AI in Game Playing: Sokoban Solver
CS 221 Project Progress Report
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1 Introduction

Sokoban is a very popular transportation puzzle game that is played extensively with its variants. Solving Sokoban is a well acknowledged area of research because it exist as a NP-Hard problem and PSPACE-Complete. The high level strategies, complicated techniques and high branching factor of this game makes it difficult to solve. The goal of our project is to develop an AI agent that can play the game effectively. In this progress report, we delve deeper into the ideas furnished in our proposal by (1) pointing out the game mechanics in a straightforward manner (2) describing the states and explaining how it is modeled in our algorithms (3) detailing our algorithms to calculate the best moves for the levels of the game (4) defining proper pruning techniques that are implemented to ameliorate the performance of the algorithm (5) providing results and comparing the developed algorithm using standard metrics of evaluation. We conclude with the next steps we plan to take up to finish this project.

2 General Game Mechanics

The rigorous rules and the confined space of the game is described exhaustively in the Section 2 (*Task Definition*) of our proposal. To briefly define the game play in simple terminology, we can say that the goal of this game if to move the crates to proper storage locations by moving in a constricted space provided in each level. The player fails to complete the game if he gets locked up in positions where the player is not able to move himself or the crates further. This game play is shown in Fig.1
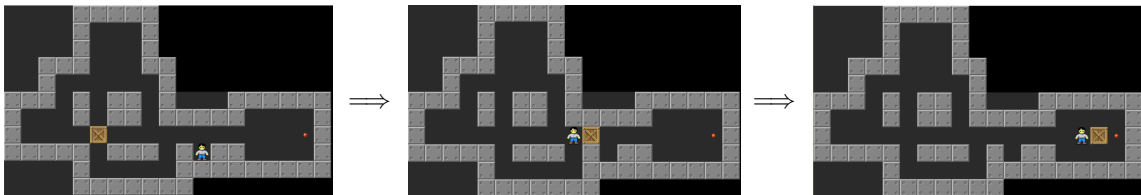

Figure 1: Game play of Sokoban

3 States and Modeling

The game of sokoban can alternatively be considered as a search problem where we essentially look out for boxes and storage locations. So intuitively, it has a valid start state and end state. The start state is the state given by the original game developer whereas the end state is the state when all crates are transferred to proper storage locations. The actions can be moving in all directions with a cost associated with it which leads to the successor state.

For modeling this game, we have a standard notation that is used to independently and distinctly distinguish between all the objects in the game. Each level (*like that in Figure 1*) is given in a unique representation. The inputs for modeling consist of the following characters: "#", " ", "$", "@", "." and "+". Each of these characters have a special attribute of the game assigned to it: **"#"** is a wall, **" "** is a free space, **"$"** is a

box, "." is a goal place, "*" is a boxes placed on a goal, "@" is for Sokoban and "+" is for Sokoban on a goal. So the level described in Figure 1 is modified as follows:

```
                    #####
                    #   #
                    #   #
                  ###   ##
                   #     #
                 ### # ## #   ######
                 #   # ## #####    #
                 #     $          .#
                 ##### ### #@##    #
                     #      #########
                     #######
```

Figure 2: Game model

4   Algorithms

Sokoban has many unique properties which are not there in other similar problems as Rubiks cube or Lloyd Fifteen Puzzle. Moves are irreversible so making a bad move can make you lose the game. Given the game mechanics, several factors are taken into account for deciding on an optimal algorithm: (1) There is a need for returning an set of moves quickly because the game is played real time also. Consequently, an algorithm that is able to produce moves in a short amount of processing time and returns a strong move is desirable. (2) The constraint added to game. Constraints in this game refers to the walls that are present inside the outer boundary and restrict the path of the player. (3) The depth to which the search is constructed. Since this search problem can lead to multitudes of states and can at times lead to infinite search depths, once concern is to know how much the algorithm can explore in the search space.

Given these considerations, we decided to evaluate considerable number of algorithms and compare them based on their space/time complexities. The strength of the algorithms in this space lies in their ability to quickly determine sequences of moves that yield relatively strong results. We have implemented four algorithms namely back-propagation, Depth First Search (DFS), Breadth First Search (BFS) and Uniform Cost Search (UCS) and shown their results.

4.1   Baseline (*Backtracking Algorithm*) and Oracle Implementation

The baseline of the project is the backtracking search algorithm and the oracle of the project is the high level human intelligence that is used to solve the game. So essentially, each level has a predefined number of moves which corresponds to the minimum moves that one can take to solve the game. This minimum number of steps forms the oracle. On the other hand, the backtracking algorithm is one of the simplest recursive algorithms and forms the baseline for our project but is seldom used widely because of its high time complexity. It just recurses to all states and finds the minimum cost in reaching the goal. The space complexity is O(D) and the time complexity is $O(b^D)$ which is very high. The implementation of Backtracking algorithm is as follows:

4.2 Depth First Search

Depth First Search (DFS) is a special case of backtracking search algorithm. The search starts from the root and proceeds to the farthest node before backtracking. The difference between this and the Backprop is that this stops the search once a goal is reached and does not care if it is not minimum. The space and time complexities, on the worst case, are the same as the baseline algorithm but stops when it finds the solution. The DFS algorithm is given below:

4.3 Breadth First Search

Breadth First Search (BFS), as the name says, explores the search space in the increasing order of the depth and the costs of traveling from one state to another is assumed to be a positive number. Typically, this algorithm is often associated with the concept of stack and queue and pushing and popping from the stack. Due to the larger states explored at shorter depths, the space complexity is very high of about $O(b^d)$ and the time complexity is $O(b^d)$. The pseudo code of BFS is:

4.4 Uniform cost Search

For any search problem, Uniform Cost Search (UCS) is the better algorithm than the previous ones. The search algorithm explores in branches with more or less same cost. This consist of a priority queue where the path from the root to the node is the stored element and the depth to a particular node acts as the priority. UCS assumes all the costs to be non negative. While the DFS algorithm gives maximum priority to maximum depth, this gives maximum priority to the minimum cumulative cost. It is implemented as:

5 Pruning Techniques

Pruning is a terminology used in machine learning and artificial intelligence that are used to reduce the size of the decision trees by removing the selected sections of the tree that provide undesirable results. We have implemented two techniques till now in the project and are described below:

5.1 Move appension

One of the primary problem in search problems is that the computation time becomes unimaginably high when the search space is big. In such cases, if we have priori knowledge about the systems, we can limit in the beginning of the search problem all the cases where we have impossible actions. For instance, the Figure 3 depicts a case where the only acceptable action is to move Up. The possible actions for any algorithm can be moving in all the directions which can be reduced to one by Move Appension where we restrict all the impossible actions.

5.2 Hashing

Hashing is a well known pruning method used to tune the algorithm to perform better. It follows the logic that decision which leads to the states that are already visited are considered as suboptimal. So all the states are stored in the hash table and at each point, a comparison is made between the current state and the stored state. If there is a match, the same action corresponding to the one in the hashing table is avoided.
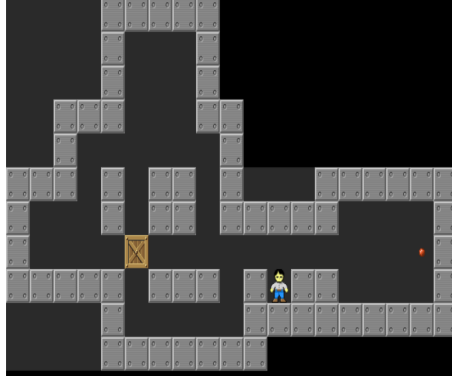
Figure 3: Move Appension Example

6 Results

7 Next Steps

We have implemented and extensively analyzed these four algorithms and compared based on their performances at different levels. Yet there are still other AI algorithms and pruning techniques which we would like to implement in the future. In particular, we plan to account for:

- implementing A* search algorithm with different heuristics. We are aiming to implement A* with Eucledian distance, Manhattan Distance and Hungarian Distance.

- exploring and trying to implement advanced search algorithms like the Depth First Search with Iterative Deepening (DFS-ID) and Breadth First Search with Iterative Deepening (BFS-ID).

- learning and trying to implement more pruning techniques like the PI corral pruning, Dead lock Tables etc.

8   Appendix 1

9 Appendix 2