Trust

What it is and how to get it

Dr. Perry Alexander

Information and Telecommunication Technology Center
Electrical Engineering and Computer Science
The University of Kansas
palexand@ku.edu



Defining Trust

Trust

"An entity can be trusted if it always behaves in the expected manner for the intended purpose [1]"



Defining Trust

Properties

- ► Unambiguous identification
- ► Unimpeded operation
- ► First-hand observation of good behavior *or* indirect experience of good behavior by a trusted third party



Necessary Capabilities for Trust

- ► Strong Identification An unambiguous, immutable identifier associated with the platform. The identifier is a protected encryption key in the TXT implementation.
- ► Reporting Configuration An unambiguous identification mechanism for software and hardware running on the platform. The mechanism is hashing in the TXT implementation



- Start with a measurer and store that is trusted
- Measure software to be launched
- ► Store the measurement
- ► Launch the new software
- ► Repeat until system boot

В

Trusted Store

Α

0

Trusted Measurer



- Start with a measurer and store that is trusted
- Measure software to be launched
- Store the measurement
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- ► Repeat until system boot

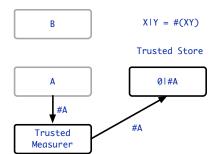


Trusted

Measurer

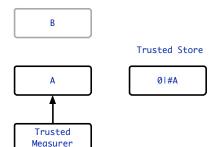


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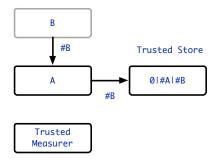


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What Can We Learn?

Assume we know 0|#A|#B is correct

- ▶ #A and #B must be correct
 - ► A and B are the correct binaries
 - ► A and B include hash and launch functions
- Measurement occurred in the right order
 - $\blacktriangleright \#(XY) \neq \#(YX)$
 - Trusted store started with 0

Chains of Trust

Trust chain starts with Trusted Measurer and Trusted Store.



Trust is a Preorder

 $T^{x}[y]$ is an homogeneous relation over actors that is true when x trusts y. $T^{x}[y]$ is a preorer:

- ► Reflexive $\forall x \cdot T^x[x]$
- ► Transitive $\forall x, y, z \cdot T^{y}[x] \wedge T^{z}[y] \Rightarrow T^{z}[x]$

The transitive property defines chains of trust.



Trusted Platform Module

The *Trusted Platform Module (TPM)* is a cryptographic coprocessor for trust.

- ► Endorsement Key (EK) factory generated asymmetric key that uniquely identifies the TPM
- ► Attestation Instance Key (AIK) TPM_CreateIdentity generated asymmetric key alias for the EK
- Storage Root Key (SRK) TPM_TakeOwnership generated asymmetric key that encrypts data associated with the TPM
- Platform Configuration Registers (PCRs) protected registers for storing and extending hashes
- ► NVRAM Non-volatile storage associated with the TPM



Endorsement Key

- Asymmetric key generated at TPM fabrication
- $ightharpoonup EK^{-1}$ is protected by the TPM
- ► EK by convention is managed by a Certificate Authority
 - ▶ Binds *EK* with a platform
 - Classic trusted third party
- Only used for encryption
- Attestation Instance Keys (AIK) are aliases for the EK
 - Used for signing
 - Authorized by the EK



Storage Root Key

- ► Asymmetric key generated by TPM_TakeOwnership
- ► *SRK*⁻¹ is protected by the TPM
- ► SRK is available for encryption
- ► Used as the root for chaining keys by wrapping
 - A wrapped key is an asymmetric key pair with it's private key sealed
 - Safe to share the entire key
 - Only usable in the presence of the wrapping key with expected PCRs



Platform Configuration Registers

Operations on PCRs

- Extension Hash a new value juxtaposed with the existing PCR value
- ► Reset Set to 0
- Set Set to a known value

Operations using PCRs

- Sealing data PCR state dependent encryption
- Wrapping keys PCR state dependent encryption of a private key
- Quote Reporting PCR values to a third party

Properties

- ► Locality Access control
- ▶ Resettable Can a PCR be reset
- Many others that we don't need yet



Roots of Trust

A *root of trust* provides a basis for transitively building trust. Roots of trust are trusted implicitly.

There are three important Roots of Trust:

- ► Root of Trust for Measurement (RTM)
- ► Root of Trust for Reporting (RTR)
- Root of Trust for Storage (RTS)



Root of Trust for Measurement

A *Root of Trust for Measurement* is trusted to take the base system measurement.

- ► A hash function called on an initial code base from a protected execution environment
- Starts the measurement process during boot
- ► In the Intel TXT process the RTM is SENTER implemented on the processor



Root of Trust for Reporting

A *Root of Trust for Reporting* is trusted to guarantee the integrity of the base system report or quote

- ► A protected key used for authenticating reports
- In the Intel TXT processes this is the TPM's Endorsement Key (EK)
- Created and bound to its platform by the TPM foundry
- ► EK⁻¹ is stored in the TPM and cannot be accessed by any entity other than the TPM
- ► EK is available for encrypting data for the TPM
- $ightharpoonup EK^{-1}$ is used for decrypting data inside the TPM
- ► Linking *EK* to its platform is done by a trusted Certificate Authority (CA)



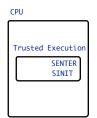
Root of Trust for Storage

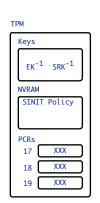
A Root of Trust for Storage is trusted to protect stored data

- ► A key stored in a protected location
- In the Intel TXT boot process this is the TPM's Storage Root Key (SRK)
- Created by TPM_TakeOwnership
- ► SRK⁻¹ is stored in the TPM and cannot be accessed by any entity other than the TPM
- SRK is available for encrypting data for the TPM
- SRK is used for protecting other keys



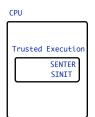
- Power-on reset, resettable
 PCRs set to -1
- SENTER called
- SENTER resets resettable PCRs to 0
- SENTER measures SINIT policy into PCR 18

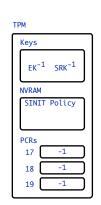






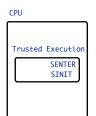
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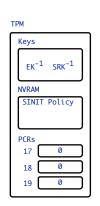






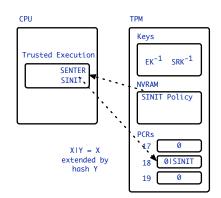
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- ▶ Power-on reset, resettable PCRs set to -1
- SENTER called
- SENTER resets resettable PCRs to 0
- SENTER measures SINIT policy into PCR 18





What We Know From Good PCR 18

A good value in PCR 18 tells us:

- ► SENTER was called Resetting PCR 18 starts measurements at 0 rather than -1
- ► SINIT was measured by SENTER Only SENTER can extend PCR 18
- SINIT uses the correct policy PCR 18 is extended with SINIT measurement policy
- ▶ SENTER ran before SINIT was measured $A \mid B \neq B \mid A$

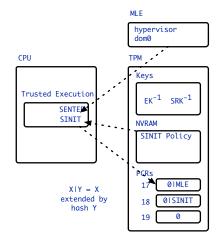
Measurement \neq Trust

Measurements must be appraised to determine trust.



Two Steps from Roots of Trust

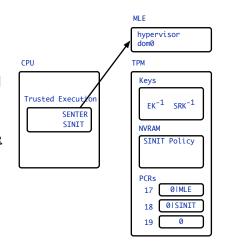
- SINIT measures the Measured Launch Environment (MLE) using measured policy
- ► SINIT returns control to SENTER
- ► SENTER invokes the MLE





Two Steps from Roots of Trust

- SINIT measures the Measured Launch Environment (MLE) using measured policy
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What We Know From Good PCRs

- ► SENTER was called Resetting PCR 18 starts measurement sequence at 0 rather than -1
- ► SINIT policy was measured by SENTER Only SENTER can extend PCR 18
- SINIT uses the correct policy PCR 18 is extended with SINIT measurement policy
- ▶ SENTER ran before SINIT $0 \mid SINIT \neq -1 \mid SINIT$
- ► SLE is good Measured by good SINIT into PCR
- ► Initial OS is good Measured by good SLE into PCR



Boot the MLE

► SENTER starts the MLE

- ► SENTER starts the hypervisor
- SENTER passes dom0 to hypervisor
- hypervisor starts dom0

dom0 constructs the Armored VP

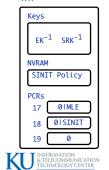
- Measures the vTPM into the TPM
- Starts the vTPM
- Measures remaining Armored VMs into the vTPM
- Starts remaining Armored VMs
- Measures Armored application into the vTPM
- Starts the Armored application



Armored VP

vTPM appraiser attester measurer application

TPM



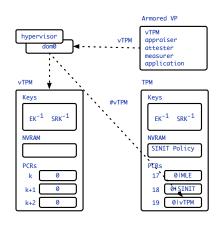
Boot the vTPM

► SENTER starts the MLE

- ► SENTER starts the hypervisor
- SENTER passes dom0 to hypervisor
- hypervisor starts dom0

dom0 constructs the Armored VP

- Measures the vTPM into the TPM
- ► Starts the vTPM
- Measures remaining Armored VMs into the vTPM
- Starts remaining Armored VMs
- Measures Armored application into the vTPM
- Starts the Armored application





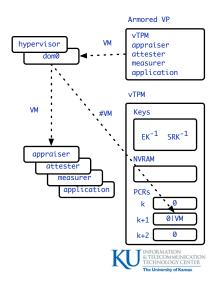
Boot the Armor Virtual Platform

► SENTER starts the MLE

- SENTER starts the hypervisor
- SENTER passes dom0 to hypervisor
- hypervisor starts dom0

dom0 constructs the Armored VP

- Measures the vTPM into the TPM
- Starts the vTPM
- Measures remaining Armored VMs into the vTPM
- Starts remaining Armored VMs
- Measures Armored application into the vTPM
- Starts the Armored application



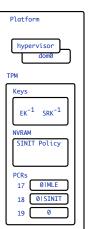
Boot Result

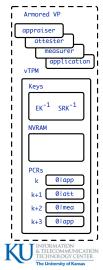
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dom0 constructs the Armored VP

- Measures the vTPM into the TPM
- Starts the vTPM
- Measures remaining Armored VMs into the vTPM
- Starts remaining Armored VMs
- Measures Armored application into the vTPM
- Starts the Armored application





Chaining Trust

► Trust is transitive

- $ightharpoonup T^{x}[y] \wedge T^{y}[z] \Rightarrow T^{x}[z]$
- Construct chains of trust
- Remember "directly observed or indirectly observed by a trusted third party"
- Roots of Trust define the "root" for trust
 - Use Roots of Trust to establish base for chain
 - RTM generates a trusted first measurement
 - ▶ RTS protects first measurement
 - ► RTR signs base quote for appraiser (eventually)
- ► Extend chains of trust by measuring before executing



[1] A. Martin et al. The ten page introduction to trusted computing. Technical Report CS-RR-08-11, Oxford University Computing Labratory, Oxford, UK, 2008.

