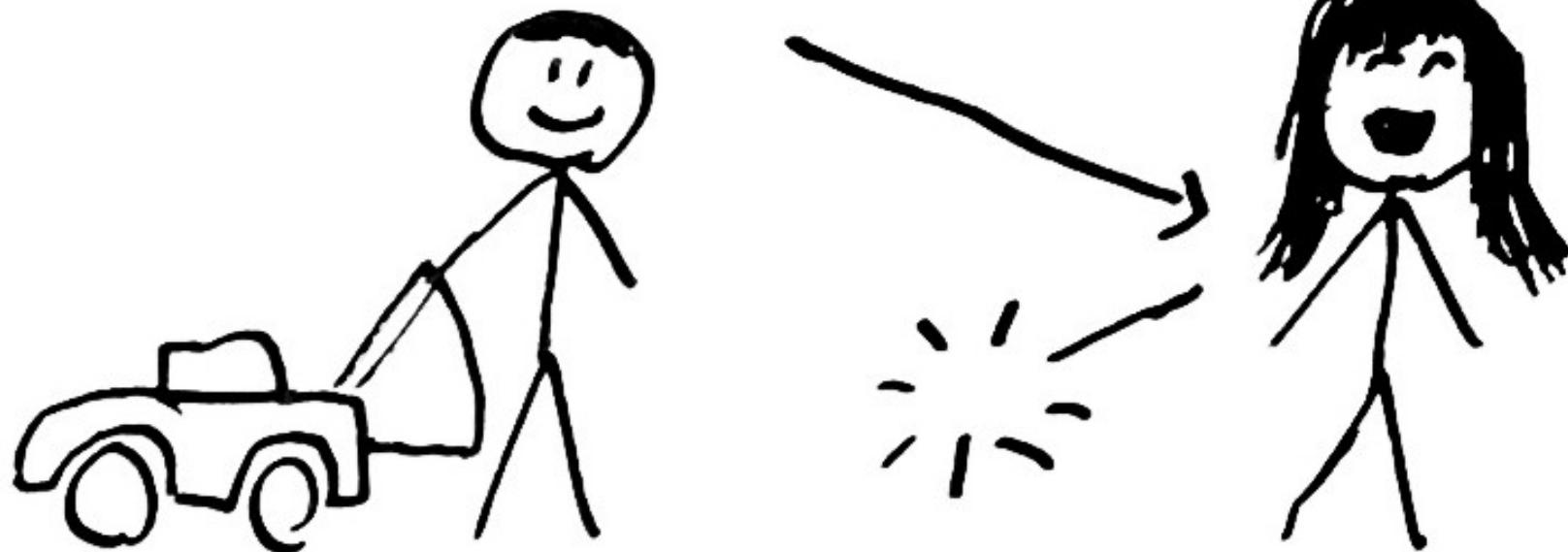


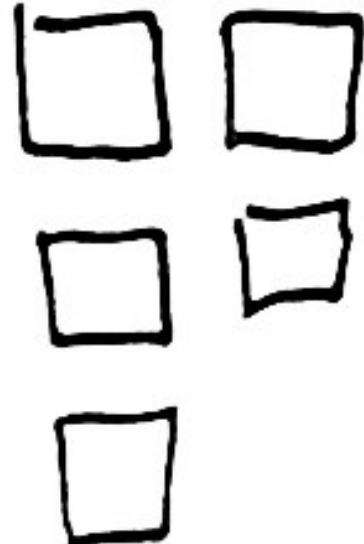
Call Me Maybe!



H

@aphyr  
Kyle Kingsbury  
"Jepsen Networks"

Backend engineer at



Factual

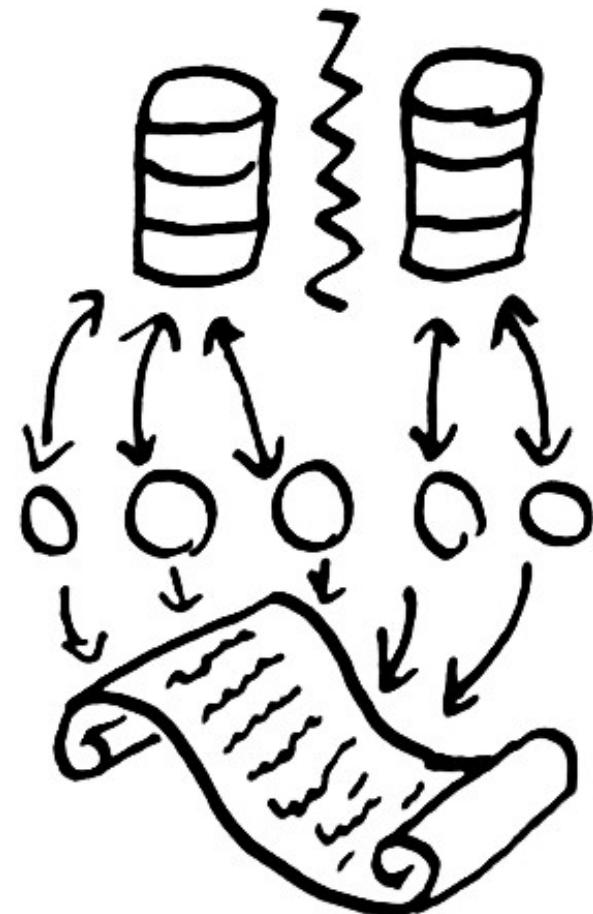
Jepsen II

Linearizable

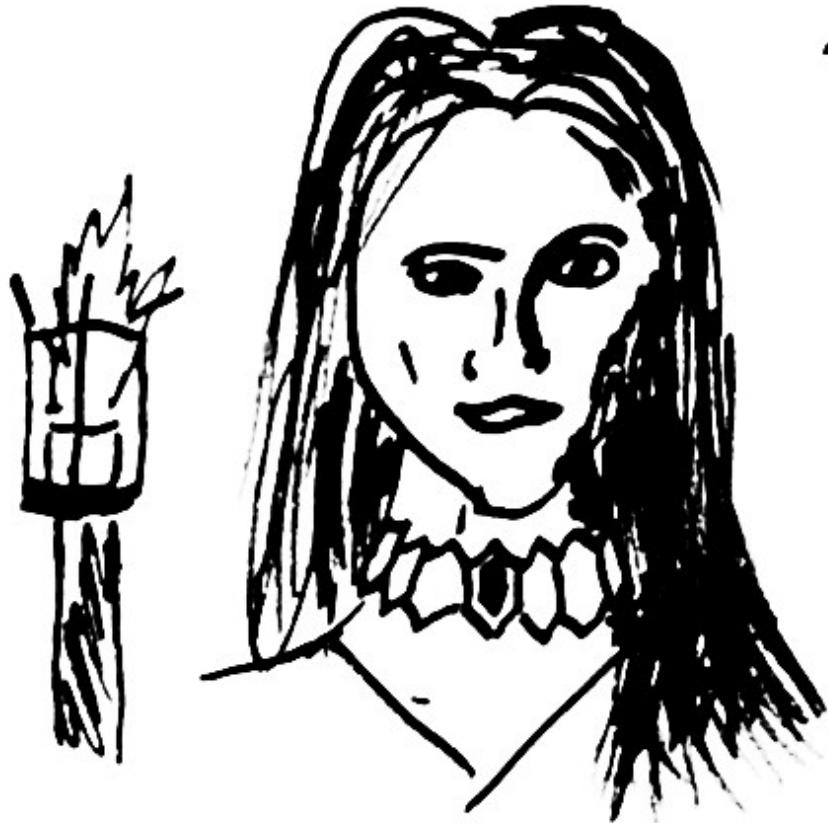
Boogaloo

# Jepsen

Measures  
Systems  
Under  
Stress

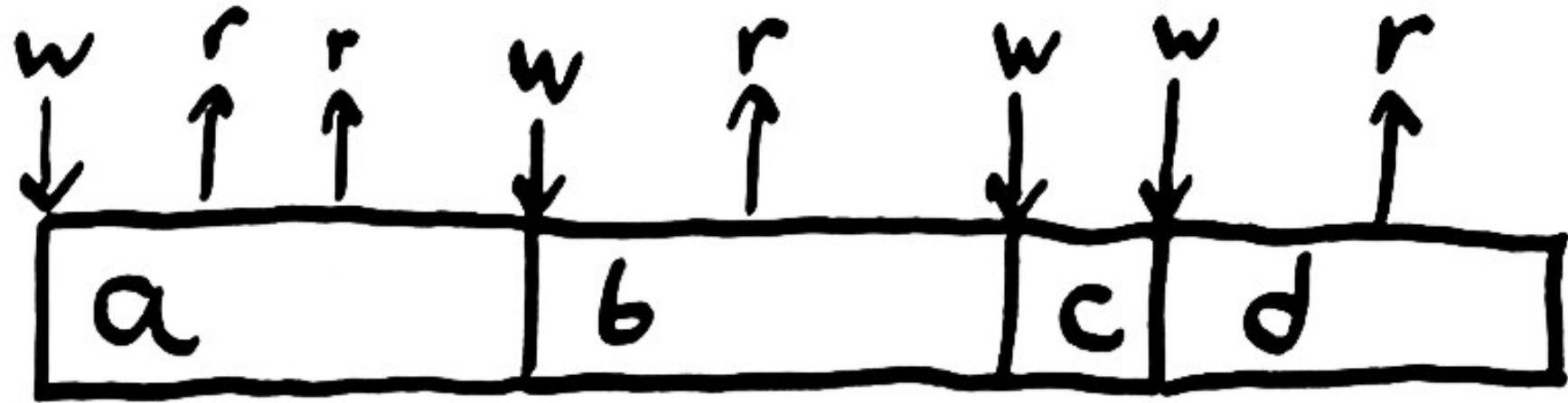


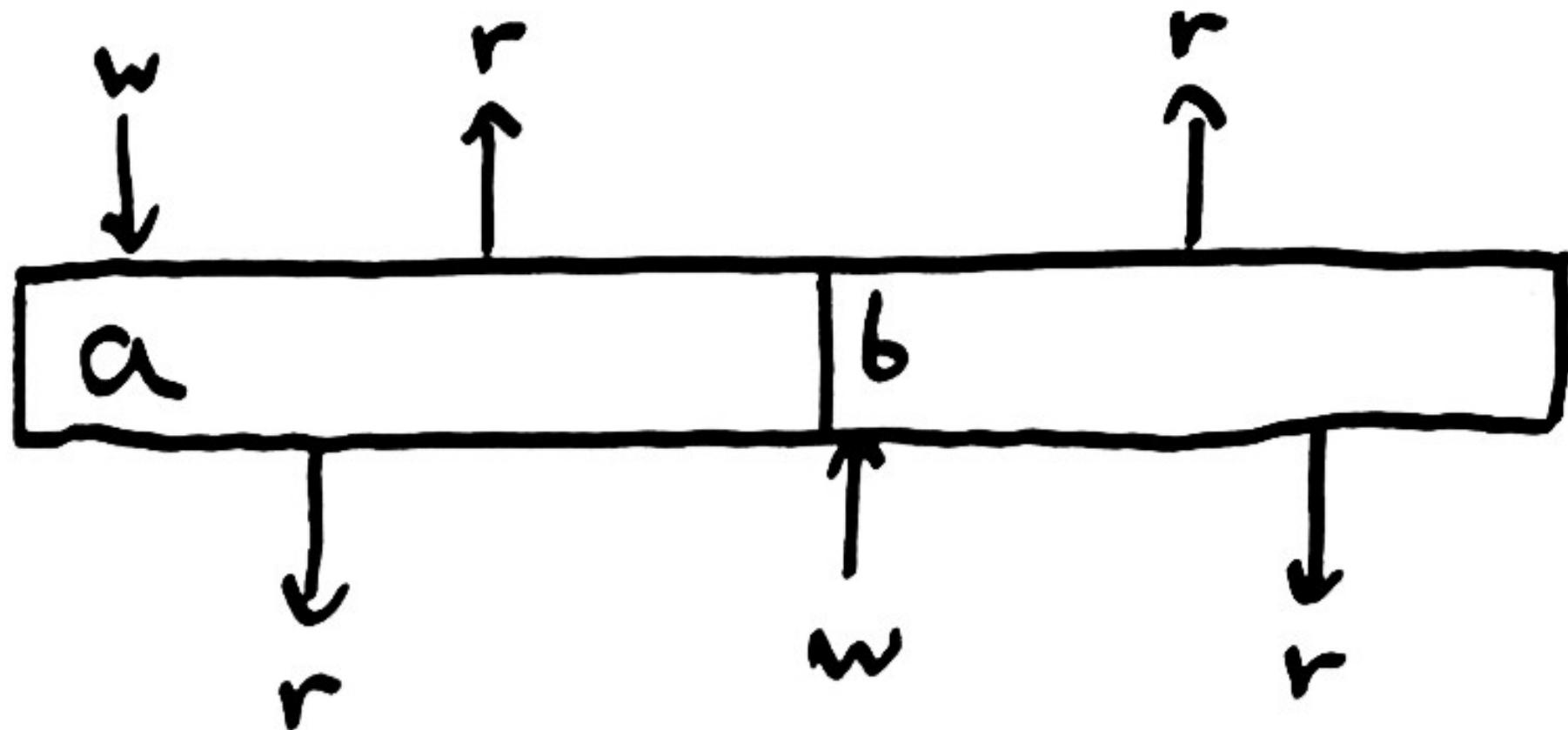
*Why Care?*

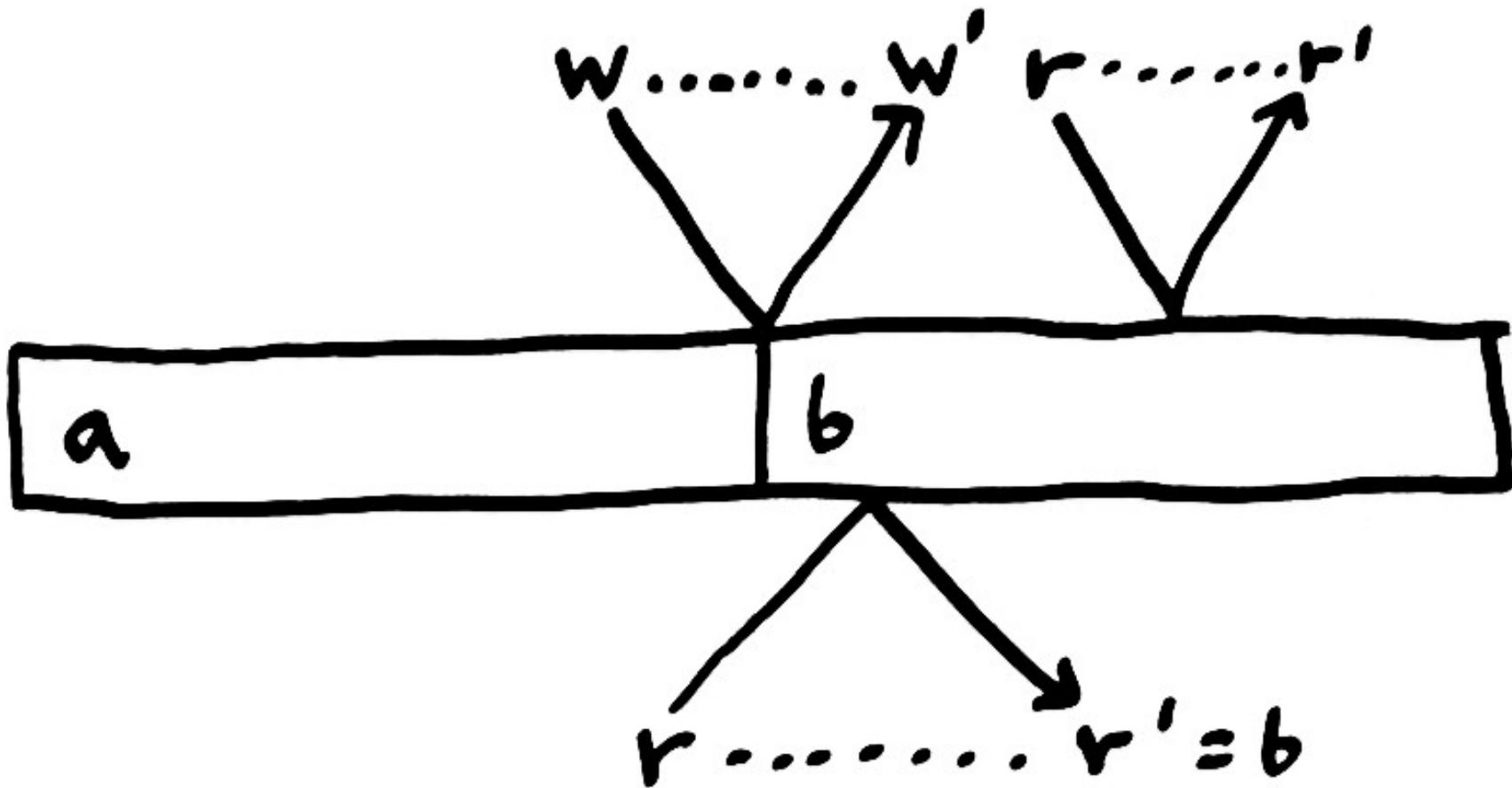


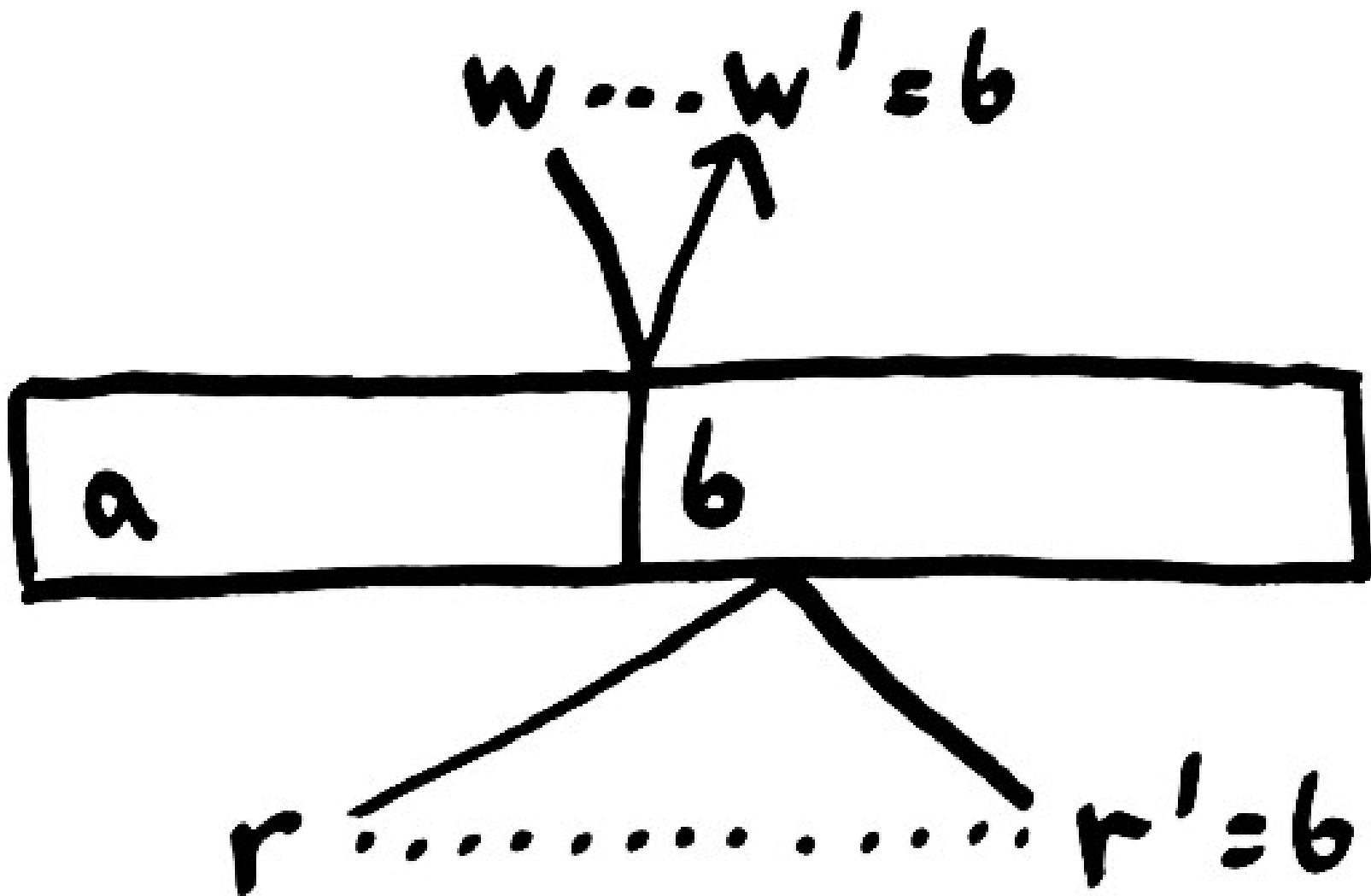
The net is  
dark, and full of  
**TERAHS**

*Intuitive Correctness*



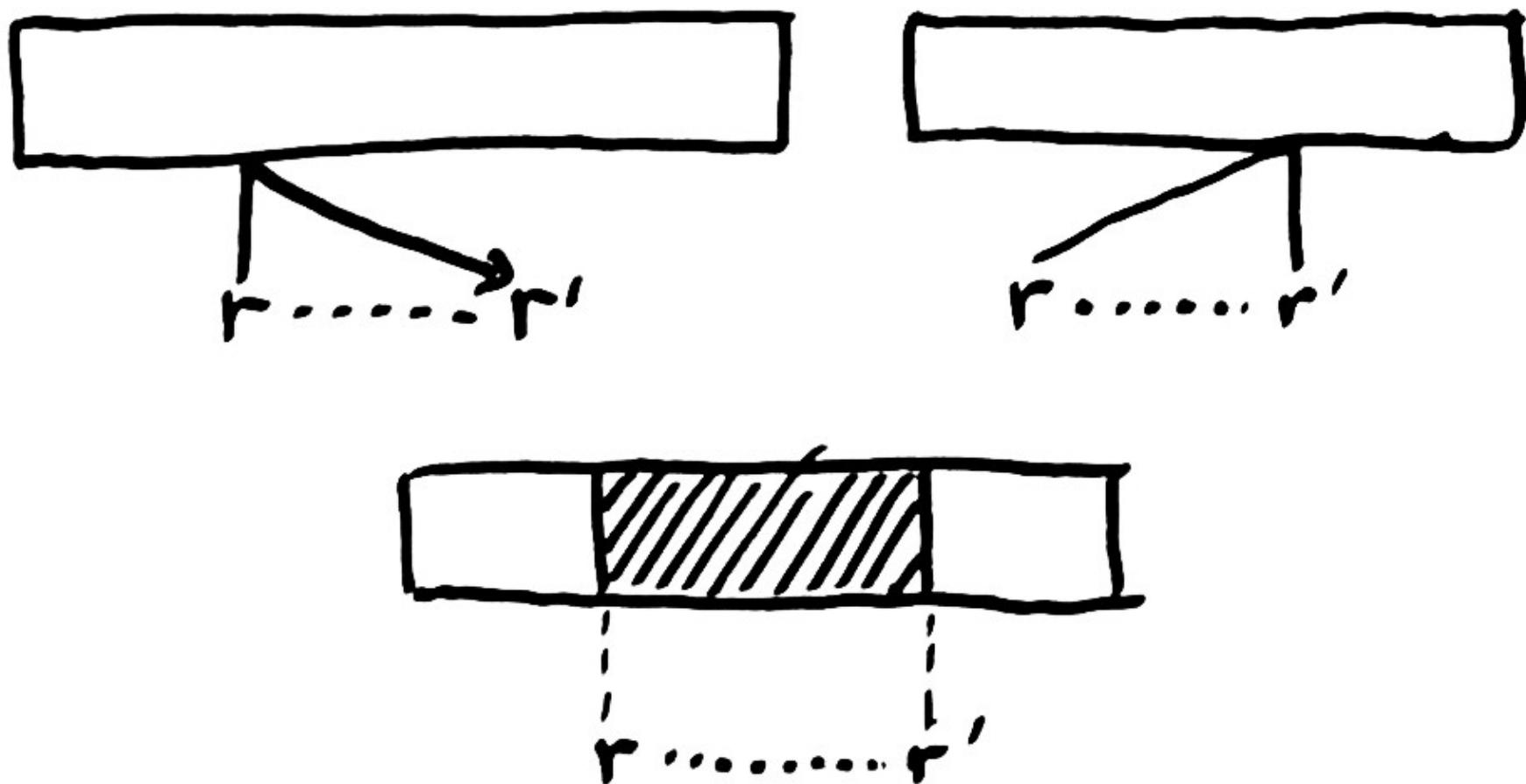




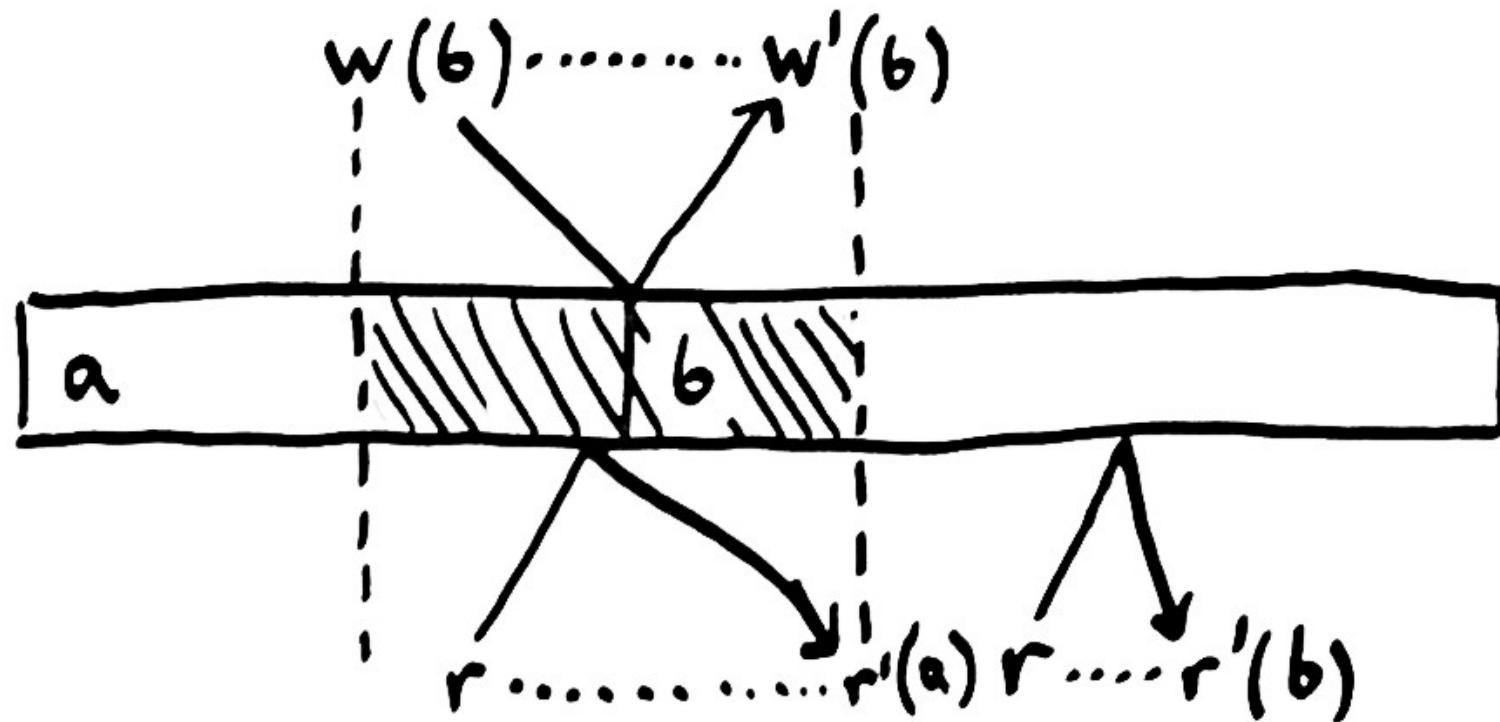


Synchronization  
delay implies  
ambiguous orders

But there are  
finite bounds!



Once operation is complete,  
it will be visible to all.



No stale reads

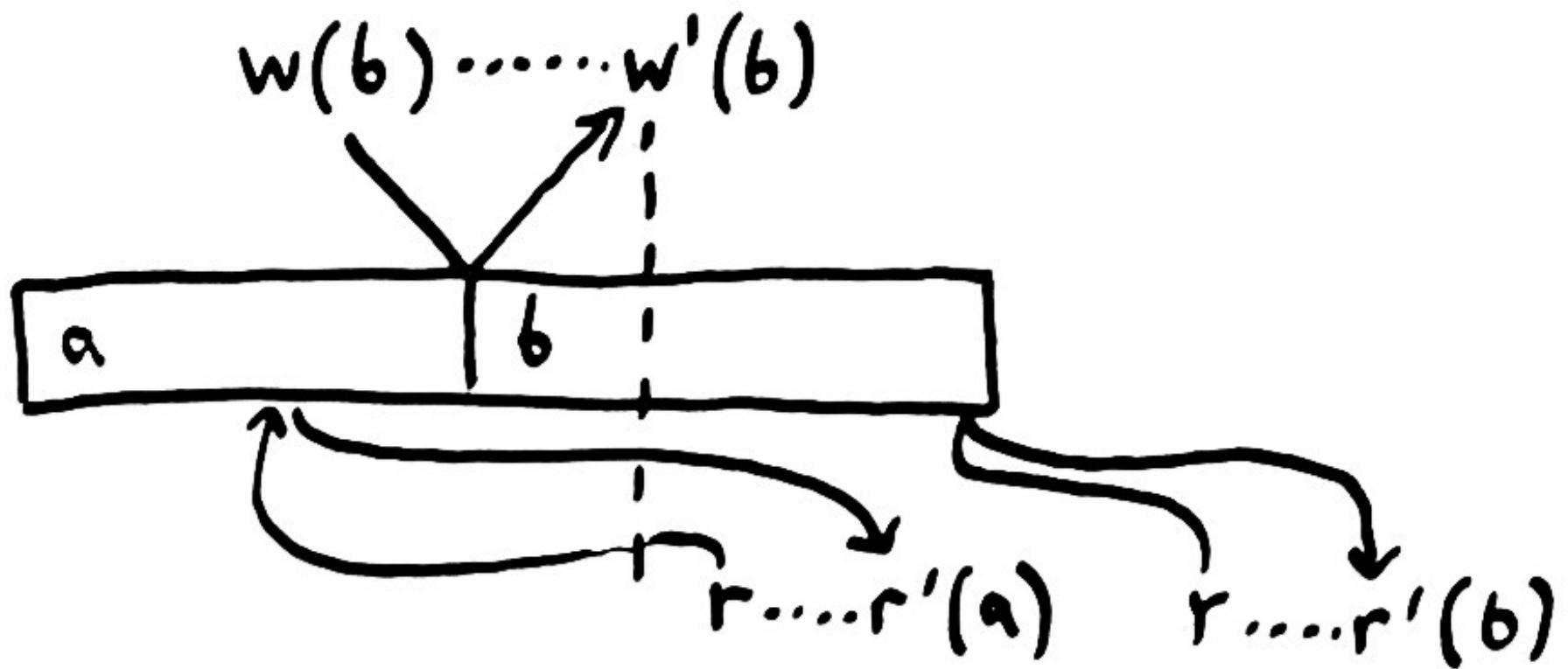
Safe state mutation

Everyone sees the  
Same order of states.

Same properties as  
atomic ops in a  
multithreaded lang.

Sequential Consistency:

Everyone agrees on  
order, but not on  
time.

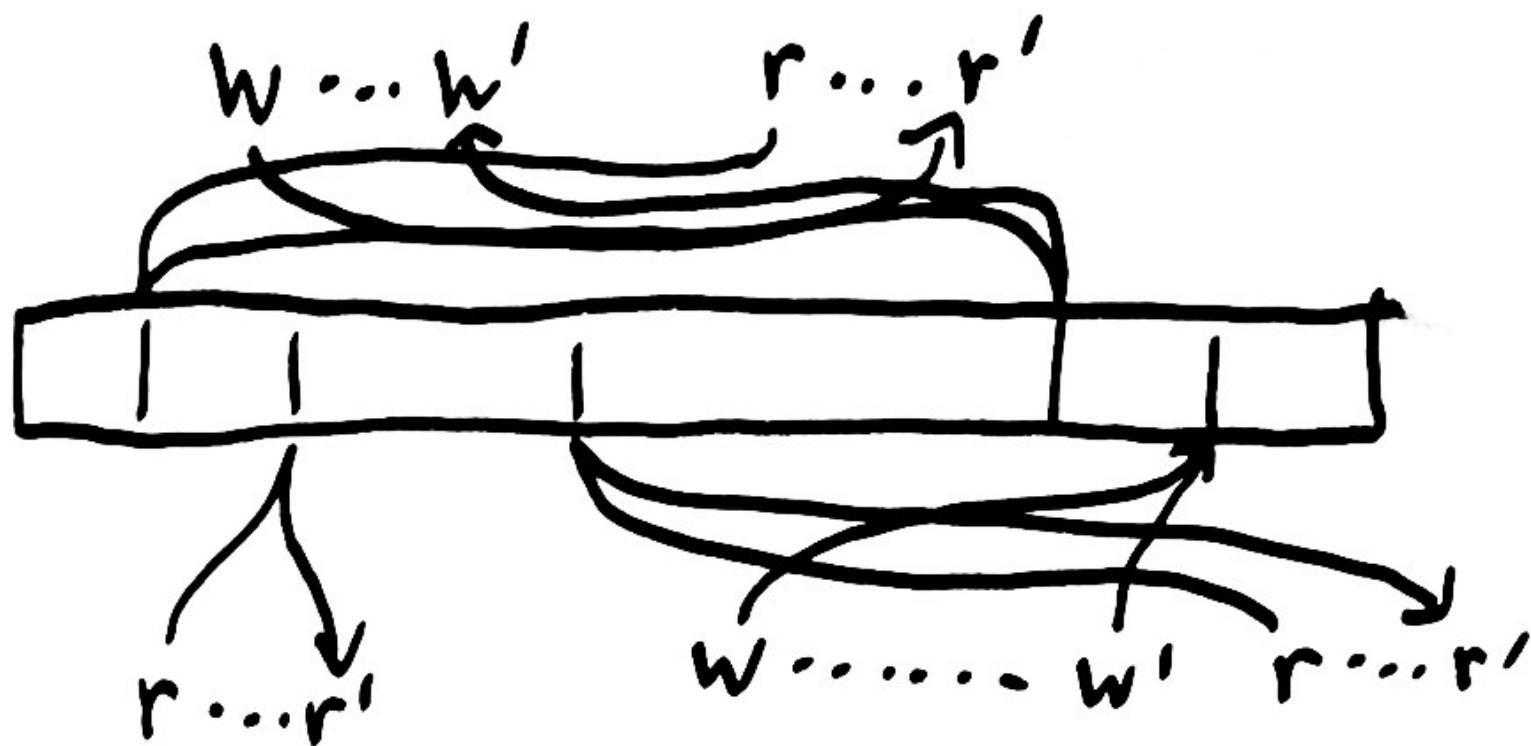


"lines can't cross"

# Serializability

"Everything happens in some order, but nobody agrees on what that order was!"

# Serializability



"ACID consistency/isolation"

"Stronger" consistency models

require more coordination,  
offer lower availability.

CAP rules out totally-available  
linearizable systems.

"Acronyms are easy  
to make shit up about."

-MRB

CAP:

C = Linearizability

A = All up nodes can  
satisfy all requests

P = Mandatory

Pick AP or CP.

Anyone who claims CA  
has no idea what they're  
doing.

Causally consistent systems

can be sticky-available.

Clients bound to specific  
servers.

Weaker models like reqd -  
uncommitted can be 100%  
available.

Strong Serializable

Linearizable

Sequential

Causal

WFR

PRAM

MR

RYW

MW

Serializable

RR

SI

CS

RC

MAV

P-CI

Infinitely many  
models!



Weaker consistency

models are less  
intuitive.

Weaker consistency

models are more

available in failure.

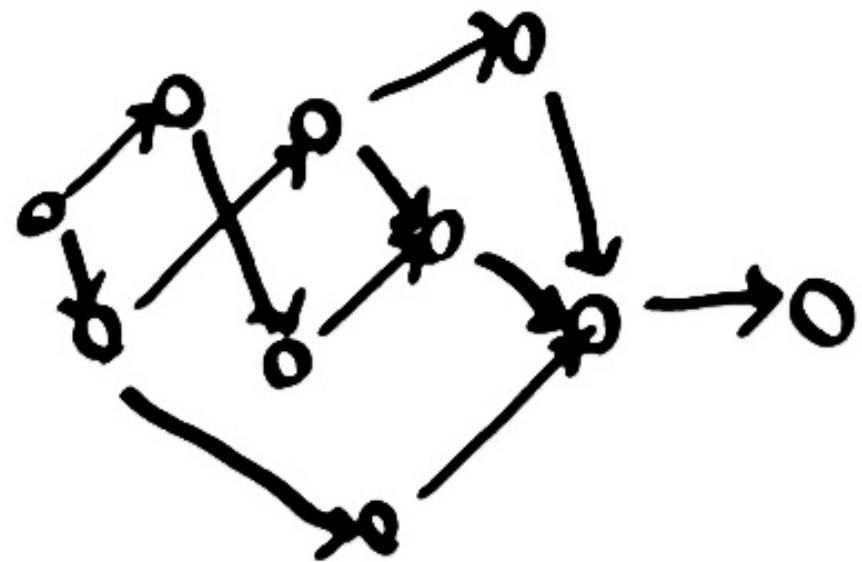
Weaker consistency

models are faster.

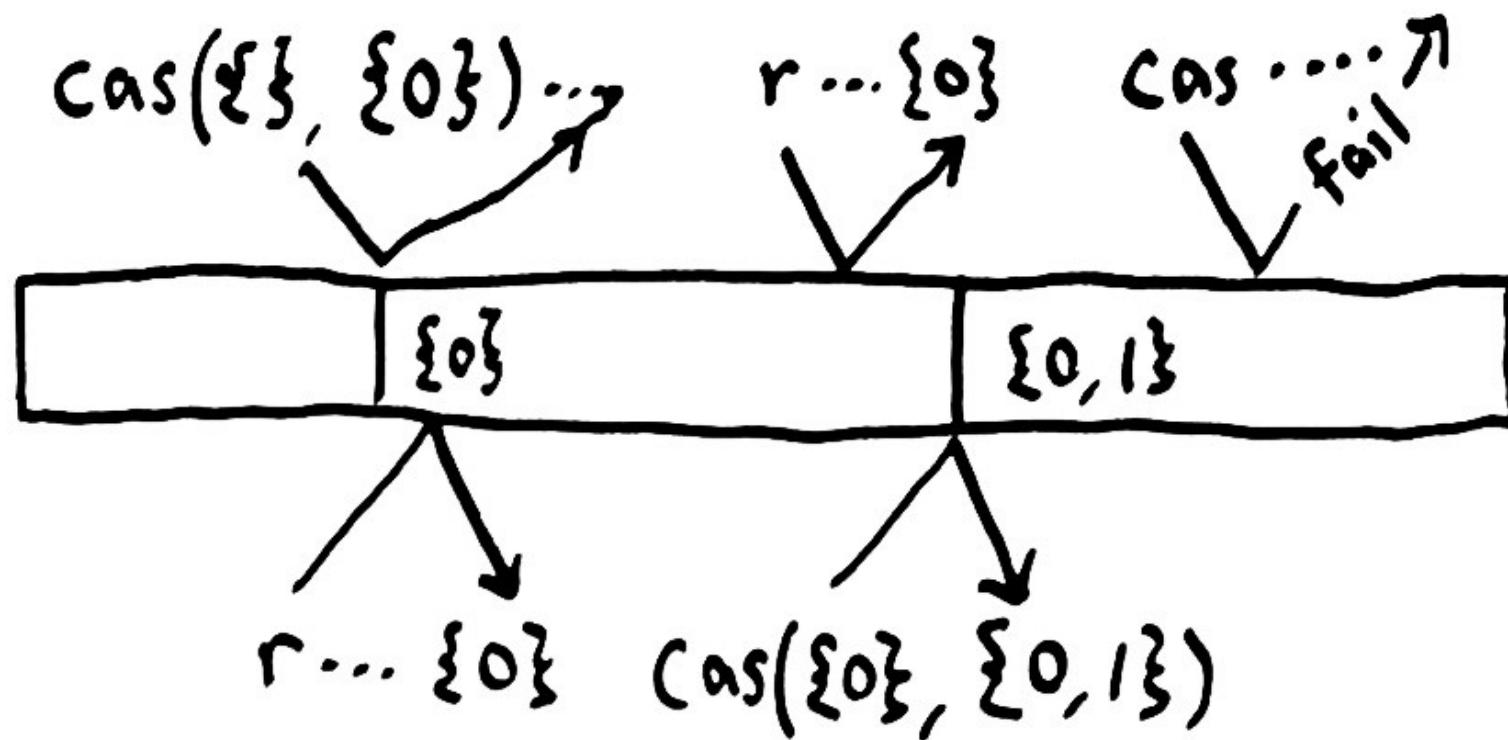
- CPU memory model

"Weak"  $\neq$  Unsafe

CRDTs



Jepsen failed to find problems  
in etcd, and other strongly  
consistent systems. Why?



CAS, CAS, CAS, CAS, ... read final  
 $+0 +1 +2 +3 \uparrow \{0, 1, 2, 3, \dots\}$   
?

set membership test implicitly  
Validates CAS correctness

What about read correctness?

Jepsen: DB automation, testing,  
failures, scheduling ops.



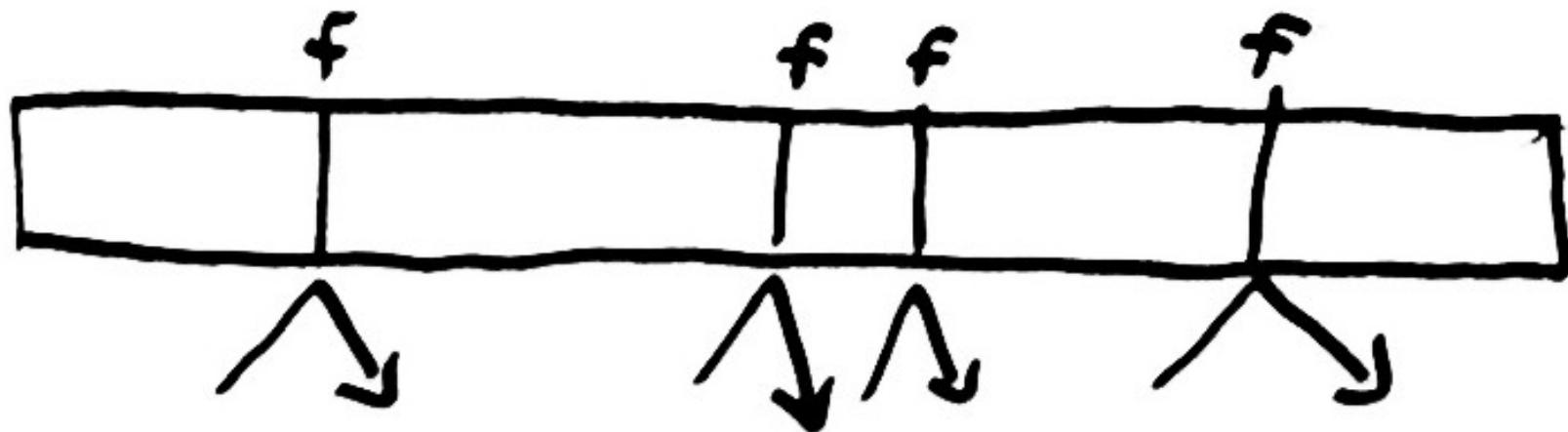
History of Ops



Knossos: Verifies Linearizability

# Model

$f(\text{state}, \text{op}) \rightarrow \text{state}'$



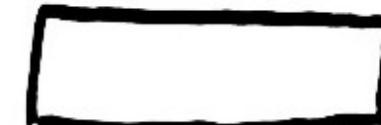
Model verifies that a  
single threaded history is OK.



But we don't know the  
singlethreaded history...

# Histories

Client 1



2



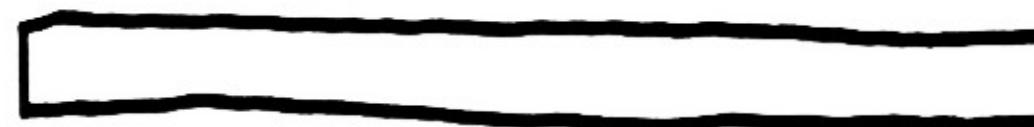
3



4



5



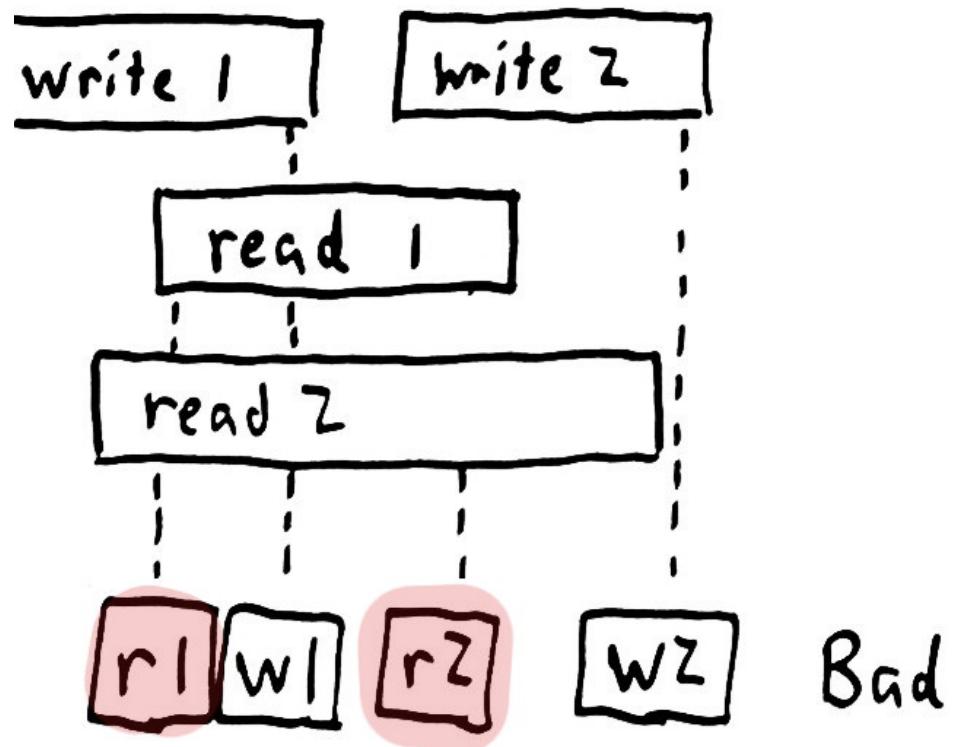
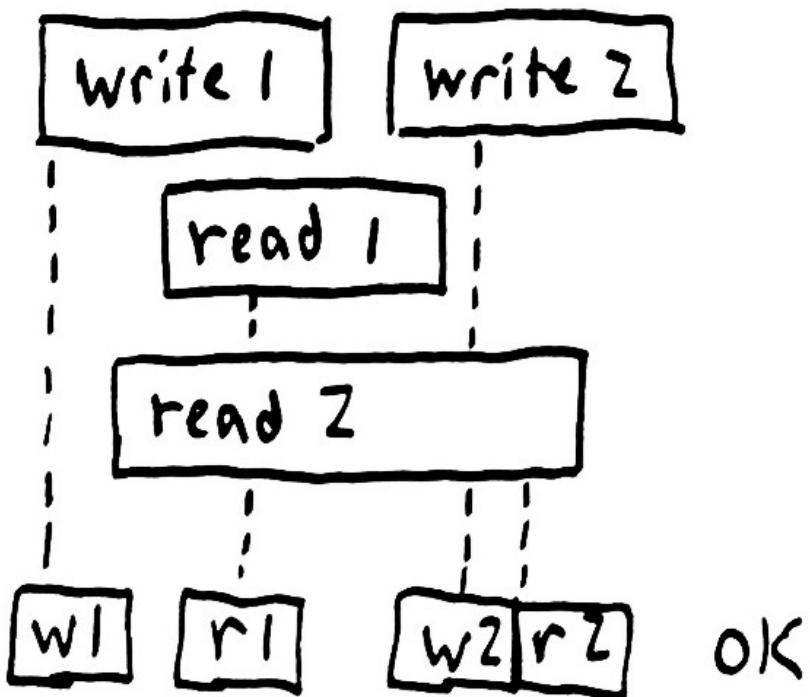
||||

Nemesis



$t$





$O(n!)$

(oh, )

# Techniques:

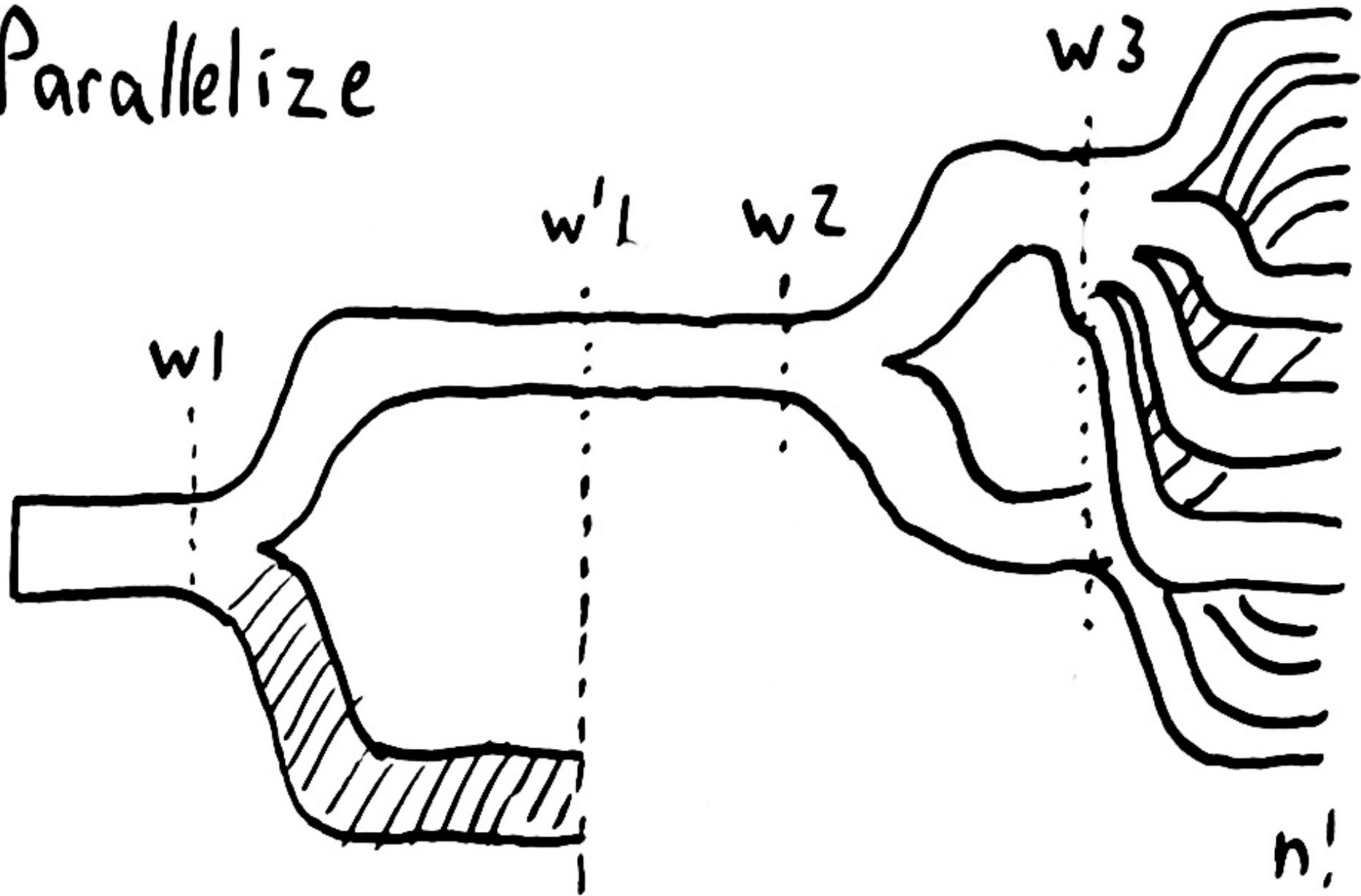
- Exploit degeneracy
- Parallelize
- Laziness + memoization
- Immutable, shared-structure representation of each world

# Degeneracy

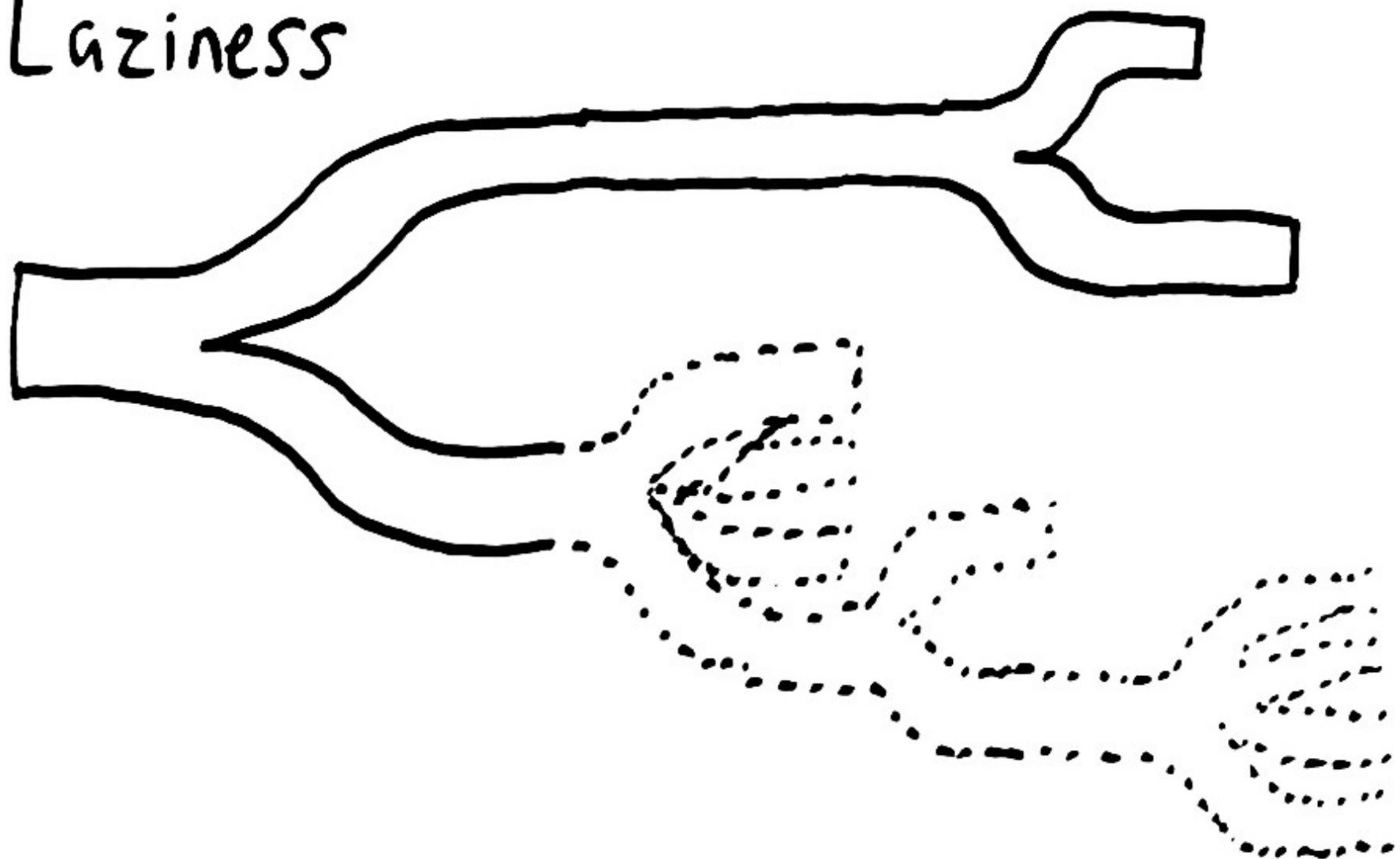
history	state	future
w'l, r'l	1	r'l, w'z, r'l, w'z
w'l	1	r'l, w'z, r'l, w'z
w'l, w'z	2	r'l, w'z, r'l, w'z
w'z,	2	w'z, r'l, w'z, r'l, w'z

The past is a fiction.

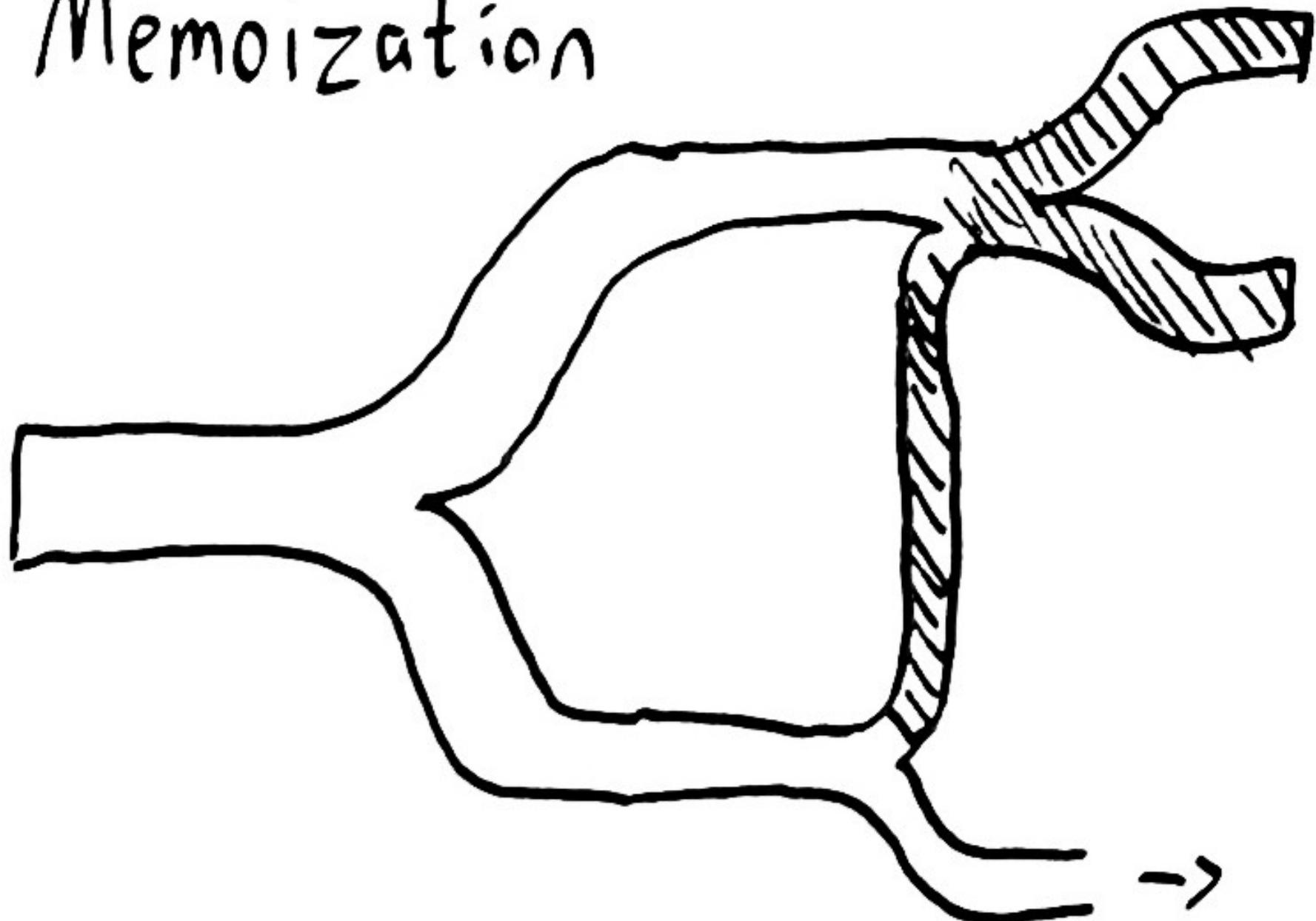
Parallelize



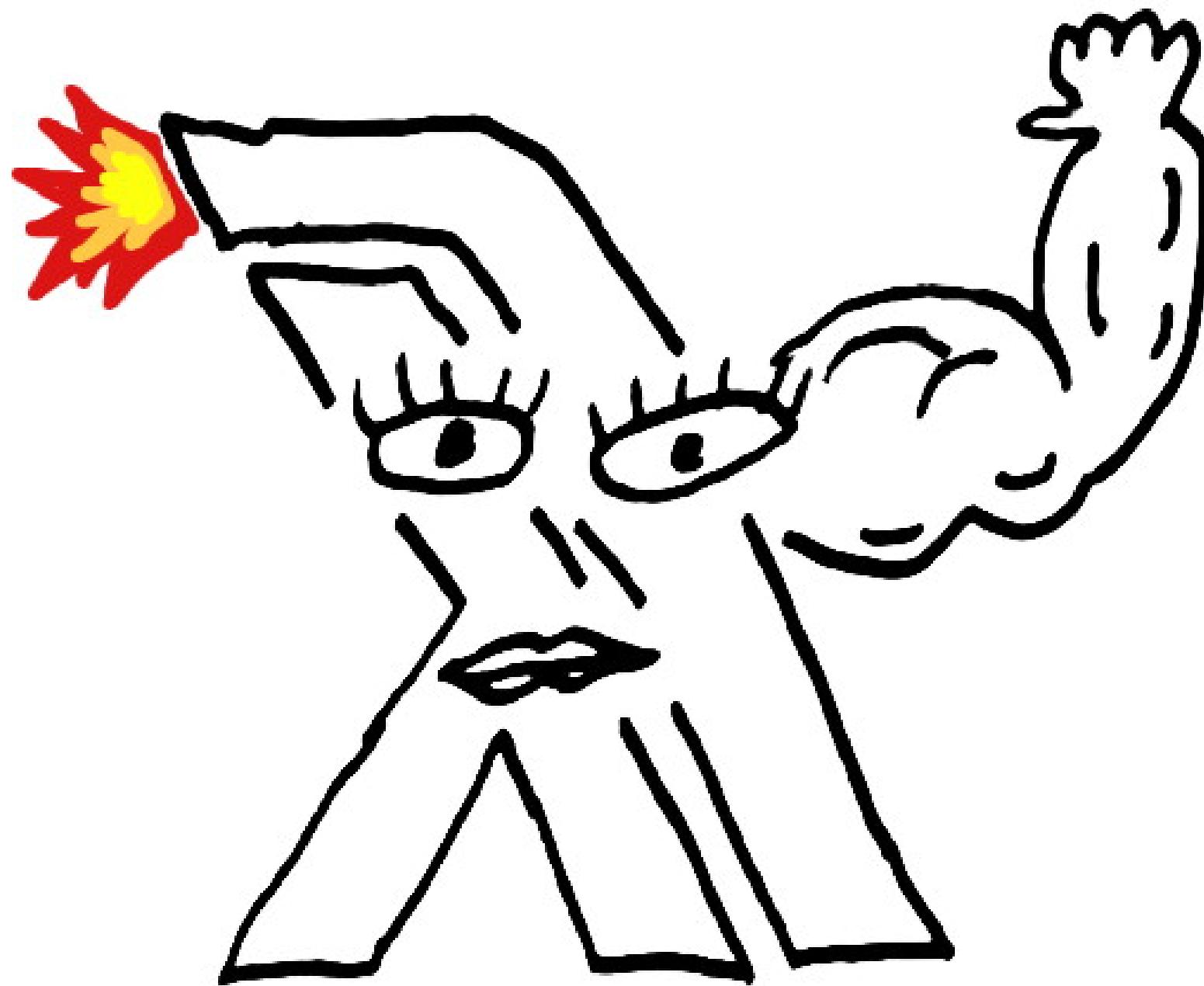
Laziness

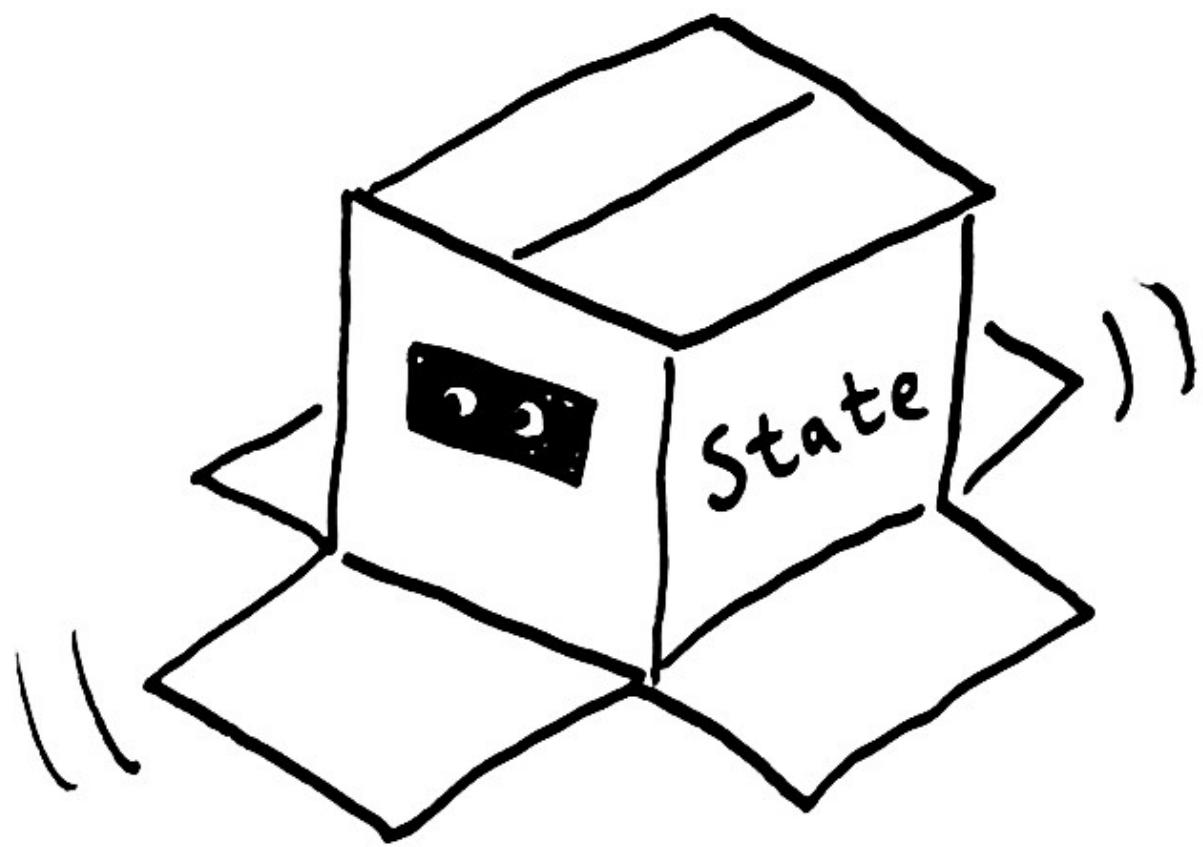


# Memoization



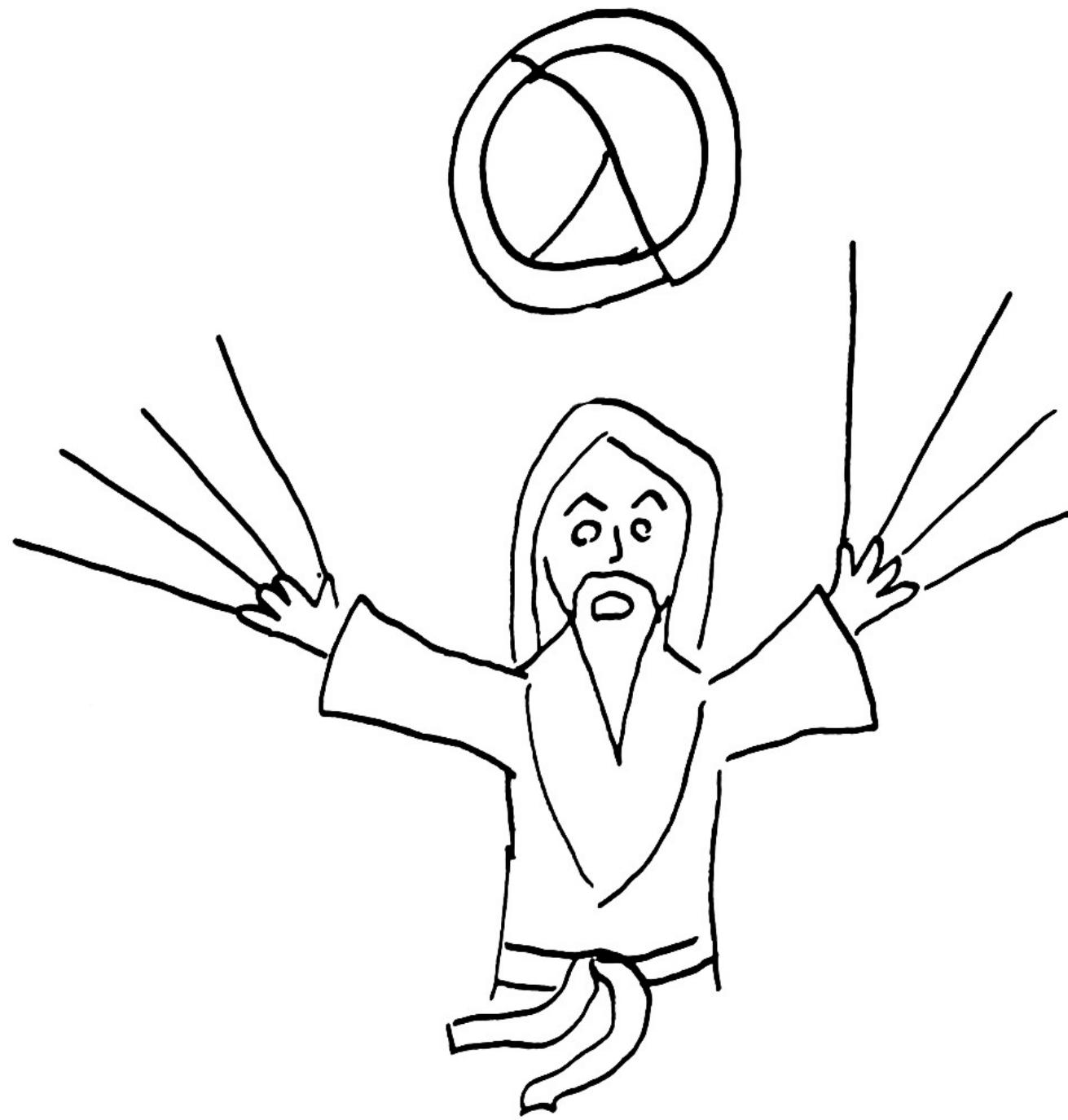
- All parallel-safe
- Execution as data
- Adapted from Margo Seltzer's talk on autoparallelizing x86





## Shared structure

Clojure's maps/vectors only store  
deltas; reduces memory overhead  
from world forks by orders of  
magnitude.



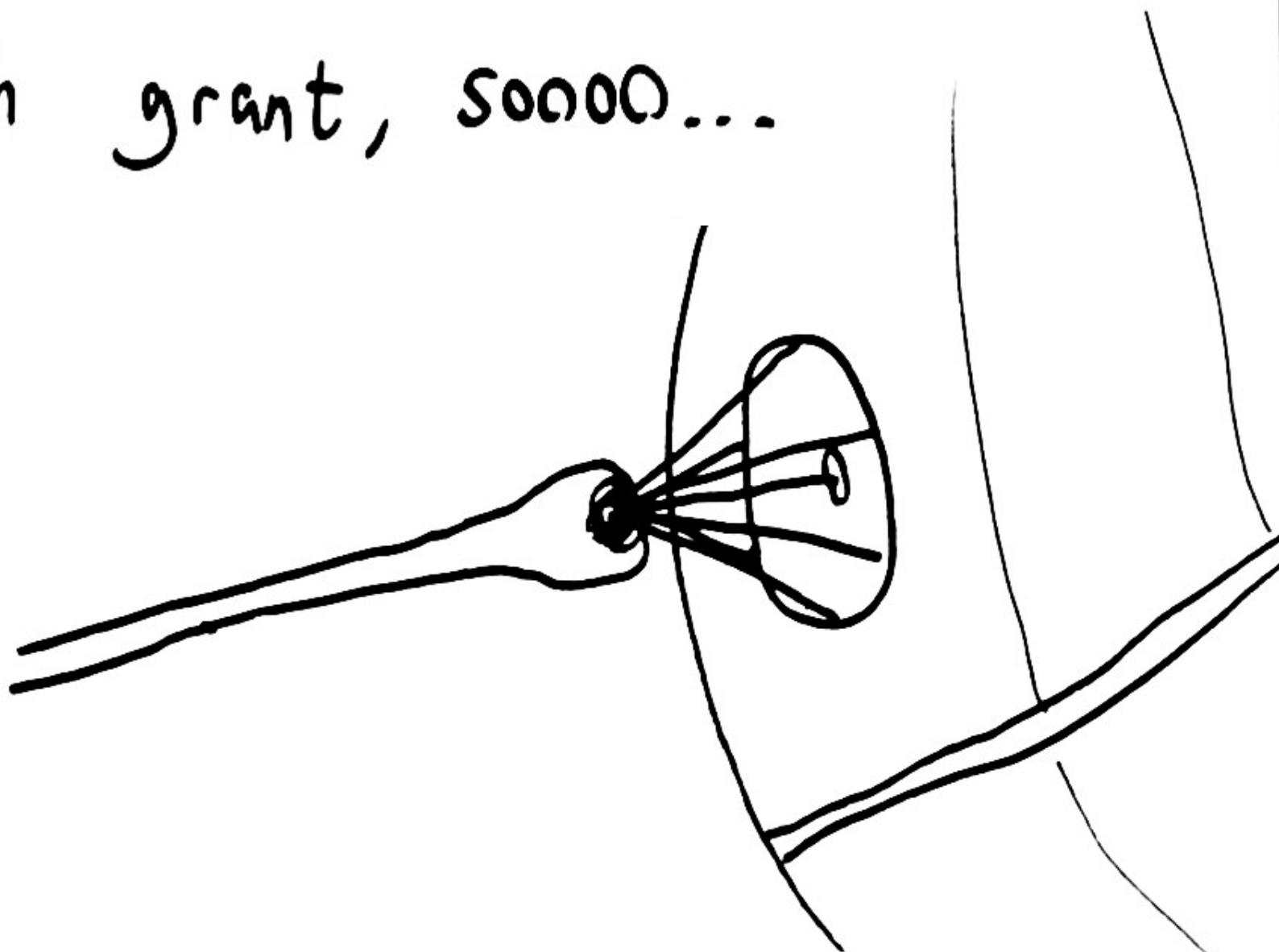


PUT A (pmap) ON IT!

Controlled use of JVM mutable  
types where perf > determinism

- Cliff Click's high-scale-lib
- Striped wrapper for j.u.c. PrioQueue
- j.u.c. atomics for CAS state
- Clojure atoms & futures too

Comcast gave me a  
research grant, sooo...



1	[				97.0%]	26	[					97.6%]
2	[				97.0%]	27	[					97.6%]
3	[				97.0%]	28	[					97.0%]
4	[				97.6%]	29	[					97.6%]
5	[				97.0%]	30	[					97.6%]
6	[				97.0%]	31	[					97.0%]
7	[				97.6%]	32	[					97.6%]
8	[				97.0%]	33	[					97.6%]
9	[				97.0%]	34	[					97.6%]
10	[				97.0%]	35	[					97.6%]
11	[				97.0%]	36	[					97.0%]
12	[				96.4%]	37	[					96.4%]
13	[				97.0%]	38	[					97.0%]
14	[				97.0%]	39	[					96.4%]
15	[				97.0%]	40	[					97.6%]
16	[				97.6%]	41	[					97.0%]
17	[				97.6%]	42	[					97.6%]
18	[				96.4%]	43	[					96.4%]
19	[				97.0%]	44	[					97.0%]
20	[				97.0%]	45	[					96.4%]
21	[				97.6%]	46	[					96.4%]
22	[				97.6%]	47	[					97.0%]
23	[				96.4%]	48	[					97.6%]
24	[				96.4%]	Mem	[					41769/129152MB]
25	[				96.4%]	Swp	[					0/5119MB]

PID	USER	PRI	NI	VIRT	RES	SHR	S	CPU%	MEM%	TIME+	Command
1702	aphyr	20	0	43.9G	1007M	27148	S	4511	0.8	33:11.91	java -classpath /home/aphyr/j
1868	aphyr	20	0	43.9G	1007M	27148	R	93.2	0.8	0:40.60	java -classpath /home/aphyr/j
2430	aphyr	20	0	43.9G	1007M	27148	R	94.4	0.8	0:38.89	java -classpath /home/aphyr/j
2432	aphyr	20	0	43.9G	1007M	27148	R	78.7	0.8	0:43.57	java -classpath /home/aphyr/j
2436	aphyr	20	0	43.9G	1007M	27148	R	93.8	0.8	0:43.06	java -classpath /home/aphyr/j
2803	aphyr	20	0	43.9G	1007M	27148	R	93.8	0.8	0:39.03	java -classpath /home/aphyr/j
2804	aphyr	20	0	43.9G	1007M	27148	R	90.2	0.8	0:38.59	java -classpath /home/aphyr/j
2805	aphyr	20	0	43.9G	1007M	27148	R	84.8	0.8	0:37.92	java -classpath /home/aphyr/j
2809	aphyr	20	0	43.9G	1007M	27148	R	80.6	0.8	0:38.95	java -classpath /home/aphyr/j

RabbitMQ

- Message Queue
- Knows about worker tasks
- Let's make a look out of it!

- Rabbit docs are very clear: partitions ruin all guarantees

=ERROR REPORT==== 10-Apr-2014::13:16:08 ===  
\*\* Node rabbit@n3 not responding \*\*  
\*\* Removing (timedout) connection \*\*

=INFO REPORT==== 10-Apr-2014::13:16:29 ===  
rabbit on node rabbit@n5 down

=ERROR REPORT==== 10-Apr-2014::13:16:45 ===  
Mnesia(rabbit@n1): \*\* ERROR \*\* mnesia\_event got  
{inconsistent\_database,  
running\_partitioned\_network, rabbit@n3}

---

« [Preventing Unbounded Buffers with RabbitMQ](#)

[Breaking things with RabbitMQ 3.3](#) »

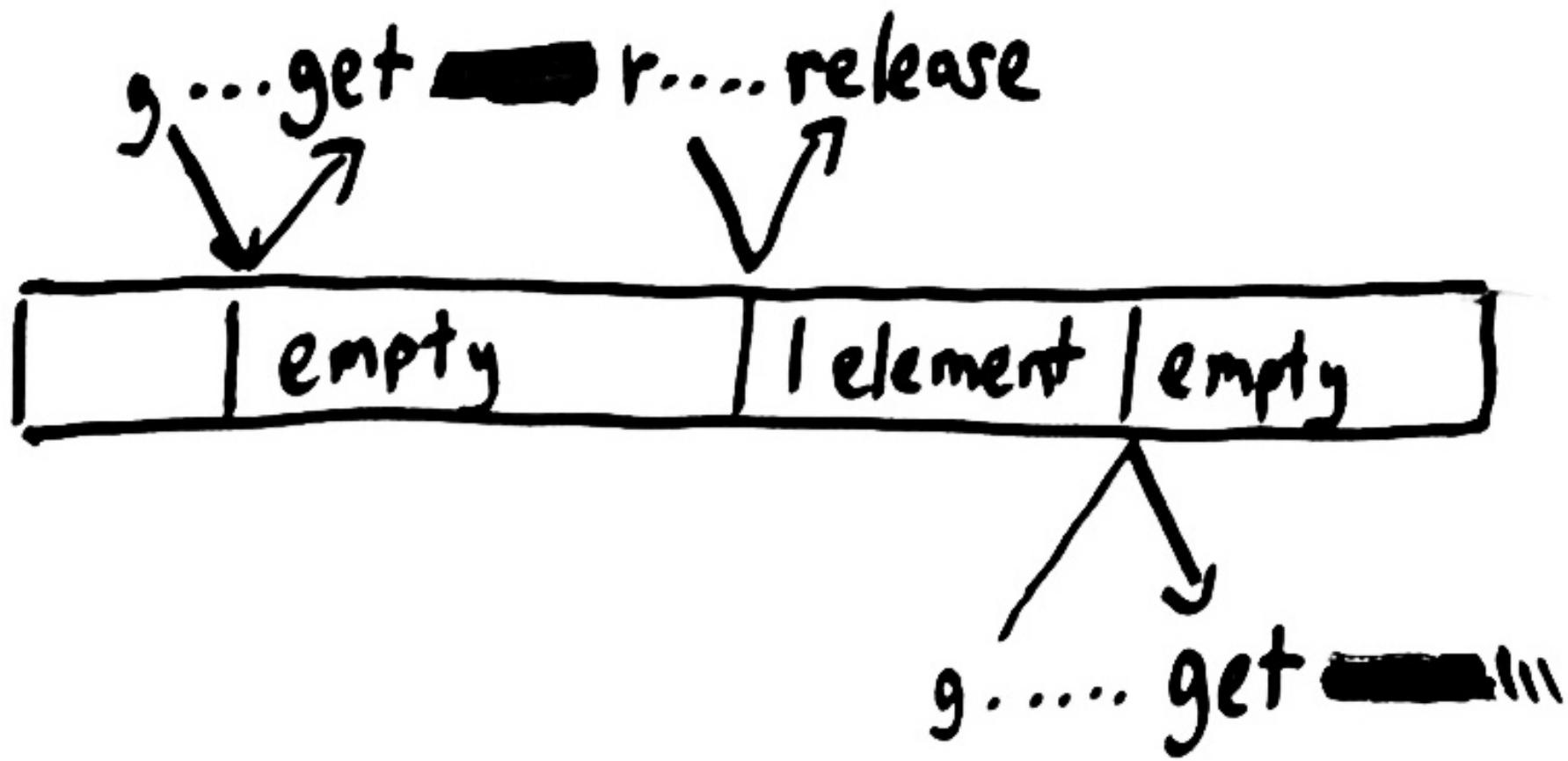
## Distributed Semaphores with RabbitMQ

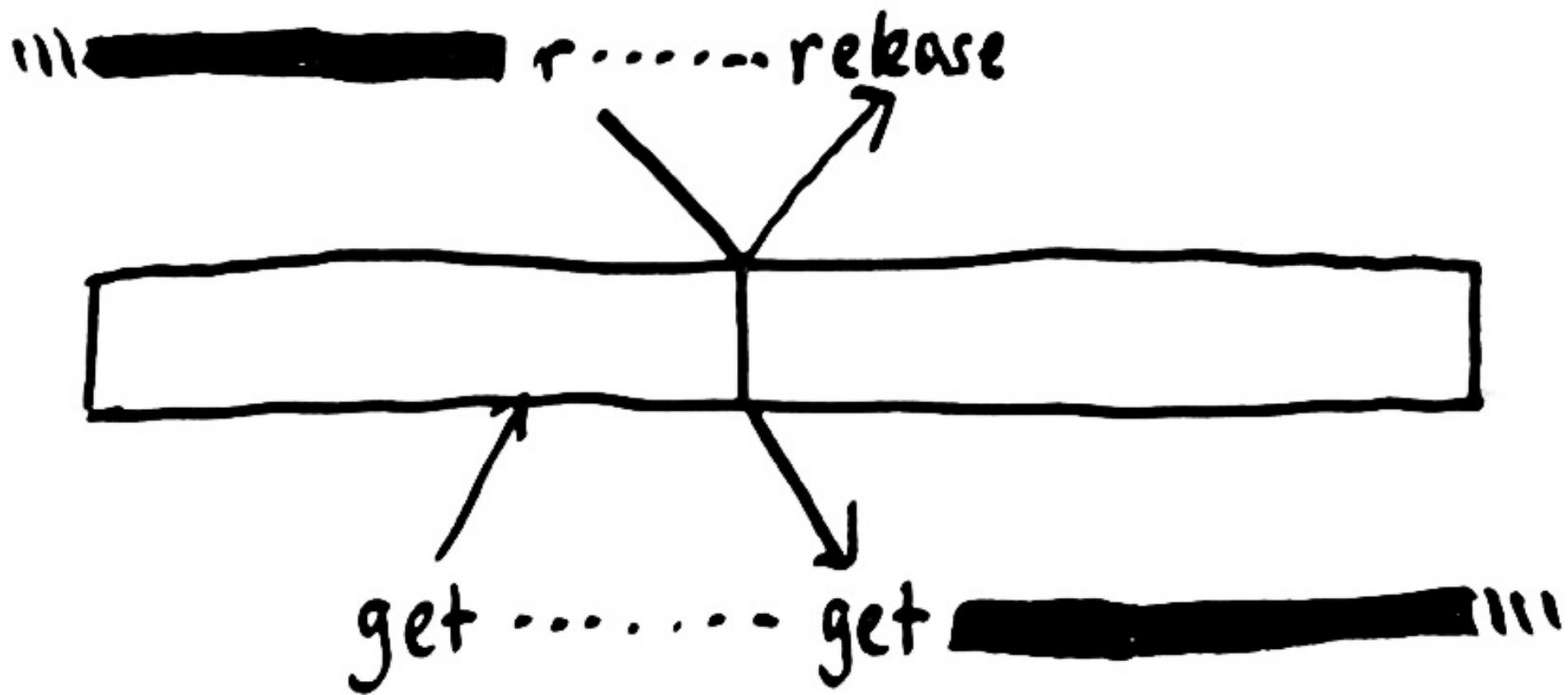
In this blog post we are going to address the problem of controlling the access to a particular resource in a distributed system. The technique for solving this problem is well known in computer science, it's called Semaphore and it was invented by Dijkstra in 1965 in his paper called "Cooperating Sequential Processes". We are going to see how to implement it using AMQP's building blocks, like consumers, producers and queues.

### The Need for Semaphores

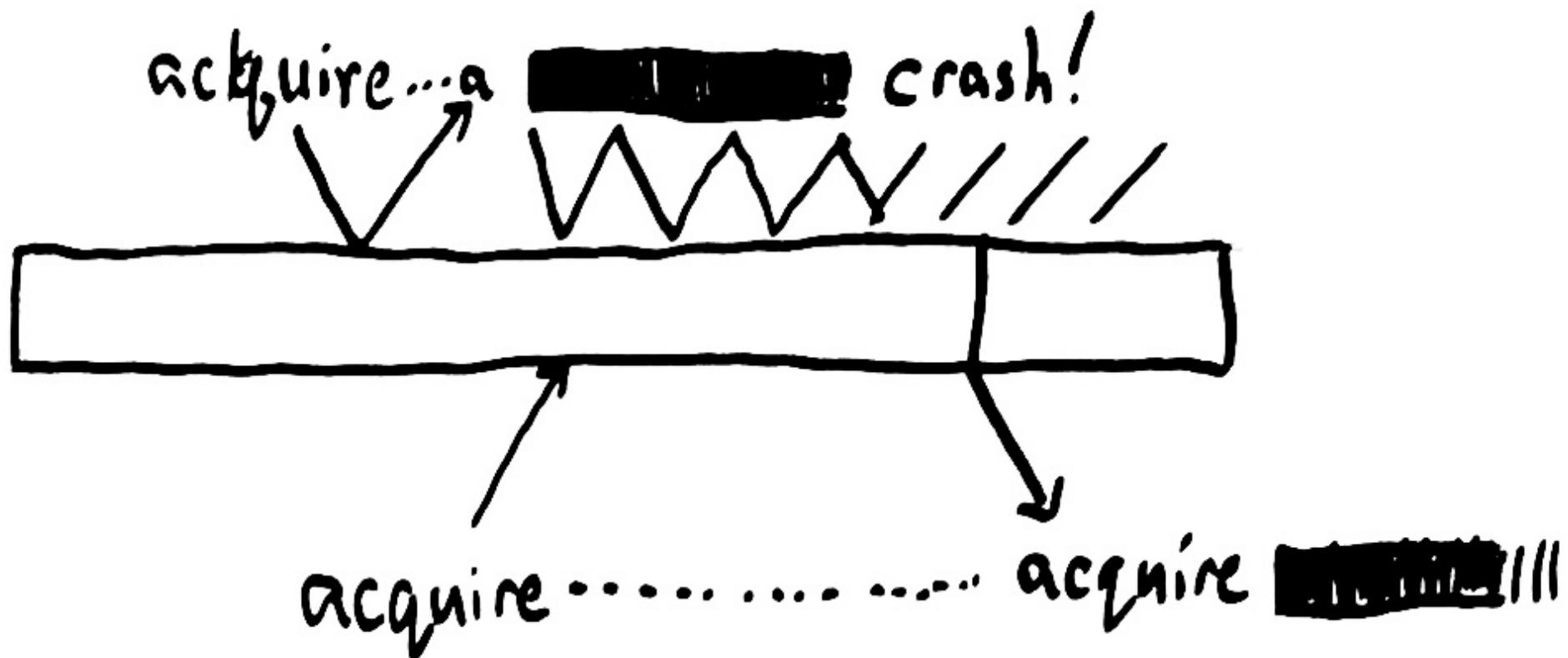
Before going into the actual solution, let's see when we might actually need something like this:

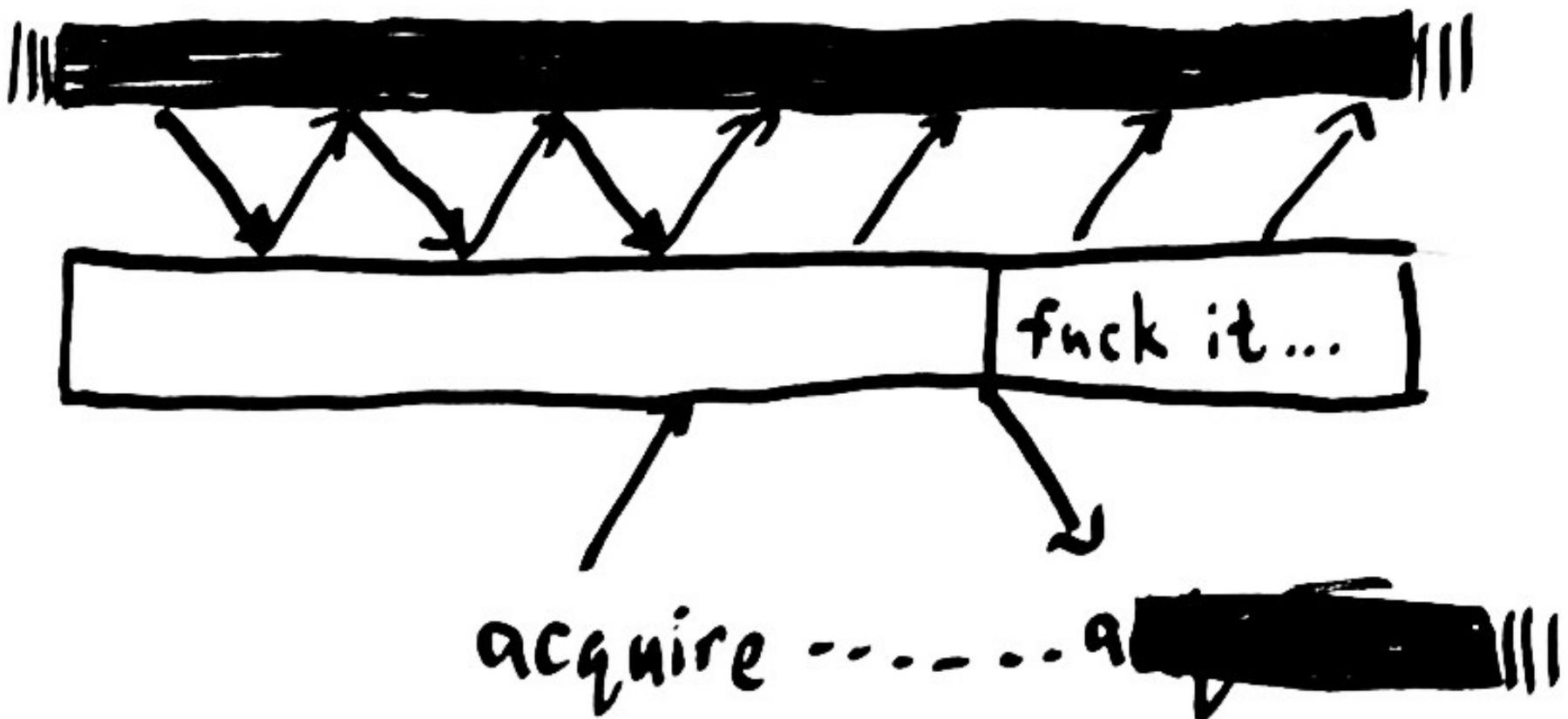
Let's say our application has many processes taking jobs from a queue and then inserting records to a database, we might need to limit how many of them do it concurrently.





But Rabbit is not  
a linearizable system!





Not linearizable. Linearizable prefix was:

...

:nemesis	:info	:stop	nil
:nemesis	:info	:stop	"fully connected"
:nemesis	:info	:start	nil
:nemesis	:info	:start	"partitioned into [(:n1 :n4) (:n3 :n5 :n2)]"
3	:invoke	:acquire	nil
2	:invoke	:acquire	1
2	:ok	:acquire	1
1	:invoke	:acquire	2

Followed by inconsistent operation:

1	:ok	:acquire	2
---	-----	----------	---

Last consistent worlds were: -----

World from fixed history:

```
1 :invoke :acquire 1
1 :invoke :release nil
2 :invoke :acquire 1
```

and current state #jepsen.model.Mutex{:locked true}

with pending operations:

```
3 :invoke :acquire nil
1 :invoke :acquire 2
0 :invoke :acquire nil
4 :invoke :acquire nil
```

---

Inconsistent state transitions:  
([{:locked true} "already held"])

Queues often sacrifice  
linearizability to provide  
fault tolerance.

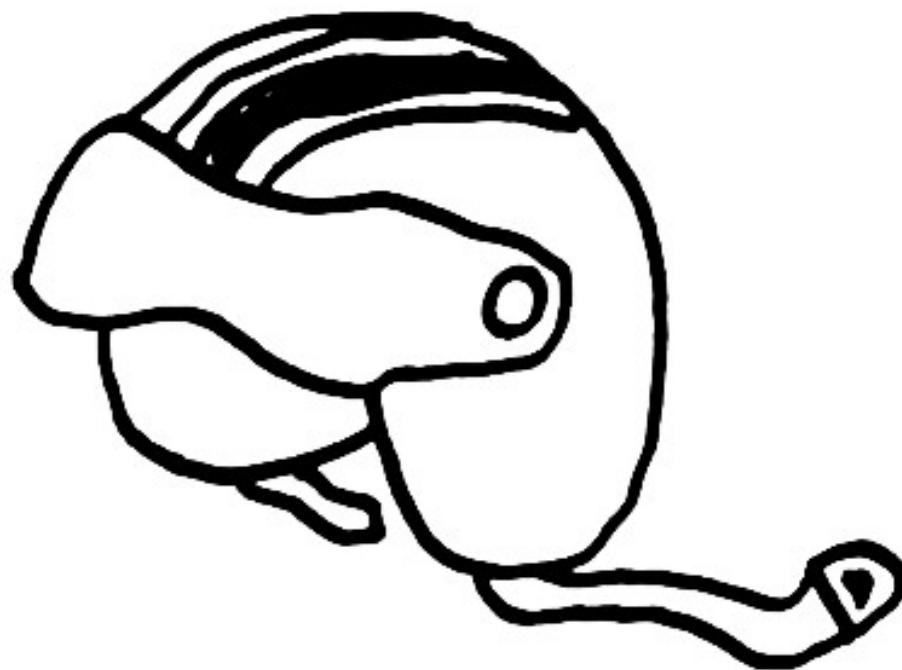
Perfectly fine queue.

Not a lock service.

e t c d

- Based on Raft
- HTTP API
- Configuration & Coordination
- Reads, writes, value &  
index-based CAS

# Race Conditions in Cluster join!



FAIL in (register-test) (etcd\_test.clj:45)  
expected: (:valid? (:results test))  
actual: false  
Not linearizable. Linearizable prefix was:  
...  
:nemesis :info :start "Cut off {:n5 #{:n4 :n1}, :n2  
#{:n4 :n1}, :n3 #{:n4 :n1}, :n1 #{:n3 :n2 :n5}, :n4  
#{:n3 :n2 :n5}}"  
2 :invoke :cas [1 4]  
4 :invoke :read 1  
4 :ok :read 1  
2 :ok :cas [1 4]

Followed by inconsistent operation:  
0 :invoke :read 1

World with fixed history:

```
...
4 :invoke :write 1
0 :invoke :read 1
3 :invoke :read 1
1 :invoke :read 1
4 :invoke :read 1
2 :invoke :cas [1 4]
```

led to state:

```
{:value 4}
```

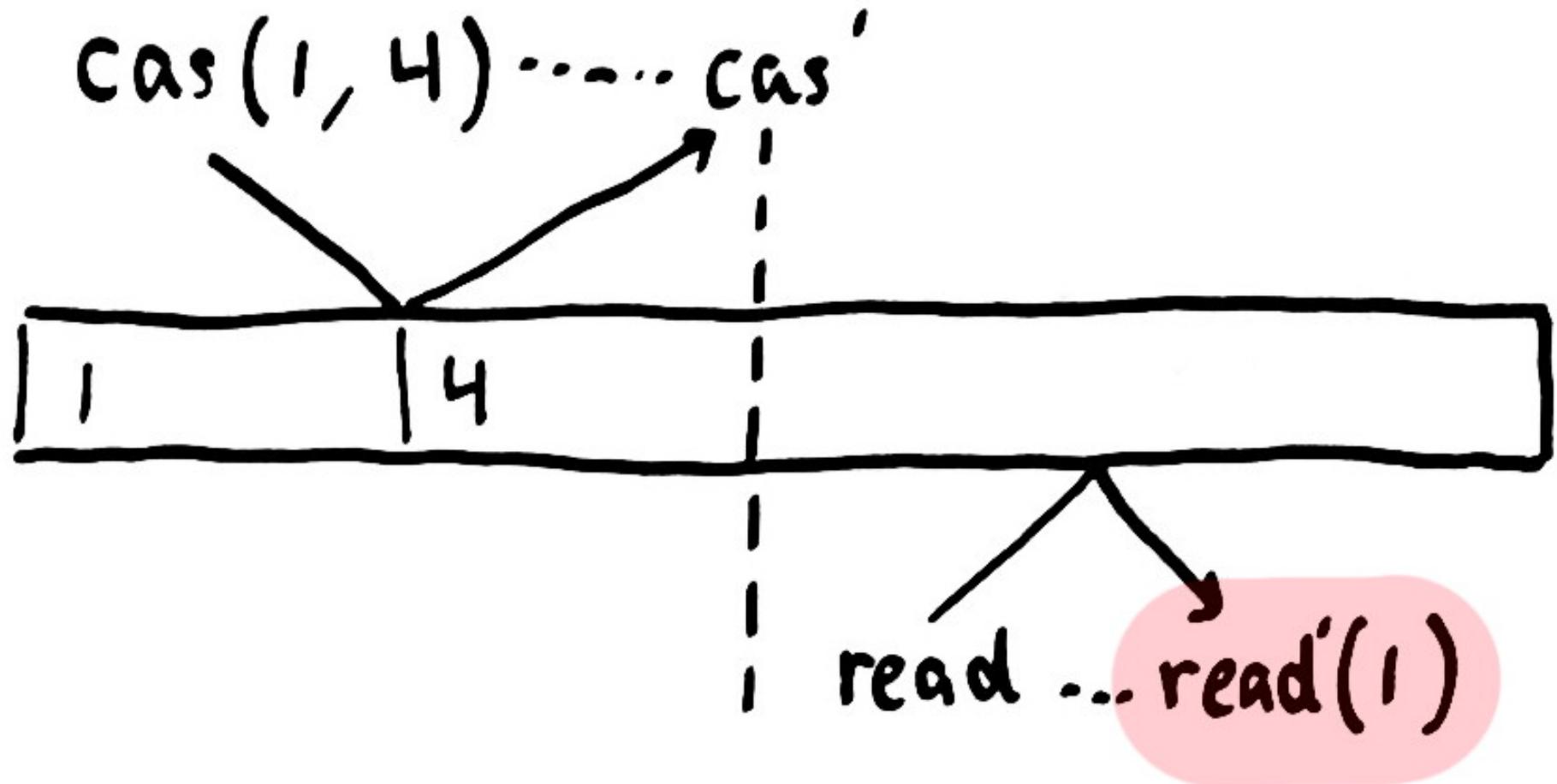
with pending operations:

(and 12928 more worlds, elided here

---

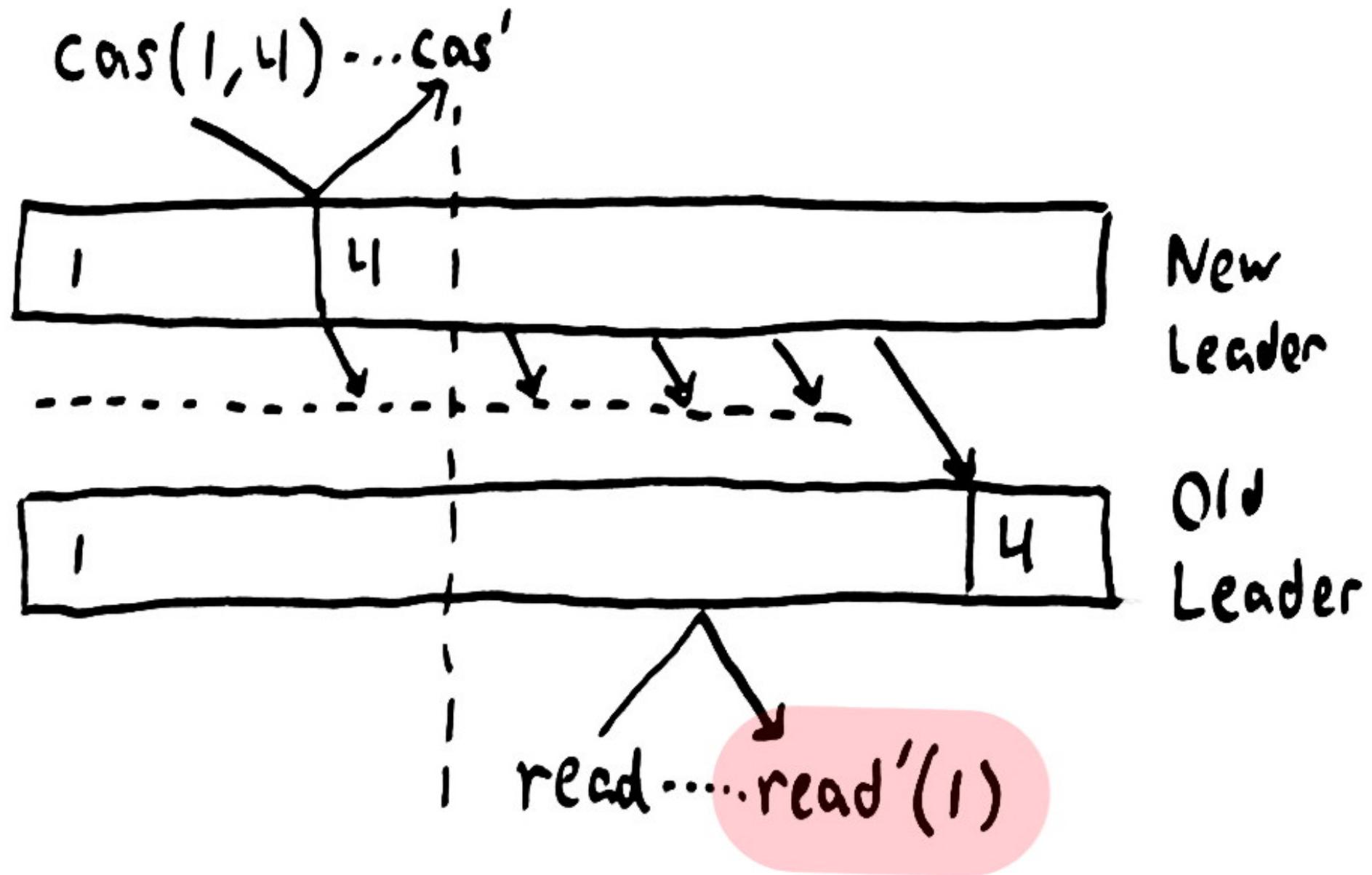
Inconsistent state transitions:

```
([{:value 4} "can't read 1 from register 4"])
```



Raft guarantees committed  
log entries are linearizable.

But etcd "consistent" reads  
don't go through the Raft  
log! They read current  
primary state & return!



Stale reads violate linearizability,  
sequential, & causal consistency.

Also RYW. Also MW. Also MR.

As part of our Consul testing, we ran it against Jepsen to determine if any consistency issues could be uncovered. In our testing, Consul gracefully recovered from partitions without introducing any consistency issues.

As part of our Consul testing, we ran it against Jepsen to determine if any consistency issues could be uncovered. In our testing, Consul gracefully recovered from partitions without introducing any consistency issues.

```
LeaderLeaseTimeout: time.Second,
```

```
LeaderLeaseTimeout: 300 * time.Millisecond,
```

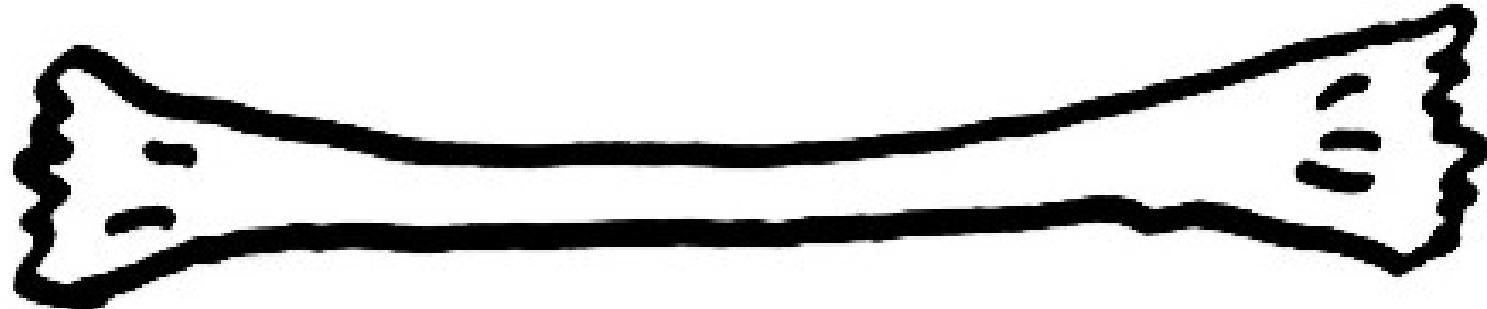
Suggestion: timeouts in a non-realtime system can't guarantee correctness, only improve odds.

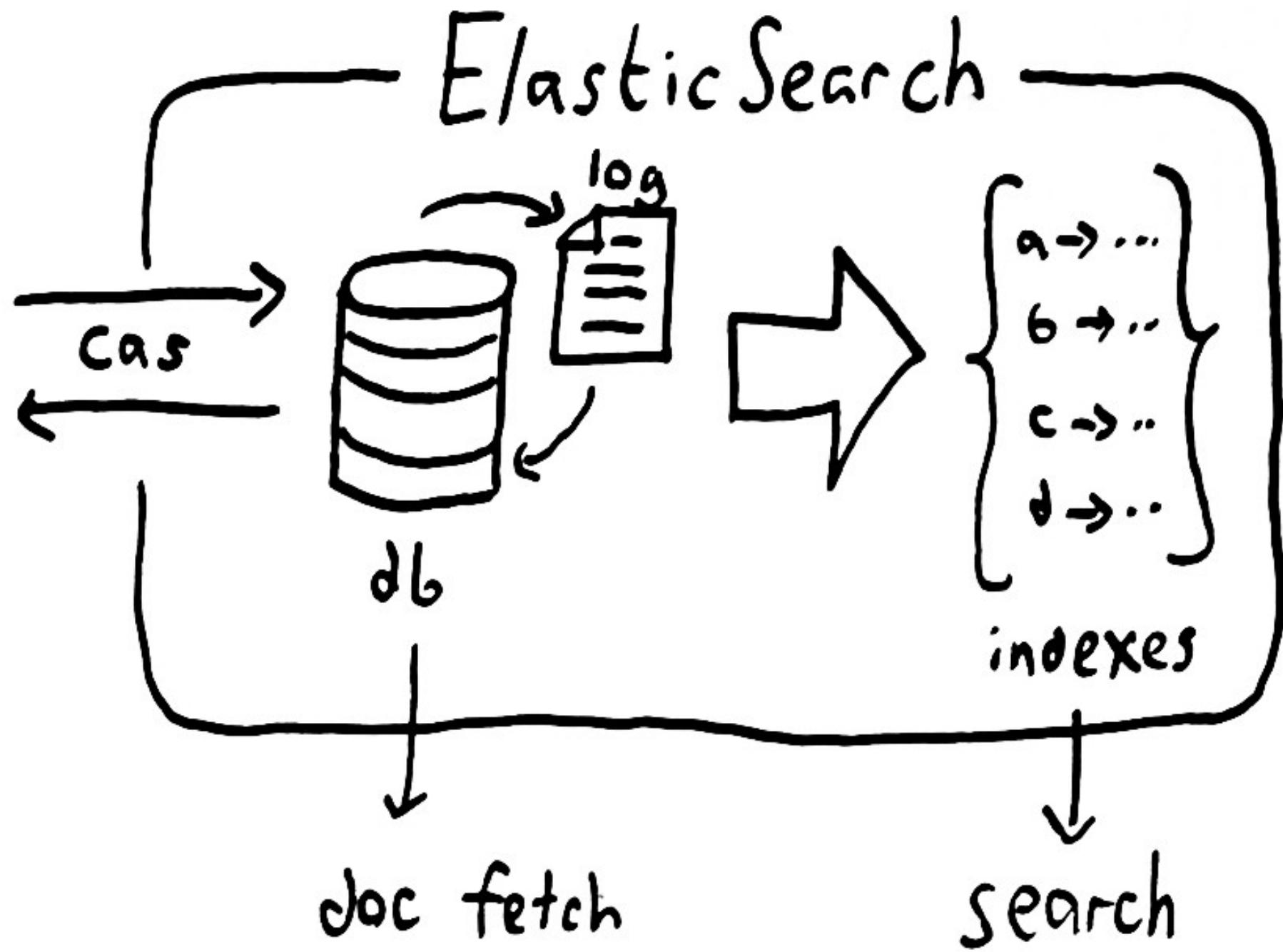
→ Elasticsearch GC partitions...

Both etcd & Consul are considering  
making reads actually consistent

- Latency cost!
- Many apps can tolerate stale reads.

ElasticSearch!





Elasticsearch offers  
atomic compare-and-set.

⇒ linearizability

[kimchy](#)

Jun 14, 2010;

[Reply](#) | [Threaded](#) | [More](#) ▾



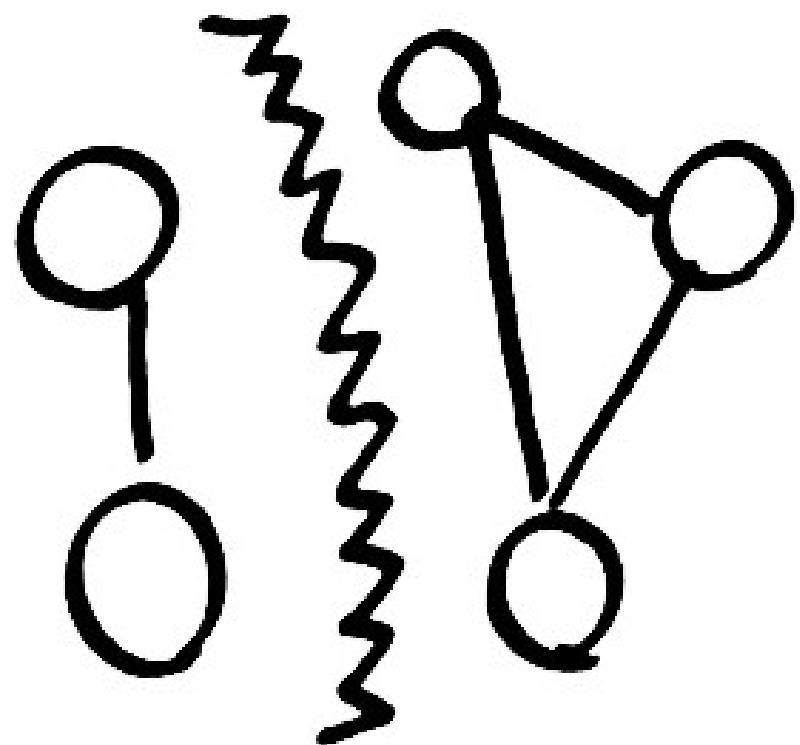
Administrator  
6265 posts

Hi,

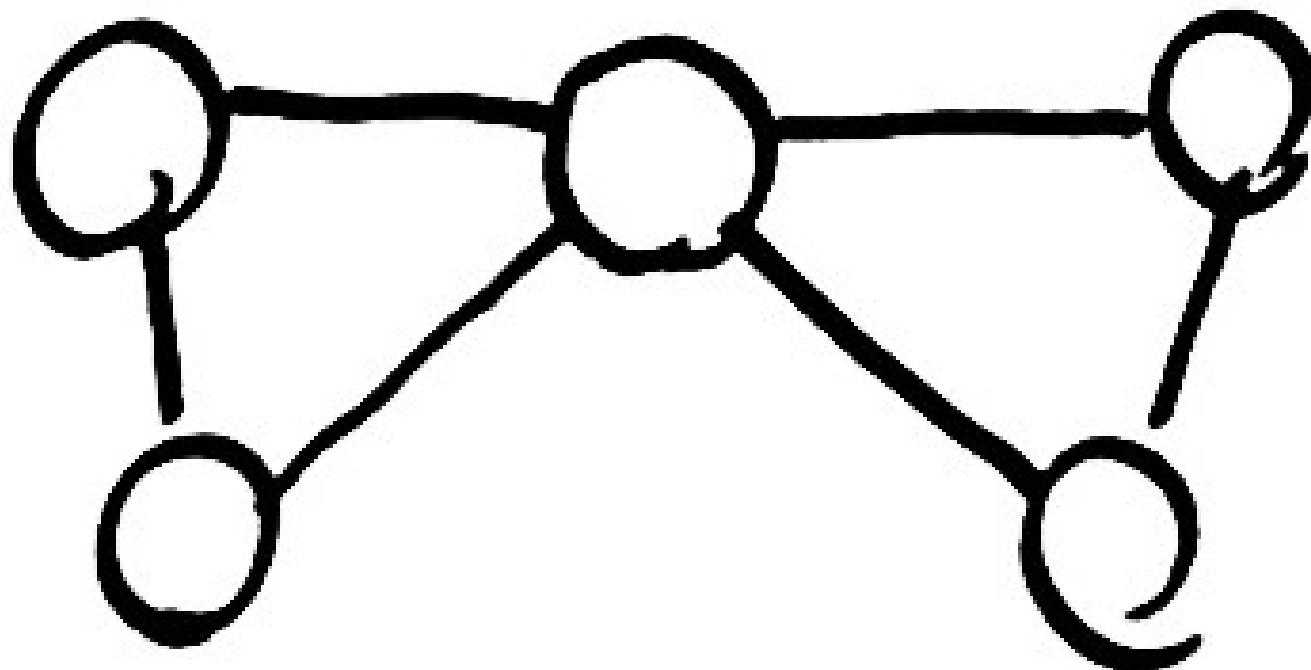
When it comes to CAP, in a very high level, elasticsearch gives up on partition tolerance. This is for several reasons:

1. I personally believe that \*within the same data center\*, network partitions very rarely happen, and when they do, its a small set (many times single) machine that gets "partitioned out of the network". When a single machine gets disconnected from the network, then that's not going to affect elasticsearch. When it comes to cross data centers, a solution that gives up on consistency can be built on top of elasticsearch, either by elasticsearch (in the future), or now, by using messaging in some form to replicate changes between two data centers.

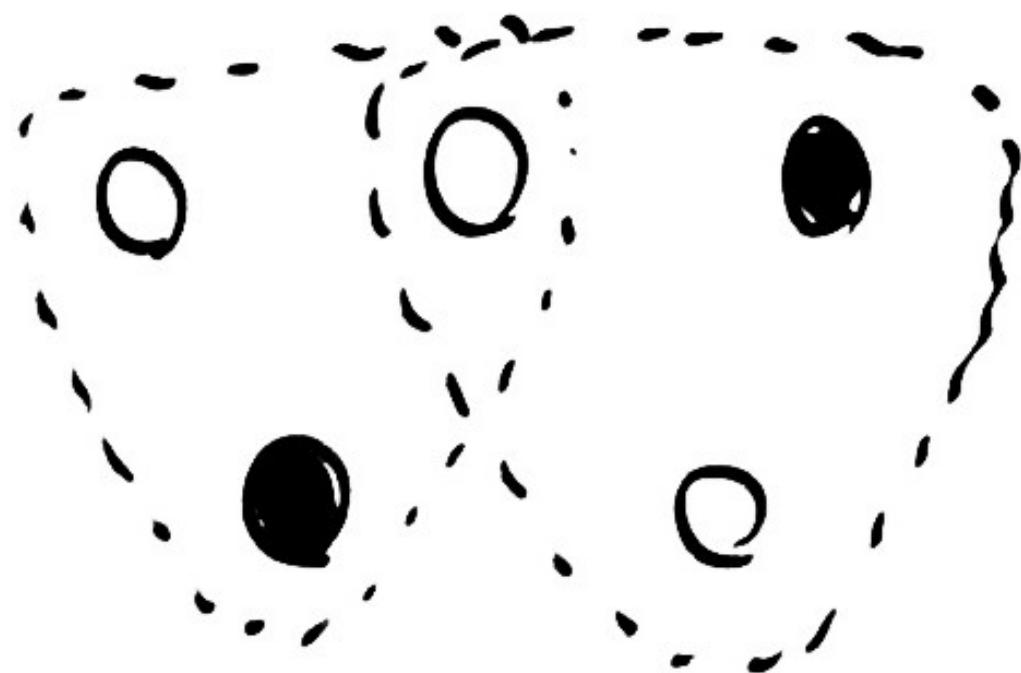
Assumes partitions are  
symmetric & complete



What about..



Every node thinks it  
has majority.



2 primaries

Short windows of write

loss during cluster transitions.

2 :ok :add 46  
1 :invoke :add 47  
4 :invoke :add 48  
4 :ok :add 48  
0 :invoke :add 49  
:nemesis :info :start "Cut off { :n5 #{:n2 :n1}, :n4 #{:n2 :n1},  
:n1 #{:n4 :n5}, :n2 #{:n4 :n5} }"  
3 :invoke :add 50  
2 :invoke :add 51  
4 :invoke :add 52  
1 :info :add :timed-out  
0 :info :add :timed-out  
3 :info :add :timed-out  
2 :info :add :timed-out  
4 :info :add :timed-out  
5 :invoke :add 53  
6 :invoke :add 54  
8 :invoke :add 55  
7 :invoke :add 56  
9 :invoke :add 57  
5 :info :add :timed-out

```
186 :invoke    :add    273
186 :fail      :add    273
173 :invoke    :add    274
173 :ok       :add    274
185 :info     :add    :timed-out
:nemesis     :info   :stop  nil
:nemesis     :info   :stop  "fully connected"
:nemesis     :info   :stop  nil
:nemesis     :info   :stop  "fully connected"
174 :invoke    :read   nil
174 :ok       :read   #{0 1 2 3 4 5 6 9 10 11 12 13 14 15 16 17 18 19
20 21 22 23 24 26 27 28 29 30 31 32 33 35 36 37 38 39 40 41 42 43
44 45 46 48 50 51 52 55 56 57 60 61 62 65 67 68 71 72 73 76 77 78
81 82 83 86 87 88 91 92 93 96 97 98 101 102 103 106 107 108 111
112 113 116 117 118 121 122 123 126 127 128 131 132 133 136 137
138 141 142 143 146 147 148 151 152 153 154 155 156 157 158 159
160 161 162 163 164 166 168 169 170 171 173 174 175 270 271 272
274}
```

```
{:valid? false,
:ok      #{0..6 9..24 26..33 35..46 48 50..52 55..57 60..62 65
           67..68 71..73 76..78 81..83 86..88 91..93 96..98
           101..103 106..108 111..113 116..118 121..123
           126..128 131..133 136..138 141..143 146..148
           151..164 166 168..171 173..175 270..272 274},
:recovered #{50..52 55..57 60..62 65 67..68 71..73 76..78 81..83
            86..88 91..93 96..98 101..103 106..108 111..113
            116..118 121..123 126..128 131..133 136..138
            141..143 146..147},
:lost       #{176..178 250..251},
:unexpected #{},
:ok-fraC    26/55,
:recovered-fraC 59/275,
:lost-fraC   1/55,
:unexpected-fraC 0}
```

```
if (result.hasConflicts()) {  
    // TODO: What should we do???
```

But wait –  
There's More!

Insert a document in

the next five minutes

and we'll lose

~~~~~  
NEW!  
~~~~~

HALF YOUR DATA!

```
"RemoteTransportException[[Cadaver]
[inet[/192.168.122.12:9300]][index]]; nested:
RemoteTransportException[[Death Adder]
[inet[/192.168.122.11:9300]][index]]; nested:
DocumentAlreadyExistsException[[jepsen-index][1]
[number][EpVU56YERB0fRqyVc-_hAg]: document already
exists];
```

The response is similar to what we saw before, except that the `_id` field has been generated for us:

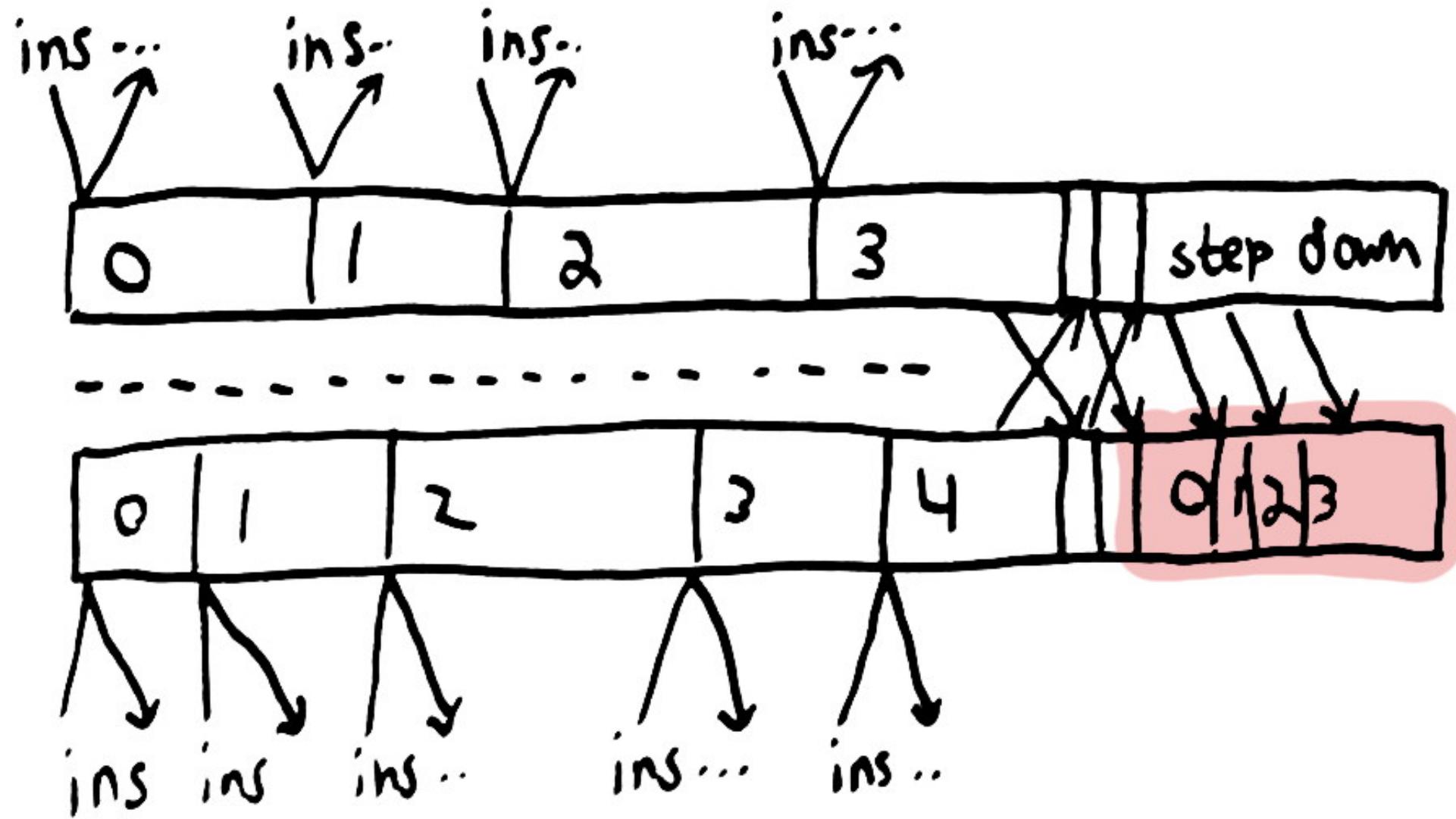
```
{
  "_index": "website",
  "_type": "blog",
  "_id": "wM00SFhDQXGZAWDf0-drSA",
  "_version": 1,
  "created": true
}
```

Auto-generated IDs are 22 character long, URL-safe, Base64-encoded string *universally unique identifiers, or UUIDs*.

```
FAIL in (register-test) (elasticsearch_test.clj:49)
expected: (:valid? (:results test))
actual: false
{:valid? false,
 :ok ...,
 :recovered #{420 438 457 475},
 :lost    #{418..419 421..427 429..437 439..444 446..456 458..463 465..474
476..481 483..493 495..500 502..503 505..506 509..510 514 516 520..521 524 527 529
531 534 536 539 542 544 547..548 552 554 557..558 562 564 567..568 572 574 577 579
582..583 587..588 592..593 597 599 602 604 607 609 612 614 617 619 622 624 626 629
631 634 636 638 642 644 647 649 652..653 657 659 662 664 667 669 671 674 676 679 681
684 686 689 691 694 696 698 701 704 706 709 711 715..716 720..721 725 727 730 732
735 737 741..742 746..747 749 752 754 757..758 762..763 767 769 772 775 777 779
782..783}
```

...

```
1706..1707 1710 1712 1716..1717 1721..1722 1726..1727 1731..1732 1736..1737
1741..1742 1746..1747 1751..1752 1756 1761 1766 1771..1772 1776..1777 1781..1782
1786..1787 1791..1792 1796..1797 1801 1806..1807 1811..1812 1816..1817 1821
1826..1827 1831..1832 1836..1837 1841..1842 1846..1847 1851..1852 1856..1857
1861..1862 1866..1867 1871..1872 1877..1878 1882..1883 1887..1888 1892..1893
1897..1898 1902..1903 1907 1912..1913 1917..1918 1922..1923 1926 1931..1932
1936..1937 1941 1946..1947 1951 1953 1957},
:unexpected      #{},
:ok-frac        1103/1961
:recovered-frac 4/1961,
:lost-frac      645/1961,
:unexpected-frac 0})
```



Multiple production user  
reports - caused by  
network failure or GC.



**brusic** commented 9 months ago

Have I mentioned how serious this problem is? Our production cluster has pretty much gone away. It is like a game of whackomole trying to kill instances that think they are the master.



**fasher** commented 9 months ago

I agree, this issue is critical please give this priority.

We gone to a single master node in our production to avoid split brains we had in the past that corrupted our index.



**anandnalya** referenced this issue 7 months ago

**NodeDoesNotExistOnMasterException Handling #3663**

Open



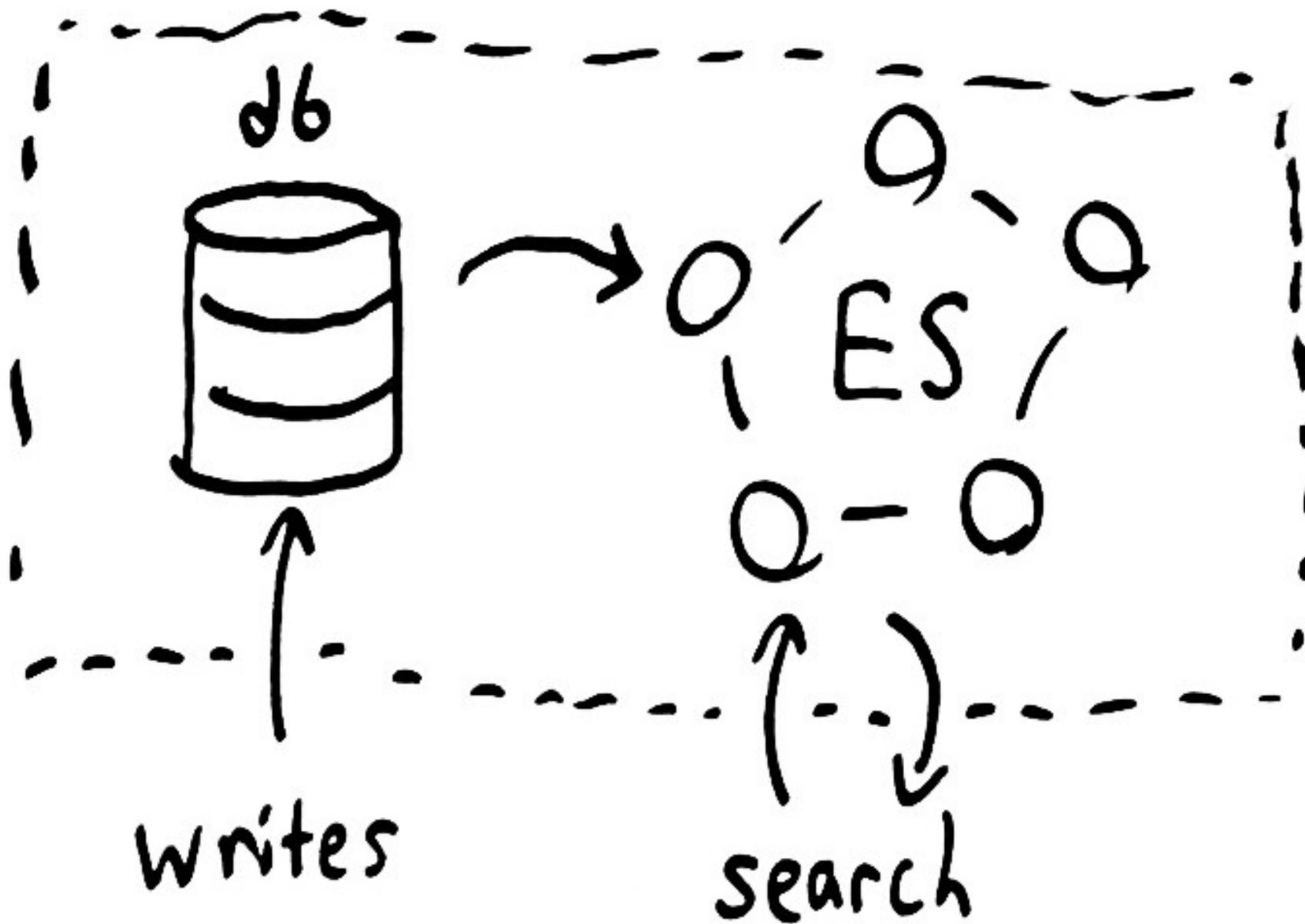
**bitsofinfo** commented 3 months ago

Any update on a timeline for a fix for this?

Zendisco is terrible.

Use zk instead!

Always treat ES as a  
derivative datastore. Keep  
Primary data somewhere  
safe, & periodically resync.



In summary...

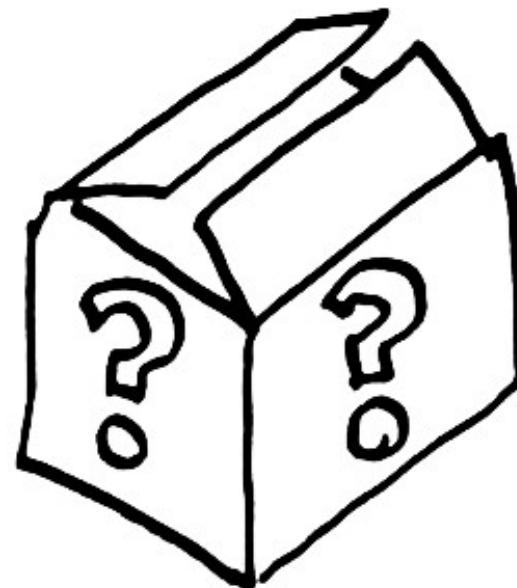
Jepsen tells you about

1 Particular

SYSTEM

Is YOUR system

safe? I dunno!



Consider...

- Runtime pauses
- Network slowness
- One-way partitions

Use iptables to simulate  
failures in prod, and record  
results!

Generate random ops,  
apply to system, compare  
recorded history to model.

Reason carefully about state  
transitions with critical impact

How much safety do  
you really need?

Losing data or stale  
reads can be preferable  
to high latencies

- Most systems have gone untested under failure
- You'd be surprised what well-designed systems will do during a network partition
- Test your systems.

Thanks!

Peter Bailis

Aaron France

Bob Poekert

Alvaro Videla

Kelly Sommers

Jacques Fuentes

Armon Dadgar

Xiang Li

Joseph Blomstedt

Tyler Schwett

Comcast + Factual + et al