

## Topics to discuss

Running Time Analysis  
How to find Time Complexity.

Eg.1

int a = 20; — 1 unit time

int b = 30; — 1 unit time

int sum = a + b; — 1 unit time

print(sum); — 1 unit time

User independent

Algorithm

Iterative

For(i → n)

{  
...  
}

Recursive

A(n)  
{  
...  
A(n/2)  
}

How much time it will take?

Total time = 4 unit of time

Time complexity =  $O(1)$  or  $O(c)$

Eg-2 : Algorithm  $\text{sum}(\underbrace{A[], n})$  input

{  
   $\text{sum} = 0;$  ——— 1 × 1 unit time

for ( $i = 1; i \leq n; i++$ ) ———  $(n+1) \times 1$  unit time

{  
   $\text{sum} = \text{sum} + A[i];$  ———  $n \times 1$

}

Total time =  $1 + (n+1) + n$

$$f(n) = 2n + 1$$

$$f(n) = O(n)$$

$$f(n) \leq c \cdot g(n)$$

Eg-3: Algorithm Sum (A[][] , B[][] , n)

{ for (i=1 ; i<=n ; i++) ——— (n+1) x 1

{ for (j=1 ; j<=n ; j++) ——— n x (n+1)

{ c[i][j] = A[i][j] + B[i][j]; ——— n x n  
}  
}  
}

$$\begin{aligned}\text{Total time} &= (n+1) + n(n+1) + n^2 \\ &= n+1 + n^2 + n + n^2\end{aligned}$$

$$f(n) = 2n^2 + 2n + 1$$

$$\text{Time Complexity} = O(n^2)$$

Eg-4: Algorithm Multiply ( $A[n][n]$ ,  $B[n][n]$ ,  $n$ )

$\{$   
for ( $i=1$  ;  $i \leq n$  ;  $i++$ ) —————  $n+1$

$\{$   
for ( $j=1$  ;  $j \leq n$  ;  $j++$ ) —————  $n \times n+1$

$\{$   
     $C[i][j] = 0$ ; —————  $n \times n$

for ( $k=1$  ;  $k \leq n$  ;  $k++$ ) —————  $n \times n \times n+1$

$\{$   
     $C[i][j] = C[i][j] + A[i][k] * B[k][j]$ ; —————  $n \times n \times n$

$\}$

$\}$

$\}$

$\}$

$$\begin{aligned} \text{Total time} &= n+1 + n(n+1) + n^2 + n^2(n+1) + n^3 \\ &= n+1 + n^2 + n + n^2 + n^3 + n^2 + n^3 \end{aligned}$$

$$f(n) = 2n^3 + 3n^2 + 2n + 1$$

$$\boxed{f(n) = O(n^3)}$$

**Follow Now**



**Start Practicing**



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