

Programming Assignment (PA) – 3 (Demo Session Simulation)

CS307- Operating Systems

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Solution Report

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```

1  #include <stdio.h>
2  #include <pthread.h>
3  #include <stdlib.h>
4  #include <unistd.h>
5  #include <string.h>
6  #include <fcntl.h>
7  #include <sys/wait.h>
8  #include <semaphore.h>
9
10 int s_inside = 0;
11 int a_inside = 0;
12 int s_free = 0;
13 int a_free = 0;
14 int s_demo = 0;
15 int a_demo = 0;
16 int g_found = 0;
17 int found_waiting = 0;
18
19 pthread_mutex_t mutex = PTHREAD_MUTEX_INITIALIZER;
20
21 sem_t student_enter;
22 sem_t student_leave;
23 sem_t group_ready_s;
24 sem_t group_ready_a;
25 sem_t demo_end_s;
26 sem_t demo_end_a;
27 sem_t found_barrier;

```

Added needed libraries also initialized needed mutex and counters, also defined needed semaphores and barrier to be used I program as global variables.

```

137 int main(int argc, char* argv[]) {
138     if (argc != 3) {
139         printf("Usage: ./demosim <assistant_count> <student_count>\n");
140         return 1;
141     }
142
143     int a_count = atoi(argv[1]);
144     int s_count = atoi(argv[2]);
145     int total = a_count + s_count;
146
147     if (a_count <= 0 || s_count != a_count * 2) {
148     }
149     else {
150         pthread_t threads[a_count + s_count];
151
152         sem_init(&student_enter, 0, 0);
153         sem_init(&student_leave, 0, 0);
154         sem_init(&group_ready_s, 0, 0);
155         sem_init(&group_ready_a, 0, 0);
156         sem_init(&demo_end_s, 0, 0);
157         sem_init(&demo_end_a, 0, 0);
158         sem_init(&found_barrier, 0, 0);
159
160         printf("My program complies with all the conditions.\n");
161
162         for (int i = 0; i < total; i++) {
163             if (i < s_count) {
164                 pthread_create(&threads[i], NULL, student, NULL);
165             }
166             else {
167                 pthread_create(&threads[i], NULL, assistant, NULL);
168             }
169         }
170
171         for (int i = 0; i < total; i++) {
172             pthread_join(threads[i], NULL);
173         }
174
175         sem_destroy(&student_enter);
176         sem_destroy(&student_leave);
177         sem_destroy(&group_ready_s);
178         sem_destroy(&group_ready_a);
179         sem_destroy(&demo_end_s);
180         sem_destroy(&demo_end_a);
181         sem_destroy(&found_barrier);
182     }
183
184     printf("The main terminates\n");
185     return 0;
186 }

```

In main function, get assistant and student counts from arguments passed to main function if available, then checks counts, if assistant count is positive integer and students count is exactly 2 times of assistant count then

program proceeds. initializes all mutexes and barrier with value of 0, creates needed assistant and student threads. Then waits for threads to finish, then deletes semaphores and barrier and main function terminates.

```

105 void* student(void* args) {
106     pthread_mutex_lock(&mutex);
107     printf("Thread ID : %ld, Role : Student, I want to enter the classroom.\n", pthread_self());
108     pthread_mutex_unlock(&mutex);
109
110     sem_wait(&student_enter);
111
112     pthread_mutex_lock(&mutex);
113     g_found_wait();
114     printf("Thread ID : %ld, Role : Student, I entered the classroom.\n", pthread_self());
115     s_inside++;
116     s_free++;
117     initial_session_condition();
118     pthread_mutex_unlock(&mutex);
119
120     sem_wait(&group_ready_s);
121
122     pthread_mutex_lock(&mutex);
123     printf("Thread ID : %ld, Role : Student, I am now participating.\n", pthread_self());
124     s_demo++;
125     end_session_condition();
126     pthread_mutex_unlock(&mutex);
127
128     sem_wait(&demo_end_s);
129
130     pthread_mutex_lock(&mutex);
131     printf("Thread ID : %ld, Role : Student, I left the classroom.\n", pthread_self());
132     s_inside--;
133     sem_post(&student_leave);
134     pthread_mutex_unlock(&mutex);
135 }

```

In student function, it first prints that, that students want to enter room. Then waits for “student_enter” semaphore to be signalled (released in other words), this semaphore will be posted by assistant. Then it will call “g_found_wait()”, which I will explain later. Then prints I am entered room and increases “s_inside” and “s_free” by 1. Then calls “initial_session_condition()”, which I will explain later. Then it waits for “group_ready_s” semaphore to be signalled. If it is signalled it means that a demo group with 2 students and 1 assistant has been assembled and then prints that it is participating and increases “s_demo” by 1. then calls “end_session_condition()”, which I will explain later. After that it waits for “demo_end_s” semaphore to be signalled, if it is then it means that its demo session is over and it is allowed to print left the room and decrease “s_inside” by 1 and signals “student_leave” semaphore, this semaphore is being waited by assistant.

Important: At first, I had implemented the enter and leave of students and assistant using while loops but while I was in an office hour TA reminded me that infinite while loop may cause busy-wait problem, so I removed loops and used semaphores to keep enter and leave rules.

```

70 void* assistant(void* args) {
71     pthread_mutex_lock(&mutex);
72     g_found_wait();
73     printf("Thread ID : %ld, Role : Assistant, I entered the classroom.\n", pthread_self());
74     sem_post(&student_enter);
75     sem_post(&student_enter);
76     sem_post(&student_enter);
77     a_free++;
78     a_inside++;
79     initial_session_condition();
80     pthread_mutex_unlock(&mutex);
81
82     sem_wait(&group_ready_a);
83
84     pthread_mutex_lock(&mutex);
85     printf("Thread ID : %ld, Role : Assistant, I am now participating.\n", pthread_self());
86     a_demo++;
87     end_session_condition();
88     pthread_mutex_unlock(&mutex);
89
90     sem_wait(&demo_end_a);
91
92     pthread_mutex_lock(&mutex);
93     printf("Thread ID : %ld, Role : Assistant, demo is over.\n", pthread_self());
94     pthread_mutex_unlock(&mutex);
95
96     sem_wait(&student_leave);
97     sem_wait(&student_leave);
98
99     pthread_mutex_lock(&mutex);
100     printf("Thread ID : %ld, Role : Assistant, I left the classroom.\n", pthread_self());
101     a_inside--;
102     pthread_mutex_unlock(&mutex);
103 }

```

In assistant function, it first calls “g_found_wait()” which I will explain later, then prints that it entered room and increases “a_free” and “a_inside” by 1. Also, signals “student_enter” 3 times which means tells 3 different students they can enter so this keeps following equation valid “s_inside < 3 * a_inside” so if equation is false then student waits for another assistant to enter and signal. then it waits for “group_ready_a” semaphore to be signalled. If it is signalled it means that a group has been assembled and it prints that it is participating and increases “a_demo” by 1. also calls “end_session_condition()” which I will explain later. Then it waits for “demo_end_a” semaphore to be signalled, if it is then it means that every person in session has started participating and now assistant can print demo is over. Then it waits for “student_leave” semaphore to be signalled twice, the reason it needs to be signalled twice is because each assistant can leave when 2 other students have left and signalled “student_leave” each of them. Then it prints it left room and decreases “a_inside” by 1 and function finishes.

```

38 void initial_session_condition() {
39     if (s_free >= 2 && a_free >= 1) {
40         g_found = 1;
41         s_free -= 2;
42         a_free--;
43         sem_post(&group_ready_a);
44         sem_post(&group_ready_s);
45         sem_post(&group_ready_s);
46     }
47     else {
48         if (found_waiting > 0) {
49             sem_post(&found_barrier);
50             found_waiting--;
51         }
52     }
53 }

```

So, what “initial_session_conditions()” does is that it first checks if inside room is there any 2 or more students free and 1 or more assistants free to create a session, if yes then it makes “g_found” as 1, here this variable is integer but it will be used like a Boolean variable so it can be either 0 or 1 throughout the flow of program. Then since now 2 students and 1 assistant are no longer free, decreases “s_free” by 2 and decreases “a_free” by 1. then it signals “group_ready_a” once and “group_ready_s” twice, so this wakes 2 students and 1 assistant up and lets them know that they can start participating. However, if the numbers do not allow to form a session group then it goes to else part and checks if there is any person waiting for a previously formed team which means if “found_waiting” is larger than 1, then is decreases the “found_waiting” and signals “found_barrier” semaphore but this semaphore is used as a barrier, so as soon as a group formed and then second group want to be created this barrier will put new group to sleep until previous group has reached a point and I will explain that point later.

```

29 void g_found_wait() {
30     while (g_found == 1) {
31         found_waiting++;
32         pthread_mutex_unlock(&mutex);
33         sem_wait(&found_barrier);
34         pthread_mutex_lock(&mutex);
35     }
36 }

```

In “g_found_wait()” it creates a loop that will run until “g_found” is 1, this means that a group has been created and groups after them are sleeping. so inside loop it increases “found_waiting” by 1 and then waits for “found_barrier” to be signalled but since it must wait, it unlocks mutex first waits and then again lock mutex, this avoids any busy-wait because if it was inside mutex block it was never going to release since barrier has never signalled before and it causes a waiting until infinity.

```

55 void end_session_condition() {
56     if (s_demo >= 2 && a_demo >= 1) {
57         s_demo -= 2;
58         a_demo--;
59         g_found = 0;
60         if (found_waiting > 0) {
61             sem_post(&found_barrier);
62             found_waiting--;
63         }
64         sem_post(&demo_end_a);
65         sem_post(&demo_end_s);
66         sem_post(&demo_end_s);
67     }
68 }
69

```

In “end_session_condition()” it tries to check if there are 2 or more students and 1 or more assistants participating in demo, if yes it means that all members of session have announced that they are participating so they are waiting to be able to leave or end demo. So, “s_demo” decreases by 2 and “a_demo” decreases by 1 also since this group has done “g_found” resets to 0 and if there is any group creating process waiting it wakes one of persons up and decreases “found_waiting” by 1. also signals “demo_end_a” once and “demo_end_s” twice. By these signals let’s 2 students to leave and 1 assistant to finish demo session.

Important: the reason I used “g_found” “found_waiting” and “found_barrier” is because at first, I had implemented program without those variable, counter and semaphore, and there was a 50% chance that program was messing order of group creation up when assistant count entered as input to program is larger than 1. For example, I will show and explain this problem in one of test runs for 2 assistants and 4 students. Lets assume assistant 1 and 2 students have entered room and since the “initial_session_condition()” can create a group, one of the members in this group starts participating, then another student enters room and since “initial_session_condition()” had signalled “group_ready_s” twice previously this 3rd student can start participating although it has entered after one of the students was started participating, so this mixes all group members with each other but after adding those 3 new variable, counter and barrier, it will make all new entered person to wait until currently created group has done and they are ready to leave room.

* Therefore, to sum up, there is 2 important aspects, first is that instead of using infinite loops for checking conditions to allow students to enter and assistants to leave I used semaphores which will not cause any busy-wait problems. Secondly, I used a complex waiting mechanism so that each group is assembled all other possible groups must wait until all group members of currently assembled group have participated, this makes sure that if there are multiple groups, they will not mix into each other.