

4.5

BOOLEAN ALGEBRA

4.5.1. Boolean Algebra

Let B be a non empty set having two binary operations $+$ and \cdot , one unary operation ' $'$ and two distinct elements 0 and 1 . B is called a Boolean Algebra if the following axioms hold :

1. For every pair of elements a, b in B
 - (i) $a + b = b + a$ (ii) $a \cdot b = b \cdot a$
(Commutative laws)
2. For all elements a, b, c in B
 - (i) $a \cdot (b + c) = (a \cdot b) + (a \cdot c)$ (ii) $a + (b \cdot c) = (a + b) \cdot (a + c)$
(Distributive laws)
3. For every element a in B
 - (i) $a + 0 = a$ (ii) $a \cdot 1 = a$
(Identity law)
4. For every element a there is an element a' in B such that
 - (i) $a + a' = 1$ (ii) $a \cdot a' = 0$
(Complement laws)

Note : (1) We designate it as " $(B, +, ., ', 0, 1)$ is a Boolean Algebra".

(2) The element 1 is called **unit element** and the element 0 is called the **zero element**.

(3) a' is called **complement** of a .

(4) We usually drop the symbol $.$ and use juxtaposition, i.e., instead of writing $a \cdot b$ we shall write ab in practice.

(5) $+$ is called 'sum' (not arithmetic sum) and \cdot is called 'product' (not arithmetic product) and ' $'$ is called complement.

(6) Conventionally the operation $+$ has precedence over the operation \cdot and \cdot has precedence over $+$. For example $x \cdot y'$ means $x \cdot (y')$ and not $(x \cdot y)'$; $x + y \cdot z$ means $x + (y \cdot z) = x + yz$.

Illustrative Examples.

(i) Let S be a set. Then its power set $P(S)$ (set of all subsets of S) is a Boolean algebra under the operations \cup (union), \cap (Intersection), ' (complement) with Φ (null set) as zero element, and S as unit.

Since $A \cup B = B \cup A$ and $A \cap B = B \cap A$ so the commutative property hold :

Since $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$
 $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$ so the distributive law hold.

Since $A \cup \Phi = A$ and $A \cap S = A$ so the Identity law is valid.

Since $A \cup A' = S$ and $A \cap A' = \Phi$ so the complement laws are hold.

So $(P(S), \cup, \cap, ', \Phi, S)$ is a Boolean Algebra.

(ii) Let $D_{40} = \{1, 2, 4, 5, 8, 10, 20, 40\}$ is the set of divisors of 40. Let \vee, \wedge and $\bar{-}$ are defined as $a \vee b = \text{LCM of } a, b; a \wedge b = \text{HCF of } a, b$ and $\bar{a} = \frac{40}{a}$. Then $(D_{40}, \vee, \wedge, \bar{-}, 1, 40)$ is not a Boolean Algebra.

Since $\text{LCM of } a, b = \text{LCM of } b, a$ and $\text{HCF of } a, b = \text{HCF of } b, a$ so $a \vee b = b \vee a$ and $a \wedge b = b \wedge a$ showing the validity of commutative property.

Since $\text{HCF of } a$ and $\text{LCM of } b, c = \text{LCM of } (\text{HCF of } a, b)$ and $(\text{HCF of } a, c)$ so $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$. Similarly, $a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$ hold. This shows that the distributive law hold.

Since $\text{LCM of } a$ and $1 = a$ and $\text{HCF of } a$ and $40 = a$ (as a is a factor of 40) so $a \vee 1 = a, a \wedge 40 = a$. This shows Identity laws are valid

But $2 \vee \bar{2} = 2 \vee 20 = \text{LCM of } 2$ and $20 = 20 \neq 40$. This shows complement laws are invalid.

Hence $(D_{40}, \vee, \wedge, \bar{-}, 1, 40)$ is not a Boolean Algebra.

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(iii) Let the set of all factors of 70,

$D_{70} = \{1, 2, 5, 7, 10, 14, 35, 70\}$. Let \vee and \wedge are defined as $a \vee b = \text{LCM of } a, b; a \wedge b = \text{HCF of } a, b; \bar{a} = \frac{70}{a}$. Then $(D_{70}, \vee, \wedge, \bar{-}, 1, 70)$ is a Boolean Algebra.

(iv) (An Important Boolean Algebra)

Let $B = \{0, 1\}$ be a set. The three operations $+, .$ and $'$ are defined as

| | | |
|---|---|---|
| + | 1 | 0 |
| 1 | 1 | 1 |
| 0 | 1 | 0 |

| | | |
|---|---|---|
| . | 1 | 0 |
| 1 | 1 | 0 |
| 0 | 0 | 0 |

| | | |
|---|---|---|
| ' | 1 | 0 |
| 0 | 0 | 1 |

B is a Boolean Algebra w.r.t the above operations.

From the very definition of the operation we see all the axioms of Boolean Algebra are valid. For instance $0+1=1=1+0$ etc. So $(B, +, ., ', 0, 1)$ is a Boolean Algebra.

(v) Let $B = \{0, 1\}$ be a Boolean Algebra as defined in the previous example. Let $B^n = B \times B \times B \times \dots \times B = \{(x_1, x_2, \dots, x_n) : \text{each } x_i \in B\}$. Then B^n is a Boolean Algebra w.r.t the operation $+, .$ and $'$ defined by

$$(x_1, x_2, \dots, x_n) + (y_1, y_2, \dots, y_n) = (x_1 + y_1, x_2 + y_2, \dots, x_n + y_n);$$

$$(x_1, x_2, \dots, x_n) \cdot (y_1, y_2, \dots, y_n) = (x_1 \cdot y_1, x_2 \cdot y_2, \dots, x_n \cdot y_n) \text{ and}$$

$$(x_1, x_2, \dots, x_n)' = (x'_1, x'_2, \dots, x'_n)$$

Since $x_i + y_i \in B, x_i \cdot y_i \in B$ and $x'_i \in B$ for all x_i, y_i in B .

So $+, .$ are binary operations and $'$ is unary operation on B^n . Let us verify the axioms of Boolean Algebra :

$$\begin{aligned} 1. \quad (i) \quad & (x_1, x_2, \dots, x_n) + (y_1, y_2, \dots, y_n) = (x_1 + y_1, x_2 + y_2, \dots, x_n + y_n) \\ & = (y_1 + x_1, y_2 + x_2, \dots, y_n + x_n) = (y_1, y_2, \dots, y_n) + (x_1, x_2, \dots, x_n) \end{aligned}$$

$$\begin{aligned}
 & \text{(ii)} (x_1, x_2, \dots, x_n) \cdot (y_1, y_2, \dots, y_n) = (x_1 \cdot y_1, x_2 \cdot y_2, \dots, x_n \cdot y_n) \\
 &= (y_1 \cdot x_1, y_2 \cdot x_2, \dots, y_n \cdot x_n) \\
 &= (y_1, y_2, \dots, y_n) \cdot (x_1, x_2, \dots, x_n)
 \end{aligned}$$

So, commutative law hold.

$$\begin{aligned}
 2. \quad & \text{(i)} (x_1, x_2, \dots, x_n) \cdot \{ (y_1, y_2, \dots, y_n) + (z_1, z_2, \dots, z_n) \} \\
 &= (x_1, x_2, \dots, x_n) \cdot \{ (y_1 + z_1, y_2 + z_2, \dots, y_n + z_n) \} \\
 &= (x_1 \cdot (y_1 + z_1), x_2 \cdot (y_2 + z_2), \dots, x_n \cdot (y_n + z_n)) \\
 &= (x_1 \cdot y_1 + x_1 \cdot z_1, x_2 \cdot y_2 + x_2 \cdot z_2, \dots, x_n \cdot y_n + x_n \cdot z_n) \\
 &= (x_1 \cdot y_1, x_2 \cdot y_2, \dots, x_n \cdot y_n) + (x_1 \cdot z_1, x_2 \cdot z_2, \dots, x_n \cdot z_n) \\
 &= (x_1, x_2, \dots, x_n) \cdot (y_1, y_2, \dots, y_n) + (x_1, x_2, \dots, x_n) \cdot (z_1, z_2, \dots, z_n)
 \end{aligned}$$

(ii) Similarly it can be shown that

$$\begin{aligned}
 & (x_1, x_2, \dots, x_n) + \{ (y_1, y_2, \dots, y_n) \cdot (z_1, z_2, \dots, z_n) \} \\
 &= \{ (x_1, x_2, \dots, x_n) + (y_1, y_2, \dots, y_n) \} \cdot \{ (x_1, x_2, \dots, x_n) + (z_1, z_2, \dots, z_n) \}
 \end{aligned}$$

So distributive law hold.

$$\begin{aligned}
 3. \quad & \text{(i)} \{ (x_1, x_2, \dots, x_n) + (0, 0, \dots, 0) \} = (x_1 + 0, x_2 + 0, \dots, x_n + 0) \\
 &= (x_1, x_2, \dots, x_n)
 \end{aligned}$$

$$\begin{aligned}
 & \text{(ii)} (x_1, x_2, \dots, x_n) \cdot (1, 1, \dots, 1) = (x_1 \cdot 1, x_2 \cdot 1, \dots, x_n \cdot 1) \\
 &= (x_1, x_2, \dots, x_n)
 \end{aligned}$$

So, the identity law hold. Here $(0, 0, \dots, 0)$ is the zero element and $(1, 1, \dots, 1)$ is unit element.

$$\begin{aligned}
 4. \quad & \text{(i)} (x_1, x_2, \dots, x_n) + (x_1, x_2, \dots, x_n)'
 \end{aligned}$$

$$= (x_1, x_2, \dots, x_n) + (x'_1, x'_2, \dots, x'_n)$$

$$= (x_1 + x'_1, x_2 + x'_2, \dots, x_n + x'_n) = (1, 1, \dots, 1) = \text{the unit element.}$$

$$\text{(ii)} (x_1, x_2, \dots, x_n) \cdot (x_1, x_2, \dots, x_n)'$$

$$= (x_1, x_2, \dots, x_n) \cdot (x'_1, x'_2, \dots, x'_n) = (x_1 \cdot x'_1, x_2 \cdot x'_2, \dots, x_n \cdot x'_n)$$

$$= (0, 0, \dots, 0) = \text{the zero element.}$$

Thus all the axioms are valid in B'' .

So B'' is a Boolean Algebra.

4.5.2. Duality in Boolean Algebra.

In a Boolean Algebra $(B, +, \cdot, 0, 1)$ the dual of any statement is defined to be the statement that is obtained by interchanging $+$ and \cdot ; 0 and 1 in the original statement. For example the dual of the statement ' $(a+0)+(1 \cdot a')=1$ ' is the statement ' $(a \cdot 1) \cdot (0+a')=0$ '.

Principle of Duality in a Boolean Algebra

If a Theorem hold in a Boolean Algebra then another Theorem is obtained which is nothing but the dual of the former.

This follows from the fact that the dual of each axiom of a Boolean Algebra also comes as an axiom.

Theorem 1. (Idempotent Laws)

In a Boolean algebra $(B, +, \cdot, ', 0, 1)$

$$\text{(i)} a + a = a \quad \text{(ii)} a \cdot a = a$$

Proof. (i) $a + a = (a+a) \cdot 1 \quad \because 1 \text{ is unit element}$

$$= (a+a) \cdot (a+a') \quad \because a' \text{ is complement of } a$$

$$= a + (a \cdot a') \text{ by distributive law}$$

$$= a + 0 = a.$$

Though (ii) follows from (i) by principle of duality we give the independent proof.

$$\begin{aligned}
 \text{(ii)} \quad a \cdot a &= (a \cdot a) + 0 \quad \because 0 \text{ is zero element} \\
 &= a \cdot a + (a \cdot a') \quad \because a' \text{ is complement.} \\
 &= a \cdot (a + a') \quad \text{by distributive property} \\
 &= a \cdot 1 = a.
 \end{aligned}$$

Theorem 2. (Boundedness Law)In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$

(i) $a + 1 = 1$ (ii) $a \cdot 0 = 0$

Proof. (i) $a + 1 = (a + 1) \cdot 1 \quad \because 1 \text{ is unit element}$

$$\begin{aligned}
 &= (a + 1) \cdot (a + a') \quad \because a' \text{ is complement of } a \\
 &= a + (1 \cdot a') \quad \text{by distributive law} \\
 &= a + (a' \cdot 1) \quad \text{by commutative law} \\
 &= a + a' = 1 \quad \text{by complement law}
 \end{aligned}$$

(ii) Similar or follows from (i) by principle of duality.

Theorem 3. (Absorption Laws)In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$

$$\begin{aligned}
 \text{(i)} \quad a + (a \cdot b) &= a \\
 \text{(ii)} \quad a \cdot (a + b) &= a
 \end{aligned}$$

[W.B.U.Tech 2007]

Proof. (i) Similar to (ii)

$$\begin{aligned}
 \text{(ii)} \quad a \cdot (a + b) &= (a + 0) \cdot (a + b) = a + (0 \cdot b) \quad \text{by distributive property} \\
 &= a + (b \cdot 0) = a + 0 = a.
 \end{aligned}$$

Theorem 4. (Associative Laws)In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$

$$\begin{aligned}
 \text{(i)} \quad (a + b) + c &= a + (b + c) \\
 \text{(ii)} \quad (a \cdot b) \cdot c &= a \cdot (b \cdot c)
 \end{aligned}$$

[W.B.U.Tech 2007]

Proof. (i) Let $x = (a + b) + c$ and $y = a + (b + c)$ we shall first prove that $a \cdot x = a \cdot y$.**BOOLEAN ALGEBRA**

$$\begin{aligned}
 \text{Now, } a \cdot x &= a \cdot ((a + b) + c) = a \cdot (a + b) + a \cdot c \\
 &= a + a \cdot c \quad \text{by Absorption Law} \\
 &= a \quad \text{by Absorption Law} \\
 \text{and } a \cdot y &= a \cdot (a + (b + c)) = (a \cdot a) + (a \cdot (b + c)) \\
 &= a + (a \cdot (b + c)) \quad \text{by idempotent Law} \\
 &= a + (a \cdot z) \quad \text{putting } z = b + c \\
 &= a \quad \text{by Absorption Law}
 \end{aligned}$$

Thus $a \cdot x = a \cdot y \quad \dots \quad (1)$

Next we shall show $a' \cdot x = a' \cdot y$

$$\begin{aligned}
 \text{Now, } a' \cdot x &= a' \cdot ((a + b) + c) \\
 &= a' \cdot (a + b) + (a' \cdot c) \quad \text{by distributive Law} \\
 &= ((a' \cdot a) + (a' \cdot b)) + (a' \cdot c) \quad \text{by distributive Law} \\
 &= (0 + (a' \cdot b)) + (a' \cdot c) = (a' \cdot b) + (a' \cdot c) \quad \text{since } 0 \text{ is zero element} \\
 &= a' \cdot (b + c) \\
 \text{and } a' \cdot y &= a' \cdot (a + (b + c)) = (a' \cdot a) + a' \cdot (b + c) \\
 &= 0 + a' \cdot (b + c) = a' \cdot (b + c)
 \end{aligned}$$

Thus $a' \cdot x = a' \cdot y \quad \dots \quad (2)$

Now, $x = 1 \cdot x = (a + a') \cdot x$

$= (a \cdot x) + (a' \cdot x) = (a \cdot y) + (a' \cdot y) \quad \text{by (1) and (2)}$

$= (a + a') \cdot y = 1 \cdot y = y$

$\therefore x = y. \text{ So, } (a + b) + c = a + (b + c).$

(ii) Similar (Follow by principle of duality from (i))

Theorem 5. (Uniqueness of Complement)Let a be any element in a Boolean Algebra $(B, +, \cdot, ', 0, 1)$. If $a + x = 1$ and $a \cdot x = 0$ then x is complement of a .

Proof. The complement of a ,

$$\begin{aligned} a' &= a' + 0 = a' + (a \cdot x) \text{ by hypothesis} \\ &= (a' + a) \cdot (a' + x) \text{ by distributive law} \\ &= 1 \cdot (a' + x) = a' + x \quad \dots \quad (1) \end{aligned}$$

$$\begin{aligned} \text{and } x &= x + 0 = x + (a \cdot a') = (x + a) \cdot (x + a') \\ &= 1 \cdot (x + a') = x + a' = a' + x \text{ by commutative law} \quad \dots \quad (2) \end{aligned}$$

From (1) and (2) we have $x = a'$ i.e., x is complement of a .

Theorem 6. (Involution Law)

In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$, $(a')' = a$ for each element a in B .

Proof. Now, $a + a' = 1$, $a \cdot a' = 0$

$$\therefore a' + a = 1, a \cdot a = 0 \text{ by commutative law}$$

$$\therefore \text{by uniqueness of complement } a = (a')'.$$

Theorem 7. In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$

- (i) $0' = 1$ (ii) $1' = 0$

Proof. By boundedness law $0 + 1 = 1$.

We have also $0 \cdot 1 = 0$ $\therefore 1$ is unit element

So, by uniqueness of complement 1 is the complement of 0, i.e. $1 = 0'$.

By law of duality $0 = 1'$ (or, pose an independent proof)

Theorem 8. (De Morgan's Law)

In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$

- (i) $(a + b)' = a' \cdot b'$ (ii) $(a \cdot b)' = a' + b'$

Proof. (i) We shall prove that

$$(a + b)' + (a' \cdot b') = 1 \text{ and } (a + b) \cdot (a' \cdot b') = 0$$

$$\begin{aligned} \text{Now, } (a + b)' + (a' \cdot b') &= (b + a)' + (a' \cdot b') \text{ by commutative law} \\ &= b + (a + a' \cdot b') \text{ by Associative law (which is shown)} \\ &= b + ((a + a') \cdot (a + b')) \text{ by distributive law} \\ &= b + (1 \cdot (a + b')) = b + (a + b') = b + (b' + a) \\ &= (b + b') + a = 1 + a = 1 \end{aligned}$$

And $(a + b) \cdot (a' \cdot b') = ((a + b) \cdot a') \cdot b'$. Since Associative law hold in Boolean Algebra

$$\begin{aligned} &= ((a \cdot a') + (b \cdot a')) \cdot b' \text{ by distributive law} \\ &= (0 + (b \cdot a')) \cdot b' \\ &= (b \cdot a') \cdot b' = (a' \cdot b) \cdot b' = a' \cdot (b \cdot b') = a' \cdot 0 = 0. \end{aligned}$$

Thus $(a + b)' + (a' \cdot b') = 1$ and $(a + b) \cdot (a' \cdot b') = 0$

So, by uniqueness of complement $(a + b)' = a' \cdot b'$.

(ii) Left to reader as exercise.

Theorem 9. In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$ the following four results are equivalent :

$$(1) a + b = b \quad (2) a \cdot b = a$$

$$(3) a' + b = 1 \quad (4) a \cdot b' = 0$$

Proof. I. Let (1) be true. $\therefore b = a + b$

So, $a \cdot b = a \cdot (a + b) = (a \cdot a) + (a \cdot b) = a + (a \cdot b)$ by Idempotent Law

$= a$ by Absorption law in Boolean Algebra.

Thus (2) is true.

II. Let (2) be true. $\therefore a \cdot b = a$

$$\therefore (a \cdot b)' = a' \quad \therefore a' + b' = a' \text{ by D' Morgans Law}$$

or, $a' + b' + b = a' + b$ or, $a' + 1 = a' + b$ by complement Law

or, $1 = a' + b$ by Boundedness Law

$\therefore a' + b = 1$. Proving (3).

III. Let (3) hold. $\therefore a' + b = 1$

$$\therefore (a' + b)' = 1' \text{ or, } (a')' \cdot b' = 1' \text{ by D'Morgans}$$

or, $a \cdot b' = 0$ by Theorem 7. This proves (4)

IV. Suppose (4) is true. Establish (1). This is left as an exercise.

Illustrative Example.

Ex. 1. Let $(B, +, \cdot, ')$ be a Boolean Algebra and $a, b, c \in B$. Prove that if $a+b=a+c$ and $a \cdot b=a \cdot c$ then $b=c$.

$$\therefore a+b=a+c \quad \therefore (a+b) \cdot b=(a+c) \cdot b$$

$$\text{or, } a \cdot b + b \cdot b = a \cdot b + c \cdot b$$

$$\text{or, } a \cdot b + b = a \cdot b + c \cdot b \text{ by Idempotent Law}$$

$$\text{or, } b = a \cdot c + b \cdot c \text{ by Absorption law and by hypothesis}$$

$$\text{or, } b = a \cdot c + b \cdot c \quad \text{or, } b = (a+b) \cdot c$$

$$\text{or, } b = (a+c) \cdot c \text{ by hypothesis}$$

$$b = a \cdot c + c \cdot c = a \cdot c + c = c. \text{ Hence } b = c.$$

Ex. 2. In a Boolean Algebra $(B, +, \cdot, ', 0, 1)$ if x, y, z belong to B simplify the expression $(x+y)(x+z)(x'y')$.

$$(x+y)(x+z)(x'y')'$$

$$= (x+y)(x+z)((x')' + (y')') \text{ by D'Morgans La}$$

$$= (x+y)(x+z)(x+y) = (x+y)(x+y)(x+z)$$

$$= (x+y)(x+z) \text{ by Idempotent Law}$$

$$= xx + xz + yx + yz = x + xz + yx + yz \text{ by Idempotent Law}$$

$$= x + yx + yz \text{ by Absorption Law}$$

$$= x + yz \text{ by Absorption Law, which is the simplified form.}$$

Note. Though it looks simple we do not know whether it can be simplified more. This answer is obtained in.

4.5.3. Boolean Ring

A Ring $(R, +, \cdot)$ with unit element is called Boolean Ring if $a^2 = a$ for all a in R (a^2 is defined as $a \cdot a$)

Example. The set $B = \{0, 1\}$ is a Boolean Ring w.r.t the compositions $+$ and \cdot and $'$ defined in the following table :

| | | |
|---|---|---|
| + | 1 | 0 |
| 1 | 1 | 1 |
| 0 | 1 | 0 |

| | | |
|---|---|---|
| ' | 1 | 0 |
| 1 | 1 | 0 |
| 0 | 0 | 0 |

| | | |
|------|---|---|
| $1'$ | = | 0 |
| $0'$ | = | 1 |

From the composition table we see all the axioms of Ring are satisfied by B . So B is a Ring with unit element which is 1.

$$\text{Moreover, } 1^2 = 1 \cdot 1 = 1 \text{ by table}$$

$$\text{and } 0^2 = 0 \cdot 0 = 0 \text{ by table}$$

So the set B is a Boolean Ring.

4.5.4. Relation with Boolean Algebra and Boolean Ring.

Theorem 1. (Boolean algebra is made into Boolean Ring)

In a Boolean algebra $(B, +, \cdot, ', 0, 1)$ we defined $a \oplus b = (a \cdot b') + (a' \cdot b)$ and $a \odot b = a \cdot b$. Then (B, \oplus, \odot) is a Boolean Ring.

Proof : We shall first show (B, \oplus, \odot) is a Ring with unit element.

I. Let $a, b, c \in B$ be arbitrary.

(i) Since B is Boolean algebra so the complements a' and b' are in B . Again since the coomposition \cdot is a binary relation so $a \cdot b'$ and $a' \cdot b$ belong to B . So $a \cdot b' + a' \cdot b \in B$ as $+$ is a binary relation on B . Thus $a \oplus b \in B$.

$$(ii) a \oplus (b \oplus c) = a \oplus ((b \cdot c') + (b' \cdot c))$$

$$= a((b \cdot c') + (b' \cdot c))' + a' \cdot ((b \cdot c') + (b' \cdot c))$$

[by definition of \oplus]

$$= a \cdot (b' + (c')') \cdot ((b')' + c') + a' \cdot ((b \cdot c') + (b' \cdot c))$$

(using D'Morgass law)

$$\begin{aligned}
 &= a \cdot ((b' + c) \cdot (b + c')) + a' \cdot (b \cdot c' + b' \cdot c) \\
 &= a \cdot (b' \cdot b + b' \cdot c' + c \cdot b + c \cdot c') + a' \cdot b \cdot c' + a' \cdot b' \cdot c \\
 &\quad (\text{using Distributive property}) \\
 &= a \cdot (0 + b' \cdot c' + c \cdot b + 0) + a' \cdot b \cdot c' + a' \cdot b' \cdot c \\
 &= a \cdot (b' \cdot c' + c \cdot b) + a' \cdot b \cdot c' + a' \cdot b' \cdot c \\
 &= a \cdot b' \cdot c' + a \cdot c \cdot b + a' \cdot b \cdot c' + a' \cdot b' \cdot c \\
 &\because a \oplus (b \oplus c) = a \cdot b \cdot c + a \cdot b' \cdot c' + a' \cdot b \cdot c' + a' \cdot b' \cdot c \quad \dots \quad (1) \\
 &\quad (\text{using commutative property})
 \end{aligned}$$

Again $(a \oplus b) \oplus c$

$$\begin{aligned}
 &= c \oplus (a \oplus b) \\
 &= c \cdot a \cdot b + c \cdot a' \cdot b' + c' \cdot a \cdot b' + c' \cdot a' \cdot b \quad \text{using (1)} \\
 &= a \cdot b \cdot c + a \cdot b' \cdot c' + a' \cdot b \cdot c' + a' \cdot b' \cdot c \quad \dots \quad (2) \\
 &\quad (\text{using commutative property})
 \end{aligned}$$

From 1 and (2) we get

$$a \oplus (b \oplus c) = (a \oplus b) \oplus c$$

(iii) We see $a \oplus 0 = a \cdot 0' + a' \cdot 0$ by definition of \oplus

$$\begin{aligned}
 &= a \cdot 1 + 0 \\
 &= a
 \end{aligned}$$

$\therefore 0$ is null element / zero element in B.

(iv) We see $a \oplus a = a \cdot a' + a' \cdot a$ by definition of \oplus

$$0 + 0 = 0$$

\therefore Additive inverse of a is a itself.

(v) Now by definition of \oplus ,

$$a \oplus b = a \cdot b' + a' \cdot b$$

$$\text{and } b \oplus a = b \cdot a' + b' \cdot a$$

$$= a' \cdot b + a \cdot b' \quad \text{using commutative property}$$

These imply $a \oplus b = b \oplus a$

Thus B is abelian group under the composition \oplus

II. Let $a, b, c \in B$ be arbitrary. By definition of \odot

$$\begin{aligned}
 \text{(i)} \quad a \odot b &= a \cdot b \in B \because \text{the composition } \cdot \text{ is a binary relation} \\
 \text{(ii)} \quad a \odot (b \odot c) &= a \cdot (b \cdot c) \\
 &= (a \cdot b) \cdot c \quad \text{by Associative law of } \cdot \\
 &= (a \odot b) \odot c \quad \text{according to definition of } \odot
 \end{aligned}$$

III. Let $a, b, c \in B$ be arbitrary.

$$\begin{aligned}
 \text{(i)} \quad a \odot (b \oplus c) &= a \cdot (b \cdot c' + b' \cdot c) \quad \text{by definition of } \odot \text{ and } \oplus \\
 &= a \cdot (b \cdot c') + a \cdot (b' \cdot c) \quad \text{by distributive property} \\
 &= a \cdot b \cdot c' + a \cdot b' \cdot c \quad \dots \quad (3)
 \end{aligned}$$

Again $(a \odot b) \oplus (a \odot c)$

$$\begin{aligned}
 &= (a \cdot b) \oplus (a \cdot c) \quad \text{by definition of } \odot \\
 &= (a \cdot b) \cdot (a \cdot c)' + (a \cdot b)' \cdot (a \cdot c) \\
 &= (a \cdot b) \cdot (a' + c') + (a' + b') \cdot (a \cdot c) \quad \text{by D'Morgan's law} \\
 &= (a \cdot b) \cdot a' + (a \cdot b) \cdot c' + a' \cdot (a \cdot c) + b' \cdot (a \cdot c) \quad \text{by distributive law} \\
 &= b \cdot (a \cdot a') + a \cdot b \cdot c' + (a' \cdot a) \cdot c + a \cdot b' \cdot c \\
 &= b \cdot 0 + a \cdot b \cdot c' + 0 \cdot c + a \cdot b' \cdot c \\
 &= a \cdot b \cdot c' + a \cdot b' \cdot c \quad \dots \quad (4)
 \end{aligned}$$

From (3) and (4) we get

$$a \odot (b \oplus c) = (a \odot b) \oplus (a \odot c) \quad \dots \quad (5)$$

$$\text{(ii)} \quad (b \oplus c) \odot a = (b \oplus c) \cdot a$$

$$= a \cdot (b \oplus c) \quad \text{by commutative property of } \cdot$$

$$= a \odot (b \oplus c)$$

$$= (a \odot b) \oplus (a \odot c) \quad \text{using (5)} \quad \dots \quad (6)$$

Thus all the axioms of Ring are satisfied the composite \oplus and \odot in B. So (B, \oplus, \odot) is a Ring.

Next we see for arbitrary element a in B ,

$a \odot 1$ (1 is the unit element in the Boolean Algebra)

$= a \cdot 1$ by definition of \odot

$= a$

Thus the ring (B, \oplus, \odot) has unit element.

Again, for arbitrary a in B ,

$$a^2 = a \odot a = a \cdot a$$

$= a$ by Idempotent law in Boolean Algebra.

So (B, \oplus, \odot) is a Boolean Ring

Theorem 2. If $(R, +, \cdot)$ is a Boolean ring prove that

$$(i) a + a = 0 \quad \forall a \in R$$

i.e. each element of R is its own additive inverse

$$(ii) a + b = 0 \Rightarrow a = b$$

(iii) R is commutative ring.

$$(i) a \in R \Rightarrow a + a \in R$$

$$\Rightarrow (a + a)^2 = a + a, \text{ by the given condition}$$

$$\Rightarrow (a + a) \cdot (a + a) = a + a$$

$$\Rightarrow (a + a) \cdot a + (a + a) \cdot a = a + a, \text{ using distributive law}$$

$$\Rightarrow (a^2 + a^2) + (a^2 + a^2) = a + a \quad [\because a \cdot a = a^2]$$

$$\Rightarrow (a + a) + (a + a) = (a + a) + 0 \quad [\because a + 0 = a]$$

$\Rightarrow a + a = 0$, (by left cancellation law for addition in R)

$$(ii) \text{ Now } a + b = 0 \Rightarrow a + b = a + a \quad [\text{ by (i)}]$$

$\Rightarrow b = a$ (by left cancellation law)

$$(iii) \text{ We have, } (a + b)^2 = a + b$$

$$\Rightarrow (a + b) \cdot (a + b) = a + b$$

$$\Rightarrow (a + b) \cdot a + (a + b) \cdot b = a + b \quad (\text{by distributive law})$$

$$\Rightarrow a^2 + b \cdot a + a \cdot b + b^2 = a + b \quad (\text{by distributive law})$$

$$\Rightarrow (a + b \cdot a) + (a \cdot b + b) = a + b \quad (\text{by given condition})$$

$$\Rightarrow (a + b) + (b \cdot a + a \cdot b) = (a + b) + 0, \text{ using commutative and associative property}$$

$$\Rightarrow b \cdot a + a \cdot b = 0, \text{ by left cancellation law}$$

$$\Rightarrow a \cdot b = b \cdot a \quad (\text{by (ii)})$$

$\therefore R$ is a commutative ring.

Theorem 3. (Boolean Ring is made into Boolean Algebra).

A Boolean Ring $(R, +, \cdot)$ is a Boolean Algebra with respect to the compositions \vee , and \wedge defined by

$a \vee b = a + b + a \cdot b; a \wedge b = a \cdot b; a' = 1 + a$ where 1 is the unit element in R .

Proof. By axioms of Ring $a + b, a \cdot b$ and $1 + a \in R$ and so $a \vee b$ and $a \wedge b$ and a' all belong to R for all element a in R . So the compositions \vee and \wedge are binary operations and the composition ' $'$ is unary composition on R .

Let us verify the axioms of Boolean Algebra :

1. For all a, b in R

$$(i) a \vee b = a + b + a \cdot b$$

$$= b + a + b \cdot a \quad [\because + \text{ and } \cdot \text{ are commutative}]$$

$$= b \vee a$$

$$(ii) a \wedge b = a \cdot b$$

$$= b \cdot a \quad [\because \cdot \text{ is commutative}]$$

2. For all elements a, b, c in R

$$(i) a \wedge (b \vee c) = a \wedge (b + c + b \cdot c)$$

$$= a \cdot (b + c + b \cdot c)$$

$$= (a \cdot b) + (a \cdot c) + a \cdot (b \cdot c) \quad [\text{by distributive law of } \cdot \text{ over } +] \quad (1)$$

$$\text{Again } (a \wedge b) \vee (a \wedge c)$$

$$= (a \cdot b) \vee (a \cdot c) \text{ by definition of } \wedge$$

$$= (a \cdot b) + (a \cdot c) + (a \cdot b) \cdot (a \cdot c)$$

$$= (a \cdot b) + (a \cdot c) + ((a \cdot b) \cdot a) \cdot c \quad [\because \cdot \text{ is Associative}]$$

$$= (a \cdot b) + (a \cdot c) + (a \cdot (a \cdot b)) \cdot c \quad [\because \cdot \text{ is commutative}]$$

$$= (a \cdot b) + (a \cdot c) + (a^2 \cdot b) \cdot c$$

$$= (a \cdot b) + (a \cdot c) + (a \cdot b) \cdot c \quad [\because R \text{ Boolean Ring}]$$

$$= (a \cdot b) + (a \cdot c) + a \cdot (b \cdot c)$$

From (1) and (2) we have

$$a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$$

$$(ii) \quad a \vee (b \wedge c) = a \wedge (b \cdot c)$$

$$= a + (b \cdot c) + a \cdot (b \cdot c)$$

$$\text{and } (a \vee b) \wedge (a \vee c)$$

$$= (a \vee b) \cdot (a \vee c) \text{ by definition of } \wedge$$

$$= (a + b + a \cdot b) \cdot (a + c + a \cdot c) \text{ by definition of } \vee$$

$$= a \cdot a + a \cdot c + a \cdot (a \cdot c) + b \cdot a + b \cdot c + b \cdot (a \cdot c) + (a \cdot b) \cdot a + (a \cdot b) \cdot c \\ + (a \cdot b) \cdot (a \cdot c)$$

[using distributive property]

$$= a^2 + a \cdot c + a^2 \cdot c + b \cdot a + b \cdot c + b \cdot (a \cdot c) + b \cdot a^2 + (a \cdot b) \cdot c + b \cdot (a^2) \cdot c \\ = a + a \cdot c + a \cdot c + b \cdot a + b \cdot c + b \cdot (a \cdot c) + b \cdot a + (a \cdot b \cdot c) + b \cdot (a \cdot c) \\ \quad [\because R \text{ is Boolean Ring}]$$

$$= a + 0 + (b \cdot c) + (b \cdot a + b \cdot a) + (b \cdot (a \cdot c) + b \cdot (a \cdot c)) + (a \cdot b \cdot c)$$

$\because a + a = 0$ in a Boolean Ring.]

$$= a + 0 + 0 + (b \cdot c) + (a \cdot b \cdot c)$$

$$= a + (b \cdot c) + a \cdot (b \cdot c)$$

By (3) and (4) we have

$$a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$$

3. (i) We see for the null element / zero element 0 in R,

$$a \vee 0 = a + 0 + a \cdot 0$$

$$= a + 0 \quad [\because a \cdot 0 = 0 \text{ in Ring}]$$

$$= a \quad \because 0 \text{ is null element}$$

$$\therefore a \vee 0 = a$$

(ii) For the unit element 1 in R,

$$a \wedge 1 = a \cdot 1 = 1$$

4. (i) We see $a \vee a' = a + a' + a \cdot a'$

$$= a + (1 + a) + a \cdot (1 + a) \text{ by definition of } a'$$

$$= 1 + (a + a) + a \cdot 1 + a \cdot a$$

$$= 1 + 0 + a + a \quad \because R \text{ is Boolean Ring}$$

$$= 1 + 0 \quad [\because a + a = 0 \text{ in Boolean Ring}]$$

$$= 1$$

$$(ii) \quad a \wedge a' = a \wedge (1 + a) = a \cdot (1 + a)$$

$$= (a \cdot 1) + (a \cdot a) = a + a^2$$

$$= a + a \quad [\because R \text{ is Boolean Ring}]$$

$$= 0 \quad [\text{by the previous theorem}]$$

Thus all the axioms of Boolean Algebra are satisfied. So $(R, \vee, \wedge, 1, 0)$ is a Boolean Algebra.

Ex. In a Boolean Ring $(R, +, \cdot)$ prove that $(a+b)^2 = a^2 + 2ab + b^2$
 $(ab \equiv a \cdot b)$.

$$\text{Solution.} \quad a^2 + 2ab + b^2 = a^2 + ab + ab + b^2$$

$$= a(a+b) + (a+b)b \text{ by distributive property}$$

$$= a(a+b) + b(a+b) \quad \because R \text{ is commutative Ring}$$

$$= (a+b)(a+b) = (a+b)^2$$

4.5.4. Boolean Function.

Boolean Expression.

Let $(B, +, \cdot, ', 0, 1)$ be a Boolean Algebra. Let x_1, x_2, \dots, x_n be n variables on B , i.e. they take values from B . An expression built up from these x_i using the operations $+$, \cdot and $'$ is called a Boolean Expression. For example $E = (xz'y + xy')' + (x' + zy')'$ is a Boolean expression of x, y, z which are variables on B . Here the operation (\cdot) is kept understood.

Boolean Function.

Let $(B, +, \cdot, ', 0, 1)$ be a Boolean Algebra. A mapping $f : B^n \rightarrow B$ is called a Boolean function of n number of variables.

Illustration.

Consider the Boolean Algebra $B = \{0, 1\}$ as cited in an earlier example.

Let $f(x, y) = x'y + y$ where x, y are variables on B . Then this f is a Boolean function of two variables.

If $x = 0, y = 1$ then $f(0, 1) = 0'1 + 1 = 11 + 1 = 1 + 1 = 1$ etc.

Note. The value of a Boolean function of some variable is a Boolean expression.

4.5.5. Simplification of a Boolean Expression Literal.

A single variable or its complement is called a literal. For example x, x', y, y' are literal of x, y etc.

Conjunctive Normal Form (Product of Sum)

A Boolean expression E is said to be in conjunctive normal form (CNF) of r variables x_1, x_2, \dots, x_r if the expression, E can be expressed as $E = E_1 E_2 \dots E_r$ where each E_i = sum of a finite number of literals of each of x_1, x_2, \dots, x_r and $E_i \neq E_j$ for all distinct i, j .

Illustration.

(i) $E = (x' + z + y)(x' + y' + z')(x + y + z)$ is a CNF of the variables x, y and z .

Here $E_1 = x' + z + y, E_2 = x' + y' + z', E_3 = x + y + z$.

But $E = (x' + z)(x' + y' + z')(x + y + z)$ is not so as the expression $x' + z$ does not contain the literal of y .

(ii) The expression $E = x + x'y$ is not a CNF.

(iii) Express the Boolean expression $(x + y + z)(xy + x'z)'$ in CNF

$$(x + y + z)(xy + x'z)' = (x + y + z) \{(xy)'(x'z)'\}$$

$$= (x + y + z) \{ (x' + y') \{ (x')' + z' \} \}$$

$$= (x + y + z) (x' + y') (x + z')$$

$$= (x + y + z) \{ (x' + y') + 0 \} \{ (x + z') + 0 \}$$

$$= (x + y + z) \{ (x' + y' + zz') (x + z' + yy') \}$$

$$= (x + y + z) \{ (x' + y' + z) (x' + y' + z') (x + z' + y) (x + z' + y') \}$$

by distributive law

(iv) Express the Boolean function $(x' + y')(x + y)$ in CNF in the variables x, y and z

$$(x' + y')(x + y) = \{ (x' + y') + zz' \} \{ (x + y) + zz' \}$$

$= (x' + y' + z) (x' + y' + z') (x + y + z) (x + y + z')$ by distributive law

(v) Express $ab + a'b'$ in CNF of the variables a, b

$$ab + a'b' = (ab + a')(ab + b')[\text{By distributive law}]$$

$$= (a' + ab)(b' + ab)$$

$$= (a' + a)(a' + b)(b' + a)(b' + b)$$

$$= 1(a' + b)(b' + a)1 = (a' + b)(b' + a) \text{ which is the required CNF.}$$

Fundamental Product.

A literal or product of two or more literals in which no two literals involve the same variable is called Fundamental Product.

For example $z'x, x, x'yz, y'$ are all fundamental product. But $xyx'z$ and xyz are not fundamental product.

Minterm.

A Fundamental Product which involves all variables is called a Minterm. For example, for these variables x, y, z the product $x'yz$ is a minterm but xy' is not a minterm.

Inclusion

Let P and Q be two fundamental products. P is said to be contained in or included in Q if every literal of P is also a literal of Q .

For example $x'z$ is contained in $x'yz$ but $x'z$ is not contained in $xy'z$.

Theorem. If a fundamental product P is contained in another fundamental product Q then $P + Q = P$.

Proof. Since P is contained in Q we can express $Q = P \cdot P_1$ where P_1 is another product.

Then $P + Q = P + P \cdot P_1 = P$ by Absorption law.

Illustration.

$$\begin{aligned} \text{(i)} \quad & xz' + z'yx = xz' + xz'y \text{ by commutative property} \\ & = (xz') + (xz')y = xz'. \end{aligned}$$

(ii) $xyz'yx$ is not a fundamental product. But it can be expressed into a fundamental product as follow :

$$xyz'yx = xyyz'x = xyz' \text{ (by idempotent law)}$$

which is a fundamental product. Note that it is a min-term also.

Sum-of-Product (Disjunctive Normal Form)

A fundamental product or the sum of two or more fundamental products none of which is contained in another is called a Sum-of-Product.

Illustration.

(i) $E = xz' + y'xz + x'y'$ is a Sum-of-Product i.e. DNF

(ii) $E_1 = yz' + xy'z + z'xy$ is not a Sum-of-Product since the fundamental product yz' is contained in $z'xy = yz'x$. Of course we can express $E_1 = yz' + yz'x + xy'z = yz' + xy'z$.

(iii) We can express any Boolean Expression e.g. $E = ((xy)'z)' + ((x'+z)(y'+z'))'$ into a Sum-of-Product form. This is shown below :

$$E = ((xy)'z)' + ((x'+z)(y'+z'))'$$

$$= \left\{ ((xy)')' + z' \right\} \{ (x'+z)' + (y'+z')' \} \text{ removing complement on parenthesis first.}$$

$$= \{ (xy) + z' \} \{ (x')'z' + (y')'(z')' \} = \{ (xy) + z' \} \{ xz' + yz \}$$

$$= \{ (x+z')(y+z') \} \{ xz' + yz \} \text{ using distributive property where ever possible}$$

$$= (xy + xz' + z'y + z'z')(xz' + yz) = xyxz' + xyyz + z'xz' + z'yz$$

$$= xxyz' + xzyy + xz'z' + yzz' = xyz' + xzy + xz' + y0$$

$$= xyz' + xyz + xz' + 0 = xy(z' + z) + xz'$$

$$= xy1 + xz' = xy + xz', \text{ which is a Sum-of-Product.}$$

[W.B.U.T 2007]

Complete Sum-of-Product (Full Disjunctive Normal Form) and Min-term

Let E be a Boolean Expression of n variables x_1, x_2, \dots, x_n . If E is a Sum-of-Product where each product P involves all the n -variables then E is called a Complete- Sum-of-Product.

Thus every product in a complete Sum-of-Product is Min-term.

For example $E(x, y, z) = xy'z + yzx + x'y'z'$ is a Complete Sum-of-Product of the three variables x, y and z . Here $xy'z$, yzx are minterms of E .

Theorem. Every non-zero Boolean expression can be expressed as a complete Sum-of-Product uniquely.

Proof. Beyond the scope of the text.

Illustrative Examples :

Ex. 1. Express the Boolean expression $z(x'y)'$ in a complete Sum-of-Product form.

Note that the given expression is not a Sum-of-Product.

$$\text{Now, } z(x'y)' = z((x')' + y') = z(x + y')$$

$= zx + zy'$, this is not a complete sum of product but it is a Sum-of-Product.

$$= zx1 + zy1$$

$$= zx(y + y') + zy'(x + x')$$

$$= zxy + zx'y' + zy'x + zy'x'$$

$$= zxy + zy'x + zy'x + zy'x'$$

$= zxy + zy'x + zy'x'$ which is a complete-Sum-of Product.

Ex. 2. Express $E = y' + z(x' + y)$ as a Full disjunctive normal form.

$$\begin{aligned} E &= y' + zx' + zy = y'11 + zx'1 + zy1 = y'(x + x')(z + z') + \\ &\quad zx'(y + y') + zy(x + x') \\ &= xyz + xy'z + xy'z' + x'yz + x'y'z + x'y'z'. \end{aligned}$$

Ex. 3. If A, B, C are subsets of the set S then express $(A^c \cup C)^c \cap (B \cap C)^c$ as union of intersection.

Since the set of all subsets of S is Boolean Algebra under the operation $+$ $\equiv \cup$. and \cdot $\equiv \cap$ so we use all the properties of Boolean Algebra in this expression.

For convenience we replace \cup by $+$ and \cap by \cdot and A^c by A' .

$$(A^c \cup C)^c \cap (B \cap C)^c = (A' + C)' \cdot (B \cdot C)' = \{(A')' \cdot C\} \cdot \{B' + C\}$$

by D'Morgans law

$$= (A \cdot C') \cdot (B' + C') = A \cdot B' + A \cdot C' + C' \cdot B' + C' \cdot C'$$

$$= A \cdot C' \cdot B' + A \cdot C' \cdot C' = A \cdot C' \cdot B' + A \cdot C'$$

$$= (A \cap C^c \cap B^c) \cup (A \cap C^c)$$

Ex. 4. Reduce the expression $E = xy + x'z' + x'y' + z'y$ to a simplified form.

We express each product into a complete Sum-of-Product:

$$xy = xy(z + z') = xyz + xyz'$$

$$x'z' = x'z'(y + y') = x'z'y + x'z'y'$$

$$x'y' = x'y'(z + z') = x'y'z + x'y'z'$$

$$z'y = z'y(x + x') = z'yx + z'yx'$$

We see the summands of $x'z'$ are $x'z'y$ and $x'z'y'$ which appear among the summands of $z'y$ and $x'y'$ respectively.

So we delete $x'z'$ and get $E = xy + x'y' + z'y$... (1)

We see there are no such Prime Implicants whose all summands appear among others.

So (1) is the required minimal form of the Boolean expression.

Ex. 5. Simplify the Boolean expression $x'y(y + z) + yz(y' + x')$

$$x'y(y + z) + yz(y' + x') = x'yy + x'yz + yzy' + yzx'$$

$$= x'y + x'yz + z0 + x'yz = x'y + x'yz + x'yz + 0$$

$$= x'y + x'yz = x'y \text{ (by absorption law)}$$

Ex. 6. Express the expression $x'yz + x'y'z' + x'y + xy'z'$ into the simplified form.

$$x'yz + x'y'z' + x'y + xy'z'$$

$$= x'yz + x'y + x'y'z' + xy'z'$$

$$= x'y + (x' + x)y'z' \text{ (by absorption and distributive law)}$$

$$= x'y + 1y'z' = x'y + y'z'.$$

Note. In the last article of this chapter we shall show another process called KARNAUGH MAP process to find the minimal expression of a Boolean function.

4.5.6. Truth Table of a Boolean Function

Let $B = \{0, 1\}$ be a Boolean Algebra w.r.t the operations $+$, \cdot and $'$ defined as

| | | | | | | |
|---|---|---|---|---|---|--|
| | 1 | 0 | | 1 | 0 | |
| 1 | 1 | 1 | 1 | 1 | 0 | |
| 0 | 1 | 0 | 0 | 0 | 0 | |

$1' = 0$ $0' = 1$

(Earlier we have shown B is a Boolean Algebra).

Let $f(x_1, x_2, \dots, x_n)$ be a Boolean function of the n variables $x_1, x_2, \dots, x_n \in B = \{0, 1\}$.

The table which shows all the possible values attained by f corresponding to all possible values taken by the variables is known as Truth Table of the Boolean function f .

Illustration.

(i) Let $f(x, y) = x' + y'$ be a Boolean function on $B = \{0, 1\}$. Then the truth table of f is

| x | y | x' | y' | $f(x' + y')$ |
|-----|-----|------|------|--------------|
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 | 0 |

We write $T(f) = 1110$

OR

$$T(0011, 0101) = 1110.$$

(ii) Construct the truth table of the Boolean function

$$f(x, y, z) = (yz + xz')(xy' + z)'.$$

First we find the minimal expression of $f(x, y, z)$.

$$\begin{aligned} f(x, y, z) &= (yz + xz') \{(xy')'z'\} \\ &= (yz + xz')((x' + (y')')z') = (yz + xz')((x' + y)z') \\ &= (yz + xz')(x'z' + yz') \\ &= yzx'z' + yzyz' + xz'x'z' + xz'yz' \quad (\text{This is not Sum-of-Product}) \\ &= yx'0 + y0 + 0z' + xyz' = 0 + 0 + 0 + xyz' = xyz'. \end{aligned}$$

This is complete sum-of-product and can not be reduced further.

So, $f(x, y, z) = xyz'$ is the minimal form.

Its Truth Table is

| x | y | z | z' | xy | xyz' |
|-----|-----|-----|------|------|--------|
| 0 | 1 | 1 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 | 1 | 1 |

$$\text{We write } T(f) = 00000001$$

$$\text{or, } T = (00001111, 10100011, 11001010) = 00000001$$

(iii) Using Truth table show that $xy' + xy + x'y = x + y$

Let the variable x, y take values on the Boolean Algebra $\{0, 1\}$ w.r.t the binary operation as mentioned earlier. If the Truth table of the two Boolean functions on LHS and RHS are same then we can claim the two sides are equal :

Truth Table of $(xy' + xy + x'y)$

| x | y | x' | y' | xy' | xy | $x'y$ | $xy' + xy$ | $xy' + xy + x'y$ |
|-----|-----|------|------|-------|------|-------|------------|------------------|
| 1 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 |
| 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 1 |

Truth Table of $x + y$

| x | y | $x + y$ |
|-----|-----|---------|
| 1 | 0 | 1 |
| 1 | 1 | 1 |
| 0 | 0 | 0 |
| 0 | 1 | 1 |

we see the Truth table of the two functions are same.
So $xy' + xy + x'y = x + y$

Construction of Boolean function from Truth Table.

This is shown by the following example :

Construct the Boolean function and simplify it given the following table:

| x | y | z | $f(x, y, z)$ |
|---|---|---|--------------|
| 1 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 0 | 1 | 1 |
| 0 | 0 | 0 | 0 |

$$\text{or, } T(11110000, 11001100, 10101010) = 01100110.$$

f gets values 1 at 2nd, 3rd, and 6th, 7th row. The terms corresponding to these rows are xyz' , $xy'z$, $x'yz'$ and $x'y'z$ respectively. The terms corresponding to other rows are 0.

So the Boolean expression is

$$\begin{aligned} & xyz' + xy'z + x'yz' + x'y'z \\ &= (x + x')yz' + (x + x')y'z = 1yz' + 1y'z = yz' + y'z \end{aligned}$$

Which is a simplified form.

i. the required Boolean function is $f(x, y, z) = yz' + y'z$.

Exercises

I. Short Answer Questions

- Prove that the set of divisors of 70 is a Boolean Algebra w.r.t the binary operation $+$ and \cdot and unary operation $'$, where $a + b = \text{lcm}$ of a, b , and $a \cdot b = \text{HCF}$ a, b and $a' = 70/a$.

2. Prove that the set of divisors of

- 55
 - 20
- Boolean algebra (operations to be introduced by you)

3. D_{110} = set of divisors of 110. Show its diagram as a Boolean algebra.

4. Prove that the power set of a set S is a Boolean Algebra w.r.t the operation \cap and \cup .

5. Construct a Boolean Algebra with only two elements

6. Prove that in a Boolean Algebra (B, \vee, \wedge) $a \vee a = a$ and $a \wedge a = a$

7. In a Boolean Algebra $(B, +, \cdot)$ prove that $a+1=1$. Write the dual of this relation

8. Prove that in a Boolean Algebra the two unit elements viz 0 and 1 are unique.

9. State and prove the 'Absorption law' in a Boolean Algebra

10. State D' Morgans law in a Boolean Algebra. Prove one of the two.

11. In a Boolean algebra under $+$ and \cdot , prove that

$$(i) a' + b = 1 \Rightarrow a + b = b.$$

$$(ii) a \cdot b' = 0 \Rightarrow a + b = b$$

12. Define fundamental product and min-term of three variables in a Boolean algebra. Give an example of a fundamental product which is not a min-term

13. In a Boolean Algebra prove that

$$(i) B + C = C + BC' \quad (ii) AB + CA = (C + A')(B + A)$$

$$(iii) (x' + y) \cdot (x + y) = y \quad (iv) xy + xyz + x'y + xy'z = y + zx$$

$$(v) b(a+c)(c'+a') + (b+c)d = (c' + abd + a'b + b'd)(c + ab + bd)$$

$$(vi) (f + g)'h + ghf' = hf'$$

$$(vii) (a + b + c) \cdot (abc)' = a \cdot b' + b \cdot c' + c \cdot a'$$

14. Write the dual of the following relations in Boolean Algebra

$$(i) \quad a + a' \cdot b = a + b \quad (ii) \quad (a \cdot 1) \cdot (0 + a') = 0$$

$$(iii) \quad (a + 1)(a + 0) = a \quad (iv) \quad (a + b)(b + c) = ac + b$$

$$(v) \quad x'y'z + x'yz' \quad (vi) \quad x'z + xy(y + y'z)$$

$$(vii) \quad (yz + y'z')x \quad (viii) \quad (a + b) \cdot (a + 1) = a + a \cdot b + b$$

$$(ix) \quad A \cdot B' = 0 \text{ iff } A \cdot B = A$$

[Hint : Replace + by . , . by +, 0 by 1 and 1 by 0 and get the result]

15. By applying D'Morgans's law find the complement of

$$(i) \quad (x + x'y) \quad (ii) \quad C(A + B') \quad (iii) \quad (y + u') \cdot (x' + z)$$

$$(iv) \quad y'z + x' \quad (v) \quad (A + B) \cdot (A' + B)$$

16. Find the complement of the expression $xy' + xz + y'z$.

17. Simplify the following Boolean expression (Just by using properties of Boolean Algebra)

$$(i) \quad (x + y)(x + z) + xyz \quad (ii) \quad a'bc + a'bc' + abc' + abc$$

$$(iii) \quad (x + y')(y + z')(y' + x')(x' + z)$$

$$(iv) \quad y'x' + y'x + yx \quad [\text{W.B.U.T. 2007}] \quad (v) \quad a' + (ab' + ac)'$$

$$(vi) \quad xy' + xy + x'y \quad (vii) \quad C(B + C)(A + B + C)$$

$$(viii) \quad A + B(B + A) + (B + A')A \quad (ix) \quad x + x'yz' + (y + z)'$$

$$(x) \quad x[y + z(xy + xz)'] \quad (xi) \quad (AB' + C) \cdot (A + B')C$$

$$(xii) \quad x'yz' + xyz' + xy'z' + x'y'z'.$$

18. Express the set expression $(A \cup B)^c \cap (C^c \cup B)$ as union of intersection with the help of the Boolean Algebra.

19. For any Boolean Algebra B , Prove that

$$(i) \quad (a + b)(b + c)(c + a) = ab + bc + ca \text{ for all } a, b, c, \text{ in } B$$

$$[\text{Hint : LHS} = (ab + ac + bb + bc)(c + a) = (ab + ac + b + bc)(c + a)]$$

$$= abc + ab + ac + ac + bc + ba + bc + abc = abc + ab + bc + ca$$

$$= ab + bc + ca (\because ab \text{ is contained in abc})$$

[W.B.U.T. 2001]

$$(ii) \quad (xy' + xz') + x' = (x' + y + z)(x' + y + z')(x' + y' + z')$$

[W.B.U.T. 2002, 2005]

$$(iii) \quad (x + y)(x' + z)(y + z) = xz + x'y + yz = (x + y)(x' + z)$$

20. Express the following expression in CNF in the variables present in the expression :

$$(i) \quad yz + x' \quad (ii) \quad (a' + b')c + (a + c)(a' + c')$$

$$(iii) \quad \{(y + z')' + x\}(z' + y + x') \quad (iv) \quad \{(a' + b) + (a' + b)'\}'$$

$$(v) \quad x' + (xz + xy)' \quad (vi) \quad (a' + b')(a + b)(a + c)$$

21. Reduce the following expression to either 0 or a fundamental product. Are the reduced form Min term ?

$$(i) \quad x \ y \ z' \ y' \ s \ t \ s \quad (vi) \quad t \ y' \ t \ z' \ y \ x$$

22. Reduce the following expression into Disjunctive Normal form (Sum-of-Product) and then to complete sum of product.

$$(i) \quad (xy' + y'z + x'y)x \quad (ii) \quad x'y + (x' + y)'$$

$$(iii) \quad (xy + x'y + y')x \quad (iv) \quad (z' + y)(y'z + x)$$

$$(v) \quad (yx' + y'x + y'z)x \quad (vi) \quad (yx')'(x' + xyz')$$

$$(vii) \quad x_1x_2' + x_3$$

$$(viii) \quad (x + y')(y + z')(z + x')(x' + y')$$

$$(ix) \quad \{(xy')'z\}'$$

[W.B.U. Tech 2005]

[W.B.U. Tech 2006]

23. Find the disjunctive normal form (Sum-of-Product) for the Boolean expression

$$(i) g(w, x, y, z) = wx\bar{y} + w\bar{y}\bar{z} + xy$$

$$(ii) (x + y + z)(xy + x'z')$$

[W.B.U. Tech B 2003]

24. Transform the following CNF into DNF :

$$(i) (x + y')(y + x')$$

$$(ii) (x' + y' + z')(y' + x' + z)(x' + y + z)(x + y' + z')(x + y + z')$$

25. Find the complements of the following Boolean function

$$(i) x'y'z' + xy'z + zyx' + xyz \quad (ii) y'x + yx'$$

$$(iii) x(y'z' + yz) \quad (iv) ab' + ac + b'c$$

26. Find $T(f)$, the truth table of the following Boolean function:

$$(i) f(x, y, z) = x'yz + xyz' \quad (ii) f = yx' + zx \quad (iii) f = x(y + x')$$

27. Find the Boolean function from the following Truth Table :

$$(i) T(00001111, 00110011, 01010101) = 00101111$$

$$(ii) T(00001111, 00110011, 01010101) = 10100110$$

$$(iii) T(11110000, 11001100, 10101010) = 11110010$$

28. Using Truth table show that $(x + y) \cdot (x + z) = x + (y \cdot z)$.

[Hint : Show that Truth Value of LHS and RHS are same]

29. Using truth table verify

$$(i) \text{idempotent law} \quad (ii) \text{D'Morgan's law} \quad (iii) x(x + y)(x + xy) = x$$

$$(iv) xy' + z + (x' + y)z' = 1 \quad (v) yz + x'y'z' + xy = y$$

30. Find the Boolean expression in CNF from the following truth table:

| x | y | $f(x, y)$ |
|---|---|-----------|
| 0 | 0 | 1 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

31. Show that the set $S = \{[0], [3]\}$ is a Boolean Ring with respect to $+$ and \cdot defined by $[a] + [b] = [a + b]$ and $[a] \cdot [b] = [ab]$ where $[x] = \{a \in z : x \equiv a \pmod{6}\}$.

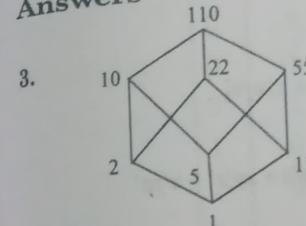
32. Test whether the Power set of a non null set S is a Boolean Ring w.r.t the compositions \cup , \cap and ' (complement).

[Hint : In fact this is not a Ring at all]

33. In a Boolean Ring prove that every element is its own inverse under addition $+$.

[Hint : Follows from a theorem]

Answers



$$7. a \cdot 0 = 0$$

$$14. (i) a \cdot (a' + b) = a \cdot b$$

$$(ii) (a + 0) + (1 \cdot a') = 1$$

$$(iii) a \cdot 0 + a \cdot 1 = a$$

$$(iv) ab + bc = (a + c)b$$

$$(v) (x' + y' + z)(x' + y + z')$$

$$(vi) (x' + z)\{x + y + y(y' + z)\}$$

$$(vii) (y + z)(y' + z') + x$$

$$(viii) a \cdot b + a \cdot 0 = a \cdot (a + b) \cdot b$$

$$(ix) A + B' = 1 \text{ iff } A + B = A$$

$$15. (i) x'y' \quad (ii) C' + A'B \quad (iii) xz' + y'u$$

$$(iv) x(y + z') \quad (v) A \cdot B + A \cdot B' + A' \cdot B'$$

$$16. (x' + y).(x + z).(y + z')$$

$$17. (i) x + yz \quad (ii) b \quad (iii) x'y'z' \quad (iv) xy + y' \quad (v) a' + bc' \quad (vi) x + y$$

$$(vii) C \quad (viii) A + B \quad (ix) x + z' \quad (x) xy \quad (xi) AC + B'C \quad (xii) z'$$

$$18. A^c \cap B^c \cap C^c$$

20. (i) $(x' + y + z)(x' + y + z')(x' + y' + z)$
(ii) $(a + b + c)(a + b' + c)(a' + b' + c')$
(iii) $(x + y' + z)(x + y' + z')(x + y + z)(x' + y + z')$
(iv) $(a + b)(a + b')(a' + b)(a' + b')$
(v) $(x' + y + z)(x' + y + z')(x' + y' + z')$
(vi) $(a' + b' + c)(a + b + c)(a + b + c')(a + b' + c)$
21. (i) 0 (ii) 0
22. (i) $xy'z' + xy'z$ (ii) $x'yz + x'y'z' + xy'z + xy'z'$
(iii) $xyz + xyz' + xy'z + xy'z'$ (iv) $xy + xz' = xyz + xyz' + xy'z'$
(v) $xy' + xy'z = xy'z' + xy'z$
(vi) $xyz' + x'y' = xyz' + x'y'z + x'y'z'$
(vii) $x_1x'_2x_3 + x_1x'_2x'_3 + x_1x_2x_3 + x'_1x_2x_3 + x'_1x_2x'_3$
(viii) $x'y'z'$ (ix) $xy + z'$
23. (i) $xy'z + xy'z' + x'y'z + x'yz$ (ii) $xy'z + xy'z' + x'yz'$
24. (i) $x'y' + xy$ (ii) $x'y'z' + xy'z$
25. (i) $x'y'z + x'yz' + xy'z' + xyz'$ (ii) $x'y' + xy$
(iii) $x' + (y + z)(y' + z')$ (iv) $(a' + b)(a + c)(b + c')$
26. (i) 00001010 (ii) 00110101 (iii) 1000
27. (i) $xz + y$, (ii) $yz' + x'z' + xy'z$, (iii) $xy + y'(x + z)$
30. $(x' + y)(x + y')$ 32. no.

II. Long Question Answers

1. Prove that the set of divisors of m is a Boolean Algebra with respect to the composition $+$, \cdot and $'$, where $a + b = LCM$ of (a, b) , $a \cdot b = HCF$ of (a, b) and $a' = \frac{m}{a}$, where m is a product of distinct prime integers.

- [Hint : Show the set is bounded, distributive complemented lattice]
2. Prove that the set of divisors of (i) 99 are not Boolean algebra
(ii) 130 are Boolean algebra (operations to be introduced by you)

3. Express the expression $(a + b)(a + b')(a' + c)$ in the smallest possible number of variables.

4. Find $T(f)$, the truth table of the following Boolean function :

- (i) $f(x, y, z) = yz'x + zx'y' + zyx$ (ii) $f(x, y, z) = (yz + z')x' + (x'y)y'z'$
(iii) $f = z'yx + xy' + y$ [W.B.U.T. 2005, 2008]
(iv) $f = y'xz' + y'z'x' + yz'x$ (v) $f = z'x'y + y'x'z$
(vi) $f = y(x' + z) + y'x$

5. Find the Disjunctive Normal form of the Boolean function having truth table :

| x | y | z | f |
|-----|-----|-----|-----|
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

6. Find the Boolean expression in CNF from the following truth table:

| (i) | x | y | z | f |
|-----|-----|-----|-----|-----|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 1 |
| 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 0 | 0 |

(ii) $T(11110000, 11001100, 10101010) = 11101001$

Answers

3. ac

4. (i) 01000101 (ii) 10101010

(iii) 00111111 (iv) 10001010 (v) 01000000 (vi) 10111100

$$5. \ xyz + xyz' + xy'z + x'y'z + x'yz' . \quad 6. \text{ (i)} \ (x' + y' + z')(x + y + z)$$

III. Multiple Choice Questions

1. Number of operations required in a Boolean Algebra is

2. Number of Unary operation required in a Boolean Algebra is

- (a) 1
 - (b) 2
 - (c) 3
 - (d) none