# GEBZE TECHNİCAL UNIVERSITY

## **CSE312**

## Operating Systems

## Homework 1 Report

#### 1. How to Run?

Open the terminal and navigate to the source where the makefile directory. Then, compile the program by typing "make" and the .iso file will be created. Then construct the virtual machine on VirtualBox Program with connecting this file. Once executed, the program will produce output as specified by strategy 1 which is asked for in the pdf file. To clean up generated files, use "make clean", which will remove the files generated.

## 2. Syscall Implementations

The system call implementation uses the **SyscallHandler** class, derived from InterruptHandler, to manage system call interrupts via the IDT. System calls are invoked using a software interrupt (int 0x80), with the **SyscallHandler** dispatching these calls based on the value in the eax register. This design treats system calls as a special type of interrupt, allowing user programs to safely request kernel services. And provided source code was already have an interrupt mechanism with 3 structure:

Interrupt Descriptor Table (IDT)

Each entry in the IDT points to an interrupt service routine (ISR). The

SetInterruptDescriptorTableEntry function is used to set up individual entries in the IDT and initialized and connected to the appropriate routines within the constructor of the

InterruptManager class. After setting up the IDT entries, the IDT is loaded into the CPU using the lidt instruction. This instruction tells the CPU where to find the IDT.

Interrupt Manager Class

The **InterruptManager** class is responsible for setting up the IDT, initializing the PIC, and managing the registration and dispatching of interrupts.

• Interrupt Handler Class

## Sets up the entry for the system call interrupt is part of the InterruptManager constructor:

```
// Set up entry for system call interrupt
SetInterruptDescriptorTableEntry(0x80, CodeSegment, &HandleInterruptRequest0x80, 0, IDT_INTERRUPT_GATE);
```

## Triggers system calls using a software interrupt (int \$0x80):

```
int myos::getPid()
 asm("int $0x80" : "=c" (pld) : "a" (SYSCALLS::GETPID));
 return pld; }
oid myos::waitpid(common::uint8_t wPid)
 asm("int $0x80" :: "a" (SYSCALLS::WAITPID), "b" (wPid)); }
void myos::sys_exit()
 asm("int $0x80" : : "a" (SYSCALLS::EXIT)); }
oid myos::sysprintf(char* str)
 asm("int $0x80" :: "a" (SYSCALLS::PRINTF), "b" (str)); }
oid myos::fork()
 asm("int $0x80" :: "a" (SYSCALLS::FORK)); }
nt myos::fork_with_pid(int pid)
 asm("int $0x80" : "=c" (pid) : "a" (SYSCALLS::FORK));
 return pid; }
 nt myos::exec(void entrypoint())
 asm("int $0x80" : "=c" (result) : "a" (SYSCALLS::EXEC), "b" ((uint32_t)entrypoint));
nt myos::addTask(void entrypoint())
 asm("int $0x80": "=c" (result): "a" (SYSCALLS::ADDTASK), "b" ((uint32_t)entrypoint));
```

return result; }

Implementation and handle the actual logic for each system call and interacts with the TaskManager to perform operations:

```
// Implementation of system call handlers
common:uint32_t InterruptHandler:sys_getpid()
{
    return interruptManager > taskManager > GetPld();
}
common:uint32_t InterruptHandler:sys_exec(common:uint32_t entrypoint)
{
    return interruptManager > taskManager > ExecTask((void (*)())entrypoint);
}
common:uint32_t InterruptHandler:sys_addTask(common:uint32_t entrypoint)
{
    return interruptManager > taskManager > AddTask((void (*)())entrypoint);
}
common:uint32_t InterruptHandler:sys_fork(CPUState* cpustate)
{
    return interruptManager > taskManager > ForkTask(cpustate);
}
bool InterruptHandler:sys_exit()
{
    return interruptManager > taskManager > ExitCurrentTask();
}
bool InterruptHandler:sys_waitpid(common:uint32_t pid)
{
    return interruptManager > taskManager > WaitTask(pid);
}
```

## Initialize the SyscallHandler and add it to the InterruptManager:

```
SyscallHandler::SyscallHandler(InterruptManager* interruptManager, uint8_t InterruptNumber)
:InterruptHandler(interruptManager, InterruptNumber + interruptManager->HardwareInterruptOffset())
{}
uint32_t SyscallHandler::HandleInterrupt(uint32_t esp) {
  CPUState* cpu = (CPUState*)esp;
  switch(cpu->eax) {
    case SYSCALLS::EXEC:
       esp = InterruptHandler::sys_exec(cpu->ebx);
    case SYSCALLS::FORK:
       cpu->ecx = InterruptHandler::sys_fork(cpu);
      return InterruptHandler::HandleInterrupt(esp);
    case SYSCALLS::PRINTF:
       printf((char*)cpu->ebx);
    case SYSCALLS::EXIT:
       if (InterruptHandler::sys_exit()) {
         return InterruptHandler::HandleInterrupt(esp); }
    case SYSCALLS::WAITPID:
      if (InterruptHandler::sys_waitpid(esp)) {
         return InterruptHandler::HandleInterrupt(esp); }
    case SYSCALLS::GETPID:
       cpu->ecx = InterruptHandler::sys_getpid();
    case SYSCALLS::ADDTASK:
       cpu->ecx = InterruptHandler::sys_addTask(cpu->ebx);
  return esp; }
```

## 3. ROUND ROBIN SCHEDULING

The scheduling part also provided by the source code. This function does:

- Triggered by a hardware timer.
- o Calls the Schedule method in the TaskManager to potentially switch processes.
- Selects the next process in the ready queue.
- Performs a context switch to the new process.
- o Prints the process table with relevant information.
- Display all entries in the process table.
- o Include details such as PID, PPID, state.

```
CPUState* TaskManager::Schedule(CPUState* cpustate)
 if (numTasks <= 0)
    return cpustate;
  if (currentTask >= 0)
    tasks[currentTask].cpustate = cpustate;
  int findTask = (currentTask + 1) % numTasks;
  while (tasks[findTask].taskState != READY)
    if (tasks[findTask].taskState == WAITING && tasks[findTask].waitPid > 0)
      int waitTaskIndex = getIndex(tasks[findTask].waitPid);
      if (waitTaskIndex > -1 && tasks[waitTaskIndex].taskState != WAITING)
         if (tasks[waitTaskIndex].taskState == FINISHED)
           tasks[findTask].waitPid = 0;
           tasks[findTask].taskState = READY;
         else if (tasks[waitTaskIndex].taskState == READY)
    findTask = (findTask + 1) % numTasks;
  currentTask = findTask;
  PrintProcessTable();
```

return tasks[currentTask].cpustate; }

## 4. LIFECYCLE (STRATEGY ONE AND KERNELMAIN)

### **Collatz Function:**

```
void collatz(int n) {
    printf("Collatz with parameter: ");
    printf("\n");
    printf("\n");
    printf("\n");
    printf(": ");
    while (n != 1) {
        if (n % 2 == 0) {
            n = 3 * n + 1;}
        printNum(n);
        printf("\n");
    }
    printf("\n");
}
```

## **Long Running Program Function:**

```
common::uint32_t long_running_program(int n){
  int result = 0;
  for(int i = 0; i < n; ++i){
    for(int j = 0; j < n; ++j){
      result += i*j; } }
  return result; }</pre>
```

## Below function strategyOne does:

- Forks three child processes to run the Collatz function.
- Forks three child processes to run the long-running program.
- Waits for all six child processes to finish before printing that all child processes have ended.

## strategyOne Function:

```
void strategyOne(){
 int childPID[5];
 printf("{{{{{{{\test Strategy 1 }}}}}}}");
 printf("\n");
 int initialPid = getPid();
 printf("Initial Process PID: ");
 printNum(initialPid);
 printf("\n");
 int collatzTest, longTest;
 int numForks = 3; // Number of times to fork to create 6 programs
   int childPid;
   int result = fork_with_pid(childPid);
      printNum(childCounter);
      printf(". Child Process with");
      printf(" PID: ");
      printNum(getPid());
      printf(" runs collatz");
      printf("\n");
      collatzTest = getPid()+5;
      sys_exit(); }
   else { childPID[i] = childCounter++; } }
   int childPid;
   int result = fork_with_pid(childPid);
      printNum(childCounter);
      printf(". Child Process with");
      printf(" PID: ");
      printNum(getPid());
      printf(" runs long running program with ");
      longTest = 1000 + i;
      printNum(longTest);
      printf("\n");
      printf("Result:");
      printNum(long_running_program(longTest));
      sys_exit(); } else { childPID[i] = childCounter++; }}
 for(int i = 0; i < 6; ++i) { waitpid(childPID[i]); }
 printf("All child processes are ended."); }
```

### kernelMain Function:

## **Necessary System Calls Test Functions:**

```
void forkTestExample() {
    //printf("((((((( Test Fork ))))))))");
    //printf("((((((( Test Fork ))))))))");
    //printf("\n");
    int parentPID=getPid();
    printf("\n");
    int childPID = fork_with_pid(getPid());
    if(childPID==0){
        printf("Forked child pid ");
        printf("\n");
        sys_exit();}
    else {
        printf("Parent pid which should be the same as before: ");
        printf("\n");
        sys_exit();
        sys_exit();
}
```

## 5. OUTPUTS

## **Sysem Calls Test and Process Table**

```
-----Welcome the ZortOS------
Parent Pid:1
Child Task 2
forkTest End Pid:2
......
PID
      PPID
              STATE
      0
              READY
      1
              FINISHED
Child Task 1
forkTest End Pid:1
PID
     PPID
              STATE
      0
              FINISHED
              FINISHED
^^^^^
```

## **Strategy One Test and Process Table**

```
--Welcome the the ZortOS---
Exec syscall is testing with PID: 1
Exec System Call Runned Properly!
Starting with parent pid: 2
{{{{{{{}}}}}}}}}}}
Initial Process PID: 3
Forked child pid 4
1. Child Process with PID: 5 runs collatz
Collatz with parameter: 10
Result: 5 16 8 4 2 1
Parent pid which should be the same as before: 2
2. Child Process with PID: 6 runs collatz
Collatz with parameter: 11
Result: 34 17 52 26 13 40 20 10 5 16 8 4 2 1
3. Child Process with PID: 7 runs collatz
Collatz with parameter: 12
Result: 6 3 10 5 16 8 4 2 1
4. Child Process with PID: 8 runs long running program with 1000
Result:392146832
5. Child Process with PID: 9 runs long running program with 1001
Result:1392146832
6. Child Process with PID: 10 runs long running program with 1002
Result:-1899817463
All child processes are ended.
```

It was hard to takle screenshots during this test but because of timer interrupts but I've managed to see chil process's were task switching by callint process table print function insde scheduling function.

PID	PPID	STATE
	Θ	RUNNING
2	1	FINISHED
3	1	FINISHED
4	1	FINISHED
5	1	FINISHED
6	1	READY
7	1	FINISHED

```
********
PID
       PPID
                STATE
       0
                RUNNING
       1
                FINISHED
       1
                FINISHED
       1
                FINISHED
       1
                FINISHED
       1
                FINISHED
       1
                FINISHED
```