

# CS 2

## Introduction to Programming Methods



## Last Time

### Dynamic programming

- how to solve optimization problems fast

#### Other Dynamic Prog Examples

Find min # of coins to make an amount

- to find the solution for 134
  - solve for all of 0, 20, 25, ..., 130
  - choose best among solutions for all solutions for (1-14)

#### Knapsack problem

- various item types of various values & weights
- one bag with limited total weight

#### Gene Sequence Alignment

#### Shortest Path

- more on this later

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## Networking

### Communication between computers

- web, email, streaming, etc



Telephony



Video



Music



Social networks



Location



Mail



Cloud computing

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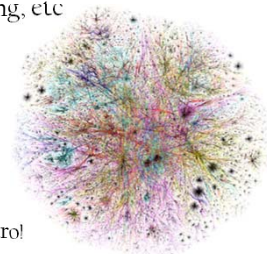
## Networking

### Communication between computers

- web, email, streaming, etc

### Internet

- complex topology
  - computers connected
  - and routers
    - traffic directing



As usual, today is just an intro!

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## Addressing

### Each machine has

- a globally unique address
  - IP address; 32 bits for now (www.caltech.edu: 50.18.115.211)
- and a domain name
  - easier for humans...
  - mapping from DN to IP may change dynamically
    - tree of DNS servers
      - keep track of domains used recently

### Addresses assigned by ICANN

- Marina del Rey

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## Circuit Switching

### Used in (old) telephone network

- when a call is made, switches are activated...
- establishing a line between two phones
  - quite rigid!
    - line busy even if no one is talking over it
    - if all lines are busy, you are stuck
  - but efficient
    - if you get a line, data can be sent reliably, with no overhead
- not quite appropriate for the internet...



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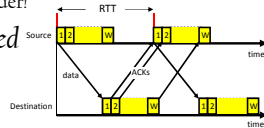
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## Packet switching

### Basic Principles

- data split into small (numbered) packets
- each packet is sent to receiver
  - so each can actually end up using a different path
- data reassembled at the receiver
  - can be received out of order!

Can be *connection-oriented* or *connectionless*



## Connection-oriented Switching

### “Virtual” Circuit Switching

- initial setup establishes route to destination
  - hops between network nodes; more on this later
- each node aware of “circuit”
  - stores connectionID to know where to forward packets
- then packets just include their connectionID
  - small header (minor overhead)
- “dedicated” line, in-order transmission
  - example: X.25, Frame-relay; also, TCP

## Connectionless Switching

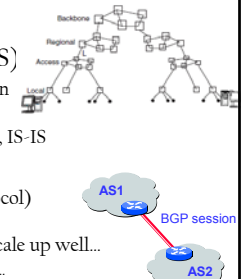
### Datagram communication

- each dispatched packet contains end address
- packets may go via different routes
  - based on target IP address
  - and which connection is active
    - routing tables dynamically updated
- example: Ethernet, IP, UDP
  - better for video conferencing or streaming
    - loss of a few packets not a showstopper

## Heterogeneity of Networks

### Two-level hierarchy

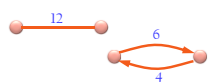
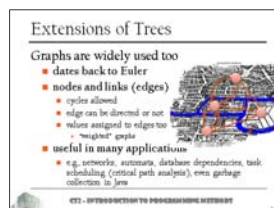
- Autonomous System (AS)
  - usually homogeneous domain
    - Sprint, AT&T, Verizon,...
  - shortest path routing: OSPF, IS-IS
- inter-domain
  - BGP (Border Gateway Protocol)
    - deal with inhomogeneity
  - the other protocols do not scale up well...
    - no guarantee of shortest path...



## Abstraction

### Networks are graphs

- extension of trees
- can be undirected
  - just a link btw nodes
- or directed
  - only in one direction
- and/or weighted
  - weight can be, e.g., RTT



## Storing Graph Connectivity

### Two basic ways of encoding graphs

- adjacency matrix
  - $M[i][j]$  indicates whether  $(i,j)$  is an edge
- linked list
  - or iterator
  - storing neighbors



### Directed and/or weighted graphs?

- same deal, small changes

## Single-Source Shortest Path

Find shortest paths from a source node

- used in OSPF (careful: only non-negative weights!)

Good news: dynamic programming at work

- if R is a node on the minimal path from P to Q, knowing the latter implies knowing the minimal path from P to R
- running time:  $O(E+V^2)$ 
  - can be made  $O(E+V \log V)$  with a priority queue



## Pseudocode

Vertices stored in V  
Source vertex = s

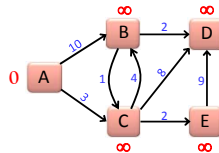
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dist[s] ← 0                                (distance to source vertex is zero)
for all v ∈ V - {s}
  do dist[v] ← ∞                            (set all other distances to infinity)
Q ← V                                       (queue Q initially contains all vertices)
while Q ≠ ∅                               (while queue is not empty)
  do u ← mindistance(Q, dist)              (select/remove element of Q with min. distance)
  for all v ∈ neighbors[u]
    do if dist[v] > dist[u] + w(u, v)        (if new shortest path found)
       then d[v] ← d[u] + w(u, v)          (set new value of shortest path)
  remove u from Q
(if desired, add traceback code by updating, e.g., an array previous[v])
return dist
    
```



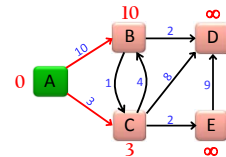
## Example I

Q: A B C D E  
d: 0 ∞ ∞ ∞ ∞



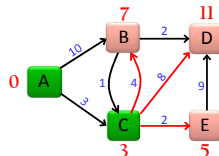
## Example II

Q: A B C D E  
d: 0 10 3 ∞ ∞



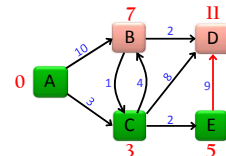
## Example III

Q: A B C D E  
d: 0 7 3 11 5



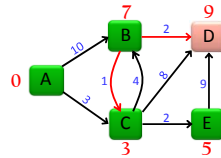
## Example IV

Q: A B C D E  
d: 0 7 3 11 5



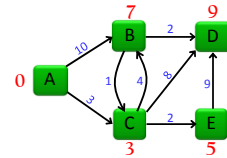
## Example V

Q: A B C D E  
d: 0 7 3 9 5



## Example VI – The End

Q: A B C D E  
d: 0 7 3 9 5



## Shortest Path in a Graph

Not just good for internet communication

- mapquest/Google map
  - roads are weighted by speed limit, average traffic, ...
- epidemiology
  - “contact” network, with weight equals probability
- robot motion planning
  - if you know a map of the environment



source: Wikipedia



## [Edsger W. Dijkstra]

May 11, 1930 – August 6, 2002

- from theoretical physicist...
- to computer scientist
  - made a case against GOTO
  - made programming a science
    - Turing award in 1972



## Graph Search

Two basic orders of search

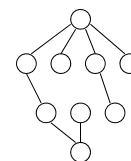
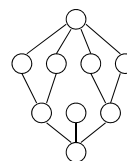
- depth first search (DFS)
  - like trees' preorder: visit node, then its neighbors
- breadth first search (BFS)
  - same order as Dijkstra if equal weights
    - i.e., level by level
  - parallel version: MapReduce
    - programming model implemented in, e.g., Hadoop
    - basically, a tree of executions, then a gathering



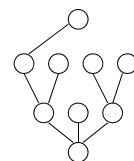
## Spanning Trees

Trees assembled during search

- BFS: short and bushy
- DFS: long and stringy



BFS from top



DFS from top



## Other Typical Graph Problems

### All-pair shortest paths

- dyn. prog. again, called Floyd-Warshall  

$$\text{shortestPath}(i, j, k) = \min(\text{shortestPath}(i, j, k-1), \text{shortestPath}(i, k, k-1) + \text{shortestPath}(k, j, k-1))$$
- if positive weights, Dijkstra's just as good
  - $O(V^3)$

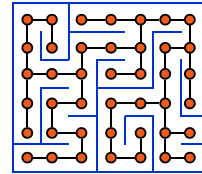
### Minimum spanning tree (e.g., power lines)

- find a subgraph of graph  $G$ 
  - forming a tree containing all vertices of  $G$
  - with min sum of weights
- Prim's algorithm (same complexity as SSSP)



## [Spanning Trees]

Cool way to make a maze too...



## Final Words on Networking

### Notion of ports

- different channels
  - IP header contains add + port

### Notion of sockets

- chat over ports
  - listen/bind,recv/send,...

### Notion of layers

DNS lookup  
FTP  
HTTP  
POP3  
Telnet

UDP  
TCP  
TCP  
TCP  
TCP

53  
21  
80  
110  
23

