

Sauer-Shelah Lemma

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Abstract

The Sauer-Shelah Lemma is a fundamental result in extremal set theory and combinatorics, that guarantees the existence of a set T of size k which is shattered by a family of sets \mathcal{F} , if the cardinality of the family is greater than some bound dependent on k . A set T is said to be shattered by a family \mathcal{F} if every subset of T can be obtained as an intersection of T with some set $S \in \mathcal{F}$. The Sauer-Shelah Lemma has found use in diverse fields such as computational geometry, approximation algorithms and machine learning. In this entry we formalize the notion of shattering and prove the generalized and standard versions of the Sauer-Shelah Lemma.

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1 Introduction

The goal of this entry is to formalize the Sauer-Shelah Lemma. The result was first published by Sauer [2] and Shelah [3] independent from one another. The proof presented in this entry is based on an article by Kalai [1].

The lemma has a wide range of applications. Vapnik and Chervonenkis [4] reproved and used the lemma in the context of statistical learning theory. For instance, the VC-dimension of a family of sets is defined as the size of the largest set the family shatters. In this context the Sauer-Shelah Lemma is a result tying the VC-dimension of a family to the number of sets in the family.

2 Definitions and lemmas about shattering

In this section we introduce the predicate *shatters* and the term for the family of sets that a family shatters *shattered-by*.

```
theory Shattering
  imports Main
begin
```

2.1 Intersection of a family of sets with a set

```
abbreviation IntF :: 'a set set  $\Rightarrow$  'a set  $\Rightarrow$  'a set set (infixl  $\cap^*$  60)
  where  $F \cap^* S \equiv ((\cap) S) ' F$ 
```

```
lemma idem-IntF:
  assumes  $\bigcup A \subseteq Y$ 
  shows  $A \cap^* Y = A$ 
```

```
proof –
  from assms have  $A \subseteq A \cap^* Y$  by blast
  thus ?thesis by fastforce
qed
```

```
lemma subset-IntF:
  assumes  $A \subseteq B$ 
  shows  $A \cap^* X \subseteq B \cap^* X$ 
  using assms by (rule image-mono)
```

```
lemma Int-IntF:  $(A \cap^* Y) \cap^* X = A \cap^* (Y \cap X)$ 
```

```
proof
  show  $A \cap^* Y \cap^* X \subseteq A \cap^* (Y \cap X)$ 
  proof
    fix  $S$ 
    assume  $S \in A \cap^* Y \cap^* X$ 
    then obtain  $a-y$  where  $A-Y0$ :  $a-y \in A \cap^* Y$  and  $A-Y1$ :  $a-y \cap X = S$  by
    blast
    from  $A-Y0$  obtain  $a$  where  $A0$ :  $a \in A$  and  $A1$ :  $a \cap Y = a-y$  by blast
    from  $A-Y1$   $A1$  have  $a \cap (Y \cap X) = S$  by fast
    with  $A0$  show  $S \in A \cap^* (Y \cap X)$  by blast
  qed
next
  show  $A \cap^* (Y \cap X) \subseteq A \cap^* Y \cap^* X$ 
  proof
```

```

    fix S
    assume  $S \in A \cap^* (Y \cap X)$ 
    then obtain a where  $A0: a \in A$  and  $A1: a \cap (Y \cap X) = S$  by blast
    from A0 have  $a \cap Y \in A \cap^* Y$  by blast
    with A1 show  $S \in (A \cap^* Y) \cap^* X$  by blast
  qed
qed

  insert distributes over  $(\cap^*)$ 

lemma insert-IntF:
  shows  $\text{insert } x \text{ ` } (H \cap^* S) = (\text{insert } x \text{ ` } H) \cap^* (\text{insert } x \text{ ` } S)$ 
proof
  show  $\text{insert } x \text{ ` } (H \cap^* S) \subseteq (\text{insert } x \text{ ` } H) \cap^* (\text{insert } x \text{ ` } S)$ 
  proof
    fix y-x
    assume  $y-x \in \text{insert } x \text{ ` } (H \cap^* S)$ 
    then obtain y where  $0: y \in (H \cap^* S)$  and  $1: y-x = y \cup \{x\}$  by blast
    from 0 obtain yh where  $2: yh \in H$  and  $3: y = yh \cap S$  by blast
    from 1 3 have  $y-x = (yh \cup \{x\}) \cap (S \cup \{x\})$  by simp
    with 2 show  $y-x \in (\text{insert } x \text{ ` } H) \cap^* (\text{insert } x \text{ ` } S)$  by blast
  qed
next
  show  $\text{insert } x \text{ ` } H \cap^* (\text{insert } x \text{ ` } S) \subseteq \text{insert } x \text{ ` } (H \cap^* S)$ 
  proof
    fix y-x
    assume  $y-x \in \text{insert } x \text{ ` } H \cap^* (\text{insert } x \text{ ` } S)$ 
    then obtain yh-x where  $0: yh-x \in (\lambda Y. Y \cup \{x\}) \text{ ` } H$  and  $1: y-x = yh-x \cap (S \cup \{x\})$  by blast
    from 0 obtain yh where  $2: yh \in H$  and  $3: yh-x = yh \cup \{x\}$  by blast
    from 1 3 have  $y-x = (yh \cap S) \cup \{x\}$  by simp
    with 2 show  $y-x \in \text{insert } x \text{ ` } (H \cap^* S)$  by blast
  qed
qed

```

2.2 Definition of shatters, VC-dim and shattered-by

abbreviation $\text{shatters} :: 'a \text{ set set} \Rightarrow 'a \text{ set} \Rightarrow \text{bool}$ (**infixl** shatters 70)
 where $H \text{ shatters } A \equiv H \cap^* A = \text{Pow } A$

definition $\text{VC-dim} :: 'a \text{ set set} \Rightarrow \text{nat}$
 where $\text{VC-dim } F = \text{Sup } \{\text{card } S \mid S. F \text{ shatters } S\}$

definition $\text{shattered-by} :: 'a \text{ set set} \Rightarrow 'a \text{ set set}$
 where $\text{shattered-by } F \equiv \{A. F \text{ shatters } A\}$

lemma $\text{shattered-by-in-Pow}$:
 shows $\text{shattered-by } F \subseteq \text{Pow } (\bigcup F)$
 unfolding shattered-by-def by blast

lemma subset-shatters :

```

    assumes  $A \subseteq B$  and  $A$  shatters  $X$ 
    shows  $B$  shatters  $X$ 
  proof -
    from  $\text{assms}(1)$  have  $A \cap^* X \subseteq B \cap^* X$  by blast
    with  $\text{assms}(2)$  have  $\text{Pow } X \subseteq B \cap^* X$  by presburger
    thus  $?thesis$  by blast
  qed

```

```

lemma supset-shatters:
  assumes  $Y \subseteq X$  and  $A$  shatters  $X$ 
  shows  $A$  shatters  $Y$ 
  proof -
    have  $h: \bigcup (\text{Pow } Y) \subseteq Y$  by simp
    from  $\text{assms}$  have  $0: \text{Pow } Y \subseteq A \cap^* X$  by auto
    from subset-IntF[OF 0, of  $Y$ ] Int-IntF[of  $Y$   $X$   $A$ ] idem-IntF[OF  $h$ ] have  $\text{Pow } Y \subseteq A \cap^* (X \cap Y)$  by argo
    with Int-absorb2[OF  $\text{assms}(1)$ ] Int-commute[of  $X$   $Y$ ] have  $\text{Pow } Y \subseteq A \cap^* Y$  by presburger
    then show  $?thesis$  by fast
  qed

```

```

lemma shatters-empty:
  assumes  $F \neq \{\}$ 
  shows  $F$  shatters  $\{\}$ 
  using  $\text{assms}$  by fastforce

```

```

lemma subset-shattered-by:
  assumes  $A \subseteq B$ 
  shows  $\text{shattered-by } A \subseteq \text{shattered-by } B$ 
  unfolding shattered-by-def using subset-shatters[OF  $\text{assms}$ ] by force

```

```

lemma finite-shattered-by:
  assumes finite  $(\bigcup F)$ 
  shows finite  $(\text{shattered-by } F)$ 
  using  $\text{assms}$  rev-finite-subset[OF - shattered-by-in-Pow, of  $F$ ] by fast

```

The following example shows that requiring finiteness of a family of sets is not enough, to ensure that *shattered-by* also stays finite.

```

lemma  $\exists F::\text{nat set set. finite } F \wedge \text{infinite } (\text{shattered-by } F)$ 
  proof -
    let  $?F = \{\text{odd} - \{ \text{True} \}, \text{odd} - \{ \text{False} \} \}$ 
    have  $0: \text{finite } ?F$  by simp

    let  $?f = \lambda n::\text{nat. } \{n\}$ 
    let  $?N = \text{range } ?f$ 
    have inj  $(\lambda n. \{n\})$  by simp
    with infinite-iff-countable-subset[of  $?N$ ] have infinite-N: infinite  $?N$  by blast
    have F-shatters-any-singleton:  $?F$  shatters  $\{n::\text{nat}\}$  for  $n$ 
    proof -

```

```

have Pow-n: Pow {n} = {{n}, {}} by blast
have 1: Pow {n} ⊆ ?F ∩* {n}
proof (cases odd n)
  case True
  from True have (odd - ' {False}) ∩ {n} = {} by blast
  hence 0: {} ∈ ?F ∩* {n} by blast
  from True have (odd - ' {True}) ∩ {n} = {n} by blast
  hence 1: {n} ∈ ?F ∩* {n} by blast
  from 0 1 Pow-n show ?thesis by simp
next
case False
from False have (odd - ' {True}) ∩ {n} = {} by blast
hence 0: {} ∈ ?F ∩* {n} by blast
from False have (odd - ' {False}) ∩ {n} = {n} by blast
hence 1: {n} ∈ ?F ∩* {n} by blast
from 0 1 Pow-n show ?thesis by simp
qed
thus ?thesis by fastforce
qed
then have ?N ⊆ shattered-by ?F unfolding shattered-by-def by force
from 0 infinite-super[OF this infinite-N] show ?thesis by blast
qed
end

```

3 Lemmas involving the cardinality of sets

In this section, we prove some lemmas that make use of the term *card* or provide bounds for it.

```

theory Card-Lemmas
  imports Main
begin

```

lemma *card-Int-copy*:

```

  assumes finite X and A ∪ B ⊆ X and ∃ f. inj-on f (A ∩ B) ∧ (A ∪ B) ∩ (f '
(A ∩ B)) = {} ∧ f ' (A ∩ B) ⊆ X
  shows card A + card B ≤ card X
proof -
  from rev-finite-subset[OF assms(1), of A] rev-finite-subset[OF assms(1), of B]
  assms(2)
  have finite-A: finite A and finite-B: finite B by blast+
  then have finite-A-Un-B: finite (A ∪ B) and finite-A-Int-B: finite (A ∩ B) by
blast+
  from assms(3) obtain f where f-inj-on: inj-on f (A ∩ B)
    and f-disjnt: (A ∪ B) ∩ (f ' (A ∩ B)) = {}
    and f-imj-in: f ' (A ∩ B) ⊆ X by blast
  from finite-A-Int-B have finite-f-imj: finite (f ' (A ∩ B)) by blast
  from assms(2) f-imj-in have union-in: (A ∪ B) ∪ f ' (A ∩ B) ⊆ X by blast

```

from *card-Un-Int*[*OF finite-A finite-B*] **have** $\text{card } A + \text{card } B = \text{card } (A \cup B) + \text{card } (A \cap B)$.
also from *card-image*[*OF f-inj-on*] **have** $\dots = \text{card } (A \cup B) + \text{card } (f \text{ ` } (A \cap B))$ **by** *presburger*
also from *card-Un-disjoint*[*OF finite-A-Un-B finite-f-img f-disjnt*] **have** $\dots = \text{card } ((A \cup B) \cup f \text{ ` } (A \cap B))$ **by** *argo*
also from *card-mono*[*OF assms(1) union-in*] **have** $\dots \leq \text{card } X$ **by** *blast*
finally show *?thesis* .
qed

lemma *card-ge-0*:
assumes $A \neq \{\}$ **and** *finite A*
shows $0 < \text{card } A$
proof –
from *assms(1)* **have** $\{\} \subset A$ **by** *blast*
from *psubset-card-mono*[*OF assms(2) this*] **show** *?thesis* **by** *force*
qed

lemma *finite-diff-not-empty*:
assumes *finite Y* **and** $\text{card } Y < \text{card } X$
shows $X - Y \neq \{\}$
proof
assume $X - Y = \{\}$
hence $X \subseteq Y$ **by** *simp*
from *card-mono*[*OF assms(1) this*] *assms(2)* **show** *False* **by** *linarith*
qed

lemma *obtain-difference-element*:
fixes $F :: 'a \text{ set set}$
assumes $2 \leq \text{card } F$
obtains x **where** $x \in \bigcup F$ $x \notin \bigcap F$
proof –
from *assms* *card-le-Suc-iff*[*of 1 F*] **obtain** $A \ F'$ **where** $0: F = \text{insert } A \ F'$ **and** $1: A \notin F'$ **and** $2: 1 \leq \text{card } F'$ **by** *auto*
from 2 *card-le-Suc-iff*[*of 0 F*] **obtain** $B \ F''$ **where** $3: F' = \text{insert } B \ F''$ **by** *auto*
from $1 \ 3$ **have** *A-noteq-B*: $A \neq B$ **by** *blast*
from $0 \ 3$ **have** *A-in-F*: $A \in F$ **and** *B-in-F*: $B \in F$ **by** *blast+*
from *A-noteq-B* **have** $(A - B) \cup (B - A) \neq \{\}$ **by** *simp*
with *A-in-F B-in-F* **that** **show** *thesis* **by** *blast*
qed

end

4 Lemmas involving the binomial coefficient

In this section we prove lemmas that use the term for the binomial coefficient *choose*.

```
theory Binomial-Lemmas
  imports Main
begin
```

```
lemma choose-mono:
```

```
  assumes  $x \leq y$ 
  shows  $x \text{ choose } n \leq y \text{ choose } n$ 
```

```
proof -
```

```
  have finite  $\{0..<y\}$  by blast
  with finite-Pow-iff[of  $\{0..<y\}$ ] have finiteness: finite  $\{K \in \text{Pow } \{0..<y\}. \text{card } K = n\}$  by simp
  from assms have  $\text{Pow } \{0..<x\} \subseteq \text{Pow } \{0..<y\}$  by force
  then have  $\{K \in \text{Pow } \{0..<x\}. \text{card } K = n\} \subseteq \{K \in \text{Pow } \{0..<y\}. \text{card } K = n\}$  by blast
  from card-mono[OF finiteness this] show ?thesis unfolding binomial-def .
qed
```

```
lemma choose-row-sum-set:
```

```
  assumes finite  $(\bigcup F)$ 
  shows  $\text{card } \{S. S \subseteq \bigcup F \wedge \text{card } S \leq k\} = (\sum i \leq k. \text{card } (\bigcup F) \text{ choose } i)$ 
```

```
proof (induction k)
```

```
  case 0
```

```
  from rev-finite-subset[OF assms] have  $S \subseteq \bigcup F \wedge \text{card } S \leq 0 \longleftrightarrow S = \{\}$  for  $S$  by fastforce
  then show ?case by simp
```

```
next
```

```
  case (Suc k)
  let  $?FS = \{S. S \subseteq \bigcup F \wedge \text{card } S \leq \text{Suc } k\}$ 
  and  $?F\text{-}Asm = \{S. S \subseteq \bigcup F \wedge \text{card } S \leq k\}$ 
  and  $?F\text{-}Step = \{S. S \subseteq \bigcup F \wedge \text{card } S = \text{Suc } k\}$ 
```

```
  from finite-Pow-iff[of  $\bigcup F$ ] assms have finite-Pow-Un-F: finite  $(\text{Pow } (\bigcup F))$  ..
  have  $?F\text{-}Asm \subseteq \text{Pow } (\bigcup F)$  and  $?F\text{-}Step \subseteq \text{Pow } (\bigcup F)$  by fast+
  with rev-finite-subset[OF finite-Pow-Un-F] have finite-F-Asm: finite  $?F\text{-}Asm$ 
and finite-F-Step: finite  $?F\text{-}Step$  by presburger+
```

```
  have  $F\text{-}Un$ :  $?FS = ?F\text{-}Asm \cup ?F\text{-}Step$  and  $F\text{-}disjoint$ :  $?F\text{-}Asm \cap ?F\text{-}Step = \{\}$ 
by fastforce+
```

```
  from card-Un-disjoint[OF finite-F-Asm finite-F-Step F-disjoint]  $F\text{-}Un$  have  $\text{card } ?FS = \text{card } ?F\text{-}Asm + \text{card } ?F\text{-}Step$  by argo
```

```
  also from Suc have  $\dots = (\sum i \leq k. \text{card } (\bigcup F) \text{ choose } i) + \text{card } ?F\text{-}Step$  by argo
  also from n-subsets[OF assms, of Suc k] have  $\dots = (\sum i \leq \text{Suc } k. \text{card } (\bigcup F) \text{ choose } i)$  by force
```

```
  finally show ?case by blast
```

```
qed
```

end

5 Sauer-Shelah Lemma

theory *Sauer-Shelah-Lemma*
imports *Main Shattering Card-Lemmas Binomial-Lemmas*
begin

5.1 Generalized Sauer-Shelah Lemma

To prove the Sauer-Shelah Lemma, we will first prove a slightly stronger fact that every family F shatters at least as many sets as $\text{card } F$. We first fix an element $x \in \bigcup F$ and consider the subfamily $F0$ of sets in the family not containing it. By induction $F0$ shatters at least as many elements of F as $\text{card } F0$. Next we consider the subfamily $F1$ of sets in the family that contain x . Again, by induction $F1$ shatters as many elements of F as its cardinality. The number of elements of F shattered by $F0$ and $F1$ sum up to at least $\text{card } F0 + \text{card } F1 = \text{card } F$. When a set $S \in F$ is shattered by only one of the two subfamilies, say $F0$, it contributes one unit to the set *shattered-by* $F0$ and to *shattered-by* F . However when the set is shattered by both subfamilies, both S and $S \cup \{x\}$ are in *shattered-by* F , so S contributes two units to *shattered-by* $F0 \cup \text{shattered-by } F1$. Therefore, the cardinality of *shattered-by* F is at least equal to the cardinality of *shattered-by* $F0 \cup \text{shattered-by } F1$, which is at least $\text{card } F$.

lemma *sauer-shelah-0*:

fixes $F :: 'a \text{ set set}$
shows $\text{finite } (\bigcup F) \implies \text{card } F \leq \text{card } (\text{shattered-by } F)$
proof (*induction F rule: measure-induct-rule[of card]*)
case (*less F*)
note $\text{finite-}F = \text{finite-UnionD}[OF \text{ less}(2)]$
note $\text{finite-sh}F = \text{finite-shattered-by}[OF \text{ less}(2)]$
show $?case$
proof (*cases 2 ≤ card F*)
case *True*
from *obtain-difference-element[OF True]*
obtain $x :: 'a$ **where** $x\text{-in-Union-}F: x \in \bigcup F$
and $x\text{-not-in-Int-}F: x \notin \bigcap F$ **by** *blast*

Define $F0$ as the subfamily of F containing sets that don't contain x .

let $?F0 = \{S \in F. x \notin S\}$
from $x\text{-in-Union-}F$ **have** $F0\text{-psubset-}F: ?F0 \subset F$ **by** *blast*
from $F0\text{-psubset-}F$ **have** $F0\text{-in-}F: ?F0 \subseteq F$ **by** *blast*
from $\text{subset-shattered-by}[OF F0\text{-in-}F]$ **have** $shF0\text{-subset-sh}F: \text{shattered-by } ?F0 \subseteq \text{shattered-by } F$.
from $F0\text{-in-}F$ **have** $Un\text{-}F0\text{-in-}Un\text{-}F: \bigcup ?F0 \subseteq \bigcup F$ **by** *blast*

$F0$ shatters at least as many sets as $\text{card } F0$ by the induction hypothesis

note $IH-F0 = \text{less}(1)[OF \text{ psubset-card-mono}[OF \text{ finite-}F \text{ } F0\text{-psubset-}F] \text{ rev-finite-subset}[OF \text{ less}(2) \text{ } Un-F0\text{-in-}Un-F]]$

Define $F1$ as the subfamily of F containing sets that contain x

let $?F1 = \{S \in F. x \in S\}$
from $x\text{-not-in-Int-}F$ **have** $F1\text{-psubset-}F$: $?F1 \subset F$ **by** *blast*
from $F1\text{-psubset-}F$ **have** $F1\text{-in-}F$: $?F1 \subseteq F$ **by** *blast*
from $\text{subset-shattered-by}[OF \text{ } F1\text{-in-}F]$ **have** $\text{sh}F1\text{-subset-sh}F$: $\text{shattered-by } ?F1 \subseteq \text{shattered-by } F$.
from $F1\text{-in-}F$ **have** $Un-F1\text{-in-}Un-F$: $\bigcup ?F1 \subseteq \bigcup F$ **by** *blast*

$F1$ shatters at least as many sets as $\text{card } F1$ by the induction hypothesis

note $IH-F1 = \text{less}(1)[OF \text{ psubset-card-mono}[OF \text{ finite-}F \text{ } F1\text{-psubset-}F] \text{ rev-finite-subset}[OF \text{ less}(2) \text{ } Un-F1\text{-in-}Un-F]]$

from $\text{sh}F0\text{-subset-sh}F$ $\text{sh}F1\text{-subset-sh}F$
have shattered-subset : $(\text{shattered-by } ?F0) \cup (\text{shattered-by } ?F1) \subseteq \text{shattered-by } F$ **by** *simp*

There is a set with the same cardinality as the intersection of $\text{shattered-by } F0$ and $\text{shattered-by } F1$ which is disjoint from their union, which is also contained in $\text{shattered-by } F$.

have $f\text{-copies-the-intersection}$:
 $\exists f. \text{inj-on } f (\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1) \wedge$
 $(\text{shattered-by } ?F0 \cup \text{shattered-by } ?F1) \cap (f \text{ ` } (\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1)) = \{\}$ \wedge
 $f \text{ ` } (\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1) \subseteq \text{shattered-by } F$

proof

have $x\text{-not-in-shattered}$: $\forall S \in (\text{shattered-by } ?F0) \cup (\text{shattered-by } ?F1). x \notin S$
unfolding shattered-by-def **by** *blast*

This set is precisely the image of the intersection under $\text{insert } x$.

let $?f = \text{insert } x$
have 0 : $\text{inj-on } ?f (\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1)$
proof
fix $X \ Y$
assume $x0$: $X \in (\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1)$ **and** $y0$: $Y \in (\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1)$
and 0 : $?f X = ?f Y$
from $x\text{-not-in-shattered } x0$ **have** $X = ?f X - \{x\}$ **by** *blast*
also from 0 **have** $\dots = ?f Y - \{x\}$ **by** *argo*
also from $x\text{-not-in-shattered } y0$ **have** $\dots = Y$ **by** *blast*
finally show $X = Y$.
qed

The set is disjoint from the union.

have 1 : $(\text{shattered-by } ?F0 \cup \text{shattered-by } ?F1) \cap ?f \text{ ` } (\text{shattered-by } ?F0 \cap \text{shattered-by } ?F1) = \{\}$

proof (*rule ccontr*)
assume (*shattered-by* ?F0 \cup *shattered-by* ?F1) \cap ?f ‘ (*shattered-by* ?F0 \cap *shattered-by* ?F1) \neq {}
then obtain *S* **where** 10: *S* \in (*shattered-by* ?F0 \cup *shattered-by* ?F1)
and 11: *S* \in ?f ‘ (*shattered-by* ?F0 \cap *shattered-by* ?F1) **by** *auto*
from 10 *x-not-in-shattered* **have** *x* \notin *S* **by** *blast*
with 11 **show** *False* **by** *blast*
qed

This set is also in *shattered-by* *F*.

have 2: ?f ‘ (*shattered-by* ?F0 \cap *shattered-by* ?F1) \subseteq *shattered-by* *F*
proof
fix *S-x*
assume *S-x* \in ?f ‘ (*shattered-by* ?F0 \cap *shattered-by* ?F1)
then obtain *S* **where** 20: *S* \in *shattered-by* ?F0
and 21: *S* \in *shattered-by* ?F1
and 22: *S-x* = ?f *S* **by** *blast*
from *x-not-in-shattered* 20 **have** *x-not-in-S*: *x* \notin *S* **by** *blast*

from 22 *Pow-insert*[*of* *x S*] **have** *Pow S-x* = *Pow S* \cup ?f ‘ *Pow S* **by** *fast*
also from 20 **have** ... = (?F0 \cap * *S*) \cup (?f ‘ *Pow S*) **unfolding** *shattered-by-def* **by** *blast*
also from 21 **have** ... = (?F0 \cap * *S*) \cup (?f ‘ (?F1 \cap * *S*)) **unfolding** *shattered-by-def* **by** *force*
also from *insert-IntF*[*of* *x S ?F1*] **have** ... = (?F0 \cap * *S*) \cup (?f ‘ ?F1 \cap * (?f *S*)) **by** *argo*
also from 22 **have** ... = (?F0 \cap * *S*) \cup (?F1 \cap * *S-x*) **by** *blast*
also from 22 **have** ... = (?F0 \cap * *S-x*) \cup (?F1 \cap * *S-x*) **by** *blast*
also from *subset-IntF*[*OF* *F0-in-F*, *of* *S-x*] *subset-IntF*[*OF* *F1-in-F*, *of* *S-x*]
have ... \subseteq (*F* \cap * *S-x*) **by** *blast*
finally have *Pow S-x* \subseteq (*F* \cap * *S-x*) .
thus *S-x* \in *shattered-by* *F* **unfolding** *shattered-by-def* **by** *blast*
qed

from 0 1 2 **show** *inj-on* ?f (*shattered-by* ?F0 \cap *shattered-by* ?F1) \wedge (*shattered-by* ?F0 \cup *shattered-by* ?F1) \cap (?f ‘ (*shattered-by* ?F0 \cap *shattered-by* ?F1)) = {} \wedge ?f ‘ (*shattered-by* ?F0 \cap *shattered-by* ?F1) \subseteq *shattered-by* *F* **by** *blast*
qed

have *F0-union-F1-is-F*: ?F0 \cup ?F1 = *F* **by** *fastforce*
from *finite-F* **have** *finite-F0*: *finite* ?F0 **and** *finite-F1*: *finite* ?F1 **by** *fastforce+*
have *disjoint-F0-F1*: ?F0 \cap ?F1 = {} **by** *fastforce*

Thus we have the following lower bound on the cardinality of *shattered-by* *F*

from *F0-union-F1-is-F* *card-Un-disjoint*[*OF* *finite-F0* *finite-F1* *disjoint-F0-F1*]
have *card F* = *card* ?F0 + *card* ?F1 **by** *argo*

```

also from IH-F0
have ...  $\leq \text{card } (\text{shattered-by } ?F0) + \text{card } ?F1$  by linarith
also from IH-F1
have ...  $\leq \text{card } (\text{shattered-by } ?F0) + \text{card } (\text{shattered-by } ?F1)$  by linarith
also from card-Int-copy[OF finite-shF shattered-subset f-copies-the-intersection]
have ...  $\leq \text{card } (\text{shattered-by } F)$  by argo
finally show ?thesis .
next

  If  $F$  contains less than 2 sets, the statement follows trivially

case False
hence  $\text{card } F = 0 \vee \text{card } F = 1$  by force
thus ?thesis
proof
  assume  $\text{card } F = 0$ 
  thus ?thesis by auto
next
  assume asm:  $\text{card } F = 1$ 
  hence  $F\text{-not-empty}$ :  $F \neq \{\}$  by fastforce
  from shatters-empty[OF F-not-empty] have  $\{\{\}\} \subseteq \text{shattered-by } F$  unfolding
shattered-by-def by fastforce
  from card-mono[OF finite-shF this] asm show ?thesis by fastforce
qed
qed
qed

```

5.2 Sauer-Shelah Lemma

The generalized version immediately implies the SauerShelah Lemma, because only $(\sum_{i \leq k}. n \text{ choose } i)$ of the subsets of an n -item universe have cardinality less than $k + 1$. Thus, when $(\sum_{i \leq k}. n \text{ choose } i) < \text{card } F$, there are not enough sets to be shattered, so one of the shattered sets must have cardinality at least $k + 1$

corollary *sauer-shelah*:

```

fixes  $F :: 'a \text{ set set}$ 
assumes finite  $(\bigcup F)$  and  $(\sum_{i \leq k}. \text{card } (\bigcup F) \text{ choose } i) < \text{card } F$ 
shows  $\exists S. (F \text{ shatters } S \wedge \text{card } S = k + 1)$ 
proof -
  let  $?K = \{S. S \subseteq \bigcup F \wedge \text{card } S \leq k\}$ 
  from finite-Pow-iff[of F] assms(1) have finite-Pow-Un: finite  $(\text{Pow } (\bigcup F))$  by
fast

  from sauer-shelah-0[OF assms(1)] assms(2) have  $(\sum_{i \leq k}. \text{card } (\bigcup F) \text{ choose } i) < \text{card } (\text{shattered-by } F)$  by linarith
  with choose-row-sum-set[OF assms(1), of k] have  $\text{card } ?K < \text{card } (\text{shattered-by } F)$  by presburger

  from finite-diff-not-empty[OF finite-subset[OF - finite-Pow-Un] this]

```

obtain S **where** $S \in \text{shattered-by } F - ?K$ **by** *blast*
then have $F\text{-shatters-}S$: F *shatters* S **and** $S \subseteq \bigcup F$ **and** $\neg(S \subseteq \bigcup F \wedge \text{card } S \leq k)$ **unfolding** *shattered-by-def* **by** *blast+*
then have $\text{card-}S\text{-ge-}Suc\text{-}k$: $k + 1 \leq \text{card } S$ **by** *simp*
from $\text{obtain-subset-with-card-}n$ [$OF \text{ card-}S\text{-ge-}Suc\text{-}k$] **obtain** S' **where** $\text{card } S' = k + 1$ **and** $S' \subseteq S$ **by** *blast*
from $\text{this}(1)$ supset-shatters [$OF \text{ this}(2) F\text{-shatters-}S$] **show** *?thesis* **by** *blast*
qed

5.3 Sauer-Shelah Lemma for hypergraphs

corollary *sauer-shelah-2*:

fixes $X :: 'a \text{ set set}$ **and** $S :: 'a \text{ set}$
assumes *finite* S **and** $X \subseteq \text{Pow } S$ **and** $(\sum_{i \leq k. \text{card } S \text{ choose } i}) < \text{card } X$
shows $\exists Y. (X \text{ shatters } Y \wedge \text{card } Y = k + 1)$
proof –
from $\text{assms}(2)$ **have** $0: \bigcup X \subseteq S$ **by** *blast*
from sum-mono [$OF \text{ choose-mono}$ [$OF \text{ card-mono}$ [$OF \text{ assms}(1) 0$]]] **have** $(\sum_{i \leq k. \text{card } (\bigcup X) \text{ choose } i}) \leq (\sum_{i \leq k. \text{card } S \text{ choose } i})$ **by** *fast*
with sauer-shelah [$OF \text{ finite-subset}$ [$OF 0 \text{ assms}(1)$]] $\text{assms}(3)$ **show** *?thesis* **by** *simp*
qed

5.4 Alternative statement of the Sauer-Shelah Lemma

We can also state the SauerShelah Lemma in terms of the *VC-dim*. If the VC dimension of F is k , then F can consist at most of $(\sum_{i \leq k. \text{card } (\bigcup F) \text{ choose } i})$ sets, which is in $\mathcal{O}(\text{card } (\bigcup F) \wedge k)$

corollary *sauer-shelah-alt*:

assumes *finite* $(\bigcup F)$ **and** $\text{VC-dim } F = k$
shows $\text{card } F \leq (\sum_{i \leq k. \text{card } (\bigcup F) \text{ choose } i})$
proof (*rule ccontr*)
assume $\neg \text{card } F \leq (\sum_{i \leq k. \text{card } (\bigcup F) \text{ choose } i})$ **hence** $(\sum_{i \leq k. \text{card } (\bigcup F) \text{ choose } i}) < \text{card } F$ **by** *linarith*
from sauer-shelah [$OF \text{ assms}(1) \text{ this}$] **obtain** S **where** F *shatters* S **and** $\text{card } S = k + 1$ **by** *blast*
from $\text{this}(1) \text{ this}(2)$ [*symmetric*] **have** $k + 1 \in \{\text{card } S \mid S. F \text{ shatters } S\}$ **by** *blast*
from cSup-upper [$OF \text{ this bdd-above-finite}$ [$OF \text{ finite-image-set}$ [$OF \text{ finite-shattered-by}$ [*unfolded shattered-by-def*, $OF \text{ assms}(1)$]]], *folded VC-dim-def*]
 $\text{assms}(2)$ **show** *False* **by** *force*
qed
end

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