### Remote Attestation for Cloud-Based Systems

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#### Clouds and Trust

- ► The promises of the cloud are substantial
  - reduced hardware and software costs
  - reduced resource consumption
  - improved availability and reliability
- ► The structure of the cloud complicates assurance
  - ▶ not under the desk
  - ambiguous and changing runtime environment
  - unknown and unknowable actors in the same environment
- ▶ Is trust possible in the cloud environment?
  - ▶ unambiguous identification
  - confirmation of uninhibited execution
  - direct or trusted indirect observation of good behavior

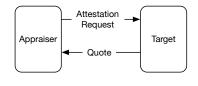
### Virtual Blinking Lights

Provide new capabilities that establish and maintain trustworthy cloud-based application deployment

- Establish trust in cloud applications
  - ► trust in cloud infrastructure
  - ► trust in user-space applications
  - ► trust in application cohorts
- Promote informed decision making
  - confirm data confidentiality
  - confirm execution and data integrity
- ► Autonomous run-time response and reconfiguration
  - respond to attack, failure, reconfiguration, and repair
  - appraisal informs response

### Semantic Remote Attestation

- ► Appraiser requests a quote
  - specifies needed information
  - provides a nonce
- ► Target gathers evidence
  - measures application
  - gathers evidence of trust
- ► Target generates quote
  - measurements and evidence
  - original nonce
  - cryptographic signature
- ► Appraiser assesses quote
  - ► good application behavior
  - ► infrastructure trustworthiness

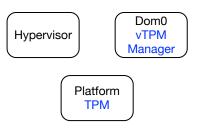


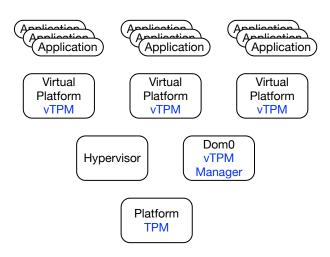
#### Trusted Platform Module

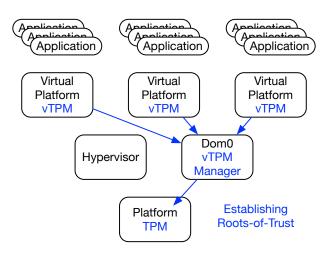
- Provides and Protects Roots of Trust
  - Storage Root Key (SRK) root of trust for storage
  - ► Endorsement Key (EK) root of trust for reporting
- Quote generation
  - ▶ high integrity quotes  $(\{|RS|\}_{AIK^-}, SML, \{|n, PCRComp|\}_{AIK^-})$
  - ▶ high integrity evidence ( $\langle E, n \rangle$ , {|| $\langle E, n \rangle$ |, PCR}<sub>AIK</sub>-
- ► Sealing data to state
  - $\{D, PCR\}_{K^+}$  will not decrypt unless PCR = current PCR
  - data is safe even in the presence of malicious machine
- Binding data to TPMs and machines
  - $(\{K^-\}_{SRK^+}, K)$   $\{D\}_{K^+}$  cannot be decrypted unless  $SRK^-$  is installed
  - $(\{J^-\}_{K^+}, \mathsf{J})$   $\{D\}_{J^+}$  cannot be decrypted unless  $K^-$  and  $SRK^-$  are installed

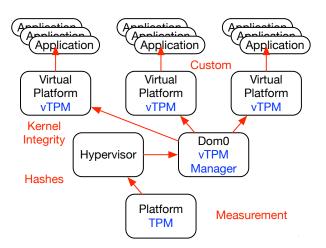
Chasing the bottom turtle

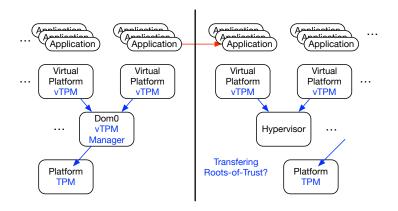
Platform TPM











## **Enabling Technologies**

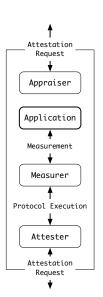
- ► Trustworthy protocol execution
  - executable and analyzable protocol representation
  - generates evidence of trustworthiness
  - negotiates attestation details
  - designed for highly focused appraisal
- Application specific measurement
  - managed and traditional execution environments
  - ► compile-time assistance for measurer synthesis
  - specialized measurement bundled with applications
- ► Lightweight trust infrastructure
  - abstract communications capability
  - migration support
  - strong identity

## Armored Application Architecture

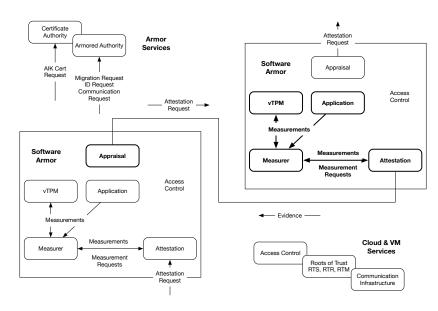
M&A targeted to an application

- Appraiser makes attestation requests
- ► Attester responds to attestation requests
- ► Measurer gathers evidence from application
- ► Influenced by the *Trusted Research*Platform and Principles of Remote

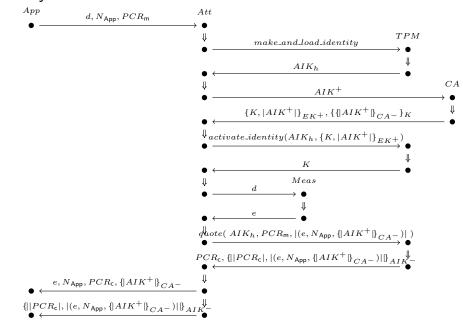
  Attestation



## System-Level Architecture



### Privacy CA Attestation



#### **EDSL** for Protocol

#### First-class protocol structures

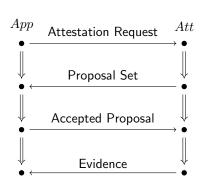
- ► First-class structure for protocols
  - encapsulates a protocol-centered computation
  - semantics provide a basis for static analysis
  - based loosely on the Reader monad
- ► Abstract communication primitives
  - ► extended RPC-style capability
  - requests remote execution
  - defines send and receive operations
  - abstracts away communication details

```
do {
    f(x);
    y <- f(x);
    send a x;
    y <- receive a
}</pre>
```

# Negotiating a Protocol

Respecting privacy

- ▶ Typical negotiation
  - request sent to Attester
  - Attester generates proposal
  - ► Appraiser selects protocol
  - Attester executes protocol
- ► Three kinds of requests
  - ▶ execute protocol 22
  - provide {OS\_config, http\_stat, firewall\_stat}
  - ► execute protocol do { ... }
- ► Three negotiation criteria
  - ► ability to satisfy the request
  - satisfaction of appraiser and attester privacy policies
  - previously obtained evidence



# **Negotiation Protocol**

Request and Select

- Requests an attestation
- Receives proposals
- ► Selects from proposals

Negotiation is a protocol that can itself be selected or negotiated

## Negotiation Results

- ► Evidence and Protocol pairs
- Satisfies privacy policy of attester
- ► Provide some or all of requested information

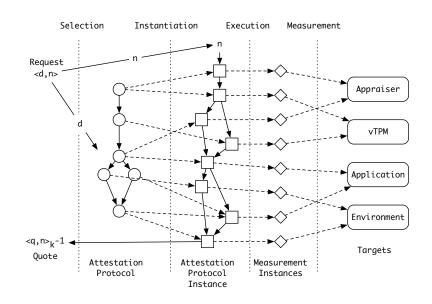
```
((ID,SIGHASH,SIGSRC),
  do { id <- getVCID;
    sig <- getSigFileEvidence;
    src <- getSigFileSrc;
    e <- createEvidence(id,sig,src);
    returnEvidence(e) })</pre>
```

#### Reified Protocol

Generated negotiation protocol code (currently by hand):

```
P = CreateChannel (AChannel "attesterChan") Target
    $ Send ANRequest (AChannel "attesterChan")
    $ Receive (Var "counterOffer") (AChannel "attesterChan")
    $ CalculateFinalRequest (Var "finalReq")
                            ANRequest
                            (Var "counterOffer")
    $ Send (Var "finalReq") (AChannel "attesterChan")
    $ Receive (Var "finalConfirmation")
              (AChannel "attesterChan")
    $ Case (Var "finalConfirmation") [(Var "finalReq")]
           (HandleFinalChoice (Var "result") (Var "finalReq")
           (Result (Var "result")))
           (Stuck "finalConf and finalReq match error")
```

### Performing Measurement and Attestation



## Single Realm Attestation

Protocol for gathering virus checker evidence

```
do { id <- getVCID;
    sig <- getSigFileEvidence;
    src <- getSigFileSrc;
    e <- createEvidence(id,sig,src);
    returnEvidence(e) }</pre>
```

and generates evidence of the form:

```
\langle (id, sig, src), \{ | (id, sig, src)|, PCRComp_0 \}_{AIK_0^-} \rangle
```

Appraisal replays the protocol up to crypto operations with known good measurements

### Multi-Realm Attestation

Nested attestation requests evidence from the signature server directly:

```
do { id <- getVCID;
    sig <- getSigFileEvidence;
    src <- getSigFileSrc;
    srcEvidence <- send src r;
    e <- createEvidence(id,sig,src,srcEvidence)
    returnEvidence(e)
}</pre>
```

and generates bundled evidence:

```
\begin{split} \text{let} & \ b \ = \langle (e), \{ |e|, PCRComp_1 \}_{AIK_1^-} \rangle \text{ in } \\ & \ \langle (id, sig, src, b), \{ |(id, sig, src, b)|, PCRComp_0 \}_{AIK_0^-} \rangle \end{split}
```

### Trusting Evidence

Why bundling is hard

- ► Trusting evidence
  - hashes and TPM quotes
  - measure and appraise the attestation infrastructure
  - ► gather evidence of good protocol execution
- Trusting bundled evidence
  - ▶ appraisers do not know the source of evidence a priori
  - ▶ no global name space for evidence sources
  - ► bundled appraisals vs bundled evidence
- ► Trusting the appraiser
  - negotiated protocols must satisfy privacy policies
  - trust may not be transitive for applications and infrastructure
  - global policy is not an answer

#### Current Status

#### Demos available

- Attestation and Appraisal development
  - CA-Based attestation protocol execution example
  - ► simple dynamic appraisal of attestation results
  - integrated negotiation protocol and attestation protocols
- Measurement development
  - ► HotSpot-based Java VM run time measurements
  - detect and report several runtime anomalies
  - standard mechanism for extending measurement capabilities
- ► Infrastructure development
  - vchan, TCP/IP and socket communication infrastructure
  - initial certificate authority implementation
  - ► language-based interface with TPM 1.2
  - ▶ integrated Berlios TPM emulator
  - ► JSON-based data exchange formats

### Ongoing Work

Goals for 2015

- Establish roots-of-trust and trust argument
  - ▶ measured launch and remeasurement of ArmoredSoftware
  - establish trust in the Xen/OpenStack infrastructure
- ► Executable protocol representation and protocol semantics
  - evidence of proper execution
  - static trust analysis
  - protocol-centered appraisal
- ► More capable measurement
  - compiler directed measurement
  - ► continuous measurement—tripping and trending
- ► Publicly available libraries and infrastructure

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