

Functional Programming

Lecture 1

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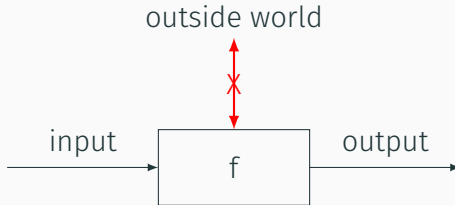
Introduction

What is functional programming?

- **Functional programming** is a programming style that prefers to structure computer programs as compositions of **pure functions**.
- It does not depend on a programming language but some languages are more suitable for functional programming than others.
- **Functional programming languages** are languages encouraging usage of pure functions.

Pure functions

A **pure function** is a function that, given the same input, will always return the same output and has no observable side effect.



No side effects = pure functions **cannot modify** and **don't depend on** any existing data structures

A **pure functional program** = a composition of pure functions

Examples of (im)pure functions

```
counter = 0
```

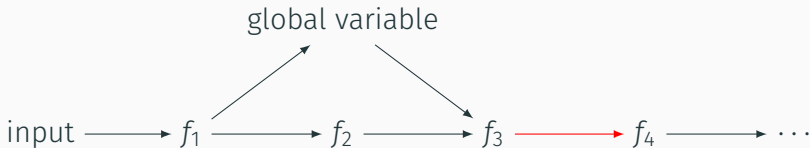
```
def pure(x, y):  
    return (x + y)/2
```

```
def do_other(x):  
    global counter  
    counter += 1  
    return x**2
```

```
def depends_on_other(x):  
    return counter + x**2
```

Advantages

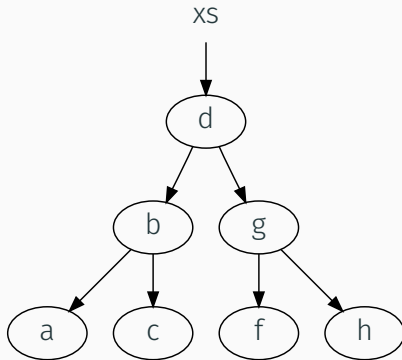
- unit testing and debugging
- simpler refactoring
- concurrency and parallelism — $f(g_1, \dots, g_n)$
- formal verification — mathematical induction, algebraic reasoning
- compiler optimization, pure functions are cachable



Consequences

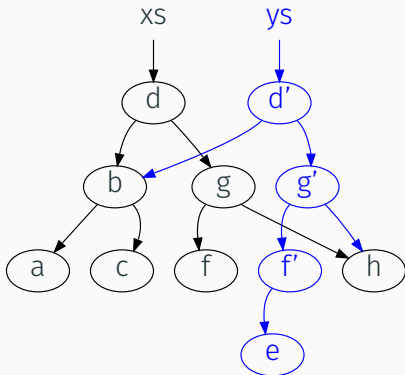
- Imperative loops updates a state in each iteration. FP uses **recursion** instead (stack holds the state).
- Data structures in pure functional programs are **immutable**.
- To modify a data structure, we need to copy it and do the desired modification.
- The code generated by functional programming languages is typically less efficient.
- To reduce the number of copying, **persistent data structures** are used.

Persistent data structures



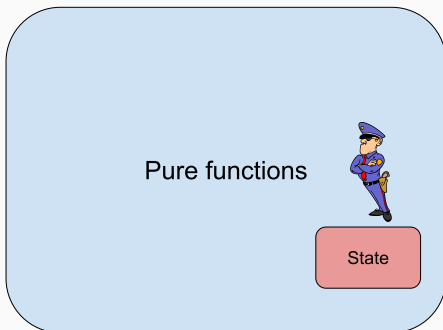
Persistent data structures

```
ys = insert ("e", xs)
```



Necessary side effects

- A pure functional program behaves like a calculator.
- Real applications need side effects. In FP, we tend to make the pure part of an app as large as possible, keeping the “unsafe” effectful code to the bare minimum.



A bit of history

Alonzo Church	Alan Turing
λ -calculus	Turing machine
Functional programming	Imperative programming
Composition of functions	Seq. of instructions changing state
Function application	Instruction execution
Recursion	Loops

Theorem

*Turing machines and λ -calculus are **equally** strong regarding computing functions.*

Organization

Course organization

- Web: <https://aicenter.github.io/FUP/>
- Lectures + Labs
- BRUTE Homework assignments (50 points) ≥ 25
 - 2x Racket
 - 2x Haskell
 - must have at least 1 point from each
 - Deadlines: -1 per day until +1 is left
- Programming exam (30 points) ≥ 16
- Theoretical oral exam (20 points) ≥ 0

What will we learn?

- **Lisp/Scheme/Racket**
 - simple syntax (directly matches λ -calculus)
 - dynamically typed
 - code-as-data (easy to write interpreters,...)
 - allows mutable data
- **λ -calculus**
- **Haskell**
 - pure functional language
 - statically typed
 - rich type system
 - strictly separates the pure core from the mutable shell

Suggested literature

[1] Harold Abelson and Gerald Jay Sussman and Julie Sussman: *Structure and Interpretation of Computer Programs*, MIT Press, 1996. <https://mitpress.mit.edu/sites/default/files/sicp/full-text/book/book.html>

[2] Raul Rojas. *A Tutorial Introduction to the Lambda Calculus*. <http://www.inf.fu-berlin.de/lehre/WS03/alpi/lambda.pdf>

[3] Graham Hutton: *Programming in Haskell*, Cambridge University Press, 2016.

Lecture notes: <https://aicenter.github.io/FUP/>

Lisp/Scheme/Racket

- **Lisp** = List processor
- **Scheme** is a dialect of Lisp (such as Common Lisp, Clojure)
- Scheme last standard from 2007 — The Revised6 Report on the Algorithmic Language Scheme (R6RS)
- **Racket** is another dialect based on R5RS (Scheme with batteries)
- **DrRacket**: **racket-lang.org**
text editor + REPL (read-evaluate-print loop)

Racket's syntax

Racket program is a collection of expressions

- Primitive expressions (literals, built-in functions)

`"Hello World!"`

- Compound expressions (built by function composition)

`(cos (+ 1 2))`

- Definitions (introduce new names and functions)

`(define (square x) (* x x))`

- Comments

`; This is a one-line comment`

`#|`

`This is`

`a block comment`

`|#`

Compound expressions

Compound expressions are built from primitive expressions by function composition.

Racket uses **prefix notation**. E.g.

$$\frac{xy^2 + 3}{x - 1}$$

```
(/ (+ (* x y y) 3)
   (- x 1))
```

S-expression

```
(fn arg1 arg2 ... argN)
```

Definitions

- Naming expressions

```
(define id exp)
```

- Defining functions

```
(define (name arg1 ... argN)  
  exp1  
  ...  
  expM)
```

- Nested definitions

```
(define (name a1 ... aN)  
  (define (fn b1 ... bM) <body-fn>)  
  <body-using-fn>)
```

Racket's semantics

Racket program is, in fact, an **expression**.

Its **evaluation** is the computation process represented by the program.

The evaluation resembles simplifying expressions we know from math.

More precisely, we subsequently evaluate subexpressions until we end up with the expression's **value**.

Evaluation strategy

```
(define (square x) (* x x))
```

```
(square (+ 3 4))
```

```
(square (+ 3 4)) => (square 7) => (* 7 7) => 49
```

```
(square (+ 3 4)) => (* (+ 3 4) (+ 3 4))  
=> (* 7 7) => 49
```

- **Evaluation strategy** defines the order of evaluating the expressions, influences program termination, not the result
- Racket's strategy is **strict** (or eager) evaluates all arguments (left to right) before evaluating the function
- Evaluation of some syntactic forms is **lazy** if, cond, and, or

Conditional expressions

```
(if test-exp then-exp else-exp)
```

```
(if (> 0 1) 1 2) => 2
```

```
(if (< 0 1) 1 (+ 3 "a")) => 1
```

```
(cond [test-exp1 exp]  
      [test-exp2 exp]  
      ...  
      [else exp])
```

```
(cond [(odd? 12) 1]  
      [(even? 12) 2]  
      [else 3])    => 2
```

Basic data types

- **Numbers:** exact $\frac{1}{2}$, inexact 3.14, complex $2 + 3i$
+, -, *, /, abs, sqrt, number?, <, >, =
- **Logical values:** #t, #f
and, or, not, boolean?
- **Strings:** "abc"
string?, substring, string-append
- **Characters:** #\A, #\@
char?, char->integer, integer->char,
list->string, string->list
- **Other types:**
symbol?, pair?, procedure?, vector?, port?

Simple debugging

- Helper print-outs

```
(begin (displayln x)  
      <do-work>)
```

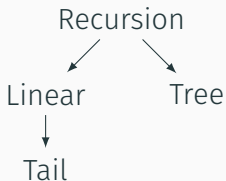
- Tracing function calls and returns

```
(require racket/trace)  
(trace fn)  
(untrace fn)
```

Recursion

Recursion

Recursive function calls itself in its body.



Indirect (mutual)



- **Linear:** makes one recursive call
- **Tree:** makes several recursive calls
- **Tail:** the result of the recursive call is the final result of the function

Examples

```
(define (loop) (loop))
```

```
(define (fact n)
  (if (<= n 1)
      1
      (* n (fact (- n 1)))))
```

```
(fact 4) => (* 4 (fact 3))
=> (* 4 (* 3 (fact 2)))
=> (* 4 (* 3 (* 2 (fact 1))))
=> (* 4 (* 3 (* 2 (* 1 1)))) => 24
```

Not space efficient. It needs $O(n)$ memory.

Example — Tail recursion

```
(define (fact n [acc 1])  
  (if (<= n 1)  
      acc  
      (fact (- n 1) (* n acc))))
```

```
(fact 4) = (fact 4 1)  
=> (fact 3 4)  
=> (fact 2 12)  
=> (fact 1 24)  
=> 24
```

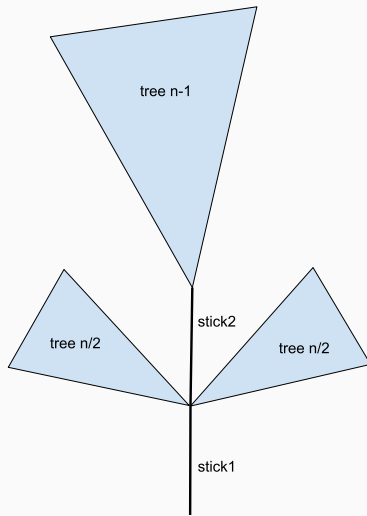
This needs $O(1)$ memory due to tail elimination.

Example – Tree recursion

Consider a tree-like fractal of a given size n and direction d in degrees generated by:

1. Draw a stick of size n in the direction d .
2. Draw the fractal of size $n/2$ in the direction $d + 60$.
3. Draw the fractal of size $n/2$ in the direction $d - 60$.
4. Draw a stick of size n in the direction d .
5. Draw the fractal of size $n - 1$ in the direction $d + 5$.

Fractal example

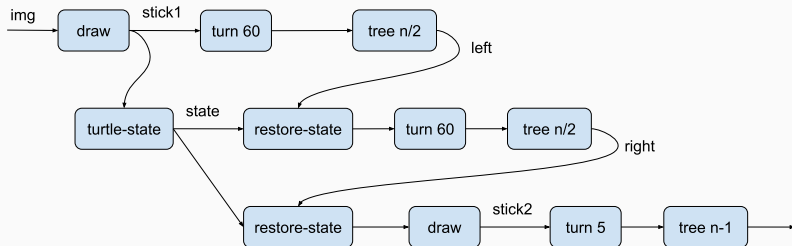


Fractal example

To draw a picture, we use the library `value-turtles`.

Its functions operates on an image together with a position and direction of a turtle.

E.g. `(draw 100 img)`



What have we learned?

- A pure function always returns the same output on a fixed input and has no side effects.
- Make the pure part of a program as large as possible, keeping the code handling the state transparent and small.
- Functional languages handle iterative computations by recursion.
- We classify recursive functions according to the number of recursive calls they make on linear-recursive and tree-recursive functions.
- Tail recursive functions are space efficient as they do not consume memory by making recursive calls.