ArmoredSoftware: Trust in the cloud

Annual Demonstration

Dr. Perry Alexander, Dr. Andrew Gill, Dr. Prasad Kulkarni, Adam Petz, Paul Kline, Justin Dawson, Jason Gevargizian, Leon Searl, Edward Komp

Information and Telecommunication Technology Center Electrical Engineering and Computer Science The University of Kansas palexand@ku.edu,andygill@ku.edu,prasadk@ku.edu

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Outline

Introduction and Project Goals
Big Picture
Implementation

Prototype demonstration and discussion
Refine big picture to current demo
Protocol Execution
Appraisal
Attestation Protocol Execution
Measurement
Communication

Short term goals and milestones

Questions and guidance



Program Goals

Virtual Blinking Lights

Trust in the Cloud

Provide new capabilities that establish and maintain trustworthy cloud-based application deployment

- Establish trust among cloud components
 - trust among cohorts of processes
 - trust among processes and environment
- Promote informed decision making
 - data confidentiality can be confirmed
 - execution and data integrity can be confirmed
- ► Autonomous run-time response and reconfiguration
 - responds to attack, failure, reconfiguration, and repair
 - response varies based on measurement



Delivery Platform

Open source, standards compliant

- Lightweight integration with existing cloud infrastructure
 - OpenStack cloud infrastructure
 - ➤ Xen+XSM VM infrastructure
 - ► Fedora, HotSpot JVM, GHC
- Trusted Computing Group standards compliant
 - ► Trusted Platform Module 1.2
 - ▶ TCG vTPM (in principle)
 - ▶ Trusted OS infrastructure
- ► Standard communication mechanisms
 - JSON structures for all exchanged data
 - vchan for on-platform communication
 - ► TCP/IP for off-platform communication



New Technologies

► Trustworthy protocol execution

- executable protocol representation
- protocol execution generates evidence of trustworthiness
- highly focused protocols
- strand space formal semantics

► Application specific measurement

- managed and traditional execution environments
- compile-time assistance for measurer synthesis
- specialized measurement bundled with applications

Attestation driven cloud application and data management

- health monitoring
- problem mitigation
- application migration
- access control



Research & Development Plan

► Development and integrate measurement capabilities

- hosted languages (Java)
- ▶ traditional compiled languages (C, C++)
- ▶ integrate with environment measurers (Xen,OpenStack,OS)

Develop attestation capabilities

- flexible, user configurable protocol representation
- measured protocol execution
- protocol execution appraisal

Develop infrastructure trust argument

- develop lightweight vTPM infrastructure supporting mobility
- launch from known roots of trust
- maintain trust evidence at run time
- maintain trust over migration



Research & Development Plan

Automated synthesis and verification

- measurer synthesis at application compile time
- automated evidence appraisal from protocols
- formal trust argument

▶ Demonstrations

- initial simple infrastructure demonstrations
- cloud-based "big data" environment demonstration
- federated trust demonstration
- demonstrations as discovered/directed

► Scale up and roll out

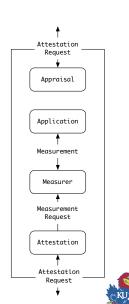
- integration with Xen, OpenStack, Linux
- installation management and packaging
- effective web presence



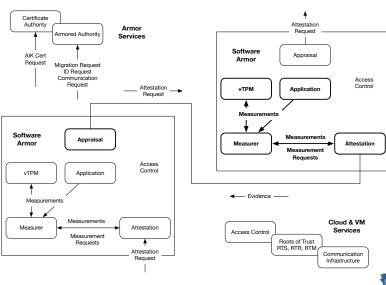
Armored Application Architecture

- ► Focus is user-space applications
- Assesses the cloud infrastructure and environment
- ► Attests to the state of its application
- High-assurance, lightweight infrastructure
- ► Influenced by the *Trusted Research*Platform and Principles of Remote

 Attestation



System-Level Architecture



What We Are Demonstrating

Execution of a CA-based Attestation Protocol

- Attestation request
- ► Protocol execution
- ► Evidence appraisal

Major architectural subsystems

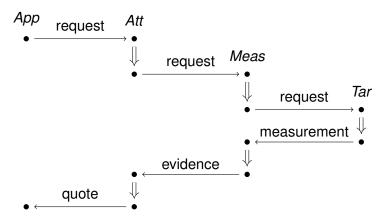
- Appraiser
- Attestation Manager
- ▶ Measurer
- Instrumented JVM
- vTPM and Certificate Authority

► Anomaly Detection

- Bad signatures and PCRs
- Bad CA certificates
- Bad quotes and AIKs
- Bad measurements

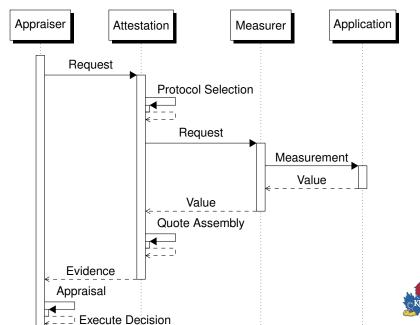


Abstract CA-Based Attestation Protocol





Abstract CA-Based Attestation Protocol





Message List Representation

 $App \rightarrow Att : d, N_{App}, PCR_{m} \text{ on } C_{AppAtt}$

 $Att \rightarrow TPM : make_and_load_identity on C_{AttTPM}$

 $TPM \rightarrow Att : AIK^+, AIK_h \text{ on } C_{TPMAtt}$

 $Att \rightarrow CA : Att, AIK^+ \text{ on } C_{AttCA}$

 $CA \rightarrow Att: \{K, |AIK|\}_{EK^+}, \{[AIK^+]_{CA^-}\}_{K^+} \text{ on } C_{CAAtt}$ $Att \rightarrow TPM: activate_identity(AIK_h, |AIK|) \text{ on } C_{AttTPM}$

 $TPM \rightarrow Att : K \text{ on } C_{TPMAtt}$ $Att \rightarrow Meas : d \text{ on } C_{AttMeas}$ $Meas \rightarrow Att : e \text{ on } C_{MeasAtt}$

 $\textit{Att} \rightarrow \textit{TPM}: \textit{quote}(\;\textit{AIK}_{h}, \textit{PCR}_{m}, |(\textit{e}, \textit{N}_{App}, [\textit{AIK}^{+}]_{\textit{CA}^{-}})|\;) \; \text{on} \; \textit{C}_{\textit{AttTPM}}$

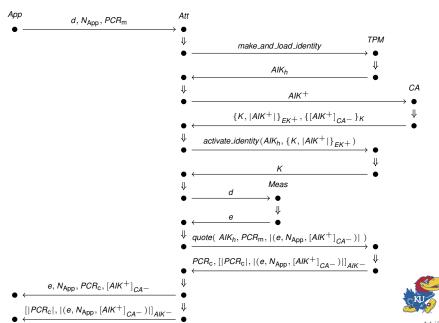
 $\mathit{TPM} \to \mathit{Att} : \mathit{PCR}_c, [|\mathit{PCR}_c|, |(e, \mathit{N}_{\mathsf{App}}, [\mathit{AIK}^+]_{\mathit{CA}^-})|]_{\mathit{AIK}^-} \ \text{on} \ \mathit{C}_{\mathit{TPMAtt}}$

 $Att \rightarrow App : e, N_{App}, PCR_c, [AIK^+]_{CA^-} \text{ on } C_{AttApp}$

 $Att o App : [|PCR_c|, |(e, N_{\mathsf{App}}, [AIK^+]_{\mathit{CA}^-})|]_{\mathit{AIK}^-} \ \mathsf{on} \ \mathit{C}_{\mathit{AttApp}}$



Strand Space Diagram Representation



Attestation Request

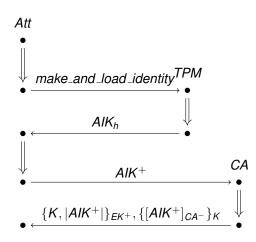


- Initiate with an attestation request
 - d abstractly defines desired evidence
 - ► *N*_{App} is the appraiser's nonce
 - ► PCR_m selects PCRs
- Attestation agent selects and executes protocol based on request



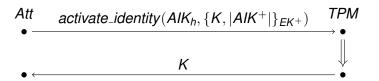
Generating and Certifying an AIK

- Request a new AIK from TPM (optional)
- ► Receive AIK handle
- ▶ Request AIK⁺ signed by CA (AIK cert)
- ► Receive *AIK* cert encrypted with session key *K*
- ► Receive *K* encrypted with public *EK*





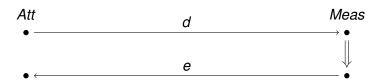
Activating the AIK



- ► Request TPM decryption of the *AIK* cert
- ▶ Receive K used to decrypt signed public AIK
- ▶ Only TPM can gain access to K
- ▶ Only TPM can obtain signed, public AIK
- ► Oddly, No manipulation of the AIK in this "activation" process



Measurement



- ► Request information from measurer
- ► Receive evidence e from measurer
- ► *d* is abstract allowing protocol reuse
- Most protocols make many requests of the measurer



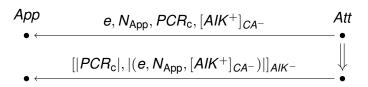
Generating a Quote

Att quote(
$$AIK_h$$
, PCR_m , $|(e, N_{App}, [AIK^+]_{CA^-})|)$ TPN_{CA^-} PCR_c , $||PCR_c|$, $|(e, N_{App}, [AIK^+]_{CA^-})||_{AIK^-}$

- Request a quote from the TPM
 - AIK identifies the signing AIK
 - ► PCR_m identifies desired PCRs
 - $|(e, N_{App}, [AIK^+]_{CA^-})|$ guarantees integrity of returned evidence
- ► Receive quote from TPM
 - ▶ PCR_c is PCR composite built from requested PCRs
 - ► $[|PCR_c|, |(e, N_{App}, [AIK^+]_{CA^-})|]_{AIK^-}$ is the signed quote



Appraisal



- Receive evidence from the attestation manager
 - evidence
 - original nonce
 - PCR composite
 - ▶ signed AIK⁺
- Receive TPM quote from the attestation manager
 - hash of all evidence
 - ► PCR composite
 - ▶ signed by AIK⁻
- ► Evaluate evidence and quote



Evidence Appraisal

Demonstration detects failure of all aspects of attestation

$$e, N_{App}, PCR_{c}, [AIK^{+}]_{CA^{-}}$$

- ► e evidence gathered from running application
- ► N_{App} prevents replay
- ▶ PCR_c evidence in the form or PCR data from the vTPM
- ► [AIK⁺]_{CA} ensures validity of AIK⁺

$$[|PCR_c|, |(e, N_{App}, [AIK^+]_{CA^-})|]_{AIK^-}$$

- ▶ PCR_c hash ensures integrity of PCR data
- ► $|(e, N_{App}, [AIK^+]_{CA^-})|$ hash ensures integrity of evidence, nonce, and signed AIK^+

Attestation Protocol Execution



3-4 Slides on Measurement



2-3 Slides on Communication Mechanisms



CA communication

Shared notion of AIKCertRequest, AIKCert, and CAResponse JSON structures.

Attester

- creates an AIKCertRequest (containing attester ID, AIK⁺) and converts to JSON
- ► JSON sent as POST request to CA running as web server

Certificate Authority

- ▶ POST body bytes \rightarrow UTF8 \rightarrow JSON \rightarrow AIKCertRequest
- ▶ looks up EK^+ associated with ID in sql database
- ► AIKCert = AIK⁺ signed with CA⁻
- ▶ generates key K and encrypts with EK⁺
- ► AIKCert encrypted with *K*
- both wrapped in a CAResponse, converted to JSON and set as response.

CA communication continued

Properties

- ► CA only responds to receiving an AIKCertRequest_{JSON}
- ► The CACert can *only* be decrypted by knowing K (and therefore EK⁻)

Appraiser Knowledge after receiving Cert:

- signature on AIK ensures it was CA who generated signature
- only an entity knowing EK⁻ could decrypt and send the CACert
 - =
- Attester is using a registered TPM



Current Status

Completed four demonstrations culminating in running an attestation protocol in response to an attestation request.

► Attestation and Appraisal development

- CA-Based attestation protocol execution example
- ▶ integration with Berlios TPM 1.2 emulator
- simple dynamic appraisal of attestation results

Measurement development

- on demand Java program measurement
- HotSpot-based Java VM run time measurements
- standard mechanism for extending measurement capabilities

▶ Communication infrastructure

- vchan, TCP/IP and socket communication infrastructure
- ► language-based interface with TPM 1.2
- JSON-based data exchange formats
- initial certificate authority API



Goals and Milestones for 2015

- ► Push to the cloud
- Establish roots of trust and trust argument
- Executable protocol representation and protocol semantics
- ► Operational, integrated vTPM prototype
- Name Server / Certificate Authority prototype
- ► More capable measurement
- ▶ Downloadable demonstration



Questions and Guidance

- ▶ What problems are interesting?
- ▶ What problem would be a nice attention grabber?
- ▶ What should we be watching and integrating with?



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