敖旭扬 PB18071477 2021 年 1 月 7 日

第一题解:

(a) 由 Gauss 顺序消元法的步骤和过程,第一次迭代后 A 的第一列完成消去时,有

$$a_{ij}^{(1)} = a_{ij} - \frac{a_{i1}}{a_{11}} a_{1j}, \quad i, j = 2, 3, \dots, n$$
 (1)

由于 A 是对称矩阵且满足 $a_{11} \neq 0$, 所以有

$$a_{ij} = a_{ji}, \quad i, j = 1, 2, 3, \dots, n$$
 (2)

所以

$$a_{ij}^{(1)} = a_{ij} - \frac{a_{i1}}{a_{11}} a_{1j} = a_{ji} - \frac{a_{1i}}{a_{11}} a_{j1} = a_{ji} - \frac{a_{j1}}{a_{11}} a_{1i} = a_{ji}^{(1)}, \quad i, j = 2, 3, \dots, n$$
 (3)

从而 $A^{(1)}$ 是对称的。

(b) 计算一个正定 (positive definite) 矩阵 LU 分解的算法如下:

```
%% 正定矩阵(Positive Definite Matrix)LU分解算法
   function [L, U] = LU_PDM(A)
 3
       [~, n] = size(A);
       L = zeros(n, n);
 4
       U = zeros(n, n);
 5
       for i = 1:n
 6
           for j = i + 1:n
 7
 8
                for k = j:n
                    A(j, k) = A(j,k)-A(i,k)*A(j,i)/A(i,i);
9
                    A(k, j) = A(j,k);
10
11
                end
12
           end
13
           for k = i:n
                U(i, k) = A(i, k);
14
                L(k, i) = A(i, k)/U(i, i);
15
16
            end
17
       end
18
   end
```

(c) 使用 Cholesky 分解解方程组 Ax = b 的 MATLAB 程序显示如下:

```
clear, clc, clf
   A = [4, -2, 4, 2; -2, 10, -2, -7; 4, -2, 8, 4; 2, -7, 4, 7];
3 \mid b = [8; 2; 16; 6];
4 \mid x = deEquations(A,b)
   %% 运行的输出结果
5
   % x =
6
7
   %
8
   %
           1
9
   %
           2
10
  1%
           1
  %
           2
11
12
13 | %% 用 Cholesky 分 解 解 方 程 组 Ax=b
14 | function [X] = deEquations(A,b)
15 \mid N = length(A);
16 L=cholesky(A);
17 | Lt=L';
18 \mid X=zeros(N, 1);
19 Y=zeros(N, 1);
20 | for j=1:N
        Y(j)=(b(j)-L(j,1:j-1)*Y(1:j-1))/L(j,j);
21
22
   end
23
   for k=N:-1:1
24
        X(k) = (Y(k) - Lt(k, k+1:N) * X(k+1:N)) / Lt(k,k);
25
   end
26
   end
27
28
   function [L]=cholesky(A)
29
   N=length(A);
30
   for i=1:N
31
        A(i,i) = sqrt(A(i,i) - A(i,1:i-1) * A(i,1:i-1)');
32
        if A(i,i) == 0
33
            ME=MException('Zero Element', 'A(%d,%d) = 0',i,i);
34
            throw(ME);
```

```
35 end
36 for j=i+1:N
37 A(j,i)=(A(j,i)-A(j,1:i-1)*A(i,1:i-1)')/A(i,i);
38 end
39 end
40 L = tril(A); %取下三角部分
41 end
```

根据上述程序在命令行输出的结果(以注释形式写在了上面的代码中), 原题所求的解为:

$$x = \begin{bmatrix} 1 \\ 2 \\ 1 \\ 2 \end{bmatrix}$$

第二题 解:

(a) 证明:

Richardson 迭代方法的迭代关系式为:

$$x^{(k+1)} = \omega I(\frac{1}{\omega}I - A)x^{(k)} + \omega Ib = (I - \omega A)x^{(k)} + \omega b \tag{4}$$

令

$$G_{\omega} = I - \omega A \tag{5}$$

则

$$x^{(k+1)} = G_{\omega}x^{(k)} + \omega b \tag{6}$$

根据课本第五章开头处对迭代法收敛性的讨论可知,上述 Richardson 迭代方法收敛的充要条件为谱半径 $\rho(G_{\omega})<1$,又 A 是正定矩阵,所以它的特征值均为正数,由已知, G_{ω} 的特征值为 $1-\omega\lambda_i$,其中 $0<\lambda_1\leq\lambda_i\leq\lambda_n$,则

$$\rho(G_{\omega}) = \max|1 - \omega \lambda_i| < 1 \tag{7}$$

则对任意 λ_i , 需要满足

$$0 < \omega \lambda_i < 2 \tag{8}$$

即

$$\omega < 2/\lambda_n \tag{9}$$

(b) 由 (a) 得,

$$\rho(G_{\omega}) = \max|1 - \omega \lambda_i| = \max(1 - \omega \lambda_1, \omega \lambda_n - 1) \tag{10}$$

应该使得收敛速度尽量快, 所以 $\rho(G_{\omega})$ 应该尽量小, 即取

$$\rho(G_{\omega}) = \min_{\omega} \max(1 - \omega \lambda_1, \omega \lambda_n - 1) \tag{11}$$

所以 ω 的最佳值为

$$\omega_b = \arg\min_{\omega} \max(1 - \omega \lambda_1, \omega \lambda_n - 1) = \frac{2}{\lambda_1 + \lambda_n}$$
(12)

且

$$\rho(G_{\omega}) = \begin{cases}
1 - \omega \lambda_1 & \omega \leq \omega_b \\
\frac{\lambda_n - \lambda_1}{\lambda_n + \lambda_1} & \omega = \omega_b \\
\omega \lambda_n - 1 & \omega \geq \omega_b
\end{cases}$$
(13)

(c) 使用 Richardson 迭代方法解 Ax = b 的 MATLAB 程序显示如下:

```
clear, clc, clf
 2 %% 构造A和b
 3 \mid D = orth(rand(5, 5));
4 \mid B = diag([1, 2, 3, 4, 5]);
 5 \mid A = D \setminus B * D;
 6 \mid b = rand(5, 1);
   xexact = A \ b; % x精确解
 7
   %% 使用多个omega迭代求解
9
   lambda1 = 1;
10
11
   lambdan = 5;
   omegab=2/(lambda1+lambdan);
12
13
   eponch=1001;
   omegalist = linspace(0.001, 0.399, eponch);
14
   count=zeros(eponch,1)';
16
   for i = 1:eponch
       count(i)=Richardson(A,b,xexact,omegalist(i));
17
18
   end
19
20 %% 输出结果
```

```
21
   [mink,idx]=min(count);
22 | omegalist(idx);
23 | fp=fopen('p2c_out.txt','w');
24 | fprintf (fp, 'bestomega = %f\tk=%d\n', omegalist(idx), mink);
25
   fprintf (fp, 'omegab = f \times d \cdot n', omegab, Richardson(A,b,
      xexact,omegab));
26 | semilogy(omegalist, count, 'k.-');
   xlabel('\omega');
27
   ylabel('迭代次数');
28
29
   %% 返回Richardson迭代方法的迭代次数
30
31
   function [k]=Richardson(A,b,xexact,omega)
32
       G = eye(5) - omega * A;
33
       x = zeros(5, 1);
34
       k = 0;
       while norm(x - xexact) > 1e-13
35
36
           x = G * x + omega * b;
37
           k = k + 1;
38
       end
   end
39
```

上述程序某一次执行输出的结果为:

```
bestomega = 0.333330 k=70
omegab = 0.333333 k=70
```

同时输出的收敛迭代次数随 ω 变化的 semilogy 图为:

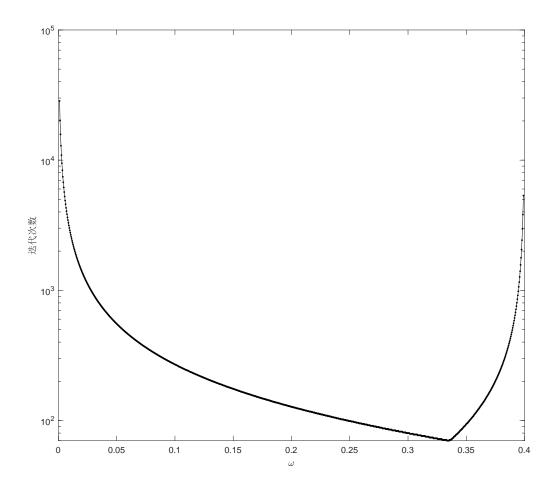


图 1: 寻找 Richardson 迭代方法的最佳 ω

这一结果验证了最佳 ω 是 ω _b, 它使得收敛速度最快。

第三题 解:

(a) $I(f) = \int_a^b f(x) dx$ 关于积分节点 $\{x_1, x_2, \cdots, x_n\}$ 的 Gauß 积分的数值积分公式为

$$I_n(f) = \sum_{i=1}^n \alpha_i f(x_i) \tag{14}$$

n=6 时,代数精度不超过 2n-1=11,所以依次取线性不相关的 $f(x)=1,x^2,\cdots,x^{10}$,即 $f_k(x)=x^{2k},\quad k=0,1,\cdots,5$,根据

$$I_n(f) = I(f) \tag{15}$$

取 a = -1, b = 1 得

$$\sum_{i=1}^{6} \alpha_i x_i^{2k} = \int_{-1}^{1} x^{2k} dx = \frac{2}{2k+1} \quad , \quad k = 0, 1, \dots, 5$$
 (16)

又根据 Gauß 的节点和积分权重关于原点对称,即

$$\alpha_i = \alpha_{i+3}, \quad x_i = -x_{i+3}, \quad i = 1, 2, 3$$
 (17)

所以式 (16) 变为

$$\sum_{i=1}^{6} \alpha_i x_i^{2k} = 2 \sum_{i=1}^{3} \alpha_i x_i^{2k} = \frac{2}{2k+1} \quad , \quad k = 0, 1, \dots, 5$$
 (18)

即

$$\sum_{i=1}^{3} \alpha_i x_i^{2k} = \frac{1}{2k+1} \quad , \quad k = 0, 1, \dots, 5$$
 (19)

写成方程组则表示为

$$\begin{cases}
\sum_{i=1}^{3} \alpha_{i} = 1 \\
\sum_{i=1}^{3} \alpha_{i} x_{i}^{2} = \frac{1}{3} \\
\sum_{i=1}^{3} \alpha_{i} x_{i}^{4} = \frac{1}{5} \\
\sum_{i=1}^{3} \alpha_{i} x_{i}^{6} = \frac{1}{7} \\
\sum_{i=1}^{3} \alpha_{i} x_{i}^{8} = \frac{1}{9} \\
\sum_{i=1}^{3} \alpha_{i} x_{i}^{10} = \frac{1}{11}
\end{cases} (20)$$

(b) 方程组 (20) 中一共有 $\langle \alpha_1, \alpha_2, \alpha_3, x_1, x_2, x_3 \rangle$ 共 6 个变量,且有

$$\frac{\partial \left(\sum_{i=1}^{3} \alpha_{i} x_{i}^{2k}\right)}{\partial \alpha_{j}} = x_{j}^{2k}$$

$$\frac{\partial \left(\sum_{i=1}^{3} \alpha_{i} x_{i}^{2k}\right)}{\partial x_{j}} = 2k\alpha_{j} x_{j}^{2k-1} \quad (k \neq 0)$$

$$\frac{\partial \left(\sum_{i=1}^{3} \alpha_{i}\right)}{\partial x_{j}} = 0$$

其中 j = 1, 2, 3; k = 0, 1, 2, 3, 4, 5

所以非线性方程组 (20) 的 Jacobian 的表达式为

$$\mathbf{J} = \frac{\partial \left(\sum_{i=1}^{3} \alpha_{i}, \sum_{i=1}^{3} \alpha_{i} x_{i}^{2}, \sum_{i=1}^{3} \alpha_{i} x_{i}^{4}, \sum_{i=1}^{3} \alpha_{i} x_{i}^{6}, \sum_{i=1}^{3} \alpha_{i} x_{i}^{8}, \sum_{i=1}^{3} \alpha_{i} x_{i}^{10}\right)}{\partial (\alpha_{1}, \alpha_{2}, \alpha_{3}, x_{1}, x_{2}, x_{3})}$$

$$= \begin{bmatrix}
1 & 1 & 1 & 0 & 0 & 0 \\
x_{1}^{2} & x_{2}^{2} & x_{3}^{2} & 2\alpha_{1} x_{1} & 2\alpha_{2} x_{2} & 2\alpha_{3} x_{3} \\
x_{1}^{4} & x_{2}^{4} & x_{3}^{4} & 4\alpha_{1} x_{1}^{3} & 4\alpha_{2} x_{2}^{3} & 4\alpha_{3} x_{3}^{3} \\
x_{1}^{6} & x_{2}^{6} & x_{3}^{6} & 6\alpha_{1} x_{1}^{5} & 6\alpha_{2} x_{2}^{5} & 6\alpha_{3} x_{3}^{5} \\
x_{1}^{8} & x_{2}^{8} & x_{3}^{8} & 8\alpha_{1} x_{1}^{7} & 8\alpha_{2} x_{2}^{7} & 8\alpha_{3} x_{3}^{7} \\
x_{1}^{10} & x_{1}^{10} & x_{3}^{10} & 10\alpha_{1} x_{1}^{9} & 10\alpha_{2} x_{2}^{9} & 10\alpha_{3} x_{3}^{9}
\end{bmatrix}$$

$$(21)$$

(c) 求解 n=6 情况下的 Gauß 积分的积分节点和积分权重的 MATLAB 程序显示如下:

```
clear, clc, clf
   x = linspace(-1, 1, 6);
   w = zeros(3, 1)';
3
4
5
   for i = 1:3
6
        syms t;
7
       Fi = @(t) alpha_fun(t, x, i);
       % 使用课本130页下方的\alpha公式初始化w
8
9
       w(i) = int(Fi, t, -1, 1);
   end
10
11
12
   x = x(1:3);
   dwdx = ones(6);
13
14
   while (max(dwdx) > 1e-10)
15
16
        J = Jacobian(w, x);
       b = cal b(w, x);
17
18
       dwdx = J \setminus b;
19
       dwdx = dwdx';
20
       w = w - dwdx(1:3);
21
       x = x - dwdx(4:6);
22
   \quad \text{end} \quad
23 | %% 輸出w和x
```

```
24 w
25
   Х
26
27
   function f = alpha_fun(t, x, i)
28
       n = length(x);
29
       fac1 = t - x;
30
       fac2 = x(i) - x;
31
       f = prod(fac1((1:n) ~= i))/prod(fac2((1:n) ~= i));
32
   end
   %% 第三题(b)中非线性方程组的Jacobian的表达式
33
   function [J] = Jacobian(w,x)
34
35
       J = [1, 1, 1, 0, 0, 0;
36
            x.^2, x .* w .* 2;
            x.^4,(x.^3).*w.*4;
37
38
           x.^6, (x.^5) .* w .* 6;
39
            x.^8, (x.^7) .* w .* 8;
            x.^10,(x.^9) .* w .* 10
40
           ];
41
42
   end
43
   function [b] = cal_b(w, x)
44
45
       b = zeros(6, 1);
46
       for i = 1:6
            b(i) = sum(w .* (x.^(2*i-2)))-1/(2*i-1);
47
48
       end
49
   \quad \text{end} \quad
```

上述程序在命令行的输出结果为

```
w =
    0.171324492379170    0.360761573048139
    0.467913934572691
x =
    -0.932469514203152    -0.661209386466265
    -0.238619186083197
```

所以积分节点和积分权重为

```
(x_1, x_2, x_3) = (-0.932469514203152, -0.661209386466265, -0.238619186083197)
(\alpha_1, \alpha_2, \alpha_3) = (0.171324492379170, 0.360761573048139, 0.467913934572691)
```

其中其他几项由式 (17) 得到。

这里初始权重的选取是根据课本(第三版)第 130 页下方的 $\alpha_i^{(n)}$ 公式计算而来(但是把其中的 $x_i^{(n)}$ 用 x 初始值代替了)。

(d) MATLAB 程序显示如下:

```
clear, clc, clf
2
3
   [w,x]=gauss_w_x(5)
4
5 | function [w,x]=gauss_w_x(n)
6 halfn=ceil(n/2);
7 | w=zeros(halfn,1)';
8
  x = che(1:n,n);
   syms t;
10 | for i=1:halfn
11
       Fi=@(t) alpha_fun(t,x,i);
       w(i)=int(Fi,t,-1,1);
12
13 | end
14 | x=x(1:floor(n/2));
15 dx=ones(halfn);
16 dw=ones(halfn);
17 \mid J = Jacobian(x, w, n);
18
   while (max([dx,dw]) > 1e-5)
19
       J = Jacobian(x, w, n);
20
       b = f(x, w, n);
21
       [dx,dw] = grad(J,b);
22
       x = x - dx;
23
       w = w - dw;
24 | end
25 end
26 %% Chebyshev点
27 \mid function [xj] = che(j,n)
```

```
28
        xj = cos(j*pi/n);
29
   end
   function [dx,dw] = grad(J,b)
30
31 \mid n=size(J,1);
32 \mid fn=floor(n/2);
33 \mid dxdw=J \setminus b;
34 | dxdw=dxdw';
35 \mid dx = dxdw(1:fn);
36 \mid dw = dxdw (fn+1:n);
37 \mid \mathbf{end} \mid
38 | function f=alpha_fun(t,x,i)
39 \mid n = length(x);
40 \mid fac1=t-x;
41 \mid fac2=x(i)-x;
   f=prod(fac1((1:n) ~= i))/prod(fac2((1:n) ~= i));
43
   end
44
   function [J] = Jacobian(x ,w,n)
45
   J=zeros(n,n);
46
   m=length(w);
47
   if mod(n,2)
48
49
        factor=ones(n,1)';
        factor(m)=1/2;
50
        xp=[x,0];
51
52
        m1zeros=zeros(m,1)';
53
        m2ones=ones(length(x),1)';
        J(1,:) = [m1zeros, m2ones].*factor;
54
        for i=2:n
55
             J(i,:) = [(xp .^ (2*i-3)) .* w .* (2*i-2) , x .^
56
                (2*i-2)].*factor;
57
        end
        J.
58
59
   else
        for i=1:n
60
             J(i,:) = [(x .^ (2*i-3)) .* w .* (2*i-2) , x .^
61
```

```
(2*i-2)];
62
        end
63 end
64 end
65
66 \mid \mathbf{function} \mid [b] = f(x, w, n)
67 \mid b = zeros(n,1);
   if mod(n,2)
68
       w=w(1:end-1);
69
70
       b(1) = sum(w) - 1;
71
       for i=2:n
72
             b(i)=sum(w .* (x .^ (2*i-2))) - 1/(2*i-1);
73
        end
74 else
75
       for i=1:n
76
             b(i)=sum(w .* (x .^ (2*i-2))) - 1/(2*i-1);
77
        end
78 end
79 \mid \texttt{end}
```