouse=a sw=4 ts=4 ai si nu wrap
nnoremap;:

// Terminal: comparing generated output to sample output
./my_program < sample.in | diff sample.out -</pre>
```

# 1.2 Template - C++

```
#include<bits/stdc++.h>
using namespace std;
typedef long long ll;
static bool DBG = 1;

ll mod(ll a, ll b) { return ((a%b)+b)%b; }

int main() {
        ios_base::sync_with_stdio(0);
        cout << fixed << setprecision(15);
        int n;
        cin >> n;
        cout << n << endl;
        return 0;
}</pre>
```

# 1.3 Template - Java

```
import java.util.*;
import java.math.*;
import java.io.*;

class modelo {
    static final double EPS = 1.e-10;
    static final boolean DBG = true;

    private static int cmp(double x, double y = 0, double tol = EPS) {
        return (x <= y + tol)? (x + tol < y)? -1 : 0 : 1;
    }

    public static void main(String[] argv) {
        Scanner s = new Scanner(System.in);
    }
}</pre>
```

# 2 2 Data Structures

## 2.1 BIT

```
template<typename T> struct BIT{
  int S;
```

```
vector<T> v;

BIT<T>(int _S){
    S = _S;
    v.resize(S+1);
}

void update(int i, T k){
    for(i++; i<=S; i+=i&-i)
        v[i] = v[i] + k;
}

T read(int i){
    T sum = 0;
    for(i++; i; i-=i&-i)
        sum = sum + v[i];
    return sum;
}

T read(int l, int r){
    return read(r) - read(l-1);
}
};</pre>
```

#### 2.2 KD Tree

```
// A straightforward, but probably sub-optimal KD-tree implmentation
// that's probably good enough for most things (current it's a
// 2D-tree)
//
// - constructs from n points in O(n lg^2 n) time
// - handles nearest-neighbor query in O(lg n) if points are well
// distributed
// - worst case for nearest-neighbor may be linear in pathological
//
//
// Sonny Chan, Stanford University, April 2009
#include <iostream>
#include <vector>
#include <limits>
#include <cstdlib>
using namespace std;
// number type for coordinates, and its maximum value
typedef long long ntype;
const ntype sentry = numeric_limits<ntype>::max();
// point structure for 2D-tree, can be extended to 3D
struct point {
   ntype x, y;
   point(ntype xx = 0, ntype yy = 0) : x(xx), y(yy) {}
};
bool operator==(const point &a, const point &b) {
   return a.x == b.x && a.y == b.y; }
// sorts points on x-coordinate
```

```
bool on_x(const point &a, const point &b) {
   return a.x < b.x; }</pre>
// sorts points on y-coordinate
bool on_y(const point &a, const point &b) {
   return a.y < b.y; }</pre>
// squared distance between points
ntype pdist2(const point &a, const point &b) {
   ntype dx = a.x-b.x, dy = a.y-b.y;
   return dx*dx + dy*dy; }
// bounding box for a set of points
struct bbox {
   ntype x0, x1, y0, y1;
   bbox(): x0(sentry), x1(-sentry), y0(sentry), y1(-sentry) {}
   // computes bounding box from a bunch of points
   void compute(const vector<point> &v) {
       for (int i = 0; i < v.size(); ++i) {</pre>
          x0 = min(x0, v[i].x); x1 = max(x1, v[i].x);
          y0 = min(y0, v[i].y); y1 = max(y1, v[i].y);
   // squared distance between a point and this bbox, 0 if inside
   ntype distance(const point &p) {
       if (p.x < x0) {
          if (p.v < v0)
                            return pdist2(point(x0, y0), p);
          else if (p.y > y1) return pdist2(point(x0, y1), p);
                             return pdist2(point(x0, p.y), p);
       else if (p.x > x1) {
          if (p.y < y0)
                            return pdist2(point(x1, y0), p);
          else if (p.y > y1) return pdist2(point(x1, y1), p);
                            return pdist2(point(x1, p.y), p);
       }
       else {
          if (p.y < y0) return pdist2(point(p.x, y0), p);</pre>
          else if (p.y > y1) return pdist2(point(p.x, y1), p);
                             return 0;
       }
// stores a single node of the kd-tree, either internal or leaf
struct kdnode {
   bool leaf;
                 // true if this is a leaf node (has one point)
   point pt;
                  // the single point of this is a leaf
   bbox bound; // bounding box for set of points in children
   kdnode *first, *second; // two children of this kd-node
   kdnode() : leaf(false), first(0), second(0) {}
   "kdnode() { if (first) delete first; if (second) delete second; }
   // intersect a point with this node (returns squared distance)
   ntype intersect(const point &p) {
```

```
return bound.distance(p);
   }
   // recursively builds a kd-tree from a given cloud of points
   void construct(vector<point> &vp) {
       // compute bounding box for points at this node
       bound.compute(vp);
       // if we're down to one point, then we're a leaf node
       if (vp.size() == 1) {
           leaf = true;
           pt = vp[0];
       } else {
           // split on x if the bbox is wider than high (not best
               heuristic...)
           if (bound.x1-bound.x0 >= bound.y1-bound.y0)
              sort(vp.begin(), vp.end(), on_x);
           // otherwise split on y-coordinate
           else sort(vp.begin(), vp.end(), on_y);
           // divide by taking half the array for each child
           // (not best performance if many duplicates in the middle)
           int half = vp.size()/2;
           vector<point> vl(vp.begin(), vp.begin()+half);
           vector<point> vr(vp.begin()+half, vp.end());
           first = new kdnode(); first->construct(vl);
           second = new kdnode(); second->construct(vr);
   }
}:
// simple kd-tree class to hold the tree and handle queries
struct kdtree {
   kdnode *root:
   // constructs a kd-tree from a points (copied here, as it sorts them)
   kdtree(const vector<point> &vp) {
       vector<point> v(vp.begin(), vp.end());
       root = new kdnode();
       root->construct(v);
    "kdtree() { delete root; }
   // recursive search method returns squared distance to nearest point
   ntype search(kdnode *node, const point &p) {
       if (node->leaf) {
           // commented special case tells a point not to find itself
             if (p == node->pt) return sentry;
              return pdist2(p, node->pt);
       }
       ntype bfirst = node->first->intersect(p);
       ntype bsecond = node->second->intersect(p);
       // choose the side with the closest bounding box to search first
       // (note that the other side is also searched if needed)
       if (bfirst < bsecond) {</pre>
           ntype best = search(node->first, p);
```

```
if (bsecond < best)</pre>
               best = min(best, search(node->second, p));
           return best;
       }
       else {
           ntype best = search(node->second, p);
           if (bfirst < best)</pre>
               best = min(best, search(node->first, p));
           return best;
       }
   }
   // squared distance to the nearest
   ntype nearest(const point &p) {
       return search(root, p);
};
// some basic test code here
int main() {
   // generate some random points for a kd-tree
   vector<point> vp;
   for (int i = 0; i < 100000; ++i) {
       vp.push_back(point(rand()%100000, rand()%100000));
   kdtree tree(vp);
   // query some points
   for (int i = 0; i < 10; ++i) {
       point q(rand()%100000, rand()%100000);
       cout << "Closest squared distance to (" << q.x << ", " << q.y <<</pre>
            << " is " << tree.nearest(q) << endl;</pre>
   }
```

#### 2.3 LCA

```
struct lca {
   int L, N;
   vector<int> depth, size, link;
   1ca(){}
   lca(const vvi &graph, int root = 0) {
       N = graph.size();
       for (L = 0; (1 << L) <= N; L++);
       depth.resize(N);
       size.resize(N);
       link.resize(L*N);
       init(root, root, graph);
   void init(int loc, int par, const vvi &graph) {
       link[loc] = par;
       for (int 1 = 1; 1 < L; 1++)
          link[1*N + loc] = link[(1-1)*N + link[(1-1)*N + loc]];
       for (int nbr : graph[loc]) {
           if (nbr == par) continue;
           depth[nbr] = depth[loc] + 1;
```

```
4
```

```
init(nbr, loc, graph);
           size[loc] += size[nbr];
       size[loc]++;
   int above(int loc, int dist) {
       for (int 1 = 0; 1 < L; 1++)
           if ((dist >> 1)&1)
              loc = link[1*N + loc];
       return loc;
   }
   int find(int u, int v) {
       if (depth[u] > depth[v]) swap(u, v);
       v = above(v, depth[v] - depth[u]);
       if (u == v) return u;
       for (int 1 = L - 1; 1 \ge 0; 1--) {
           if (link[1*N + u] != link[1*N + v])
              u = link[1*N + u], v = link[1*N + v];
       return link[u];
   }
};
```

## 2.4 Lazy Segment Tree

```
template<typename T, typename U> struct seg_tree_lazy {
   int S, H;
   T zero;
   vector<T> value;
   U noop;
   vector<bool> dirty;
   vector<U> prop;
   seg_tree_lazy<T, U>(int _S, T _zero = T(), U _noop = U()) {
       zero = _zero, noop = _noop;
       for (S = 1, H = 1; S < \_S;) S *= 2, H++;
       value.resize(2*S, zero);
       dirty.resize(2*S, false);
       prop.resize(2*S, noop);
   void set_leaves(vector<T> &leaves) {
       copy(leaves.begin(), leaves.end(), value.begin() + S);
       for (int i = S - 1; i > 0; i--)
          value[i] = value[2 * i] + value[2 * i + 1];
   void apply(int i, U &update) {
       value[i] = update(value[i]);
       if(i < S) {
          prop[i] = prop[i] + update;
          dirty[i] = true;
   }
   void rebuild(int i) {
       for (int 1 = i/2; 1; 1 /= 2) {
          T combined = value[2*1] + value[2*1+1]:
          value[1] = prop[1](combined);
   }
```

```
void propagate(int i) {
       for (int h = H; h > 0; h--) {
           int 1 = i \gg h;
           if (dirty[1]) {
              apply(2*1, prop[1]);
              apply(2*l+1, prop[l]);
              prop[1] = noop;
              dirty[1] = false;
       }
   void upd(int i, int j, U update) {
       i += S, i += S;
       propagate(i), propagate(j);
       for (int 1 = i, r = j; 1 \le r; 1 \ne 2, r \ne 2) {
           if((1&1) == 1) apply(1++, update);
           if((r\&1) == 0) apply(r--, update);
       rebuild(i), rebuild(j);
   T query(int i, int j){
       i += S, j += S;
       propagate(i), propagate(j);
       T res_left = zero, res_right = zero;
       for(; i <= j; i /= 2, j /= 2){
           if((i\&1) == 1) res_left = res_left + value[i++];
           if((j&1) == 0) res_right = value[j--] + res_right;
       return res_left + res_right;
};
```

## 2.5 Segment Tree

```
template<typename T> struct seg_tree {
   int S;
   T zero:
   vector<T> value;
   seg_tree<T>(int _S, T _zero = T()) {
       S = _S, zero = _zero;
       value.resize(2*S+1, zero);
   void set_leaves(vector<T> &leaves) {
       copy(leaves.begin(), leaves.end(), value.begin() + S);
       for (int i = S - 1; i > 0; i--)
          value[i] = value[2 * i] + value[2 * i + 1];
   void upd(int i, T v) {
      i += S:
      value[i] = v;
       while(i>1){
          value[i] = value[2*i] + value[2*i+1];
       }
   T query(int i, int j) {
       T res_left = zero, res_right = zero;
```

```
for(i += S, j += S; i <= j; i /= 2, j /= 2){
    if((i&1) == 1) res_left = res_left + value[i++];
    if((j&1) == 0) res_right = value[j--] + res_right;
}
return res_left + res_right;
}
};</pre>
```

#### 2.6 Union Find

```
// (struct) also keeps track of sizes
struct union find {
       vector<int> P,S;
       union find(int N) {
              P.resize(N), S.resize(N, 1);
              for(int i = 0; i < N; i++) P[i] = i;</pre>
       int rep(int i) {return (P[i] == i) ? i : P[i] = rep(P[i]);}
       bool unio(int a, int b) {
              a = rep(a), b = rep(b);
              if(a == b) return false;
              P[b] = a;
              S[a] += S[b];
              return true;
       }
// (Shorter) union-find set: the vector/array contains the parent of each
int find(vector <int>& C, int x){return (C[x]==x) ? x : C[x]=find(C,
    C[x]); //C++
int find(int x){return (C[x]==x)?x:C[x]=find(C[x]);} //C
```

# 3 3 Graph

### 3.1 2-SAT

```
struct two_sat {
   int N;
   vector<vector<int>> impl;

   two_sat(int _N) {
      N = _N;
      impl.resize(2 * N);
   }

   void add_impl(int var1, bool neg1, int var2, bool neg2) {
      impl[2 * var1 + neg1].push_back(2 * var2 + neg2);
      impl[2 * var2 + !neg2].push_back(2 * var1 + !neg1);
   }

   void add_clause(int var1, bool neg1, int var2, bool neg2) {
      add_impl(var1, !neg1, var2, neg2);
   }

   void add_clause(int var1, bool neg1) {
      add_clause(var1, neg1, var1, neg1);
   }
}
```

```
int V, L, C;
   stack<int> view;
   int dfs(int loc) {
       visit[loc] = V:
       label[loc] = L++;
       int low = label[loc];
       view.push(loc);
       in_view[loc] = true;
       for (int nbr : impl[loc]) {
           if(!visit[nbr]) low = min(low, dfs(nbr));
           else if(in_view[nbr]) low = min(low, label[nbr]);
       if(low == label[loc]) {
           while (true) {
              int mem = view.top();
              comp[mem] = C;
              in_view[mem] = false;
              view.pop();
              if(mem == loc) break;
          C++;
       }
       return low;
   vector<int> visit, label, comp, in_view;
   void reset(vector<int> &v) {
       v.resize(2 * N);
       fill(v.begin(), v.end(), 0);
   bool consistent() {
       V = 0, L = 0, C = 0;
       reset(visit), reset(label), reset(comp), reset(in_view);
       for (int i = 0: i < 2 * N: i++) {
           if(!visit[i]) {
              V++;
              dfs(i);
       for (int i = 0; i < N; i++)</pre>
           if(comp[2 * i] == comp[2 * i + 1]) return false;
       return true;
};
```

#### 3.2 Eulerian Path

```
struct Edge;
typedef list<Edge>::iterator iter;

struct Edge {
  int next_vertex;
  iter reverse_edge;
  Edge(int next_vertex) : next_vertex(next_vertex) {}
};
```

```
const int max_vertices = 100;
int num_vertices;
list<Edge> adj[max_vertices];
                                    // adjacency list
vector<int> path;
void find_path(int v) {
       while(adj[v].size() > 0) {
              int vn = adj[v].front().next_vertex;
              adj[vn].erase(adj[v].front().reverse_edge);
              adj[v].pop_front();
              find_path(vn);
       path.push_back(v);
}
void add_edge(int a, int b) {
       adj[a].push_front(Edge(b));
       iter ita = adj[a].begin();
       adj[b].push_front(Edge(a));
       iter itb = adj[b].begin();
       ita->reverse_edge = itb;
       itb->reverse_edge = ita;
```

## 3.3 Heavy Light

```
template<typename T> struct heavy_light {
   lca links;
   seg_tree<T> st;
   vector<int> preorder, index, jump;
   heavy_light(const vvi &graph, int root) {
       links = lca(graph, 0);
       st = seg_tree<T>(graph.size());
       index.resize(graph.size()), jump.resize(graph.size());
       dfs(root, root, graph);
   void dfs(int loc, int par, int lhv, const vvi &graph) {
       jump[loc] = lhv;
       index[loc] = preorder.size();
       preorder.push_back(loc);
       vector<int> ch = graph[loc];
       sort(ch.begin(), ch.end(), [&](int i, int j) {
          return links.size[i] > links.size[j]; });
       if (loc != par) ch.erase(ch.begin());
       for (int c = 0; c < ch.size(); c++)</pre>
          dfs(ch[c], loc, c ? ch[c] : lhv, graph);
   void assign(int loc, T value) {
       st.upd(index[loc], value);
   T __sum(int u, int r) {
       T res;
       while (u != r) {
          int go = max(index[r] + 1, index[jump[u]]);
          res = res + st.query(go, index[u]);
          u = links.link[preorder[go]];
       return res;
```

```
}
T sum(int u, int v) {
    int r = links.find(u, v);
    return st.query(index[r], index[r]) + __sum(u, r) + __sum(v, r);
}
};
```

#### 3.4 Poset Width

```
// requires bipartite graph (4.1)
vector<int> width(vector<vector<int>> poset) {
   int N = poset.size();
   bipartite_graph g(N, N);
   for (int i = 0; i < N; i++) {
       for (int j : poset[i])
          g.edge(j, i);
   g.matching();
   vector<bool> vis[2];
   vis[false].resize(2 * N, false);
   vis[true].resize(2 * N, false);
   for (int i = 0; i < N; i++) {</pre>
       if (g.match[i] != -1) continue;
       if (vis[false][i]) continue;
       queue<pair<bool, int>> bfs;
       bfs.push(make_pair(false, i));
       vis[false][i] = true;
       while (!bfs.empty()) {
           bool inm = bfs.front().first;
           int loc = bfs.front().second;
           bfs.pop();
           for (int nbr : g.adj[loc]) {
              if (vis[!inm][nbr]) continue;
              if ((g.match[loc] == nbr) ^ inm) continue;
              vis[!inm][nbr] = true;
              bfs.push(make_pair(!inm, nbr));
           }
       }
   vector<bool> inz(2 * N, false);
   for (int i = 0; i < 2 * N; i++)
       inz[i] = vis[true][i] || vis[false][i];
   vector<bool> ink(N, false);
   for (int i = 0; i < N; i++)</pre>
       if (!inz[i])
           ink[i] = true;
   for (int i = N; i < 2 * N; i++)</pre>
       if (inz[i])
          ink[i - N] = true:
   vector<int> res;
   for (int i = 0; i < N; i++) {</pre>
       if (!ink[i])
          res.push_back(i);
   return res;
```

#### 3.5 SCC

```
struct edge{int e, nxt;};
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x) {
 int i;
 for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
 stk[++stk[0]]=x;
void fill_backward(int x) {
 int i;
 v[x]=false;
 group_num[x]=group_cnt;
 for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e]) fill_backward(er[i].e);
void add_edge(int v1, int v2) { //add edge v1->v2
 e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
 er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC() {
 int i:
 stk[0]=0:
 memset(v, false, sizeof(v));
 for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);</pre>
 group_cnt=0;
 for(i=stk[0];i>=1;i--) if(v[stk[i]]){group_cnt++;
      fill_backward(stk[i]);}
```

## 3.6 Topological Sort

```
// This function uses performs a non-recursive topological sort.
//
// Running time: O(|V|^2). If you use adjacency lists (vector<map<int> >),
               the running time is reduced to O(|E|).
    INPUT: w[i][j] = 1 if i should come before j, 0 otherwise
    OUTPUT: a permutation of 0,...,n-1 (stored in a vector)
//
            which represents an ordering of the nodes which
//
            is consistent with w
// If no ordering is possible, false is returned.
typedef double TYPE;
typedef vector<TYPE> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool TopologicalSort (const VVI &w, VI &order){
 int n = w.size();
 VI parents (n);
```

```
queue<int> q;
order.clear();
for (int i = 0; i < n; i++){
   for (int j = 0; j < n; j++)
      if (w[j][i]) parents[i]++;
   if (parents[i] == 0) q.push (i);
}
while (q.size() > 0){
   int i = q.front();
   q.pop();
   order.push_back (i);
   for (int j = 0; j < n; j++) if (w[i][j]){
      parents[j]--;
      if (parents[j] == 0) q.push (j);
   }
}
return (order.size() == n);
}</pre>
```

# 4 4 Combinatorial Optimization

## 4.1 Bipartite Graph

```
struct bipartite_graph {
   int A, B;
   vector<vector<int>> adj;
   bipartite_graph(int _A, int _B) {
       A = A, B = B;
       adj.resize(A + B);
   void edge(int i, int j) {
       adj[i].push_back(A+j);
       adj[A+j].push_back(i);
   vector<int> visit, match;
   bool augment(int loc, int run) {
       if(visit[loc] == run) return false;
       visit[loc] = run;
       for (int nbr : adj[loc]) {
          if (match[nbr] == -1 || augment(match[nbr], run)) {
              match[loc] = nbr, match[nbr] = loc;
              return true:
       }
       return false;
   int matching() {
       visit = vector<int>(A+B, -1);
       match = vector<int>(A+B, -1);
       int ans = 0;
       for (int i = 0; i < A; i++)
          ans += augment(i, i);
       return ans;
```

```
vector<bool> vertex_cover() {
       vector<bool> res(A + B, false);
       queue<int> bfs;
       for (int i = 0; i < A; i++) {
           if (match[i] == -1) bfs.push(i);
           else res[i] = true:
       while (!bfs.empty()) {
           int loc = bfs.front();
          bfs.pop();
           for (int nbr : adj[loc]) {
              if (res[nbr]) continue;
              res[nbr] = true;
              int loc2 = match[nbr];
              if (loc2 == -1) continue;
              res[loc2] = false;
              bfs.push(loc2);
       }
       return res;
   }
};
```

### 4.2 Max Flow - Dinic

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
//
// Running time:
     O(|V|^2 |E|)
//
// INPUT:
      - graph, constructed using AddEdge()
//
      - source and sink
// OUTPUT:
      - maximum flow value
      - To obtain actual flow values, look at edges with capacity > 0
        (zero capacity edges are residual edges).
struct Edge {
 int from, to, cap, flow, index;
 Edge(int from, int to, int cap, int flow, int index) :
   from(from), to(to), cap(cap), flow(flow), index(index) {}
 11 rcap() { return cap - flow; }
};
struct Dinic {
 int N;
 vector<vector<Edge> > G;
 vector<vector<Edge *> > Lf;
 vector<int> laver;
 vector<int> Q;
 Dinic(int N) : N(N), G(N), Q(N) {}
  void AddEdge(int from, int to, int cap) {
   if (from == to) return;
```

```
G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
 G[to].push_back(Edge(to, from, 0, 0, G[from].size() - 1));
11 BlockingFlow(int s, int t) {
 layer.clear(); layer.resize(N, -1);
 laver[s] = 0:
 Lf.clear(); Lf.resize(N);
 int head = 0, tail = 0;
 Q[tail++] = s;
 while (head < tail) {</pre>
   int x = Q[head++];
   for (int i = 0; i < G[x].size(); i++) {</pre>
     Edge &e = G[x][i]; if (e.rcap() <= 0) continue;</pre>
     if (layer[e.to] == -1) {
       layer[e.to] = layer[e.from] + 1;
       Q[tail++] = e.to;
     if (layer[e.to] > layer[e.from]) {
       Lf[e.from].push_back(&e);
 if (layer[t] == -1) return 0;
 11 \text{ totflow} = 0;
 vector<Edge *> P;
 while (!Lf[s].empty()) {
   int curr = P.empty() ? s : P.back()->to;
   if (curr == t) { // Augment
     11 amt = P.front()->rcap();
     for (int i = 0; i < P.size(); ++i) {</pre>
       amt = min(amt, P[i]->rcap());
     totflow += amt;
     for (int i = P.size() - 1; i >= 0; --i) {
       P[i]->flow += amt:
       G[P[i]->to][P[i]->index].flow -= amt;
       if (P[i]->rcap() <= 0) {</pre>
         Lf[P[i]->from].pop_back();
         P.resize(i);
       }
   } else if (Lf[curr].empty()) { // Retreat
     P.pop_back();
     for (int i = 0; i < N; ++i)
       for (int j = 0; j < Lf[i].size(); ++j)</pre>
         if (Lf[i][j]->to == curr)
           Lf[i].erase(Lf[i].begin() + j);
   } else { // Advance
     P.push_back(Lf[curr].back());
 return totflow;
11 GetMaxFlow(int s, int t) {
 11 \text{ totflow} = 0;
 while (ll flow = BlockingFlow(s, t))
   totflow += flow;
 return totflow;
```

};

## 4.3 Min Cost Matching

```
// Min cost bipartite matching via shortest augmenting paths
// This is an O(n^3) implementation of a shortest augmenting path
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
// second.
   cost[i][j] = cost for pairing left node i with right node j
   Lmate[i] = index of right node that left node i pairs with
   Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
 int n = int(cost.size());
 // construct dual feasible solution
 VD u(n);
 VD v(n);
 for (int i = 0; i < n; i++) {</pre>
   u[i] = cost[i][0];
   for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
 for (int j = 0; j < n; j++) {
   v[i] = cost[0][i] - u[0];
   for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);</pre>
 // construct primal solution satisfying complementary slackness
 Lmate = VI(n, -1);
 Rmate = VI(n, -1);
 int mated = 0;
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++) {
     if (Rmate[j] != -1) continue;
     if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
      Lmate[i] = j;
      Rmate[j] = i;
      mated++;
      break;
 VD dist(n):
 VI dad(n);
 VI seen(n);
 // repeat until primal solution is feasible
```

```
while (mated < n) {</pre>
 // find an unmatched left node
 int s = 0:
 while (Lmate[s] != -1) s++;
 // initialize Dijkstra
 fill(dad.begin(), dad.end(), -1);
 fill(seen.begin(), seen.end(), 0);
 for (int k = 0; k < n; k++)
   dist[k] = cost[s][k] - u[s] - v[k];
 int j = 0;
 while (true) {
   // find closest
   j = -1;
   for (int k = 0; k < n; k++) {
             if (seen[k]) continue;
             if (j == -1 || dist[k] < dist[j]) j = k;</pre>
   seen[j] = 1;
   // termination condition
   if (Rmate[j] == -1) break;
   // relax neighbors
   const int i = Rmate[j];
   for (int k = 0; k < n; k++) {
             if (seen[k]) continue;
             const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
             if (dist[k] > new_dist) {
              dist[k] = new_dist;
               dad[k] = j;
 // update dual variables
 for (int k = 0; k < n; k++) {
   if (k == j || !seen[k]) continue;
   const int i = Rmate[k];
   v[k] += dist[k] - dist[j];
   u[i] -= dist[k] - dist[j];
 u[s] += dist[j];
 // augment along path
 while (dad[j] >= 0) {
   const int d = dad[j];
   Rmate[j] = Rmate[d];
   Lmate[Rmate[j]] = j;
   j = d;
 Rmate[j] = s;
 Lmate[s] = j;
 mated++;
double value = 0;
for (int i = 0; i < n; i++)</pre>
 value += cost[i][Lmate[i]];
return value;
```

#### 4.4 Min Cost Max Flow

```
// Min cost max flow algorithm using an adjacency matrix. If you
// want just regular max flow, setting all edge costs to 1 gives
// running time O(|E|^2 |V|).
// Running time: O(\min(|V|^2 * totflow, |V|^3 * totcost))
//
// INPUT: cap -- a matrix such that cap[i][j] is the capacity of
//
               a directed edge from node i to node j
//
         cost -- a matrix such that cost[i][j] is the (positive)
//
                cost of sending one unit of flow along a
                directed edge from node i to node j
         source -- starting node
         sink -- ending node
//
//
// OUTPUT: max flow and min cost; the matrix flow will contain
          the actual flow values (note that unlike in the MaxFlow
//
          code, you don't need to ignore negative flow values -- there
//
          shouldn't be any)
typedef vector<11> v11;
typedef vector<vll> vvll;
const 11 INF = 1LL << 60;</pre>
struct MCMF {
       int N:
       vll found, dad, dist, pi;
       vvll cap, flow, cost;
       MCMF(int N) : N(N), cap(N, vll(N)), flow(cap), cost(cap),
                                     dad(N), found(N), pi(N), dist(N+1) {};
   void add_edge(int from, int to, ll ca, ll co) {
              cap[from][to] = ca; cost[from][to] = co; }
   bool search(int source, int sink) {
              fill(found.begin(), found.end(), 0);
              fill(dist.begin(), dist.end(), INF);
              dist[source] = 0;
              while(source != N) {
                      int best = N:
                      found[source] = 1;
                      for(int k = 0; k < N; k++) {
                             if(found[k]) continue;
                             if(flow[k][source]) {
                                    11 val = dist[source] + pi[source] -
                                         pi[k] - cost[k][source];
                                     if(dist[k] > val) {
                                            dist[k] = val;
                                            dad[k] = source;
                             if(flow[source][k] < cap[source][k]) {</pre>
                                     11 val = dist[source] + pi[source] -
                                         pi[k] + cost[source][k];
                                     if(dist[k] > val) {
                                            dist[k] = val;
                                            dad[k] = source;
                             if(dist[k] < dist[best]) best = k;</pre>
```

```
source = best;
              for(int k = 0; k < N; k++)
                     pi[k] = min((11)(pi[k] + dist[k]), INF);
              return found[sink]:
   pair<11,11> mcmf(int source, int sink) {
              11 totflow = 0, totcost = 0;
              while(search(source, sink)) {
                     11 amt = INF;
                     for(int x = sink; x != source; x = dad[x])
                             amt = min(amt, (ll)(flow[x][dad[x]] != 0 ?
                                    flow[x][dad[x]] : cap[dad[x]][x] -
                                        flow[dad[x]][x]));
                     for(int x = sink; x != source; x = dad[x]) {
                             if(flow[x][dad[x]] != 0) {
                                flow[x][dad[x]] -= amt;
                                 totcost -= amt * cost[x][dad[x]];
                             } else {
                                flow[dad[x]][x] += amt;
                                totcost += amt * cost[dad[x]][x];
                     totflow += amt;
              return {totflow, totcost};
   }
};
```

#### 4.5 Min Cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
//
// Running time:
      O(|V|^3)
//
11
// INPUT:
      graph, constructed using AddEdge()
// OUTPUT:
// - (min cut value, nodes in half of min cut)
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
 int N = weights.size();
 VI used(N), cut, best_cut;
 int best_weight = -1;
 for (int phase = N-1; phase >= 0; phase--) {
   VI w = weights[0];
   VI added = used;
   int prev, last = 0;
   for (int i = 0; i < phase; i++) {</pre>
     prev = last;
     last = -1;
     for (int j = 1; j < N; j++)
```

```
if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
if (i == phase-1) {
    for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];
    for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];
    used[last] = true;
    cut.push_back(last);
    if (best_weight == -1 || w[last] < best_weight) {
        best_cut = cut;
        best_weight = w[last];
    }
} else {
    for (int j = 0; j < N; j++)
        w[j] += weights[last][j];
        added[last] = true;
}
}
return make_pair(best_weight, best_cut);</pre>
```

# 5 5 Geometry

### 5.1 Convex Hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if
    REMOVE REDUNDANT is
// #defined.
// Running time: O(n log n)
    INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise,
    starting
            with bottommost/leftmost point
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
 Тх, у;
 PT() {}
 PT(T x, T y) : x(x), y(y) {}
 bool operator<(const PT &rhs) const { return make_pair(y,x) <</pre>
      make_pair(rhs.v,rhs.x); }
 bool operator==(const PT &rhs) const { return make_pair(y,x) ==
      make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE_REDUNDANT
bool between(const PT &a, const PT &b, const PT &c) {
 return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x) <= 0 &&
      (a.y-b.y)*(c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
```

```
sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
   while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >=
        0) up.pop_back();
   while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <=</pre>
        0) dn.pop_back();
   up.push_back(pts[i]);
   dn.push_back(pts[i]);
 for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
 if (pts.size() <= 2) return;</pre>
 dn.clear();
 dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
   if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
   dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
   dn[0] = dn.back();
   dn.pop_back();
 pts = dn;
#endif
```

## 5.2 Delaunay

```
// Slow but simple Delaunay triangulation. Does not handle
// degenerate cases (from O'Rourke, Computational Geometry in C)
//
// Running time: O(n^4)
//
// INPUT: x[] = x-coordinates
//
           y[] = y-coordinates
// OUTPUT: triples = a vector containing m triples of indices
                     corresponding to triangle vertices
typedef double T;
struct triple {
   int i, j, k;
   triple() {}
   triple(int i, int j, int k) : i(i), j(j), k(k) {}
vector<triple> delaunayTriangulation(vector<T>& x, vector<T>& y) {
       int n = x.size();
       vector<T> z(n);
       vector<triple> ret;
       for (int i = 0; i < n; i++)</pre>
          z[i] = x[i] * x[i] + y[i] * y[i];
       for (int i = 0; i < n-2; i++) {
          for (int j = i+1; j < n; j++) {
                     for (int k = i+1; k < n; k++) {
```

```
if (j == k) continue;
                          double xn = (y[j]-y[i])*(z[k]-z[i]) -
                              (y[k]-y[i])*(z[j]-z[i]);
                          double yn = (x[k]-x[i])*(z[j]-z[i]) -
                              (x[j]-x[i])*(z[k]-z[i]);
                          double zn = (x[j]-x[i])*(y[k]-y[i]) -
                              (x[k]-x[i])*(y[j]-y[i]);
                          bool flag = zn < 0;</pre>
                          for (int m = 0; flag && m < n; m++)</pre>
                              flag = flag && ((x[m]-x[i])*xn +
                                             (y[m]-y[i])*yn +
                                             (z[m]-z[i])*zn <= 0);
                          if (flag) ret.push_back(triple(i, j, k));
          }
       }
       return ret;
int main() {
   T xs[]={0, 0, 1, 0.9};
   T ys[]={0, 1, 0, 0.9};
   vectorT > x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
   vector<triple> tri = delaunayTriangulation(x, y);
   //expected: 0 1 3
   //
              0 3 2
   int i;
   for(i = 0; i < tri.size(); i++)</pre>
       printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
   return 0;
```

## 5.3 Geometry

```
// C++ routines for computational geometry.
double INF = 1e100:
double EPS = 1e-12;
struct PT {
 double x, y;
 PT() {}
 PT(double x, double y) : x(x), y(y) {}
 PT(const PT \&p) : x(p.x), y(p.y) {}
 PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
 PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
 PT operator * (double c) const { return PT(x*c, y*c); }
 PT operator / (double c) const { return PT(x/c, y/c); }
};
double dot(PT p, PT q) { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream & os, const PT & p) {
 os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y,p.x); }
PT RotateCW90(PT p) { return PT(p.y,-p.x); }
```

```
PT RotateCCW(PT p, double t) {
 return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
 return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
 double r = dot(b-a,b-a);
 if (fabs(r) < EPS) return a;</pre>
 r = dot(c-a, b-a)/r;
 if (r < 0) return a;
 if (r > 1) return b;
 return a + (b-a)*r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                        double a, double b, double c, double d) {
 return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
 return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
 return LinesParallel(a, b, c, d)
     && fabs(cross(a-b, a-c)) < EPS
     && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
 if (LinesCollinear(a, b, c, d)) {
   if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
     dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
   if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-b, d-b) > 0)
     return false;
   return true;
 if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
 if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
 return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
```

```
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
   b=b-a; d=c-d; c=c-a;
   assert(dot(b, b) > EPS && dot(d, d) > EPS);
   return a + b*cross(c, d)/cross(b, d):
}
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
   b=(a+b)/2;
    c=(a+c)/2;
   return ComputeLineIntersection(b, b+RotateCW90(a-b), c,
             c+RotateCW90(a-c));
}
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
   bool c = 0;
   for (int i = 0; i < p.size(); i++){</pre>
        int j = (i+1)%p.size();
        if ((p[i].y <= q.y && q.y < p[j].y ||
           p[j].y \le q.y && q.y < p[i].y) &&
           q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y - p[i].y) / (p[j].y - p[i].y - p[i].y) / (p[j].y - p[i]
                    p[i].y))
           c = !c;
   }
   return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
   for (int i = 0: i < p.size(): i++)</pre>
        if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
           return true;
        return false;
}
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
    vector<PT> ret;
   b = b-a;
    a = a-c;
    double A = dot(b, b);
    double B = dot(a, b);
    double C = dot(a, a) - r*r;
    double D = B*B - A*C;
    if (D < -EPS) return ret;</pre>
    ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
   if (D > EPS)
        ret.push_back(c+a+b*(-B-sqrt(D))/A);
    return ret;
```

```
// compute intersection of circle centered at a with radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
 vector<PT> ret;
 double d = sqrt(dist2(a, b));
 if (d > r+R \mid d+min(r, R) < max(r, R)) return ret;
 double x = (d*d-R*R+r*r)/(2*d);
 double y = sqrt(r*r-x*x);
 PT v = (b-a)/d;
 ret.push_back(a+v*x + RotateCCW90(v)*y);
 if (y > 0)
   ret.push_back(a+v*x - RotateCCW90(v)*y);
 return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
 double area = 0:
 for(int i = 0; i < p.size(); i++) {</pre>
   int j = (i+1) % p.size();
   area += p[i].x*p[j].y - p[j].x*p[i].y;
 return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
 PT c(0.0):
 double scale = 6.0 * ComputeSignedArea(p);
 for (int i = 0; i < p.size(); i++){</pre>
   int j = (i+1) % p.size();
   c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
 return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
 for (int i = 0; i < p.size(); i++) {</pre>
   for (int k = i+1; k < p.size(); k++) {</pre>
     int j = (i+1) % p.size();
     int 1 = (k+1) % p.size();
     if (i == 1 || j == k) continue;
     if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
       return false;
 return true;
// Plane distance between parallel planes aX + bY + cZ + d1 = 0 and
```

```
// aX + bY + cZ + d2 = 0 is abs(d1 - d2) / sqrt(a*a + b*b + c*c)
// distance from point (px, py, pz) to line (x1, y1, z1)-(x2, y2, z2)
// (or ray, or segment; in the case of the ray, the endpoint is the
// first point)
public static final int LINE = 0;
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1, double y1, double z1,
   double x2, double y2, double z2, double px, double py, double pz,
   int type) {
  double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-y2) + (z1-z2)*(z1-z2);
 double x, y, z;
  if (pd2 == 0) {
   x = x1;
   y = y1;
   z = z1;
  else {
   double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (pz-z1)*(z2-z1)) /
        pd2;
   x = x1 + u * (x2 - x1);
   y = y1 + u * (y2 - y1);
   z = z1 + u * (z2 - z1);
   if (type != LINE && u < 0) {</pre>
     x = x1;
     y = y1;
     z = z1;
   if (type == SEGMENT && u > 1.0) {
     x = x2;
     y = y2;
     z = z2;
 return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)*(z-pz);
int main() {
 // expected: (-5,2)
 cerr << RotateCCW90(PT(2,5)) << endl;</pre>
 // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
  // expected: (5,2) (7.5,3) (2.5,1)
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
      << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "</pre>
      << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr \lt\lt DistancePointPlane(4,-4,3,2,-2,5,-8) \lt\lt endl;
  // expected: 1 0 1
  cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
      << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "</pre>
      << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
  // expected: 0 0 1
  cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
      << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
      << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;</pre>
  // expected: 1 1 1 0
```

```
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "</pre>
    << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
    << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
    << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) <</pre>
    endl;
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;</pre>
vector<PT> v;
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "
    << PointInPolygon(v, PT(2,0)) << " "</pre>
    << PointInPolygon(v, PT(0,2)) << " "</pre>
    << PointInPolygon(v, PT(5,2)) << " "</pre>
    << PointInPolygon(v, PT(2,5)) << endl;</pre>
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
    << PointOnPolygon(v, PT(2,0)) << " "
    << PointOnPolygon(v, PT(0,2)) << " "
    << PointOnPolygon(v, PT(5,2)) << " "</pre>
    << PointOnPolygon(v, PT(2,5)) << endl;</pre>
// expected: (1,6)
            (5,4)(4,5)
//
//
            blank line
11
            (4,5) (5,4)
11
            blank line
//
            (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = \{ PT(0,0), PT(5,0), PT(1,1), PT(0,5) \};
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;</pre>
cerr << "Centroid: " << c << endl;</pre>
```

## 6 6 Numerics

### 6.1 Euclid

```
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the
// algorithms described here work on nonnegative integers.
typedef vector<int> VI;
typedef pair<int,int> PII;
// return a % b (positive value)
int mod(int a, int b) {
 return ((a%b)+b)%b;
// computes gcd(a,b)
int gcd(int a, int b) {
 while(b){a%=b; tmp=a; a=b; b=tmp;}
 return a;
// computes lcm(a,b)
int lcm(int a, int b) {
 return a/gcd(a,b)*b;
// returns d = gcd(a,b); finds x,y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
 int xx = y = 0;
 int yy = x = 1;
 while (b) {
   int q = a/b;
   int t = b; b = a%b; a = t;
   t = xx; xx = x-q*xx; x = t;
   t = yy; yy = y-q*yy; y = t;
 return a;
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
 int x, y;
 VI solutions;
 int d = extended_euclid(a, n, x, y);
 if (!(b%d)) {
   x = mod (x*(b/d), n);
   for (int i = 0; i < d; i++)</pre>
     solutions.push_back(mod(x + i*(n/d), n));
 return solutions;
// computes b such that ab = 1 (mod n), returns -1 on failure
int mod_inverse(int a, int n) {
 int x, y;
 int d = extended_euclid(a, n, x, y);
 if (d > 1) return -1;
 return mod(x,n);
// Chinese remainder theorem (special case): find z such that
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y).
// Return (z,M). On failure, M = -1.
PII chinese_remainder_theorem(int x, int a, int y, int b) {
 int s. t:
 int d = extended_euclid(x, y, s, t);
 if (a%d != b%d) return make_pair(0, -1);
```

```
return make_pair(mod(s*b*x+t*a*y,x*y)/d, x*y/d);
// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm_i (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &x, const VI &a) {
 PII ret = make_pair(a[0], x[0]);
 for (int i = 1; i < x.size(); i++) {</pre>
   ret = chinese_remainder_theorem(ret.second, ret.first, x[i], a[i]);
   if (ret.second == -1) break;
 return ret;
// computes x and y such that ax + by = c; on failure, x = y = -1
void linear_diophantine(int a, int b, int c, int &x, int &y) {
 int d = gcd(a,b);
 if (c%d) {
   x = y = -1;
 } else {
   x = c/d * mod_inverse(a/d, b/d);
   y = (c-a*x)/b;
int main() {
 // expected: 2
 cout << gcd(14, 30) << endl;
 // expected: 2 -2 1
 int x, y;
 int d = extended_euclid(14, 30, x, y);
  cout << d << " " << x << " " << y << endl;
  // expected: 95 45
 VI sols = modular linear equation solver(14. 30. 100):
 for (int i = 0; i < (int) sols.size(); i++) cout << sols[i] << " ";</pre>
  cout << endl;</pre>
 // expected: 8
 cout << mod_inverse(8, 9) << endl;</pre>
  // expected: 23 56
             11 12
 int xs[] = {3, 5, 7, 4, 6};
 int as[] = \{2, 3, 2, 3, 5\};
 PII ret = chinese_remainder_theorem(VI (xs, xs+3), VI(as, as+3));
  cout << ret.first << " " << ret.second << endl;</pre>
 ret = chinese_remainder_theorem (VI(xs+3, xs+5), VI(as+3, as+5));
  cout << ret.first << " " << ret.second << endl;</pre>
  // expected: 5 -15
 linear_diophantine(7, 2, 5, x, y);
  cout << x << " " << y << endl;
```

#### 6.2 FFT

```
namespace fft {
    struct cnum {
        double a, b;
        cnum operator+(const cnum &c) { return { a + c.a, b + c.b }; }
```

```
cnum operator-(const cnum &c) { return { a - c.a, b - c.b }; }
       cnum operator*(const cnum &c) { return { a*c.a - b*c.b, a*c.b +
            b*c.a }: }
       cnum operator/(double d) { return { a / d, b / d }; }
   };
   const double PI = 2 * atan2(1, 0);
   int deg;
   vector<int> rev;
   void set_degree(int _deg) {
       assert(__builtin_popcount(_deg) == 1);
       deg = _deg;
       rev.resize(deg);
       for (int i = 1, j = 0; i < deg; i++) {
           int bit = deg / 2;
           for (; j >= bit; bit /= 2)
              j -= bit;
           j += bit;
           rev[i] = j;
   }
   void transform(vector<cnum> &poly, bool invert) {
       if(deg != poly.size()) set_degree(poly.size());
       for (int i = 1; i < deg; i++)</pre>
           if(rev[i] > i)
               swap(poly[i], poly[rev[i]]);
       for (int len = 2; len <= deg; len *= 2) {
           double ang = 2 * PI / len * (invert ? -1 : 1);
           cnum base = { cos(ang), sin(ang) };
           for (int i = 0; i < deg; i += len) {</pre>
               cnum w = \{1, 0\};
               for (int j = 0; j < len / 2; j++) {
                  cnum u = poly[i+j];
                  cnum v = w * poly[i+j+len/2];
                  polv[i+j] = u + v;
                  poly[i+j+len/2] = u - v;
                  w = w * base:
           }
       if(invert) {
           for (int i = 0; i < deg; i++)</pre>
               poly[i] = poly[i] / double(deg);
   }
};
```

### 6.3 Gauss-Jordan

```
// Gauss-Jordan elimination with full pivoting.
//
// Uses:
// (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
// (3) computing determinants of square matrices
//
// Running time: O(n^3)
```

```
// INPUT: a[][] = an nxn matrix
//
            b[][] = an nxm matrix
//
// OUTPUT: X
                  = an nxm matrix (stored in b[][])
//
            A^{-1} = an nxn matrix (stored in a[][])
            returns determinant of a[][]
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
const double EPS = 1e-10;
T GaussJordan(VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
 T \det = 1:
  for (int i = 0; i < n; i++) {</pre>
   int pj = -1, pk = -1;
   for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
     for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
       if (p_j == -1 \mid | fabs(a[j][k]) > fabs(a[p_j][p_k])) { p_j = j; p_k = k;}
   if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl;</pre>
        exit(0): }
   ipiv[pk]++;
   swap(a[pj], a[pk]);
   swap(b[pj], b[pk]);
   if (pj != pk) det *= -1;
   irow[i] = pj;
   icol[i] = pk;
   T c = 1.0 / a[pk][pk];
   det *= a[pk][pk];
   a[pk][pk] = 1.0;
   for (int p = 0; p < n; p++) a[pk][p] *= c;
   for (int p = 0; p < m; p++) b[pk][p] *= c;
   for (int p = 0; p < n; p++) if (p != pk) {
     c = a[p][pk];
     a[p][pk] = 0;
     for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
     for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p \ge 0; p--) if (irow[p] != icol[p]) {
   for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
 return det;
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \};
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {</pre>
   a[i] = VT(A[i], A[i] + n);
```

```
b[i] = VT(B[i], B[i] + m);
 }
 double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;</pre>
  // expected: -0.233333 0.166667 0.133333 0.0666667
             0.166667 0.166667 0.333333 -0.333333
  //
             0.05 -0.75 -0.1 0.2
  //
  cout << "Inverse: " << endl;</pre>
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < n; j++)
     cout << a[i][j] << ' ';
   cout << endl;</pre>
    expected: 1.63333 1.3
  //
             -0.166667 0.5
             2.36667 1.7
  //
  //
             -1.85 - 1.35
  cout << "Solution: " << endl;</pre>
 for (int i = 0; i < n; i++) {
   for (int j = 0; j < m; j++)
     cout << b[i][j] << ' ';
   cout << endl;</pre>
}
```

#### 6.4 Matrix

```
template<typename T> struct matrix {
   int N:
   vector<T> dat;
   matrixT> (int _N, T fill = T(0), T diag = T(0)) {
       N = N;
       dat.resize(N * N, fill);
       for (int i = 0; i < N; i++)</pre>
           (*this)(i, i) = diag;
   T& operator()(int i, int j) {
       return dat[N * i + j];
   matrix<T> operator *(matrix<T> &b){
       matrix<T> r(N);
       for(int i=0; i<N; i++)</pre>
           for(int j=0; j<N; j++)</pre>
              for(int k=0; k<N; k++)</pre>
                  r(i, j) = r(i, j) + (*this)(i, k) * b(k, j);
       return r;
   }
   matrix<T> pow(ll expo){
       if(!expo) return matrix<T>(N, T(0), T(1));
       matrix<T> r = (*this * *this).pow(expo/2);
       return expo&1 ? r * *this : r;
   friend ostream& operator<<(ostream &os, matrix<T> &m){
       os << "{";
```

```
for(int i=0; i<m.N; i++){</pre>
           if(i) os << "},\n ";
           os << "{";
           for(int j=0; j<m.N; j++){</pre>
               if(j) os << ", ";
               os << setw(10) << m(i, j) << setw(0);
       }
       return os << "}}";</pre>
};
struct mll {
    const int MOD;
    ll val:
    mll(ll val = 0) {
       val = _val % MOD;
       if (val < 0) val += MOD;</pre>
    mll operator+(const mll &o) {
       return mll((val + o.val) % MOD);
    mll operator*(const mll &o) {
       return mll((val * o.val) % MOD);
    friend ostream& operator << (ostream &os, mll &m) {
       return os << m.val;</pre>
};
```

## 6.5 Primes

```
// O(sqrt(x)) Exhaustive Primality Test
#define EPS 1e-7
typedef long long LL;
bool IsPrimeSlow (LL x) {
  if(x<=1) return false;</pre>
  if(x<=3) return true;
  if (!(x\%2) || !(x\%3)) return false;
  LL s=(LL)(sqrt((double)(x))+EPS);
  for(LL i=5;i<=s;i+=6)</pre>
   if (!(x%i) || !(x%(i+2))) return false;
 return true:
// Primes less than 1000:
//
       2
             3
                   5
                        7
                             11
                                   13
                                         17
                                               19
                                                    23
                                                          29
                                                                31
                                                                      37
//
            43
                  47
                             59
                                         67
                                              71
                                                    73
                                                          79
      41
                       53
                                   61
                                                                      89
//
           101
                 103
                       107
                            109
                                  113
                                        127
                                             131
                                                   137
                                                         139
                                                               149
                                                                     151
//
     157
           163
                 167
                       173
                            179
                                  181
                                        191
                                             193
                                                   197
                                                         199
                                                                     223
                                                               211
//
     227
           229
                 233
                       239
                            241
                                  251
                                        257
                                              263
                                                    269
                                                         271
                                                               277
                                                                     281
//
     283
           293
                 307
                       311
                            313
                                  317
                                        331
                                             337
                                                    347
                                                         349
                                                                     359
                                                               353
//
     367
           373
                 379
                       383
                            389
                                             409
                                                   419
                                                         421
                                  397
                                        401
                                                               431
                                                                     433
//
           443
                 449
                       457
                            461
                                  463
                                        467
                                             479
                                                   487
                                                                     503
//
     509
           521
                 523
                      541
                            547
                                  557
                                        563
                                             569
                                                   571
                                                         577
                                                               587
                                                                     593
//
     599
           601
                 607
                      613
                            617
                                  619
                                        631
                                             641
                                                   643
                                                         647
                                                               653
                                                                     659
//
           673
                 677
                       683
                                        709
                                              719
                                                   727
                                                         733
     661
                            691
                                  701
                                                               739
                                                                     743
//
     751
           757
                 761
                      769
                            773
                                  787
                                        797
                                              809
                                                   811
                                                         821
                                                               823
                                                                     827
     829
           839
                 853
                      857
                            859
                                  863
                                        877
                                             881
                                                   883
                                                         887
                                                               907
                                                                     911
```

### 6.6 Reduced Row Echelon Form

```
// Reduced row echelon form via Gauss-Jordan elimination
// with partial pivoting. This can be used for computing
// the rank of a matrix.
// Running time: O(n^3)
// INPUT: a[][] = an nxn matrix
// OUTPUT: rref[][] = an nxm matrix (stored in a[][])
            returns rank of a[][]
typedef vector<double> VD;
typedef vector<VD> VVD;
const double EPSILON = 1e-7;
// returns rank
int rref (VVD &a){
 int i, j, r, c;
 int n = a.size();
  int m = a[0].size();
 for (r=c=0;c< m;c++){
   for (i=r+1; i < n; i++) if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
   if (fabs(a[j][c]) < EPSILON) continue;</pre>
   for (i=0;i<m;i++) swap(a[j][i],a[r][i]);</pre>
   double s = a[r][c];
   for (j=0;j<m;j++) a[r][j] /= s;</pre>
   for (i=0;i<n;i++) if (i != r){</pre>
     double t = a[i][c];
     for (j=0;j<m;j++) a[i][j] -= t*a[r][j];</pre>
   r++;
 return r;
```

## 6.7 Simplex

```
// Two-phase simplex algorithm for solving linear programs of the form
//
// maximize c^T x
// subject to Ax <= b
// x >= 0
//
// INPUT: A -- an m x n matrix
// b -- an m-dimensional vector
// c -- an n-dimensional vector
// x -- a vector where the optimal solution will be stored
```

```
// OUTPUT: value of the optimal solution (infinity if unbounded
          above, nan if infeasible)
//
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then. call Solve(x).
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
 int m, n;
 VI B, N;
 VVD D:
 LPSolver(const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2, VD(n + 2)) {
   for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] =
   for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n] = -1; D[i][n + 1]
        = b[i]; }
   for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
   N[n] = -1; D[m + 1][n] = 1;
 void Pivot(int r, int s) {
   for (int i = 0; i < m + 2; i++) if (i != r)
     for (int j = 0; j < n + 2; j++) if (j != s)
       D[i][j] = D[r][j] * D[i][s] / D[r][s];
   for (int j = 0; j < n + 2; j++) if (j != s) D[r][j] /= D[r][s];
   for (int i = 0; i < m + 2; i++) if (i != r) D[i][s] /= -D[r][s];
   D[r][s] = 1.0 / D[r][s];
   swap(B[r], N[s]);
  bool Simplex(int phase) {
   int x = phase == 1 ? m + 1 : m;
   while (true) {
     int s = -1;
     for (int j = 0; j \le n; j++) {
       if (phase == 2 && N[j] == -1) continue;
       if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j] <
           N[s]) s = j;
     if (D[x][s] > -EPS) return true;
     int r = -1;
     for (int i = 0; i < m; i++) {</pre>
       if (D[i][s] < EPS) continue;</pre>
       if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] ||
         (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s]) && B[i] <
             B[r]) r = i;
     if (r == -1) return false;
     Pivot(r, s);
 DOUBLE Solve(VD &x) {
   int r = 0;
   for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n + 1]) r = i;
```

```
if (D[r][n + 1] < -EPS) {
     Pivot(r, n);
     if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return</pre>
          -numeric_limits<DOUBLE>::infinity();
     for (int i = 0; i < m; i++) if (B[i] == -1) {
       int s = -1:
       for (int j = 0; j \le n; j++)
         if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i][s] && N[j] <
             N[s]) s = j;
       Pivot(i, s);
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity();
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n + 1];
    return D[m][n + 1];
};
int main() {
  const int m = 4;
  const int n = 3;
 DOUBLE _A[m][n] = {
    \{6, -1, 0\},\
    \{-1, -5, 0\},\
    { 1, 5, 1 },
    \{-1, -5, -1\}
 DOUBLE _b[m] = \{ 10, -4, 5, -5 \};
 DOUBLE _{c}[n] = \{ 1, -1, 0 \};
 VVD A(m);
  VD b(_b, _b + m);
  VD c(_c, _c + n);
 for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n);
 LPSolver solver(A, b, c);
  VD x;
 DOUBLE value = solver.Solve(x);
  cerr << "VALUE: " << value << endl; // VALUE: 1.29032</pre>
  cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1</pre>
  for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
  cerr << endl;</pre>
 return 0;
```

# 7 7 String

### 7.1 Aho-Corasick

```
namespace aho_corasick {
  const int SIGMA = 2;
  const int TOTL = 1e7 + 100;

struct node {
   int link[SIGMA];
   int suff, dict, patt;
   node() {
```

```
suff = 0, dict = 0, patt = -1;
       memset(link, 0, sizeof(link));
   }
   // link[]: contains trie links + failure links
   // suff: link to longest proper suffix that exists in the trie
   // dict: link to longest suffix that exists in the dictionary
   // patt: index of this node's word in the dictionary
int tail = 1:
vector<node> trie(TOTL);
vector<string> patterns;
void add_pattern(string &s) {
   int loc = 0;
   for (char c : s) {
       int &nloc = trie[loc].link[c-'a'];
       if (!nloc) nloc = tail++;
       loc = nloc;
   }
   trie[loc].dict = loc;
   trie[loc].patt = patterns.size();
   patterns.push_back(s);
void calc_links() {
   queue<int> bfs({0});
   while (!bfs.empty()) {
       int loc = bfs.front(); bfs.pop();
       int fail = trie[loc].suff;
       if (!trie[loc].dict)
          trie[loc].dict = trie[fail].dict;
       for (int c = 0; c < SIGMA; c++) {
          int &succ = trie[loc].link[c];
          if (succ) {
              trie[succ].suff = loc ? trie[fail].link[c] : 0;
              bfs.push(succ);
          } else succ = trie[fail].link[c];
       }
   }
void match(string &s, vector<bool> &matches) {
   int loc = 0;
   for (char c : s) {
       loc = trie[loc].link[c-'a'];
       for (int dm = trie[loc].dict; dm; dm =
           trie[trie[dm].suff].dict) {
          if (matches[trie[dm].patt]) break;
          matches[trie[dm].patt] = true;
       }
   }
}
```

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### 7.2 KMP

```
template<typename T> struct kmp {
  int M;
  vector<T> needle;
```

```
vector<int> succ;
   kmp(vector<T> _needle) {
       needle = _needle;
       M = needle.size();
       succ.resize(M + 1);
       succ[0] = -1, succ[1] = 0;
       int cur = 0;
       for (int i = 2; i <= M; ) {</pre>
           if (needle[i-1] == needle[cur]) succ[i++] = ++cur;
           else if (cur) cur = succ[cur];
           else succ[i++] = 0;
       }
   }
   vector<bool> find(vector<T> &haystack) {
       int N = haystack.size(), i = 0;
       vector<bool> res(N);
       for (int m = 0; m + i < N; ) {
           if (i < M && needle[i] == haystack[m + i]) {</pre>
              if (i == M - 1) res[m] = true;
              i++;
           } else if (succ[i] != -1) {
              m = m + i - succ[i];
              i = succ[i];
           } else {
              i = 0;
              m++:
       return res;
};
```

## 7.3 Suffix Arrays

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
//
          of substring s[i...L-1] in the list of sorted suffixes.
//
          That is, if we take the inverse of the permutation suffix[],
          we get the actual suffix array.
struct SuffixArray {
 const int L;
 string s;
 vector<vector<int> > P;
 vector<pair<int,int>,int> > M;
 SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L,
      0)), M(L) {
   for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
   for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {</pre>
     P.push_back(vector<int>(L, 0));
     for (int i = 0; i < L; i++)
```

```
M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ?</pre>
           P[level-1][i + skip] : -1000), i);
     sort(M.begin(), M.end());
     for (int i = 0; i < L; i++)
       P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ?
            P[level][M[i-1].second] : i;
 }
 vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and
      s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
   int len = 0:
   if (i == j) return L - i;
   for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
     if (P[k][i] == P[k][j]) {
       i += 1 << k;
       j += 1 << k;
       len += 1 << k;
   return len;
};
int main() {
 // bobocel is the 0'th suffix
 // obocel is the 5'th suffix
 // bocel is the 1'st suffix
     ocel is the 6'th suffix
       cel is the 2'nd suffix
         el is the 3'rd suffix
         l is the 4'th suffix
 SuffixArray suffix("bobocel");
 vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
 for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
 cout << endl;</pre>
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
```

## 8 8 Misc

## 8.1 IO

```
int main() {
    // Ouput a specific number of digits past the decimal point,
    // in this case 5
    cout.setf(ios::fixed); cout << setprecision(5);
    cout << 100.0/7.0 << endl;
    cout.unsetf(ios::fixed);
    // Output the decimal point and trailing zeros
    cout.setf(ios::showpoint);
    cout << 100.0 << endl;
}</pre>
```

```
cout.unsetf(ios::showpoint);
// Output a '+' before positive values
cout.setf(ios::showpos);
cout << 100 << " " << -100 << endl;
cout.unsetf(ios::showpos);
// Output numerical values in hexadecimal
cout << hex << 100 << " " << 1000 << " " << 10000 << endl;</pre>
```

## 8.2 Longest Increasing Subsequence

```
// Given a list of numbers of length n, this routine extracts a
// longest increasing subsequence.
// Running time: O(n log n)
// INPUT: a vector of integers
// OUTPUT: a vector containing the longest increasing subsequence
typedef vector<int> VI;
typedef pair<int,int> PII;
typedef vector<PII> VPII;
#define STRICTLY_INCREASNG
VI LongestIncreasingSubsequence(VI v) {
 VPII best:
 VI dad(v.size(), -1);
 for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
   PII item = make_pair(v[i], 0);
   VPII::iterator it = lower_bound(best.begin(), best.end(), item);
   item.second = i;
#else
   PII item = make_pair(v[i], i);
   VPII::iterator it = upper_bound(best.begin(), best.end(), item);
#endif
   if (it == best.end()) {
     dad[i] = (best.size() == 0 ? -1 : best.back().second);
     best.push_back(item);
   } else {
     dad[i] = dad[it->second];
     *it = item;
 }
 for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
 reverse(ret.begin(), ret.end());
 return ret;
```

## 8.3 Regular Expressions - Java

```
// Code which demonstrates the use of Java's regular expression libraries.
// This is a solution for
//
// Loglan: a logical language
// http://acm.uva.es/p/v1/134.html
```

```
// In this problem, we are given a regular language, whose rules can be
// inferred directly from the code. For each sentence in the input, we
// determine whether the sentence matches the regular expression or not.
// code consists of (1) building the regular expression (which is fairly
// complex) and (2) using the regex to match sentences.
import java.util.*;
import java.util.regex.*;
public class LogLan {
   public static String BuildRegex (){
       String space = " +";
       String A = "([aeiou])";
       String C = "([a-z\&\&[^aeiou]])";
       String MOD = (g' + A + )';
       String BA = "(b" + A + ")";
       String DA = "(d" + A + ")";
       String LA = "(1" + A + ")";
       String NAM = "([a-z]*" + C + ")";
       String PREDA = "(" + C + C + A + C + A + "|" + C + A + C + C + A +
       String predstring = "(" + PREDA + "(" + space + PREDA + ")*)";
       String predname = "(" + LA + space + predstring + "|" + NAM + ")";
       String preds = "(" + predstring + "(" + space + A + space +
           predstring + ")*)";
       String predclaim = "(" + predname + space + BA + space + preds +
           "|" + DA + space +
          preds + ")":
       String verbpred = "(" + MOD + space + predstring + ")";
       String statement = "(" + predname + space + verbpred + space +
           predname + "|" +
          predname + space + verbpred + ")";
       String sentence = "(" + statement + "|" + predclaim + ")";
       return "^" + sentence + "$";
   public static void main (String args[]){
       String regex = BuildRegex();
       Pattern pattern = Pattern.compile (regex);
       Scanner s = new Scanner(System.in);
       while (true) {
           // In this problem, each sentence consists of multiple lines,
               where the last
           // line is terminated by a period. The code below reads lines
          // encountering a line whose final character is a '.'. Note
               the use of
                s.length() to get length of string
```

```
s.charAt() to extract characters from a Java string
       // s.trim() to remove whitespace from the beginning and end
           of Java string
       // Other useful String manipulation methods include
       //
       //
            s.compareTo(t) < 0 if s < t, lexicographically</pre>
       // s.indexOf("apple") returns index of first occurrence of
           "apple" in s
       // s.lastIndexOf("apple") returns index of last occurrence
           of "apple" in s
       // s.replace(c,d) replaces occurrences of character c with d
            s.startsWith("apple) returns (s.indexOf("apple") == 0)
       // s.toLowerCase() / s.toUpperCase() returns a new
           lower/uppercased string
       //
       //
            Integer.parseInt(s) converts s to an integer (32-bit)
            Long.parseLong(s) converts s to a long (64-bit)
            Double.parseDouble(s) converts s to a double
       String sentence = "";
       while (true) {
          sentence = (sentence + " " + s.nextLine()).trim();
          if (sentence.equals("#")) return;
          if (sentence.charAt(sentence.length()-1) == '.') break;
       }
       // now, we remove the period, and match the regular expression
       String removed_period = sentence.substring(0,
           sentence.length()-1).trim();
       if (pattern.matcher (removed_period).find()){
          System.out.println ("Good");
       } else {
          System.out.println ("Bad!");
}
```