

## 0.1 Digit: the basic building block

As the basic building block of numerals, we will demonstrate how to choose a suitable representation for digits in this section.

### 0.1.1 Fin

To represent a digit, we use a datatype that is conventionally called *Fin* which can be indexed to have some exact number of inhabitants.

```
data Fin : ℕ → Set where
  zero : {n : ℕ} → Fin (suc n)
  suc  : {n : ℕ} (i : Fin n) → Fin (suc n)
```

The definition of **Fin** looks the same as  $\mathbb{N}$  on the term level, but different on the type level. The index of a **Fin** increases with every **suc**, and there can only be at most  $n$  of them before reaching **Fin (suc n)**. In other words, **Fin n** would have exactly  $n$  inhabitants.

### 0.1.2 Definition

**Digit** is simply just a synonym for **Fin**, indexed by the number of digits **d** of a system. Since the same digit may represent different values in different numeral systems, it is essential to make the context clear.

```
Digit : ℕ → Set
Digit d = Fin d
```

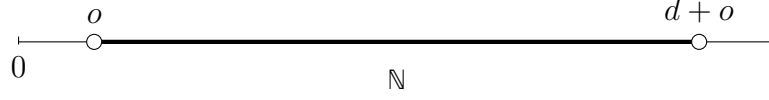
Ordinary binary digits for example can thus be represented as:

```
Binary : Set
Binary = Digit 2

零 : Binary
零 = zero

一 : Binary
一 = suc zero
```

### 0.1.3 Digit Assignment



Digits are assigned to  $\mathbb{N}$  together with the offset  $o$  of a system, ranging from  $o$  to  $d + o$ .

```
Digit-to $\mathbb{N}$  :  $\forall \{d\} \rightarrow \text{Digit } d \rightarrow \mathbb{N} \rightarrow \mathbb{N}$ 
Digit-to $\mathbb{N}$   $x \ o = \text{to}\mathbb{N} \ x + o$ 
```

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However, not all natural numbers can be converted to digits. The value has to be in a certain range, between  $o$  and  $d + o$ . Values less than  $o$  are increased to  $o$ . Values greater than  $d + o$  are prohibited by the supplied upper-bound.

```
Digit-from $\mathbb{N}$  :  $\forall \{d\}$ 
   $\rightarrow (n \ o : \mathbb{N})$ 
   $\rightarrow (\text{upper-bound} : d + o \geq n)$ 
   $\rightarrow \text{Digit } (\text{suc } d)$ 
Digit-from $\mathbb{N}$   $\{d\} \ n \ o \ \text{upper-bound}$  with  $n \div o \leq? d$ 
Digit-from $\mathbb{N}$   $\{d\} \ n \ o \ \text{upper-bound} \mid \text{yes } p = \text{from}\mathbb{N}\leq (s \leq s \ p)$ 
Digit-from $\mathbb{N}$   $\{d\} \ n \ o \ \text{upper-bound} \mid \text{no } \neg p = \text{contradiction } p \ \neg p$ 
  where  $p : n \div o \leq d$ 
         $p = \text{start}$ 
           $n \div o$ 
           $\leq \{ \div n\text{-mono } o \ \text{upper-bound} \}$ 
           $(d + o) \div o$ 
           $\approx \{ m + n \div n \equiv m \ d \ o \}$ 
           $d$ 
           $\square$ 
```

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properties

$$\mathbb{N} \begin{array}{c} \xrightarrow{\text{Digit-from}\mathbb{N}} \\ \xleftarrow{\text{Digit-to}\mathbb{N}} \end{array} \text{Digit } d$$

<sup>1</sup>  $\text{to}\mathbb{N} : \forall \{n\} \rightarrow \text{Fin } n \rightarrow \mathbb{N}$   
converts from  $\text{Fin } n$  to  $\mathbb{N}$ .

<sup>2</sup>  $\text{from}\mathbb{N}\leq : \forall \{m \ n\} \rightarrow m < n \rightarrow \text{Fin } n$   
converts from  $\mathbb{N}$  to  $\text{Fin } n$  given the number is small enough.

**Digit-from $\mathbb{N}$ -to $\mathbb{N}$**  states that the value of a natural number should remain the same, after converting back and forth between **Digit** and  $\mathbb{N}$ .

```

Digit-from $\mathbb{N}$ -to $\mathbb{N}$  :  $\forall \{d\ o\}$ 
   $\rightarrow (n : \mathbb{N})$ 
   $\rightarrow$  (lower-bound :  $o \leq n$ )
   $\rightarrow$  (upper-bound :  $d + o \geq n$ )
   $\rightarrow$  Digit-to $\mathbb{N}$  (Digit-from $\mathbb{N}$  {d} n o upper-bound)  $o \equiv n$ 
Digit-from $\mathbb{N}$ -to $\mathbb{N}$  {d} {o} n lb ub with  $n \div o \leq? d$ 
Digit-from $\mathbb{N}$ -to $\mathbb{N}$  {d} {o} n lb ub | yes q =
  begin
    to $\mathbb{N}$  (from $\mathbb{N} \leq (s \leq s\ q)) + o$ 
     $\equiv$  ( cong (  $\lambda\ x \rightarrow x + o$ ) (to $\mathbb{N}$ -from $\mathbb{N} \leq (s \leq s\ q))$  )
       $n \div o + o$ 
     $\equiv$  (  $m \div n + n \equiv m\ lb$  )
      n
  ■
Digit-from $\mathbb{N}$ -to $\mathbb{N}$  {d} {o} n lb ub | no  $\neg q$  = contradiction q  $\neg q$ 
  where   q :  $n \div o \leq d$ 
          q = +n-mono-inverse o (
            start
               $n \div o + o$ 
             $\approx$  (  $m \div n + n \equiv m\ lb$  )
              n
             $\leq$  ( ub )
              d + o
            □)

```

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Digits have a upper-bound and a lower-bound after evaluation.

```

Digit-upper-bound :  $\forall \{d\} \rightarrow (o : \mathbb{N}) \rightarrow (x : \text{Digit } d) \rightarrow \text{Digit-to}\mathbb{N}\ x\ o < d + o$ 
Digit-upper-bound {d} o x = +n-mono o (bounded x)

Digit-lower-bound :  $\forall \{d\} \rightarrow (o : \mathbb{N}) \rightarrow (x : \text{Digit } d) \rightarrow \text{Digit-to}\mathbb{N}\ x\ o \geq o$ 
Digit-lower-bound {d} o x =  $m \leq n + m\ o$  (to $\mathbb{N}\ x$ )

```

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<sup>3</sup> to $\mathbb{N}$ -from $\mathbb{N} \leq$  :  $\forall \{m\ n\} (m < n : m < n) \rightarrow \text{to}\mathbb{N} (\text{from}\mathbb{N} \leq m < n) \equiv m$   
states that a number should remain the same after converting back and forth.

<sup>4</sup> bounded :  $\forall \{n\} (i : \text{Fin } n) \rightarrow \text{to}\mathbb{N}\ i < n$   
a property about the upper-bound of a **Fin** n.

## 0.1.4 Operations

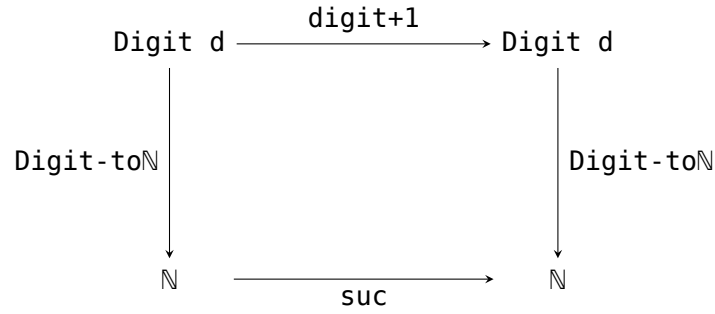
### Increment

To increment a digit, the digit must not be *the greatest*.

```
digit+1 : ∀ {d}
  → (x : Digit d)
  → (¬greatest : ¬ (Greatest x))
  → Fin d
digit+1 x ¬greatest =
  fromℕ≤ {suc (toℕ x)} (≤Λ≡=< (bounded x) ¬greatest)
```

Where  $\leq \wedge \equiv < (\text{bounded } x) \neg \text{greatest} : \text{suc } (\text{to}\mathbb{N} \ x) < d$ .

### properties



A digit taking these two routes should result in the same  $\mathbb{N}$ .

```
digit+1-toℕ : ∀ {d o}
  → (x : Digit d)
  → (¬greatest : ¬ (Greatest x))
  → Digit-toℕ (digit+1 x ¬greatest) o ≡ suc (Digit-toℕ x o)
digit+1-toℕ {d} {o} x ¬greatest =
  begin
    Digit-toℕ (digit+1 x ¬greatest) o
  ≡< cong (λ w → w + o) (toℕ-fromℕ≤ (≤Λ≡=< (bounded x) ¬greatest)) >
    suc (Digit-toℕ x o)
  ■
```

### Increase then Subtract

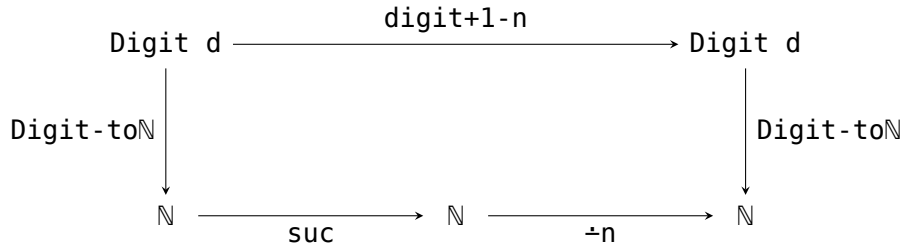
Increases a digit and then subtract it by  $n$ . This function is useful for implementing *carrying*. When the digit to increase is already the greatest, we have to subtract it by an amount (usually the base) after the increment.

```

digit+1-n : ∀ {d}
  → (x : Digit d)
  → Greatest x
  → (n : ℕ)
  → n > 0
  → Digit d
digit+1-n x greatest n n>0 =
  fromℕ≤ (digit+1-n-lemma x greatest n n>0)

```

properties



A digit taking these two routes should result in the same  $\mathbb{N}$ .

```

digit+1-n-toℕ : ∀ {d o}
  → (x : Digit d)
  → (greatest : Greatest x)
  → (n : ℕ)
  → (n>0 : n > 0)
  → n ≤ d
  → Digit-toℕ (digit+1-n x greatest n n>0) o ≡ suc (Digit-toℕ x o) ÷ n
digit+1-n-toℕ {zero} {o} () greatest n n>0 n≤d
digit+1-n-toℕ {suc d} {o} x greatest n n>0 n≤d =
  begin
    toℕ (digit+1-n x greatest n n>0) + o
  ≡( cong (λ w → w + o) (toℕ-fromℕ≤ (digit+1-n-lemma x greatest n n>0)) )
    suc (toℕ x) ÷ n + o
  ≡( +-comm (suc (toℕ x) ÷ n) o )
    o + (suc (toℕ x) ÷ n)
  ≡( sym (+-÷-assoc o {suc (toℕ x)} {n}) (
    start
      n
    ≤( n≤d )
      suc d
    ≈( sym greatest )
      suc (toℕ x)
    □)
  )
  (o + suc (toℕ x)) ÷ n
  ≡( cong (λ w → w ÷ n) (+-comm o (suc (toℕ x))) )
    suc (toℕ x) + o ÷ n
  ■)

```

### 0.1.5 Special Digits

#### The Greatest Digit

**constructions** The greatest digit of a system is constructed by converting the index  $d$  to  $\text{Fin}$ .

```
greatest-digit :  $\forall d \rightarrow \text{Digit } (\text{suc } d)$ 
greatest-digit  $d = \text{from}\mathbb{N} \ d$ 
```

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**predicates** Judging whether a digit is the greatest by converting it to  $\mathbb{N}$ . This predicate also comes with a decidable version.

```
Greatest :  $\forall \{d\} (x : \text{Digit } d) \rightarrow \text{Set}$ 
Greatest  $\{d\} \ x = \text{suc } (\text{to}\mathbb{N} \ x) \equiv d$ 

Greatest? :  $\forall \{d\} (x : \text{Digit } d) \rightarrow \text{Dec } (\text{Greatest } x)$ 
Greatest?  $\{d\} \ x = \text{suc } (\text{to}\mathbb{N} \ x) \stackrel{?}{=} d$ 
```

**properties** Converting from the greatest digit to  $\mathbb{N}$  should result in  $d + o$ .

```
greatest-digit-to $\mathbb{N}$  :  $\forall \{d \ o\}$ 
   $\rightarrow (x : \text{Digit } (\text{suc } d))$ 
   $\rightarrow \text{Greatest } x$ 
   $\rightarrow \text{Digit-to}\mathbb{N} \ x \ o \equiv d + o$ 
greatest-digit-to $\mathbb{N}$   $\{d\} \ \{o\} \ x \ \text{greatest} = \text{cancel-suc } ($ 
  begin
     $\text{suc } (\text{Digit-to}\mathbb{N} \ x \ o)$ 
   $\equiv ( \text{refl } )$ 
     $\text{suc } (\text{to}\mathbb{N} \ x) + o$ 
   $\equiv ( \text{cong } (\lambda w \rightarrow w + o) \ \text{greatest} )$ 
     $\text{suc } d + o$ 
   $\blacksquare)$ 
```

A digit is the greatest if and only if it is greater than or equal to all other digits. This proposition is proven by induction on both of the compared digits.

```
greatest-of-all :  $\forall \{d\} (o : \mathbb{N}) \rightarrow (x \ y : \text{Digit } d)$ 
   $\rightarrow \text{Greatest } x$ 
```

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<sup>5</sup>  $\text{from}\mathbb{N} : \forall \{n\} \rightarrow \text{Fin } n \rightarrow \mathbb{N}$   
construct the greatest possible  $\text{Fin } n$  when given an index  $n$ .

```

    → Digit-to $\mathbb{N}$  x o ≥ Digit-to $\mathbb{N}$  y o
greatest-of-all o zero    zero    refl    = ≤-refl
greatest-of-all o zero    (suc ()) refl
greatest-of-all o (suc x) zero    greatest
    = +n-mono o {zero} {suc (to $\mathbb{N}$  x)} z≤n
greatest-of-all o (suc x) (suc y) greatest
    = s≤s (greatest-of-all o x y (cancel-suc greatest))

```

## The Carry

A carry is a digit that is transferred to a more significant digit to compensate the “loss” of the original digit.

**constructions** The carry is defined as the greater of these two values:

- the least digit of a system
- the digit that is assigned to 1

In case that the least digit is assigned to 0, rendering the carry useless. Since the least digit is determined by the offset  $o$ , the value of the carry is defined as follows.

```

carry :  $\mathbb{N} \rightarrow \mathbb{N}$ 
carry o = 1  $\sqcup$  o

```

And then we construct the carry by converting `carry o` to `Digit`:

```

carry-digit :  $\forall d o \rightarrow 2 \leq \text{suc } d + o \rightarrow \text{Digit } (\text{suc } d)$ 
carry-digit d o proper =
  Digit-from $\mathbb{N}$ 
    (carry o)
    o
    (carry-upper-bound {d} proper)

```

**properties** The value of the carry should remain the same after converting back and forth.

```

carry-digit-to $\mathbb{N}$  :  $\forall d o$ 
  → (proper :  $2 \leq \text{suc } (d + o)$ )
  → Digit-to $\mathbb{N}$  (carry-digit d o proper) o  $\equiv$  carry o
carry-digit-to $\mathbb{N}$  d o proper
  = Digit-from $\mathbb{N}$ -to $\mathbb{N}$ 

```

```

(carry o)
(m≤nℓm o 1)
(carry-upper-bound {d} proper)

```

The carry also have an upper-bound and a lower-bound, similar to that of `Digit`.

```

carry-lower-bound : ∀ {o} → carry o ≥ o
carry-lower-bound {o} = m≤nℓm o 1

carry-upper-bound : ∀ {d o} → 2 ≤ suc d + o → carry o ≤ d + o
carry-upper-bound {d} {zero} proper = ≤-pred proper
carry-upper-bound {d} {suc o} proper = n≤m+n d (suc o)

```