

# Sauer-Shelah Lemma

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## Abstract

The Sauer-Shelah Lemma is a fundamental result in extremal set theory and combinatorics, that guarantees the existence of a set  $T$  of size  $k$  which is shattered by a family of sets  $\mathcal{F}$ , if the cardinality of the family is greater than some bound dependent on  $k$ . A set  $T$  is said to be shattered by a family  $\mathcal{F}$  if every subset of  $T$  can be obtained as an intersection of  $T$  with some set  $S \in \mathcal{F}$ . The Sauer-Shelah Lemma has found use in diverse fields such as computational geometry, approximation algorithms and machine learning. In this entry we formalize the notion of shattering and prove the generalized and the standard version of the Sauer-Shelah Lemma.

## Contents

<b>1</b>	<b>Definitions and lemmas about shattering</b>	<b>1</b>
1.1	Intersection of a family of sets with a set . . . . .	1
1.2	Definition of shattering . . . . .	3
<b>2</b>	<b>Lemmas involving the cardinality of sets</b>	<b>4</b>
<b>3</b>	<b>Lemmas involving the binomial coefficient</b>	<b>6</b>
<b>4</b>	<b>Sauer-Shelah Lemma</b>	<b>7</b>
4.1	Generalized Sauer-Shelah Lemma . . . . .	7
4.2	Sauer-Shelah Lemma . . . . .	10
4.3	Sauer-Shelah Lemma for hypergraphs . . . . .	11
4.4	Alternative statement of the Sauer-Shelah Lemma . . . . .	11

## 1 Definitions and lemmas about shattering

```
theory Shattering
  imports Main
begin
```

## 1.1 Intersection of a family of sets with a set

**abbreviation**  $\text{IntF} :: 'a \text{ set set} \Rightarrow 'a \text{ set} \Rightarrow 'a \text{ set set}$  (infixl  $\cap^*$  60)  
**where**  $F \cap^* S \equiv ((\cap) S) ' F$

**lemma** *idem-IntF*:

**assumes**  $\bigcup A \subseteq Y$

**shows**  $A \cap^* Y = A$

**proof** –

**from** *assms* **have**  $A \subseteq A \cap^* Y$  **by** *blast*

**thus** *?thesis* **by** *fastforce*

**qed**

**lemma** *subset-IntF*:

**assumes**  $A \subseteq B$

**shows**  $A \cap^* X \subseteq B \cap^* X$

**using** *assms* **by** (*rule image-mono*)

**lemma** *Int-IntF*:  $(A \cap^* Y) \cap^* X = A \cap^* (Y \cap X)$

**proof**

**show**  $A \cap^* Y \cap^* X \subseteq A \cap^* (Y \cap X)$

**proof**

**fix**  $S$

**assume**  $S \in A \cap^* Y \cap^* X$

**then obtain**  $a\text{-}y$  **where**  $A\text{-}Y0$ :  $a\text{-}y \in A \cap^* Y$  **and**  $A\text{-}Y1$ :  $a\text{-}y \cap X = S$  **by**

*blast*

**from**  $A\text{-}Y0$  **obtain**  $a$  **where**  $A0$ :  $a \in A$  **and**  $A1$ :  $a \cap Y = a\text{-}y$  **by** *blast*

**from**  $A\text{-}Y1$   $A1$  **have**  $a \cap (Y \cap X) = S$  **by** *fast*

**with**  $A0$  **show**  $S \in A \cap^* (Y \cap X)$  **by** *blast*

**qed**

**next**

**show**  $A \cap^* (Y \cap X) \subseteq A \cap^* Y \cap^* X$

**proof**

**fix**  $S$

**assume**  $S \in A \cap^* (Y \cap X)$

**then obtain**  $a$  **where**  $A0$ :  $a \in A$  **and**  $A1$ :  $a \cap (Y \cap X) = S$  **by** *blast*

**from**  $A0$  **have**  $a \cap Y \in A \cap^* Y$  **by** *blast*

**with**  $A1$  **show**  $S \in (A \cap^* Y) \cap^* X$  **by** *blast*

**qed**

**qed**

insert distributes over IntF

**lemma** *insert-IntF*:

**shows**  $\text{insert } x ' (H \cap^* S) = (\text{insert } x ' H) \cap^* (\text{insert } x S)$

**proof**

**show**  $\text{insert } x ' (H \cap^* S) \subseteq (\text{insert } x ' H) \cap^* (\text{insert } x S)$

**proof**

**fix**  $y\text{-}x$

**assume**  $y\text{-}x \in \text{insert } x ' (H \cap^* S)$

**then obtain**  $y$  **where**  $0$ :  $y \in (H \cap^* S)$  **and**  $1$ :  $y\text{-}x = y \cup \{x\}$  **by** *blast*

```

    from 0 obtain  $yh$  where 2:  $yh \in H$  and 3:  $y = yh \cap S$  by blast
    from 1 3 have  $y-x = (yh \cup \{x\}) \cap (S \cup \{x\})$  by simp
    with 2 show  $y-x \in (\text{insert } x \text{ ` } H) \cap * (\text{insert } x \text{ ` } S)$  by blast
  qed
next
  show  $\text{insert } x \text{ ` } H \cap * (\text{insert } x \text{ ` } S) \subseteq \text{insert } x \text{ ` } (H \cap * S)$ 
  proof
    fix  $y-x$ 
    assume  $y-x \in \text{insert } x \text{ ` } H \cap * (\text{insert } x \text{ ` } S)$ 
    then obtain  $yh-x$  where 0:  $yh-x \in (\lambda Y. Y \cup \{x\}) \text{ ` } H$  and 1:  $y-x = yh-x \cap (S \cup \{x\})$  by blast
    from 0 obtain  $yh$  where 2:  $yh \in H$  and 3:  $yh-x = yh \cup \{x\}$  by blast
    from 1 3 have  $y-x = (yh \cap S) \cup \{x\}$  by simp
    with 2 show  $y-x \in \text{insert } x \text{ ` } (H \cap * S)$  by blast
  qed
qed

```

## 1.2 Definition of shattering

**abbreviation**  $\text{shatters} :: 'a \text{ set } \text{set} \Rightarrow 'a \text{ set} \Rightarrow \text{bool}$  (**infixl**  $\text{shatters}$  70)  
 where  $H \text{ shatters } A \equiv H \cap * A = \text{Pow } A$

**definition**  $\text{VC-dim} :: 'a \text{ set } \text{set} \Rightarrow \text{nat}$   
 where  $\text{VC-dim } F = \text{Sup } \{\text{card } S \mid S. F \text{ shatters } S\}$

**definition**  $\text{shattered-by} :: 'a \text{ set } \text{set} \Rightarrow 'a \text{ set } \text{set}$   
 where  $\text{shattered-by } F \equiv \{A. F \text{ shatters } A\}$

**lemma**  $\text{shattered-by-in-Pow}$ :  
 shows  $\text{shattered-by } F \subseteq \text{Pow } (\bigcup F)$   
 unfolding  $\text{shattered-by-def}$  by blast

**lemma**  $\text{subset-shatters}$ :  
 assumes  $A \subseteq B$  and  $A \text{ shatters } X$   
 shows  $B \text{ shatters } X$   
**proof** –  
 from  $\text{assms}(1)$  have  $A \cap * X \subseteq B \cap * X$  by blast  
 with  $\text{assms}(2)$  have  $\text{Pow } X \subseteq B \cap * X$  by presburger  
 thus  $?thesis$  by blast  
**qed**

**lemma**  $\text{supset-shatters}$ :  
 assumes  $Y \subseteq X$  and  $A \text{ shatters } X$   
 shows  $A \text{ shatters } Y$   
**proof** –  
 have  $h: \bigcup (\text{Pow } Y) \subseteq Y$  by simp  
 from  $\text{assms}$  have 0:  $\text{Pow } Y \subseteq A \cap * X$  by auto  
 from  $\text{subset-IntF}[OF \ 0, \text{ of } Y] \ \text{Int-IntF}[of \ Y \ X \ A] \ \text{idem-IntF}[OF \ h]$  have  $\text{Pow } Y \subseteq A \cap * (X \cap Y)$  by argo

```

with Int-absorb2[OF assms(1)] Int-commute[of X Y] have  $\text{Pow } Y \subseteq A \cap^* Y$ 
by presburger
  then show ?thesis by fast
qed

```

```

lemma shatters-empty:
  assumes  $F \neq \{\}$ 
  shows  $F \text{ shatters } \{\}$ 
using assms by fastforce

```

```

lemma subset-shattered-by:
  assumes  $A \subseteq B$ 
  shows  $\text{shattered-by } A \subseteq \text{shattered-by } B$ 
unfolding shattered-by-def using subset-shatters[OF assms] by force

```

```

lemma finite-shattered-by:
  assumes finite ( $\bigcup F$ )
  shows finite ( $\text{shattered-by } F$ )
  using assms rev-finite-subset[OF - shattered-by-in-Pow, of F] by fast

```

The following example shows that requiring finiteness of a family of sets is not enough

```

lemma  $\exists F::\text{nat set set. finite } F \wedge \text{infinite } (\text{shattered-by } F)$ 
proof –
  let  $?F = \{\text{odd} - \{ \text{True} \}, \text{odd} - \{ \text{False} \}\}$ 
  have  $0: \text{finite } ?F$  by simp

  let  $?f = \lambda n::\text{nat. } \{n\}$ 
  let  $?N = \text{range } ?f$ 
  have inj ( $\lambda n. \{n\}$ ) by simp
  with infinite-iff-countable-subset[of ?N] have infinite-N: infinite  $?N$  by blast
  have F-shatters-any-singleton:  $?F \text{ shatters } \{n::\text{nat}\}$  for  $n$ 
  proof –
    have  $\text{Pow-}n: \text{Pow } \{n\} = \{\{n\}, \{\}\}$  by blast
    have  $1: \text{Pow } \{n\} \subseteq ?F \cap^* \{n\}$ 
    proof (cases odd n)
      case True
      from True have  $(\text{odd} - \{ \text{False} \}) \cap \{n\} = \{\}$  by blast
      hence  $0: \{\} \in ?F \cap^* \{n\}$  by blast
      from True have  $(\text{odd} - \{ \text{True} \}) \cap \{n\} = \{n\}$  by blast
      hence  $1: \{n\} \in ?F \cap^* \{n\}$  by blast
      from  $0\ 1\ \text{Pow-}n$  show ?thesis by simp
    next
      case False
      from False have  $(\text{odd} - \{ \text{True} \}) \cap \{n\} = \{\}$  by blast
      hence  $0: \{\} \in ?F \cap^* \{n\}$  by blast
      from False have  $(\text{odd} - \{ \text{False} \}) \cap \{n\} = \{n\}$  by blast
      hence  $1: \{n\} \in ?F \cap^* \{n\}$  by blast
      from  $0\ 1\ \text{Pow-}n$  show ?thesis by simp
  qed

```

```

    qed
    thus ?thesis by fastforce
  qed
  then have ?N  $\subseteq$  shattered-by ?F unfolding shattered-by-def by force
  from 0 infinite-super[OF this infinite-N] show ?thesis by blast
qed

end

```

## 2 Lemmas involving the cardinality of sets

```

theory Card-Lemmas
  imports Main
begin

```

```

lemma card-diff:
  assumes finite A
  shows card A = card (A - B) + card (A  $\cap$  B)
proof -
  from assms have fin0: finite (A - B) and fin1: finite (A  $\cap$  B) by blast+
  have A-equ: A = (A - B)  $\cup$  (A  $\cap$  B) and disjoint: (A - B)  $\cap$  (A  $\cap$  B) = {}
  by blast+
  from card-Un-disjoint[OF fin0 fin1 disjoint] A-equ show ?thesis by argo
qed

```

```

lemma card-Int-copy:
  assumes finite X and A  $\cup$  B  $\subseteq$  X and  $\exists f. \text{inj-on } f \text{ } (A \cap B) \wedge (A \cup B) \cap (f \text{ ` } (A \cap B)) = \{\}$ 
  shows card A + card B  $\leq$  card X
proof -
  from rev-finite-subset[OF assms(1), of A] rev-finite-subset[OF assms(1), of B]
  assms(2)
  have finite-A: finite A and finite-B: finite B by blast+
  then have finite-A-Un-B: finite (A  $\cup$  B) and finite-A-Int-B: finite (A  $\cap$  B) by
  blast+
  from assms(3) obtain f where f-inj-on: inj-on f (A  $\cap$  B) and f-disjnt: (A  $\cup$ 
  B)  $\cap$  (f ` (A  $\cap$  B)) = {} and f-imj-in: f ` (A  $\cap$  B)  $\subseteq$  X by blast
  from finite-A-Int-B have finite-f-img: finite (f ` (A  $\cap$  B)) by blast
  from assms(2) f-imj-in have union-in: (A  $\cup$  B)  $\cup$  f ` (A  $\cap$  B)  $\subseteq$  X by blast

  from card-Un-Int[OF finite-A finite-B] have card A + card B = card (A  $\cup$  B)
  + card (A  $\cap$  B) .
  also from card-image[OF f-inj-on] have ... = card (A  $\cup$  B) + card (f ` (A  $\cap$ 
  B)) by presburger
  also from card-Un-disjoint[OF finite-A-Un-B finite-f-img f-disjnt] have ... =
  card ((A  $\cup$  B)  $\cup$  f ` (A  $\cap$  B)) by argo
  also from card-mono[OF assms(1) union-in] have ...  $\leq$  card X by blast
  finally show ?thesis .

```

qed

lemma *card-ge-0*:

assumes  $A \neq \{\}$  and *finite*  $A$

shows  $0 < \text{card } A$

proof –

from *assms*(1) have  $\{\} \subset A$  by *blast*

from *psubset-card-mono*[*OF assms*(2) *this*] show ?thesis by *force*

qed

lemma *finite-diff-not-empty*:

assumes *finite*  $Y$  and  $\text{card } Y < \text{card } X$

shows  $X - Y \neq \{\}$

proof

assume  $X - Y = \{\}$

hence  $X \subseteq Y$  by *simp*

from *card-mono*[*OF assms*(1) *this*] *assms*(2) show *False* by *linarith*

qed

lemma *obtain-difference-element*:

fixes  $F :: 'a \text{ set set}$

assumes  $2 \leq \text{card } F$

obtains  $x$  where  $x \in \bigcup F$   $x \notin \bigcap F$

proof –

from *assms* *card-le-Suc-iff*[*of* 1  $F$ ] obtain  $A$   $F'$  where 0:  $F = \text{insert } A \ F'$  and 1:  $A \notin F'$  and 2:  $1 \leq \text{card } F'$  by *auto*

from 2 *card-le-Suc-iff*[*of* 0  $F$ ] obtain  $B$   $F''$  where 3:  $F' = \text{insert } B \ F''$  by *auto*

from 1 3 have *A-noteq-B*:  $A \neq B$  by *blast*

from 0 3 have *A-in-F*:  $A \in F$  and *B-in-F*:  $B \in F$  by *blast*+

from *A-noteq-B* have  $(A - B) \cup (B - A) \neq \{\}$  by *simp*

with *A-in-F* *B-in-F* that show *thesis* by *blast*

qed

end

### 3 Lemmas involving the binomial coefficient

theory *Binomial-Lemmas*

imports *Main*

begin

lemma *choose-mono*:

assumes  $x \leq y$

shows  $x \text{ choose } n \leq y \text{ choose } n$

proof –

have *finite*  $\{0..<y\}$  by *blast*

with *finite-Pow-iff*[*of*  $\{0..<y\}$ ] have *finiteness*: *finite*  $\{K \in \text{Pow } \{0..<y\}. \text{card } K = n\}$  by *simp*

**from** *assms* **have**  $\text{Pow } \{0..<x\} \subseteq \text{Pow } \{0..<y\}$  **by** *force*  
**then** **have**  $\{K \in \text{Pow } \{0..<x\}. \text{card } K = n\} \subseteq \{K \in \text{Pow } \{0..<y\}. \text{card } K = n\}$  **by** *blast*  
**from** *card-mono[OF finiteness this]* **show** *?thesis unfolding binomial-def .*  
**qed**

**lemma** *choose-row-sum-set:*

**assumes** *finite*  $(\bigcup F)$   
**shows**  $\text{card } \{S. S \subseteq \bigcup F \wedge \text{card } S \leq k\} = (\sum i \leq k. \text{card } (\bigcup F) \text{ choose } i)$   
**proof** (*induction k*)  
**case** 0  
**from** *rev-finite-subset[OF assms]* **have**  $S \subseteq \bigcup F \wedge \text{card } S \leq 0 \longleftrightarrow S = \{\}$  **for** *S* **by** *fastforce*  
**then** **show** *?case* **by** *simp*  
**next**  
**case** (*Suc k*)  
**let** *?FS* =  $\{S. S \subseteq \bigcup F \wedge \text{card } S \leq \text{Suc } k\}$   
**and** *?F-Asm* =  $\{S. S \subseteq \bigcup F \wedge \text{card } S \leq k\}$   
**and** *?F-Step* =  $\{S. S \subseteq \bigcup F \wedge \text{card } S = \text{Suc } k\}$   
  
**from** *finite-Pow-iff[of  $\bigcup F$  assms]* **have** *finite-Pow-Un-F*: *finite*  $(\text{Pow } (\bigcup F))$  **..**  
**have** *?F-Asm*  $\subseteq \text{Pow } (\bigcup F)$  **and** *?F-Step*  $\subseteq \text{Pow } (\bigcup F)$  **by** *fast+*  
**with** *rev-finite-subset[OF finite-Pow-Un-F]* **have** *finite-F-Asm*: *finite* *?F-Asm*  
**and** *finite-F-Step*: *finite* *?F-Step* **by** *presburger+*  
  
**have** *F-Un*: *?FS* = *?F-Asm*  $\cup$  *?F-Step* **and** *F-disjoint*: *?F-Asm*  $\cap$  *?F-Step* =  $\{\}$   
**by** *fastforce+*  
**from** *card-Un-disjoint[OF finite-F-Asm finite-F-Step F-disjoint]* *F-Un* **have** *card* *?FS* = *card* *?F-Asm* + *card* *?F-Step* **by** *argo*  
**also** **from** *Suc* **have**  $\dots = (\sum i \leq k. \text{card } (\bigcup F) \text{ choose } i) + \text{card } ?F\text{-Step}$  **by** *argo*  
**also** **from** *n-subsets[OF assms, of Suc k]* **have**  $\dots = (\sum i \leq \text{Suc } k. \text{card } (\bigcup F) \text{ choose } i)$  **by** *force*  
**finally** **show** *?case* **by** *blast*  
**qed**  
  
**end**

## 4 Sauer-Shelah Lemma

**theory** *Sauer-Shelah-Lemma*

**imports** *Main Shattering Card-Lemmas Binomial-Lemmas*  
**begin**

### 4.1 Generalized Sauer-Shelah Lemma

**lemma** *sauer-shelah-0:*

**fixes** *F* :: '*a* set set  
**shows** *finite*  $(\bigcup F) \implies \text{card } F \leq \text{card } (\text{shattered-by } F)$   
**proof** (*induction F rule: measure-induct-rule[of card]*)

**case** (*less F*)  
**note** *finite-F* = *finite-UnionD*[*OF less(2)*]  
**note** *finite-shF* = *finite-shattered-by*[*OF less(2)*]  
**show** ?*case*  
**proof** (*cases 2 ≤ card F*)  
**case** *True*  
**from** *obtain-difference-element*[*OF True*] **obtain** *x :: 'a* **where** *x-in-Union-F*:  
*x ∈ ∪ F* **and** *x-not-in-Int-F*: *x ∉ ∩ F* **by** *blast*

Define F0 as the subfamily of F containing those sets that don't contain  
 x

**let** ?*F0* = {*S ∈ F. x ∉ S*}  
**from** *x-in-Union-F* **have** *F0-psubset-F*: ?*F0* ⊂ *F* **by** *blast*  
**from** *F0-psubset-F* **have** *F0-in-F*: ?*F0* ⊆ *F* **by** *blast*  
**from** *subset-shattered-by*[*OF F0-in-F*] **have** *shF0-subset-shF*: *shattered-by* ?*F0*  
 ⊆ *shattered-by F* .  
**from** *F0-in-F* **have** *Un-F0-in-Un-F*: ∪ ?*F0* ⊆ ∪ *F* **by** *blast*

F0 shatters at least as many sets as |F0| by the induction hypothesis

**note** *IH-F0* = *less(1)*[*OF psubset-card-mono*[*OF finite-F F0-psubset-F*] *rev-finite-subset*[*OF less(2) Un-F0-in-Un-F*]]

Define F1 as the subfamily of F containing those sets that contain x

**let** ?*F1* = {*S ∈ F. x ∈ S*}  
**from** *x-not-in-Int-F* **have** *F1-psubset-F*: ?*F1* ⊂ *F* **by** *blast*  
**from** *F1-psubset-F* **have** *F1-in-F*: ?*F1* ⊆ *F* **by** *blast*  
**from** *subset-shattered-by*[*OF F1-in-F*] **have** *shF1-subset-shF*: *shattered-by* ?*F1*  
 ⊆ *shattered-by F* .  
**from** *F1-in-F* **have** *Un-F1-in-Un-F*: ∪ ?*F1* ⊆ ∪ *F* **by** *blast*

F1 shatters at least as many sets as |F1| by the induction hypothesis

**note** *IH-F1* = *less(1)*[*OF psubset-card-mono*[*OF finite-F F1-psubset-F*] *rev-finite-subset*[*OF less(2) Un-F1-in-Un-F*]]

**from** *shF0-subset-shF shF1-subset-shF* **have** *shattered-subset*: (*shattered-by* ?*F0*) ∪ (*shattered-by* ?*F1*) ⊆ *shattered-by F* **by** *simp*

There is a set with the same cardinality as the intersection of *shattered-by* {*S ∈ F. x ∉ S*} and *shattered-by* {*S ∈ F. x ∈ S*} which is disjoint from their union, which is also contained in *shattered-by F*.

**have** *f-copies-the-intersection*:  
 ∃ *f. inj-on f* (*shattered-by* ?*F0* ∩ *shattered-by* ?*F1*) ∧  
 (*shattered-by* ?*F0* ∪ *shattered-by* ?*F1*) ∩ (*f* ‘ (*shattered-by* ?*F0* ∩ *shattered-by* ?*F1*)) = {} ∧  
*f* ‘ (*shattered-by* ?*F0* ∩ *shattered-by* ?*F1*) ⊆ *shattered-by F*  
**proof**  
**have** *x-not-in-shattered*: ∀ *S* ∈ (*shattered-by* ?*F0*) ∪ (*shattered-by* ?*F1*). *x ∉ S*  
**unfolding** *shattered-by-def* **by** *blast*



This set is precisely the image of the intersection under *insert x*.

```

let ?f = insert x
have 0: inj-on ?f (shattered-by ?F0  $\cap$  shattered-by ?F1)
proof
  fix X Y
  assume x0: X  $\in$  (shattered-by ?F0  $\cap$  shattered-by ?F1) and y0: Y  $\in$ 
(shattered-by ?F0  $\cap$  shattered-by ?F1)
  and 0: ?f X = ?f Y
  from x-not-in-shattered x0 have X = ?f X - {x} by blast
  also from 0 have ... = ?f Y - {x} by argo
  also from x-not-in-shattered y0 have ... = Y by blast
  finally show X = Y .
qed

```

The set is disjoint from the union.

```

have 1: (shattered-by ?F0  $\cup$  shattered-by ?F1)  $\cap$  ?f ‘ (shattered-by ?F0  $\cap$ 
shattered-by ?F1) = {}
proof (rule ccontr)
  assume (shattered-by ?F0  $\cup$  shattered-by ?F1)  $\cap$  ?f ‘ (shattered-by ?F0  $\cap$ 
shattered-by ?F1)  $\neq$  {}
  then obtain S where 10: S  $\in$  (shattered-by ?F0  $\cup$  shattered-by ?F1) and
11: S  $\in$  ?f ‘ (shattered-by ?F0  $\cap$  shattered-by ?F1) by auto
  from 10 x-not-in-shattered have x  $\notin$  S by blast
  with 11 show False by blast
qed

```

This set is also in *shattered-by F*.

```

have 2: ?f ‘ (shattered-by ?F0  $\cap$  shattered-by ?F1)  $\subseteq$  shattered-by F
proof
  fix S-x
  assume S-x  $\in$  ?f ‘ (shattered-by ?F0  $\cap$  shattered-by ?F1)
  then obtain S where 20: S  $\in$  shattered-by ?F0 and 21: S  $\in$  shattered-by
?F1 and 22: S-x = ?f S by blast
  from x-not-in-shattered 20 have x-not-in-S: x  $\notin$  S by blast

  from 22 Pow-insert[of x S] have Pow S-x = Pow S  $\cup$  ?f ‘ Pow S by fast
  also from 20 have ... = (?F0  $\cap^*$  S)  $\cup$  (?f ‘ Pow S) unfolding shat-
tered-by-def by blast
  also from 21 have ... = (?F0  $\cap^*$  S)  $\cup$  (?f ‘ (?F1  $\cap^*$  S)) unfolding
shattered-by-def by force
  also from insert-IntF[of x S ?F1] have ... = (?F0  $\cap^*$  S)  $\cup$  (?f ‘ ?F1  $\cap^*$ 
(?f S)) by argo
  also from 22 have ... = (?F0  $\cap^*$  S)  $\cup$  (?F1  $\cap^*$  S-x) by blast
  also from 22 have ... = (?F0  $\cap^*$  S-x)  $\cup$  (?F1  $\cap^*$  S-x) by blast
  also from subset-IntF[OF F0-in-F, of S-x] subset-IntF[OF F1-in-F, of S-x]
have ...  $\subseteq$  (F  $\cap^*$  S-x) by blast
  finally have Pow S-x  $\subseteq$  (F  $\cap^*$  S-x) .
  thus S-x  $\in$  shattered-by F unfolding shattered-by-def by blast
qed

```

**from** 0 1 2 **show**  $\text{inj-on } ?f \text{ (shattered-by } ?F0 \cap \text{shattered-by } ?F1) \wedge$   
 $(\text{shattered-by } ?F0 \cup \text{shattered-by } ?F1) \cap (?f'(\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1)) = \{\}$   $\wedge$   
 $?f'(\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1) \subseteq \text{shattered-by } F$  **by** blast  
**qed**

**have**  $F0\text{-union-}F1\text{-is-}F: ?F0 \cup ?F1 = F$  **by** fastforce  
**from** finite- $F$  **have** finite- $F0$ : finite  $?F0$  **and** finite- $F1$ : finite  $?F1$  **by** fastforce+  
**have** disjoint- $F0$ - $F1$ :  $?F0 \cap ?F1 = \{\}$  **by** fastforce

Thus we have the following lower bound on the cardinality of *shattered-by*  $F$

**from**  $F0\text{-union-}F1\text{-is-}F$  card-Un-disjoint[ $OF$  finite- $F0$  finite- $F1$  disjoint- $F0$ - $F1$ ]

**have**  $\text{card } F = \text{card } ?F0 + \text{card } ?F1$  **by** argo  
**also from** IH- $F0$   
**have**  $\dots \leq \text{card } (\text{shattered-by } ?F0) + \text{card } ?F1$  **by** linarith  
**also from** IH- $F1$   
**have**  $\dots \leq \text{card } (\text{shattered-by } ?F0) + \text{card } (\text{shattered-by } ?F1)$  **by** linarith  
**also from** card-Int-copy[ $OF$  finite-sh $F$  shattered-subset f-copies-the-intersection]  
**have**  $\dots \leq \text{card } (\text{shattered-by } F)$  **by** argo  
**finally show**  $?thesis$  .

**next**

If  $F$  contains less than 2 sets, the statement follows trivially

**case** False  
**hence**  $\text{card } F = 0 \vee \text{card } F = 1$  **by** force  
**thus**  $?thesis$   
**proof**

**assume**  $\text{card } F = 0$   
**thus**  $?thesis$  **by** auto

**next**

**assume**  $asm: \text{card } F = 1$   
**hence**  $F\text{-not-empty}: F \neq \{\}$  **by** fastforce

**from** shatters-empty[ $OF$   $F\text{-not-empty}$ ] **have**  $\{\{\}\} \subseteq \text{shattered-by } F$  **unfolding**  
 $\text{shattered-by-def}$  **by** fastforce

**from** card-mono[ $OF$  finite-sh $F$  this]  $asm$  **show**  $?thesis$  **by** fastforce

**qed**

**qed**

**qed**

## 4.2 Sauer-Shelah Lemma

**corollary** sauer-shelah:

**fixes**  $F :: 'a \text{ set set}$

**assumes** finite  $(\bigcup F)$  **and**  $(\sum_{i \leq k}. \text{card } (\bigcup F) \text{ choose } i) < \text{card } F$

**shows**  $\exists S. (F \text{ shatters } S \wedge \text{card } S = k + 1)$

**proof** –

**let**  $?K = \{S. S \subseteq \bigcup F \wedge \text{card } S \leq k\}$   
**from** *finite-Pow-iff*[*of F*] *assms(1)* **have** *finite-Pow-Un*: *finite* (*Pow* ( $\bigcup F$ )) **by**  
*fast*  
  
**from** *sauer-shelah-0*[*OF assms(1)*] *assms(2)* **have**  $(\sum_{i \leq k}. \text{card } (\bigcup F) \text{ choose } i) < \text{card } (\text{shattered-by } F)$  **by** *linarith*  
**with** *choose-row-sum-set*[*OF assms(1)*, *of k*] **have**  $\text{card } ?K < \text{card } (\text{shattered-by } F)$  **by** *presburger*  
  
**from** *finite-diff-not-empty*[*OF finite-subset*[*OF - finite-Pow-Un*] *this*]  
**obtain** *S* **where**  $S \in \text{shattered-by } F - ?K$  **by** *blast*  
**then have** *F-shatters-S*: *F shatters S* **and**  $S \subseteq \bigcup F$  **and**  $\neg(S \subseteq \bigcup F \wedge \text{card } S \leq k)$  **unfolding** *shattered-by-def* **by** *blast+*  
**then have** *card-S-ge-Suc-k*:  $k + 1 \leq \text{card } S$  **by** *simp*  
**from** *obtain-subset-with-card-n*[*OF card-S-ge-Suc-k*] **obtain** *S'* **where**  $\text{card } S' = k + 1$  **and**  $S' \subseteq S$  **by** *blast*  
**from** *this(1)* *supset-shatters*[*OF this(2)* *F-shatters-S*] **show** *?thesis* **by** *blast*  
**qed**

### 4.3 Sauer-Shelah Lemma for hypergraphs

**corollary** *sauer-shelah-2*:

**fixes** *X* :: 'a set set **and** *S* :: 'a set  
**assumes** *finite S* **and**  $X \subseteq \text{Pow } S$  **and**  $(\sum_{i \leq k}. \text{card } S \text{ choose } i) < \text{card } X$   
**shows**  $\exists Y. (X \text{ shatters } Y \wedge \text{card } Y = k + 1)$   
**proof** –  
**from** *assms(2)* **have** *0*:  $\bigcup X \subseteq S$  **by** *blast*  
**from** *sum-mono*[*OF choose-mono*[*OF card-mono*[*OF assms(1)* *0*]]] **have**  $(\sum_{i \leq k}. \text{card } (\bigcup X) \text{ choose } i) \leq (\sum_{i \leq k}. \text{card } S \text{ choose } i)$  **by** *fast*  
**with** *sauer-shelah*[*OF finite-subset*[*OF 0 assms(1)*]] *assms(3)* **show** *?thesis* **by**  
*simp*  
**qed**

### 4.4 Alternative statement of the Sauer-Shelah Lemma

**corollary** *sauer-shelah-alt*:

**assumes** *finite* ( $\bigcup F$ ) **and**  $\text{VC-dim } F = k$   
**shows**  $\text{card } F \leq (\sum_{i \leq k}. \text{card } (\bigcup F) \text{ choose } i)$   
**proof** (*rule ccontr*)  
**assume**  $\neg \text{card } F \leq (\sum_{i \leq k}. \text{card } (\bigcup F) \text{ choose } i)$  **hence**  $(\sum_{i \leq k}. \text{card } (\bigcup F) \text{ choose } i) < \text{card } F$  **by** *linarith*  
**from** *sauer-shelah*[*OF assms(1)* *this*] **obtain** *S* **where** *F shatters S* **and**  $\text{card } S = k + 1$  **by** *blast*  
**from** *this(1)* *this(2)*[*symmetric*] **have**  $k + 1 \in \{\text{card } S \mid S. F \text{ shatters } S\}$  **by**  
*blast*  
**from** *cSup-upper*[*OF this bdd-above-finite*[*OF finite-image-set*[*OF finite-shattered-by*[*unfolded shattered-by-def*, *OF assms(1)*]]], *folded VC-dim-def*]  
*assms(2)* **show** *False* **by** *force*  
**qed**

**end**