

Trust

What it is and how to get it

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Trust

“An entity can be trusted if it always behaves in the expected manner for the intended purpose ?”

Properties

- ▶ Unambiguous identification
- ▶ Unimpeded operation
- ▶ First-hand observation of good behavior *or* indirect experience of good behavior by a trusted third party

- ▶ *Strong Identification* — An unambiguous, immutable identifier associated with the platform. The identifier is a protected encryption key in the TXT implementation.
- ▶ *Reporting Configuration* — An unambiguous identification mechanism for software and hardware running on the platform. The mechanism is hashing in the TXT implementation

$T^x[y]$ is an homogeneous relation over actors that is true when x *trusts* y . $T^x[y]$ is a preorer:

- ▶ Reflexive - $\forall x \cdot T^x[x]$
- ▶ Transitive - $\forall x, y, z \cdot T^y[x] \wedge T^z[y] \Rightarrow T^z[x]$

The transitive property defines *chains of trust*.

A *root of trust* provides a basis for transitively building trust. Roots of trust are trusted implicitly.

There are three important Roots of Trust:

- ▶ Root of Trust for Measurement (RTM)
- ▶ Root of Trust for Reporting (RTR)
- ▶ Root of Trust for Storage (RTS)

Root of Trust for Measurement

A *Root of Trust for Measurement* is trusted to take the base system measurement.

- ▶ Typically a hash function called on an initial code base from a protected execution environment
- ▶ In the Intel TXT process the RTM is SENTER implemented on the processor
- ▶ SENTER resets and measures SINIT into PCR 18 in the TPM¹
- ▶ SENTER runs securely on the Intel processor ²
- ▶ Access control prevents access to PCR 18 by other instructions

¹'Measures X into PCR Y' informally means extending PCR Y with a hash of X.

²AMD has a similar command. ARM does not.

A *Root of Trust for Reporting* is trusted to authenticate the base system report or quote

- ▶ Typically a protected key used for authenticating reports
- ▶ In the Intel TXT processes this is the TPM's Endorsement Key (EK)
- ▶ Created and bound to its platform by the TPM foundry
- ▶ EK^{-1} is stored in the TPM and cannot be accessed by any entity other than the TPM
- ▶ EK is available for encrypting data for the TPM
- ▶ EK^{-1} is used for decrypting data
- ▶ Binding of EK to its platform is maintained by a trusted Certificate Authority (CA)

A *Root of Trust for Reporting* is trusted to protect the base stored data

- ▶ Typically a key stored in a trusted location
- ▶ In the Intel TXT boot process this is the TPM's Storage Root Key (SRK)
- ▶ Created by TPM_TakeOwnership
- ▶ SRK^{-1} is stored in the TPM
- ▶ SRK^{-1} cannot be accessed directly and can only be used by the TPM

- ▶ Trust is transitive
 - ▶ $T^x[y] \wedge T^y[z] \Rightarrow T^x[z]$
 - ▶ Construct chains of trust
 - ▶ Remember “directly observed or indirectly observed by a trusted third party”
- ▶ Roots of Trust define the basis for trust
 - ▶ Use Roots of Trust to establish base for chain
 - ▶ RTM generates a trusted first measurement
 - ▶ RTS protects first measurement
 - ▶ rtr signs base quote for appraiser

1. Power On

- ▶ PCRs reset to -1

2. SENTER is called (RTM)

- ▶ SENTER resets PCRs to 0
- ▶ SENTER hashes SINIT into PCR 18 (RTS)

3. SINIT is called (Trusted)

- ▶ Review access control modeling objectives
 - ▶ modeling platform MAC
 - ▶ modeling local access control
- ▶ Overview access control policy definition
 - ▶ design and modeling assumptions
 - ▶ platform boot policy definition
 - ▶ local policy definitions
- ▶ Overview models
 - ▶ domain and system models
 - ▶ communication model
 - ▶ theorems and status
- ▶ Identify next steps
 - ▶ runtime and moving beyond the SVP line
 - ▶ adding M&A detail

Access Control Modeling Objectives

What we're about here

Reporting joint work with Geoffrey Brown, Indiana University (submitted) in which we verify two physical layer protocols.

- ▶ Biphase Mark Protocol (BMP)
- ▶ 8N1 Protocol

These protocols are used in data transmission for CDs, Ethernet, and Tokenring, etc. as well as UARTs.

- ▶ Correctness is reasonably difficult to prove due to many real-time constraints.
- ▶ Many previous formal modeling/verification efforts for these protocols.

Some normal text goes here
just for introduction

- ▶ Appraisal
- ▶ Measurement
- ▶ Attestation
- ▶ vTPM

Why is this column getting
higher?

Maybe it's not
Center alignment seems best.

I like this for two column test
and graphics

Getting higher???

Big Picture

Armor Architecture

Introduction to \LaTeX

Beamer is a \LaTeX class for creating presentations that are held using a projector...”

This is a definition

Not really a proof.

1. This is a step



Not really a proof.

1. This is a step
2. This is another step



Not really a proof.

1. This is a step
2. This is another step
3. This is a third step
4. This is a third step
5. This is a third step
6. This is a third step



- ▶ Item 1 followed by a pause

- ▶ Item 1 followed by a pause
- ▶ Item 3 followed by a pause

- ▶ Item 1 followed by a pause
- ▶ Item 2 followed by a pause
- ▶ Item 3 followed by a pause

- ▶ BMP has been verified in PVS twice and required
 - ▶ 37 invariants and 4000 individual proof directives (initially) in the one effort
 - ▶ 5 hours just to *check* the proofs in the other effort
 - ▶ A formal specification and verification of an independent real-time model in both efforts
- ▶ BMP has been verified in (the precursor to) ACL2 by J. Moore and required
 - ▶ A significant conceptual effort to fit the problem in the logic, arguably omitting some salient features of the model
 - ▶ The statement and proof of many antecedent results
 - ▶ J. Moore reports this as one of his “best ideas” in his career

Not Your Father's Theorem-Prover

The verifications are carried out in the SAL infinite-state bounded model-checker that combines SAT-solving and SMT decision procedures to *prove* safety properties about infinite-state models.

- ▶ Theorem-proving efforts took multiple engineer-months if not years to complete.
- ▶ Our initial effort in SAL consumed about *two engineer-days*.
...and we found a significant bug in a UART application note.

Parameterized Timing Constraints

SMT allows for *parameterized* proofs of correctness. The following are example constraints from the BMP verification:

TIME: TYPE = REAL;

TPERIOD: TIME = 16;

TSAMPLE: INTEGER = 23;

TSETTLE: {x: TIME |
 0 ≤ x
 AND (x + TPERIOD < TSAMPLE)
 AND (x + TSAMPLE + 1 < 2 * TPERIOD)};

TSTABLE: TIME = TPERIOD - TSETTLE;

ERROR: {x: TIME |
 (0 ≤ x)
 AND (TPERIOD + TSETTLE < TSAMPLE*(1-x))
 AND (TSAMPLE*(1+x) + (1+x) + TSETTLE < 2 * TPERIOD)};

RSAMPMAX: TIME = TSAMPLE * (1 + ERROR);

RSAMPMIN: TIME = TSAMPLE * (1 - ERROR);

RSCANMAX: TIME = 1 + ERROR;

RSCANMIN: TIME = 1 - ERROR;

- ▶ Parser
- ▶ Simulator
- ▶ Symbolic model-checker (BDDs)
- ▶ Witness symbolic model-checker
- ▶ Bounded model-checker
- ▶ Infinite-state bounded model-checker
- ▶ Future releases include:
 - ▶ Explicit-state model-checker
 - ▶ MDD-based symbolic model-checking

All of which are “state-of-the-art”

Please direct your attention to the whiteboard.

Timeout Automata³ (Semantics)

An *explicit* real-time model.

- ▶ Vocabulary:
 - ▶ A set of state variables.
 - ▶ A *global clock*, $c \in \mathbb{R}^{0\leq}$.
 - ▶ A set of *timeout* variables T such that for $t \in T$, $t \in \mathbb{R}^{0\leq}$.
- ▶ Construct a transition system $\langle S, S^0, \rightarrow \rangle$:
 - ▶ States are mappings of all variables to values.
 - ▶ Transitions are either *time transitions* or *discrete transitions*.
 - ▶ Time transitions are enabled if the clock is less than all timeouts. Updates clock to least timeout.
 - ▶ Discrete transitions are enabled if the clock equals some timeout. Updates state variables and timeouts.

³B. Dutertre and M. Sorea. Timed systems in SAL. *SRI TR*, 2004.

Disjunctive Invariants

Even with k -induction, getting a sufficiently strong invariant is still hard! *Disjunctive invariants* help. A disjunctive invariant can be built iteratively from the counterexamples returned for the hypothesized invariant being verified.

```
t0: THEOREM system |-  
    G( ( (phase = Settle)  
        AND (rstate = tstate + 1)  
        AND (rclk - tclk - TPERIOD > 0)  
        AND (tclk + TPERIOD + TSTABLE - rclk > 0))  
    OR  
      ( (phase = Stable)  
        AND (rstate = tstate + 1)  
        AND (rclk - tclk - TSETTLE > 0)  
        AND (tclk + TPERIOD - rclk > 0)  
        AND (rdata = tdata))  
      .  
      .  
      .
```