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# Graph Algorithms with Hostile Partners

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#### **Abstract**

A short description of the project goes here.

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## **Chapter 1**

## Introduction

#### **Chapter 2**

### **Dominating sets**

We begin by listing some definitions.

**Definition.** The Dominating set, D, of a graph G = (E, V) is any subset of V such that every vertex in V is adjacent to at least one vertex in D.

**Definition.** The Dominating number,  $\gamma(G)$ , of a graph G = (E, V) is the size of the smallest dominating set of G.

**Definition.** Independent set, maximum independent set, independence number  $\alpha(G)$ 

#### 2.1 min size dominating set

**Lemma 2.1.** *Let G be a graph.* 

$$\gamma(G) \ge \alpha(G)$$

*Proof.* Let X be a minimum dominating set in some graph G = (V, E). By definition of dominating set vertex in V is adjacent to at least one vertex in

Recall that  $\chi(G)$  is the chromatic number of the graph G.

**Theorem 2.2** (Willis 2011 3.1). *For any graph G* = (V, E) [2]

$$\alpha(G) \le \frac{|V|}{\chi(G)}$$

Recall that  $\Delta(G)$  is the maximum degree of any vertex in G.

**Theorem 2.3** (Balakrishnan 2012 10.3.2). For any graph G with n vertices,

$$\left\lceil \frac{n}{1 + \Delta(G)} \right\rceil \le \gamma(G) \le n - \Delta(G)$$

[1]

It is obvious that in the case when  $\gamma(G)$  is known that  $\gamma(G) > \gamma_g(G)$ .

**Theorem 2.4.** Let G be a graph if x is a tight upper bound for the domination number  $(\gamma(G))$  then

$$\gamma_{g}(G) \geq x$$

*Proof.* Let *G* be a graph where  $\gamma(G) = x$ 

Thus for G we are unable to find a dominating set with < x vertices.

Therefore there cannot be a winning strategy for Alice with < *x* vertices.

Therefore  $\gamma_g(G) \ge x$ 

**Theorem 2.5.** *Let* G *be a graph with* n *vertices, such that*  $n \geq 4$ . *Then,* 

$$\gamma_g(G) > \left\lfloor \frac{n}{2} \right\rfloor$$

*Proof.* Worst case is *G* is minimally connected, i.e. *G* is a tree. Thus by 2.3

$$\gamma(G) \ge \left\lceil \frac{n}{1 + \Delta(G)} \right\rceil$$

**Theorem 2.6.** Let G be a graph with n vertices. Then,

$$\gamma_g(G) \leq \frac{2n}{3}$$

## Bibliography

- [1] BALAKRISHNAN, R. A Textbook of Graph Theory (Universitext). Springer, sep 2012.
- [2] WILLIS, W. Bounds for the independence number of a graph.