

Topics to discuss

- ① Theta Notation (Θ -notation)
- ② Examples

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Theta Notation

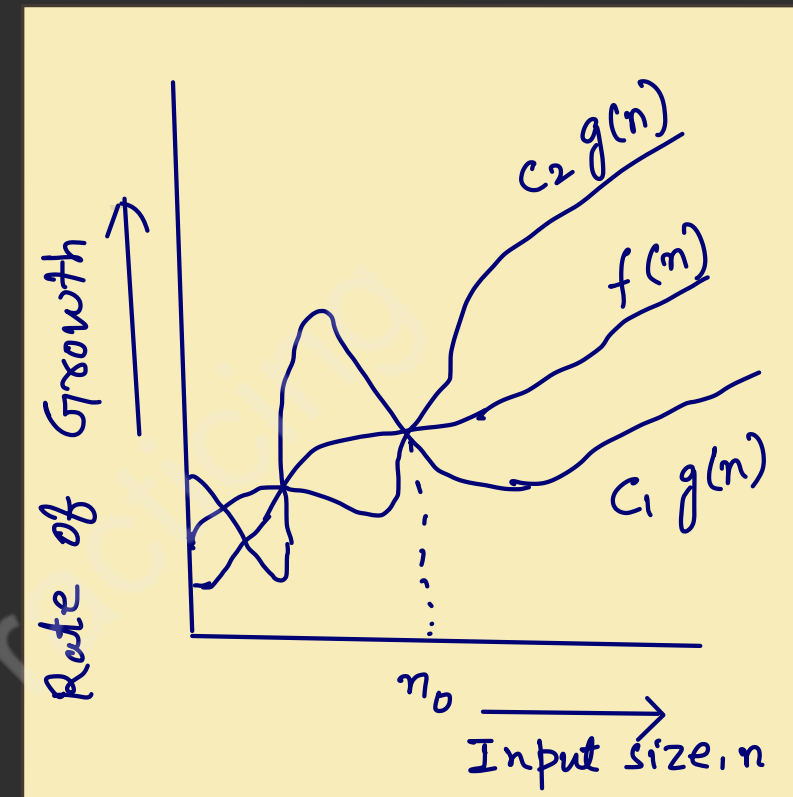
This notation decides whether the upper and lower bounds of a given function (algorithm) are the same. The average running time of an algorithm is always between the lower bound and the upper bound.

It is defined as $\Theta(g(n)) = \{ f(n) : \text{there exist positive constants } c_1, c_2 \text{ and } n_0 \text{ such that } 0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n) \text{ for all } n \geq n_0 \}$.

$g(n)$ is asymptotic tight bound for $f(n)$.

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$\Theta(g(n))$ is the set of functions with the same order of growth as $g(n)$.



① Find Θ bound for $f(n) = \frac{n^2}{2} - \frac{n}{2}$

Solution :- $f(n) = \frac{n^2}{2} - \frac{n}{2}$

By definition, $c_1 g(n) \leq f(n) \leq c_2 g(n)$
 $\frac{1}{5} n^2 \leq \frac{n^2}{2} - \frac{n}{2} \leq n^2$

$c_1 = \frac{1}{5}$, $c_2 = 1$
 $g(n) = n^2$

$f(n) = \Theta(g(n))$

$f(n) = \Theta(n^2)$

② Prove $n \neq \Theta(n^2)$

Solution :-

$$\text{Given, } f(n) = n \\ g(n) = n^2$$

We know,

$$c_1 g(n) \leq f(n) \leq c_2 g(n)$$

$$n^2 \leq n \leq n^2 \leftarrow \boxed{n=1}$$

This is not true.

Hence, $n \neq \Theta(n^2)$ (proved)

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