Abstractions and Object-Orientation for PDE Solvers

Knut-Andreas Lie

Dept. of Informatics, University of Oslo

Abstractions for PDE Solvers

For PDEs, the discretization consists of

- The unknown u as a scalar or vector *field*
- The field is defined over a grid

Primarily a field consists of values at discrete grid points, but may also contain interpolation rules (e.g., for FEM methods).

- Industry PDE simulators: 50 000+ code lines, but not all is number crunching
- Maintainability important
- Should be easy to extend
- Should be easy to use/understand
- Abstractions close to mathematical language are needed

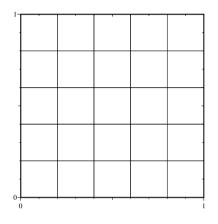
Grid and Field Abstractions

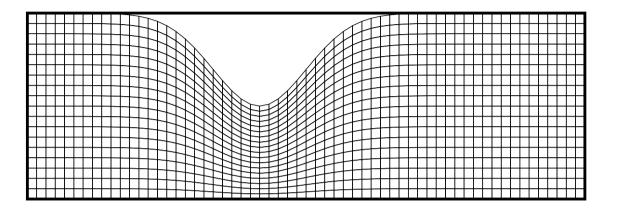
Consider the PDE (describing e.g., heat conduction):

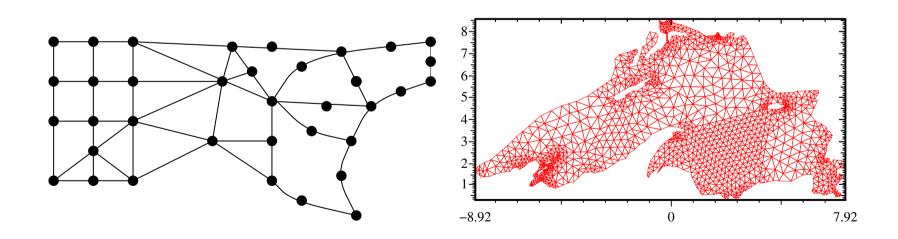
$$-\nabla \cdot [\lambda(\mathbf{x})\nabla u(\mathbf{x})] = f(\mathbf{x}), \quad \mathbf{x} \in \Omega$$

- Principal mathematical quantities:
 - scalar fields: $\lambda(\mathbf{x})$, $u(\mathbf{x})$, $f(\mathbf{x})$
 - domain: Ω
- Principal numerical quantities:
 - scalar fields, arbitrary representation: $\lambda(\mathbf{x})$, $f(\mathbf{x})$ (explicit functions, discrete fields)
 - scalar fields, finite difference type: discrete u
 - grid: discrete Ω

How Does a Grid Look Like?







Programming Considerations

- Obvious ideas:
 - collect grid information in a grid class
 - collect field information in a field class
- Gain:
 - shorter code, closer to the mathematics
 - finite difference methods: minor
 - finite element methods: important
 - big programs: fundamental
 - possible to write code that is (almost) independent of the number of space dimensions (i.e., easy to go from 1D to 2D to 3D!)

Grids and Fields for FDM

Assume a finite difference method (FDM):

- Field represented by FieldLattice:
 - a grid of type GridLattice
 - a set of point values, MyArray
 - MyArray is a class implementing user-friendly arrays in one and more dimensions (i.e., an extension of our earlier MyArray)
- Grid represented by GridLattice
 - lattice with uniform partition in d dimensions
 - initialization from input string, e.g.,

```
d=1 domain: [0,1], index [1:20]

d=3 [0,1]x[-2,2]x[0,10]
 indices [1:20]x[-20:20]x[0:40]
```

Programming with a GridLattice Class

```
GridLattice q: // declare an empty grid
g.scan("d=2\lfloor [0,1]x[0,2] \lfloor [1:10]x[1:40]"); // initialize g
const int i0 = g.getBase(1); // start of first index
const int j0 = g.getBase(2); // start of second index
const int in = g.getMaxI(1); // end of first index
const int jn = g.getMaxI(2); // end of second index
int i.i:
for (i = i0; i \le in; ++i)
   for (i = i0; i \le in; ++i)
        std::cout << "grid_point_(" << i << ', ' << j
                 << ") has coordinates (" << q.getPoint(1,i)</pre>
                 << ', ' << g.getPoint(2, j ) << ")\n";
// other tasks:
const int nx = g.getDivisions (1); // number of grid cells in x dir
const int ny = g.getDivisions (2); // number of grid cells in y dir
const int dx = g.Delta(1); // grid spacing in x dir
const int dy = g.Delta (2); // grid spacing in y dir
```

The GridLattice Class

Data representation:

- Trivial in our case, use: $x_{\min}, x_{\max}, y_{\min}, y_{\max}$, and n_x, n_y .
- For unstructured grids: a carefully designed data structure is vital to obtain code efficiency

```
class GridLattice
{
    static const int MAX_DIMENSIONS = 2;

    double min[MAX_DIMENSIONS]; // min coordinate values in each dimension
    double max[MAX_DIMENSIONS]; // max coordinate values in each dimension
    int division[MAX_DIMENSIONS]; // number of points in each dimension
    int dimensions; // number of dimensions
```

The static keyword means that MAX_DIMENSIONS is a "global" constant shared by all GridLattice objects.

Member functions:

- Constructors
- Initialization through the scan function
- Accessors access to internal data structure

```
double Delta(int dimension) const;
double getPoint(int dimension, int index);

int getBase(int dimension) const; // get base values
int getMaxI(int dimension) const; // upper limit of array

void scan(const std::string& init_string ); // scan parameters from string

friend std::ostream& operator<<(std::ostream&, GridLattice&);
};</pre>
```

Mutators, i.e., functions for setting internal data members, are not implemented here. Examples: setBase(..), setMaxI(..), etc.

The friend keyword enables operator« to access private data in the GridLattice object.

```
double GridLattice:: xMin(int dimension) const
{ return min[dimension - 1]; }
double GridLattice:: xMax(int dimension) const
{ return max[dimension - 1]; }
inline int GridLattice :: getDivisions(int i) const
{ return division [i −1]; }
int GridLattice :: getNoPoints() const {
    int return value = 1;
   for(int i = 0; i = dimensions; ++i)
       return_value *= division[i];
   return return value;
```

Inline functions means that the function call can be optimized away by the compiler.

Nested inline functions:

```
inline double GridLattice:: Delta(int dimension) const {
   return (max[dimension-1] - min[dimension-1]) / double(division[dimension-1]);
}
inline double GridLattice:: getPoint(int dimension, int index) {
   return min[dimension-1] + (Delta(dimension) * (index - 1));
}
```

Some compilers do not allow nested inlines. To come around this, we can use a preprocessor macro:

Typical call: g.scan("d=2 [0,1]x[0,2] [1:10]x[1:40]");

```
void GridLattice :: scan(const string& init_string)
    using namespace std; // allows dropping std:: prefix
    istrstream is ( init_string .c_str ());
    // ignore "d="
    is.ignore(1, 'd'); is.ignore(1, '=');
    // get the dimensions
    is >> dimensions;
    if (dimensions < 1 || dimensions > MAX_DIMENSIONS) {
       cerr << "GridLattice :: scan()_-__illegal_dimensions_"
            << dimensions << endl;
       cerr << "___MAX_DIMENSIONS_is_set_to_"
            << MAX_DIMENSIONS << endl;
        exit (1);
```

```
GridLattice:: GridLattice(int nx, int ny, double xmin, double xmax,
                          double ymin, double ymax) {
   dimensions = 2;
   max[0] = xmax; max[1] = ymax;
   min[0] = xmin; min[1] = ymin;
   division [0] = nx; division [1] = ny;
GridLattice :: GridLattice () {
   dimensions = 2;
    int i;
   for (i = 1; i \le MAX_DIMENSIONS; ++i) {
       min[i] = 0; max[i] = 1; division [i] = 2;
```

We have chosen to initialize internal data also for the empty constructor to avoid errors.

The FieldLattice Class

```
class FieldLattice {
protected:
                         grid_lattice;
   Handle<GridLattice>
   Handle< MyArray<real> > grid_point_values;
                          fieldname:
   std:: string
public:
    // make a field from a grid and a fieldname:
    FieldLattice (GridLattice & g, const std::string & fieldname);
         MyArray<real>& values() { return *grid point values; }
   const MyArray<real>& values() const { return *grid point values; }
         GridLattice& grid () { return * grid_lattice ; }
   const GridLattice& grid() const { return * grid lattice; }
   std:: string name() const
                                       { return fieldname; }
   void values(MyArray<real>& new_array);
};
```

- Inline functions are obtained by implementing the function body inside the class declaration.
- We use a parameter real, which equals float or double (by default).
- The Handle<> construction is a smart pointer, implementing reference counting and automatic deallocation (garbage collection).
- Using a Handle<GridLattice> object instead of a
 GridLattice object, means that a grid can be shared among several fields.

Example: For the wave equation um, u, and up may share the same grid (and possibly also Hm, H, and Hp).

Smart Pointers

Observations:

- Dynamic memory in C/C++ means that pointers are needed
- Lack of garbage collection (automatic clean-up of memory that is no longer in use) means that manual deallocation is important
- Codes with memory leakage slowly eat up the memory and slow down computations
- How does one determine when memory is no longer in use?
 Example: Suppose 5 fields point to the same grid when can we safely remove the grid object?
- Pointers are bug no. 1 in C/C++ ...

Smart Pointers — Reference Counting

Solution to problems:

- Avoid explicit deallocation
- Introduce reference counting: count the number of times a smart pointer "uses a pointer" (i.e., points to a given object).

Advantages:

- negligible overhead
- automatic garbage collection
- several fields can safely share one grid

```
FieldLattice :: FieldLattice (GridLattice & g, const std:: string & name )
  grid lattice .rebind(&g);
  // allocate the grid point values array:
  if ( grid_lattice ->getNoSpaceDim() == 1)
    grid_point_values.rebind(new MyArray<real>(grid_lattice->getDivisions(1)));
 else if ( grid lattice ->getNoSpaceDim() == 2)
    grid point values.rebind(new MyArray<real>(
       grid lattice ->getDivisions(1), grid lattice ->getDivisions(2)));
 else
    ; // three-dimensional fields are not yet supported...
 fieldname = name_;
void FieldLattice :: values(MyArray<real>& new array)
{ grid_point_values.rebind(&new_array); }
```

Simulator Classes

- The PDE solver is a class itself
 - ⇒ easy to combine solvers (systems of PDEs)
 - ⇒ easy to extend/modify solver
 - ⇒ enables coupling to optimization, automatic parameter analysis etc.
- Typical look (for a static problem):

```
class MySim
{
  protected:
    // grid and field objects
    // PDE dependent parameters
  public:
    void scan();    // read input and init
    void solveProblem();
    void resultReport();
};
```

Heat. C Revisited

```
#ifndef Heat1D1 h IS INCLUDED
#define Heat1D1 h IS INCLUDED
class Heat1D1
           // data items visible in subclasses
protected:
 MatTri<real> A; // the coefficient matrix
 MyArray<real> b; // the right—hand side
 Handle<GridLattice> grid; // 1D grid
 FieldLattice u; // the discrete solution
                  alpha; // parameter in the test problem
 real
public:
 Heat1D1() {}
 ~Heat1D1() {}
 void scan ();
              // read input, set size of A, b and u
 void solveProblem (); // compute A, b; solve Au=b
 void resultReport (); // write and plot results
};
#endif
```

Reading Input

```
void Heat1D1:: scan ()
 char gridstr [30];
  int n; // no of intervals in the domain (0,1)
 CommandLineArgs::read("-n", n, 5); // e.g ., ./ heat1d -n 20
  grid.rebind (new GridLattice());
  sprintf ( gridstr , "d=1_[0,1]_[1:%d]",n);
  grid—>scan(gridstr);
 u.rebind (new GridLattice(*grid, "u"));
 A.redim(n); // set size of tridiagonal matrix A
 b.redim(n); // set size of vector b
 CommandLineArgs::read ("-a", alpha, 0.0);
```

CommandLineArgs::read is described closer in [R&L]

Solving the Problem

```
void Heat1D1:: solveProblem () {
 A. fill (0.0); b. fill (0.0);
  const int n = b.size (); // alternative : grid->getMaxI(1)
  const real h = grid—>Delta(1);
  real x: int i:
  i = 1; A (1,1) = 1; A (1,2) = 0; b (1) = 0;
 for (i = 2; i < n; i++)
   x = grid - getPoint(1,i);
   A(i-1,i) = 1; A(i,i) = -2; A(i,i+1) = 1;
   b(i) = h*h*(alpha+1)*pow(x,alpha);
 x = grid -> getPoint(1,n);
 A(n-1,n) = 2; A(n,n) = -2;
  b(n) = -2*h + h*h*(alpha+1)*pow(x,alpha);
 A.factLU();
 A.forwBack(b,u.values());
```

The Main Program

```
#include <Heat1D1.h>

int main(int argc, const char* argv[])
{
    Heat1D1 simulator;
    simulator.scan ();
    simulator.solveProblem ();
    simulator.resultReport ();
    return 0; // success
}
```

So far, very little has been gained, but let us also consider the nonlinear case

Nonlinear Heat Conduction

Heat conduction typically depends upon the temperature

$$-\frac{d}{dx}\left(\frac{\lambda(u)}{dx}\frac{du}{dx}\right) = f(x), \quad 0 < x < 1$$

This is a *nonlinear differential equation*.

Straightforward discretization gives

$$\lambda(u_{i+\frac{1}{2}})(u_{i+1} - u_i) - \lambda(u_{i-\frac{1}{2}})(u_i - u_{i-1}) = -h^2 f_i$$

• Since $\lambda(u_{i+\frac{1}{2}})$ will depend upon u_i and u_{i+1} , this is a set of nonlinear algebraic equations

$$\mathbf{A}(\mathbf{u})\mathbf{u} = \mathbf{b}(\mathbf{u})$$

i.e., the coefficients are unknown quantities....

Sucessive Substitution

Idea: use the following simple algorithm

```
Guess a solution {f u}^0. Repeat {
m solve}~{f A}({f u}^n){f u}^{n+1}={f b}({f u}^n) until difference of {f u}^n and {f u}^{n+1} is small
```

Each equation is now a *linear* equation with variable coefficients.

Advantage: may reuse previous code by inserting evaluations of the coefficient $\lambda(u_{i+\frac{1}{2}})$

Disadvantage: slow convergence

Implementation

Reuse old program HeatTri with:

- Loop around system generation and solution
- Two arrays uk and ukm
- Initial guess in ukm
- New auxiliary variables (for iteration etc.)
- Function lambda to evaluate $\lambda(u)$
- Update A [mcp] and b for each step
- Check for termination upon convergence
- Set ukm equal uk before new iteration

Implementation in Pseudocode

```
// Iteration loop
Repeat
   // Assemble matrix and left—hand side
  for i=1:n
      // Left boundary
      // Right boundary
      // Interior
     else
        I1 = lambda(u(i-1));
        12 = lambda(u(i));
        13 = lambda(u(i+1));
        A(i,i-1) = 0.5*(11 + 12);
        A(i,i) = -0.5*(I1 + 2*I2 + I3);
        A(i,i-1) = 0.5*(12 + 13);
  end
   // Solve linear system
```

Nonlinear Heat Conduction, cont'd

We can implement different $\lambda()$ functions objects (or *functors*) in a class hierarchy.

First define the base class:

```
class LambdaFunc : public Handleld
{
public:
    virtual real lambda (real u)=0;
    virtual real exactSolution (real x)=0;
    virtual void scan () =0;
    virtual string formula ()=0;
};
```

No data and all functions pure virtual, thus defining only *an interface* (providing access all functions).

Nonlinear Heat Conduction, cont'd

Then a particular realisation in a subclass:

```
class Lambda1 : public LambdaFunc
  real m:
public:
 Lambda1() {}
  virtual real lambda (real u)
         { return pow(u,m); }
  virtual real exactSolution (real x)
         { return pow(x,1/(m+1)); }
  virtual void scan ()
         { CommandLineArgs::read("-m",m,0.0); }
  virtual String formula () { return "u^m"; }
};
```

Notice that the parameter m is a member of the *subclass* and not of the base class

The Heat1D NL Class

```
class Heat1D NL {
                    // data items visible in subclasses
protected:
 MatTri<real> A; // the coefficient matrix
 MyArray<real> b; // the right—hand side
 Handle<GridLattice> grid; // 1D grid
             uk; // the discrete solution u^k
  FieldLattice
  FieldLattice ukm; // the discrete solution u^{k-1}
                  epsilon; // tolerance in nonlinear iteration
 real
 Handle<LambdaFunc> lambda; // specific lambda function
public:
 Heat1Dn1() {}
 ~Heat1Dn1() {}
 void scan ();
               // read input, set size of A, b, and u
 void makeSystem(); // compute A, b;
 void solveSystem(); // solve Au^k=b
 void solveProblem (); // nonlinear iteration loop
 void resultReport ();  // write numerical error ( if possible)
};
```

Using λ in the Code

Picking a specific realisation

```
int lambda_tp;
CommandLineArgs::read("-N", lambda_tp, 1);
if (lambda_tp == 1)
    lambda.rebind (new Lambda1());
else if (lambda_tp == 2)
    lambda.rebind (new Lambda2());
else
```

```
Initialize function-specific parameters: lambda->scan();

Evaluating λ:
lambda->lambda(ukm.values()(i-1));

Computing the exact solution (if available):
lambda->exactSolution(grid->getPoint(1,i));
```

Wave Equation Revisited

We wish to solve:

$$\frac{\partial^2 u}{\partial t^2} = \nabla \cdot (H(x)\nabla u), \qquad u(x,0) = I(x), \quad u_t(x,0) = 0$$

The solution algorithm was presented in Lecture 4. What are the natural objects here?

- a grid object of type GridLattice
- three fields of type FieldLattice for $u_{i,j}^+$, $u_{i,j}$, and $u_{i,j}^-$
- ullet a field of type <code>FieldLattice</code> for the depth $H_{i,j}$
- a time integration parameter object TimePrm
- functor hierarchies for the known functions H(x) and I(x)

Working with class hierarchies and virtual functions is really the point that qualifies this to be called object-oriented programming as opposed to "programming with objects"...

The TimePrm Class

```
class TimePrm
   double time ; // current time value
   double delta; // time step size
   double stop; // stop time
   int
          timestep; // time step counter
public:
   TimePrm(double start, double delta, double stop)
        : time (start), delta(delta), stop(stop)
    { initTimeLoop(); }
   double time() { return time_; }
   double Delta() { return delta; }
   void initTimeLoop() { time = 0; timestep = 0; }
    bool finished () { return (time >= stop) ? true : false; }
   void increaseTime() { time_ += delta; ++timestep; }
    int getTimeStepNo() { return timestep; }
};
```

The Wave2D1 Class

```
class Wave2D1 {
   Handle<GridLattice> grid; // lattice grid here 1D grid
   Handle<FieldLattice> up,u,um; // solution u at time levels I+1, I, and I-1
   Handle<FieldLattice> lambda; // variable coefficient (depth)
   Handle<FieldLattice> tmp; // variable coefficient (depth)
   Handle<TimePrm> tip; // time integr. prms: dt, t_stop
   Handle<WaveFunc> H, I; // function for depth and initial surface
   void timeLoop(); // perform time stepping
   void plot(bool initial );  // dump fields to file , plot later
   void WAVE(FieldLattice& up, const FieldLattice& u,
            const FieldLattice& um, real a, real b, real c);
                 // set initial conditions
   void setIC ();
   void setH ();
                // load H into lambda for efficiency
   real calculateDt(int func); // calculate optimal timestep
public:
                 // read input and initialize
   void scan();
   void solveProblem(); // start the simulation
};
```

The Wave2D1 Class, cont'd

```
void Wave2D1:: solveProblem () {
 setIC ();
           // set initial conditions
 timeLoop(); // run the algorithm
void Wave2D1:: setIC () {
  const int nx = grid -> getMaxI(1);
  const int ny = grid ->getMaxl(2);
  // set initial surface elevation
  MyArray<real>& uv = u->values();
  for (int i = 1; i <= ny; i++)
     for (int i = 1; i <= nx; i++)
         uv(i, j) = I->valuePt(grid->getPoint(1, i), grid->getPoint(2, j));
  // set the help variable um:
  WAVE (*um, *u, *um, 0.5, 0.0, 0.5);
```

The Wave2D1 Class, cont'd

```
void Wave2D1:: timeLoop ()
  tip ->initTimeLoop();
  plot(true);
  while(!tip ->finished()) {
      tip ->increaseTime();
      WAVE (*up, *u, *um, 1, 1, 1);
      // move handles (get ready for next step):
      tmp = um; um = u; u = up; up = tmp;
      plot(false);
```

The Wave2D1 Class, cont'd

```
void Wave2D1:: scan () {
 // create the grid ...
 grid.rebind(new GridLattice());
 grid—>scan(CommandLineArgs::read("—grid",
     "d=2, [-10,10]x[-10,10], [1:30]x[1:30]"));
 cout << (*arid) << endl:
 // create new fields ...
       rebind(new FieldLattice(*grid, "up"));
 up.
       rebind(new FieldLattice(*grid, "u"));
        rebind(new FieldLattice(*grid, "um"));
 um.
 lambda.rebind(new FieldLattice(*grid, "lambda"));
  // select the appropriate I and H
  int func = CommandLineArgs::read("-func", 1);
  if \{\text{func} == 1\}
      H.rebind(new GaussianBell('H'));
      I.rebind(new GaussianBell('U'));
```

```
else {
    H.rebind(new Flat());
    1.rebind(new Plug('U'));
    initialize the parameters in the functions.
H->scan();
I->scan();
// set H field and compute optimal dt
setH();
tip .rebind(new TimePrm(0, calculateDt(func),
    CommandLineArgs::read("-tstop", 30.0)));
```

Object-Oriented Implementation of FEMs

Very much to be gained from using object orientation (e.g., as in Diffpack):

- based on generic FEM tools for
 - coordinate mappings
 - numerical integration
 - assembling of linear systems
 - solution of linear systems
- implement only problem specific part
- heavy code reuse
- inheritance and virtual functions are crucial
- great flexibility, e.g., code can be made independent of number of space dimensions

Example: the Poisson Equation

Model problem: $-\nabla \cdot [k(x)\nabla u(x)] = f(x), \quad x \in \Omega$ Weak formulation:

$$A_{i,j} = \int_{\Omega} k(x) \nabla N_i(x) \nabla N_j(x) d\Omega, \quad b_i = \int_{\Omega} f(x) N_i(x) d\Omega.$$

Elementwise computation in local coordinates

$$A_{i,j}^{(e)} = \int_{\Omega_e} k \nabla N_i \cdot \nabla N_j \, d\Omega_e = \int_{\tilde{\Omega}} k \nabla N_i \cdot \nabla N_j \, \det J \, d\xi_1 \cdots d\xi_d$$
$$b_i^{(e)} = \int_{\Omega_e} f N_i \, d\Omega_e = \int_{\tilde{\Omega}} f N_i \, \det J \, d\xi_1 \cdots d\xi_d$$

Numerical integration

$$A_{i,j}^{(e)} \approx \sum_{k=1}^{n_I} I_{i,j}(\xi_k) w_k, \quad b_i^{(e)} \approx \sum_{k=1}^{n_I} K_i(\xi_k) w_k.$$

Diffpack FEM Applications

- FEM solvers in Diffpack are represented as separate classes.
- Such classes are derived from the skeleton application (base class) FEM.
- The derived simulator class must provide problem-specific data:
 - the PDE in terms of the discrete integrands $I_{i,j}(\xi)w$ and $K_i(\xi)w$,
 - handling of boundary values,
 - for time-dependent problems: Handling of the initial value.
- These problem-specific details are provided by virtual functions defined by class FEM.

The Function integrands

void integrands (ElmMatVec& elmat, const FiniteElement& fe);

- The object fe is preloaded (from class FEM) with the current integration point ξ_k
- Moreover, it contains N_i , ∇N_i , $\det J \cdot w_k$ and other useful quantities, such as the global coordinates x corresponding to ξ_k
- The user must provide evaluation of the functions f(x) and k(x)
- The function integrands should add the contributions $I_{i,j}w_k$ and K_iw_k to the elemental matrix (Mat(real) structure elmat.A), and the elemental vector (Vec(real) structure elmat.b)

Example: Poisson0

Consider the test problem in 2D, $\Omega = (0,1) \times (0,1)$. The simulator class Poisson0 consists of:

- scan geometry input, allocates grid and some other large objects.
- fillessBC sets the essential boundary conditions u=0 on $\partial\Omega_E$.
- integrands samples the element-based integrands.

 Assembly is handled automatically by FEM::makeSystem.
- solveProblem entry point for the simulator.

The Poisson⁰ Class

```
#ifndef Poisson0 h IS INCLUDED
#define Poisson0 h IS INCLUDED
#include <FEM.h> // FEM algorithms, class FieldFE, GridFE
#include <DegFreeFE.h> // degree of freedom book—keeping
#include <LinEqAdmFE.h> // linear systems: storage and solution
class Poisson0 : public FEM {
protected:
 // general data:
 Handle(GridFE) grid; // pointer to a finite element grid
 Handle(DegFreeFE) dof; // trivial book-keeping for a scalar PDE
 Handle(FieldFE) u; // finite element field, primary unknown
 Vec(real) linsol; // solution of the linear system
 Handle(LinEqAdmFE) lineq; // linear system: storage and solution
 void fillEssBC (); // set boundary conditions u=g
 virtual void integrands(ElmMatVec& elmat, const FiniteElement& fe);
```

The Poisson 0 Class, cont'd

```
public:
 Poisson0 ();
 ~Poisson0 () {}
 void scan (); // read and initialize data
 void solveProblem (); // main driver routine
 void resultReport (); // write comparison with analytical sol.
  real f(real x, real y); // source term in the PDE
  real k(real x, real y); // coefficient in the PDE
  real g(real x, real y); // essential boundary conditions
};
#endif
```

Poisson0::scan

```
void Poisson0:: scan ()
{
    // extract input from the command line:
    int nx, ny; // number of nodes in x— and y—direction
    initFromCommandLineArg ("—nx", nx, 6); // read nx, default : nx=6
    initFromCommandLineArg ("—ny", ny, 6);
    String elm_tp;
    initFromCommandLineArg ("—elm", elm_tp, "ElmB4n2D");
```

- Read information about the number of nodes in the x- and ydirections.
- Read information about the element type.

Poisson0::scan, cont'd

- Construct input strings to the box preprocessor, specifying the domain, the partition, and the element type.
- Call the box preprocessor to generate the grid.

Poisson0::scan, cont'd

```
u.rebind (new FieldFE (*grid,"u")); // allocate, set name="u"

dof.rebind (new DegFreeFE (*grid, 1)); // 1 unknown per node

lineq.rebind (new LinEqAdmFE()); // Ax=b system and solvers

linsol .redim (grid—>getNoNodes()); // redimension linsol

lineq—>attach (linsol); // use linsol as x in Ax=b

// banded Gaussian elimination is the default solver in lineq
}
```

- Set u to point to a new finite element field over the grid.
- Set dof to point to a new DegFreeFE object.
- Set lineq to point to a new LinEqAdmFE object.
- Redimension the unknown vector in the linear system and attach it to lineq.

Poisson:FillEssBC

```
void Poisson0:: fillEssBC ()
 dof—>initEssBC ();
                    // init for assignment below
 const int nno = grid->getNoNodes(); // no of nodes
 Ptv(real) x;
                       // a nodal point
 for (int i = 1; i \le nno; i++) {
   // is node i subjected to any boundary indicator?
   if (grid—>boNode (i)) {
     x = grid->getCoor (i); // extract coor. of node i
     dof->fillEssBC (i, g(x(1), x(2))); // u=g at boundary nodes
 dof—>printEssBC (s_o, 2); // debug output
```

Initialize assignment of essential boundary conditions

Loop through all nodes and if a node is on the boundary, set essential boundary condition

Poisson0::integrands

The integrands function:

- 1. Evaluate the Jacobian determinant times the integration weight.
- 2. Find the global coordinates of the current integration point.
- 3. Evaluate the f and k functions at this global point.
- 4. For i = 1 to the number of basis functions (element nodes)
 - For j = 1 to the number of basis functions
 - Add the appropriate value to the *elemental* matrix.
 - Add the appropriate value to the elemental vector.

Poisson0::integrands, cont'd

```
void Poisson0::integrands (ElmMatVec& elmat, const FiniteElement& fe)
 // find the global coord. xy of the current integration point:
Ptv(real) xy = fe.getGlobalEvalPt();
const real x = xy(1); const real y = xy(2); // 2D specific
const real f_value = f(x,y);
                               // 2D specific
const real k value = k(x,y);
                                              // 2D specific
int i, j;
const int nbf = fe.getNoBasisFunc(); // = no of nodes in element
const real detJxW = fe.detJxW(); // Jacobian * intgr. weight
for (i = 1; i \le nbf; i++)
  for (i = 1; i \le nbf; i++)
    elmat.A(i,j) += k_value*(fe.dN(i,1)*fe.dN(j,1) // 2D specific
                           + fe.dN(i,2)*fe.dN(i,2))*detJxW;
  elmat.b(i) += fe.N(i)*f value*detJxW;
```

Problem-Dependent Functions

Specific test case:

$$f(x,y) = -2x(x-1) - 2y(y-1),$$

$$g(x,y) = 0, \quad k(x,y) = 1.$$

Analytical solution: u(x,y) = x(x-1)y(y-1)

The functions f, g, and k are then:

```
real Poisson0:: f (real x, real y)
{ return -2*x*(x-1)-2*y*(y-1); }

real Poisson0:: g (real /*x*/, real /*y*/)
{ return 0; }

real Poisson0:: k (real /*x*/, real /*y*/)
{ return 1; }
```

Solving the Problem

```
void Poisson0:: solveProblem () // main routine of class Poisson0
 fillEssBC ();
               // set essential boundary conditions
 makeSystem (*dof, *lineq); // assembly algorithm from class FEM
  linsol . fill (0.0); // init start vector (good habit)
 lineq->solve(); // solve linear system
 dof->vec2field (linsol, *u); // load solution (linsol) into u
int main (int argc, const char* argv[])
  initDiffpack (argc, argv);
 Poisson0 simulator; // make a simulator object
 simulator.scan (); // read and initialize data
 simulator.solveProblem (); // main routine: compute the solution
 simulator.resultReport (); // compare with exact solution
 return 0;
```

The FEM:: makeSystem algorithm:

initialize global linear system:

set
$$A_{i,j} = 0$$
 for $i, j = 1, ..., n$, set $b_i = 0$ for $i = 1, ..., n$ loop over all elements:

for
$$e=1,\ldots,m$$

$$\text{set } \tilde{A}_{r,s}^{(e)}=0,\,r,s=1,\ldots,n_e,\qquad\text{set } \tilde{b}_r^{(e)}=0,\,r=1,\ldots,n_e$$
 loop over numerical integration points:

for
$$k = 1, \ldots, n_I$$

evaluate $\tilde{N}_r(\xi_k)$, derivatives of \tilde{N}_r wrt. ξ and x, J

integrands:

for
$$r=1,\ldots,n_e$$

$$\tilde{A}_{r,s}^{(e)}:=\tilde{A}_{r,s}^{(e)}+\frac{d\tilde{N}_r}{dx}\frac{\tilde{N}_s}{dx}\det J\ w_k$$

$$\tilde{b}_r^{(e)}:=\tilde{b}^{(e)}+f(x^{(e)}(\xi_k))N_r\det J\ w_k$$

fillEssBc:

for
$$r = 1, \ldots, n_e$$

if node r has an essential boundary condition then modify $\tilde{A}_{r,s}^{(e)}$ and $\tilde{b}_{r}^{(e)}$ due to this condition

assemble:

for
$$r=1,\ldots,n_e$$

$$for \ s=1,\ldots,n_e$$

$$A_{q(e,r),q(e,s)}:=A_{q(e,r),q(e,s)}+\tilde{A}_{r,s}^{(e)}$$

$$b_{q(e,r)}:=b_{q(e,r)}+\tilde{b}_r^{(e)}$$

Summary

Concluding remarks on FEM:

- The simulator class provides only functionality that is closely related to the concrete problem.
- This functionality is part of a larger framework implemented in the base class FEM.
- The use of inheritance and virtual functions is crucial in this software architecture.

Software abstractions:

- Brings the implementation closer to the mathematics
- Hides unnecessary details
- FDM: nice, but sometimes trivial and obfuscating
- FVM/FEM: important, in particular for unstructured grids

In this course: since we mostly study trivial examples, it might be hard to see the benefit