Multiplayer Game Programming

Lecture 3: TCP and Sockets

ITP 484

Problems with UDP

- Offers no guarantees over the underlying network layer, which means it too is Unreliable
 - Segments might arrive out of order
 - Segments aren't guaranteed to arrive at all

TCP: Transmission Control Protocol

- "Reliable"
 - Segments are received in order
 - Segments are "guaranteed" to arrive

```
struct TCPHeader
{
  mDestinationPort;
  uint16 t
  uint32 t
              mSequenceNumber;
  uint32 t
              mAcknowledgementNumber;
  unsigned int mDataOffset: 4; //in 32bit words
  unsigned int mUnused: 4;
  uint8 t
              mFlags;
  uint16 t
              mReceiveWindow;
  uint16 t mChecksum;
  uint16 t
              mUrgentPointer;
};
```

How does reliability work

- Receiver acknowledges segments
- Sender resends unacknowledged segment
- Sequence number increases by size of segment
- Acknowledgement number is sequence number of last received segment + the size of the segment's data

Examples

- Normal
- Lost segment in middle of stream
- Lost segment at end of stream
- Delayed ACK
 - Up to 500 ms or every other packet
- Lost ACK

Requirements

- Buffer for segments
 - So they can be resent (by sender)
 - So they can be processed in order (by receiver)
- Stateful Connection
 - Sequence number
 - Acknowledgement number
 - Buffers
 - Timers

How much to send per segment?

- Maximum Segment Size (MSS)
 - Don't Fragment (DF) Flag on Network Layer, causes segment to get dropped if too big

```
struct TCPHeader
 uint32 t mAcknowledgementNumber;
 unsigned int mDataOffset: 4;
 unsigned int mUnused: 4;
 uint8 t
         mFlags;
 uint16 t
         mReceiveWindow;
 uint16 t mChecksum;
 };
```

Flags

- CWR: Congestion Window Reduced
- ECE: Explicit Congestion Notification (ECN)capable
- URG: Urgent pointer field is significant
- ACK: Acknowledgment field is significant
- PSH: Push function
- RST: Reset the connection
- SYN: Synchronize the sequence number
- FIN: Finished!

Handshake and Shutdown

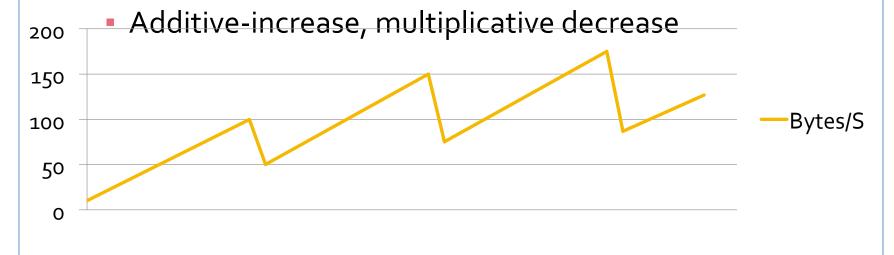
- Three way to establish
 - SYN->SYN,ACK->ACK
 - "Phantom Byte" sequence number increase
- Four way to shutdown
 - FIN->ACK.....FIN->ACK
 - Receiver can finish sending or other activities

Encouraging Reliability

- Flow Control
 - Prevents receiver's data buffer getting full
 - Receive Window: maximum unack'd data
 - Max bandwidth?
 - Receive Window / Round Trip Time

Encouraging Reliability

- Congestion Control
 - Prevents network getting full
 - Congestion Window: maximum unack'd data
 - Up when no packets lost, Down when packets lost



Disadvantages of TCP

- Must establish a connection
- Higher latency than UDP
- No way to send packets unreliably

Sockets

- Berkeley Socket API
- Originally from BSD UNIX
- Ported in some way to almost every language and operating system

Operating System Differences

- POSIX / Berkeley Sockets
 - <sys/socket.h>
 - Socket is a file descriptor (int)
- Windows
 - <Winsock2.h>
 - Socket is a SOCKET

Creating a socket

Binding a socket to an address

```
int bind(
    _In_ SOCKET s,
    _In_ const struct sockaddr *name,
    _In_ int namelen );
```

 Actual type of sockaddr can change based on protocol family, etc.

SOCKADDR

```
struct sockaddr
     ushort sa family;
     char sa data[14];
};
struct sockaddr in
     short sin family;
     u short sin port;
     struct in addr sin addr;
     char sin zero[8];
};
```

IN_ADDR IPv4 address

```
typedef struct in addr
      union
             struct
                   u char s b1, s b2, s b3, s b4;
             } S un b;
             struct
                   u short s w1,s_w2;
             } S un w;
             u long S addr;
      } S un;
 IN ADDR;
```

INADDR_ANY will allow you to bind to all local interfaces on the host

Creating an address from a string

- nodename: dns name or ip address
- Servname: port number or service name
- hints: can specify address family, etc
- res: results. Must be deallocated with freeaddrinfo

Byte Order

- Network Byte Order is Big Endian
- Your platform is probably Little Endian
- Use conversion functions when manually filling in or reading address structures

```
u_short htons( u_short hostshort );
u_long htonl( u_long hostlong);
u_short ntohs( u_short netshort);
u_long ntohl(u_long netlong);
```

Datagram Transmission

```
int sendto(
   _In_ SOCKET s,
   _In_ const char *buf,
   _In_ int len,
   _In_ int flags,
   _In_ const struct sockaddr *to,
   _In_ int tolen );
```

Returns number of bytes sent or a negative error number

Datagram Reception

Returns number of bytes sent or a negative error number

Stream Connections

```
int listen(
   _In_ SOCKET s,
   _In_ int backlog );
```

- backlog: number of connections that can be pending acceptance
- Returns o for no error

Stream Connections

```
SOCKET accept(
_In_ SOCKET s,
_Out_ struct sockaddr
*addr,
_Inout_ int *addrlen );
```

Returns a new socket, and its address info

Stream Connections

```
int connect(
   _In_ SOCKET s,
   _In_ const struct sockaddr *name,
   _In_ int namelen );
```

Blocking

 Sending, receiving, connecting and accepting all block by default if there's no space, data, pending connection available or just not enough time to complete

Pass FIONBIO as cmd and argp that points to the number 1 to enable nonblocking mode

In nonblocking mode, many sockets will return what looks like and error, but when you call WSAGetLastError you'll see the error is just WSAEWOULDBLOCK. This means there's no result yet because the operation hasn't completed yet

Blocking with Select

Returns number of ready sockets or negative error. Sets are modified by function

Select Socket Set

```
FD ZERO(fd set* set )
Initializes set
FD SET(SOCKET s, fd set* set )
Add socket to set
FD ISSET(SOCKET s, fd set* set )
Tests if a socket is set
FD CLR(SOCKET s, fd set* set )
Removes socket from set
```

Initialization and Shutdown