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Graph Algorithms with Hostile Partners

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Abstract

A short description of the project goes here.

Contents

1	Introduction			
2	Dominating sets	3		
	2.1 min size dominating set	3		

Figures

Chapter 1

Introduction

Chapter 2

Dominating sets

We begin by listing some definitions.

Definition. The Dominating set, D, of a graph G = (E, V) is any subset of V such that every vertex in V is adjacent to at least one vertex in D.

Definition. The Dominating number, $\gamma(G)$, of a graph G = (E, V) is the size of the smallest dominating set of G.

Definition. Independent set, maximum independent set, independence number $\alpha(G)$

2.1 min size dominating set

Lemma 2.1. *Let G be a graph.*

$$\gamma(G) \ge \alpha(G)$$

Proof. Let X be a minimum dominating set in some graph G = (V, E). By definition of dominating set vertex in V is adjacent to at least one vertex in

Recall that $\chi(G)$ is the chromatic number of the graph G.

Theorem 2.2 (Willis 2011 3.1). *For any graph G* = (V, E) [4]

$$\alpha(G) \le \frac{|V|}{\chi(G)}$$

Recall that $\Delta(G)$ is the maximum degree of any vertex in G.

Theorem 2.3 (Balakrishnan 2012 10.3.2). [2] For any graph G with n vertices,

$$\left\lceil \frac{n}{1 + \Delta(G)} \right\rceil \le \gamma(G) \le n - \Delta(G)$$

It is obvious that in the case when $\gamma(G)$ is known that $\gamma(G) > \gamma_g(G)$.

Theorem 2.4. (Ore 1962) [3] For any graph G with n vertices,

$$\gamma(G) \leq \frac{n}{2}$$

Theorem 2.5. Let G be a graph. If x is a tight upper bound for the domination number, $\gamma(G)$, then

$$\gamma_{g}(G) \geq x$$

Proof. Let *G* be a graph where $\gamma(G) = x$. Thus for *G* we are unable to find a dominating set with < x vertices. Therefore there cannot be a winning strategy for Alice with < x vertices. Therefore $\gamma_g(G) \ge x$

Theorem 2.6. Let G be a graph with n vertices, such that $n \geq 4$. Then,

$$\gamma_g(G) \ge \left\lfloor \frac{n}{2} \right\rfloor$$

Proof. By combination of theorems 2.4 and 2.5 we get $\gamma_g(G) \geq \lfloor \frac{n}{2} \rfloor$

Thereom 2.6 is also proved in Alona, Baloghc, Bollobas, and Szabo 2002 [1]. The trivial upper bound is n.

Theorem 2.7. *Let G be a graph with n vertices. Then,*

$$\gamma_g(G) \le \left\lceil \frac{2n}{3} \right\rceil$$

Proof. A dominating set on a spanning tree in a dominating set in the parent graph. Thus for any graph, *G*, it suffices to show we have a winning strategy for a spanning tree of *G*. let *T* be a spanning tree of *G*. The winning strategy for Alice is the greedy strategy as follows. Let *D* be the current dominating set in *T* i.e. neighbours of all selected vertices.

- 1. Pick any vertex, v, not in D with a maximal number of neighbours not in D. That is maximise the set $\{x: x \in N(v) \land v \notin D\}$.
- 2. repeat until you have a dominating set.

worst case path graph requires twice the minimum of the path graph??? with no opponent this will give n/3 thus at worst with the opponent it will take 2n/3 At worst Alice will add two vertices to

Theorem 2.8. *Given p players then,*

$$\gamma_{gp}(G) \geq p\gamma(G)$$

$$\gamma_{gp}(G) \le p\gamma_{g2}(G) \le p\left\lceil \frac{2n}{3}\right\rceil$$

Chapter 3

Colouring

3.1 title

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- [4] WILLIS, W. Bounds for the independence number of a graph. Virginia Commonwealth University.