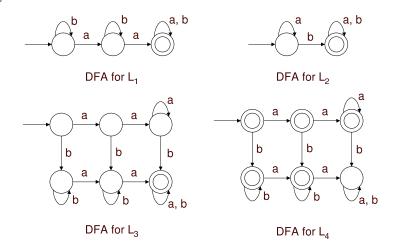
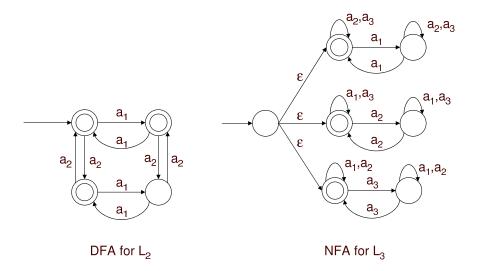
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Solutions to Written Assignment 1

1. DFAs for languages



2. DFAs and NFAs for languages



3. Regular expressions

- (a) $(0+1)^*0(0+1)^*11$
- (b) $\epsilon + 0 + 1 + (00 + 10 + 11)(0 + 1)^*$
- (c) 0*10*(0*10*10*)*
- $(\mathrm{d})\ (00^*10^*10^*) + (0^*100^*10^*) + (0^*10^*100^*)$

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- 4. Regular expression and description from DFA
 - The language of all strings that begin and end with a 0 and have an even, non-zero number of 1s (that is, 2k 1s, for some $k \ge 1$).
 - 00*10*10*(10*10*)*0

5. Lexical analysis

(a) An even length string of 0s prints all as, while an odd length string of 0s will have one c at the end (because of the maximal munch rule). Thus, strings of 0s generate the language a^*c ?. Interspersed 1s generate bs, so the full language is

$$(a^*c?b^+)^*a^*c?$$

A common mistake might be to incorrectly account for the priority between the rules for 0 and 00.

- (b) Every input string can be divided into three sections (any of which might be empty)
 - An initial string of 1s
 - A section where a string of 0s is followed by a string of 1s, and this pattern repeats (i.e., the language $((0+)(1+))^*$)
 - A final string of 0s

Thus, this specification can print any string in the language $c^*a^*b^*$.

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