

Trust

What it is and how to get it

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Trust

“An entity can be trusted if it always behaves in the expected manner for the intended purpose [1]”

Properties

- ▶ Unambiguous identification
- ▶ Unimpeded operation
- ▶ First-hand observation of good behavior *or* indirect experience of good behavior by a trusted third party

Necessary Capabilities for Trust

- ▶ *Strong Identification* — An unambiguous, immutable identifier associated with the platform. The identifier is a protected encryption key in the TXT implementation.
- ▶ *Reporting Configuration* — An unambiguous identification mechanism for software and hardware running on the platform. The mechanism is hashing in the TXT implementation

Chaining Measurements

- Start with a measurer and store that is trusted

B

A

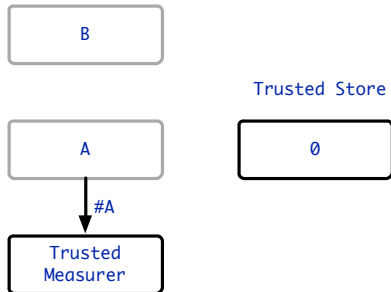
Trusted
Measurer

Trusted Store

0

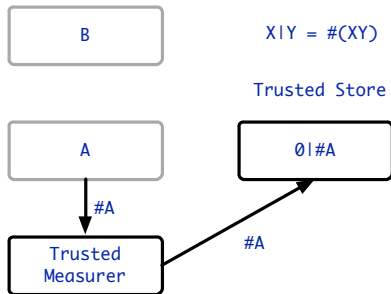
Chaining Measurements

- ▶ Start with a measurer and store that is trusted
- ▶ Measure software to be launched



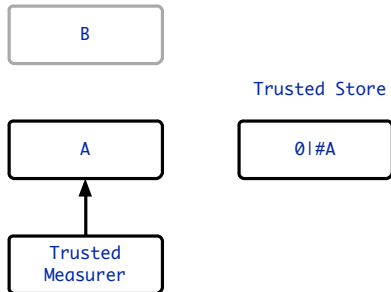
Chaining Measurements

- ▶ Start with a measurer and store that is trusted
- ▶ Measure software to be launched
- ▶ Store the measurement



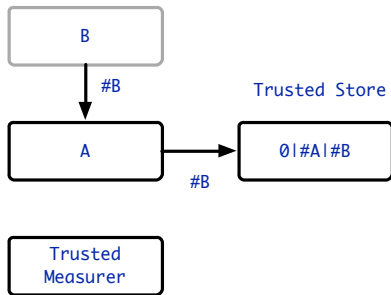
Chaining Measurements

- ▶ Start with a measurer and store that is trusted
- ▶ Measure software to be launched
- ▶ Store the measurement
- ▶ Launch the new software



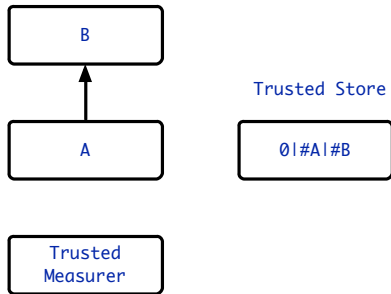
Chaining Measurements

- ▶ Start with a measurer and store that is trusted
- ▶ Measure software to be launched
- ▶ Store the measurement
- ▶ Launch the new software
- ▶ Repeat until system boot



Chaining Measurements

- ▶ Start with a measurer and store that is trusted
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What Do We Know?

Assume we know $0|\#A|\#B$ is correct

- ▶ $\#A$ and $\#B$ must be correct
 - ▶ A and B are the correct binaries
 - ▶ A and B include hash and launch functions
- ▶ Measurement occurred in the right order
 - ▶ $\#(XY) \neq \#(YX)$
 - ▶ Trusted store started with 0

But Why Trust B?

A chain exists from Trusted Measurer and Trusted Store to B

- ▶ Trusted Measurer and Trusted Store are trusted *a priori*
- ▶ A is trusted to be A because its measurement is:
 - ▶ Correct
 - ▶ Taken by a trusted party (Trusted Measurer)
 - ▶ Stored by a trusted party (Trusted Store)
- ▶ B is trusted to be B because its measurement is:
 - ▶ Correct
 - ▶ Taken by a trusted party (A)
 - ▶ Stored by a trusted party (Trusted Store)
 - ▶ If A's ability to measure B were compromised, #A would be wrong
- ▶ and so on and so on...

$T^x[y]$ is an homogeneous relation over actors that is true when x *trusts* y . $T^x[y]$ is a preorer:

- ▶ Reflexive - $\forall x \cdot T^x[x]$
- ▶ Transitive - $\forall x, y, z \cdot T^x[y] \wedge T^y[z] \Rightarrow T^z[x]$

Trusted Boot builds evidence supporting these chains

The *Trusted Platform Module (TPM)* is a cryptographic coprocessor for trust.

- ▶ Endorsement Key (EK) — factory generated asymmetric key that uniquely identifies the TPM
- ▶ Attestation Instance Key (AIK) — TPM_CreateIdentity generated asymmetric key alias for the EK
- ▶ Storage Root Key (SRK) — TPM_TakeOwnership generated asymmetric key that encrypts data associated with the TPM
- ▶ Platform Configuration Registers (PCRs) — protected registers for storing and extending hashes
- ▶ NVRAM — Non-volatile storage associated with the TPM

- ▶ Asymmetric key generated at TPM fabrication
- ▶ EK^{-1} is protected by the TPM
- ▶ EK by convention is managed by a Certificate Authority
 - ▶ Binds EK with a platform
 - ▶ Classic trusted third party
- ▶ Only used for encryption
- ▶ Attestation Instance Keys (AIK) are aliases for the EK
 - ▶ Used for signing
 - ▶ Authorized by the EK

- ▶ Asymmetric key generated by TPM_TakeOwnership
- ▶ SRK^{-1} is protected by the TPM
- ▶ SRK is available for encryption
- ▶ Used as the root for chaining keys by *wrapping*
 - ▶ A wrapped key is an asymmetric key pair with its private key sealed
 - ▶ Safe to share the entire key
 - ▶ Only usable in the presence of the wrapping key with expected PCRs

Platform Configuration Registers

► Operations on PCRs

- Extension — Hash a new value juxtaposed with the existing PCR value
- Reset — Set to 0
- Set — Set to a known value

► Operations using PCRs

- Sealing data — PCR state dependent encryption
- Wrapping keys — PCR state dependent encryption of a private key
- Quote — Reporting PCR values to a third party

► Properties

- Locality — Access control
- Resettable — Can a PCR be reset
- Many others that we don't need yet

A *root of trust* provides a basis for transitively building trust. Roots of trust are trusted implicitly.

There are three important Roots of Trust:

- ▶ Root of Trust for Measurement (RTM)
- ▶ Root of Trust for Reporting (RTR)
- ▶ Root of Trust for Storage (RTS)

Root of Trust for Measurement

A *Root of Trust for Measurement* is trusted to take the base system measurement.

- ▶ A hash function called on an initial code base from a protected execution environment
- ▶ Starts the measurement process during boot
- ▶ In the Intel TXT process the RTM is SENTER implemented on the processor

Root of Trust for Reporting

A *Root of Trust for Reporting* is trusted to guarantee the integrity of the base system report or quote

- ▶ A protected key used for authenticating reports
- ▶ In the Intel TXT processes this is the TPM's Endorsement Key (EK)
- ▶ Created and bound to its platform by the TPM foundry
- ▶ EK^{-1} is stored in the TPM and cannot be accessed by any entity other than the TPM
- ▶ EK is available for encrypting data for the TPM
- ▶ EK^{-1} is used for decrypting data inside the TPM
- ▶ Linking EK to its platform is done by a trusted Certificate Authority (CA)

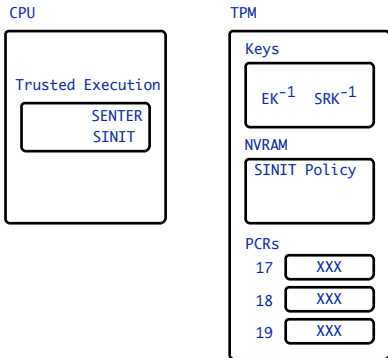
A *Root of Trust for Storage* is trusted to protect stored data

- ▶ A key stored in a protected location
- ▶ In the Intel TXT boot process this is the TPM's Storage Root Key (SRK)
- ▶ Created by TPM_TakeOwnership
- ▶ SRK^{-1} is stored in the TPM and cannot be accessed by any entity other than the TPM
- ▶ SRK is available for encrypting data for the TPM
- ▶ SRK is used for protecting other keys

One Step from Roots of Trust

Roots of trust are used to build a trusted system from boot.

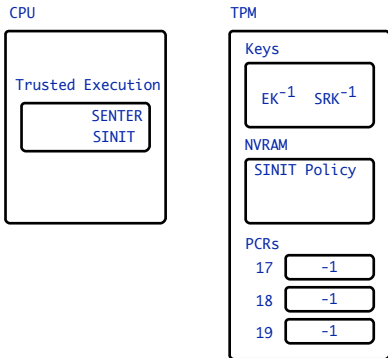
- Power-on reset



One Step from Roots of Trust

Roots of trust are used to build a trusted system from boot.

- ▶ Power-on reset
- ▶ Resettable PCRs set to -1

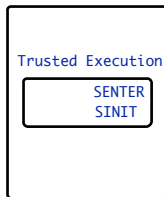


One Step from Roots of Trust

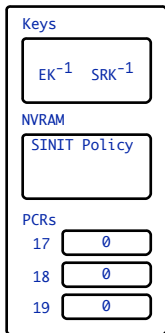
Roots of trust are used to build a trusted system from boot.

- ▶ Power-on reset
- ▶ Resettable PCRs set to -1
- ▶ SENTER called, resets resettable PCRs to 0

CPU



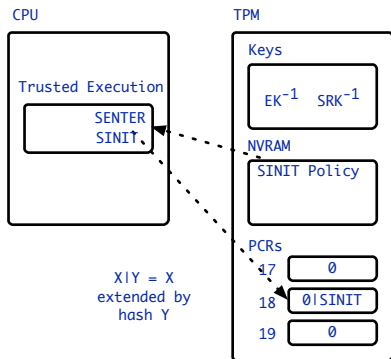
TPM



One Step from Roots of Trust

Roots of trust are used to build a trusted system from boot.

- ▶ Power-on reset
- ▶ Resettable PCRs set to -1
- ▶ SENTER called, resets resettable PCRs to 0
- ▶ SENTER measures SINIT policy into PCR 18



What We Know From Good PCR 18

A good value in PCR 18 tells us:

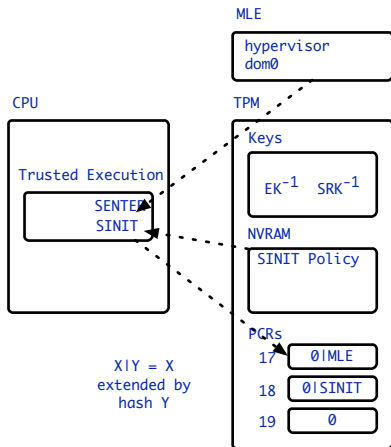
- ▶ SENTER was called — Resetting PCR 18 starts measurements at 0 rather than -1
- ▶ SINIT was measured by SENTER — Only SENTER can extend PCR 18
- ▶ SINIT uses the correct policy — PCR 18 is extended with SINIT measurement policy
- ▶ SENTER ran before SINIT was measured — $A \mid B \neq B \mid A$

Measurement \neq Trust

Measurements must be appraised to determine trust.

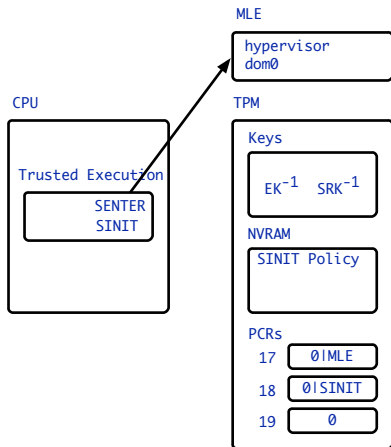
Two Steps from Roots of Trust

- ▶ SINIT measures the Measured Launch Environment (MLE) using measured policy
- ▶ SINIT returns control to SENTER



Two Steps from Roots of Trust

- ▶ SINIT measures the Measured Launch Environment (MLE) using measured policy
- ▶ SINIT returns control to SENTER
- ▶ SENTER invokes the MLE



What We Know From Good PCRs

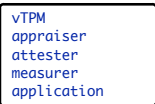
- ▶ SENTER was called — Resetting PCR 18 starts measurement sequence at 0 rather than -1
- ▶ SINIT policy was measured by SENTER — Only SENTER can extend PCR 18
- ▶ SINIT uses the correct policy — PCR 18 is extended with SINIT measurement policy
- ▶ SENTER ran before SINIT — $0 \mid SINIT \neq -1 \mid SINIT$
- ▶ SLE is good — Measured by good SINIT into PCR
- ▶ Initial OS is good — Measured by good SLE into PCR

Boot the MLE

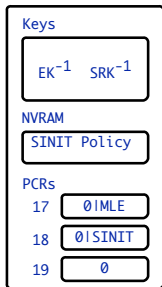
- ▶ **SENDER starts the MLE**
 - ▶ SENTER starts the hypervisor
 - ▶ SENTER passes dom0 to hypervisor
 - ▶ hypervisor starts dom0
- ▶ **dom0 constructs the Armored VP**
 - ▶ Measures the vTPM into the TPM
 - ▶ Starts the vTPM
 - ▶ Measures remaining Armored VMs into the vTPM
 - ▶ Starts remaining Armored VMs
 - ▶ Measures Armored application into the vTPM
 - ▶ Starts the Armored application



Armored VP

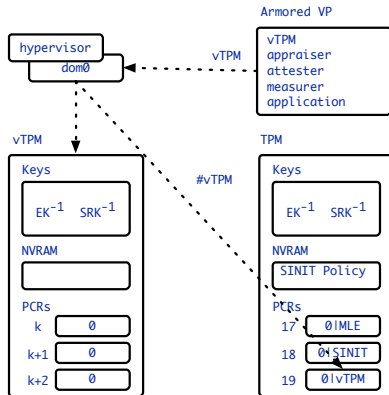


TPM



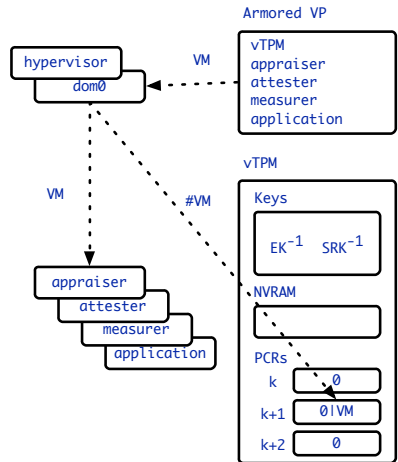
Boot the MLE

- ▶ SENTER starts the MLE
 - ▶ SENTER starts the hypervisor
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 - ▶ Measures Armored application into the vTPM
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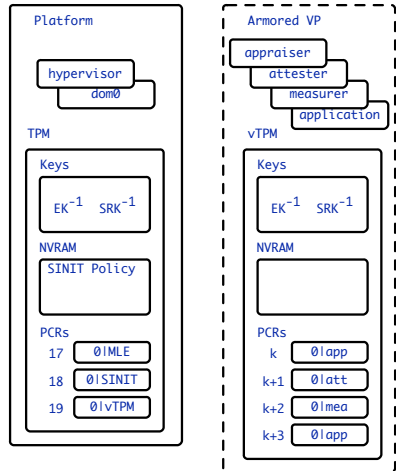


Boot the MLE

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 - ▶ Starts remaining Armored VMs
 - ▶ Measures Armored application into the vTPM
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 - ▶ Starts the vTPM
 - ▶ Measures remaining Armored VMs into the vTPM
 - ▶ Starts remaining Armored VMs
 - ▶ Measures Armored application into the vTPM
 - ▶ Starts the Armored application



Chaining Trust (Reprise)

- ▶ Trust is transitive
 - ▶ $T^x[y] \wedge T^y[z] \Rightarrow T^x[z]$
 - ▶ Construct evidence trust chains
 - ▶ Remember “directly observed or indirectly observed by a trusted third party”
- ▶ Roots of Trust define the “root” for trust
 - ▶ Use Roots of Trust to establish base for chain
 - ▶ RTM is the Trusted Measurer
 - ▶ RTS is the Trusted Store
 - ▶ RTR is the Trusted Reporter (coming soon...)
- ▶ Extend chains of trust by measuring before executing

- [1] A. Martin et al. The ten page introduction to trusted computing. Technical Report CS-RR-08-11, Oxford University Computing Laboratory, Oxford, UK, 2008.