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Algorithm 6 No-U-Turn Sampler with Dual Averaging
    Given \theta^0, \delta, \mathcal{L}, M, M^{\text{adapt}}:
    Set \epsilon_0 = \text{FindReasonableEpsilon}(\theta), \mu = \log(10\epsilon_0), \bar{\epsilon}_0 = 1, \bar{H}_0 = 0, \gamma = 0.05, t_0 = 10, \kappa = 0.75.
    for m = 1 to M do
         Sample r^0 \sim \mathcal{N}(0, I).
         Resample u \sim \text{Uniform}([0, \exp\{\mathcal{L}(\theta^{m-1} - \frac{1}{2}r^0 \cdot r^0\}])
Initialize \theta^- = \theta^{m-1}, \ \theta^+ = \theta^{m-1}, \ r^- = r^0, r^+ = r^0, j = 0, \theta^m = \theta^{m-1}, n = 1, s = 1.
         while s = 1 do
              Choose a direction v_i \sim \text{Uniform}(\{-1,1\}).
              if v_i = -1 then
                   \theta^-, r^-, -, -, \theta', n', s', \alpha, n_\alpha \leftarrow \text{BuildTree}(\theta^-, r^-, u, v_i, j, \epsilon_{m-1}\theta^{m-1}, r^0).
              else
                   -, -, \theta^+, r^+, \theta', n', s', \alpha, n_\alpha \leftarrow \text{BuildTree}(\theta^+, r^+, u, v_i, j, \epsilon_{m-1}, \theta^{m-1}, r^0).
              end if
              if s' = 1 then
                   With probability \min\{1, \frac{n'}{n}\}, set \theta^m \leftarrow \theta'.
              end if
              n \leftarrow n + n'.
              s \leftarrow s' \mathbb{I}[(\theta^+ - \theta^-) \cdot r^- > 0] \mathbb{I}[(\theta^+ - \theta^-) \cdot r^+ > 0].
              j \leftarrow j + 1.
         end while
         if m \leq M^{\text{adapt}} then
              Set \bar{H}_m = \left(1 - \frac{1}{m+t_0}\right) \bar{H}_{m-1} + \frac{1}{m+t_0} (\delta - \frac{\alpha}{n_0}).
              Set \log \epsilon_m = \mu - \frac{\sqrt{m}}{\epsilon} \bar{H}_m, \log \bar{\epsilon}_m = m^{-\kappa} \log \epsilon_m + (1 - m^{-\kappa}) \log \bar{\epsilon}_{m-1}.
         else
              Set \epsilon_m = \bar{\epsilon}_{Madapt}.
         end if
    end for
    function BuildTree(\theta, r, u, v, j, \epsilon, \theta^0, r^0)
    if j = 0 then
         Base case—take one leapfrog step in the direction v.
         \theta', r' \leftarrow \text{Leapfrog}(\theta, r, v\epsilon).
         n' \leftarrow \mathbb{I}[u \le \exp\{\mathcal{L}(\theta') - \frac{1}{2}r' \cdot r'\}].
         s' \leftarrow \mathbb{I}[u < \exp\{\Delta_{\max} + \bar{\mathcal{L}}(\theta') - \frac{1}{2}r' \cdot r'\}].
         return \theta', r', \theta', r', \theta', n', s', \min\{1, \exp\{\mathcal{L}(\theta') - \frac{1}{2}r' \cdot r' - \mathcal{L}(\theta^0) + \frac{1}{2}r^0 \cdot r^0\}\}.1
    else
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Recursion—implicitly build the left and right subtrees.

 $s' \leftarrow s'' \mathbb{I}[(\theta^+ - \theta^-) \cdot r^- \ge 0] \mathbb{I}[(\theta^+ - \theta^-) \cdot r^+ > 0]$ 

With probability  $\frac{n''}{n'+n''}$ , set  $\theta' \leftarrow \theta''$ . Set  $\alpha' \leftarrow \alpha' + \alpha''$ ,  $n'_{\alpha} \leftarrow n'_{\alpha} + n''_{\alpha}$ .

return  $\theta^-, r^-, \theta^+, r^+, \theta', n', s', \alpha', n'_{\alpha}$ .

if s' = 1 then if v = -1 then

 $n' \leftarrow n' + n''$ 

else

end if

end if

 $\theta^-, r^-, \theta^+, r^+, \theta', n', s', \alpha', n'_{\alpha} \leftarrow \text{BuildTree}(\theta, r, u, v, j - 1, \epsilon, \theta^0, r^0).$ 

 $\theta^-, r^-, -, -, \theta'', n'', s'', \alpha'', n''_{\alpha} \leftarrow \text{BuildTree}(\theta^-, r^-, u, v, j - 1, \epsilon, \theta^0, r^0).$ 

 $-, -, \theta^+, r^+, \theta'', n'', s'', \alpha'', n''_{\alpha} \leftarrow \text{BuildTree}(\theta^+, r^+, u, v, j - 1, \epsilon, \theta^0, r^0).$