

Sorting
=

Algorithm
=

① Bubble Sort

Bubble sort is the simplest sorting algorithm that works by repeatedly swapping the adjacent elements if they are in wrong order.

7	11	9	2	17	4	6
0	1	2	3	4	5	6

1st pass

7	11	9	2	17	4	6
0	1	2	3	4	5	6

5 comp.
5 pass

7	9	11	2	17	4	6
0	1	2	3	4	5	6

5 swap

7	9	2	11	17	4	6
0	1	2	3	4	5	6

swap

7	9	2	11	17	4	6
0	1	2	3	4	5	6

7	9	2	11	17	4	6
0	1	2	3	4	5	6

7	9	2	11	17	4	6
0	1	2	3	4	5	6

2nd pass

7	9	2	11	17	4	6
0	1	2	3	4	5	6

sorted

4 comp.
4 pass

7	9	2	11	17	4	6
0	1	2	3	4	5	6

4 swap

7	9	2	11	17	4	6
0	1	2	3	4	5	6

3 comp.
3 pass

7	9	2	11	17	4	6
0	1	2	3	4	5	6

3 swap

7	9	2	11	17	4	6
0	1	2	3	4	5	6

17

~~3rd pass~~ 0 1 2 3 4 5
 7 2 9 4 11 17

2 comparison
3 possible swap

swap 2 7 9 4 11 17 → 0, 1
 2 7 9 4 11 17 → 1, 2
 2 7 4 9 11 17 → 2, 3

~~4th pass~~ 0 1 2 3 4 5
 2 7 4 9 11 17

2 comparison
3 possible swap

swap 2 11 7 4 9 11 17 → 0, 1
 2 4 7 9 11 17 → 1, 2

~~5th pass~~ 0 1 2 3 4 5
 2 4 7 9 11 17

1 comparison
2 possible swap

swap 2 4 7 9 11 17 → 0, 1

↓ sorted array

no. of passes = len. of array - 1.

Total no. of Comparision

$= = = = =$

$$1 + 2 + 3 + 4 + \dots + (n-1) = n(n-1)$$

$\frac{n(n-1)}{2}$

$$= O(n^2)$$

Checking for stability

7, 8, 2, 7*

7, 8, 2, 7*

7, 2, 8, 7*

7, 2, 7*, 8

2, 7, 7*, 8

==

order is same as the order in input array. So, Bubble sort is stable sorting algorithm.

Bubble sort is not a Adaptive in nature. But we can make Bubble sort as Adaptive. (Having Benefits of sorted array).

II program for Bubble Sort

```
void Bubble_Sort (int Arr[], int size) {  
    for (int j=0; j < size - 1; j++) {  
        for (int i=0; i < size - j - 1; i++) {  
            if (Arr[i] > Arr[i+1])  
                swap (Arr[i], Arr[i+1]);  
        }  
    }  
}
```

⑪ Insertion Sort

17	12	3	4	1	
0	1	2	3	4	

1st pass sorted $i \rightarrow$ unsorted \leftarrow

temp
12

17	12	3	4	1	
0	1	2	3	4	

17	12	3	4	1	
0	1	2	3	4	

2nd pass $i \rightarrow$ \leftarrow temp

3

17	12	3	4	1	
0	1	2	3	4	

7, 12, 4, 1
7, 12, 4, 1
3, 7, 12, 4, 1

3	7	12	4	1	
0	1	2	3	4	

3rd pass

4

3	7	12	4	1	
0	1	2	3	4	

3, 7, 12, 1
3, 7, 12, 1

3	4	7	12	1	
0	1	2	3	4	

$j \leftarrow$

4th pass
=

temp

1

i

j

3 4 7 12

3, 4, 7, 12

3, 4, 7, 12

3, 4, 7, 12

0, 3, 4, 7, 12

1 3 4 7 12

0 1 2 3 4

→ sorted array.

No. of pass = len - 1 = 7 - 1 .

$$\begin{aligned}
 & \text{Total number of Comparison/Swap} \\
 &= 1 + 2 + 3 + 4 + \dots + (n-1) \\
 &= \frac{n(n-1)}{2} \\
 &= O(n^2)
 \end{aligned}$$

for Best Case - In that

1	3	4	7	9
---	---	---	---	---

$$\begin{aligned}
 \text{Total Comparison} &= 1 + 1 + 1 + 1 = 4 \\
 &= n-1
 \end{aligned}$$

Analysis

$$= O(n)$$

(i) Total pass = $n-1$

Total Comparison = $O(n^2)$

(ii) Best Case $\rightarrow O(n)$ (when array is sorted)

(iii) Stable \rightarrow Yes

(iv) Adaptive \rightarrow Yes

// program for insertion sort

```
void Insertion-Sort (int Arr[], int size) {
```

```
    for (int i=1; i<size; i++) {
```

```
        int temp = Arr[i];
```

```
        int j = i-1;
```

```
        while (j >= 0 && Arr[j] > temp) {
```

```
            Arr[j+1] = Arr[j];
```

```
i = i + 1; } // outer loop } // inner loop ;
```

```
int =
```

```
(i) = Arr[j+1] = temp;
```

```
and (i) = Arr[j+1] = temp;
```

```
if (i < size) {
```

```
    Arr[i] = temp;
```

```
else {
```

```
    break;
```

Selection Sort

= finding min. element.

8	0	7	1	3	
0	1	2	3	4	

→ next sorted

1st pass:

0	8	7	1	3	
0	1	2	3	4	

4 Comparison

2nd pass:

0	1	7	8	3	
8	0	1	2	3	4

3 Comparison

3rd pass:

0	1	3	8	7	
0	1	2	3	4	

2 Comparison

4th pass:

0	1	3	7	8	
0	1	2	3	4	

Comparison

only one
element left, so
no need to do 5th
pass.

Total Comparison = $1 + 2 + 3 + \dots + (n-1)$

$$= \frac{n(n-1)}{2}$$

$$\text{max-sweeps} = n-1$$

stability

1	8	7	8	8*	1
---	---	---	---	----	---

1	8	8*	7
---	---	----	---

1	8	9	10
---	---	---	----

1	7	8	8*	8
---	---	---	----	---

↑ sorted

so, it is not a stable algorithm.

Adaptive \rightarrow NO.

1	8	7	6	5	4	3	2	1	0
---	---	---	---	---	---	---	---	---	---

With some starting \rightarrow

W. breaking

selection sort

using

// program for selection sort

```
void selectionSort (int arr[], int size) {
```

```
    for (int i=0; i<size-1; i++) {
```

```
        int min = arr[i];
```

```
        for (int j=i+1; j<size; j++) {
```

```
            if (min > arr[j]) {
```

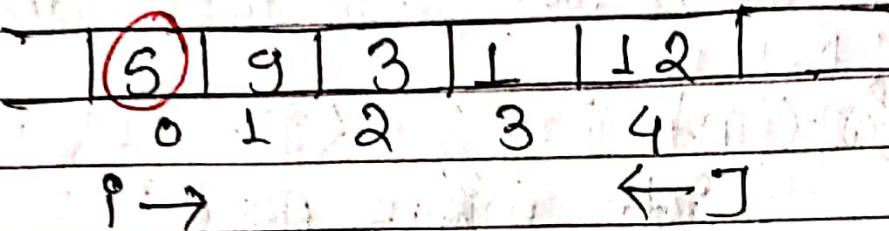
```
                min = arr[j];
```

```
        swap (arr[i], min);
```

```
}
```

3

Quick Sort



partition 1 | pivot | partition 2
values $<$ pivot values \geq pivot

Algo:- partitioning

- ① $i = \text{low}$
- ② $j = \text{high}$
- ③ $\text{pivot} = \text{low}$
- ④ $i++$ until element $>$ pivot is found.
- ⑤ $j--$ until element $<$ pivot is found.
- ⑥ swap $A[i]$ & $A[j]$ & repeat ④ & ⑤ until $j \leq i$.
- ⑦ swap pivot and $A[j]$.

5	9	3	1	12
---	---	---	---	----

0 1 2 3 4

↑
low ↑
high

high > low → so swap (high, low)

5	1	3	9	12
---	---	---	---	----

0 1 2 3 4

↑
low ↑
high

5	1	3	9	12
---	---	---	---	----

0 1 2 3 4

↑
from low

low > high

so, swap (pivot, high).

3	1	5	9	12
---	---	---	---	----

0 1 2 3 4

← →

II program for partition in quick sort.

```

int partition(int arr[], int start, int end) {
    int pivot = arr[start];
    int LB = start + 1;
    int UB = end;

    while (UB > LB) {
        while (arr[LB] <= pivot) {
            LB++;
        }
        while (arr[UB] > pivot) {
            UB--;
        }
        if (UB > LB) {
            swap(arr[LB], arr[UB]);
        }
    }
    swap(arr[start], arr[UB]);
    return UB;
}

```

II program for Quick Sort

```
void quickSort (int Arr[], int LB, int UB)
```

```
    int position;
    if (LB < UB) {
        position = partition(Arr, LB, UB);
        quickSort (Arr, LB, position - 1);
        quickSort (Arr, position + 1, UB);
    }
}
```

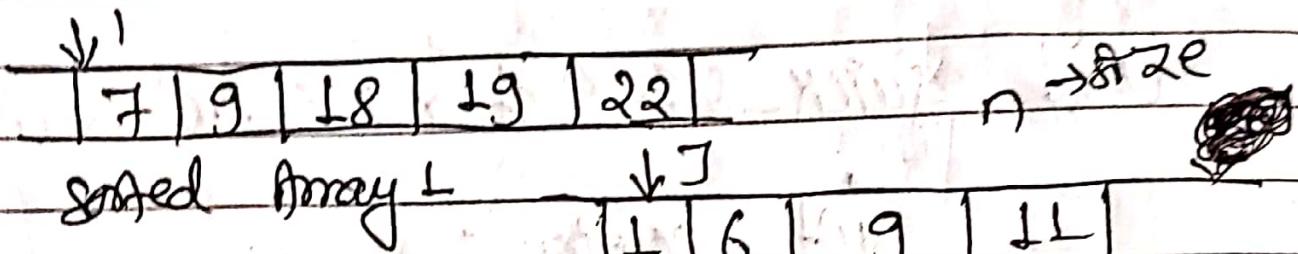
Analysis

- ① Best case $\rightarrow O(n \log n)$
- ② Worst Case $\rightarrow O(n^2)$
- ③ Stable $\rightarrow NO$

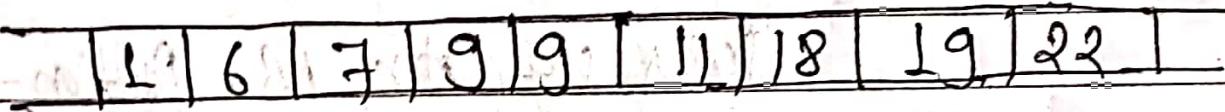
merge sort

$i = j = k = 0$

$m \rightarrow \text{size}$



$(m+n)$



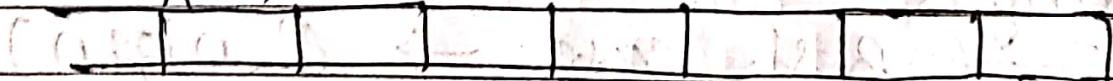
$i \rightarrow m(\text{size})$

$J \rightarrow n(\text{size})$

A

B

$K \rightarrow$



C

$(m+n)$

void merge (int A[], int B[], int C[], int m, n) {
 int i, j, k;
 i = j = k = 0;

```

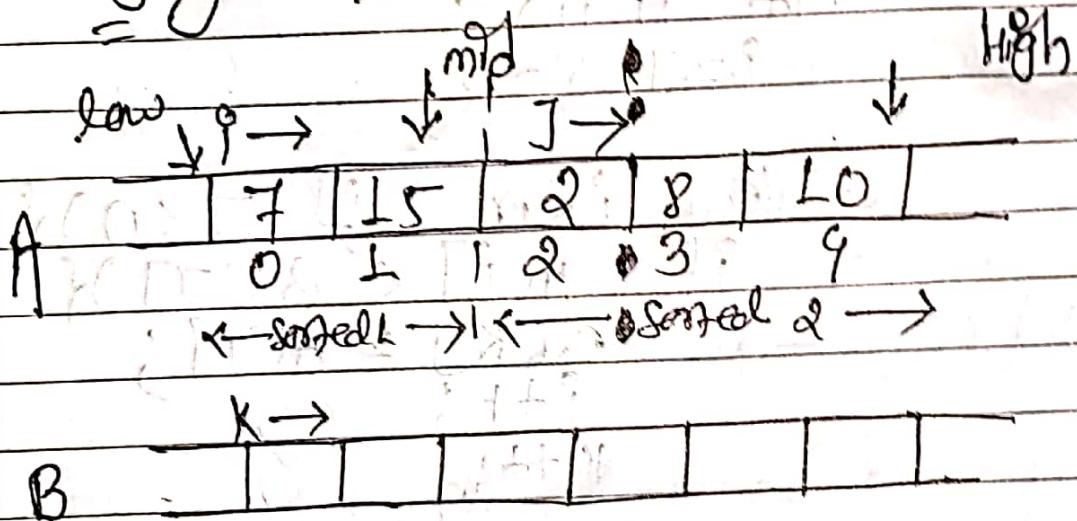
while ( i < m && j < n ) {
    if ( A[ i ] < B[ j ] ) {
        C[ k ] = A[ i ];
        i++;
        K++;
    } else {
        C[ k ] = B[ j ];
        J++;
        K++;
    }
}

```

```
while (i < m) { // Copy all remaining  
    C[i] = A[i]; // element from  
    i++; // A to C.  
    k++;  
}
```

```
while (j < n) { // copy all remaining  
    C[K] = B[j];   -ng element from  
    j++;           B to C.  
    K++;  
}
```

merging in a single array



```
void merge (int A[], int mid, int low, int high) {
    int i, j, k;
    int B[high + 1];
    i = low;
    j = mid + 1;
    K = low;
```

```
while (i <= mid && j <= high) {
    if (A[i] < A[j]) {
        B[K] = A[i];
        i++;
        K++;
    }
```

else {

```
        B[K] = A[j];
        j++;
        K++;
    }
```

}

while ($i \leq mid$) { // Copy all remaining
 $B[k] = A[i]$; elements from A
 $k++$;
 $i++$;

} // Copying

while ($j \leq high$) { // Copy all remaining
 $B[k] = A[j]$; elements from B
 $k++$;
 $j++$;

// Copy All B to A.

for ($i = low$; $i \leq high$; $i++$) {

$A[i] = B[i]$;
 }

}

11. merge sort program.

```
void mergeSort (int A[], int low, int high)
```

```
if (low < high) {
```

```
    mid = (low + high) / 2;
```

```
    mergeSort (A, low, mid);
```

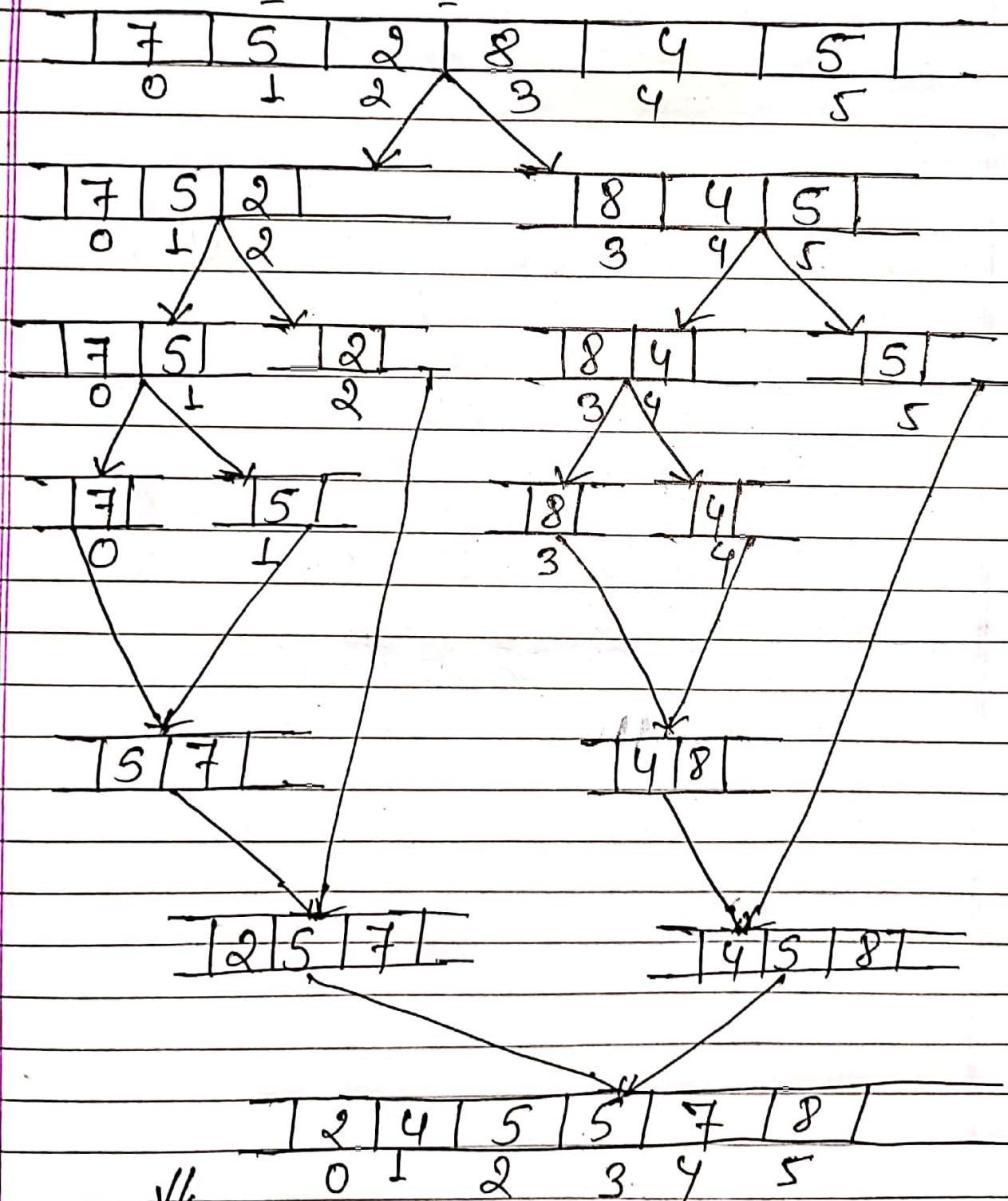
```
    mergeSort (A, mid+1, high);
```

```
    merge (A, mid, low, high);
```

O(2^{10})
O(1)

$\{ \{f + g\} \cdot d \} = f \cdot d + g \cdot d$
 $\{ f \cdot g \} \cdot d = f \cdot d \cdot g$

merge sort



↓
sorted array.