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LE4.5

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LE4.5.1: 2-input functions

4/4 points (ungraded)

Consider the 2×2 K-map needed to hold the truth table for a 2-input Boolean function.

If the truth table for the 2-input function contained only a single "1" in the output column, what is the maximum number of prime implicants that could be circled in the corresponding K-map?

Maximum number of prime-implicants:



If the truth table for the 2-input function contained exactly two "1s" in the output column, what is the maximum number of prime implicants that could be circled in the corresponding K-map?

Maximum number of prime-implicants:



If the truth table for the 2-input function contained exactly three "1s" in the output column, what is the maximum number of prime implicants that could be circled in the corresponding K-map?

Maximum number of prime-implicants:



If you only had a supply of 2-input NAND gates to build a circuit, what is the *minimum* number of 2-input NANDs you would need to implement any arbitrary 2-input Boolean function? Hint: think about your answers to the questions above and what they imply about a minimal sum-of-products expression for an arbitrary 2-input function. Then think about how to implement a sum-of-products circuits using only NANDs.

Minimum number of 2-input NANDs needed:



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LE4.5.2: Minimal Sum Of Products

3/3 points (ungraded)

The 3-input boolean function $G(A, B, C)$ computes $\overline{A} \cdot \overline{C} + A \cdot \overline{B} + \overline{B} \cdot \overline{C}$.

A) How many 1's are there in the output column of G's 8-row truth table?

☐ 3

☒ 4

☐ 5

☐ 6

☐ none of the above



B) A minimal sum-of-products expression for G is:

☐ $A \cdot \overline{B} \cdot C + \overline{A}$

☐ $A \cdot B \cdot C + A \cdot \overline{B} \cdot \overline{C}$

☒ $A \cdot \overline{B} + \overline{A} \cdot \overline{C}$

- ☐ $\overline{A} \cdot \overline{C} + \overline{B} \cdot \overline{C}$
- ☐ all of the above
- ☐ none of the above



C) There's good news and bad news: the bad news is that the stockroom only has G gates. The good news is that it has as many as you need. Using only combinational circuits built from G gates, one can implement

- ☐ only functions with 3 inputs or less
- ☒ any function (G is universal)
- ☐ only functions with the same truth table as G
- ☐ only inverting functions
- ☐ only non-inverting functions



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💬	question and answer are misconducts	2	▼
the question asks for the minimum number of 2 input nand and the answer is we need at most 5 2-input nand ! i can make an expres...			
💬	Minimum number of 2-input NANDs needed?	2	▼
I think the minimum number of NANDs is 9. As to a 2-input Boolean function, say $f(A,B)$, then $f(A,B) = \text{not}(A)*\text{not}(B) + AB$ would cost ...			
💬	G(A,B,0) == Nand ?	2	▼
I can see that if I only take the terms where C=0 in the Truth table of G — We get NAND's truth table. How can I simplify $**G = \sim C...$			
✓	Why just $A.\sim B + \sim A.\sim C$?	2	▼
For part B of question LE4.5.2, could someone explain to me why the answer is just $A.\sim B + \sim A.\sim C$ but not $A.\sim B + \sim A.\sim C$ and $\sim A.\sim C + \sim B...$			
💬	[STAFF] Minimum number of 2-input NANDs needed to implement arbitrary 2-input Boolean function	8	▼
It seems to me the minimum number of 2-input NANDs needed to implement any arbitrary 2-input Boolean function is only 3. Reason...			

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